



SO 1: List the characteristics of segmentation





Segmentation

- ❑ Memory-management scheme that supports user view of memory.
- ❑ A program is a collection of segments. A segment is a logical unit such as:

main program,
procedure,
function,
method,
object,
local variables, global variables,
common block,
stack,
symbol table, arrays



Programs are a collection of logical modules

```
static unsigned int blk_flush_policy(unsigned int fflags, struct request *rq)
{
    unsigned int policy = 0;

    if (blk_rq_sectors(rq))
        policy |= REQ_FLUSH_DATA;

    return policy;
}

static unsigned int blk_flush_int_req(struct request *rq)
{
    return 1 & ffs(rq->flush_seq);
}

static void blk_flush_restore_request(struct request *rq)
{
    rq->bio = rq->biotail;
    rq->end_io = rq->flush->end_io;
}

static bool blk_flush_queue_rq(struct request *rq, bool add_front)
{
    if (rq->rq->qq) {
        struct request_queue *qq = rq->rq;

        blk_rq_add_to_request_list(rq, add_front);
        blk_rq_vic_request_list(rq);
        return false;
    }
}

static bool blk_flush_complete_req(struct request *rq,
                                   struct blk_flush_queue *fq,
                                   unsigned int seq, int error)
{
    struct request_queue *qq = rq->rq;
    struct list_head *pending = &fq->flush_queue[0]->flush_pending_list;
    bool queued = false, kicked;

    BUG_ON(rq->flush_seq < seq);
    rq->flush_seq |= seq;

    static void flush_end_io(struct request *flush_rq, int error)
    {
        struct request_queue *qq = flush_rq->rq;
        struct list_head *pending;
        bool kicked = false;
    }
}
```

Logical modules : such as global data, stack, heap, functions, classes, namespaces, etc.

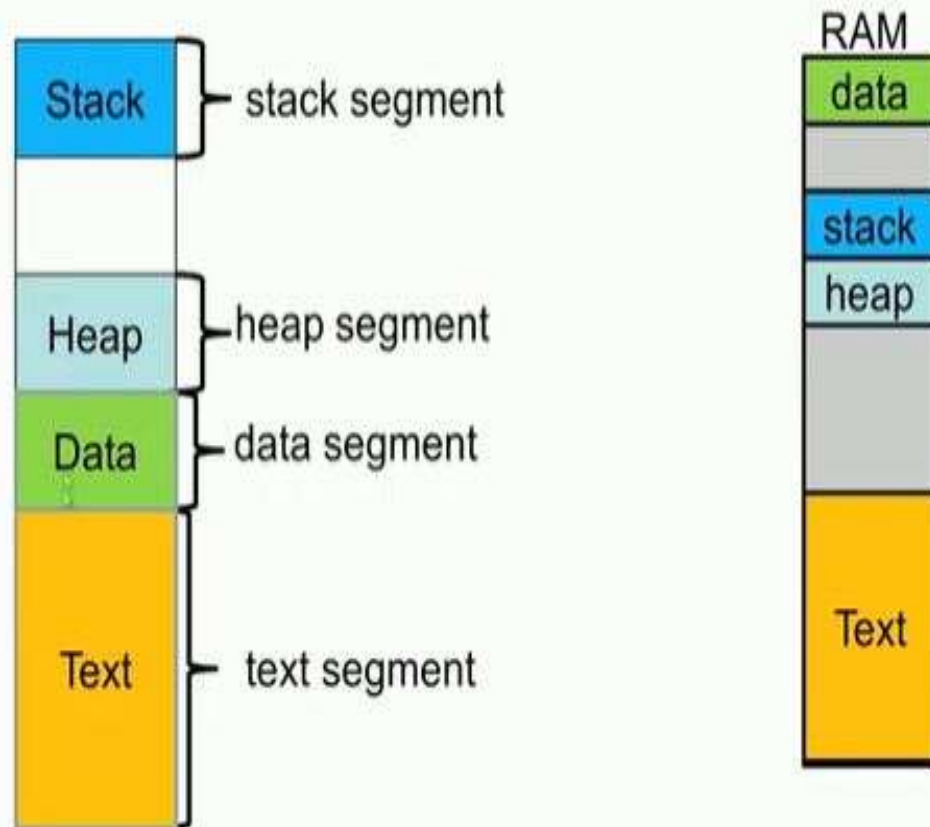
Virtual memory does not split programs into logical modules, instead splits programs into fixed size blocks.

Segmentation can be used to split program into segments that are more logical.

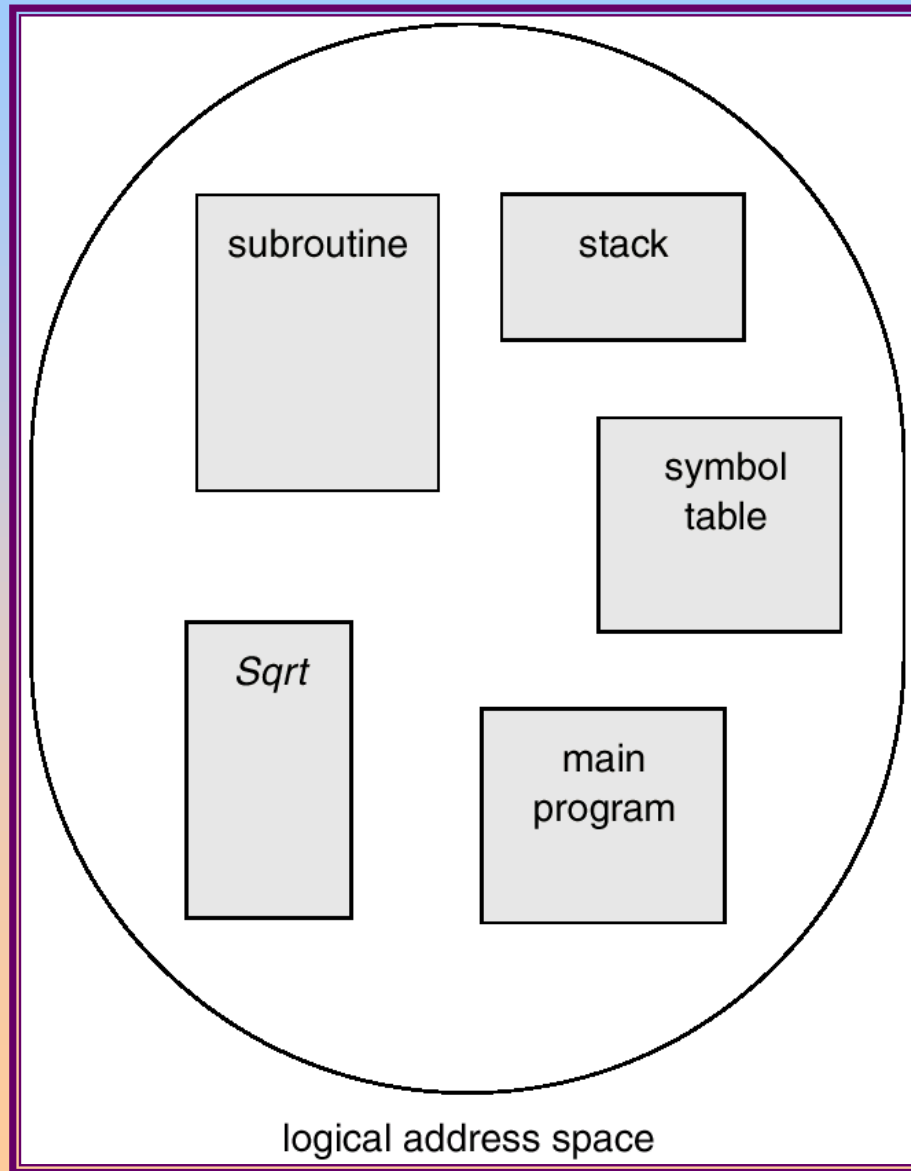
The segment size could range from a few bytes to the maximum size (4GB in 32 bit Intel machines)



Commonly Used Segments (an example)

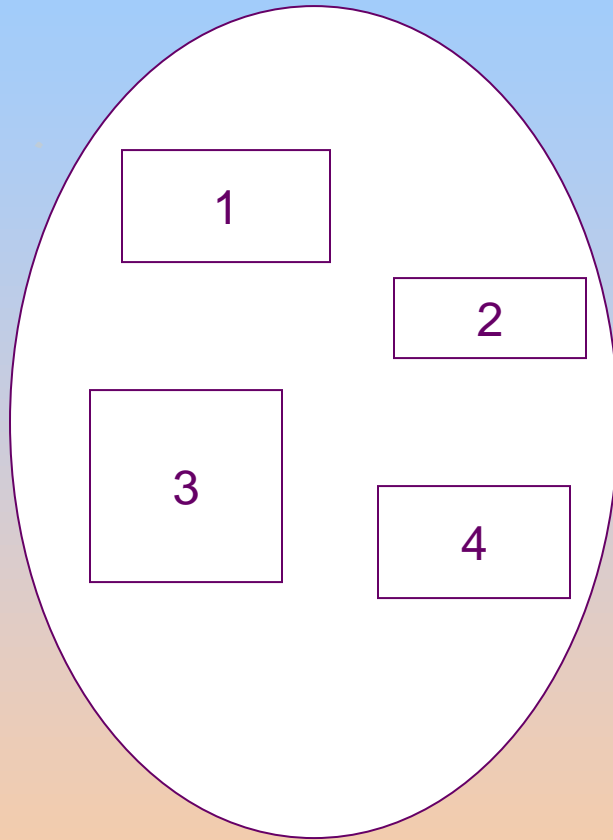


User's View of a Program

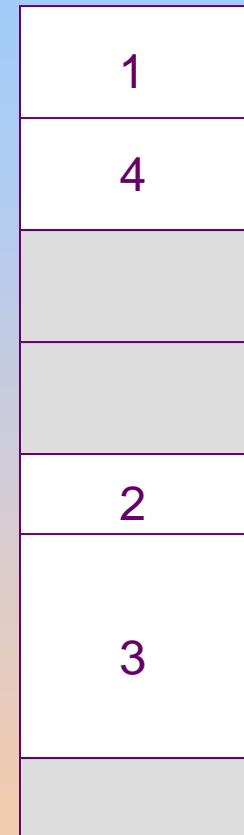




Logical View of Segmentation



user space

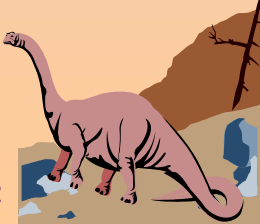


physical memory space





SO 2: Illustrate the segmentation architecture



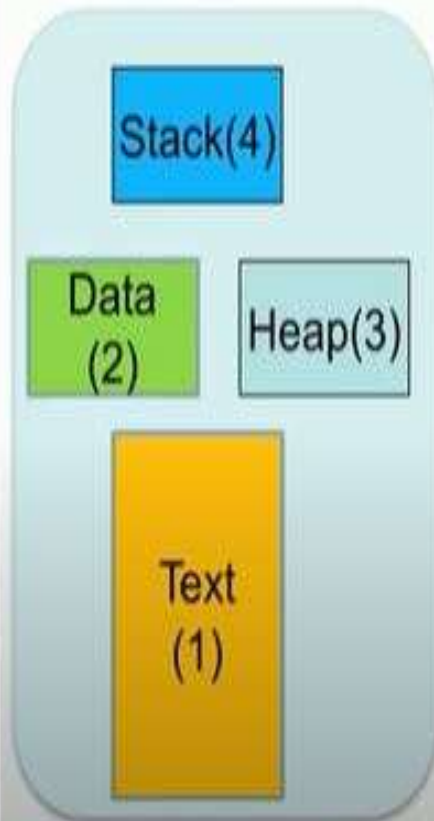
Segmentation Architecture

- Logical address consists of a two tuple:
 - $\langle \text{segment-number}, \text{offset} \rangle$,
- Segment table – maps two-dimensional physical addresses; each table entry has:
 - base – contains the starting physical address where the segments reside in memory.
 - limit – specifies the length of the segment.
- Segment-table base register (STBR) points to the segment table's location in memory.



Address Mapping with Segmentation

(logical view)



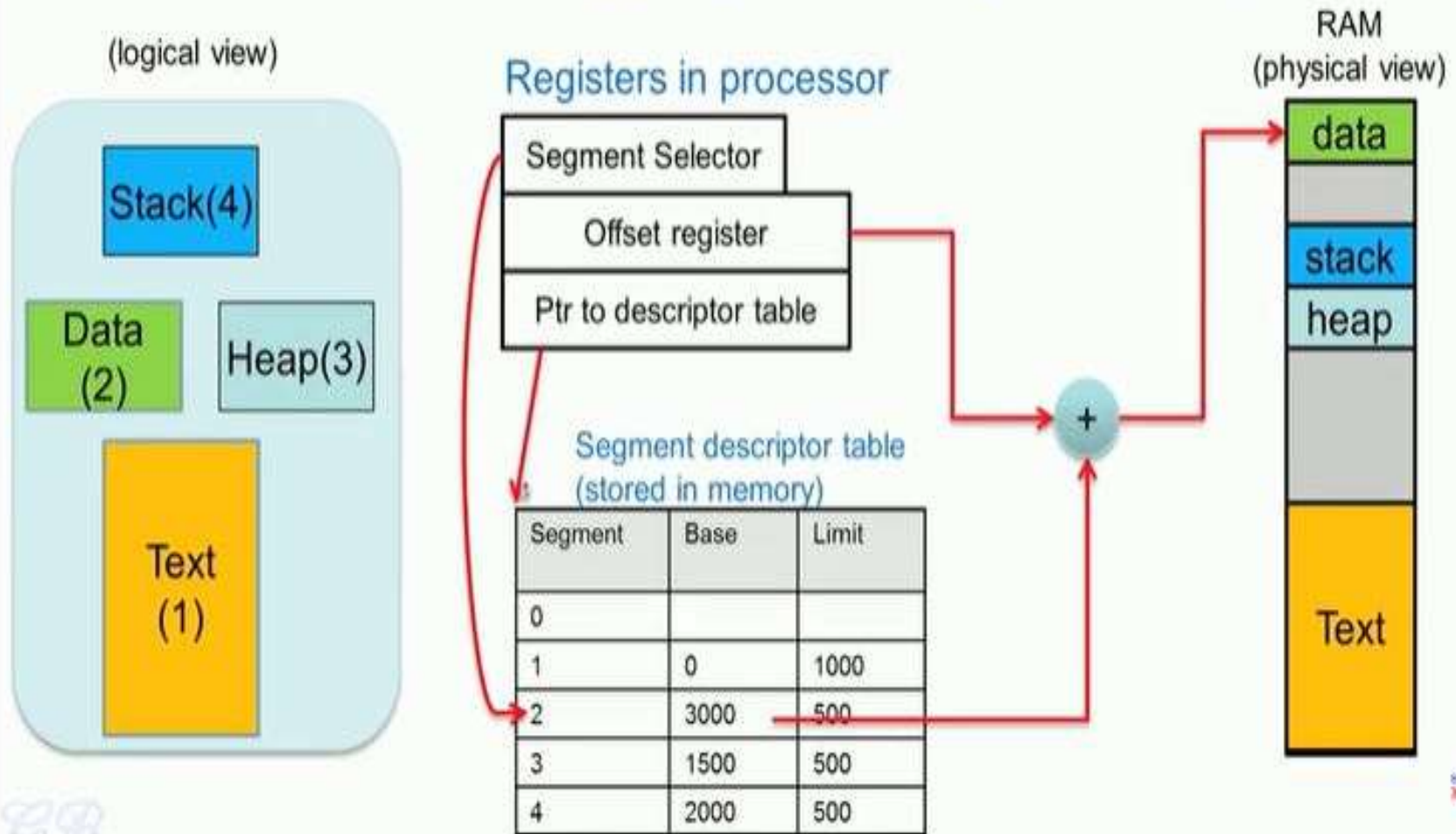
Segment descriptor table
(stored in memory)

Segment	Base	Limit
0		
1	0	1000
2	3000	500
3	1500	500
4	2000	500

RAM
(physical view)

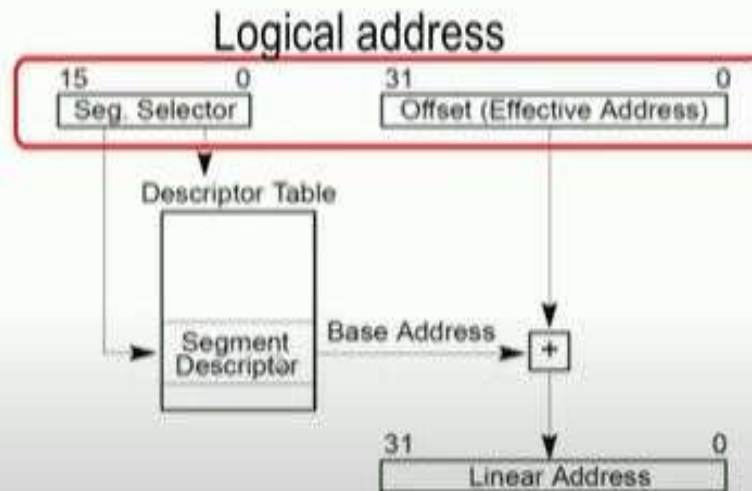


Address Mapping with Segmentation

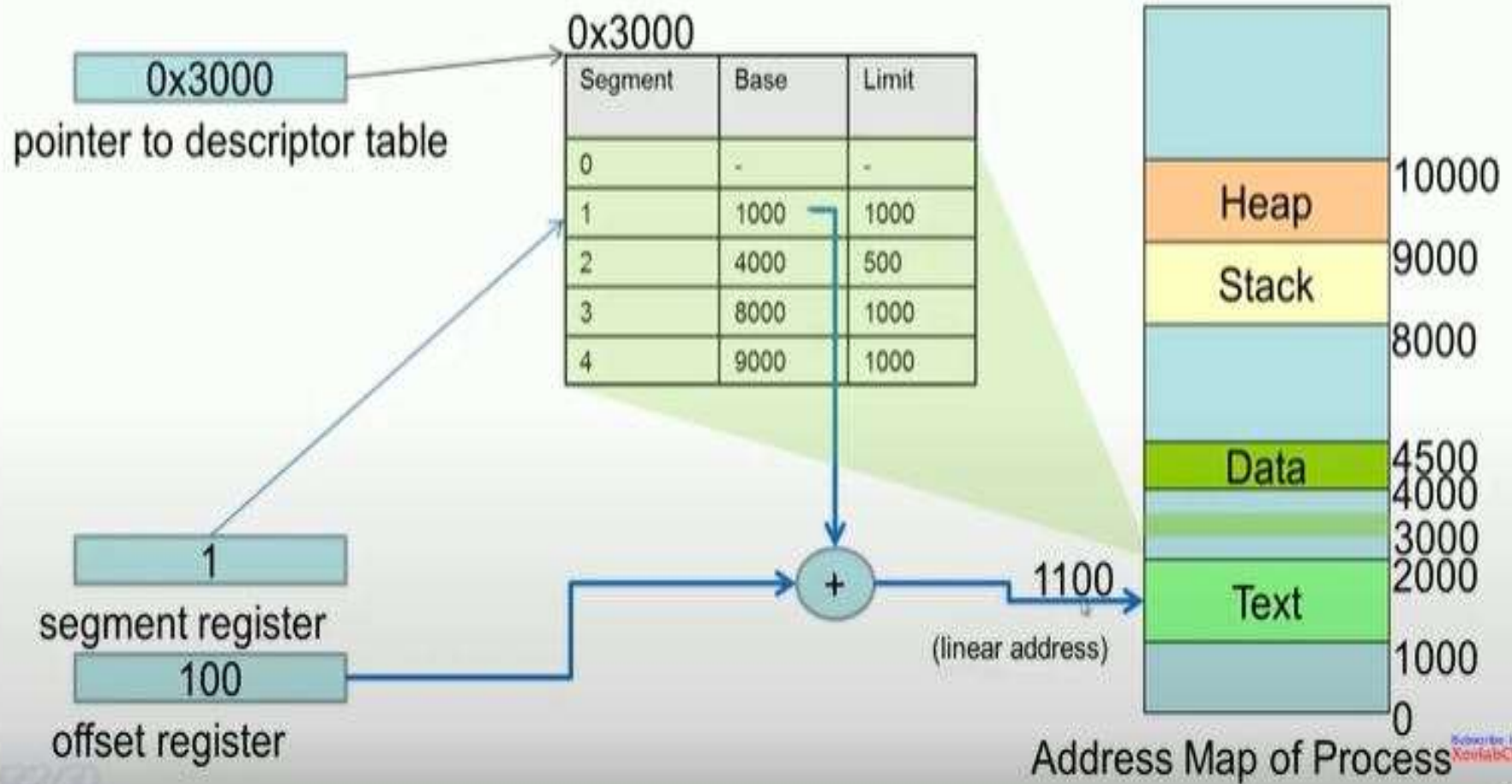


Segmentation

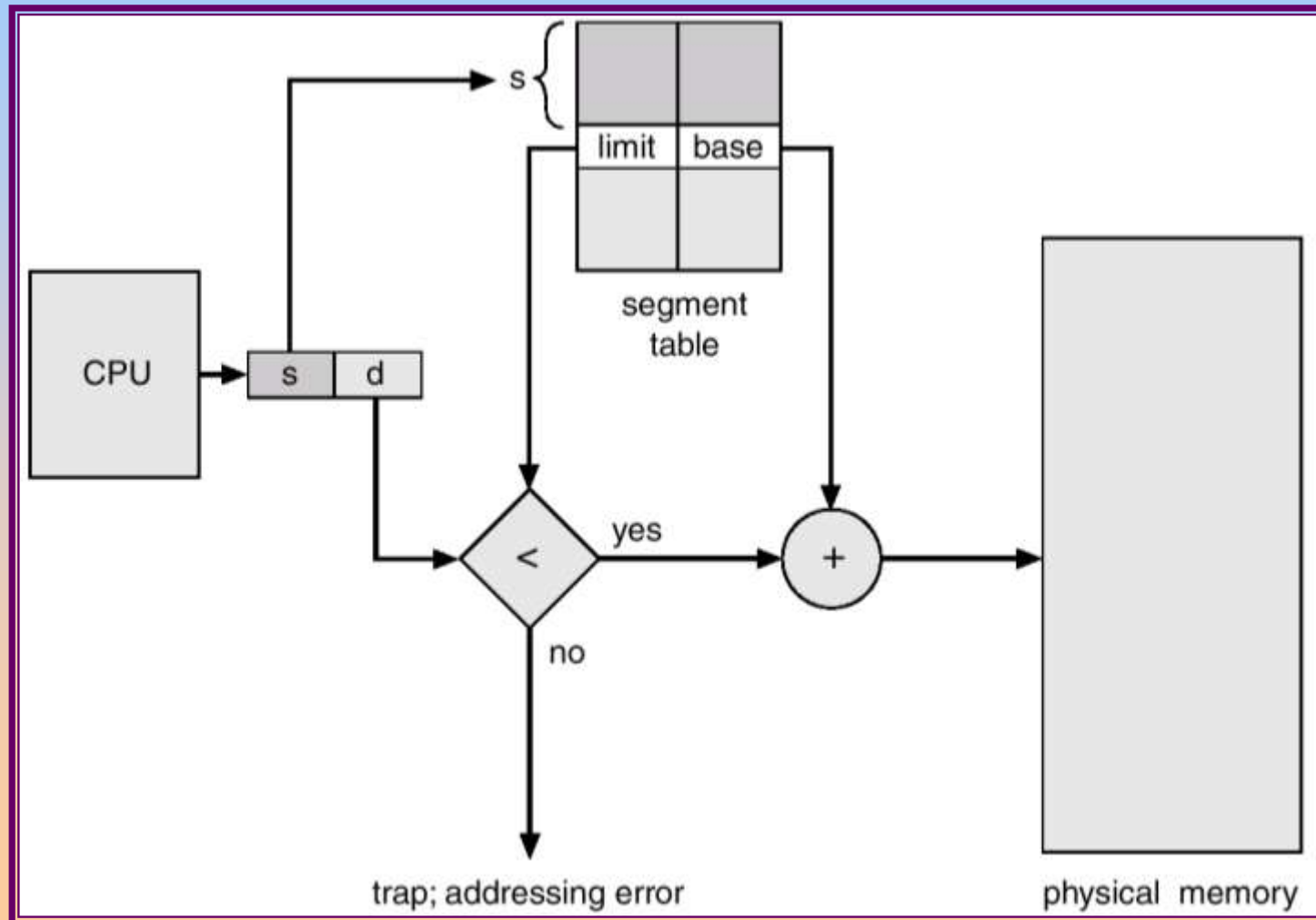
(*logical to linear address*)



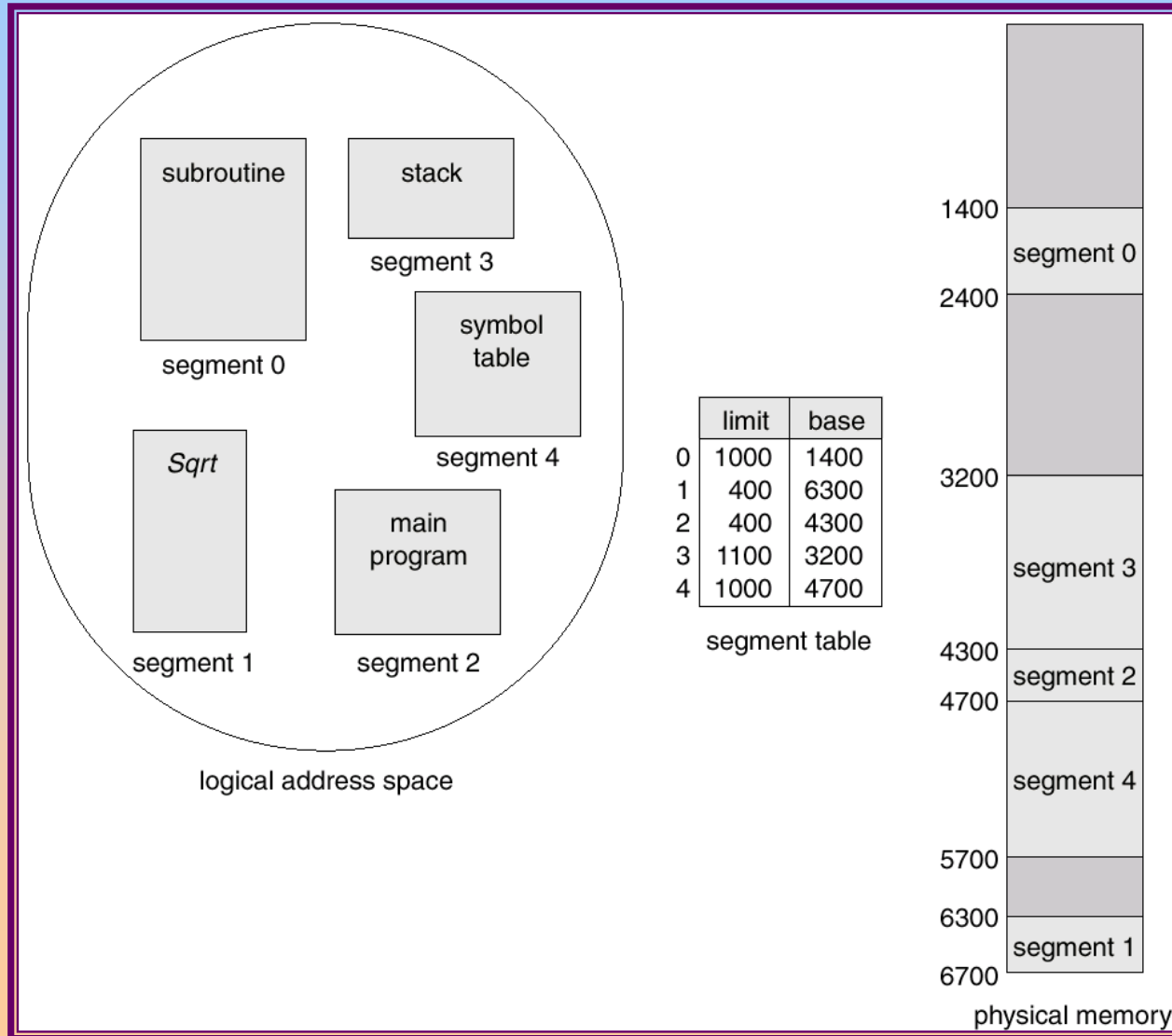
Example



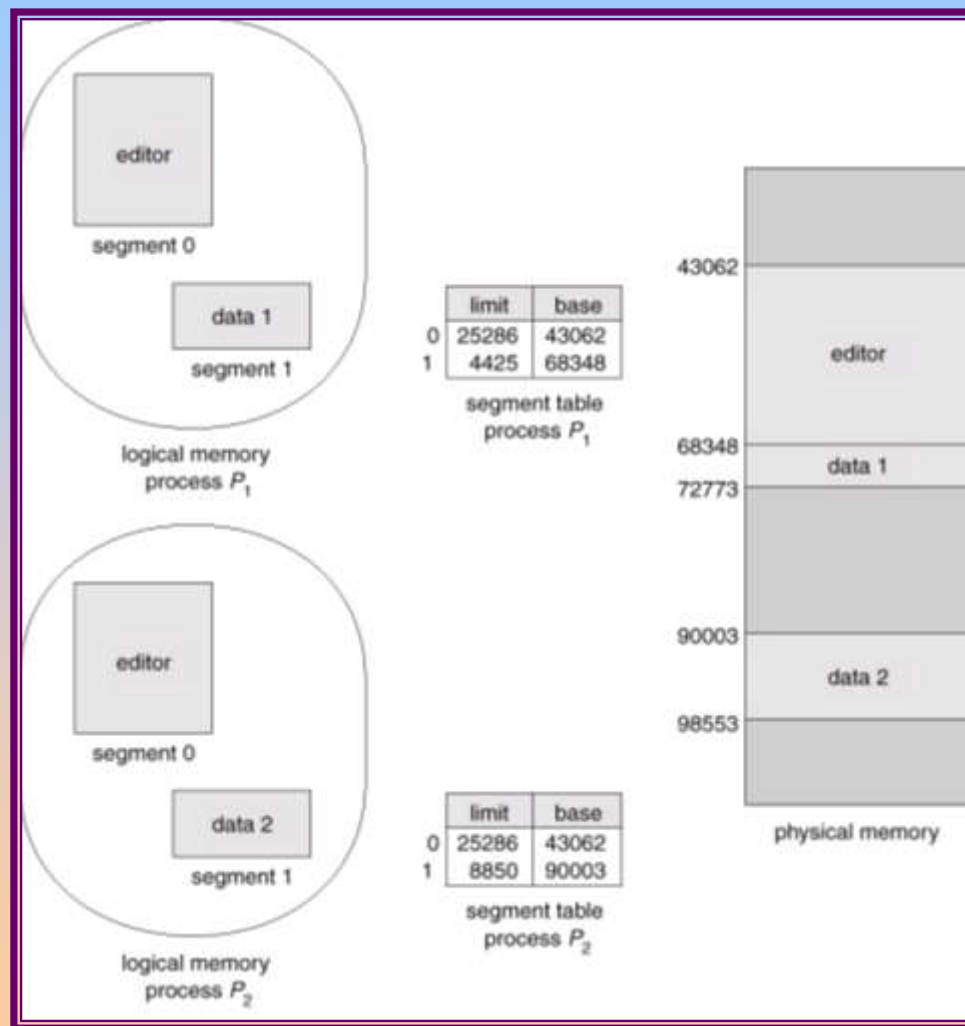
Segmentation Hardware



Example of Segmentation

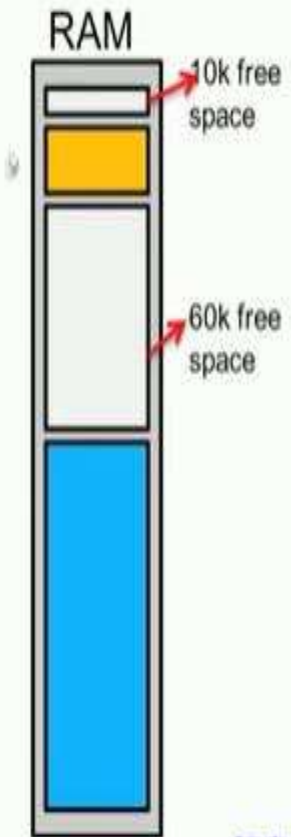


Sharing of Segments



Segmentation and Fragmentation

- Can lead to fragmentation
 - memory space is available but not contiguous



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