State-Space Formulation

- Intelligent agents: problem solving as search
- Search consists of
 - state space
 - operators
 - start state
 - goal states
- The search graph
- A Search Tree is an effective way to represent the search process
- There are a variety of search algorithms, including
 - Depth-First Search
 - Breadth-First Search
 - Others which use heuristic knowledge (in future lectures)

Uninformed search strategies

uninformed (or blind) search:

While searching you have no clue whether one non-goal state is better than any other. Your search is blind. You don't know if your current exploration is likely to be fruitful.

common blind strategies:

- Breadth-first search
- Uniform-cost search
- Depth-first search
- Iterative deepening search

General Tree Search

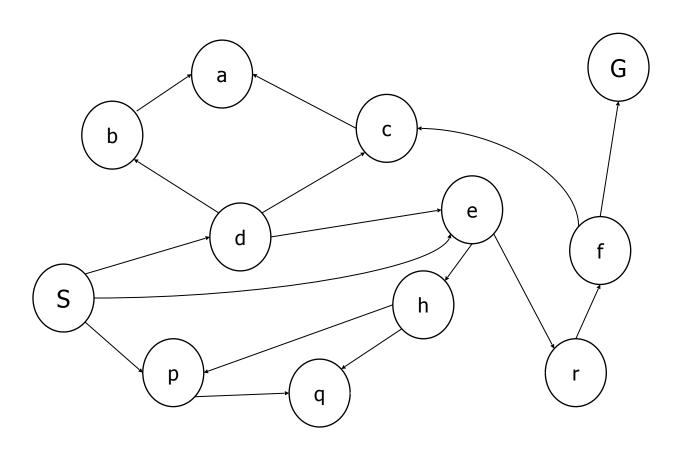
```
function TREE-SEARCH( problem, strategy) returns a solution, or failure initialize the search tree using the initial state of problem loop do

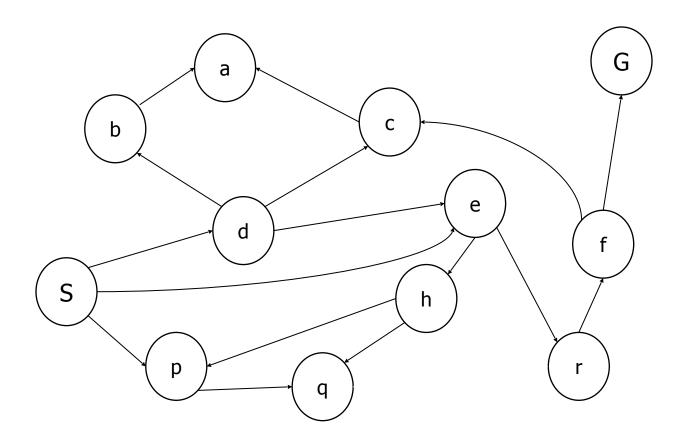
if there are no candidates for expansion then return failure choose a leaf node for expansion according to strategy

if the node contains a goal state then return the corresponding solution else expand the node and add the resulting nodes to the search tree end
```

- Important details:
 - Fringe (frontier)
 - Expansion
 - Exploration strategy
- Main question: which fringe nodes to explore?

Graph representation of the landscape explored

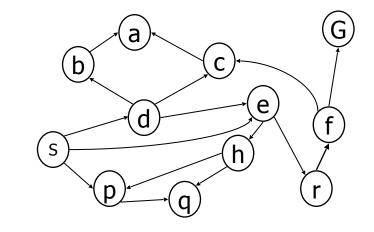


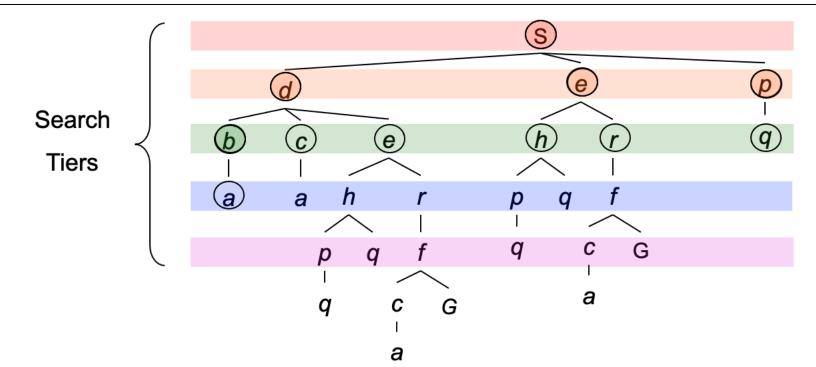


Breadth-First Search

Strategy: expand a shallowest node first

Implementation: Fringe is a FIFO queue





```
function BREADTH-FIRST-SEARCH(problem) returns a solution, or failure

node ← a node with STATE = problem.INITIAL-STATE, PATH-COST = 0

if problem.GOAL-TEST(node.STATE) then return SOLUTION(node)

frontier ← a FIFO queue with node as the only element

explored ← an empty set

loop do

if EMPTY?(frontier) then return failure

node ← POP(frontier) /* chooses the shallowest node in frontier */

add node.STATE to explored

for each action in problem.ACTIONS(node.STATE) do

child ← CHILD-NODE(problem, node, action)

if child.STATE is not in explored or frontier then

if problem.GOAL-TEST(child.STATE) then return SOLUTION(child)

frontier ← INSERT(child, frontier)
```

Figure 3.11 Breadth-first search on a graph.

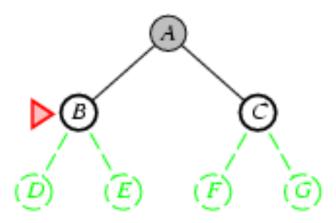
- Expand shallowest unexpanded node
- Fringe: nodes waiting in a queue to be explored, also called OPEN
- Implementation:
 - *fringe* is a first-in-first-out (FIFO) queue, i.e., new successors go at end of the queue.

Is A a goal state? $(B) \qquad (C)$

- Expand shallowest unexpanded node
- Implementation:
 - fringe is a FIFO queue, i.e., new successors go at end

Expand: fringe = [B,C]

Is B a goal state?

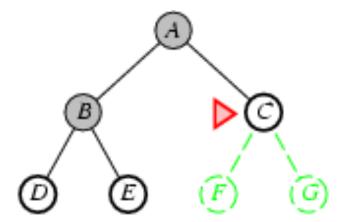


Expand shallowest unexpanded node

- Implementation:
 - fringe is a FIFO queue, i.e., new successors go at end

Expand: fringe=[C,D,E]

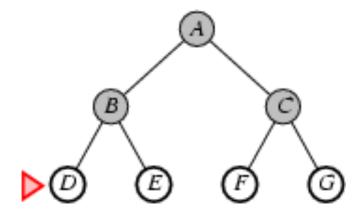
Is C a goal state?



- Expand shallowest unexpanded node
- Implementation:
 - fringe is a FIFO queue, i.e., new successors go at end

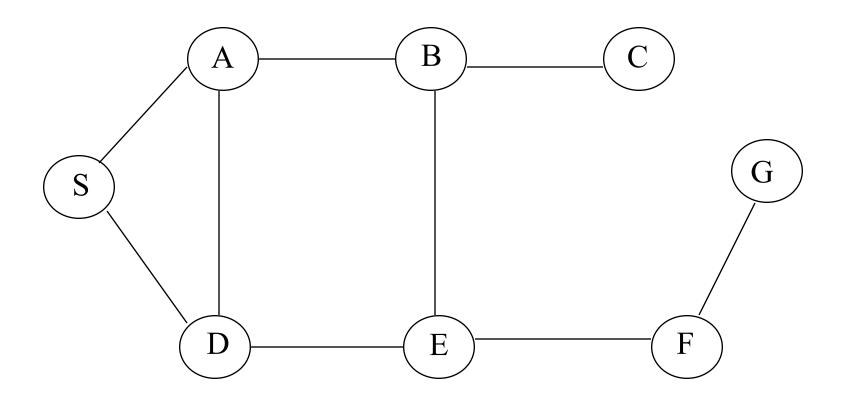
Expand: fringe=[D,E,F,G]

Is D a goal state?



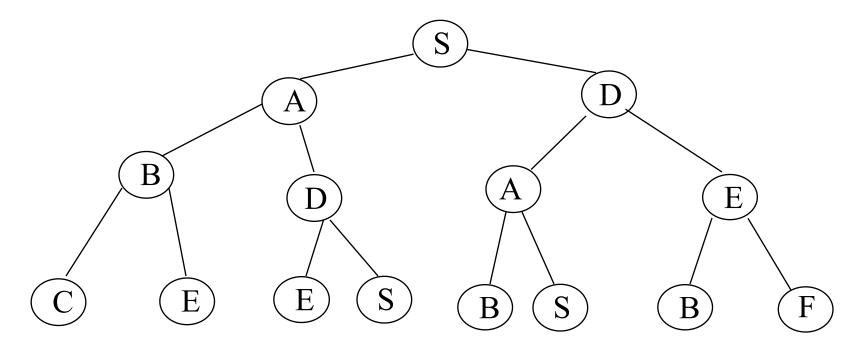
BFS for 8 puzzle

Example: Map Navigation



S = start, G = goal, other nodes = intermediate states, links = legal transitions

Breadth First Search Tree (BFS)



Here BFS is implemented as a tree search with only parent node not added as a child node.

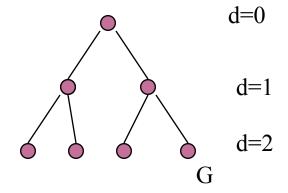
What is the Complexity of Breadth-First Search?

Time Complexity

- assume (worst case) that there is 1
 goal leaf at the RHS
- so BFS will expand all nodes

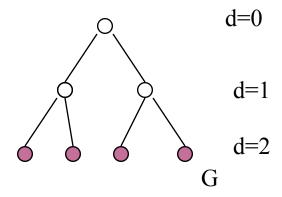
=
$$1 + b + b^2 + \dots + b^d$$

= $O(b^{d+1})$



Space Complexity

- how many nodes can be in the queue (worst-case)?
- at depth d there are bd unexpanded nodes in the Q = O (bd)
- Time and space of number of generated nodes is O (b^{d+1})



Examples of Time and Memory Requirements for tree search version of Breadth-First Search

Depth of Solution	Nodes Expanded	Time	Memory
0	1	1 millisecond	100 bytes
2	111	0.1 seconds	11 kbytes
4	11,111	11 seconds	1 megabyte
8	108	31 hours	11 giabytes
12	1012	35 years	111 terabytes

Assuming b=10, 1000 nodes/sec, 100 bytes/node

Breadth-First Search (BFS) Properties

- Complete (with find a solution in a finite number of steps if one exists)
- Solution Length: optimal (assuming unit cost per move)
- (Can) expand each node once (if checks for duplicates)
- Search Time: O(bd) which is the size of the state space
- Memory Required: O(bd)
- Drawback: requires space proportional to the statespace (Search time is unavoidable)