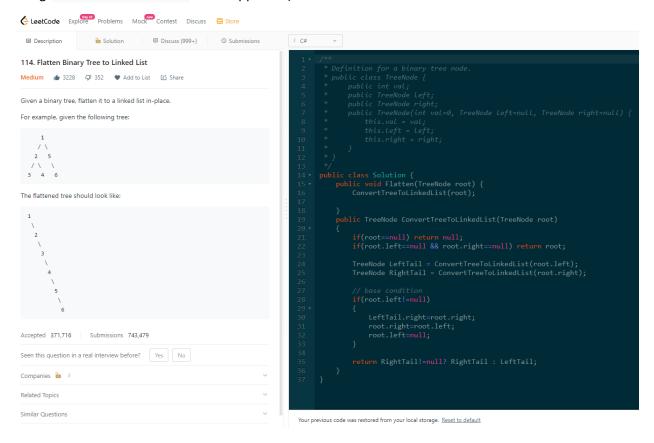
// https://leetcode.com/problems/flatten-binary-tree-to-linked-list/ (check last solution with O(1) space
using 'Morris Traversal' based approach)



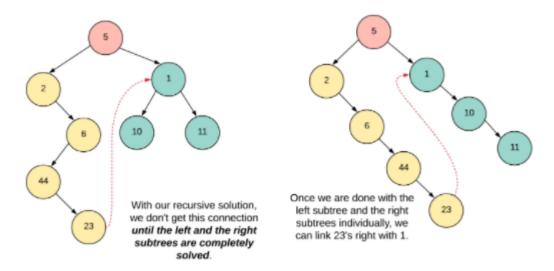
```
/**
```

- * Definition for a binary tree node.
- * public class TreeNode {
- * public int val;
- * public TreeNode left;
- * public TreeNode right;
- * public TreeNode(int val=0, TreeNode left=null, TreeNode right=null) {
- * this.val = val;
- * this.left = left;
- * this.right = right;
- * }
- * }

```
*/
public class Solution {
  public void Flatten(TreeNode root) {
    ConvertTreeToLinkedList(root);
  }
  public TreeNode ConvertTreeToLinkedList(TreeNode root)
  {
    if(root==null) return null;
    if(root.left==null && root.right==null) return root;
    TreeNode LeftTail = ConvertTreeToLinkedList(root.left);
    TreeNode RightTail = ConvertTreeToLinkedList(root.right);
    // base condition
    if(root.left!=null)
       LeftTail.right=root.right;
      root.right=root.left;
      root.left=null;
    }
    return RightTail!=null? RightTail: LeftTail;
  }
}
```

// Morris Traversal based soln

With recursion, we only re-wire the connections for the "current node" once we are already done processing the left and the right subtrees *completely*. Let's see what that looks like in a figure.



However, the postponing of rewiring of connections on the current node until the left subtree is done, is basically what recursion is. Recursion is all about postponing decisions until something else is completed. In order for us to be able to postpone stuff, we need to use the stack. However, in our current approach we want to get rid of the stack altogether. So, we will have to come up with a greedy way that will be costlier in terms of time, but will be space efficient in achieving the same results.

For a current node, we will check if it has a left child or not. If it does, we will find the last node in the rightmost branch of the subtree rooted at this left child.

Once we find this "rightmost" node, we will hook it up with the right child of the current node.

Let's look at this idea on our current sample tree.

Let's say our current node is the root node of the tree. This node does have a left child. So, we will find the final node in the rightmost branch



Once we do find this

class Solution {

public void flatten(TreeNode root) {

```
// Handle the null scenario
if (root == null) {
  return;
}
TreeNode node = root;
while (node != null) {
  // If the node has a left child
  if (node.left != null) {
    // Find the rightmost node
    TreeNode rightmost = node.left;
    while (rightmost.right != null) {
       rightmost = rightmost.right;
    }
    // rewire the connections
    rightmost.right = node.right;
    node.right = node.left;
    node.left = null;
  }
  // move on to the right side of the tree
  node = node.right;
}
```

}