MINORS OF TREE DISTANCE MATRICES

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ABSTRACT. We prove an identity that relates the principal minors of the distance matrix of a tree, on one hand, to a combinatorial expression involving counts of rooted spanning forests of the underlying tree. This generalizes a result of Graham and Pollak. A variant of this identity applies to the case of edge-weighted trees.

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1. Introduction

Suppose G = (V, E) is a tree with n vertices. Let D denote the distance matrix of G. In [4], Graham and Pollak proved that

eq:full-det

(1)
$$\det D = (-1)^{n-1} 2^{n-2} (n-1).$$

This identity is remarkable in that the result does not depend on the tree structure, beyond the number of vertices. The identity (I) was motivated by a problem in data communication, and inspired much further research on distance matrices.

inspired much further research on distance matrices. The main result of this paper is to generalize (I) by replacing det D with any of its principal minors. For a subset $S \subset V(G)$, let D[S] denote the submatrix consisting of the S-indexed rows and columns of D.

thm:main

Theorem 1. Suppose G is a tree with n vertices, and distance matrix D. Let $S \subset V(G)$ be a nonempty subset of vertices. Then

(2)
$$\det D[S] = (-1)^{|S|-1} 2^{|S|-2} \left((n-1) \kappa(G; S) - \sum_{\mathcal{F}_2(G; S)} (\deg^o(F, *) - 2)^2 \right),$$

where $\kappa(G; S)$ is the number of S-rooted spanning forests of G, $\mathcal{F}_2(G; S)$ is the set of (S, *)-rooted spanning forests of G, and $\deg^o(F, *)$ denotes the outdegree of the *-component of F.

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For definitions of (S,*)-rooted spanning forests and other terminology, see Section 2. Note that the quantity $\deg^o(F,*)$ satisfies the bounds

$$1 \le \deg^o(F, *) \le |S|.$$

When S = V is the full vertex set, the set of V-rooted spanning forests is a singleton, consisting of the subgraph with no edges, so $\kappa(G; V) = 1$; and moreover the set $\mathcal{F}_2(G; V)$ of (V, *)-rooted spanning forests is empty. Thus (2) recovers the Graham–Pollak identity (II) when S = V.

1.1. Weighted trees. If $\{\alpha_e : e \in E\}$ is a collection of positive edge weights, the α -distance matrix D_{α} is defined by setting the (u, v)-entry to the sum of the weights α_e along the unique path from u to v. The following weighted version of (II) is satisfied by the weighted distance matrix,

eq:w-full-det

(3)
$$\det D_{\alpha} = (-1)^{n-1} 2^{n-2} \sum_{e \in E} \alpha_e \prod_{e \in E} \alpha_e$$

 $(3) \qquad \det D_{\alpha} = (-1)^{n-1} 2^{n-2} \sum_{e \in E} \alpha_e \prod_{e \in E} \alpha_e,$ which was proved by Bapat–Kirkland–Neumann [I]. The weighted identity (3) reduces to (II) when taking all unit weights, $\alpha_e = 1$. We also prove the following weighted version of our main theorem.

thm:w-main

Theorem 2. Suppose G = (V, E) is a finite, weighted tree with edge weights $\{\alpha_e : e \in E\}$, and corresponding weighted distance matrix $D = D_{\alpha}$. For any nonempty subset $S \subset V$, we have

eq:w-main

(4)
$$\det D[S] = (-1)^{|S|-1} 2^{|S|-2} \left(\sum_{E(G)} \alpha_e \sum_{\mathcal{F}_1(G;S)} w(\overline{T}) - \sum_{\mathcal{F}_2(G;S)} (\deg^o(F,*) - 2)^2 w(F) \right).$$

where $\mathcal{F}_1(G;S)$ is the set of S-rooted spanning forests of G, $\mathcal{F}_2(G;S)$ is the set of (S,*)-rooted spanning forests of G, $w(\overline{T})$ and w(F) denote the α -weights of the forests T and F, and $\deg^o(F,*)$ is the outdegree of the *-component of F, as above.

Theorem $\frac{\text{thm:w-main}}{2 \text{ reduces}}$ to Theorem $\frac{\text{thm:main}}{1 \text{ when}}$ taking all unit weights, $\alpha_e = 1$.

1.2. Applications. Suppose we fix a tree distance matrix D. It is natural to ask, how do the expressions det D[S] vary as we vary the vertex subset S? To our knowledge there is no nice behavior to answer this question, but as S varies there is nice behavior of the ratios $\det D[S]/\cot D[S]$ which we describe here.

Given a matrix A, let cof A denote the sum of cofactors of A, i.e.

$$\operatorname{cof} A = \sum_{i=1}^{|S|} \sum_{j=1}^{|S|} (-1)^{i+j} \det A_{ij}.$$

If A is invertible, then cof A is related to the sum of entries of the matrix inverse A^{-1} by a factor of det A, i.e. cof $A = (\det A)(\mathbf{1}^{\mathsf{T}}A^{-1}\mathbf{1})$. In [2], Bapat and Sivasubramanian showed the following identity for the sum of cofactors of a distance submatrix D[S] of a tree.

eq:cof-trees

(5)
$$\operatorname{cof} D[S] = (-2)^{|S|-1} \sum_{T \in \mathcal{F}_1(G;S)} w(\overline{T}).$$

Using the Bapat–Sivasubramanian identity (b), an immediate corollary to Theorem 2 is the following result: Suppose G = (V, E) is a finite, weighted tree with edge weights $\{\alpha_e : e \in E\}$. Let $D = D_{\alpha}$ denote the weighted distance matrix of G. For any nonempty subset $S \subset V$, we have

eq:det-cof

(6)
$$\frac{\det D[S]}{\cot D[S]} = \frac{1}{2} \left(\sum_{e \in E} \alpha_e - \frac{\sum_{F \in \mathcal{F}_2(G;S)} w(F) k(F,*)^2}{\sum_{T \in \mathcal{F}_1(G;S)} w(\overline{T})} \right)$$

where $k(F,*) = \deg^o(F,*)_{cof} - 2$. The expression (I.2) satisfies a monotonicity condition as we vary the vertex set $S \subset V(G)$.

thm:monotonic

Theorem 3 (Monotonicity of normalized principal minors). If $A, B \subset V(G)$ are nonempty subsets with $A \subset B$, then

$$\frac{\det D[A]}{\cot D[A]} \le \frac{\det D[B]}{\cot D[B]}.$$

We remark that the calculation of $\det D[S]$ is related to the following quadratic optimization problem: for all vectors $\mathbf{m} \in \mathbb{R}^S$,

optimize objective function: $\mathbf{m}^{\intercal}D[S]\mathbf{m}$

with constraint: $\mathbf{1}^{\mathsf{T}}\mathbf{m} = 1$.

This result can be shown using Lagrange multipliers; for details, see Section $\frac{\texttt{sec:optimization}}{3}$

Theorem 4 (Bounds on principal minor ratios). Suppose G = (V, E) is a finite, weighted tree with distance matrix D.

(1) If $S \subset V(G)$ is nonempty,

$$0 \le \frac{\det D[S]}{\cot D[S]} \le \frac{1}{2} \sum_{E(G)} \alpha_e.$$

(2) If conv(S,G) denotes the subtree of G consisting of all paths between points of $S \subset V(G)$,

$$\frac{\det D[S]}{\operatorname{cof} D[S]} \le \frac{1}{2} \sum_{E(\operatorname{conv}(S,G))} \alpha_e.$$

(3) If γ is a simple path between vertices $s_0, s_1 \in S$, then

$$\frac{1}{2} \sum_{e \in \gamma} \alpha_e \le \frac{\det D[S]}{\cot D[S]}$$

Corollary 5 (Nonsingular minors). Let G be a finite tree with (weighted) distance matrix D, and let $S \subset V(G)$ be a subset of vertices. If $|S| \ge 2$ then $\det D[S] \ne 0$.

1.3. Previous work. A formula for the inverse matrix D^{-1} was found by Graham and Lovász in 3.

1.4. **Notation.** G a finite graph, loops and parallel edges allowed, possibly disconnected

E(G) edge set of G

V(G) vertex set of G

 $\mathcal{F}_1(G;S)$ the set of S-rooted spanning forests of G

 $\mathcal{F}_2(G;S)$ the set of (S,*)-rooted spanning forests of G

2. Graphs and matrices

For background on enumeration problems for graphs and trees, see Moon $\begin{bmatrix} \frac{poon}{5} \end{bmatrix}$. \diamondsuit choose other reference? \diamondsuit

Given a graph G=(V,E) with edge weights $\{\alpha_e:e\in E\}$, for any edge subset $A\subset E$ we define the weight of A as

$$w(A) = \prod_{e \in A} \alpha_e.$$

We define the co-weight of A as

$$w(\overline{A}) = \prod_{e \notin A} \alpha_e.$$

By abuse of notation, if H is a subgraph of G, we use H to also denote its subset of edges E(H), so e.g. $w(\overline{H}) = w(\overline{E(H)})$.

:graphs-matrices

2.1. **Spanning trees and forests.** A spanning tree of a graph G is a subgraph which is connected, has no cycles, and contains all vertices of G. A spanning forest of a graph G is a subgraph which has no cycles and contains all vertices of G. Let $\kappa(G)$ denote the number of spanning forests of G, and let $\kappa_T(G)$ denote the number of r-component spanning forests.

Given a set of vertices $S = \{v_1, v_2, \dots, v_k\}$, an S-rooted spanning forest of G is a spanning forest which has exactly one vertex v_i in each connected component. An (S, *)-rooted spanning forest of G is a spanning forest which has |S| + 1 components, where |S| components each contain one vertex of S, and the additional component is disjoint from S. We call the component disjoint from S the floating component, following terminology in [?].

Let $\mathcal{F}_1(G; S)$ denote the set of S-rooted spanning forests of G, and let $\mathcal{F}_2(G; S)$ denote the set of (S, *)-rooted spanning forests of G.

Let

$$\kappa_k(v_1|v_2|\cdots|v_k)$$

denote the number of k-component spanning trees which have a vertex v_i in each component. If $S = \{v_1, \ldots, v_k\}$, then $\kappa_k(v_1|\cdots|v_k) = \kappa(G/S)$.

If u, v, w are vertices, then let

$$\kappa_2(uv|w)$$

denote the number of two-forests which have u, v in one component and w in the other component.

Example 6. Suppose G is the tree with unit edge lengths shown below.



Let S be the set of three leaf vertices. Then $\mathcal{F}_1(G;S)$ contains 11 forests, while $\mathcal{F}_2(G;S)$ contains 19 forests. These are shown in Figures 1 and 2, respectively.

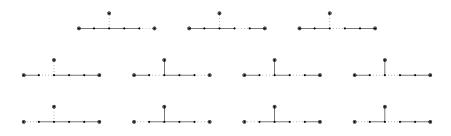


FIGURE 1. Forests in $\mathcal{F}_1(G; S)$.

fig:1-forests

2.2. **Laplacian matrix.** Given a graph G = (V, E), let $L \in \mathbb{R}^{V \times V}$ denote the *Laplacian matrix* of G. If G is a weighted graph with edge weights $\alpha_e \in \mathbb{R}_{>0}$ for $e \in E$, let L denote the weighted Laplacian matrix of G.

Definition 7 (Weighted Laplacian matrix). Given a graph G = (V, E) and edge weights $\{\alpha_e : e \in E\}$, the weighted Laplacian matrix $L_{\alpha} \in \mathbb{R}^{V \times V}$ is defined by

$$(L_{\alpha})_{v,w} = \begin{cases} 0 & \text{if } v \neq w \text{ and } (v,w) \notin E \\ -\alpha_e^{-1} & \text{if } v \neq w \text{ and } (v,w) = e \in E \\ \sum_{e \in N(v)} \alpha_e^{-1} & \text{if } v = w. \end{cases}$$

Given $S \subset V$, let $L[\overline{S}]$ denote the matrix obtained from L by removing the rows and columns indexed by S. For any graph G, let $\kappa(G)$ denote the number of spanning trees of G. The following theorem is due to Kirchhoff.

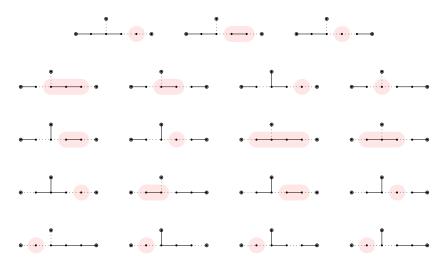


FIGURE 2. Forests in $\mathcal{F}_2(G;S)$.

fig:2-forests

thm:matrix-tree

sec:tree-splits

Theorem 8 (All-minors matrix tree theorem). Let G = (V, E) be a finite graph.

(a) and let L denote the Laplacian matrix of G. Then for any nonempty vertex set $S \subset V$,

$$\det L[\overline{S}] = \kappa(G; S).$$

(b) Let L_{α} denote the weighted Laplacian matrix of G, for edge weights $\{\alpha_e\}$. Then for any nonempty vertex set $S \subset V$,

$$\det L[\overline{S}] = \sum_{T \in \mathcal{F}_1(G;S)} w(T)^{-1} = \sum_{T \in \mathcal{F}_1(G;S)} w(\overline{T}) \prod_{e \in E} \alpha_e^{-1}.$$

Note that $\kappa(G; S)$ is also the number of spanning trees of the quotient graph G/S, which "glues together" all vertices in S as a single vertex.

The following result is due to Bapat–Sivasubramanian. Recall that cof M denotes the *sum of cofactors* of M, i.e.

$$cof M = \sum_{i=1}^{n} \sum_{j=1}^{n} (-1)^{i+j} \det M[\overline{i}, \overline{j}].$$

Theorem 9 (Distance matrix cofactor sums [2]). Given a tree G, let D be the distance matrix of G, and L the Laplacian matrix. Let $S \subset V(G)$ be a nonempty subset of vertices of G. Then

$$\operatorname{cof} D[S] = (-2)^{|S|-1} \det L[\overline{S}].$$

2.3. Tree splits and tree distance. In this section we describe the tree splits associated to a tree, and use their associated indicator functions to give an expression for the tree distance.

Given a tree G=(V,E) and an edge $e\in E$, the edge deletion $G\setminus e$ contains two connected components. Using the implicit orientation on $e=(e^+,e^-)$, we let $(G\setminus e)^+$ denote the component that contains endpoint e^+ , and let $(G\setminus e)^-$ denote the other component. For any $e\in E$ and $v\in V$, we let $(G\setminus e)^v$ denote the component of $G\setminus e$ containing v, respectively $(G\setminus e)^{\overline{v}}$ for the component not containing v.

Tree splits can be used to express the path distance between vertices in a tree. Given an edge $e \in E$ and vertices $v, w \in V$, let

$$\delta(e; v, w) = \begin{cases} 1 & \text{if } e \text{ separates } v \text{ from } w, \\ 0 & \text{otherwise.} \end{cases}$$

In other words, $\delta(e; v, w) = 1$ if the vertices are in different components of the split $G \setminus e$, and $\delta(e; v, w) = 0$ if they are in the same component. Note that $\delta(e; v, v) = 0$ for any e and v.

We have the following perspectives on the function $\delta(e; v, w)$:

- If we fix e and v, then $\delta(e; v, -) : V(G) \to \{0, 1\}$ is the indicator function for the component $(G \setminus e)^{\overline{v}}$ of the tree split $G \setminus e$ not containing v.
- On the other hand if we fix v and w, then $\delta(-; v, w) : E(G) \to \{0, 1\}$ is the indicator function for the unique $v \sim w$ path in G.

Proposition 10 (Weighted tree distance). For a tree G = (V, E) with weights $\{\alpha_e : e \in E\}$, the weighted distance function satisfies

$$d_{\alpha}(v, w) = \sum_{e \in E} \alpha_e \, \delta(e; v, w).$$

For an unweighted tree, we can express the tree distance d(v, w) as the unweighted sum

$$d(v,w) = \sum_{e \in E(G)} \delta(e; v, w).$$

2.4. Outdegree of rooted forest. Given a rooted forest F in $\mathcal{F}_1(G; S)$ and $s \in S$, let F(s) denote the s-component of F. We define the outdegree $\deg^o(F, s)$ as the number of edges which join F(s) to a different component; i.e.

eq:outdeg

(7)
$$\deg^{o}(F,s) = \#\{e = (a,b) \in E : a \in F(s), b \notin F(s)\}.$$

If F is a forest in $\mathcal{F}_2(G; S)$, let $\deg^o(F, *)$ denote the outdegree of the floating component.

lem:outdeg-sum

Lemma 11. Suppose G is a tree and $H \subset G$ is a (nonempty) connected subgraph. Then

$$\sum_{v \in V(H)} (2 - \deg(v)) = 2 - \deg^{o}(H).$$

Proof. This is straightforward to check by induction on |V(H)|, with base case |V(H)| = 1: if $H = \{v\}$ consists of a single vertex, then $\deg^o(H) = \deg(v)$.

2.5. Miscellaneous. We use the fact that the submatrix D[S] has nonzero determinant.

Lemma 12 (Bapat [?, Lemma 8.15]). Suppose D is the (weighted) distance matrix of a tree with n vertices. Then D has one positive eigenvalue and n-1 negative eigenvalues.

Cauchy interlacing, c.f. Horn–Johnson [7, Theorem 4.3.17]

Proposition 13 (Cauchy interlacing). Suppose M is a symmetric real matrix and M[i] is a principal submatrix. Then the eigenvalues of M_i interlace the eigenvalues of M.

$$\lambda_1 \le \mu_1 \le \lambda_2 \le \dots \le \mu_{n-1} \le \lambda_n$$
.

Proposition 14. Suppose D is the distance matrix of a tree G = (V, E) and $S \subset V$ is a subset of size $|S| \geq 2$. Then

- (i) D[S] has one positive and |S|-1 negative eigenvalues;
- (ii) $\det D[S] \neq 0$.

 \square

Remark 15. A key step in the main proof is the use of the map

$$E(G) \times \mathcal{F}_1(G; S) \mapsto S \sqcup \{\text{error}\}$$

defined by

$$(e,T) \mapsto \begin{cases} s & \text{if } e \in T(s), \\ \text{error} & \text{if } e \notin T. \end{cases}$$

rop:distance-sum

Remark 16. For a given spanning forest $F \in \mathcal{F}_2(G; S)$, there are exactly $\deg^o(F, *)$ -many choices of pairs $(T, e) \in \mathcal{F}_1(G; S) \times E(G)$ such that $F = T \setminus e$. Consider the "deletion" map

$$E(G) \times \mathcal{F}_1(G;S) \to \mathcal{F}_2(G;S) \sqcup \mathcal{F}_1(G;S)$$

defined by

$$(e,T) \mapsto \begin{cases} T \setminus e & \text{if } e \in T, \\ T & \text{if } e \notin T. \end{cases}$$

For a forest F in $\mathcal{F}_2(G; S)$, the preimage under this map has $\deg^o(F, *)$ elements. There is an associated "union" map

$$E(G) \times \mathcal{F}_2(G;S) \longrightarrow \mathcal{F}_1(G;S) \sqcup \mathcal{F}_2(G;S)$$

defined by

$$(e,F) \mapsto \begin{cases} F \cup e & \text{if } e \in \partial F, \\ F & \text{if } e \notin \partial F \end{cases}$$

For a forest T in $\mathcal{F}_1(G;S)$, the preimage under this map has |E(T)|-many elements.

3. Optimization: Quadratic programming

Proposition 17. If D[S] is a principal submatrix of a distance matrix indexed by S, then

$$\frac{\det D[S]}{\cot D[S]} = \max\{\mathbf{m}^{\intercal}D[S]\mathbf{m} : \mathbf{m} \in \mathbb{R}^{S}, \ \mathbf{1}^{\intercal}\mathbf{m} = 1\}$$

where $\operatorname{cof} D[S]$ denotes the sum of cofactors of D[S].

Proposition 18. If D[S] is a principal submatrix of a distance matrix indexed by S, then

$$\frac{\det D[S]}{\operatorname{cof} D[S]} = \max\{\mathbf{m}^{\mathsf{T}} D\mathbf{m} : \mathbf{m} \in \mathbb{R}^{V}, \, \mathbf{1}^{\mathsf{T}} \mathbf{m} = 1, \, \mathbf{m}_{v} = 0 \, \text{if } v \notin S\}$$

where cof D[S] denotes the sum of cofactors of D[S].

The gradient of the objective function is $2D[S]\mathbf{m}$, and the gradient of the constraint is 1. By the theory of Lagrange multipliers, the optimal solution \mathbf{m}^* is a vector satisfying

$$D[S]\mathbf{m}^* = \lambda \mathbf{1}$$
 for some $\lambda \in \mathbb{R}$.

The constant λ is in fact the optimal objective value, since

$$(\mathbf{m}^*)^{\mathsf{T}}D[S]\mathbf{m}^* = (D[S]\mathbf{m}^*)^{\mathsf{T}}\mathbf{m}^* = \lambda(\mathbf{1}^{\mathsf{T}}\mathbf{m}^*) = \lambda.$$

The above computation uses the fact that D[S] is a symmetric matrix, and the given constraint $\mathbf{1}^{\mathsf{T}}\mathbf{m} = 1$.

On the other hand, assuming D[S] is invertible we have $\mathbf{m}^* = \lambda(D[S]^{-1}\mathbf{1})$, so that

$$1 = \mathbf{1}^{\mathsf{T}} \mathbf{m}^* = \lambda (\mathbf{1}^{\mathsf{T}} D[S]^{-1} \mathbf{1}) = \lambda \frac{\operatorname{cof} D[S]}{\det D[S]}.$$

Thus the optimal objective value is $\lambda = \frac{\det D[S]}{\cot D[S]}$.

sec:optimization

4. Proofs

In this section we prove our main result, Theorem 2. Theorem I follows by setting all edge weights to one.

Outline of proof: given a subset $S \subset V$ and distance submatrix D[S], we will

- (i) Find vector $\mathbf{m} \in \mathbb{R}^S$ such that $D[S]\mathbf{m} = \lambda \mathbf{1} \in \mathbb{R}^S$.
- (ii) Compute the sum of entries of \mathbf{m} , i.e. $\mathbf{1}^{\mathsf{T}}\mathbf{m}$.
- (iii) Using (i), relate the sum $\mathbf{1}^{\mathsf{T}}\mathbf{m}$ to the sum of entries of the inverse matrix $D[S]^{-1}$:

$$\mathbf{1}^{\mathsf{T}}\mathbf{m} = \lambda(\mathbf{1}^{\mathsf{T}}D[S]^{-1}\mathbf{1}) = \lambda \frac{\operatorname{cof}D[S]}{\det D[S]}.$$

where $\operatorname{cof} D[S]$ is the sum of cofactors of D[S].

(iv) Use known expression for cof D[S] to compute

$$\det D[S] = \lambda(\operatorname{cof} D[S]) \left(\mathbf{1}^{\mathsf{T}} \mathbf{m}\right)^{-1}.$$

The interesting part of this expression will turn out to be in the constant λ .

Example 19. Suppose G is a tree consisting of three paths joined at a central vertex. Let S consist of the central vertex, and the three endpoints of the paths. The corresponding submatrix of the distance matrix is

$$D[S] = \begin{bmatrix} 0 & a & b & c \\ a & 0 & a+b & a+c \\ b & a+b & 0 & b+c \\ c & a+c & b+c & 0 \end{bmatrix}.$$

Following the steps outlined above:

- (i) The vector $\mathbf{m} = \begin{bmatrix} -1\\1\\1\\1 \end{bmatrix}$ satisfies $D[S]\mathbf{m} = (a+b+c)\mathbf{1}$
- (ii) The sum of entries of **m** is $\mathbf{1}^{\mathsf{T}}\mathbf{m} = 2$.
- (iii) We have

$$2 = \mathbf{1}^\intercal \mathbf{m} = \lambda (\mathbf{1}^\intercal D[S]^{-1} \mathbf{1}) = \lambda \frac{\operatorname{cof} D[S]}{\det D[S]}.$$

(iv) The cofactor sum cof D[S] is -8abc, so the determinant is

$$\det D[S] = \lambda \frac{\cot A}{\mathbf{1}^{\mathsf{T}} \mathbf{m}} = (a+b+c)(-8abc)\frac{1}{2} = -4(a+b+c)abc.$$

4.1. Warmup case: S = V.

Proposition 20. Let G = (V, E) a tree, and consider the vector $\mathbf{m} \in \mathbb{R}^V$ defined by

$$\mathbf{m}_v = 2 - \deg v$$
 for each $v \in V$.

Then $\mathbf{1}^{T}\mathbf{m} = \sum_{v \in V} (2 - \deg v) = 2$.

Proof. For any graph, $\sum_{v \in V} \deg v = 2|E|$. Since G is a tree, |E| = |V| - 1.

Proposition 21. Let **m** be the vector defined above, and let D be the distance matrix of G. Then $D\mathbf{m} = \lambda \mathbf{1}$ for some constant λ .

Proof. It suffices to show that for each edge e, with endpoints (e^+, e^-) , we have

$$(D\mathbf{m})_{(e^+)} = (D\mathbf{m})_{(e^-)}.$$

-distance-warmup

We compute

$$(D\mathbf{m})_{(e^{+})} - (D\mathbf{m})_{(e^{-})} = \sum_{v \in V} (d(v, e^{+}) - d(v, e^{-}))(2 - \deg v)$$

$$= \sum_{v \in (G \setminus e)^{-}} \alpha_{e}(2 - \deg v) - \sum_{v \in (G \setminus e)^{+}} \alpha_{e}(2 - \deg v)$$

$$= \alpha_{e} \left(\sum_{v \in (G \setminus e)^{-}} (2 - \deg v) - \sum_{v \in (G \setminus e)^{+}} (2 - \deg v) \right)$$

since

(8)

$$d(v, e^+) - d(v, e^-) = \begin{cases} \alpha_e & \text{if } v \text{ is closer to } e^- \text{ than } e^+, \\ -\alpha_e & \text{if } v \text{ is closer to } e^+ \text{ than } e^-. \end{cases}$$
 For each sum in (7), we apply Lemma lem: obtain

$$\sum_{v \in (G \setminus e)^-} (2 - \deg v) = (2 - \deg^o((G \setminus e)^-) = 1,$$

since each component of $(G \setminus e)$ has outdegree one. The same identity applies to the sum over $(G \setminus e)^+$, so $(D\mathbf{m})_{(e^+)} = (D\mathbf{m})_{(e^-)}$ as desired.

4.2. General case: $S \subset V$. Fix a tree G = (V, E) and a nonempty subset $S \subset V$.

dfn:m-vector

eq:12-1

Definition 22. Let $\mathbf{m} = \mathbf{m}(G; S) \in \mathbb{R}^S$ be defined by

eq:m-vector

(9)
$$\mathbf{m}_{v} = \sum_{T \in \mathcal{F}_{1}(G;S)} (2 - \deg^{o}(T,v)) w(\overline{T}) \quad \text{for each } v \in S.$$

where $\deg^o(T,v)$ is the outdegree of the v-component of T, (6).

Let 1 denote the all-ones vector.

Proposition 23. For **m** defined above, $\mathbf{1}^{\mathsf{T}}\mathbf{m} = 2\sum_{T \in \mathcal{F}_1(G;S)} w(\overline{T})$.

Proof. We have

$$\mathbf{1}^{\mathsf{T}}\mathbf{m} = \sum_{s \in S} \mathbf{m}_s = \sum_{s \in S} \left(\sum_{T \in \mathcal{F}_1(G;S)} (2 - \deg^o(T,s)) w(\overline{T}) \right)$$

$$= \sum_{T \in \mathcal{F}_1(G;S)} w(\overline{T}) \left(\sum_{s \in S} \sum_{v \in T(s)} 2 - \deg(v) \right)$$

$$= \sum_{T \in \mathcal{F}_1} w(\overline{T}) \left(\sum_{v \in V} 2 - \deg(v) \right) = \sum_{T \in \mathcal{F}_1} w(\overline{T}) \cdot 2.$$

In the second line we apply Lemma II and exchange the outer summations. To obtain the third line, we observe that the vertex sets of T(s) for $s \in S$ form a partition of V, since T is an S-rooted spanning forest. Finally we again apply Lemma II for the last equality, as $\deg^o(G) = 0$.

Corollary 24. If G is a graph with unit edge weights $\alpha_e = 1$, then the vector \mathbf{m} defined in (8)satisfies $\mathbf{1}^{\mathsf{T}}\mathbf{m} = 2 \kappa(G; S)$.

Theorem 25. With $\mathbf{m} = \mathbf{m}(G; S)$ defined as in (8), $D[S]\mathbf{m} = \lambda \mathbf{1}$ for the constant

$$\lambda = \sum_{F_2(G:S)} \alpha_e \sum_{F_2(G:S)} w(\overline{T}) - \sum_{F_2(G:S)} (2 - \deg^o(F, *))^2 w(F).$$

Proof. For $e \in E$ and $v, w \in V$, let $\delta(e; v, w)$ denote the function defined in Section Section Section Section For any $v \in S$, we have

$$(D[S]\mathbf{m})_{v} = \sum_{s \in S} d(v, s)\mathbf{m}_{s}$$

$$= \sum_{s \in S} \left(\sum_{e \in E(G)} \alpha_{e} \,\delta(e; v, s)\right) \left(\sum_{T \in \mathcal{F}_{1}(G; S)} (2 - \deg^{o}(T, s))w(\overline{T})\right)$$

$$= \sum_{T \in \mathcal{F}_{1}} w(\overline{T}) \sum_{e \in E} \alpha_{e} \left(\sum_{s \in S} \delta(e; v, s)(2 - \deg^{o}(T, s))\right)$$

$$= \sum_{T \in \mathcal{F}_{1}} w(\overline{T}) \sum_{e \in E} \alpha_{e} \left(\sum_{s \in S} \delta(e; v, s) \sum_{u \in T(s)} (2 - \deg(u))\right).$$

$$= \sum_{T \in \mathcal{F}_{1}} w(\overline{T}) \sum_{e \in E} \alpha_{e} \left(\sum_{s \in S} \delta(e; v, s) \sum_{u \in T(s)} (2 - \deg(u))\right).$$

where in the last equality, we apply Lemma II to the subgraph H = T(s).

We introduce additional notation to handle the double sum in parentheses in (9). Each S-rooted spanning tree T naturally induces a surjection $\pi_T: V \to S$, defined by

$$\pi_T(u) = s$$
 if and only if $u \in T(s)$.

Using this notation,

$$(D[S]\mathbf{m})_v = \sum_{T \in \mathcal{F}_1} w(\overline{T}) \sum_{e \in E} \alpha_e \left(\sum_{u \in V} (2 - \deg(u)) \delta(e; v, \pi_T(u)) \right)$$

We will compare the above expression with the one obtained after replacing $\delta(e; v, \pi_T(u))$ with $\delta(e; v, u)$. From Lemma II again, for any $v \in V$ and $e \in E$ we have

$$\sum_{v \in V} (2 - \deg(u))\delta(e; v, u) = 2 - \deg^{o}((G \setminus e)^{\overline{v}}) = 1.$$

Thus

$$\underbrace{ \begin{array}{c} \mathbf{eq:10} \end{array}} \quad (12) \qquad \qquad \underbrace{\sum_{T \in \mathcal{F}_1} w(\overline{T}) \sum_{e \in E} \alpha_e} = \sum_{T \in \mathcal{F}_1} w(\overline{T}) \sum_{e \in E} \alpha_e \left(\sum_{u \in V} (2 - \deg(u)) \delta(e; v, u) \right)$$

After subtracting equation ($\stackrel{|eq:10}{(4.2)}$ from ($\stackrel{|eq:9}{(4.2)}$,

$$(D[S]\mathbf{m})_v - \sum_{T \in \mathcal{F}_1} w(\overline{T}) \sum_{e \in E} \alpha_e = \sum_{T \in \mathcal{F}_1} w(\overline{T}) \sum_{e \in E} \alpha_e \sum_{u \in V} (2 - \deg(u)) \left(\delta(e; v, \pi_T(u)) - \delta(e; v, u)\right)$$

When $e \in E$ and $v \in V$ are fixed, $u \mapsto \delta(e; v, u)$ is the indicator function of one component of the principal cut $G \setminus e$. We have

$$\underline{\text{eq:delta-diff}} \quad (13) \qquad \delta(e; v, \pi_T(u)) - \delta(e; v, u) = \begin{cases} 0 & \text{if } \delta(e; \pi_T(u), u) = 0 \\ 1 & \text{if } \delta(e; \pi_T(u), u) = 1 \text{ and } \delta(e; v, \pi_T(u)) = 1 \\ -1 & \text{if } \delta(e; \pi_T(u), u) = 1 \text{ and } \delta(e; v, \pi_T(u)) = 0 \end{cases}$$

Now consider varying u over all vertices, when e, T, and v are fixed. We have the following three cases:

Case 1: if $e \notin T$, then u and $\pi_T(u)$ are on the same side of the principal cut $G \setminus e$, for every vertex u. In this case $\delta(e; v, \pi_T(\cdot)) - \delta(e; v, \cdot) = 0$.

Case 2: if $e \in T(s_0)$ and s_0 is separated from v by e, then $\delta(e; v, \pi_T(\cdot)) - \delta(e; v, \cdot)$ is the indicator function for the floating component of $T \setminus e$. See Figure 3, left.

Case 3: if $e \in T(s_0)$ and s_0 is on the same component as v from e, then $\delta(e; v, \pi_T(\cdot)) = \delta(e; v, \cdot) = \delta(e; v,$

FIGURE 3. Edge $e \in T(s_0)$ with $\delta(e; v, s_0) = 1$ (left) and $\delta(e; v, s_0) = 0$ (right). The floating component of $T \setminus e$ is highlighted.

fig:e-delete-fro

Thus when multiplying the term (4.2) by $(2-\deg(u))$ and summing over all vertices u, we obtain

$$\sum_{u \in V} (2 - \deg(u)) (\delta(e; v, \pi_T(u)) - \delta(e; v, u)) = \begin{cases} 0 & \text{if } e \not\in T, \\ 2 - \deg^o(T \setminus e, *) & \text{if } e \in T(s_0) \text{ and } \delta(e; v, s_0) = 1, \\ -(2 - \deg^o(T \setminus e, *)) & \text{if } e \in T(s_0) \text{ and } \delta(e; v, s_0) = 0. \end{cases}$$

Thus

$$\begin{array}{ll} \boxed{ \begin{array}{ll} \mathbf{eq:1} \end{array} } & (14) & (D[S]\mathbf{m})(v) - \sum_{T \in \mathcal{F}_1} w(\overline{T}) \sum_{e \in E} \alpha_e \\ \\ & = \sum_{T \in \mathcal{F}_1} w(\overline{T}) \sum_{s_0 \in S} \left(\sum_{\substack{e \in T(s_0) \\ \delta(e;v,s_0) = 1}} \alpha_e(2 - \deg^o(T \setminus e, *)) - \sum_{\substack{e \in T(s_0) \\ \delta(e;v,s_0) = 0}} \alpha_e(2 - \deg^o(T \setminus e, *)) \right) \\ \\ & = \sum_{T \in \mathcal{F}_1} w(\overline{T}) \sum_{e \in T} \alpha_e(2 - \deg^o(T \setminus e, *)) \left(\sum_{\substack{e \in T(s_0) \\ e \in T(s_0)}} 1 - \sum_{\substack{e \in T(s_0) \\ e \in T(s_0)}} 1 \right). \end{aligned}$$

We now rewrite the above expression in terms of $\mathcal{F}_2(G; S)$, observing that the deletion $T \setminus e$ is an (S, *)-rooted spanning forest of G, if $e \in T$, and that the corresponding weights satisfy

$$w(F) = \alpha_e \cdot w(\overline{T})$$
 if $F = T \setminus e$.

Thus

$$\begin{pmatrix}
|eq:1\\ |T|| = \sum_{F \in \mathcal{F}_2} w(F)(2 - \deg^o(F, *)) \sum_{T \in \mathcal{F}_1} \sum_{s_0 \in S} \left(\sum_{\substack{e \in T(s_0)\\ \delta(e; v, s_0) = 1}} \mathbb{1}(F = T \setminus e) - \sum_{\substack{e \in T(s_0)\\ \delta(e; v, s_0) = 0}} \mathbb{1}(F = T \setminus e) \right)$$

$$= \sum_{F \in \mathcal{F}_2} w(F)(2 - \deg^o(F, *)) \left(\#\{T \in \mathcal{F}_1 : F = T \setminus e \text{ for some } e \in T(s_0), \ \delta(e; v, s_0) = 1\} \right)$$

$$- \#\{T \in \mathcal{F}_1 : F = T \setminus e \text{ for some } e \in T(s_0), \ \delta(e; v, s_0) = 0\} \right)$$

Next, we note that $F = T \setminus e$ is equivalent to $T = F \cup e$, and in particular this only occurs when we choose the edge e to be in the floating boundary $\partial F(*)$:

FIGURE 4. Edge $e \in \partial F(*)$ with $\delta(e; v, F(*)) = 0$ (left) and $\delta(e; v, F(*)) = 1$ (right). The floating component F(*) is highlighted.

Now for any $e \notin F$, let $\delta(e; v, F(*)) = \delta(e; v, x)$ for any $x \in F(*)$, i.e.

$$\delta(e; v, F(*)) = \begin{cases} 1 & \text{if } e \text{ lies on path from } v \text{ to } F(*), \\ 0 & \text{otherwise.} \end{cases}$$

The condition $F = T \setminus e$ for some $e \in T(s_0)$ with $\delta(e; v, s_0) = 1$ (resp. $\delta(e; v, s_0) = 0$) is equivalent to $T = F \cup e$ for some $e \in \partial F(*)$ with $\delta(e; v, F(*)) = 0$ (resp. $\delta(e; v, F(*)) = 1$). Thus

$$(\stackrel{\text{leq}:1}{\text{IO}}) = \sum_{F \in \mathcal{F}_2} w(F)(2 - \deg^o(F, *)) \Bigg(\#\{e \in \partial F(*) : \delta(e; v, F(*)) = 0\}$$

$$- \#\{e \in \partial F(*) : \delta(e; v, F(*)) = 1\} \Bigg).$$

Finally, we observe that for any forest F in $\mathcal{F}_2(G;S)$, there is exactly one edge e in the boundary $\partial F(*)$ of the floating component which satisfies $\delta(e;v,F(*))=1$, namely the unique boundary edge on the path from the floating component F(*) to v. The previous expression ($\overline{10}$) simplifies as

$$\#\{e \in \partial F(*) : \delta(e; v, F(*)) = 1\} = 1$$
 and $\#\{e \in \partial(F, *) : \delta(e; v, F(*)) = 0\} = \deg^{o}(F, *) - 1.$

Thus

$$(\stackrel{\text{leq}: 1}{\text{IO}}) = \sum_{F \in \mathcal{F}_2} w(F)(2 - \deg^o(F, *)) \Big((\deg^o(F, *) - 1) - (1) \Big)$$

$$= -\sum_{F \in \mathcal{F}_2} w(F)(2 - \deg^o(F, *))^2.$$

as desired.



Figure 5. Components rooted in $S(G \setminus e)^{\overline{v}}$.

Finally we can prove our main theorem.

Proof of Theorem
$$2. \overline{\text{Since}}$$
.

Remark 26. It is worth observing that depending on the chosen subset $S \subset V$, the distances appearing in the submatrix D[S] may ignore a large part of the ambient tree G. We could instead replace G by the subtree $\operatorname{conv}(S,G)$ consisting of the union of all paths between vertices in S, which we call the *convex hull* of $S \subset G$. To apply formula (2) or (4) "efficiently," we should replace G on the right-hand side with the subtree $\operatorname{conv}(S,G)$. However, the formulas as stated are true even without this replacement due to cancellation of terms.

5. Physical interpretation

If we consider G as a network of wires with each edge containing a unit resistor, which is grounded at all nodes in S, then \mathbf{m}_S has an interpretation as current flow: it records the currents flowing to S when current is added on $V \setminus S$ in the amount $2 - \deg v$ for each $v \notin S$.

5.1. Alternate proof. Let 1 denote the all-ones vector. When we choose a subset $S \subset V(G)$, we no longer have a single "obvious" replacement for \mathbf{m} inside \mathbb{R}^S . Instead, we can take an average over S-rooted spanning forests.

In the outline above, our first goal is to find a "special" vector $\mathbf{m} \in \mathbb{R}^S$ satisfying $D[S]\mathbf{m} = \lambda \mathbf{1}$. We can approach this first goal as follows: consider \mathbb{R}^S inside the larger vector space $\mathbb{R}^V = \mathbb{R}^S \oplus \mathbb{R}^{V \setminus S}$, and let π_S denote the projection from \mathbb{R}^V to \mathbb{R}^S . We wish to find vectors $\mathbf{n}_i \in \mathbb{R}^V$ satisfying $\pi_S(D\mathbf{n}_i) = \lambda_i \mathbf{1}$. OR, $D\mathbf{n}_i = \lambda_i \mathbf{1}^S \oplus (-)$. By finding sufficiently many such vectors \mathbf{n}_i , we can hope to find a linear combination that lies inside $\mathbb{R}^S \oplus \{0\}$.

prop:n-vector

Proposition 27. Suppose $v \in V \setminus S$. For each $s_j \in S$, let $\mu(v, s_j) = current$ flowing to s_j when G is grounded at S and one unit of current enters G at v. Explicitly,

$$\begin{split} \mu(v,s) &= \frac{\# \ of \ S\text{-}rooted \ spanning forests of } G \ whose \ s_{j}\text{-}component \ contains } v}{\# \ of \ S\text{-}rooted \ spanning forests of } G \\ &= \frac{\sum_{\mathcal{F}_{1}(G/S)} \mathbb{1}(v \in T(s))}{\kappa(G/S)} \\ &= \frac{\kappa_{r}(s_{1}|\cdots|s_{j}v|\cdots|s_{r})}{\kappa_{r}(s_{1}|\cdots|s_{r})} \end{split}$$

Consider the vector $\mathbf{n} = \mathbf{n}(G; S, v) \in \mathbb{R}^V$ defined by

$$\mathbf{n}_v = 1, \quad \mathbf{n}_s = -\mu(v, s) \text{ if } s \in S, \quad \mathbf{n}_w = 0 \text{ if } w \notin S \cup v$$

Then $D\mathbf{n}$ is constant on S, i.e. $\pi_S(D\mathbf{n}) = \lambda \mathbf{1}$ for some λ .

Proof sketch. For any $s, s' \in S$, consider tracking the value of $D\mathbf{n}$ along path from s to s'. The value of $D\mathbf{n}$ changes according to current flow in the corresponding network, i.e. $D\mathbf{n}$ records electrical potential. By assumption S is grounded, so $D\mathbf{n}$ takes the same value at s and s'.

Theorem 28. Let G be a tree, S a nonempty subset of vertices, and D[S] the corresponding submatrix of the distance matrix. Suppose $\mathbf{m} = \mathbf{m}(G; S) \in \mathbb{R}^S$ is defined by (8);

$$\mathbf{m}(G;S)_v = \sum_{T \in \mathcal{F}_1(G;S)} w(\overline{T}) \sum_{u \in T(v)} (2 - \deg u) = \sum_{T \in \mathcal{F}_1(G;S)} w(\overline{T}) (2 - \deg^o(T,v)).$$

Then $D[S]\mathbf{m} = \lambda \mathbf{1}$ for some constant λ .

Proof. The vector $\mathbf{m} = \mathbf{m}(G; S)$ can be expressed as a linear combination

$$\mathbf{m}(G; S) = \kappa(G; S) \left(\sum_{v \in V} (2 - \deg v) \delta(v) - \sum_{v \in V \setminus S} (2 - \deg v) \mathbf{n}(G; S, v) \right)$$
$$= \kappa(G; S) \left(\mathbf{m}(G; V) - \sum_{v \in V \setminus S} (2 - \deg v) \mathbf{n}(G; S, v) \right)$$

♦ TODO: elaborate on this equation of Proposition Proposition 21 we know that $D\mathbf{m}(G; V)$ is constant on V, and from Proposition 27 we know that $D\mathbf{n}(G; S, v)$ is constant on S. Hence by linearity, $D\mathbf{m}(G; S)$ is constant on S.

Proposition 29. Let G = (V, E) be a tree, and $S \subset V$. Suppose we label $S = \{s_1, \ldots, s_r\}$ and $V \setminus S = \{t_1, \ldots, t_{n-r}\}$. For each $t_i \in V \setminus S$, consider $\mathbf{f}_i \in \mathbb{R}^V$ defined by

Example 30. If

$$D[S \cup t] = \begin{bmatrix} 0 & a & b & c \\ a & 0 & a+b & a+c \\ b & a+b & 0 & b+c \\ c & a+c & b+c & 0 \end{bmatrix}$$

then

$$\begin{bmatrix} 0 & a & b & c \\ a & 0 & a+b & a+c \\ b & a+b & 0 & b+c \\ c & a+c & b+c & 0 \end{bmatrix} \begin{bmatrix} ab+ac+bc \\ -bc \\ -ac \\ -ab \end{bmatrix} = \begin{bmatrix} -3abc \\ -abc \\ -abc \\ -abc \end{bmatrix}$$

6. Examples

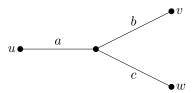
Example 31. Suppose G is a tree consisting of three paths joined at a central vertex. Let S consist of the central vertex, and the three endpoints of the paths. The corresponding minor of the distance matrix is

$$D[S] = \begin{bmatrix} 0 & a & b & c \\ a & 0 & a+b & a+c \\ b & a+b & 0 & b+c \\ c & a+c & b+c & 0 \end{bmatrix} \sim \begin{bmatrix} 0 & a & b & c \\ a & -a & a & a \\ b & b & -b & b \\ c & c & c & c & -c \end{bmatrix} \sim \begin{bmatrix} 0 & a & b & c \\ a & -2a & 0 & 0 \\ b & 0 & -2b & 0 \\ c & 0 & 0 & -2c \end{bmatrix}.$$

The determinant is

$$\det D[S] = -4(a+b+c)abc.$$

Example 32. Suppose Γ is a tripod with lengths a, b, c and corresponding leaf vertices u, v, w.



Let $S = \{u, v, w\}$. Then

$$D[S] = \begin{bmatrix} 0 & a+b & a+c \\ a+b & 0 & b+c \\ a+c & b+c & 0 \end{bmatrix}.$$

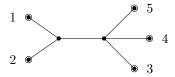
and

$$\det D[S] = 2(a+b)(a+c)(b+c) = 2((a+b+c)(ab+ac+bc) - abc).$$

The "special vector" that satisfies $D[S]\mathbf{m} = \lambda \mathbf{1}$ in this example is

$$\mathbf{m} = \begin{bmatrix} a(b+c) & b(a+c) & c(a+b) \end{bmatrix}^{\mathsf{T}}.$$

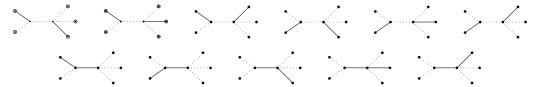
Example 33. Suppose G is the tree with unit edge lengths shown below, with five leaf vertices.



Let S denote the set of five leaf vertices. Then

$$D[S] = \begin{bmatrix} 0 & 2 & 3 & 3 & 3 \\ 2 & 0 & 3 & 3 & 3 \\ 3 & 3 & 0 & 2 & 2 \\ 3 & 3 & 2 & 0 & 2 \\ 3 & 3 & 2 & 2 & 0 \end{bmatrix}.$$

There are 11 forests in $\mathcal{F}_1(G;S)$:

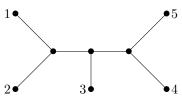


There are 6 forests in $\mathcal{F}_2(G;S)$:

and

$$\det D[S] = 368 = (-1)^4 2^3 \left(6 \cdot 11 - \left(3 \cdot 1^2 + 2 \cdot 2^2 + 1 \cdot 3^2 \right) \right)$$

Example 34. Suppose G is the tree with unit edge lengths shown below, with five leaf vertices and three internal vertices.



Let S denote the set of five leaf vertices. Then

$$D[S] = \begin{bmatrix} 0 & 2 & 3 & 4 & 4 \\ 2 & 0 & 3 & 4 & 4 \\ 3 & 3 & 0 & 3 & 3 \\ 4 & 4 & 3 & 0 & 2 \\ 4 & 4 & 3 & 2 & 0 \end{bmatrix}$$

and

$$\det D[S] = 864 = (-1)^4 2^3 \left(7 \cdot 21 - (14 \cdot 1^2 + 4 \cdot 2^2 + 1 \cdot 3^2)\right)$$

7. Further work

It is natural to ask whether these results for trees may be generalized to arbitrary finite graphs. See [6], [?].

7.1. **Symanzik polynomials.** We note that the expression in the main theorem \Diamond cite \Diamond is related to Symanzik polynomials, which we recall here.

Given a graph G=(V,E), the first Symanzik polynomial is the homogeneous polynomial in $\underline{x}=\{x_e:e\in E\}$

$$\psi_G(\underline{x}) = \sum_{T \in \mathcal{F}_1(G)} \prod_{e \notin T} x_e.$$

while the second Symanzik polynomial is

$$\varphi_G(p; \underline{x}) = \sum_{F \in \mathcal{F}_2(G)} \left(\sum_{v \in F_1} p(v) \right)^2 \prod_{e \notin F} x_e.$$

Theorem statement:

$$\det D[S] = (-1)^{|S|-1} 2^{|S|-2} \left(\sum_{E(G)} \alpha_e \sum_{\mathcal{F}_1(G;S)} w(\overline{T}) - \sum_{\mathcal{F}_2(G;S)} (\deg^o(F,*) - 2)^2 w(F) \right).$$

$$\frac{\det D[S]}{\cot D[S]} = \frac{1}{2} \left(\sum_{e \in E} \alpha_e - \frac{\sum_{F \in \mathcal{F}_2(G;S)} w(F) (\deg^o(F,*) - 2)^2}{\sum_{T \in \mathcal{F}_1(G;S)} w(\overline{T})} \right)$$

In terms of Symanzik polynomials, let ψ and φ denote the first and second Symanzik polynomials of the quotient graph G/S.

$$\frac{\det D[S]}{\cot D[S]} = \frac{1}{2} \left(\sum_{e \in E} \alpha_e - \frac{\varphi_{G/S}(p; \underline{\alpha})}{\psi_{G/S}(\underline{\alpha})} \right)$$

with momentum function $p(v) = \deg^{o}(v) - 2$ for $v \notin S$.

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