

Trees 3 : Binary Search Trees (BST)

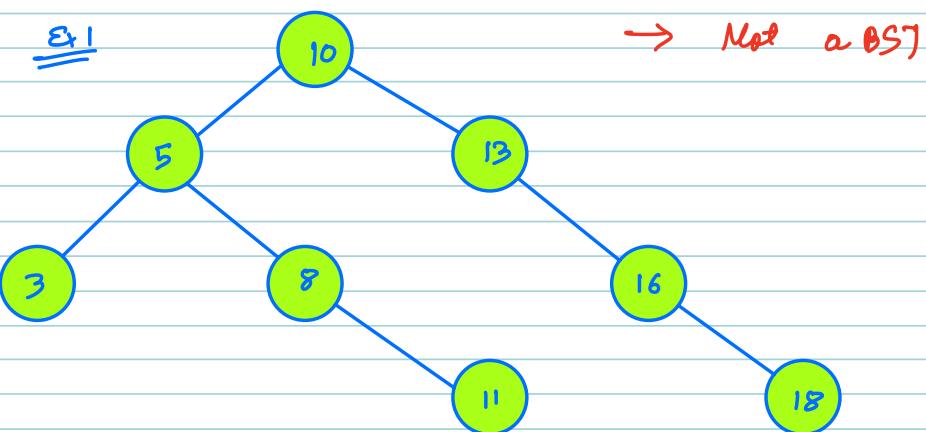
TODAY's AGENDA

- BST Basics + searching
- Insert / Search / Delete
- Construct BST from sorted array
- Is BST()

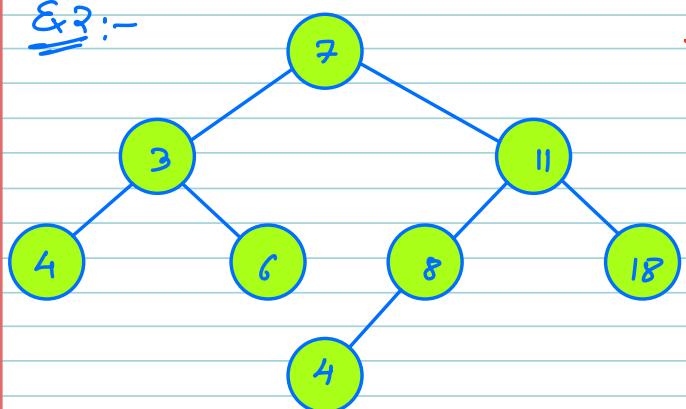
BINARY SEARCH TREES (BST)

↳ A Binary Tree is called BST if For all nodes

All element in LST < Node < All element in RST

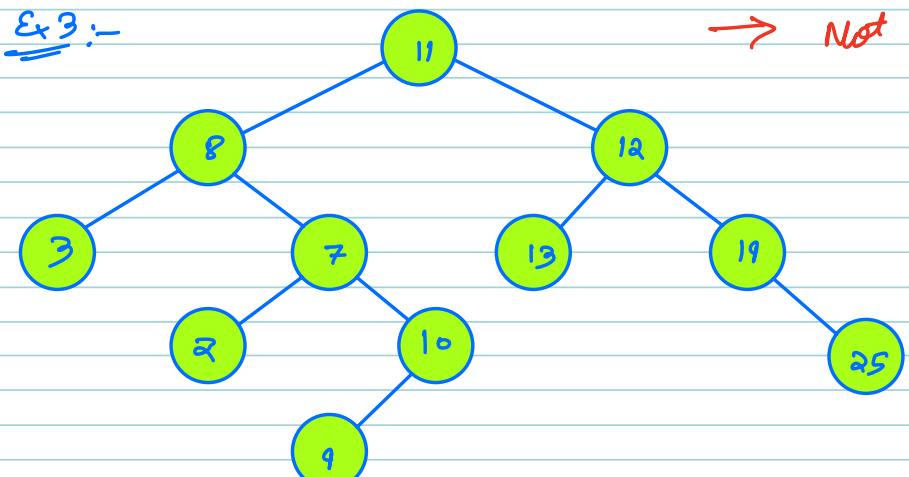


Ex 3 :-



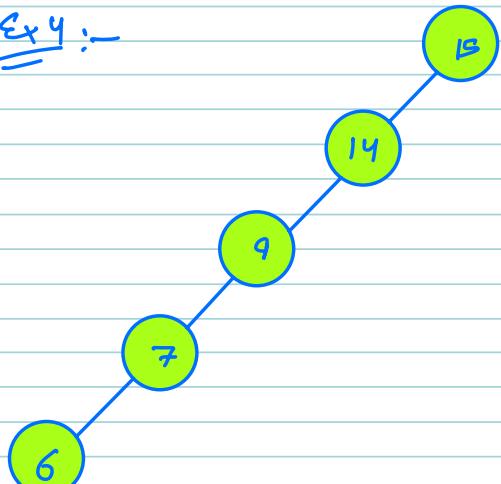
→ Not a BST

Ex 3 :-



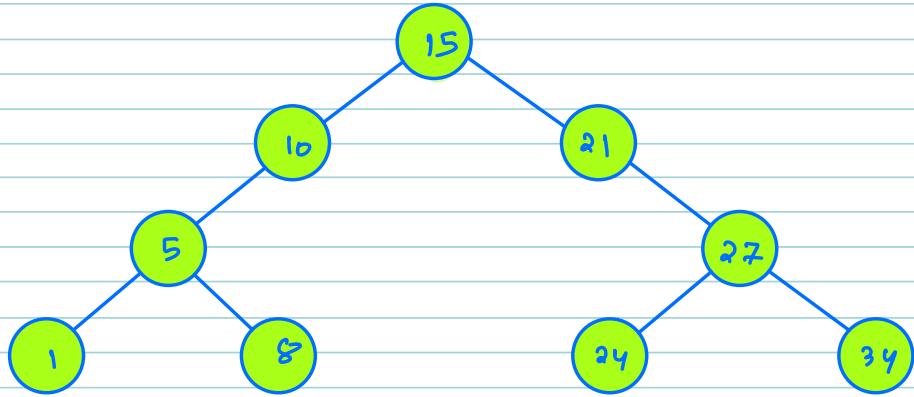
→ Not a BST

Ex 4 :-



→ It is a BST

⇒ PROPERTY OF BST



INORDER :- 1 5 8 10 15 21 24 27 34

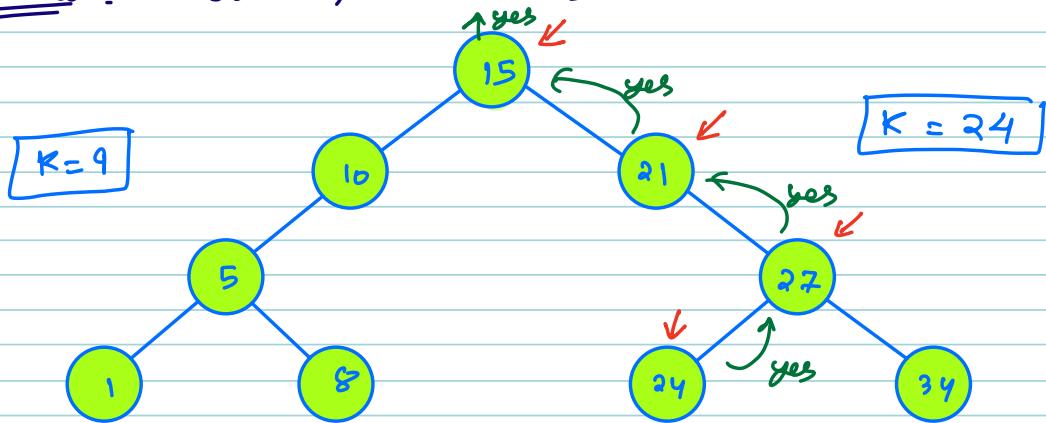
↙ observation:- Inorder of a BST is sorted.
(L D R)



Ques 1 :- what is a Binary Search Tree (BST) ?

↳ A tree where for a node x , everything on the left has data $\leq x$ & on the right $> x$.

Question:- Searching in a BST



Brute force

Do traversal & check for $\text{Node} == K$.

$$\begin{aligned} \text{TC} &\rightarrow O(N) \\ \text{SC} &\rightarrow O(1) \end{aligned}$$

\Rightarrow OPTIMIZATION

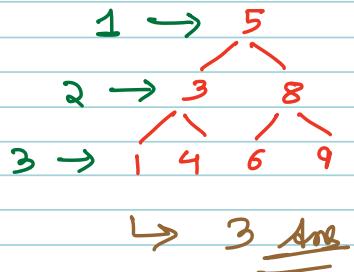
\hookrightarrow use the property that tree is a BST.

i.e. \rightarrow If a Node $>$ element, then move Left.

\rightarrow else if Node $<$ element, then move right

\rightarrow else Got data.

Ques 2 :- what is the no. of nodes you need to visit to find the number '1' in the following BST?

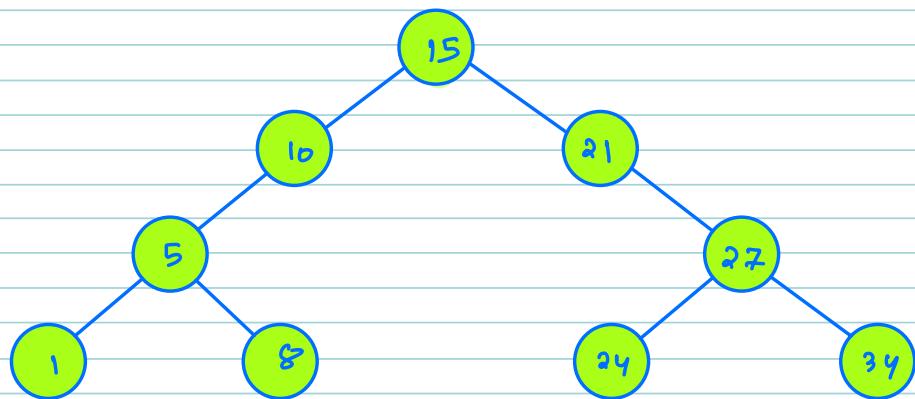


PSEUDO CODE

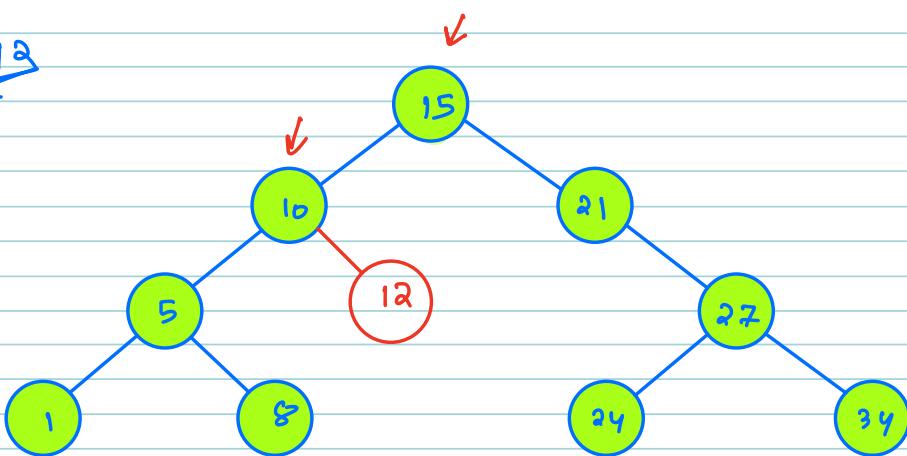
```
boolean Search ( Node root, int k ) {  
    if ( root == Null ) { return False; }  
    if ( root. data == k ) {  
        return true;  
    } else if ( root. data < k ) {  
        return search ( root. right, k );  
    } else {  
        return search ( root. Left, k );  
    }  
}
```

TC - $O(H)$
SC - $O(H)$

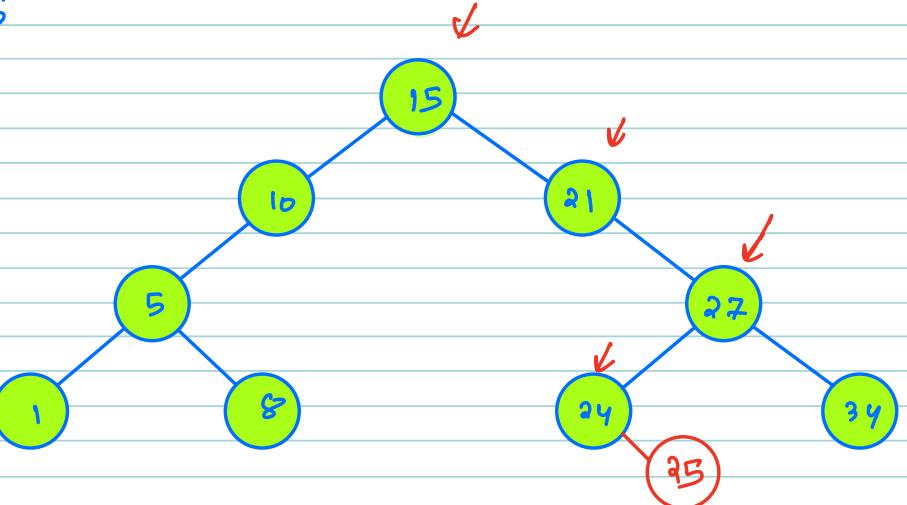
Question :- Insertion in a BST



$K = 12$



$K = 25$



IDEA :- ↳ If element > Node, call will be made to right to insert & expect root of final BST in return.

↳ If element < Node, call will be made to left to insert & expect root of final BST in return.

↳ If Node is null, create Node with value to insert & return it.

PSEUDO CODE

Mode Insert (Mode root , int k) {

 if (root == Null) {

 return new Node (k);

 if (root. data > k) {

 root.left = Insert (root. Left , k);

 } else {

 root.right = Insert (root. right , k);

 return root;

$K = 25$

DRY RUN

Node

Insert (15, 25)

\downarrow 15K
 $=15K = 25$

Node Insert (Node root, int k){

if C.root == null {
return new Node (k);
}

if C.root.data > K {
root.left = Insert (root.left, K);
}

else {
root.right = Insert (root.right, K);
}

} return root;

\uparrow 21K

Node Insert (Node root, int k){

if C.root == null {
return new Node (k);
}

if C.root.data > K {
root.left = Insert (root.left, K);
}

else {
root.right = Insert (root.right, K);
}

} return root;

\uparrow 27K

Node Insert (Node root, int k){

if C.root == null {
return new Node (k);
}

if C.root.data > K {
root.left = Insert (root.left, K);
}

else {
root.right = Insert (root.right, K);
}

} return root;

\uparrow 24K
 $=24K = 25$

Node Insert (Node root, int k){

if C.root == null {
return new Node (k);
}

if C.root.data > K {
root.left = Insert (root.left, K);
}

else {
root.right = Insert (root.right, K);
}

} return root;

\uparrow 25K

$=NULL = 25$

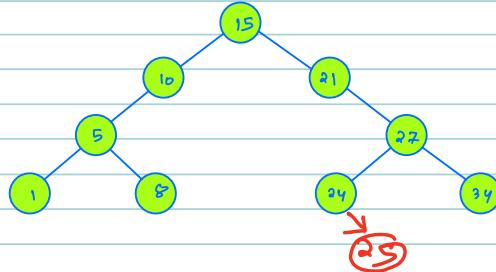
Node Insert (Node root, int k){

if C.root == null {
return new Node (k);
}

if C.root.data > K {
root.left = Insert (root.left, K);
}

else {
root.right = Insert (root.right, K);
}

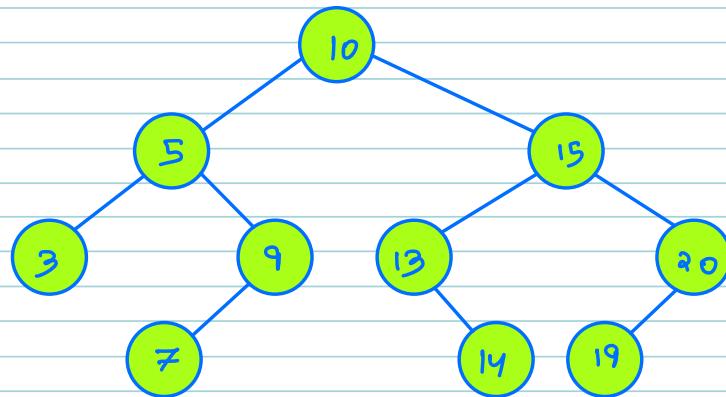
} return root;



Quiz 3 :- where does the node with the smallest value resides in a BST?

→ We need to keep on going left till we get null. And there we will find the smallest Node.

Question :- Min in BST



Idea :- We Need to go as left as possible, Not even single right as it will make the ans increased.

PSEUDO CODE

temp = root

while (temp.Left != null) {

 temp = temp.Left;

}

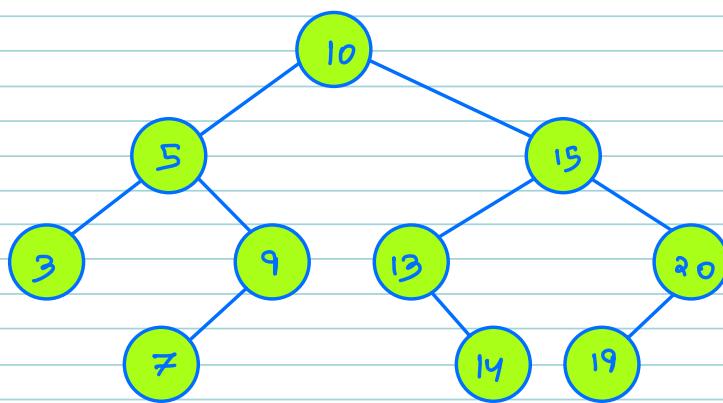
return temp.

TC → O(n)

SC → O(1)

Question

Max in B.S.T.



IDEA:- We Need to go as right as possible,
Not even single left as it will make the
Ans decreased.

PSEUDO CODE

temp = root

while (temp.right != null) {

| temp = +temp.right;
| }
| }

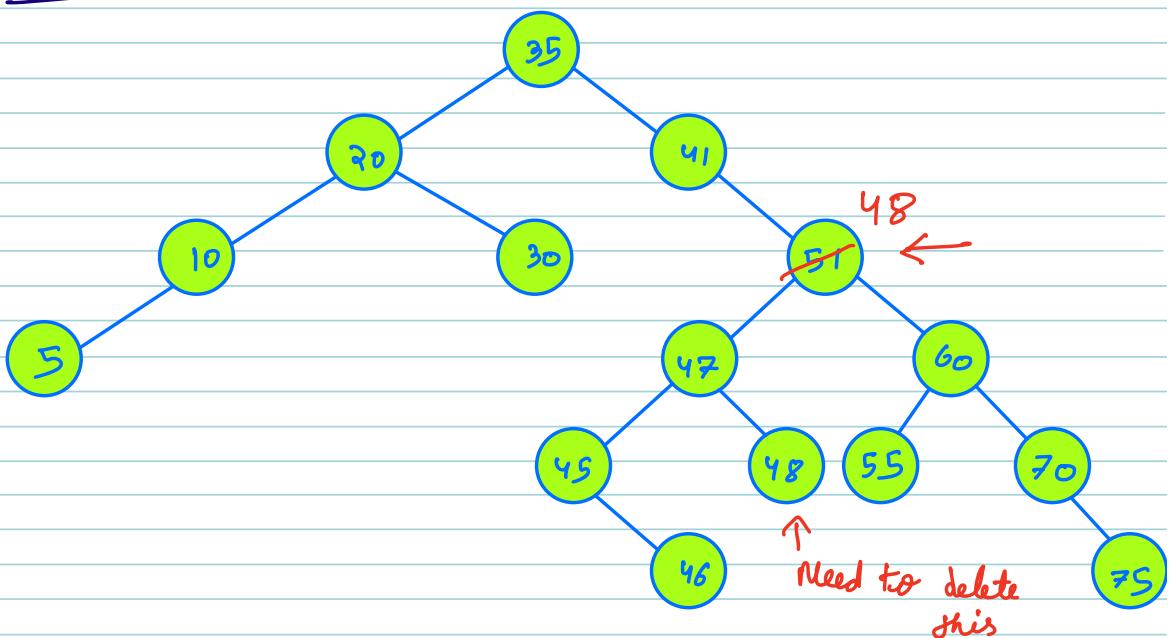
return temp.

TC $\rightarrow O(1)$

SC $\rightarrow O(1)$

10:22 ← Resume time

question :- Delete a node in BST



$K = 5$ } 0 children
 $K = 30$ }

$K = 10$ } 1 child
 $K = 41$ }

$K = 51 \rightarrow 2 \text{ children}$

Case I → 0 children

IDEA :- → If $\text{Node} > K$, move left & give responsibility of Deletion to Left & In return expect root of final LST with deleted Node.

→ If $\text{Node} < K$, move right & give responsibility of Deletion to right & In return expect root of final RST with deleted Node.

→ If Node.data == K, return Null;

Case II → 1 child

IDEA:- → If Node > K, move left & give responsibility of Deletion to left & In return expect root of final LST with deleted Node.

→ If Node < K, move right & give responsibility of Deletion to right & In return expect root of final RST with deleted Node.

→ If Node.data == K, then return the Non Null child.

Case III → 2 children

IDEA:- → If Node > K, move left & give responsibility of Deletion to left & In return expect root of final LST with deleted Node.

→ If Node < K, move right & give responsibility of Deletion to right & In return expect root of final RST with deleted Node.

Q If I want to replace 51 with any element in subtrees then which will be the best element?

↳ Max in LST, because It will be greater than all the element on left & smaller than all the element on right.

Alternative :- Min in RST

* Post Selecting Node from LST, delete A Selected Node from LST.

PSEUDO CODE

Node delete (Node root, int k) {

```
if (root.data >= k) {  
    root.left = delete (root.left, k);  
}  
else if (root.data < k) {  
    root.right = delete (root.right, k);  
}  
else {  
    if (root.left == null && root.right == null)  
        return null;  
    else if (root.left == null)  
        return root.right;  
    else if (root.right == null)  
        return root.left;
```

Case I

Case II

```

    'else {'
        int v = max C root. left);
        root. data = v;
        root. left = delete (root. left, v);
    }
    return root;
}

```

$TC \rightarrow O(H)$

$SC \rightarrow O(H)$

BALANCED BST

↳ It's there so that the height of the tree can be constrained to $(\log N)$ because when BST is balanced the height is of $(\log N)$ order.

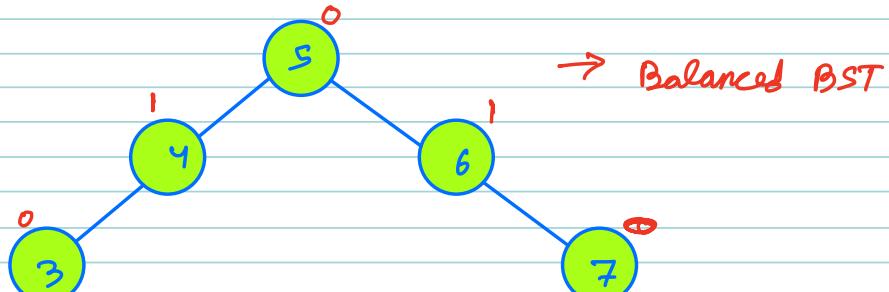
Ques :- what is the purpose of balancing a binary search tree?

↳ To maintain efficient search, insert & delete operation.

Question :- Construct a Balanced BST using sorted array.

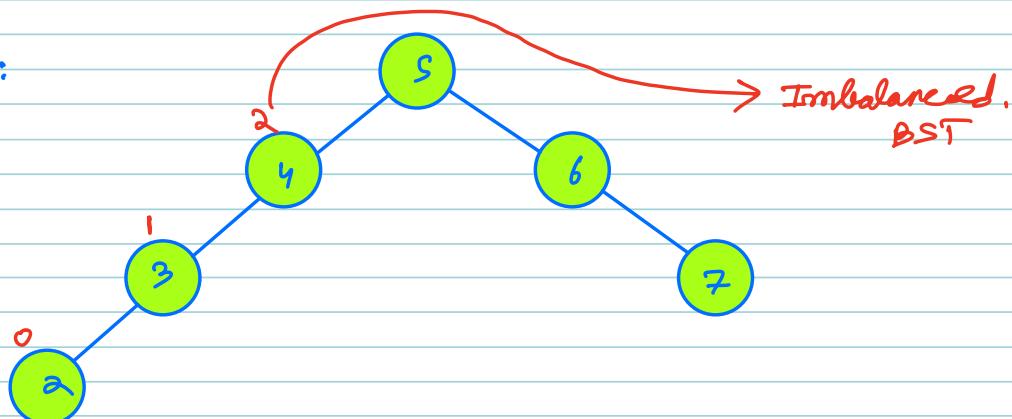
Balanced BST $\rightarrow | \text{LST height} - \text{RST height} | \leq 1$

Ex1 :-



\rightarrow Balanced BST

Ex2 :-

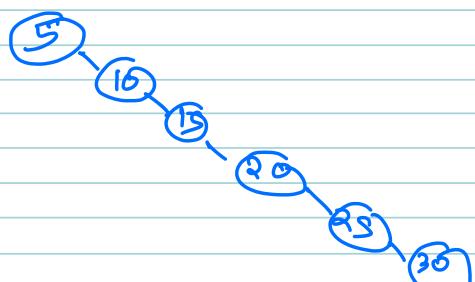


\rightarrow Unbalanced BST

Ex :- 5, 10, 15, 20, 25, 30

\hookrightarrow Create Balanced BST out of given array.

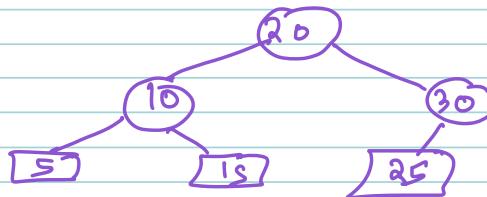
Idea :- Call Insert Method on BST for every Nob.



Problem :- A Tree formed is not a Balanced BST

IDEA? :- 5, 10, 15, 20, 25, 30

Q which Node will be best to be treated as root of Balanced BST?



Q How we should write the code?
↳ Recursion

PSEUDO CODE

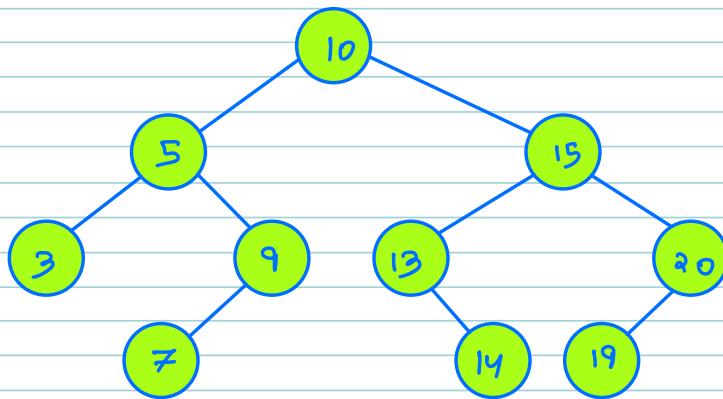
```
Node Create [int] arr, int s, int e){  
    if (s > e) { return null; }  
    m = s +  $\frac{e-s}{2}$ ;
```

```
    Node root = new Node (arr[m]);  
    root.left = Create (arr, s, m-1);  
    root.right = Create (arr, m+1, e);  
    return root;
```

$$TC \rightarrow O(N)$$

$$SC \rightarrow O(\log N)$$

Question :- Given BT, check whether it is a BST or not?



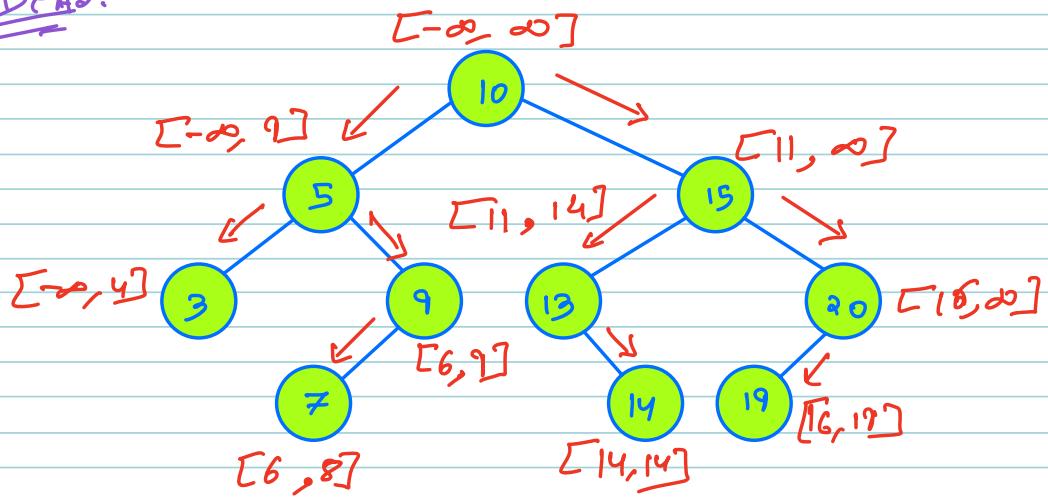
IDEA 1 :- Create ArrayList of Inorder for the given tree. And check if it is sorted or Not.

$$TC = O(N)$$

$$SC \leftarrow O(1)$$

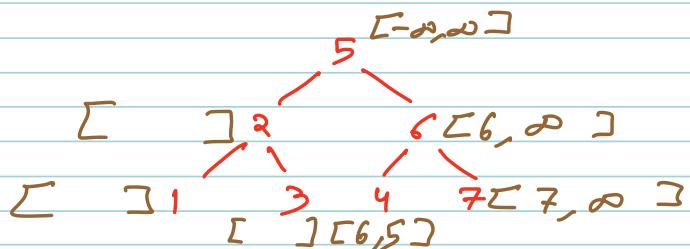
↳ Code :- HW [must to try]

I DEAD.



- Steps :-
- ① when going left update right (i.e. Max)
 - ② when going right update left (i.e. Min).
 - ③ For tree to be a BST, every Node should be in a given range.

Ques :- check whether the given BT is BST or not?



↳ Ans No

PSEUDO CODE

```
boolean is BST (Node root, int s, int e){  
    if (root == NULL) { return True; }  
    if (s <= root.data & & root.data <= e) {  
        l = is BST (root.left, s, root.data-1);  
        if (l == False) {  
            return False;  
        }  
        r = is BST (root.right, root.data+1, e);  
        if (r == False) {  
            return False;  
        }  
        return True;  
    }  
    return False;  
}
```

TC $\rightarrow O(N)$
SC $\rightarrow O(H)$