

Graph 3 : MST (Prims Algorithm) & Dijkstra Algorithm

AGENDA

→ PRIM's ALGO

→ DIJKSTRA ALGO

Question:- Given N Islands & Cost of Construction of bridge b/w Multiple pair of Islands. Find the min cost of construction required such that it is possible to travel from one Island to another Island via bridges. If not possible return -1.

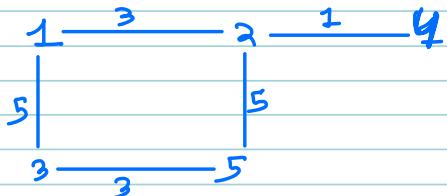
Ex:-

$N=5$

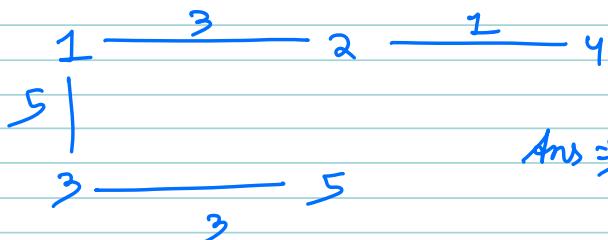
$$\begin{array}{c} 1 \xrightarrow{3} 2 \\ | \quad \quad | \\ 1 \xrightarrow{5} 3 \\ | \quad \quad | \\ 2 \xrightarrow{1} 4 \end{array}$$

$$2 \xrightarrow{5} 5$$

$$3 \xrightarrow{3} 5$$

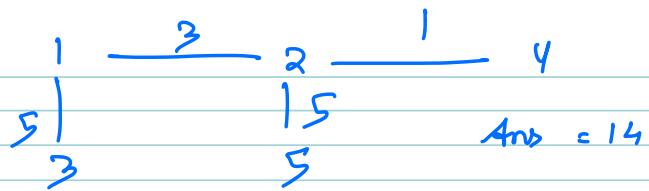


TRY1:-



Ans \Rightarrow 12

TQY2 :-



Q How Many Edges I need given N Nodes

$$\hookrightarrow N-1$$

$$\begin{aligned} \text{Ex:- } N &= 5 \\ \text{edges} &= N-1 = 4 \end{aligned}$$

⇒ SPANNING TREE

↳ A Spanning Tree is a Subgraph of a graph that includes all the nodes, maintains connectivity & has no cycle.

Note:- It does not give Shortest Distance from one point to other ^(weight)

⇒ MINIMUM SPANNING TREE

Tree generated from a Connected graph such that all nodes are connected & sum of weight of all selected edges is Minimum

SAMPLE PROBLEM STATEMENT

Scenario:

Suppose Flipkart has N local distribution centers spread across a large metropolitan city. These centers need to be interconnected for efficient movement of goods. However, building and maintaining roads between these centers is costly. Flipkart's goal is to minimize these costs while ensuring every center is connected and operational.

Goal: You are given number of centers and possible connections that can be made with their cost. Find minimum cost of constructing roads between centers such that it is possible to travel from one center to any other via roads.

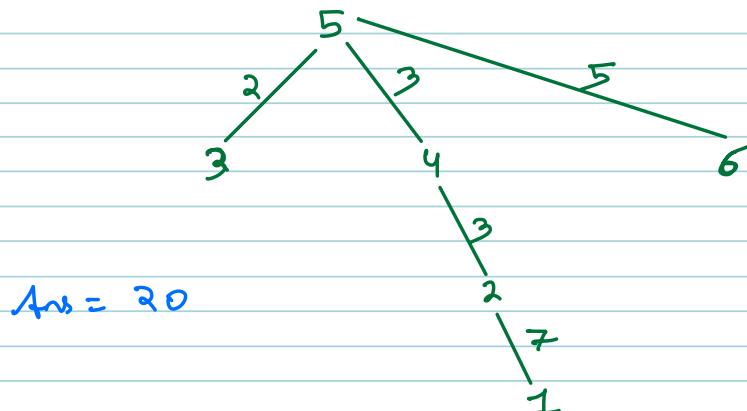
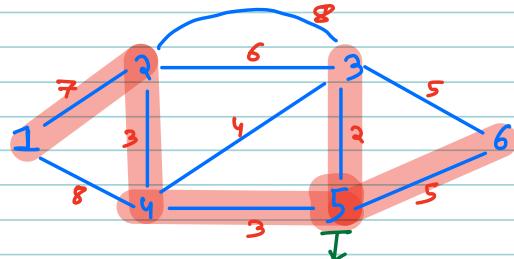
⇒ ALGORITHM WHICH HELPS WITH MST

① Prims Algo

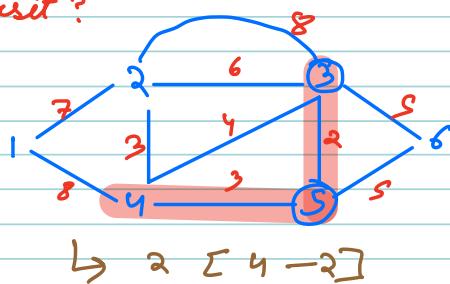
② Kruskal Algo → (DSA 4.2)

PRIMS ALGORITHM

Goal :- Find Minimum Spanning Tree



Ques 1 :- which is the next node, we need to visit?

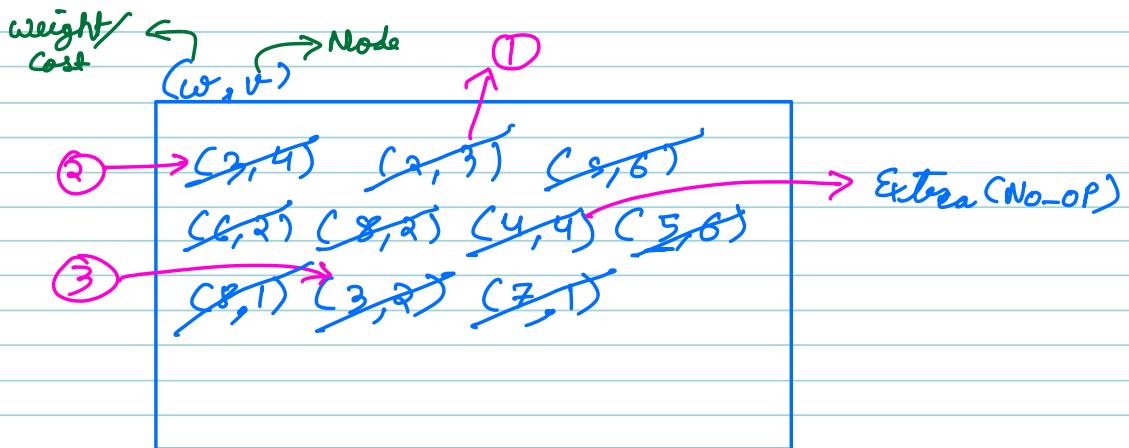
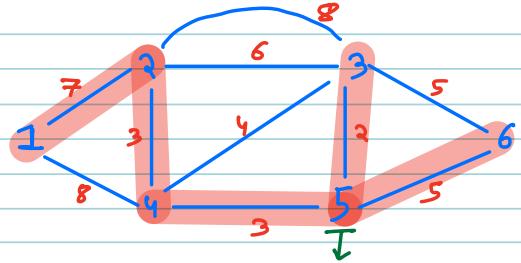


$\hookrightarrow 2 [4 - 2]$

Ques 2 :- what is the most suitable data structure for implementing prim's algorithm?

$\hookrightarrow \text{Heap}$.

\Rightarrow one more iteration



Q why $(4, 4)$ got added when there is another way to reach 4 with lesser weight
↳ 4 was not visited when we have inserted $(4, 4)$

PSEUDO CODE

Step 1 :- Select any node & add all its neighbour in Min heap.

Step 2 :- while ($Mh.\text{size}() > 1$) {

 (w, v) ← pair $p = Mh.\text{getMin}();$
 $Mh.\text{removeMin}();$
 if ($Vis[\Gamma[w]] == \text{true}$) {
 |
 | continue;
 }
 ans $\leftarrow ans + w;$
 $Vis[\Gamma[w]] = \text{true};$
 for ((w, v) in $adj[\Gamma[w]]$) {
 |
 | if ($Vis[\Gamma[v]] != \text{true}$) {
 | |
 | | $Mh.\text{add}(w, v);$
 | }
 }
 return ans;
 }

TC

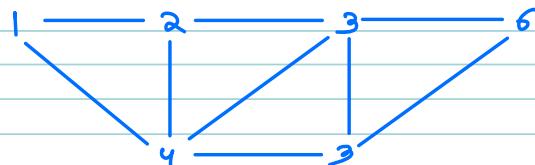
Max Size of Heap = Σ

$$TC \rightarrow \varepsilon \log \varepsilon$$

SC $\Rightarrow O(\epsilon + N + \text{adj List}) \xrightarrow{\quad} O(N + \epsilon)$
 $\Leftrightarrow O(N + \epsilon)$

question:- Find min Edges to travel from u to v .

Ex:-



$u = 1$ }
 $v = 6$ } shortest path

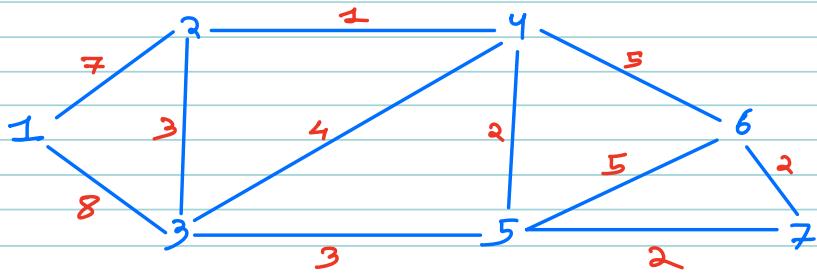
Ans = 3
[Simple BFS]

\Rightarrow The BFS will fail when edges have weight with min weight.

\hookrightarrow This can be solved Using DIJKSTRA

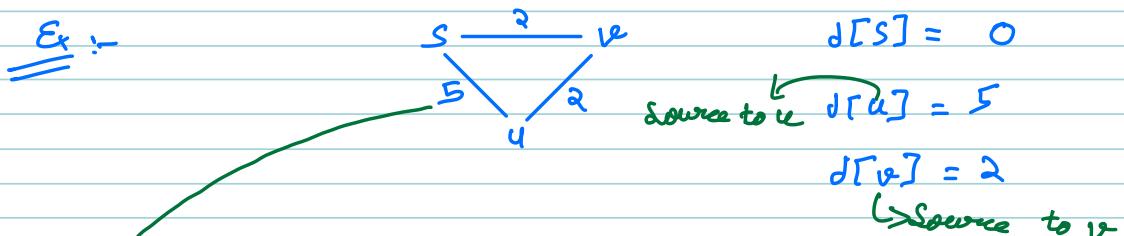
DJIKSTRA ALGO

question :- There are n cities in a Country, you are living in city 1. Find min distance to reach every other city from City 1.



⇒ Single Source shortest path → Dijkstra

$d[i] \rightarrow d[i]$ min dis to reach from source to i^{th} node.

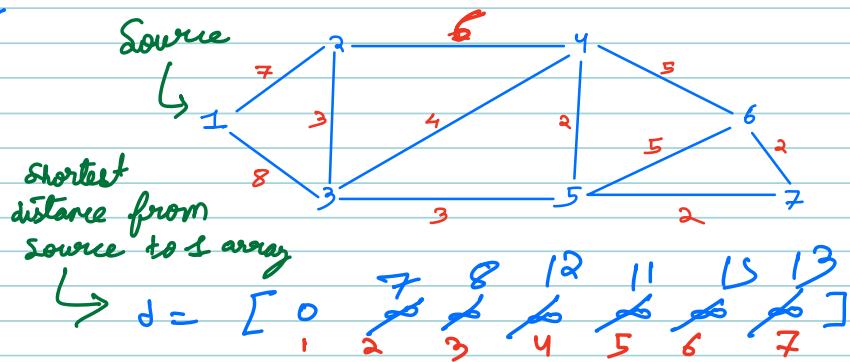


if $d[S-u] > d[S-v] + d[v-u]$

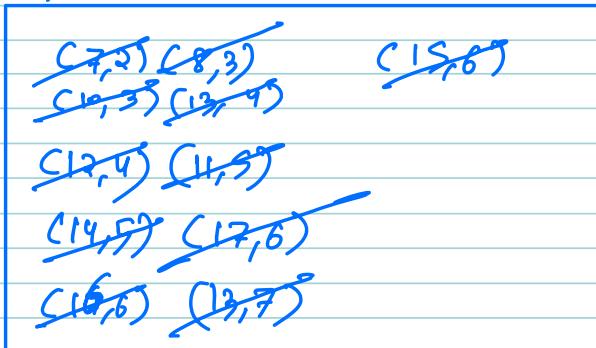
then $d[S-u] = d[S-v] + d[v-u]$

Relaxing an edge.

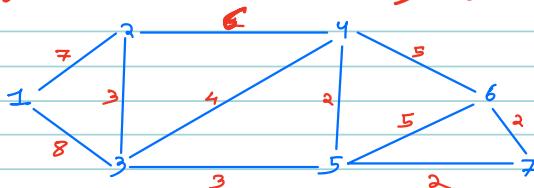
Ex



(5,4)



Ques 3 :- The current states are, Given Graph



min heap = [(10,3) (13,4) (11,5) (12,6)]

$$d[] = [1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7] \\ d[] = [0 \ 7 \ 8 \ -\infty \ -\infty \ -\infty \ -\infty]$$

what would be the next step, which is going to take place?

→ pick (10,3) & Don't update the distance array.

PSEUDO CODE

$\text{dist} \rightarrow \{ \text{TODO} [\text{all } \infty] \}$

$\text{dis}[s] = 0;$

Add all the neighbours of source to mh.
 $\hookrightarrow C(d, u)$

while $C \text{ mh.size}(c) > 0 \}$
 (d, u) pair $p = \text{mh.getmin}()$
if $C \text{ dis}[u] < d \}$
| Continue;
|
|
|

$\text{dis}[u] = d;$

for $C(w, v)$ of neighbours $\} S$

$\text{nbr} = \{ w, v \};$

$\text{new_distance} = \text{dis}[u] + \omega;$
if $C \text{ new_distance} < \text{dis}[v] \}$
| mh.add($\{ \text{new_distance}, v \}$);
|
|

s.

i

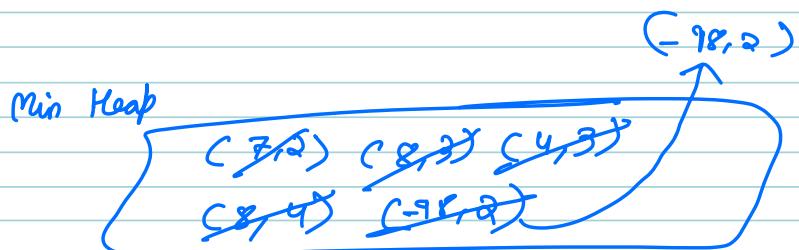
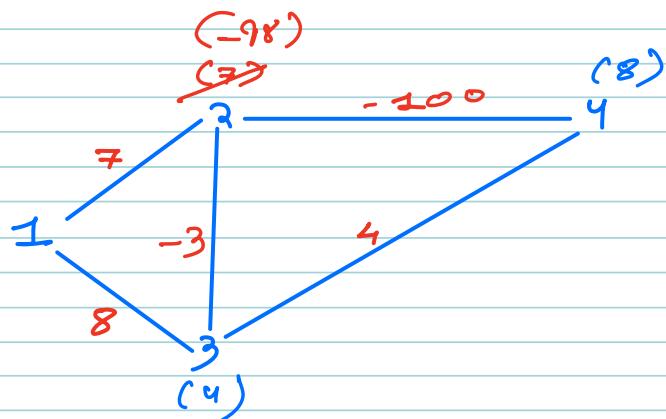
get $\text{dis};$

$T_C \rightarrow O(E \log E)$

$SC \rightarrow O(N+E)$

IMP NOTE

DIJKSTRA WON'T WORK OVER -ve NOS.



* For -ve nos \Rightarrow Bellman Ford Work

