

# IT352: INFORMATION ASSURANCE AND SECURITY

## LAB 4 – Fiat Shamir Algorithm

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### Introduction:

One of the most well-known protocols identification using zero knowledge proof protocol is proposed by Amos Fiat and Adi Shamir, whose resistance is based on the difficulty of finding square root modulo a sufficiently large composite number  $n$ , the factorization is unknown.

### Steps involved in the Algorithm:

1. Select the module sufficiently large number  $n = p * q$ , factored that difficult.  $p, q$  are primes and kept secret
2. Client chooses a **secret  $s$**  belonging to interval  $(1, n-1)$  and relatively prime with  $n$ .
3. Public key is then calculated as  $V = s^2 \bmod n$ .
4. The resulting  $v$  registered as a trust center's public key of Client, and the value of  $s$  is the secret of Client. It is the knowledge of this secret  $s$  necessary to prove to the Client side without disclosure for  $t$  rounds.
5. Each round consists of:
  - a. Client chooses a **random  $r$**  in the interval  $(1, n-1)$  and **sends  $x = r^2 \bmod n$  to Server**
  - b. Server randomly selects a bit  $e$  aka **challenge (c) ( 0 or 1)** and sends it to Client.
  - c. Client computes  **$y = x * v^e \bmod n$**  and sends it back to Server.
  - d. Server checks the equality  **$y^2 = x * v^e \bmod n$** . If it is true, it proceeds to the next round of the protocol, otherwise the proof is not accepted.