

Enrichment Reading

In this page we give some suggestions for further reading for each lecture. These readings are entirely optional and some of them may be very challenging.

- Lecture 1:
 - For more about the algorithm "PIERCE" see [this paper](#). Vlado Keselj has the [best bound](#) known.
 - Jon Bentley's two books, *Programming Pearls* and *More Programming Pearls* are essential reading for any programmer.
- Lecture 2:
 - Many people have studied the subrange sum problem, creating efficient algorithms and generalizations. For example, look at Brodal and Jorgensen, A linear time algorithm for the k maximal sums problem, MFCS 2007, pp. 442-453 and Fan, Lee, et al., An optimal algorithm for maximum-sum segment and its application in bioinformatics, CIAA 2003, pp. 252-257.
- Lecture 3:
 - There are classic papers on the RAM model, where it is proved that the log-cost model is polynomially-equivalent to a Turing machine. See, for example, Stephen A. Cook and Robert A. Reckhow, [Time-bounded random access machines](#), *Journal of Computer Systems Science* 7 (1973), 354-375.
- Lecture 4:
 - For more about the discrete logarithm, this [survey](#) by Andrew Odlyzko, although a little out of date, is useful.
 - For more about the convex hull problem, see [this paper](#), which discusses an algorithm by Waterloo's Timothy Chan.
- Lecture 5:
 - Probably the best reference on algorithms to enumerate combinatorial objects is still the old book of Nijenhuis and Wilf, *Combinatorial Algorithms*.
 - Frank Ruskey has a page [here](#) to generate combinatorial objects of many types.
- Lecture 6:
 - It's actually possible to multiply two n-bit numbers in $O(n \log n \log \log n)$ time, but the algorithm is not simple. The book of Crandall and Pomerance, *Prime Numbers - A Computational Perspective*, is probably the best place to start.
 - In 2007, Martin Fürer found an even faster way to multiply numbers [here](#). But again, the algorithm is not simple.
 - Recently, Umans and Cohn found a [new approach](#) to matrix multiplication that may eventually lead to much faster algorithms.
- Lecture 8:
 - For more about the number of possible ways to parenthesize things, see [this paper by Guy and Selfridge](#) with the amusing title, "The Nesting and Roosting Habits of The Laddered Parenthesis".
 - Our dynamic programming algorithm for finding the optimal multiplication order uses $\Theta(n^3)$ time if there are n matrices. But this can be improved to $O(n \log n)$. See [this paper](#) by Hu and Shing in *SIAM J. Comput.* **13** (1984), 228-251.
- Lecture 10:
 - In 1997, Cummings and Smyth found a [better algorithm](#) for abelian squares than the one we presented here, in that it does not depend on the alphabet size, but unfortunately it is still $\Omega(n^2)$. The lower bound "proof" in the paper is apparently incorrect and we don't know how to repair it.

- Writing $x = 200a + 100b + 25c + 10d + 5e + f$ where a, b, c, d, e, f are integers ≥ 0 , and such that $a+b+c+d+e+f$ is minimized, is an example of an *integer program*. Unfortunately, integer programming in general is NP-hard, so there is no general fast way known to solve systems like this.
- [This paper](#) on change-making might amuse you.
- Lecture 11:
 - One of the best sources for papers about Egyptian fractions is the book of Richard Guy, *Unsolved Problems in Number Theory*. See Section D11.
 - For more about Egyptian fractions, see [David Eppstein's page](#).
- Lecture 12:
 - For the story of Finck and his little-known analysis of the Euclidean algorithm, see [here](#). (J. Shallit, Origins of the analysis of the Euclidean algorithm, *Historia Mathematica* **21** (1994), 401-419.)
- Lecture 12:
 - The [Mersenne twister](#) algorithm is a good way to generate pseudorandom numbers.
- Lectures 13-14:
 - For a different perspective on some of our lower bound arguments, see Eric Bach, Energy arguments in the theory of algorithms, **104** (1997), 831-837. Perhaps you can read it [here](#).
- Lectures 15-16:
 - Prim's algorithm dates from 1957. In 1986, Gabow, Galil, and Spencer found a faster algorithm for minimum spanning tree: it runs in $O(|E| \log^* (|E|/|V|))$ time, where \log^* is the slowly-growing function that gives the number of logs needed to reduce the argument to < 1 . See [here](#). In 1990, Fredman and Willard found a minimum-spanning tree algorithm that runs in linear time, but it depends on operations on the binary representation of the edge weights. See [here](#). In 1996, Karger, Klein, and Tarjan found a randomized algorithm that runs in $O(|E|)$ expected time and uses only comparisons between the edge weights. See [here](#).
- Lecture 17:
 - Shortest-path algorithms continue to attract attention. At the FOCS 2010 conference, Peres et al. presented a paper showing you can find shortest paths in n -vertex graphs where the edge weights are chosen uniformly and independently at random from $[0, 1]$ in about $O(n^2)$ expected time. See [here](#) for a video of the talk.
- Lectures 19-23:
 - The lower bound on extended regular expression equivalence is due to Meyer and Stockmeyer, and can be found [here](#).
 - [This page](#) maintained by Gerhard Woeginger is a sad list of people who have claimed to resolve P versus NP. Needless to say, none of their proofs have been accepted by the general computer science community.
 - The standard reference for NP-completeness is Garey and Johnson, *Computers and Intractability*. Johnson wrote a column for many years in the *Journal of Algorithms*, and they can be found [here](#).
- Lectures 23-24:
 - For more about Hilbert's 10th problem, see *Hilbert's Tenth Problem* by Yuri Matijasevich, Davis Centre library, QA242.M4213 1993.
 - For more about Gödel's proof, see *Gödel's Proof* by Nagel and Newman, Davis Centre library, QA9.65.N34 2001.
 - For a biography of Gödel, see Rebecca Goldstein, *Incompleteness*.
 - For more about the life of Alan Turing, see *Alan Turing: The Enigma* by Hodges, Davis Centre library, QA29.T87H64x 1983 or [this website](#) maintained by Hodges.

- For more about FRACTRAN, and the prime-producing program, see Guy's article, "Conway's Prime Producing Machine" *Math. Mag.* **56** (1983), 26-33. You may be able to access it [here](#).