# CS4450/7450 AoPL, Chapter 6: Computational Strategies Principles of Programming Languages

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#### **Announcements**

• We're continuing with William Cook's online textbook, *Anatomy of Programming Languages*. It is available here. We're in Chapter 6.

#### Outline for section 1

Error Checking

- 2 Mutable State
  - Pure Functional Operations on Memory

## Outline for section 2

Error Checking

- Mutable State
  - Pure Functional Operations on Memory

#### Mutable State

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Mutable state means that the state of a program changes or mutates: that a variable can be assigned a new value or a part of a data structure can be modified.

Ex: The values of x and i change over time:

```
x = 1;
for (i = 2; i <= 5; i = i + 1) {
  x = x * i;
}
```

#### Addresses

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Operation	Meaning
mutable e	Creates a mutable cell with initial value given by e
@e	Accesses contents stored at address e
a = e	Updates contents at address a to be value of expression e

# Example

```
x = mutable 1;
for (i = mutable 2; @i <= 5; i = @i + 1) {
  x = @x * @i;
}</pre>
```

# New Values and Expressions

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Step	Memory
start	[]

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start	[]
x = mutable 1;	[1]

Step	Memory
start	
x = mutable 1;	[1]
<pre>i = mutable 2;</pre>	[1,2]

Step	Memory
start	[]
x = mutable 1;	[1]
i = mutable 2;	[1,2] [2,2]
x = @x * @i;	[2,2]

Step	Memory
start	
x = mutable 1;	[1]
i = mutable 2;	[1,2]
x = @x * @i;	[2,2]
i = @i + 1;	[2,3]

Step				Memory
start				[]
x =	mut	ak	ole 1;	[1]
<u>i</u> =	mut	ak	ole 2;	[1,2]
x =	@ X	*	@i;	[2,2]
i =	@i	+	1;	[2,3]
x =	@ X	*	@i;	[6,3]
i =	@i	+	1;	[6,4]
x =	@ X	*	@i;	[24,4]
i =	@i	+	1;	[24,5]
x =	@ X	*	@i;	[120,5]
i =	@i	+	1;	[120,6]

# Operations on Memory

## Accessing Memory:

```
access i mem = mem !! i
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  let (before, _ : after) = splitAt addr mem in
  before ++ [val] ++ after
```

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    before ++ [val] ++ after

ghci> :t splitAt
  splitAt :: Int -> [a] -> ([a], [a])
ghci> splitAt 3 "abcdefg"
  ("abc", "defg")
ghci> splitAt 0 "abcdefg"
  ("", "abcdefg")
```

• Currently:

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evaluate :: Exp -> Env -> Value
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evaluate :: Exp -> Env -> Memory -> (Value, Memory)
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evaluate :: Exp -> Env -> Memory -> (Value, Memory)
```

Generalizing:

```
type Stateful a = Memory -> (a, Memory)
evaluate :: Exp -> Env -> Memory -> Stateful Value
```

### **Evaluation Rules**

```
evaluate (Mutable e) env mem =
  let
          (ev, mem') = evaluate e env mem
  in
          (AddressV (length mem'), mem' ++ [ev])
```

#### **Evaluation Rules**

# Evaluation Rules (cont'd)

```
evaluate (Assign a e) env mem =
  let
     (AddressV i, mem') = evaluate a env mem
  in
     let.
        (ev, mem'') = evaluate e env mem'
     in
        (ev, update i ev mem'')
evaluate (Binary op a b) env mem =
  let.
     (av, mem') = evaluate a env mem
  in
     let.
        (bv, mem'') = evaluate b env mem'
     in
        (binary op av bv, mem'')
```

# A pattern is forming...

• Recall:

```
evaluate (Mutable e) env mem =
  let (ev, mem') = evaluate e env mem
  in (AddressV (length mem'), mem' ++ [ev])
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evaluate (Mutable e) env mem =
  let (ev, mem') = evaluate e env mem
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• The pattern:

```
\ mem ->
  let
     (val, mem') = first-part mem
in
     next-part val mem'
```

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Recall:

```
evaluate (Mutable e) env mem =
  let (ev, mem') = evaluate e env mem
  in (AddressV (length mem'), mem' ++ [ev])
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• The pattern:

```
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let
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```

Generalizing as a higher-order function: