

# Carry Forward & Subarrays

⇒ Intro about Contest

$$\Rightarrow \begin{array}{c} 9:00 \longrightarrow 10:30 \\ \hline 1.5 \text{ hour} \end{array} \qquad \begin{array}{c} 10:30 \longrightarrow 11:30 \\ \hline \text{Doubt Session} \end{array}$$

⇒ 3 questions [ 2 questions to pass ]

⇒ contest will be Monday [ re-attempt ]

Question 1 :- Given a string 's' of lowercase characters, return the count of pairs  $(i, j)$  such that  $i < j$  &  $s[i]$  is 'a' &  $s[j]$  is 'g'.

Ex:-  $s = "a^1 b^2 e^3 g^4 a^5 g^6"$

Ans = 3

Quiz 1

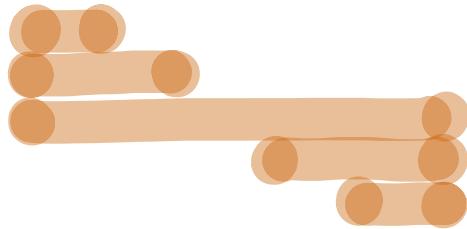
$s = "a^1 c^2 g^3 d^4 g^5 a^6 g^7"$



Ans = 4

Quiz 2

$s = " b \ c \ a \ g \ g \ a \ a \ g "$



Ans = 5

### Brute force

For each 'a' we can check how many 'g' are present on right. And for this we can use two nested loops.

### PSEUDO CODE

```

int ans = 0;
for (i=0; i < n; i++) {
    if (s[i] == 'a') {
        for (j=i+1; j < n; j++) {
            if (s[j] == 'g')
                ans++;
    }
}
return ans;
    
```

$$TC \rightarrow O(N * N)$$

$$SC \rightarrow O(1)$$

## ↳ OPTIMIZATION

### Observation

- ① For each 'g'. Count the nos. of 'a' on the left.
- ② We can maintain a counter variable for storing the no. of 'a' on the left while iterating from left to right.

### DRY RUN

	a	c	b	a	g	K	a	g	g		
counta =	1	1	1	2	2	2	3	3	3		
ans =	0	0	0	0	2	2	2	5	8		

### PSEUDO CODE

```
int ans=0;
int counta=0;

for (i=0; i<n; i++){
    if (s[i] == 'a'){
        counta++;
    }
    if (s[i] == 'g'){
        ans += counta;
    }
}
```

return ans;

TC  $\rightarrow$  O(N)

SC  $\rightarrow$  O(1)

\* This technique is known as carry forward

## # SUBARRAYS

A Subarray is a contiguous part of an array. It is formed by deleting a range of elements from the array. A Subarray can have one or more elements & must be a contiguous part of the original array.

Ex:- arr[] = [4, 1, 2, 3, -1, 6, 9, 8, 12]

[2 3 -1 6] ✓

[9] ✓

[4, 1, 2, 3, -1, 6, 9, 8, 12] ✓

[4, 12] X

[1, 2, 6] X

[3 2 14] X

Qniz 3 :-

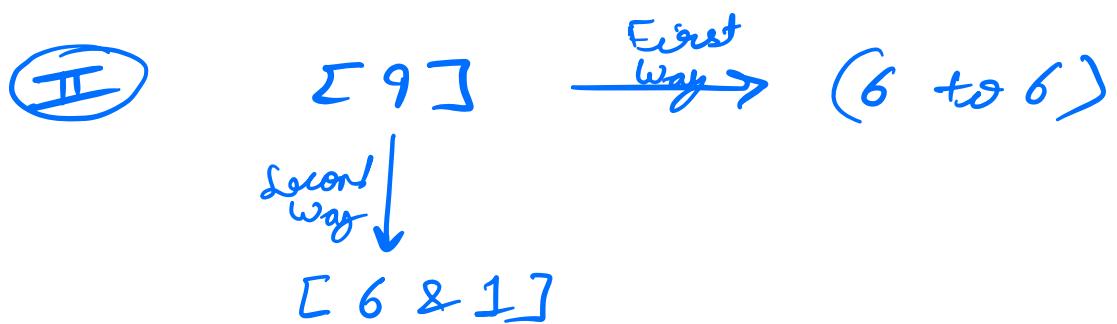
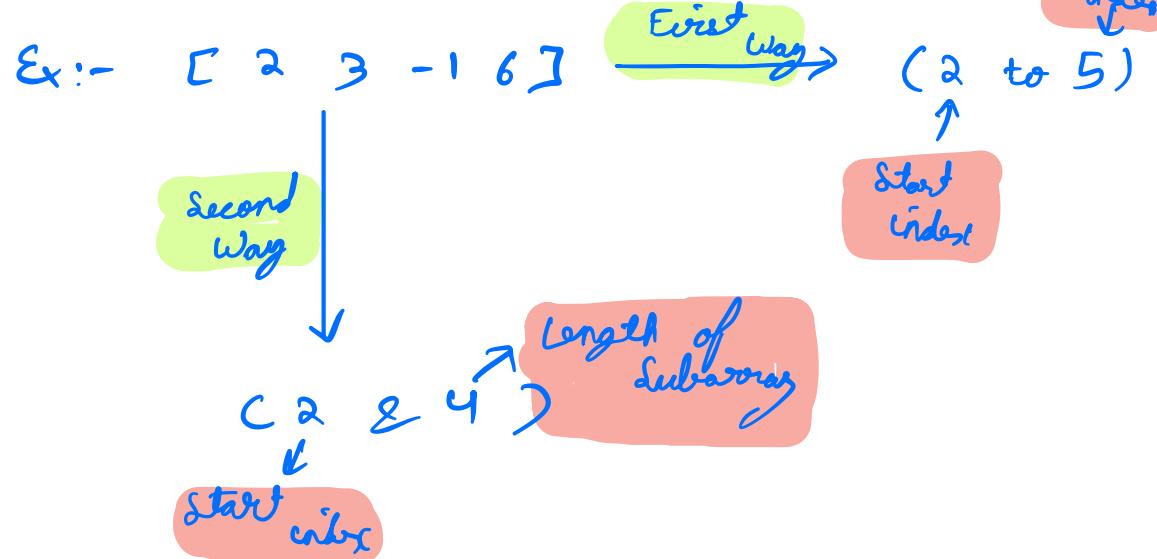
A[] = {2, 4, 1, 6, -3, 7, 8, 4}

Ans  $\Rightarrow$  {7 8 4}

# Representation of an Subarray.

↳ Two Way

$$\text{arr}[] = [4^0, 1^1, 2^2, 3^3, -1^4, 6^5, 9^6, 8^7, 12^8]$$



Ques 4 How many subarrays of the following array starts from index 0.

$$[4^0, 2^1, 10^2, 3^3, 12^4, -2^5, 15^6]$$

Ans

		index
$[4]$		$(0, 0)$
$[4 \ 2]$		$(0, 1)$
$[4 \ 2 \ 10]$		$(0, 2)$
$[4 \ 2 \ 10 \ 3]$		$(0, 3)$
$[4 \ 2 \ 10 \ 3 \ 12]$		$(0, 4)$
$[4 \ 2 \ 10 \ 3 \ 12 \ -2]$		$(0, 5)$
$[4 \ 2 \ 10 \ 3 \ 12 \ -2 \ 15]$		$(0, 6)$

Quiz 5 :- How many subarrays of the following array starts from index 1.

$[4^0, 2, 10, 3, 12, -2, 15]$

Ans

	index
$[2]$	$(1, 1)$
$[2, 10]$	$(1, 2)$
$[2, 10, 3]$	$(1, 3)$
$[2, 10, 3, 12]$	$(1, 4)$
$[2, 10, 3, 12, -2]$	$(1, 5)$
$[2, 10, 3, 12, -2, 15]$	$(1, 6)$

## # Formula to count total subarrays

arr []

index  $\geq 0, 1, 2, 3, \dots, n-2, n-1$

Subarrays starting with 0

$(0, 0)$
$(0, 1)$
$(0, 2)$
$\vdots$
$(0, n-1)$

Total  $n$

Subarray starting with 1

$(1, 1)$
$(1, 2)$
$(1, 3)$
$\vdots$
$(1, n-1)$

$n-1$

Subarray starting with 2

$(2, 2)$
$(2, 3)$
$(2, 4)$
$\vdots$
$(2, n-1)$

$n-2$

Subarray starting with  $(n-1)$

$(n-1, n-1)$
--------------

1

## Total Subarrays

$$\Rightarrow N + (N-1) + (N-2) + \dots + 1$$

$$\Rightarrow \frac{N * (N+1)}{2}$$

Ex:-  $[1 \ 2 \ 3]$

Total Subs :-

$\begin{matrix} [1] \\ [1, 2] \\ [1, 2, 3] \\ [2] \\ [2, 3] \\ [3] \end{matrix}$

(6)

By formula

$$\frac{N * (N+1)}{2}$$

$$\frac{3 * 4}{2} = 6$$

Question 2:- Given an array of integers & 2 indices, a start & end index. we need to print the subarray of the array that starts from start index & ends at the end index (both inclusive)

Ex:-  $[5^{\circ} \ 4 \ 3 \ 1 \ 2]$

$$\begin{matrix} S = 1 \\ E = 3 \end{matrix}$$

Ans:  $[4 \ 3 \ 1]$

## Brute Force

→ We can print from  $S$  to  $E$

## PSEUDO CODE

$T.C \rightarrow O(N)$   
 $S.C \rightarrow O(1)$

```
void printSubarray ( A[], S, E ) {
    for ( i = S; i <= E; i++ ) {
        Print ( A[i] );
    }
}
```

Question 3 :- Print all the subarrays of the given array.

Ex:-  $[1^0, 2^0, 3^2]$

Subarrays :-

$[10]$   
 $[10 20]$   
 $[10 20 30]$   
 $[20]$   
 $[20 30]$   
 $[30]$

### OBSERVATION

- ① Exist of all we need find all possible S & E. For getting all the possible S & E, we can use nested loops.

### PSEUDO CODE

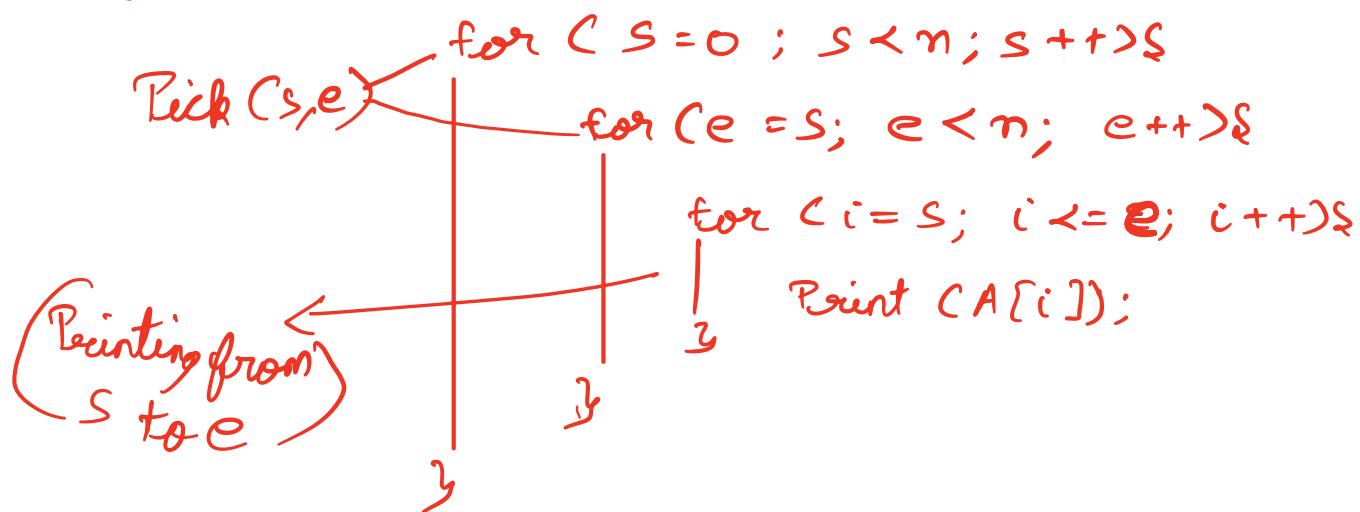
```
for (s=0; s<n; s++) {  
    for (e=s; e<n; e++) {  
        print (s, e);  
    }  
}
```

### DRY RUN

<u>s</u>	<u>e</u>
0	0
0	1
1	1
2	2

② And we just print it. Need to iterate from  $s$  to  $e$ .

### PSEUDO CODE



$$TC \rightarrow O(N^3)$$
$$SC \rightarrow O(1)$$

Question 4 Given an array of  $N$  integers, return the length of the smallest subarray which contains both maximum & minimum elements of the array.

Ques 6 which is the length of the smallest subarray which has both, the Max & Min of the array?

Ex:- 2, 2, 6, 4, 5, 1, 5, 2, 6, 4, 1

$$\text{Max} = 6$$

$$\text{Min} = 1$$

[ 6 4 1 ]

(3)

AB

## Brute force

→ check all the possible Subarray & the one with minimum length containing Max & Min will be the Ans.

\* C HW)

$$TC \rightarrow O(N^3)$$

$$SC \rightarrow O(1)$$

## ↳ OPTIMIZATION

### Observation

① we need only 1 Max & 1 Min in Ans.

② We will be having only Max & Min on corners

→ Max on Left & Min on Right

→ Min on Left & Max on Right

③ Basically we are looking for Subarray which starts with max & end on closest possible min or vice-versa

④ We can keep a track of last closest Max & Min to check smallest possible Sub.

## DRY RUN

$A[] = \{ 2, 2, 6, 4, 5, 1, 5, 2, 6, 4, 1 \}$

Let,  
 $\text{last\_Min\_idx} = -1$   
 $\text{last\_Max\_idx} = -1$   
 $\text{ans} = \text{INT\_MAX}$   
 $\text{min Value} = 1$   
 $\text{max Value} = 6$

$i$	$A[i]$	$\text{last\_min\_idx}$	$\text{last\_max\_idx}$	$\text{ans}$
0	2	-1	-1	8
1	2	-1	-1	8
2	6	-1	-1	8
3	4	1	1	8
4	5	1	1	8
5	1	5	5	$5 - 2 + 1 = 4$
6	5	5	5	4
7	2	5	5	4
8	6	5	8	$8 - 5 + 1 = 4$
9	4	5	8	4
10	1	10	8	$10 - 8 + 1 = 3$

## PSEUDO CODE

```
int Max // TODO  
int Min // TODO  
last_Min-idx = -1, last_Max-idx = -1  
ans = ∞  
  
for (i = 0; i < n; i++) {  
    if (arr[i] == max) {  
        if (last_min-idx != -1) {  
            ans = min(ans, i - last_min-idx + 1);  
        }  
        last_Max-idx = i;  
    }  
    else if (arr[i] == min) {  
        if (last_max-idx != -1) {  
            ans = min(ans, i - last_max-idx + 1);  
        }  
        last_Min-idx = i;  
    }  
  
return ans;
```

↳ EDGE CASE :- 888 [ all elements same ]

if (max == min) return 1;

TC  $\rightarrow O(N)$

SC  $\rightarrow O(1)$

## # Next Class Content

- SLIDING WINDOW
- CONTRIBUTION TECHNIQUE
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