

Review (Contd)

Subset problem:

Items: $1, \dots, n$
Weights: w_1, \dots, w_n

Find a subset of maximal weight, under W .

We want a subset $\{i_1, \dots, i_k\}$ s.t. $\sum_{j=1}^k w_{i_j} \leq W$ and

Either item 1 is a part of Opt

$\sum_j w_{i_j}$ is maximized & maximal

$\hookrightarrow \text{Opt}([2, n], W - w_1)$

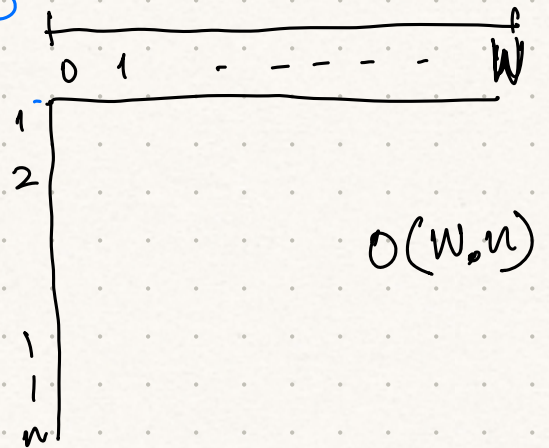
or not.

$\hookrightarrow \text{Opt}([2, n], W)$

2-param DP

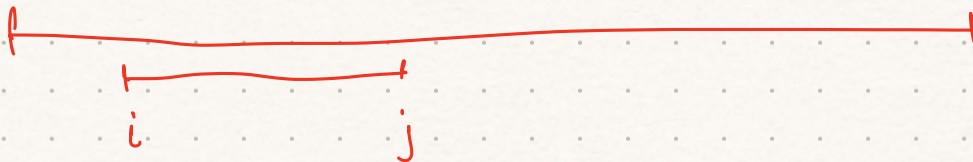
We assumed that the wts are int.

$\text{Opt}([i, n], \tilde{W})$



Qn: Is there a subset that sums up to A .

Contiguous



Array Partitioning Problem:
(Sundar Vishwanathan's notes, IITB)



Contiguous/partition

Input is a Penalty matrix $n \times n$

Each partition involves a cost c .

Want: Divide the array into parts s.t each part has contiguous elements and total partition cost is minimized.

→ If there is only one part → $P(1, n)$

→ Sum of Penalties + cost of parts

→ $P(1, i) + c + P(i+1, n)$

→ $i_1, \dots, i_k \rightarrow c \cdot N + P(1, i_1) + P(i_1+1, i_2) + \dots + P(i_k+1, n)$

$$\text{Cost}([1, n]) = \min \left\{ \min_k \left\{ \text{Cost}([1, k]) + c + \text{Cost}([k+1, n]) \right\} \right\}$$

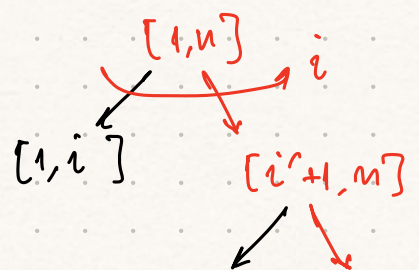
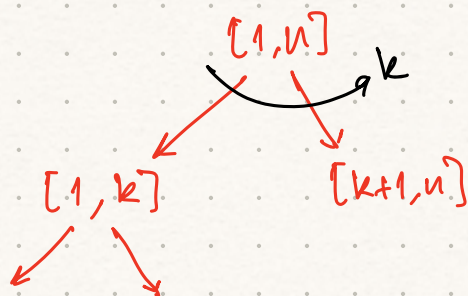
↙ Say there is an i^* in opt part s.t $[1, i^*]$ is a part ~ $[i^*+1, k^*], [k^*+1, n]$ ← space $O(n^2)$ time $O(n^2 \cdot n)$

$$\text{MinCost}([1, n]) = \min_{1 \leq i \leq n} \left\{ \min \left\{ \underline{P[1, i]} + c + \text{MinCost}[i+1, n] \right\} \right\},$$

Space $O(n)$
time $O(n^2)$

$$\text{MinCost}([i, n])$$

$$= \min \left(\left\{ P[i, k] + c + \text{MinCost}[k+1, n] \mid i \leq k < n \right\} \cup \{P[i, n]\} \right)$$



Interval Scheduling:

Before

← Max no. of jobs that can be scheduled w/o overlaps.

Jobs/Reqs J_1, \dots, J_n

Start s_1 s_n

Finish f_1 f_n

Now

← Want: The duration of scheduling to be maximized as well.

→ Sort your reqs in incr order of start times.

J_1, \dots, J_n index of first non-overlapping job

$|J_1| + \text{Opt}(\text{First}[1], n)$

$\text{Opt}(2, n)$

$$\text{Opt}(1, n) = \max \begin{cases} \text{Opt}(2, n) \\ |f_1 - s_1| + \text{Opt}(\text{First}[1], n) \end{cases}$$

$$\text{Opt}(i, n) = \max \begin{cases} \text{Opt}(i+1, n) \\ |f_i - s_i| + \text{Opt}(\text{First}[i], n) \end{cases}$$