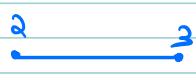
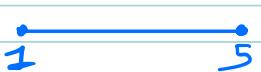


Arrays - Interview Problems

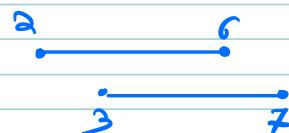
MERGE INTERVALS

↳ An Interval is defined by Start & End time
where start \leq end.

Ex:- $[s \ s'] \rightarrow [2 \ 3]$ 

$\hookrightarrow [1 \ 5] \rightarrow [1 \ 5]$ 

↳ Let's say we have been given a list of intervals. We have to merge if they overlap.



$\hookrightarrow [2, 7]$

$\Rightarrow (2, 8) \quad (4, 6) \Rightarrow [2, 8]$ 

$\Rightarrow (3, 7) \quad (4, 10) \Rightarrow [3, 10]$ 

$\Rightarrow (3, 6) \quad (6, 10) \Rightarrow [3, 10]$ 

$\Rightarrow (2, 5) \quad (8, 10) \Rightarrow X$ 

$\Rightarrow (5, 8) \quad (1, 3) \Rightarrow X$ 

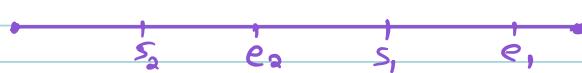
observation

① what will be the conditions for Not overlap?

I_1
 (s_1, e_1)

I_2
 (s_2, e_2)

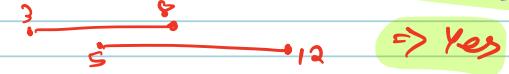
Case I  if $(s_2 > e_1)$

Case II  if $(s_1 > e_2)$

② what will be the condition if intervals overlap.

↳ If overlap \Rightarrow the ans will be

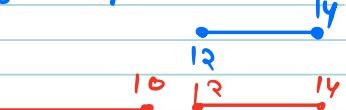
$[\min(s_1, s_2), \max(e_1, e_2)]$

Ques 1 :- $[3, 8] [5, 12]$ 

Ques 2 :- $[6, 10] [8, 15]$ 

Question :- Given a sorted list of overlapping interval, sorted based on start time, merge all overlapping intervals & return sorted list.

Ex:- $\{ (0, 2) (1, 4) (5, 6) (6, 8) (7, 10) (8, 9) (12, 14) \}$

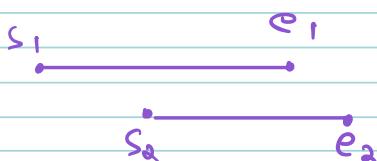


$$\text{Ans} \Rightarrow \{(0, 4), (5, 10), (12, 14)\}$$

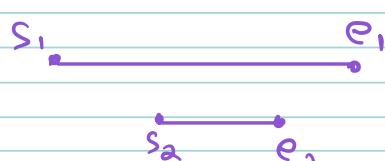
Observation

① If two intervals overlap, then condition will be.

$$I_1, \\ (s_1, e_1)$$



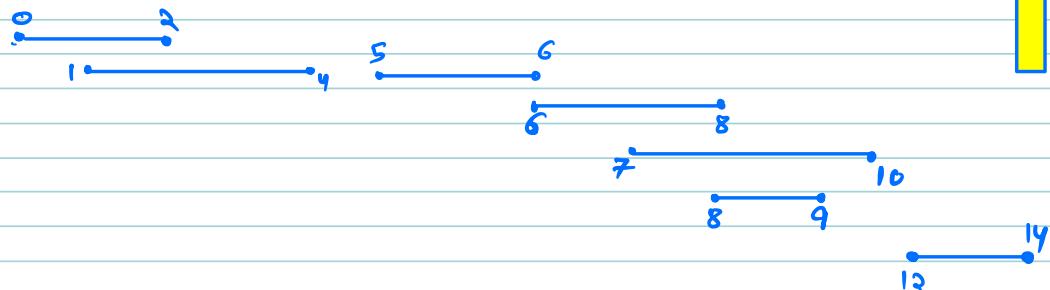
$$I_2, \\ (s_2, e_2) \quad [s_1 \leq s_2]$$



condition \Rightarrow if $(s_2 \leq e_1)$

DRY RUN

$$\text{Ex:- } \{(0, 2), (1, 4), (5, 6), (6, 8), (7, 10), (8, 9), (12, 14)\}$$



curr	Next Interval	After Margin	Answer List
(0, 2)	(1, 4)	(0, 4)	_____
(0, 4)	(5, 6)	_____	(0, 4)
(5, 6)	(6, 8)	(5, 8)	_____

$(5, 8)$	$(7, 10)$	$(5, 10)$	$\underline{\quad}$
$(5, 10)$	$(8, 9)$	$(5, 10)$	$\underline{\quad}$
$(5, 10)$	$(12, 14)$	$\underline{\quad}$	$(0, 4)(5, 10)$
$(12, 14)$	$\underline{\quad}$	$\underline{\quad}$	$(0, 4)(5, 10)$
			$(12, 14)$

* We can say there is No overlap b/c current & next Available then we will move to Ans.

* When we are able to merge then nothing goes to Final Ans because we are not sure whether anything from future will merge or not.

PSEUDO CODE

$\text{if, int[]} \text{ arr} \rightarrow \text{int}.$
$\text{over here, Interval[]} \text{ arr} \rightarrow \text{arr}[0].S$
$\text{arr}[0].E$

$\text{Interval[]} \text{ arr}; \quad // \text{ Given.}$

$\text{list } <\text{Interval}> \text{ ans};$

$\text{curr_S} = \text{arr}[0].S;$
 $\text{curr_E} = \text{arr}[0].E;$

$\text{for } (i=1; i < \text{arr.length}; i++) \{$

$\text{if } (\text{arr}[i].S \leq \text{curr.E}) \&$

$// \text{overlapping}$

$\text{curr_E} = \text{Max}(\text{curr_E}, \text{arr}[i].E);$

$\} \text{ else } \{$

$// \text{Non-overlapping.}$

$\text{Interval temp}(\text{curr_S}, \text{curr_E});$

$\text{ans. add}(\text{temp});$

$\text{curr_S} = \text{arr}[i].S;$

$\text{curr_E} = \text{arr}[i].E;$

3

Interval temp(curr_S, curr_E);

ans. add(temp);

return ans;

$$TC = O(N)$$

$$SC = O(1) \leftarrow$$

Given we
can ignore the
ans space.

Question :- Given a sorted list of non overlapping intervals based on start time, insert a new interval such that the final list of intervals is also sorted & non-overlapping.
Print the Intervals.

Ex:-

N=9

Given New
Interval
(12, 22)

Ans

(1, 3)¹

(4, 7)

(12, 22)

(4, 7)²

(10, 14)

(12, 22) : (10, 22) ——

(16, 19)

(10, 22) : (10, 22) ——

(21, 24)

(10, 22) : (10, 24) ——

4 (27, 30)

(10, 24)

(10, 24)³

5 (32, 35)

6 (38, 41)

7 (43, 50)

List

↓

——

——

——

——

——

——

——

——

——

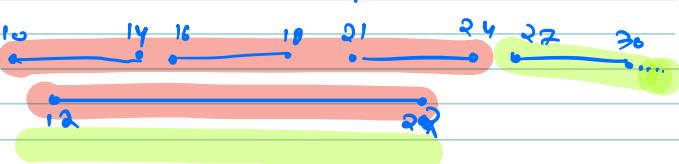
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Observation

- ① We need to move interval in Ans if its ending before given interval.
- ② If the current interval overlaps with given interval then we will move ahead without pushing anything into Ans.
- ③ Continuous step ② till you get interval which is larger than given interval. Now you can push given & all other into Ans.

Ex 2

$N = 5$

Given New
Interval

Ans

(1, 5)

(12, 24)

(1, 5)

(8, 10)

(12, 24)

(8, 10)

(11, 14)

(12, 24) : (11, 24) —

(15, 20)

(11, 24) : (11, 24) —

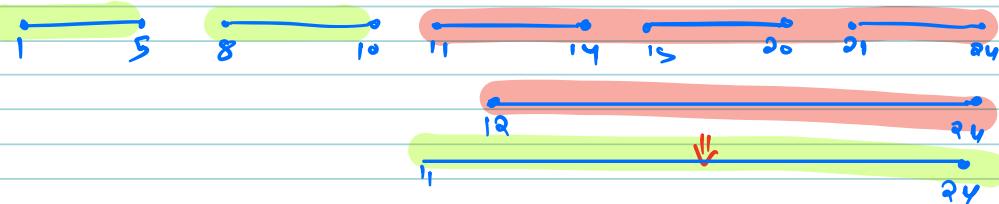
(21, 24)

(11, 24) : (11, 24) —

—

(11, 24)

(11, 24)



PSEUDO CODE

```
(ns, ne) → new Interval  
for ( i=0 ; i < arr.length ; i++ )  
    cInterval = arr[i];
```

```

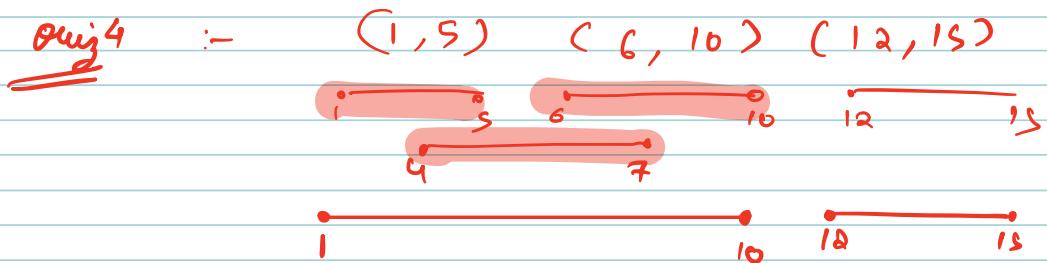
    // Non overlapping
    if ( ns > cInterval.e ) {
        Print( cInterval )
    } else if ( ne < cInterval.s ) {
        print( ns, ne );
        for ( j = i; j < arr.length; j++ ) {
            Print( arr[j].s, arr[j].e );
        }
        return;
    } else {
        ns = Min( cInterval.s, ns );
        ne = Max( cInterval.e, ne );
        Print( ns, ne );
    }
}

when new lies before current → 2
when overlapping → 3

```

<u>Ex:-</u>	$N=9$	Given Algo	Ans
	$(1, 3)$	Interval $(12, 22)$	$(1, 3)$
	$(4, 7)$	$(12, 22)$	$(4, 7)$
	$(10, 14)$	$(12, 22) \cup (10, 22)$	—
	$(16, 19)$	$(10, 22) \cup (16, 19)$	—
	$(21, 24)$	$(10, 22) \cup (21, 24)$	—
	$(27, 30)$	$(10, 24)$	$(10, 24)$
	$(30, 35)$		
	$(38, 41)$		
	$(43, 50)$		

TC - $O(N)$
 SC - $O(1)$



Question :- Given an unsorted array of integers, Find first missing Natural Number. $\rightarrow \{1, 2, 3, \dots\}$

Ex 1 $\text{arr}[5] = \{3, -2, 1, 2, 7\} \Rightarrow \text{Ans} = 4$

Ex 2 $\text{arr}[7] = \{-9, 2, 6, 4, -8, 1, 3\} \Rightarrow \text{Ans} = 5$

Ex 3 $\{-2, 4, -1, -6, 3, 7, 8, 4, -3\} \Rightarrow \text{Ans} = 1$

Ex 4 $\{1, 0, -5, -6, 4, 2\} \Rightarrow \text{Ans} = 3$

******* Ex 5 $\{1, 2, 5, 6, 4, 3\} \Rightarrow \text{Ans} = 7$

Ex 6 $\{-4, 8, 3, -1, 0\} \Rightarrow \text{Ans} = 1$

Ex 7 $\{4, 2, 1, 3\} \Rightarrow \text{Ans} = 5$

Ques 5 :- $[5, 3, 1, -1, -2, -4, 7, 2] \Rightarrow \text{Ans} = 4$

Observations

\hookrightarrow ① If length is N , then ans will be definite in range of $[1, N+1]$

Brute force

Solution 1 :- Since ans is in range [1 to N+1] so, I can search missing no. by traversing repeatedly over array.

TC - $O(N^2)$ SC - $O(1)$

Solution 2 :- sort the array \rightarrow travel

TC - $O(N \log N)$ SC - $O(1)$

Solution 3 :- Use set, insert all element in set & then search from 1 to N+1.

TC - $O(N)$ SC - $O(N)$

OPTIMIZATION

Hint: We need to use concept that ans will be in range of [1, N+1]

Ex:- $\{4, -2, 8, 9, 1, 3\}$ index $\rightarrow A[i]-1$

index	element	Bresenham Marked at index	Actual Array
-------	---------	---------------------------	--------------

0	4	3	$\{4, 2, 8, -9, 1, 3\}$
1	2	1	$\{4, -2, 8, -9, 1, 3\}$
2	8	7	X
3	9	8	X
4	1	0	$\{-4, -2, 8, -9, 1, 3\}$
5	3	2	$\{-4, -2, 8, -9, 1, 3\}$

$\rightarrow \{ -4, -2, -8, -9, 1, 3 \} \rightarrow Ans = 5$

1 to 7

\Rightarrow with negatives

Ex :- $\{1, -5, -3, 4, 2\}$

index	element	Presence marked at index	Actual array
0	1	0	$\{-1, 8, 8, 4, 2\}$
1	-5	-1	X
2	-3	—	X
3	4	3	X
4	2	1	$\{-1, 8, 8, 4, 2\}$

$\Rightarrow \{-1, \boxed{0}, \boxed{-5}, -3, 4, 2\}$

IDEA :-

If $(A[i] \leq 0) \{$

$A[i] = N+2;$ // Can't be Ans.

↳

Final :- $\{-1, -8, 8, -8, 4, 2\}$
 $\hookrightarrow \text{Ans} = 3$

PSEUDO CODE

for (i=0; i < N; i++) {

 if ($A[i] \leq 0 \{$

$A[i] = N+2;$

 }

}

```

for (c i= 0 ; i < N; c++) {
    ele = abs (A[i]);
    if (ele >= 1 && ele <= N) {
        A[ele-1] = -1 * abs(A[ele-1]);
    }
}
for (c i= 0; i < N; c++) {
    if (A[i] > 0) return i+1;
}
return N+1;

```

$$\boxed{\begin{array}{l} TC = O(N^2) \\ SC = O(1) \end{array}}$$

Ex:- 0, 4, 1, 2, 8, 3, 1, 5, 4, 5, 3, 3