

CS4450/7450  
LYAHGG (Chapter 3)  
Types and Type Classes

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## Let the type be your guide

What is the type of `foo`?

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```
ghci> foo "A Connecticut Yankee in King  
        Arthur's Court"  
"ACYKAC"
```

# Type Systems

Haskell has “static types with inference”

- **Type Checking:** given an expression  $e$  and a type  $t$ , check whether  $e :: t$ . E.g.,

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- **Static Types.** a type system for which the types of expressions are known at *compile-time*. I.e., the type of every expression is known by inspecting its code—and not by running it.

# Type Variables

Reintroducing what we called “parametric polymorphism”

The following type means that, for all types `a` and `b`, the function `fst` can be applied.

```
ghci> :t fst  
fst :: (a, b) -> a
```



# Type Instances

Given:

```
ghci> :t fst
fst :: (a, b) -> a
```

“Instances” of  $(a, b) \rightarrow a$  determine how `fst` can be applied:

```
(Int, Char) -> Int           -- fst (99, 'A')
([Char], Float) -> [Char]    -- fst ("Hey ", 3.14)
(Int -> Int, Bool) -> Int -> Int -- fst (id, True)
                               ⋮
```

`fst`, `fst`, `fst` all refer to **the same code**.

# Type Classes

The **following** is a *type constraint*:

```
ghci> :t (==)
(==)  :: (Eq a) => a -> a -> Bool
```

It means that `(==)` can be applied only at types in the `Eq` class.

There are many predefined classes in Haskell, including `Ord`, `Show`, `Enum`, `Num`, etc.