

- i) The existing new efforts to make computer think. Machines with mind in the full and literal sense.
 - ii) The automation of activities that we associate with human thinking.
 - iii) The art of creating machines that perform functions that require intelligence when performed by people.
 - iv) The study of how to make computers do things at which, at a moment, people are better. (mimic human behavior)
 - v) The study of computations that make it possible to perceive, reason and act.
 - vi) The branch of computer science that is concerned with the automation of intelligent behavior.
- The pre-requisite for K.S is to know about the problem or have knowledge about it. Then:
- I] Representation of knowledge [rules
various techniques like predicate logic,
Syntactic etc. are available.
[Prolog works on predicate logic]
 - II] Reasoning, Problem Solving, decision making

III] Action → Valid.

When a tester is unable to distinguish whether the answer is coming from a human (expert) or a machine, the machine is said to have passed the Turing Test.

31/7/19 AI problem

which requires some sense to be solved.

AI technique

There is no specific technique available. Any technique which gives a solution to the AI problem can be classified as AI technique.

Steps:

→ Define the problem

Input state → goal state

Representation

Analyze

No. of input & goal states

Identify the pattern (how can it be solved based on the problems solved previously)

Technique

⇒

1	2	3
4	6	8
7	5	

Initial

1	2	3
4	5	6
7	8	

i) calculate the no. of tiles which are out of place.

Initial

3

Move 8 down 3

Goal

0

Move 5 right 3

0

0

Here, the no. is same

but, if the no. increases,

that move can be dropped.

ii) calculate the no. of steps by which each tile is away from its goal.

Initial

$$0 + 0 + 0 + 0 + 1 + 1 + 0 + 2 \\ = 4$$

Goal

0

Move 8 down

$$0 + 0 + 0 + 0 + 1 + 1 + 0 + 1 \\ = 3$$

0

Move 5 right

$$0 + 0 + 0 + 0 + 2 + 1 + 0 + 2 \\ = 5$$

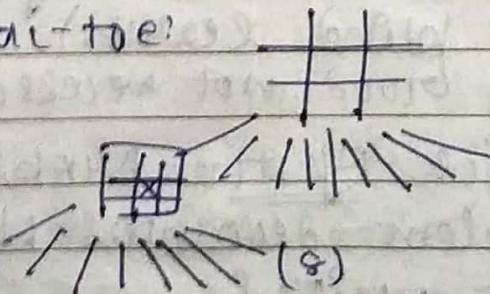
0

Hence, moving 5 right is not feasible.
as it does not lead to the goal state
further.

Problem state space:

(Space in which all the states are available)

Tic-tac-toe:



(9)

/ / / / / (8)

This problem has multiple goal states

9 layers

state space search - searching for the exact space which can lead to the goal state(s).

Production system (Expert system)
Involves production rules and solution.

- Define the problem (I & B states)
- Represent
 - Rules → set of rules/ action.
valid step.
- Solution → Rule Applier
(which rule to be applied which can take us nearer to goal state)

Control Strategies

Control Strategy

It should be such that:

- i) give you motion. - the task does not stop and give a new state with every move.
- ii) systematic - The move should be able to answer the question of why the particular move was taken.
It is also called informed search.
(Opposite - blind search)
- iii) efficient - Good, not necessarily best, solution.

characteristic of the problem

- Is the problem decomposable into smaller independent units?
- Can solution steps be ignored or undone?

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ignore - at any point of time, we can ignore all the steps so far and start from the beginning.

undo - go back to the (immediate) previous state. Problems in which this can be done are called recoverable problems.

[Unrecoverable problems are the ones in which step cannot be undone]

- Is universe (of the problem) predictable? Can the result be predicted based on the steps taken.
- Is a good solution absolute or relative? Whether the goal depends on no. of states or does the answer is what matters in the end.
Playing chess - relative - depends on opponent

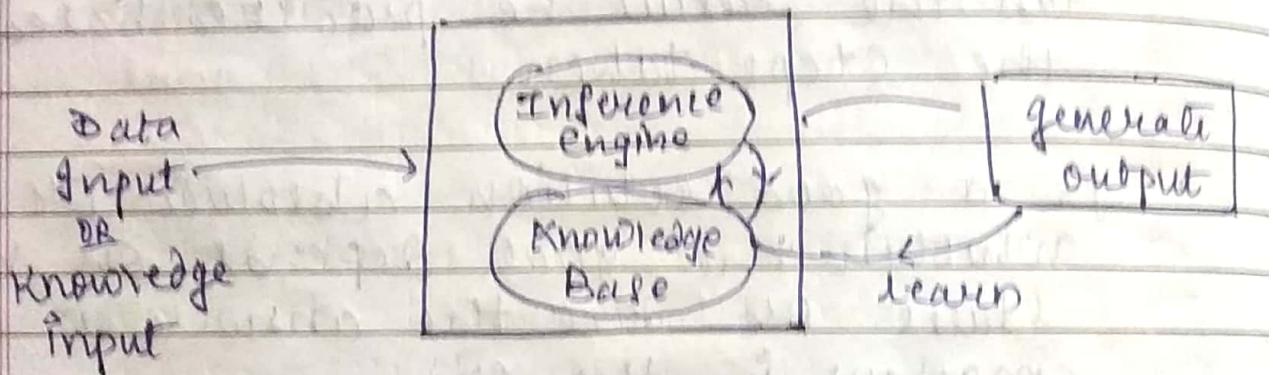
8-tile puzzle - absolute

Best path problem: Travelling Salesman problem.

- Is the solution a state or path?
- What is the role of knowledge?
Knowledge is important only to constrain the search for solution. It is required even to be able to identify or recognize a solution.

- Does the task require interaction with person?
 - Ex - where a computer is given a problem and it generates the result without any further involvement from the user, it is called problem-in solution-DIL.
 - Ex (where interaction req.) medical diagnosis.

Production System



Production System Rules

can be categorized as deductive or abductive inference rules.

Recognize-Act cycle for production system

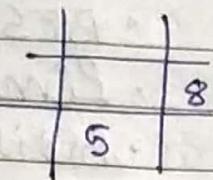
1. Match
2. Conflict Resolution (Strategy to choose which rule to apply)
3. Apply
4. Check

Classes

- i) Monotonic PS - Application of one rule should not affect the application of any other rule at any

given time.

Ex:



2 Rules: \downarrow or \rightarrow

Application of \downarrow will result into $5 \rightarrow$ never not being applied in the next step. Thus, it is not monotonic.

ii) Partially commutative PS - Application of particular seq.

leads from state x to y , then any permutation of that seq. will also lead from state x to y .

\Rightarrow

Monotonic

Non-monotonic

Partially commutative

NOT partially commu.

- \Rightarrow • Missionaries & cannibals problem
- Tower of Hanoi
- Block world
- Monkey Banana
- Cryptarithmetic \rightarrow SEND (no. aligned)
+ MORE to each
MONEY letter &
the addition
should stand
true)
- 8 - tile puzzle .
- Tic-tac-toe .

Problem solving by Search

Informed search

Uninformed
Search /
Blind Search

VIS/BS

- LS
 - Heuristic search
 - Hill climbing → simple
↳ steepest ^{ascent} HC
 - Best first search
 - A* (star)
 - A* optimal
 - Branch & bound
 - Randomized
- DFS
- BFS
- Limited DFS
- Bidirectional search
- Iterative deepening DFS

Evaluation of algorithm

- Correctness
- Completeness
- Optimality
- Time - complexity
- Space - complexity

If the algorithm can guarantee to arrive at the correct result (solution) if there exists any, it is said to be complete.

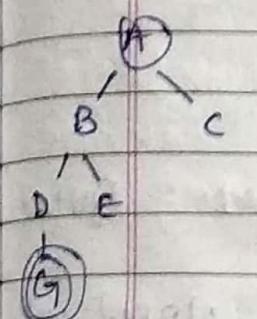
The algorithm not only gives the solution but, gives the best one, if it is optimal.

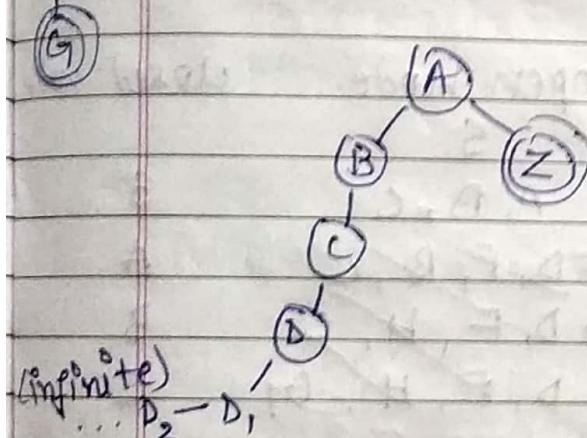
DFS

1. In the initial state, if goal state quit and return success.
2. Otherwise, do the following until success or failure is signaled.
 - a. Generate successors, for the ^{new} initial state. If there are no more successors, signal failure.

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- Page No. _____
Date _____
- b. Call D-F-S with E as initial state.
 - c. If success is returned, signal success,
otherwise continue in this loop.
- Data structure - Stack.
Also, Open node Used (Visited) node

	A B, C D, E, C G, E, C	A B D G
--	---	------------------------------



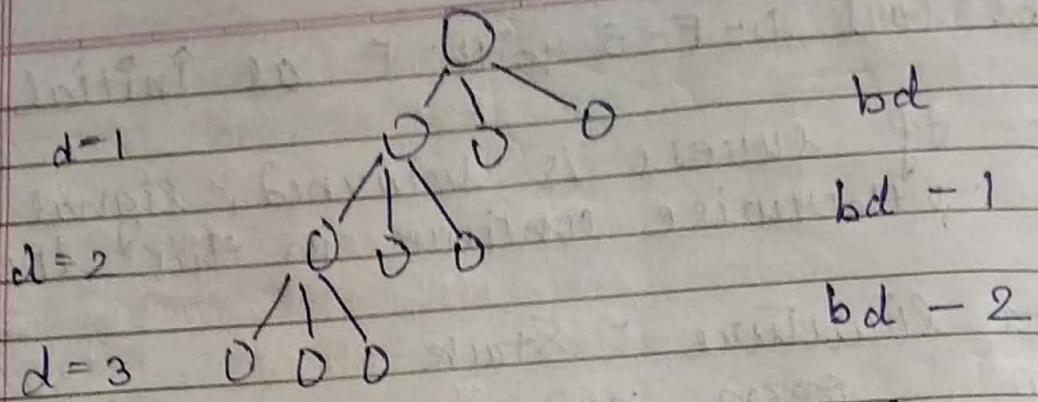
DFS does not guarantee to find the solution (can enter into loop which can lead away from goal). Hence, it is not complete.

It is not optimal either as it does not always guarantee to find the correct solution.

Branching factor - NO. of branches generated per node.

8/7/19 Let b be the branching factor and d be the depth of the tree.
 \therefore Time complexity : $O(b^d)$

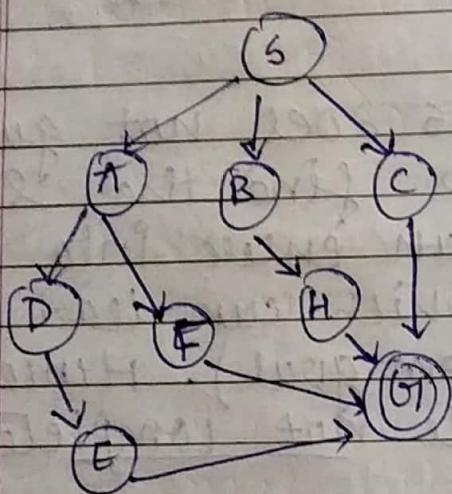
Space complexity : Let $b = 3$
then, tree :



$\therefore \text{Space complexity} = O(b^d)$

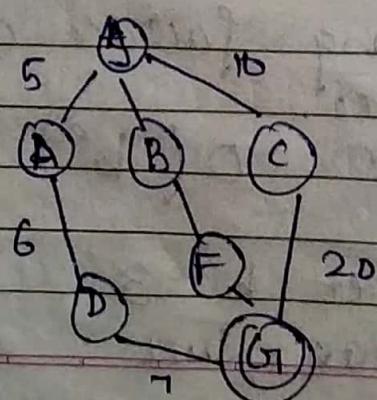
Breadth First Search

Data Structure : Queue



open node	closed node
S	
A, B, C	S
B, C, D, F	A
D, F, H	B
G	C

BFS can be complete if the branching factor is limited or finite.
If the depth is infinite, there exists no solution. Hence, completeness does not come into picture at all.



DFS: S - A - D - G

BFS: S - C - G

In this case, DFS is better.

But, if we have uniform cost.

- BFS always finds the nearest solution
- ∴ BFS is optimal only when we have uniform cost graph

Time complexity : $O(b^d)$

Space complexity : $O(b^d)$

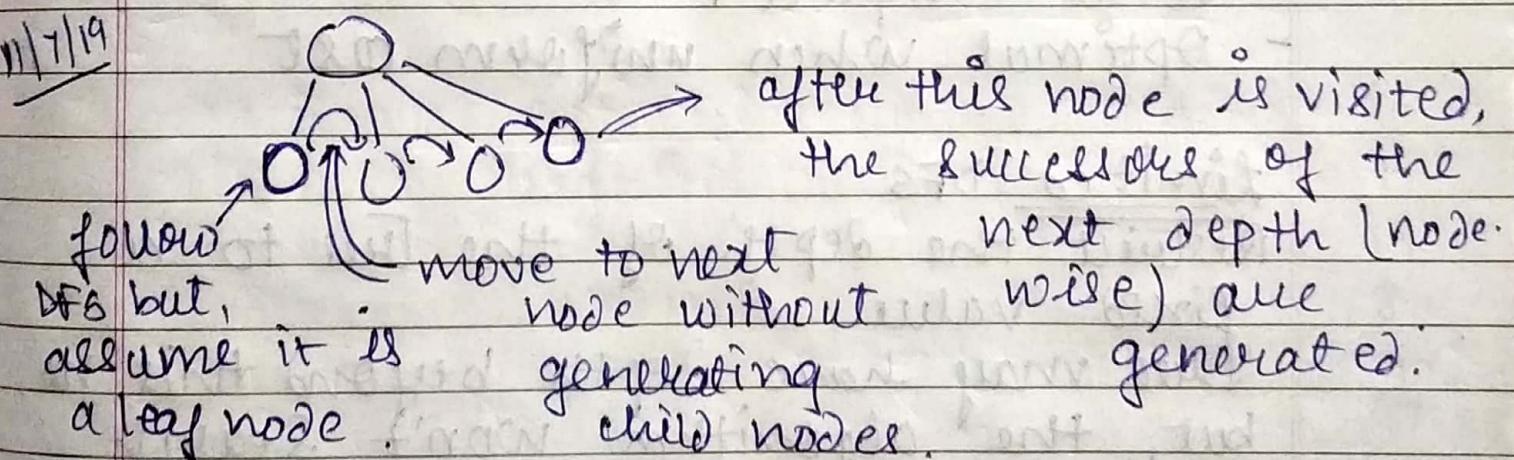
→ DFS is better than BFS in terms of space complexity (time complexity is same)

Search progress is better in BFS than DFS.

Iterative Deepening DFS :

- work progress like BFS
- space complexity DFS.

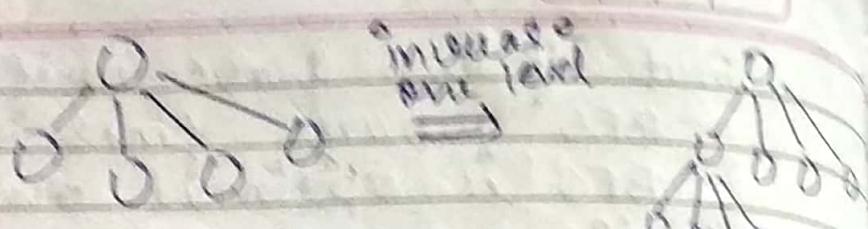
(Exhausts every level before moving to next)



Thus, we are iteratively moving in depth of the tree.

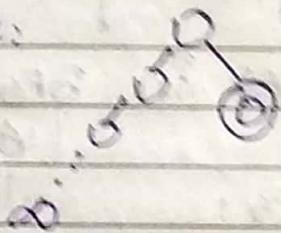
Time complexity :
Assume $b=4$.

remove the nodes which aren't goal nodes from the memory.

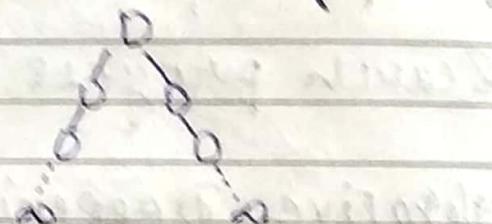


We make sure that no level is left unvisited helps to solve the infinite loop in terms of depth. (cannot work in entire tree is infinite)

Ex:



works



cannot work

- It'll complete.
- Optimal when uniform cost

Limited DFS

Recruit the depth of the tree to a fixed value.

You may have depth beyond this no. but, the algorithm won't search beyond the fixed level.

Not complete as the solution may be at the very next level. Complete only if the solution lies inside the limited depth.

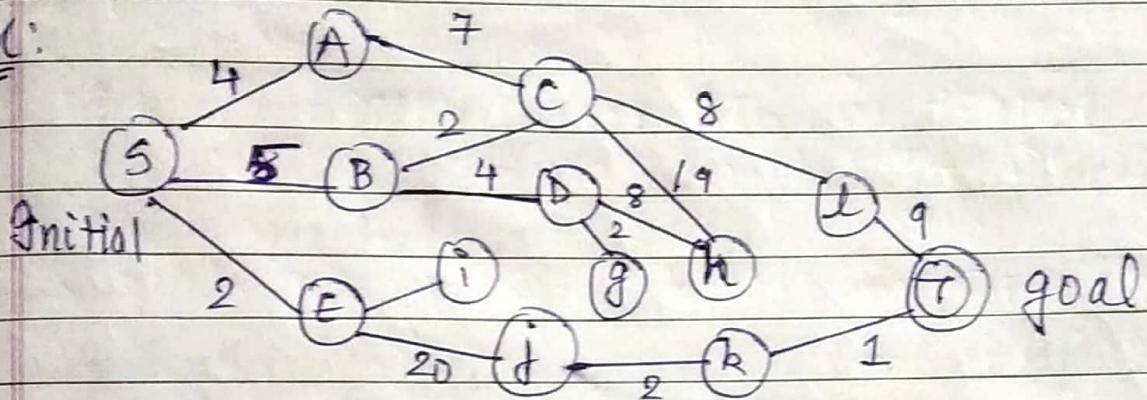
Not optimal either as if the solution does not exist, there is no point in finding the best one.

Time complexity = $O(b^l)$

Space complexity = $O(bl)$

where, l = limited depth.

Ex:



DFS:

open nodes

- S
- A, B, E
- C, B, E
- H, L, B, E
- L, B, E
- T, B, E
- B, E

closed nodes

S	
A	4
C	7
H	19
L	8
T	9
28	

BFS:

open nodes

- S
- A, B, E
- B, E, C, A
- E, C, D
- C, D, I, J
- I, J, H, L

closed nodes

S	
A	4
B	
E	
C	

(we need
to consider
node C from
B as well to

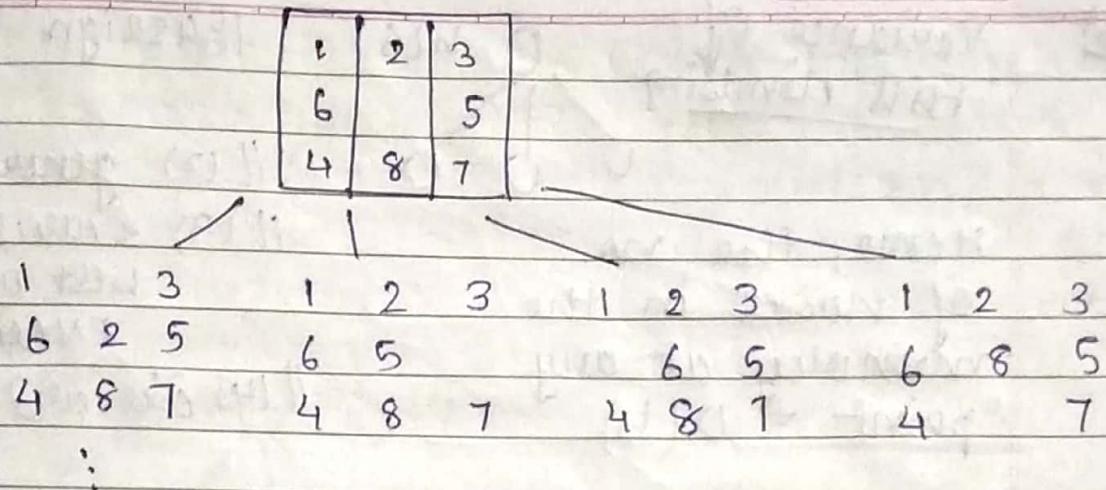
get the optimal
result because we don't have uniform cost here)

i, j, h, l, g
j, h, l, g
h, l, g, k
l, g, k
g, k, T
k, T
T, E

D
o
i
h
l
g
k
T

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- Heuristic / Informed Search
Every node has some information regarding the goal state, i.e. how good or bad a node is which on visiting might lead us nearer or far from the goal state.

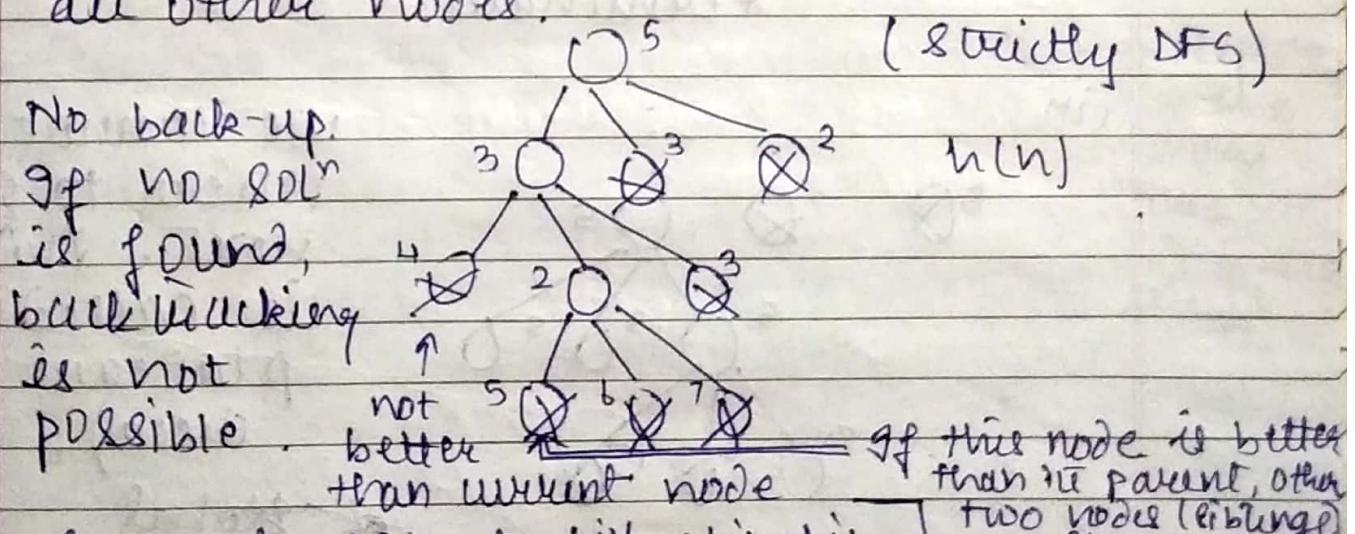


heuristic
funct.
 f_1 = no. of tiles at incorrect places
 f_2 = no. of tiles at correct places
 f_3 = manhattan distance

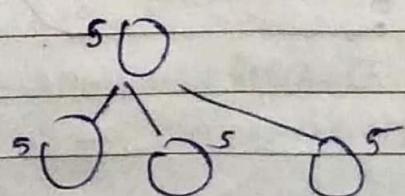
A heuristic funct. is said to be optimal if it gives the exact value for an answer.

Simple hill climbing

choose the best way available and discard all other nodes.



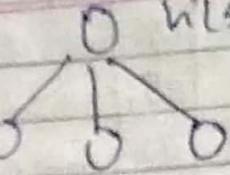
Steepest ascent hill climbing



cannot go any further.

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Variance of
Hill climbing



Hence, the no.
of nodes in the
memory at any
point = $O(b)$

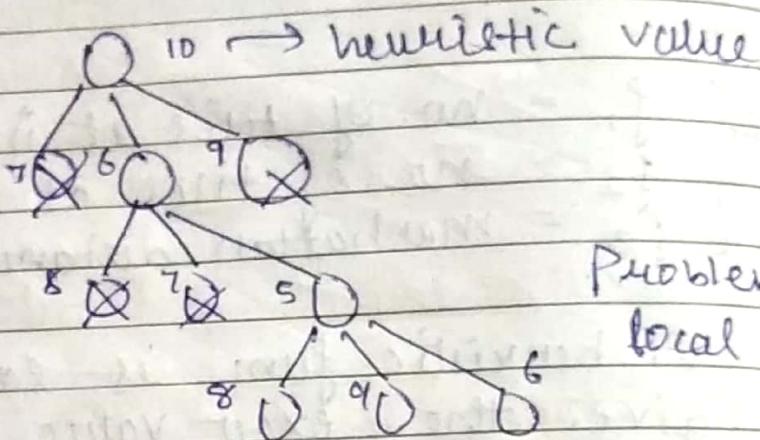
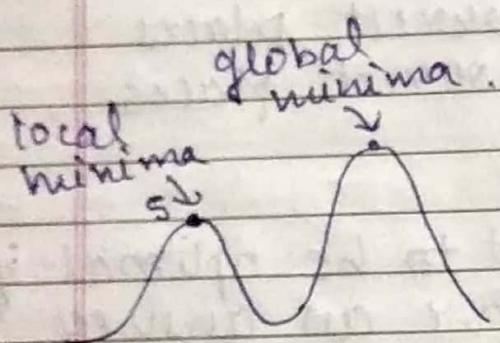
Date _____
Time _____
Phase assign heuristic
value

II (2) generate successor

II (3) Search for the
best option at
every level

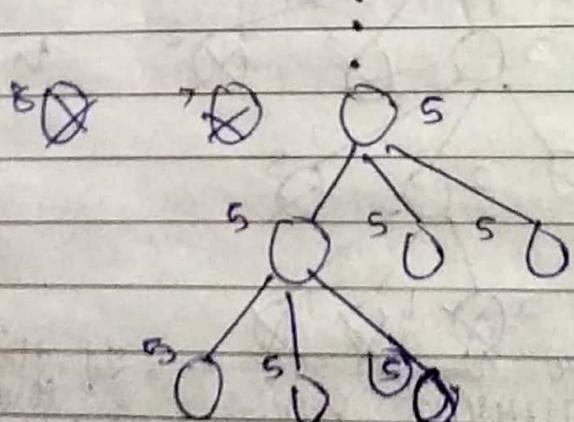
II (4) Discard other nodes

Ex: (i)

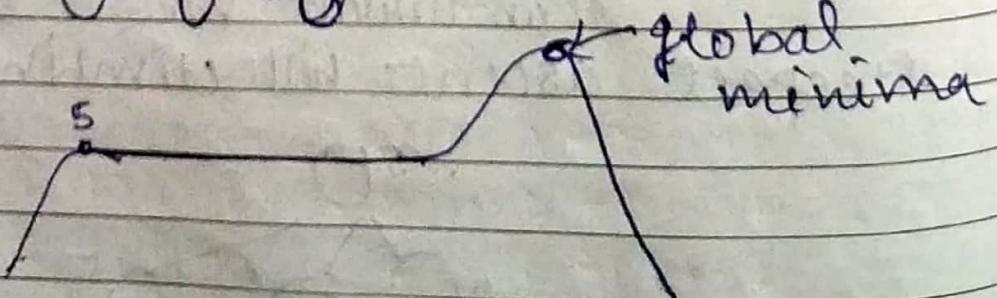


Since, no node has a
heuristic value better than
(5), the algorithm / procedure
terminates.

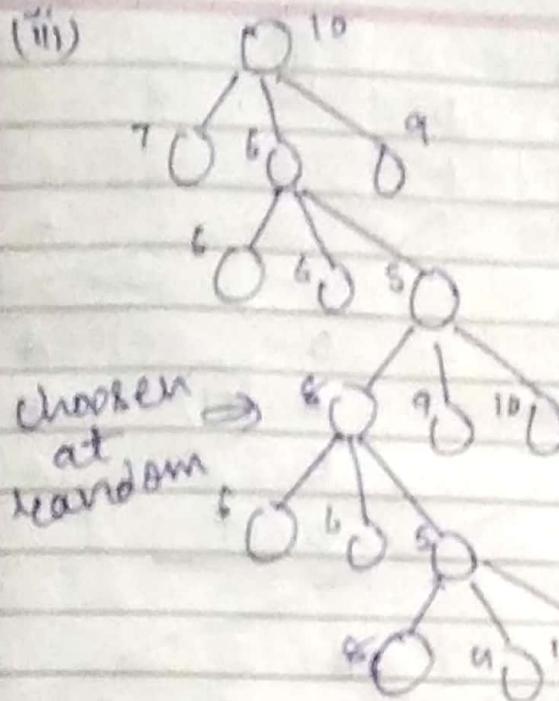
(ii)



we cannot decide
where to go
next. This
is called
plateau (flat
area).



(iii)



When we cannot find a better solution, a node is chosen at random and others discarded. This problem is called ridge global minima.

Solutions:

i) Backtracking - A mechanism to store more nodes rather than discarding them. This increases the space complexity but backtracking can be executed.

ii) Random jump - Apply a random func to explore any node in the memory

Block World Problem

+1 A
+1 D (5)
+1 B
+1 C +1

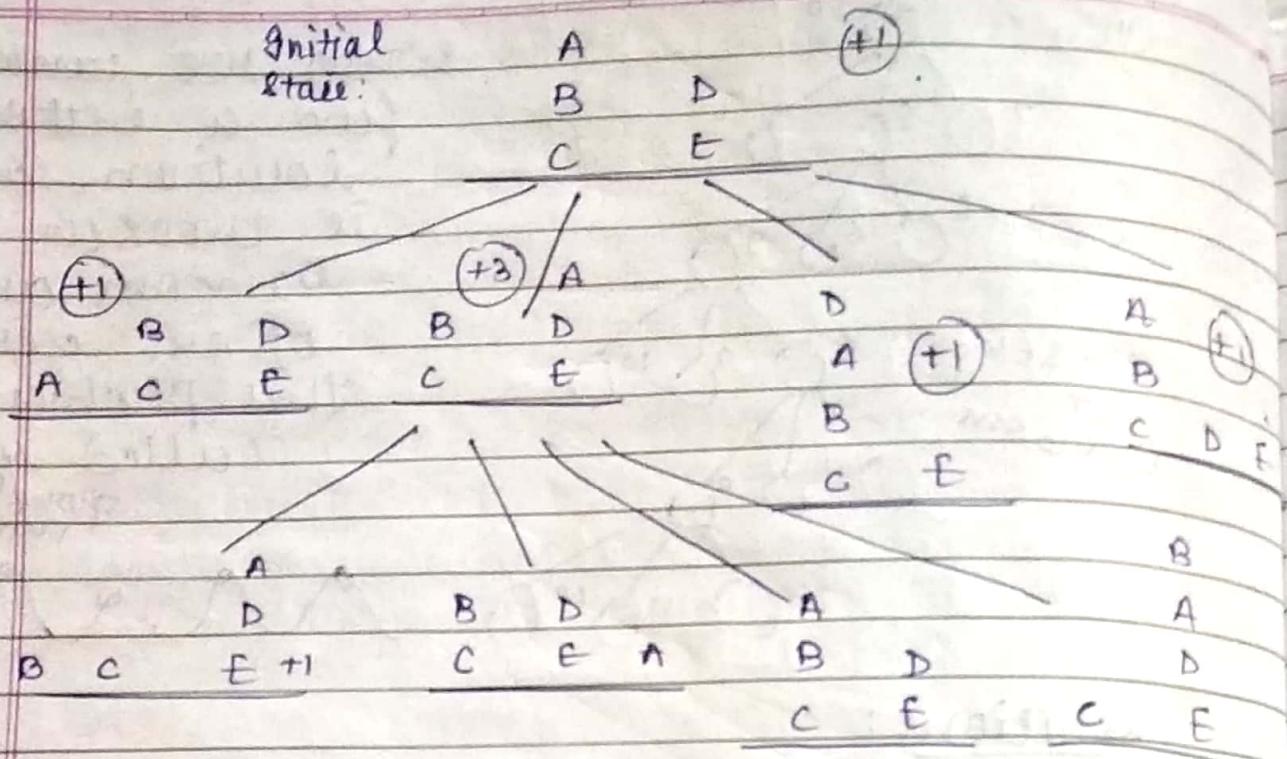
goal state

-1 A (+1)
+1 B D -1
+1 C E +1

initial state

(because we didn't discard them)

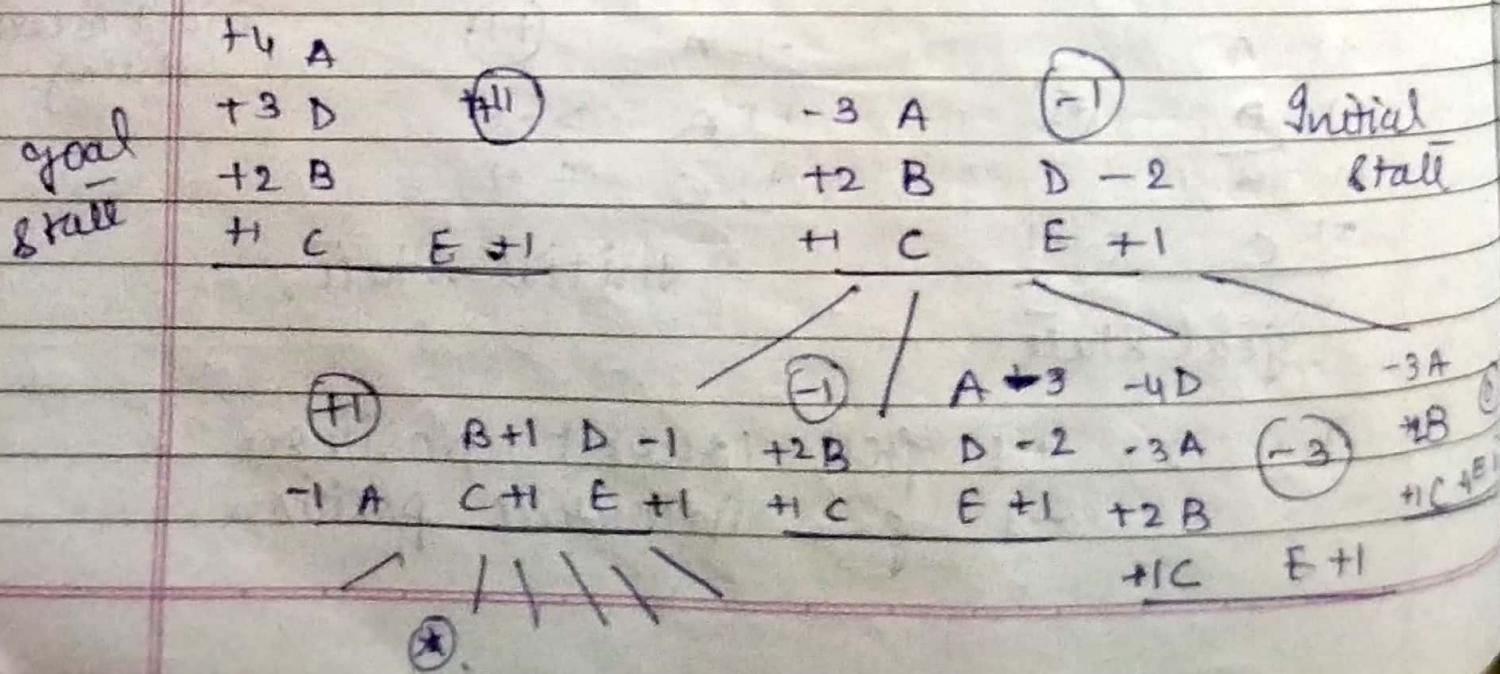
h_1 : +1 for correct position
-1 for incorrect position



Since, we don't get a better solution than +3, the procedure terminates. Hence, the heuristic func. must be incorrect or the algorithm is incorrect.

h_2 : Add +1 for every block in correct structure.

-1 for every block for incorrect structure.



$$\begin{array}{ccccccccc}
 & \textcircled{*} & \textcircled{-1} & \textcircled{-3} & \textcircled{-2} & \textcircled{+1} & \textcircled{+6} \\
 \textcircled{F2} & A & D & -2B & D-2 & -2D & D & B \\
 B & C & E & -A + C & E + 1 & -B - A + C & E + 1 & A C E \\
 \hline
 A C E D & C & E & \underline{-A + C} & \underline{E + 1} & \underline{-B - A + C} & \underline{E + 1} & \underline{A C E} \\
 \end{array}$$

$$\begin{array}{cccccc}
 +3 & D & \textcircled{f5} & A & -2D & +2B \\
 +2B & A-2 & D & \textcircled{f11} & -A + C & E + 1 \\
 +1C & E + 1 & B & . & +2B & -2D \\
 \hline
 C & E & \underline{-A + C + 1E} & & &
 \end{array}$$

goal

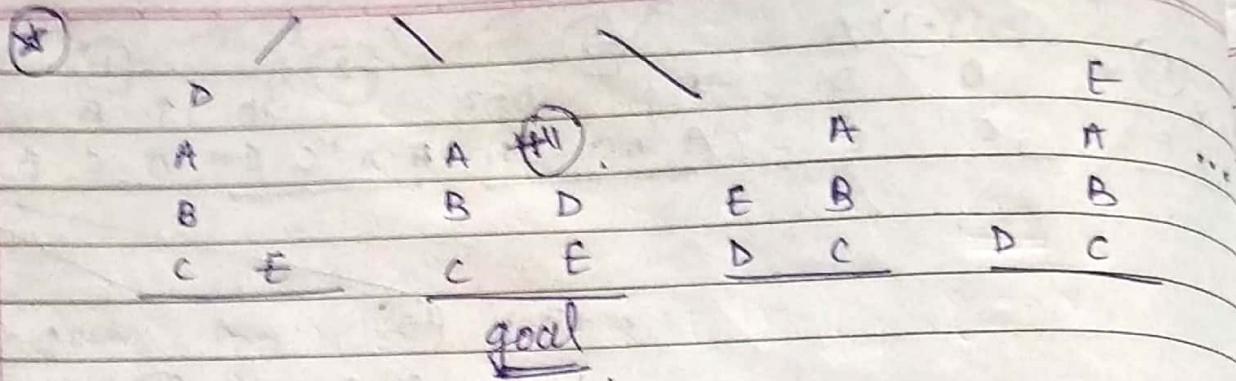
$$\begin{array}{cccccc}
 \Rightarrow & -4A & & +3A & \textcircled{+9} \\
 -3D & \textcircled{f3} & & +2B & D + 2 \\
 +2B & & & +1C & E + 1 \\
 +1C & E + 1 & & +1C & E + 1
 \end{array}$$

goal

$$\begin{array}{cccccc}
 -3D & \textcircled{D} & -3D & \textcircled{-1} & A & -4 \\
 +2B & & +2B & & D & -3 \\
 -A + C & E + 1 & +1C & E + 1 & B & +2 \\
 \hline
 C + 1
 \end{array}$$

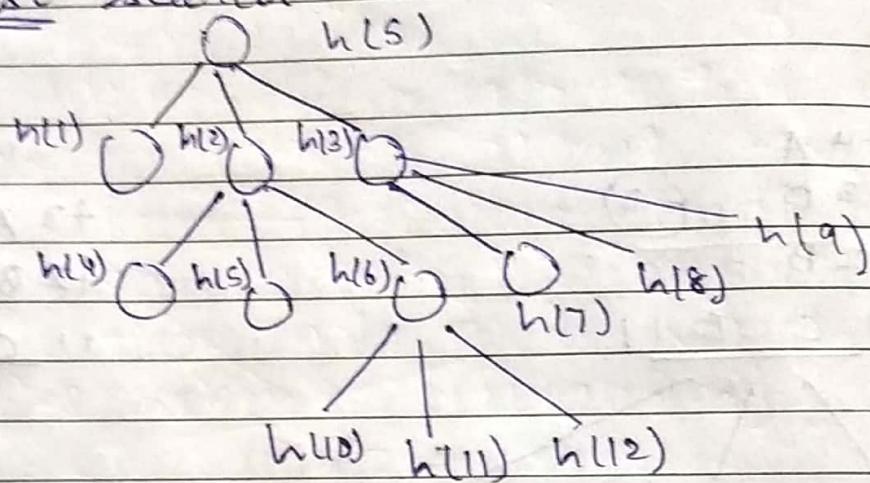
$$\begin{array}{cccccc}
 +2B & A-1 & -4A & \textcircled{-1} & D-3 & -3D \\
 -1D + C & E + 1 & -3D & \textcircled{-3} & +2B & \textcircled{D} \\
 \hline
 +2B & +C & E + 1 & +1C & E + 1 & A-1
 \end{array}$$

$$\begin{array}{cccccc}
 +2B & A-1 & -4A & \textcircled{-1} & D-3 & -3D \\
 -1D + C & E + 1 & -3D & \textcircled{-3} & +2B & \textcircled{D} \\
 \hline
 +2B & +C & E + 1 & +1C & E + 1 & A-1 \\
 +1C & E + 1 & -1D + C & E + 1 & -1D + C & E + 1 \\
 \hline
 -1A & -1D & +2B & \textcircled{+2} & +1C & E + 1 \\
 \hline
 \end{array}$$



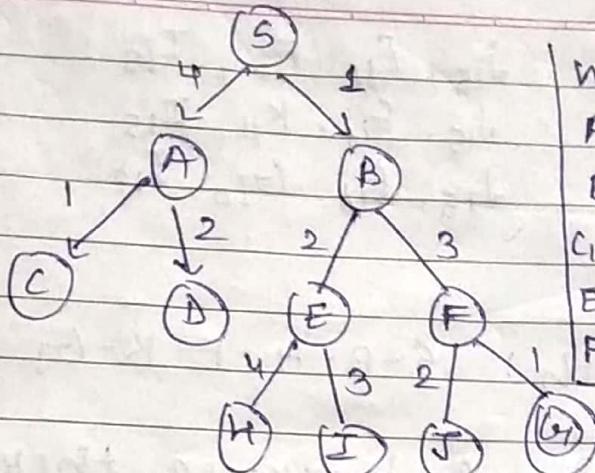
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Best first search

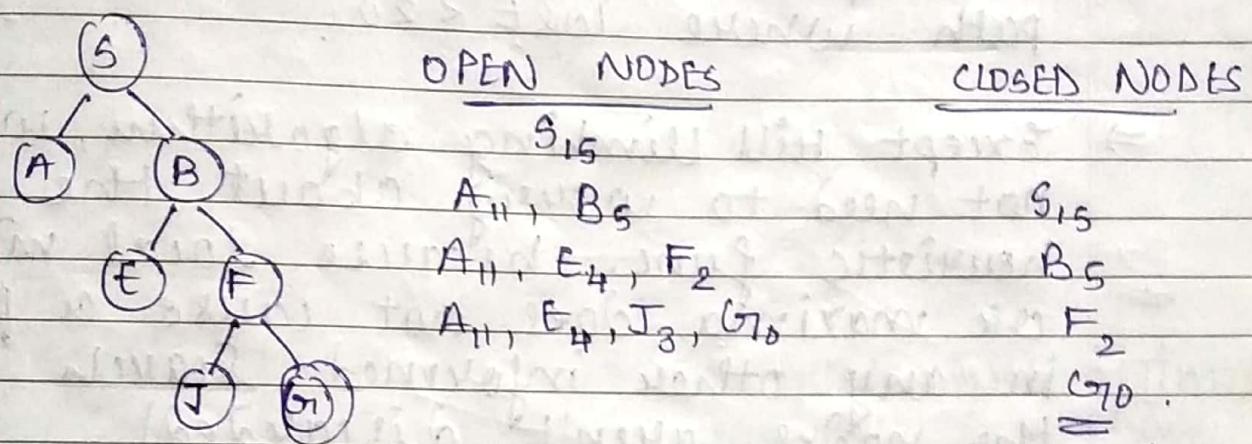


- Select the best node so far based on its heuristic value considering the entire tree.
- Here, we start with $h(5)$ generate successor $(h(1), h(2), h(3))$. Let's say $h(2)$ is the best PDL^h generate successor of $h(2)$ $(h(4), h(5))$. Now, nodes in mem: $h(1), h(3), h(4), h(5)$. Let's say the best PDL^h among them is $h(3)$. Generate successor for $h(3)$ $(h(7), h(8), h(9))$. Now, nodes in mem: $h(1), h(4), h(5), h(6), h(7), h(8), h(9)$ and so on.

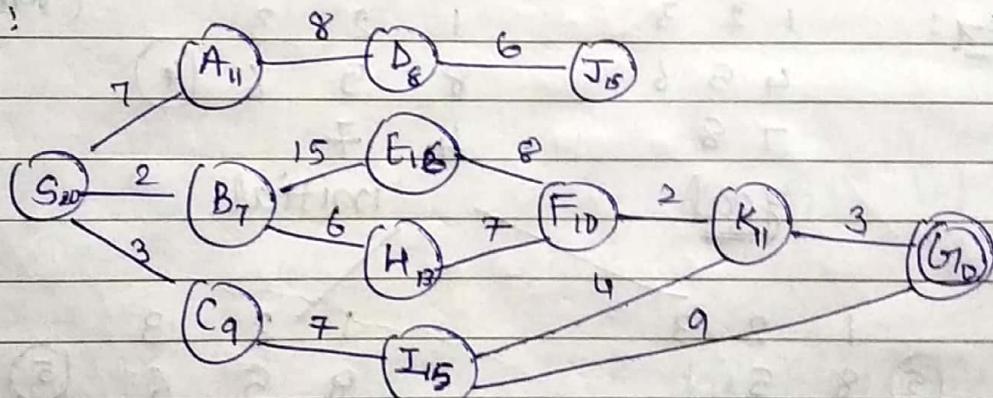
Ex:



node	$h(n)$	node	$h(n)$
A	11	H	7
B	5	I, J	3
C, D	9	S	15
E	4	G	0
F	2		



Ex:



OPEN NODE

CLOSED NODES

S_{20}

A_{11}, B_7, C_9

$A_{11}, E_{16}, H_{13}, C_9$

$A_{11}, E_{16}, H_{13}, I_{15}$

$D_8, E_{16}, H_{13}, I_{15}$

$I_{15}, E_{16}, H_{13}, I_{15}$

S_{20}

B_7

C_9

A_{11}

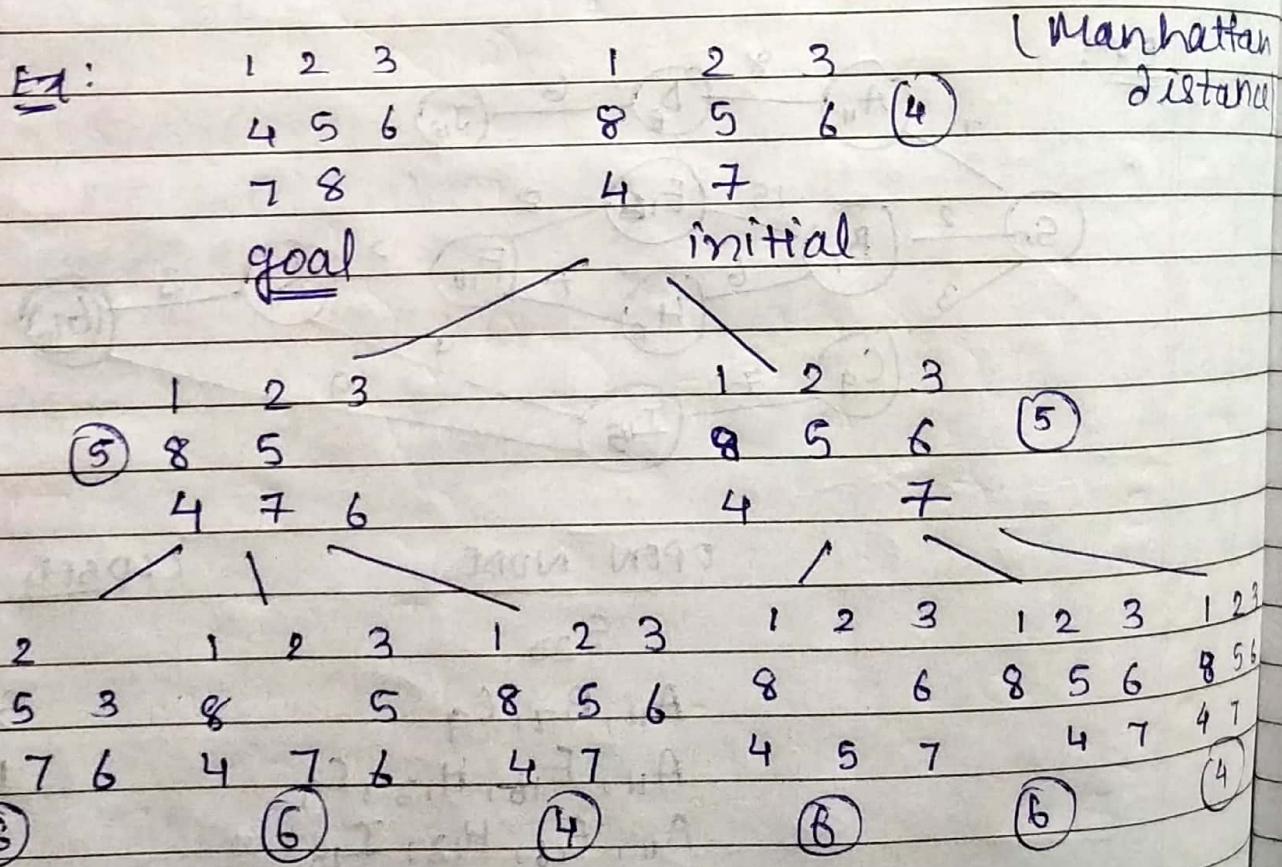
D_8

$J_{15}, E_{16}, F_{10}, I_{15}$ H_{13}
 $J_{15}, E_{16}, K_{11}, I_{15}$ F_{10}
 $J_{15}, E_{16}, G_{70}, I_{15}$ K_{11}
 G_{70}

Solⁿ path: S-B-H-F-K-G₇ (20)

Not optimal because there exist path where cost < 20.

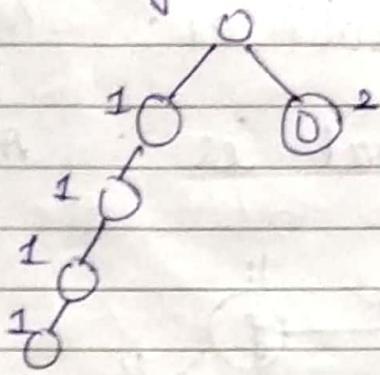
⇒ Except Hill Climbing algorithm, we do not need to worry about the heuristic func. because local minima or maxima does not cause a problem in any other informed search (as the nodes aren't discarded).



Time complexity for Best First Search:
 $O(b^d)$

Space complexity = $O(b^d)$

Depth can be infinite but, branching factor is fixed.



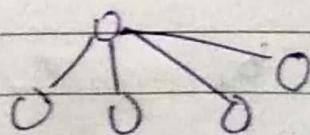
thus, it is not complete nor optimal.
 [depends on the heuristic function.
 Better the heuristic func., better
 the algorithm]

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A* algorithm

(*-algo. can be made optimal by applying
 a no. of constraints)

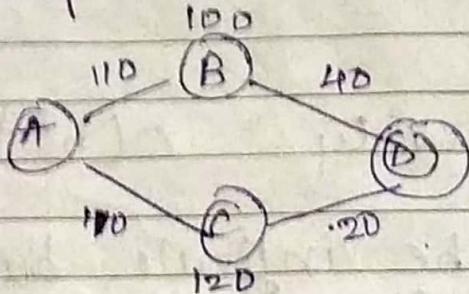
Instead of considering the heuristic value
 only (for every node), we take into
 consideration the actual cost too.



$$f(n) = h(n) + g(n)$$

heuristic value actual value

considering heuristic value only:



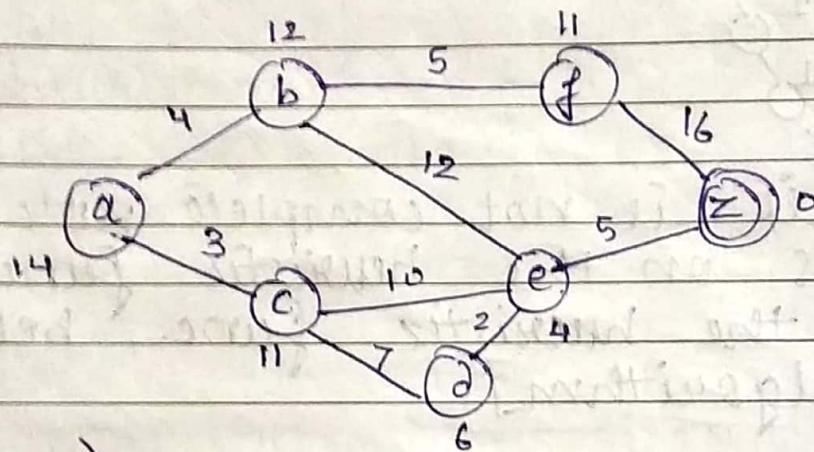
We obtain SLDⁿ as: $A - B - D$

$$\begin{array}{r} 11D \quad 4D \\ \hline = 150 \end{array}$$

but, optimal is : $A - C - D$

$$\begin{array}{r} 11D \quad 2D \\ \hline = 130 \end{array}$$

Ex:



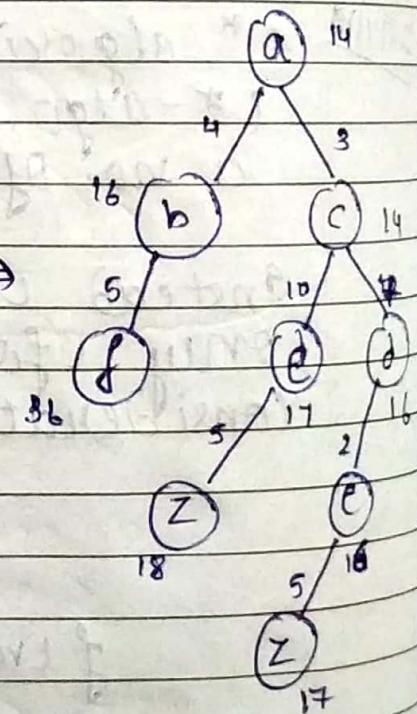
not removed
from mem.

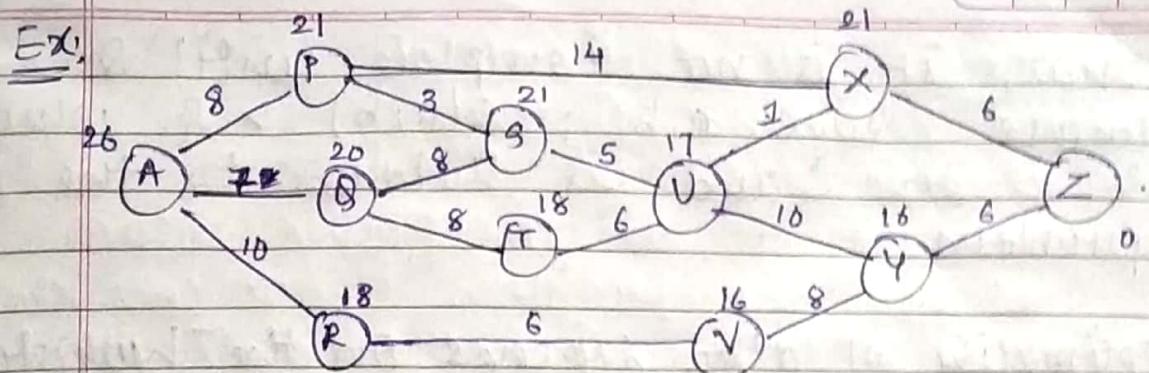
$$f(a) = 14$$

$$\begin{aligned} f(b) &= g(b) + h(b) \\ &= 4 + 12 \\ &= 16 \end{aligned}$$

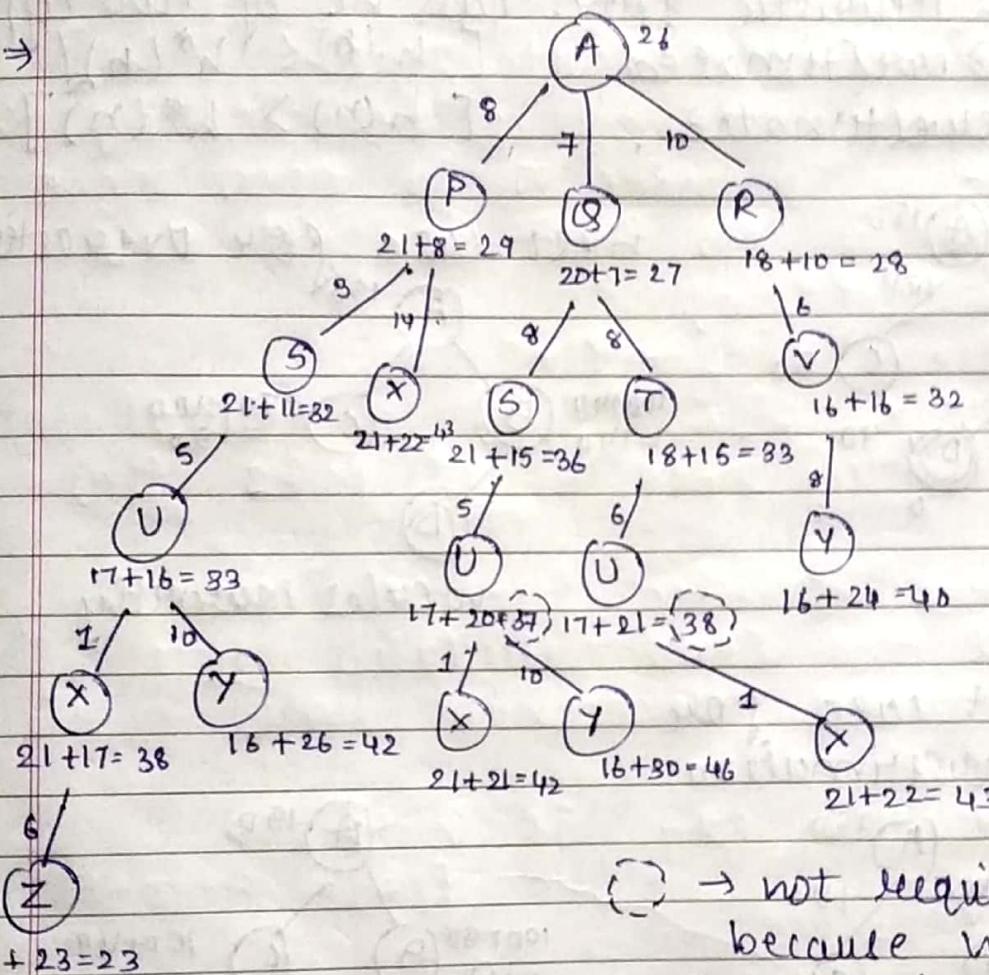
$$\begin{aligned} f(c) &= g(c) + h(c) \\ &= 3 + 11 \\ &= 14 \end{aligned}$$

(same as b,
so follow
alphabetical
order).



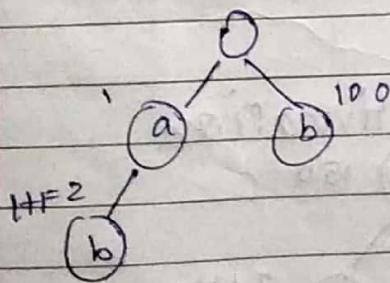


→



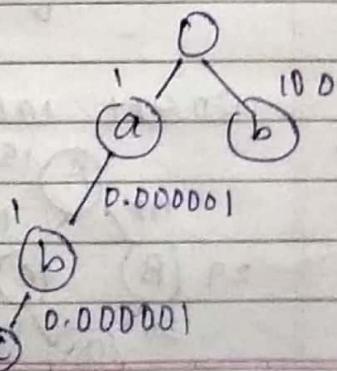
○ → not required to explore
because we already
have a lower cost path
till U than 37 or 38

25/11/19



$$1 + 100 = 1$$

then, we get a solution

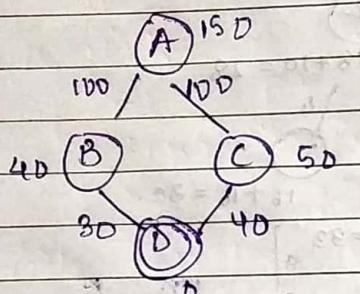


Here, it goes into an
infinite loop

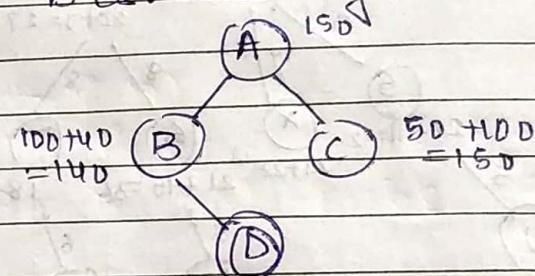
Thus, it is not complete until & unless edge a, b ; $e(a, b) > 0$ when, & it is +ve and is related to the problem.

Optimality of algo depends on the heuristic func. Heuristic func. can be of two types:

- Underestimated [$h(n) \leq h^*(n)$]
- Overestimated [$h(n) > h^*(n)$]



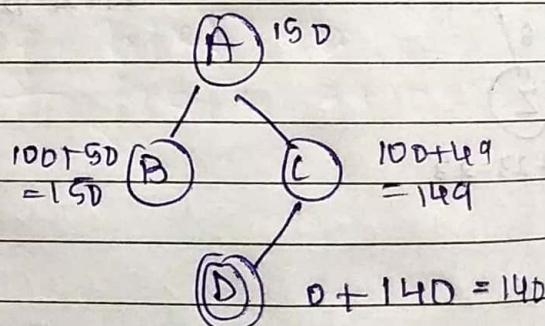
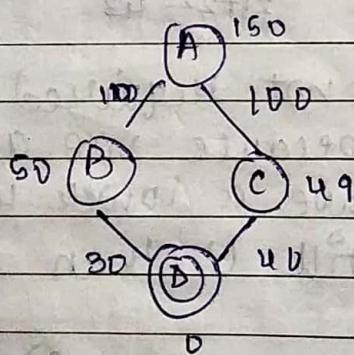
Best case for overestimation



Ex:

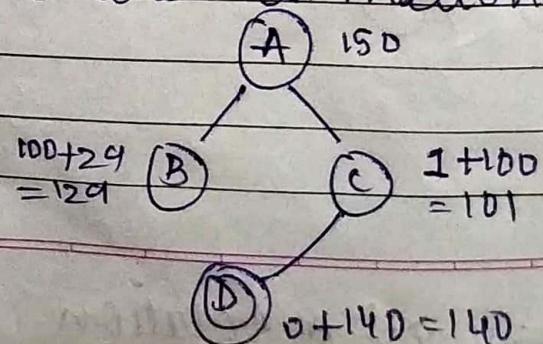
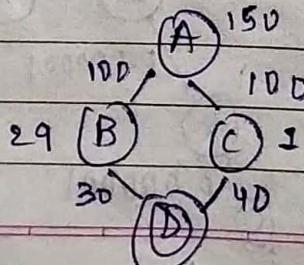
would correctly

Worst case for overestimation:



Ex:

Worst case for underestimation



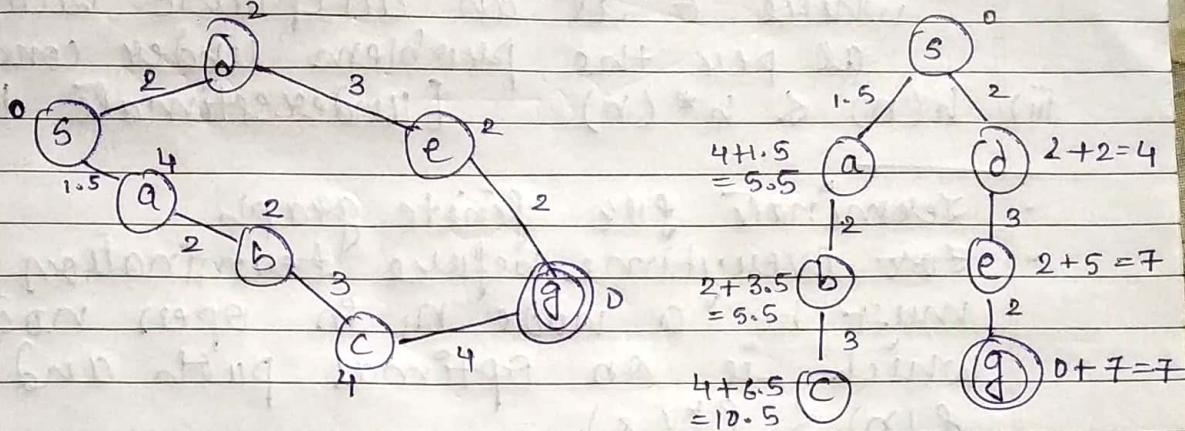
Ex:

$$T+14 = 21$$

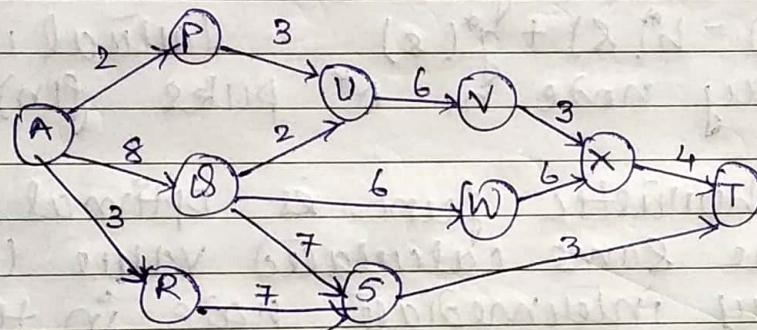
$$0+18 = 18$$

Since, $140 > 129$, the algo. goes on to explore B and thus, gets the optimal path. Hence, we should go for underestimation.

Ex:



Ex

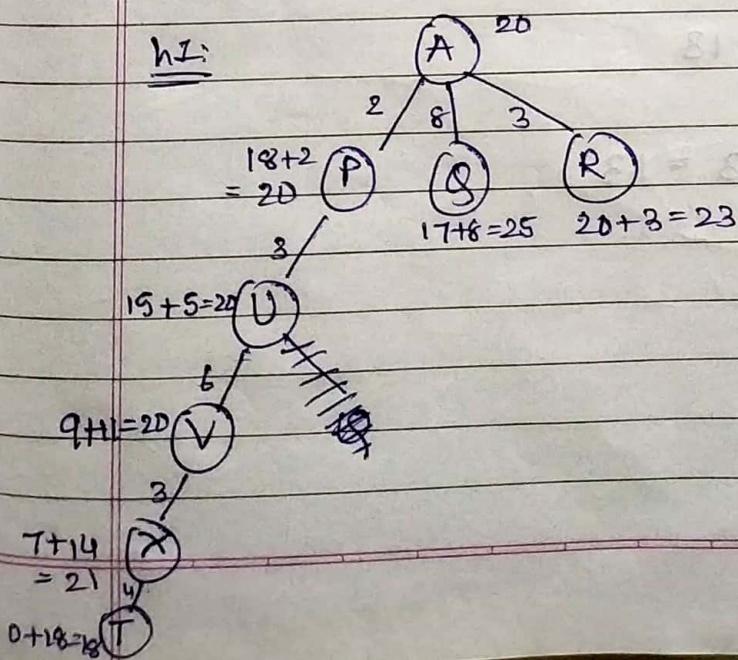


h_1, h_2

A	20	20
P	18	8
Q	17	14
R	20	12
S	10	3
V	15	12
W	20	2
X	7	3
T	0	0

[$h_1 \rightarrow$ Overestimated heuristic func]
[$h_2 \rightarrow$ Underestimated H.F]

h_1 :



conditions for optimality:

- i) finite bounding factor.
- ii) for each edge, cost $c(a,b) > \epsilon$; (ϵ depth)
where ϵ is an acceptable value
as per the problem under consideration
- iii) $h(n) \leq h^*(n)$ [Underestimate HF]

- Terminate for finite graph
- For everytime before termination, there must be a node n_i in open node which is on optimal path and
 $f(n) \leq f^*(s)$
 $f(n) = h(n) + g(n)$
 $f^*(s) = h^*(s) + g^*(s)$ (actual cost)
- For every node n , A* picks $f(n) \leq f^*(s)$

If the heuristic func. is optimal, it gives the same calculated value ($f(n)$) for every intermediate node in the optimal path.

In previous ex: if $h(R) = 10$

$$f(R) = 8 + 10 = 18$$

$$h(S) = 3$$

$$f(S) = 3 + 10 = 13$$

$$h(T) = 0$$

$$f(T) = 3 + 7 + 3 = 13$$