

CS4450: Simple While Language

Bill Harrison

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Announcements

- ▶ Today: While language static semantics + transition semantics

State Machines

Transition Semantics:

- ▶ Define the meaning of a language with transition rules.
Execute input program p in state m_0 :
- ▶ $(p, m_0) \rightarrow (p_1, m_1) \rightarrow \dots \rightarrow (p_n, m_n) \rightarrow \dots$

The Simple Imperative Language with Loops

(Abstract Syntax of While)

$I \in \text{Identifier}$

$N \in \text{Numeral}$

$B ::= \text{true} \mid \text{false} \mid B \text{ and } B \mid B \text{ or } B \mid \text{not } B \mid E < E \mid E = E$

$E ::= N \mid I \mid E + E \mid E * E \mid E - E \mid - E$

$C ::= \text{skip} \mid C ; C \mid I := E \mid$
 $\quad \text{if } B \text{ then } C \text{ else } C \text{ fi} \mid \text{while } B \text{ do } C \text{ od}$

While Abstract Syntax in Haskell

```
data E = N Number
      | I Ident
      | Add E E | Mult E E | Subt E E | Inv E
```

```
data B = T | F | Conj B B | Disj B B
      | Negt B | LTC E E | EQL E E
```

```
data C = Nop
      | Seq C C
      | IfElse B C C | While B C
```

The Memory

(Memory maps I to Z)

$lookup\ m\ i = \langle \text{current value of } i \rangle$

$m[i \mapsto v] = \langle \text{new memory s.t. } i \text{ is bound to } v \rangle$

$$lookup\ m[i \mapsto v]\ j = \begin{cases} v & i \text{ is } j \\ lookup\ m\ j & \text{otherwise} \end{cases}$$

While Semantics in Haskell

```
type Ident  = String
type Number = Int
type Memory = [(Ident, Number)]

--
-- memory look-up
--
lkup :: Memory -> Ident -> Number
lkup ((x, n):ms) x' = if x == x' then n else lkup ms x'
lkup [ ] _          = error (x ++ " is unbound!")
```

E and B semantics

$$\begin{aligned} \text{ev}_E \ n \ m &= n \\ \text{ev}_E \ i \ m &= \text{lookup } i \ m \\ \text{ev}_E \ -e \ m &= -(\text{ev}_E \ e \ m) \end{aligned}$$

...

$$\begin{aligned} \text{ev}_B \ \mathbf{true} \ m &= \text{true} \\ \text{ev}_B \ (e_1 = e_2) \ m &= (\text{ev}_E \ e_1 \ m == \text{ev}_E \ e_2 \ m) \\ &\dots \end{aligned}$$

E and B semantics

$$\begin{aligned} \text{ev}_E \ n \ m &= n \\ \text{ev}_E \ i \ m &= \text{lookup } i \ m \\ \text{ev}_E \ -e \ m &= -(\text{ev}_E \ e \ m) \\ &\dots \\ \text{ev}_B \ \mathbf{true} \ m &= \text{true} \\ \text{ev}_B \ (e_1 = e_2) \ m &= (\text{ev}_E \ e_1 \ m == \text{ev}_E \ e_2 \ m) \\ &\dots \end{aligned}$$

Question: what are the types of ev_E and ev_B ?

Transition Semantics of While

$$(i := e, m) \rightarrow m[i \mapsto \text{ev}_E \ e \ m] \qquad (\text{skip}, m) \rightarrow m$$

$$\frac{(c_1, m) \rightarrow (c'_1, m')}{(c_1; c_2, m) \rightarrow (c'_1; c_2, m')} \qquad \frac{(c_1, m) \rightarrow m'}{(c_1; c_2, m) \rightarrow (c_2, m')}$$

$$\frac{\text{ev}_B \ b \ m = \text{true}}{(\text{if } b \text{ then } c_1 \text{ else } c_2 \text{ fi}, m) \rightarrow (c_1, m)}$$

$$\frac{\text{ev}_B \ b \ m = \text{false}}{(\text{if } b \text{ then } c_1 \text{ else } c_2 \text{ fi}, m) \rightarrow (c_2, m)}$$

Transition Semantics of While (cont'd)

$$\frac{ev_B \ b \ m = true}{(\mathbf{while} \ b \ \mathbf{do} \ c \ \mathbf{od}, m) \rightarrow (c; \mathbf{while} \ b \ \mathbf{do} \ c \ \mathbf{od}, m)}$$

$$\frac{ev_B \ b \ m = false}{(\mathbf{while} \ b \ \mathbf{do} \ c \ \mathbf{od}, m) \rightarrow m}$$

Exercise

Formulate the transition semantics for While in Haskell.

Hint:

1. Define each of E, B, and C as **data** declarations;
2. Define ev_E and ev_B as Haskell functions;
3. Define the Memory data type next;
4. Finally, define the transitions for C as a function of type:
 $(C, Memory) \rightarrow Memory$.