

CS4450/7450  
AoPL, Chapter 6: Computational Strategies  
Principles of Programming Languages

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# Announcements

- We're continuing with William Cook's online textbook, *Anatomy of Programming Languages*. It is available [here](#). We're in Chapter 6.

# Outline for section 1

## 1 Error Checking

## 2 Mutable State

- Pure Functional Operations on Memory

# Outline for section 2

1 Error Checking

2 Mutable State

- Pure Functional Operations on Memory

# Mutable State

## Mutable State

Mutable state means that the state of a program changes or mutates: that a variable can be assigned a new value or a part of a data structure can be modified.

Ex: The values of `x` and `i` change over time:

```
x = 1;
for (i = 2; i <= 5; i = i + 1) {
    x = x * i;
}
```

# Addresses

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Operation	Meaning
<code>mutable e</code>	Creates a mutable cell with initial value given by <code>e</code>
<code>@e</code>	Accesses contents stored at address <code>e</code>
<code>a = e</code>	Updates contents at address <code>a</code> to be value of expression <code>e</code>

# Example

```
x = mutable 1;  
for (i = mutable 2; @i <= 5; i = @i + 1) {  
    x = @x * @i;  
}
```



# New Values and Expressions

```
data Value = IntV  Int
           | BoolV Bool
           | ClosureV String Exp Env
           | AddressV Int           -- new
deriving (Eq, Show)

type Memory = [Value]
```

# New Values and Expressions

```
data Value = IntV  Int
           | BoolV Bool
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           | AddressV Int          -- new
deriving (Eq, Show)

type Memory = [Value]

data Exp = ...
        | Mutable    Exp          -- mutable e
        | Access     Exp          -- @a
        | Assign     Exp Exp      -- a = e
```

# Illustrating Memory Operations

Step	Memory
start	[]

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<code>x = mutable 1;</code>	[1]

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start	[]
x = mutable 1;	[1]
i = mutable 2;	[1,2]

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x = @x * @i;	[2,2]

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Step	Memory
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x = mutable 1;	[1]
i = mutable 2;	[1,2]
x = @x * @i;	[2,2]
i = @i + 1;	[2,3]

# Illustrating Memory Operations

Step	Memory
start	[]
x = mutable 1;	[1]
i = mutable 2;	[1,2]
x = @x * @i;	[2,2]
i = @i + 1;	[2,3]
x = @x * @i;	[6,3]
i = @i + 1;	[6,4]
x = @x * @i;	[24,4]
i = @i + 1;	[24,5]
x = @x * @i;	[120,5]
i = @i + 1;	[120,6]



# Operations on Memory

Accessing Memory:

```
access i mem = mem !! i
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access i mem = mem !! i
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## Updating Memory:

```
update :: Int -> Value -> Memory -> Memory
```

```
update addr val mem =
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```
  let (before, _ : after) = splitAt addr mem in  
    before ++ [val] ++ after
```

# Operations on Memory

## Accessing Memory:

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access i mem = mem !! i
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## Updating Memory:

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update :: Int -> Value -> Memory -> Memory
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update addr val mem =
```

```
    let (before, _ : after) = splitAt addr mem in  
    before ++ [val] ++ after
```

```
ghci> :t splitAt
```

```
    splitAt :: Int -> [a] -> ([a], [a])
```

```
ghci> splitAt 3 "abcdefg"
```

```
    ("abc", "defg")
```

```
ghci> splitAt 0 "abcdefg"
```

```
    ("", "abcdefg")
```

# Whither evaluate?

- Currently:

`evaluate :: Exp -> Env -> Value`

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- Now a stateful computation:

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evaluate :: Exp -> Env -> Memory -> (Value, Memory)
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- Now a stateful computation:

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evaluate :: Exp -> Env -> Memory -> (Value, Memory)
```

- Generalizing:

```
type Stateful a = Memory -> (a, Memory)  
evaluate :: Exp -> Env -> Memory -> Stateful Value
```

# Evaluation Rules

```
evaluate (Mutable e) env mem =  
  let  
    (ev, mem') = evaluate e env mem  
  in  
    (AddressV (length mem'), mem' ++ [ev])
```



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```
evaluate (Access a) env mem =  
  let  
    (AddressV i, mem') = evaluate a env mem  
  in  
    (access i mem', mem')
```

# Evaluation Rules

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evaluate (Mutable e) env mem =  
  let  
    (ev, mem') = evaluate e env mem  
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    (AddressV (length mem'), mem' ++ [ev])
```

```
evaluate (Access a) env mem =  
  let  
    (AddressV i, mem') = evaluate a env mem  
  in  
    (access i mem', mem')
```

## Evaluation Rules (cont'd)

```
evaluate (Assign a e) env mem =  
  let  
    (AddressV i, mem') = evaluate a env mem  
  in  
    let  
      (ev, mem'') = evaluate e env mem'  
    in  
      (ev, update i ev mem'')
```

```
evaluate (Binary op a b) env mem =  
  let  
    (av, mem') = evaluate a env mem  
  in  
    let  
      (bv, mem'') = evaluate b env mem'  
    in  
      (binary op av bv, mem'')
```

## A pattern is forming...

- Recall:

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  let (ev, mem') = evaluate e env mem  
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- Recall:

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evaluate (Mutable e) env mem =  
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- The pattern:

```
\ mem ->  
  let  
    (val, mem') = first-part mem  
  in  
    next-part val mem'
```

## A pattern is forming...

- Recall:

```
evaluate (Mutable e) env mem =
  let (ev, mem') = evaluate e env mem
  in (AddressV (length mem'), mem' ++ [ev])
```

- The pattern:

```
\ mem ->
  let
    (val, mem') = first-part mem
  in
    next-part val mem'
```

- Generalizing as a higher-order function:

```
(>>=) :: Stateful a -> (a -> Stateful b) -> Stateful b
first-part >>= next-part =
  \ mem ->
    let
      (val, mem') = first-part mem
    in
      next-part val mem'
```









