CSC 611: Analysis of Algorithms

Lecture 5

Divide and Conquer

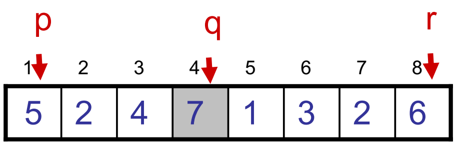
Divide-and-Conquer

- **Divide** the problem into a number of subproblems
  - Similar sub-problems of smaller size
- **Conquer** the sub-problems
  - Solve the sub-problems recursively
  - Sub-problem size small enough ⇒ solve the problems in straightforward manner
- **Combine** the solutions to the sub-problems
  - Obtain the solution for the original problem

Merge Sort Approach

- To sort an array $A[p \,..\, r]$:
- **Divide**
  - Divide the n-element sequence to be sorted into two subsequences of $n/2$ elements each
- **Conquer**
  - Sort the subsequences recursively using merge sort
  - When the size of the sequences is 1 there is nothing more to do
- **Combine**
  - Merge the two sorted subsequences

Merge Sort



Alg.: MERGE-SORT(A, p, r)

if $p < r$ ▷ Check for base case

  then $q \leftarrow \lfloor (p + r)/2 \rfloor$ ▷ Divide

    MERGE-SORT(A, p, q) ▷ Conquer

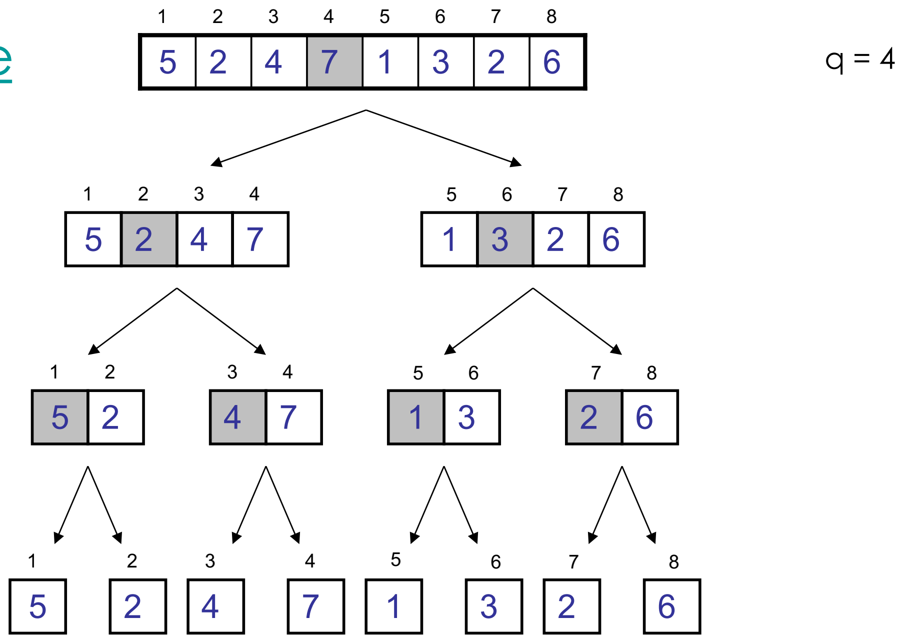
    MERGE-SORT(A, $q + 1$, r) ▷ Conquer

    MERGE(A, p, q, r) ▷ Combine

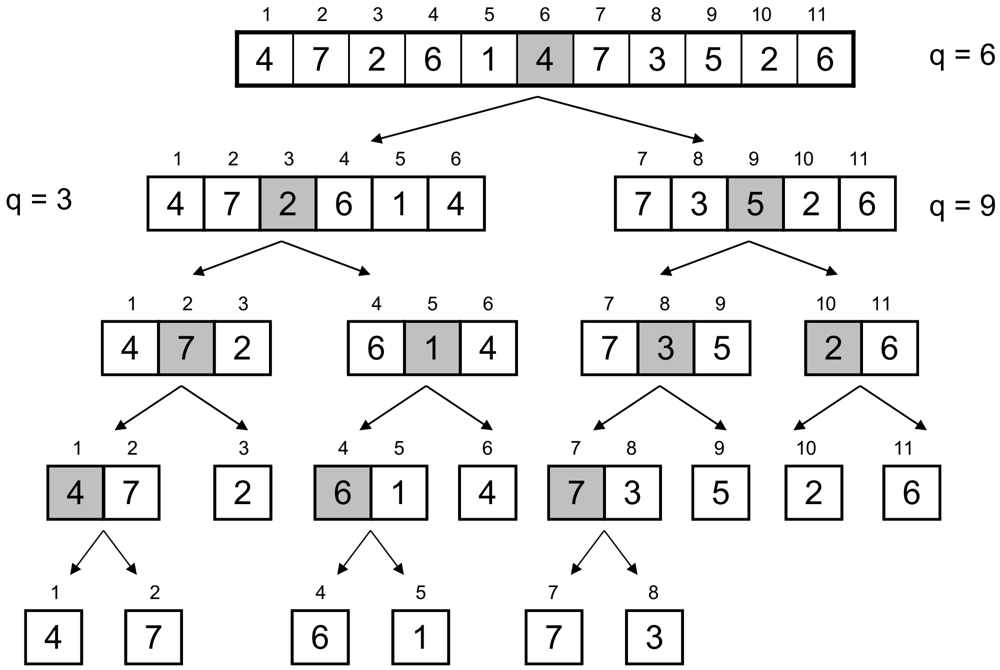
- Initial call: MERGE-SORT(A, 1, n)

Example – n Power of 2

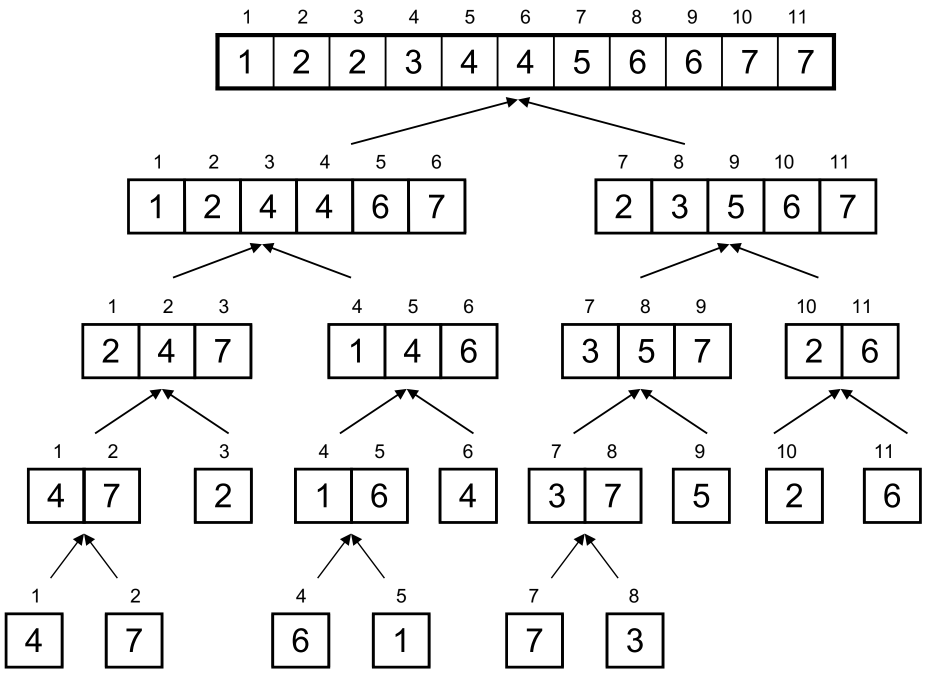
Example



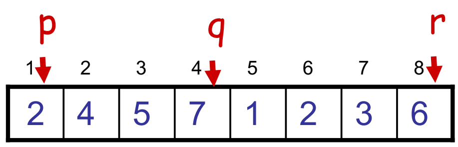
Example – n Not a Power of 2



Example – n Not a Power of 2



Merging

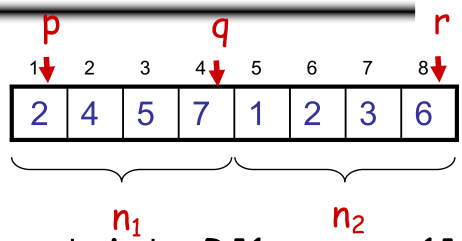


- **Input:** Array *A* and indices **p, q, r** such that **p ≤ q < r**
  - Subarrays *A*[**p . . q**] and *A*[**q + 1 . . r**] are sorted
- **Output:** One single sorted subarray *A*[**p . . r**]

Merging

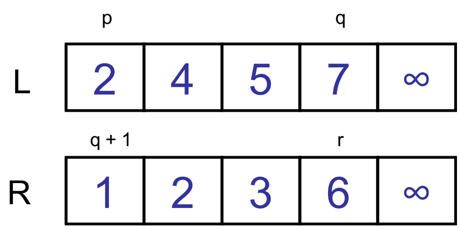
- Idea for merging:
  - Two piles of sorted cards
    - Choose the smaller of the two top cards
    - Remove it and place it in the output pile
  - Repeat the process until one pile is empty
  - Take the remaining input pile and place it face-down onto the output pile

Merge - Pseudocode



Alg.: MERGE(A, p, q, r)

1. Compute n_1 and n_2
2. Copy the first n_1 elements into $L[1 \mathrel{..} n_1 + 1]$ and the next n_2 elements into $R[1 \mathrel{..} n_2 + 1]$
3. $L[n_1 + 1] \leftarrow \infty$; $R[n_2 + 1] \leftarrow \infty$
4. $i \leftarrow 1$; $j \leftarrow 1$
5. **for** $k \leftarrow p$ **to** r
6. **do if** $L[i] \leq R[j]$
7. **then** $A[k] \leftarrow L[i]$
8. $i \leftarrow i + 1$
9. **else** $A[k] \leftarrow R[j]$
10. $j \leftarrow j + 1$



Running Time of Merge

- Initialization (copying into temporary arrays):
  - $\Theta(n_1 + n_2) = \Theta(n)$
- Adding the elements to the final array (the **for** loop):
  - n iterations, each taking constant time $\Rightarrow \Theta(n)$
- Total time for Merge:
  - $\Theta(n)$

Analyzing Divide and Conquer Algorithms

- The recurrence is based on the three steps of the paradigm:
  - $T(n)$ – running time on a problem of size n
  - **Divide** the problem into a subproblems, each of size n/b: takes $D(n)$
  - **Conquer** (solve) the subproblems: takes $aT(n/b)$
  - **Combine** the solutions: takes $C(n)$

$$T(n) = \begin{cases} \Theta(1) & \text{if } n \leq c \\ aT(n/b) + D(n) + C(n) & \text{otherwise} \end{cases}$$

MERGE – SORT Running Time

- **Divide:**
  - compute q as the average of p and r: $D(n) = \Theta(1)$
- **Conquer:**
  - recursively solve 2 subproblems, each of size $n/2$ $\Rightarrow 2T(n/2)$
- **Combine:**
  - MERGE on an n-element subarray takes $\Theta(n)$ time $\Rightarrow C(n) = \Theta(n)$

$$T(n) = \begin{cases} \Theta(1) & \text{if } n = 1 \\ 2T(n/2) + \Theta(n) & \text{if } n > 1 \end{cases}$$

Solve the Recurrence

$$T(n) = \begin{cases} c & \text{if } n = 1 \\ 2T(n/2) + cn & \text{if } n > 1 \end{cases}$$

Use Master's Theorem:

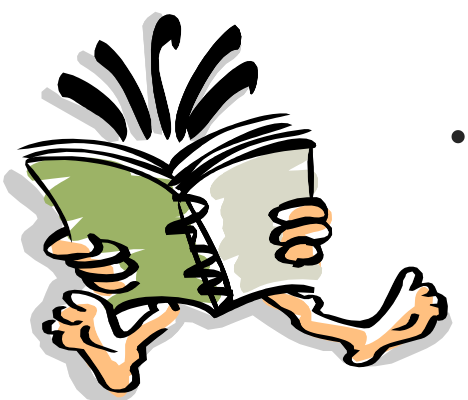
Compare n with $f(n) = cn$

Case 2: $T(n) = \Theta(n\lg n)$

Merge Sort - Discussion

- Running time insensitive of the input
- Advantages:
  - Guaranteed to run in $\Theta(n\lg n)$
- Disadvantage
  - Requires extra space ≈N
- Applications
  - Maintain a large ordered data file
  - How would you use Merge sort to do this?

Readings



- Chapter 7