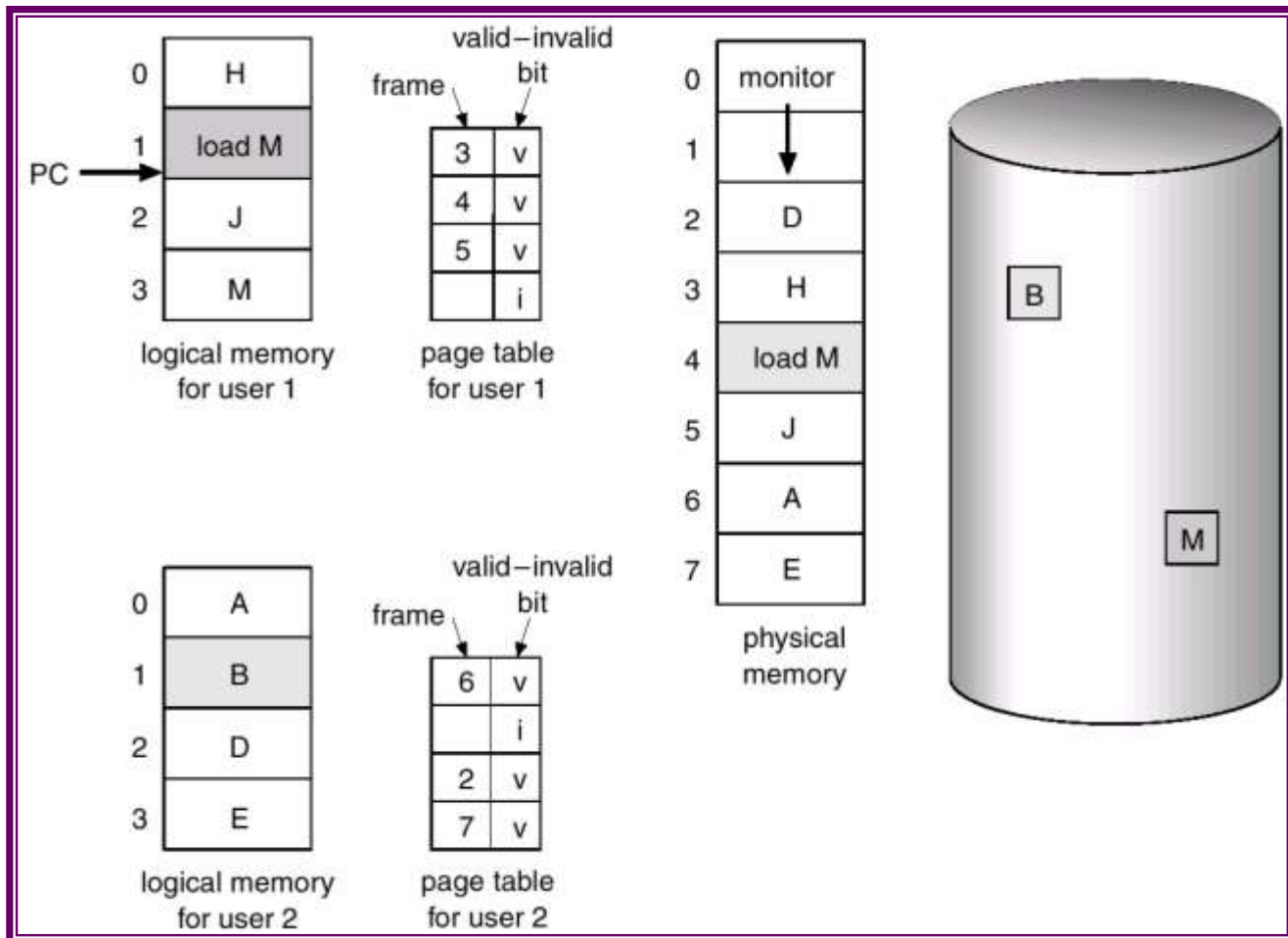




# SO 1: Identify the need for page replacement



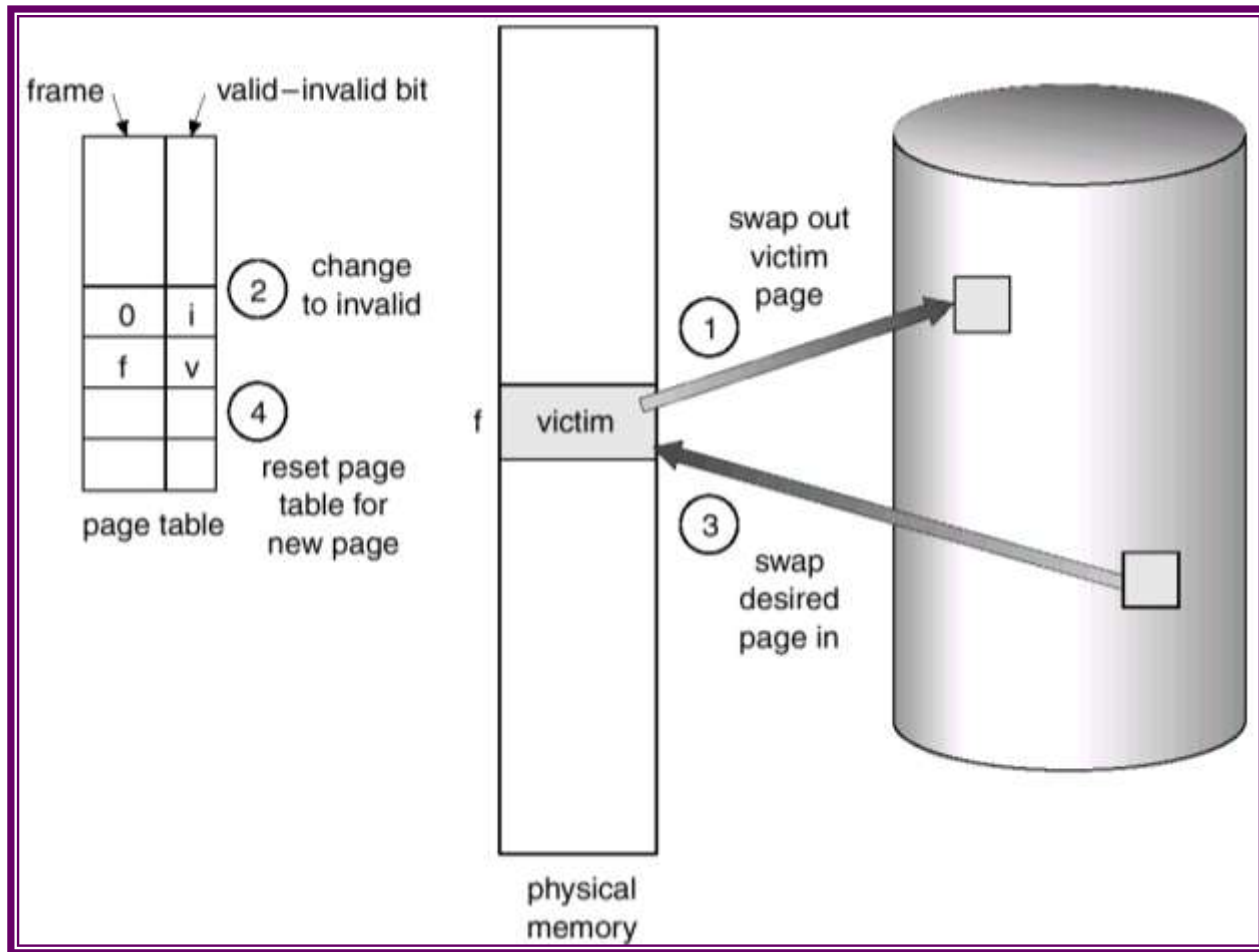
# Need For Page Replacement

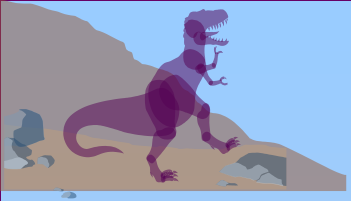


# Basic Page Replacement

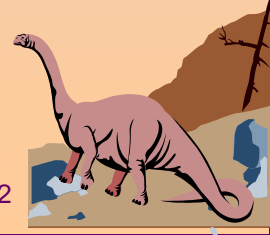
1. Find the location of the desired page on disk.
2. Find a free frame:
  - If there is a free frame, use it.
  - If there is no free frame, use a page replacement algorithm to select a *victim* frame.
3. Read the desired page into the (newly) free frame.  
Update the page and frame tables.
4. Restart the process.

# Page Replacement





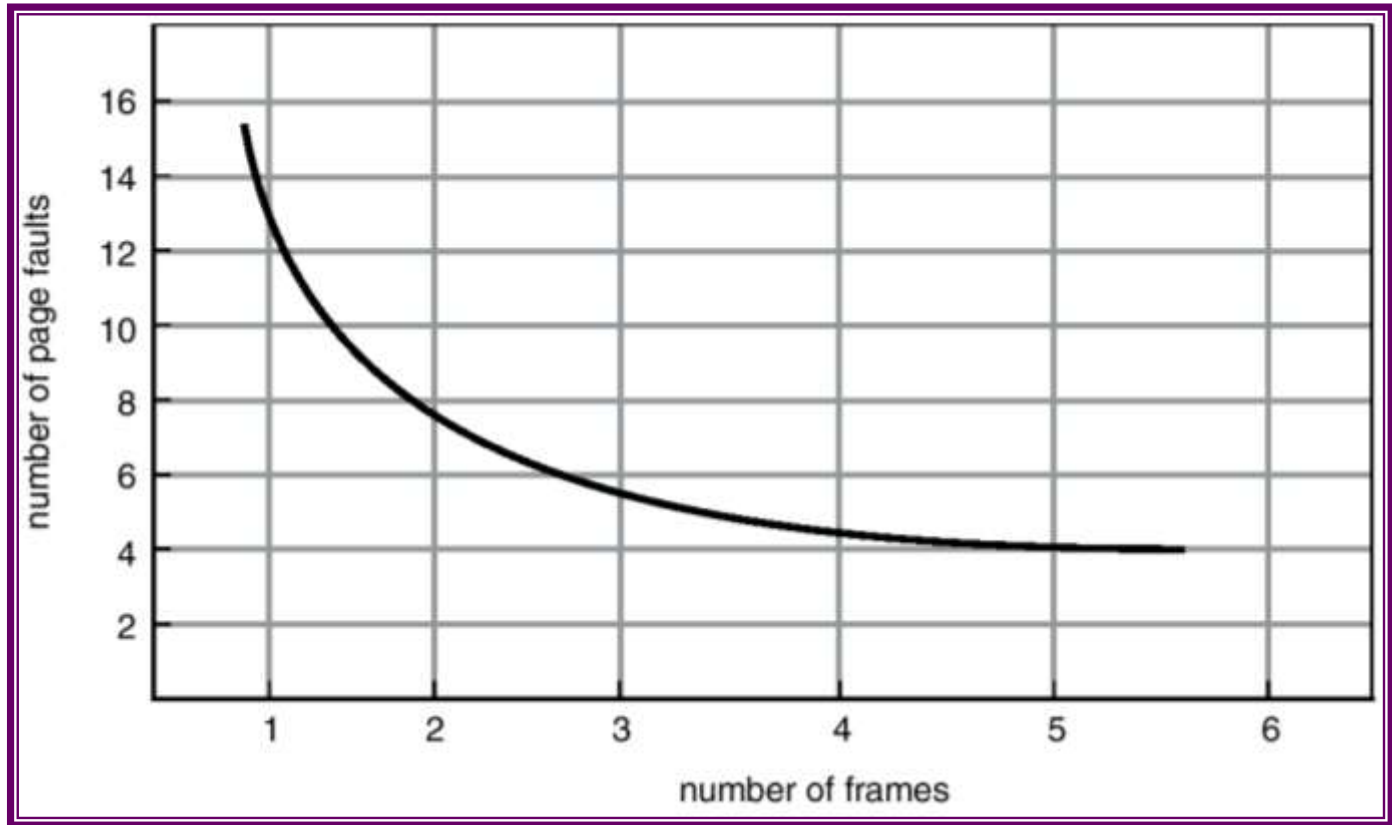
## SO 2: Demonstrate the page replacement algorithms



# Page Replacement Algorithms

- Want lowest page-fault rate.
- Evaluate algorithm by running it on a particular string of memory references (reference string) and computing the number of page faults on that string.
- In all our examples, the reference string is  
1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5.

# Graph of Page Faults Versus The Number of Frames



# First-In-First-Out (FIFO) Algorithm

Reference string: 1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5

3 frames (3 pages can be in memory at a time per process)

|   |   |   |   |               |
|---|---|---|---|---------------|
| 1 | 1 | 4 | 5 |               |
| 2 | 2 | 1 | 3 | 9 page faults |
| 3 | 3 | 2 | 4 |               |

4 frames

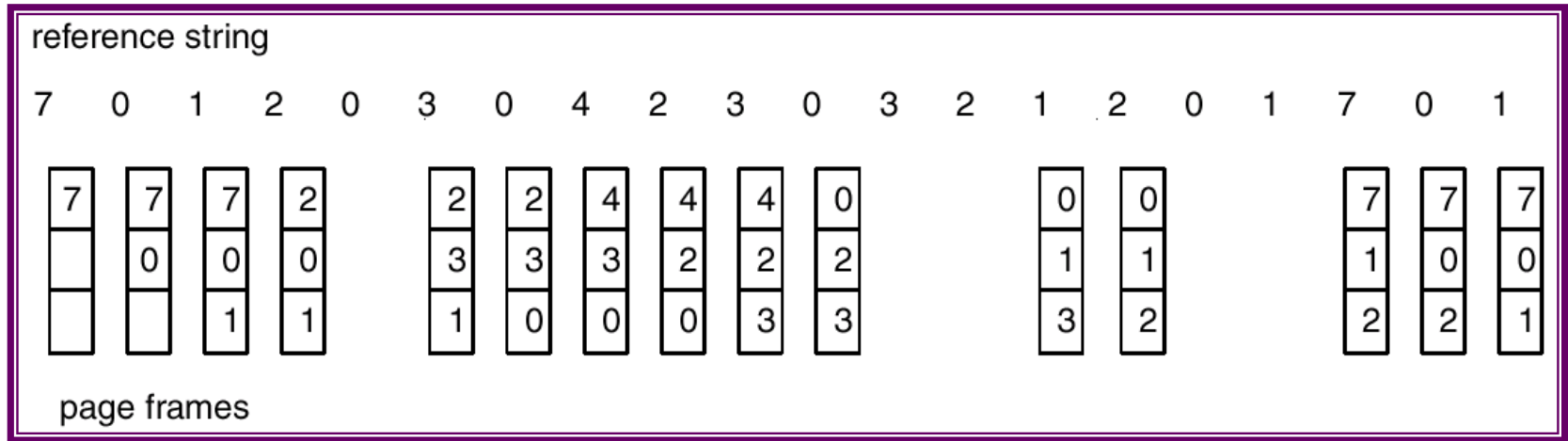
|   |   |   |   |                |
|---|---|---|---|----------------|
| 1 | 1 | 5 | 4 |                |
| 2 | 2 | 1 | 5 | 10 page faults |
| 3 | 3 | 2 |   |                |
| 4 | 4 | 3 |   |                |

FIFO Replacement – Belady's Anomaly

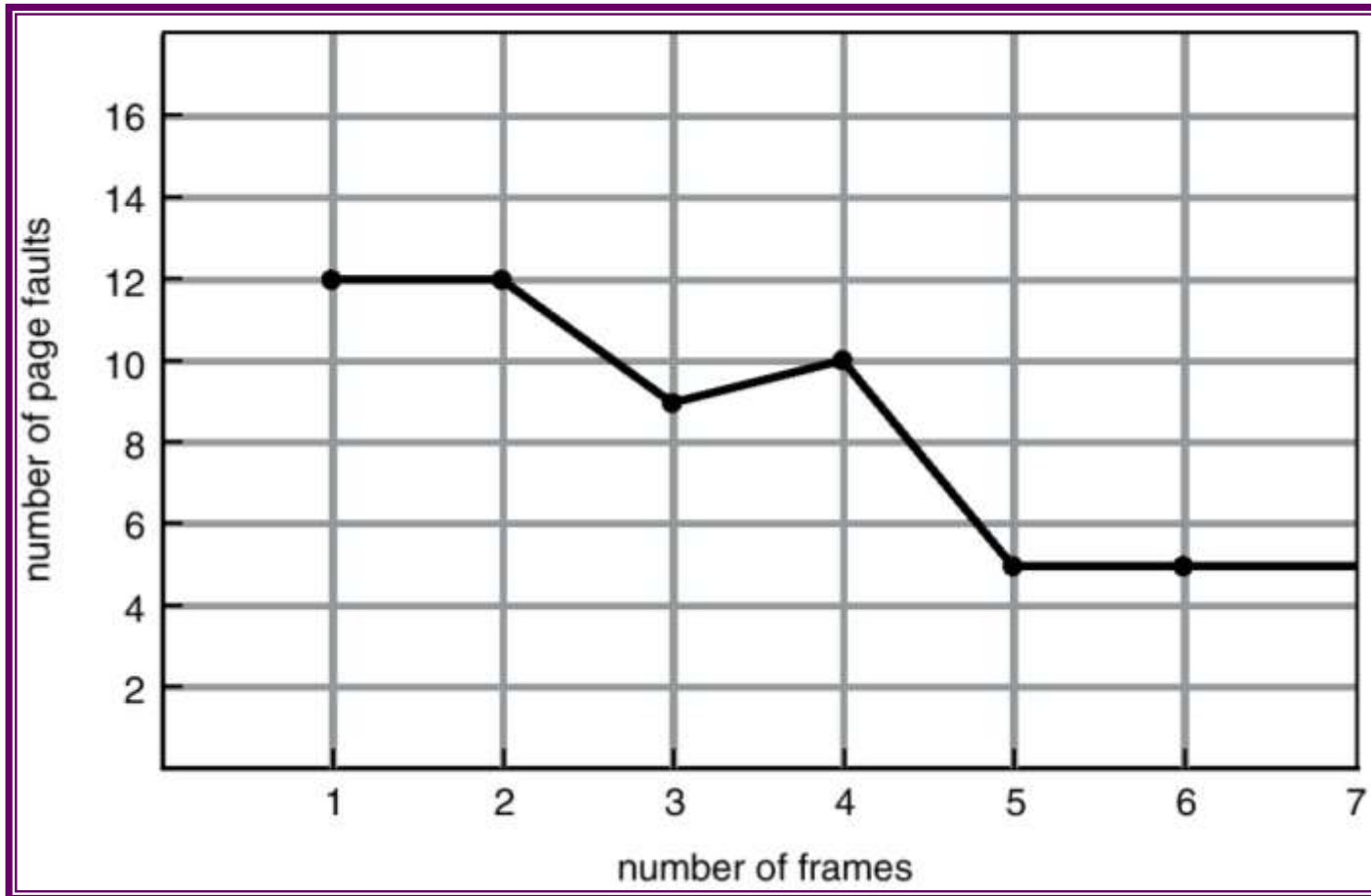
more frames  $\Rightarrow$  less page faults



# FIFO Page Replacement



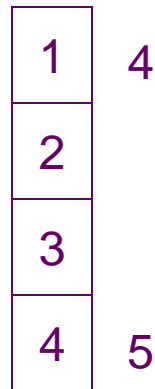
# FIFO Illustrating Belady's Anamoly



# Optimal Algorithm

- ❑ Replace page that will not be used for longest period of time.
- ❑ 4 frames example

1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5



6 page faults

- ❑ How do you know this?
- ❑ Used for measuring how well your algorithm performs.

# Optimal Page Replacement

reference string

7 0 1 2 0 3 0 4 2 3 0 3 2 1 2 0 1 7 0 1

|   |   |   |   |  |   |  |   |  |   |  |   |  |  |  |  |  |   |  |  |
|---|---|---|---|--|---|--|---|--|---|--|---|--|--|--|--|--|---|--|--|
| 7 | 7 | 7 | 2 |  | 2 |  | 2 |  | 2 |  | 2 |  |  |  |  |  | 7 |  |  |
|   | 0 | 0 | 0 |  | 0 |  | 4 |  | 0 |  | 0 |  |  |  |  |  | 0 |  |  |
|   |   | 1 | 1 |  | 3 |  | 3 |  | 3 |  | 1 |  |  |  |  |  | 1 |  |  |

page frames

# Least Recently Used (LRU) Algorithm

- Reference string: 1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5

|   |   |   |
|---|---|---|
| 1 | 5 |   |
| 2 |   |   |
| 3 | 5 | 4 |
| 4 | 3 |   |

- Counter implementation
  - Every page entry has a counter; every time page is referenced through this entry, copy the clock into the counter.
  - When a page needs to be changed, look at the counters to determine which are to change.

# LRU Page Replacement

reference string

7 0 1 2 0 3 0 4 2 3 0 3 2 1 2 0 1 7 0 1

|   |   |   |   |  |   |  |   |   |   |   |  |  |   |  |   |  |   |  |  |
|---|---|---|---|--|---|--|---|---|---|---|--|--|---|--|---|--|---|--|--|
| 7 | 7 | 7 | 2 |  | 2 |  | 4 | 4 | 4 | 0 |  |  | 1 |  | 1 |  | 1 |  |  |
|   | 0 | 0 | 0 |  | 0 |  | 0 | 0 | 3 | 3 |  |  | 3 |  | 0 |  | 0 |  |  |
|   |   | 1 | 1 |  | 3 |  | 3 | 2 | 2 | 2 |  |  | 2 |  | 2 |  | 7 |  |  |

page frames



□ Reference strings:

1, 2, 3, 4, 1, 2, 5, 3, 2, 1, 3, 4, 5, 2, 1

Frame size : 4

2, 3, 4, 8, 7, 6, 8, 4, 4, 5, 6, 7, 8, 4, 4, 3, 3, 7, 8, 9

Frame size : 3

