

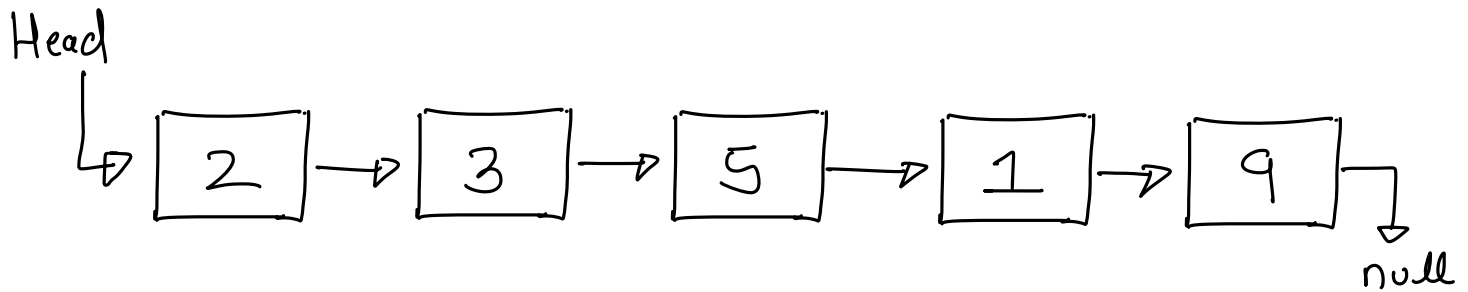
→ Linked list questions are usually not tough algorithmically.

→ Coding LL questions can be tricky

→ Always think about Edge Cases

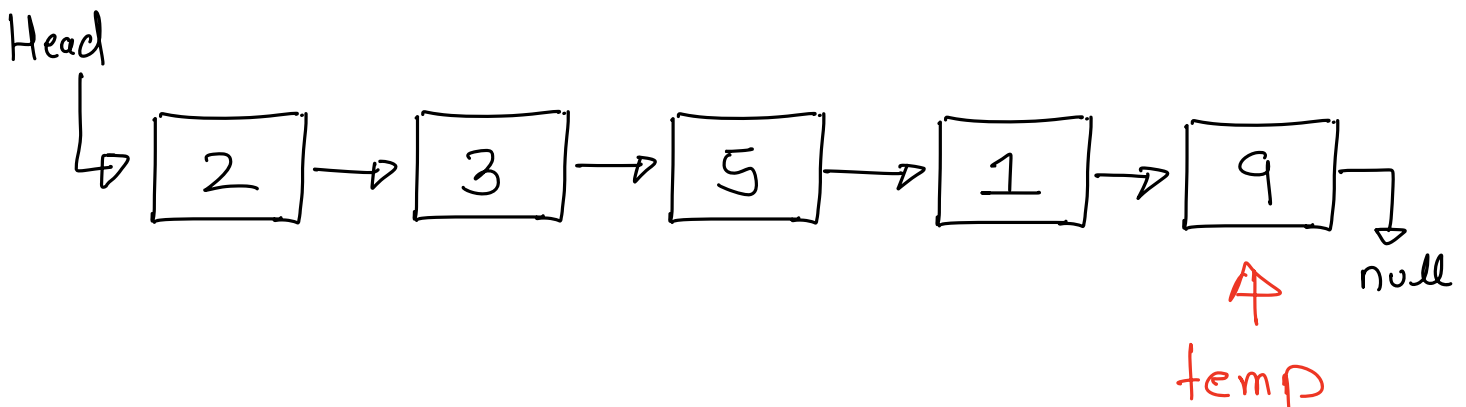
- 1) When linked list is null
- 2) When linked list has just one node
- 3) Problem specific edge cases.

Q₁ Search for an element **K** in given linked list.
Head node will be given. **Distinct elements** present.



K = 1

→ true



Node temp = head;

K = 9

while (temp != null) {

if (temp.val == K)
return true;

temp = temp.next;

}

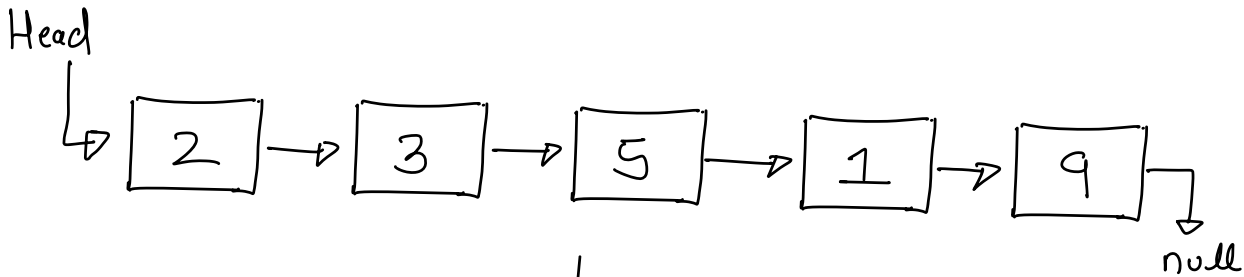
return false;

Tc: $O(n)$

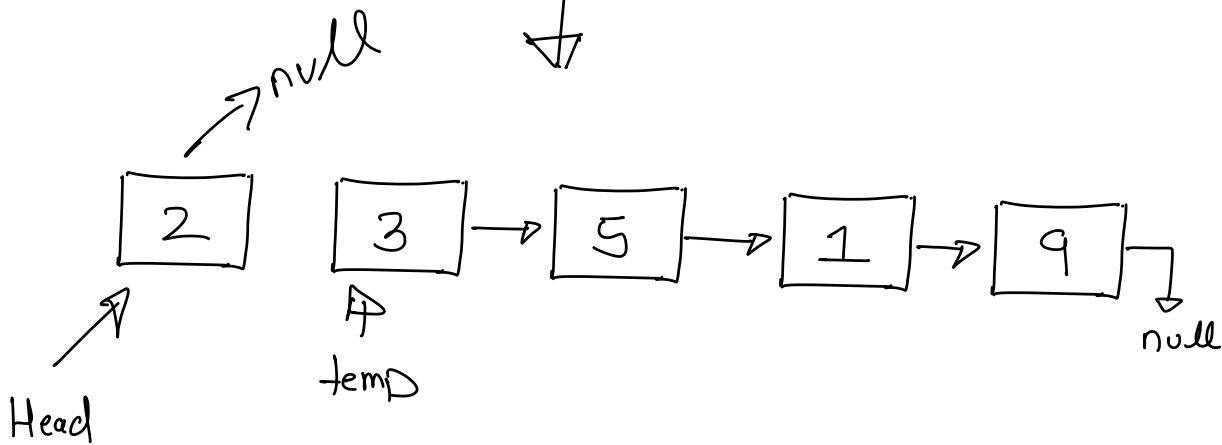
Sc: $O(1)$

Q2 Delete Node

→ at Head



Deletion.



```
if (head == null)
    return null;
```

```
Node temp = head
```

```
temp = temp.next
```

```
head.next = null
```

```
head = temp
```

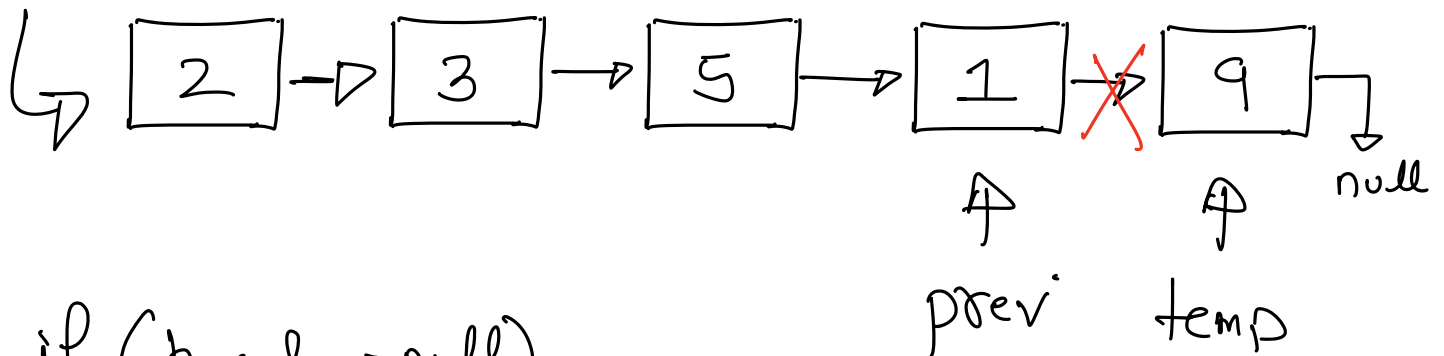
```
return head;
```

$T_c : O(1)$

$Sc : O(1)$

→ at tail.

Head



```
if (head == null)
    return null;
```

```
if (head.next == null)
    return null;
```

```
Node temp = head
```

```
Node prev = null
```

```
while (temp.next != null) {
```

```
    prev = temp;
```

```
    temp = temp.next;
```

```
}
```

```
prev.next = null
```

```
del(temp) → optional. Depends on how memory  
management works in your language.  
return head;
```

$TC: O(n)$

$SC: O(1)$

Node temp = head;

while (temp.next.next != null) {

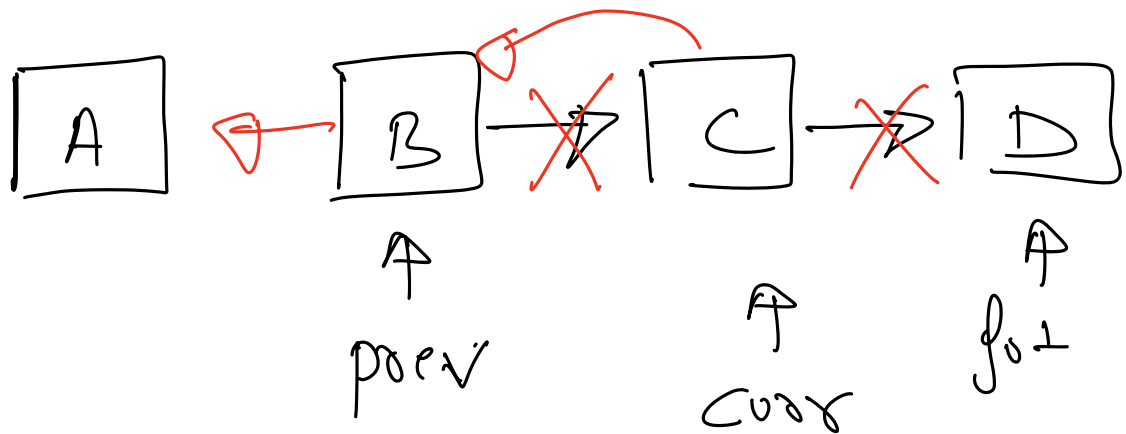
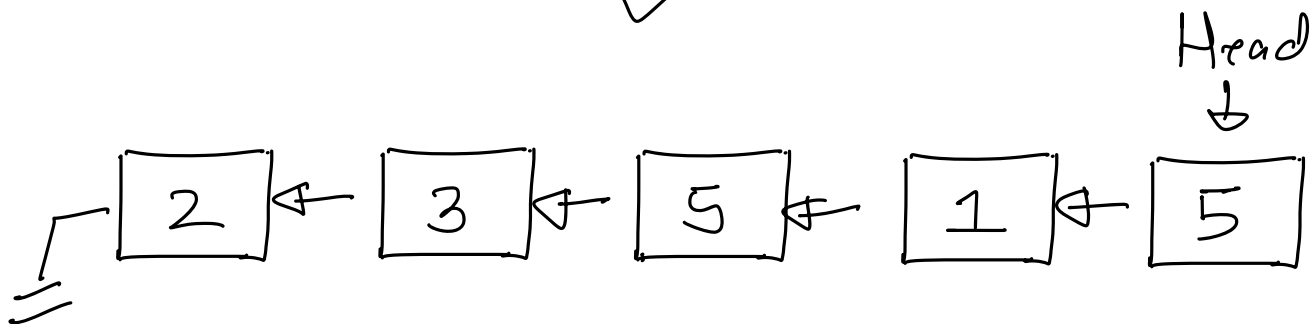
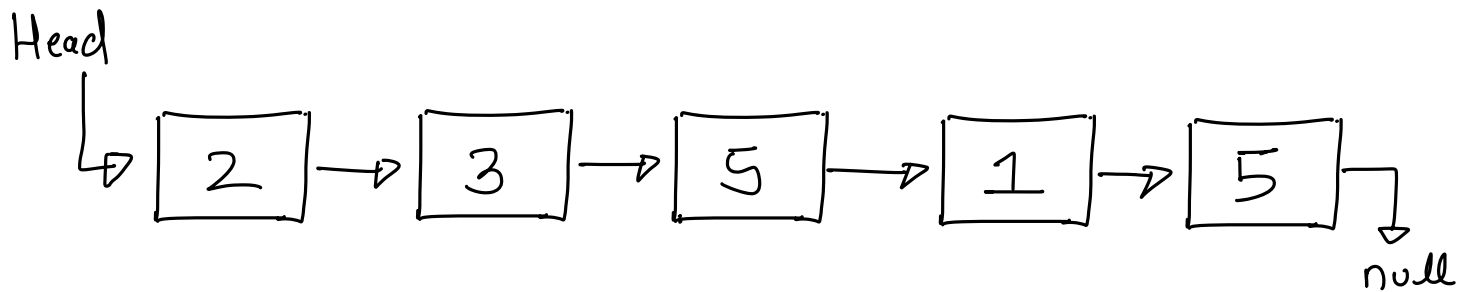
temp = temp.next;

}

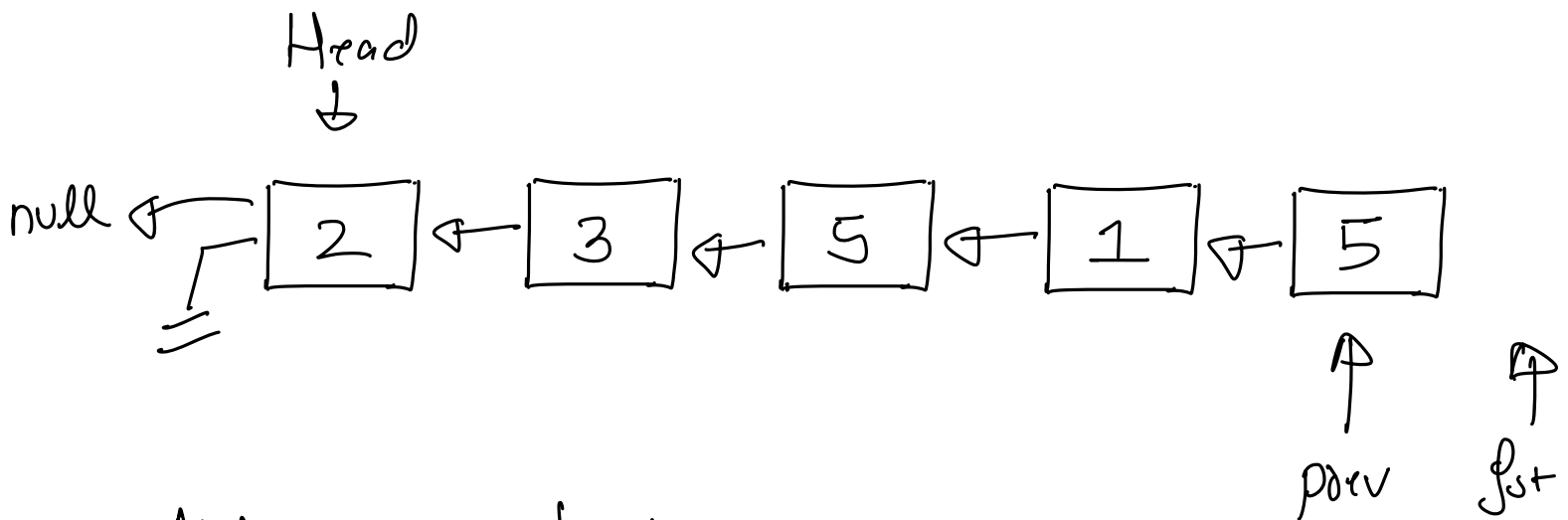
temp.next = null

10:18pm

★ Q4 Reverse the linked list. SC: $O(1)$



$next = curr.next$
 $curr.next = prev$
 $prev = curr$
 $curr = next$



Node curr \Rightarrow head

Node prev \Rightarrow null

Node fnt \Rightarrow null

while (curr != null)

fnt = curr.next

curr.next = prev

prev = curr

curr = fnt

3

return prev;

Tc: $O(n)$

Sc: $O(1)$

~~curr.next != null~~

fnt != null

curr != null

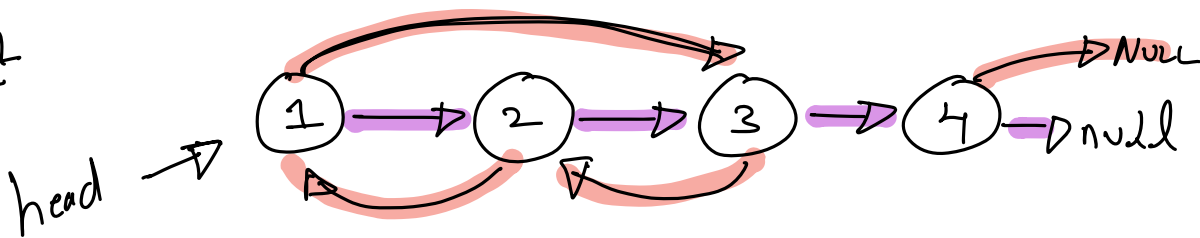
Q5 Given a Linked list. Every node has 2 pointers.

1) next pointer : As usual.

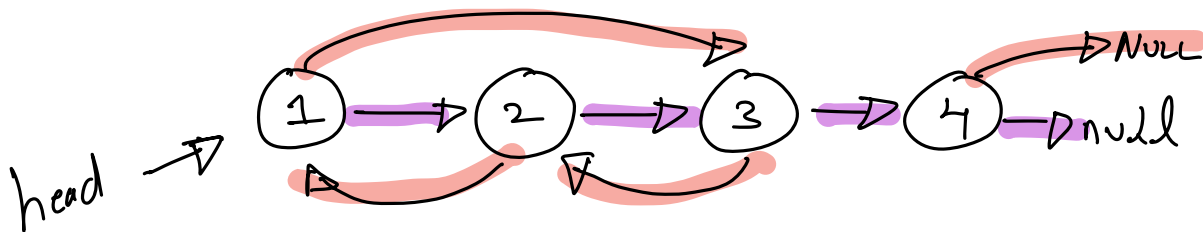
2) random pointer : Can point to any node or null

Make a clone of this. Brand new n nodes with the same structure.

Ex 1



↓ new set of n nodes.




```
class Node {
```

```
    int val;
```

```
    Node next;
```

```
    Node random;
```

```
Node (int val) {
```

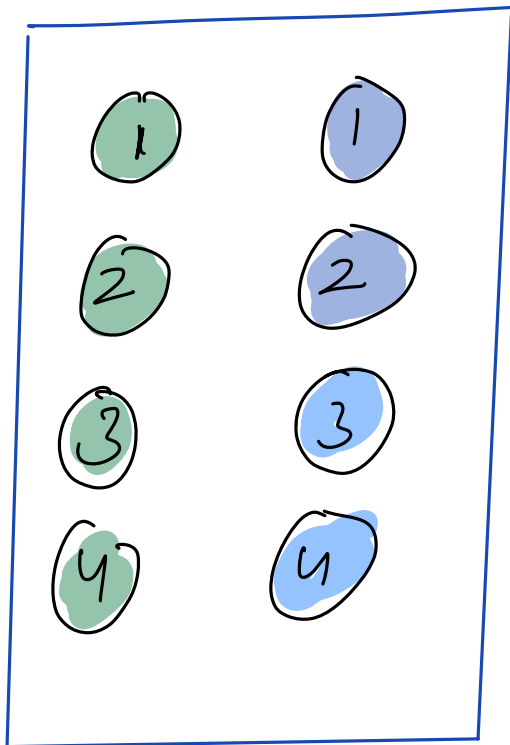
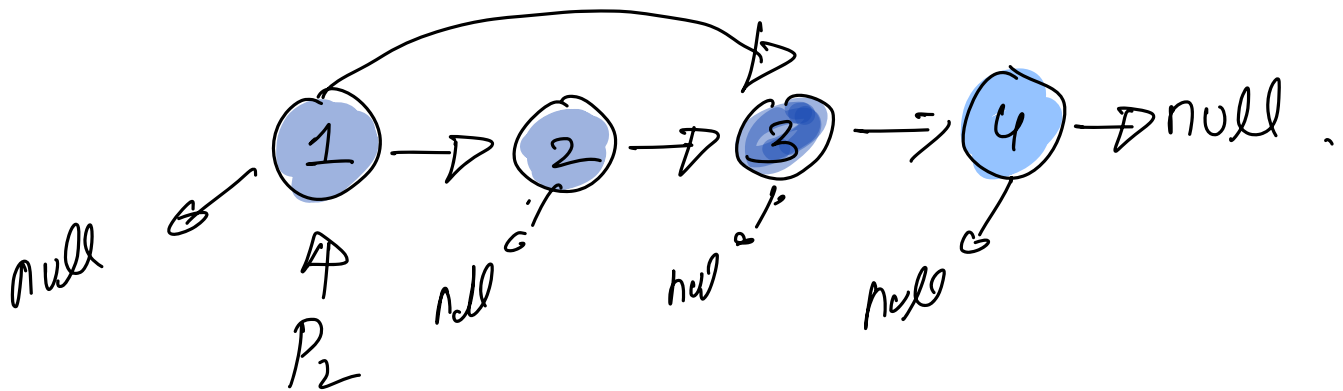
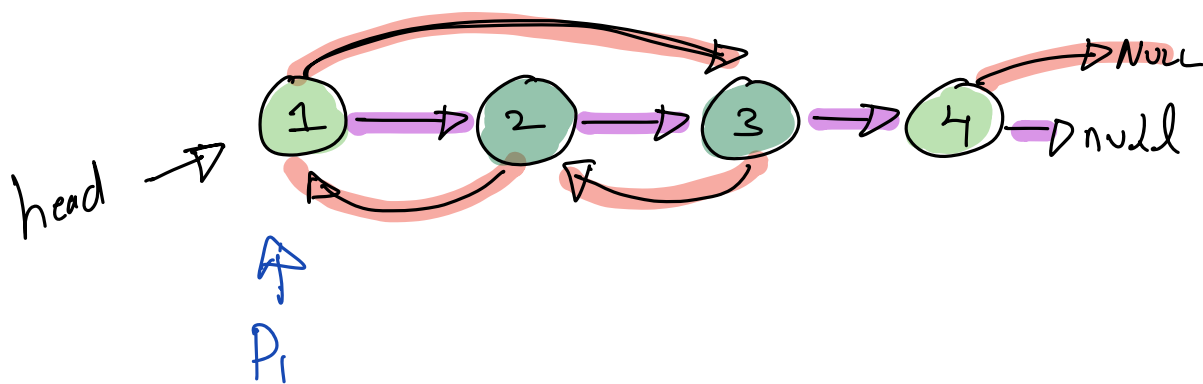
```
    this.val = val;
```

```
    this.next = null;
```

```
    this.random = null;
```

```
}
```

Approach 1:

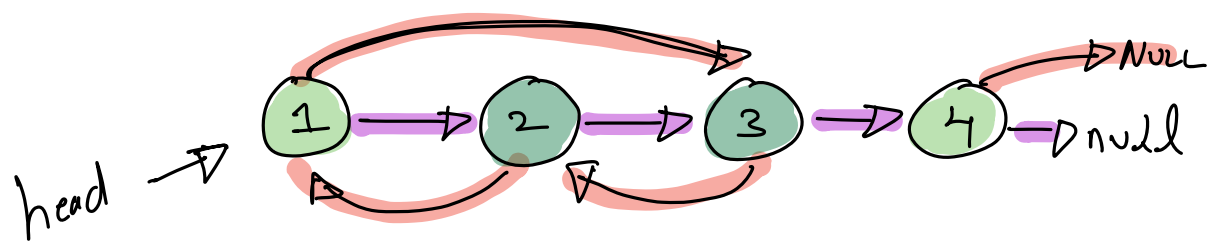


$p2.random \Rightarrow$
 $\Rightarrow mp.get(p1.random)!$
=====

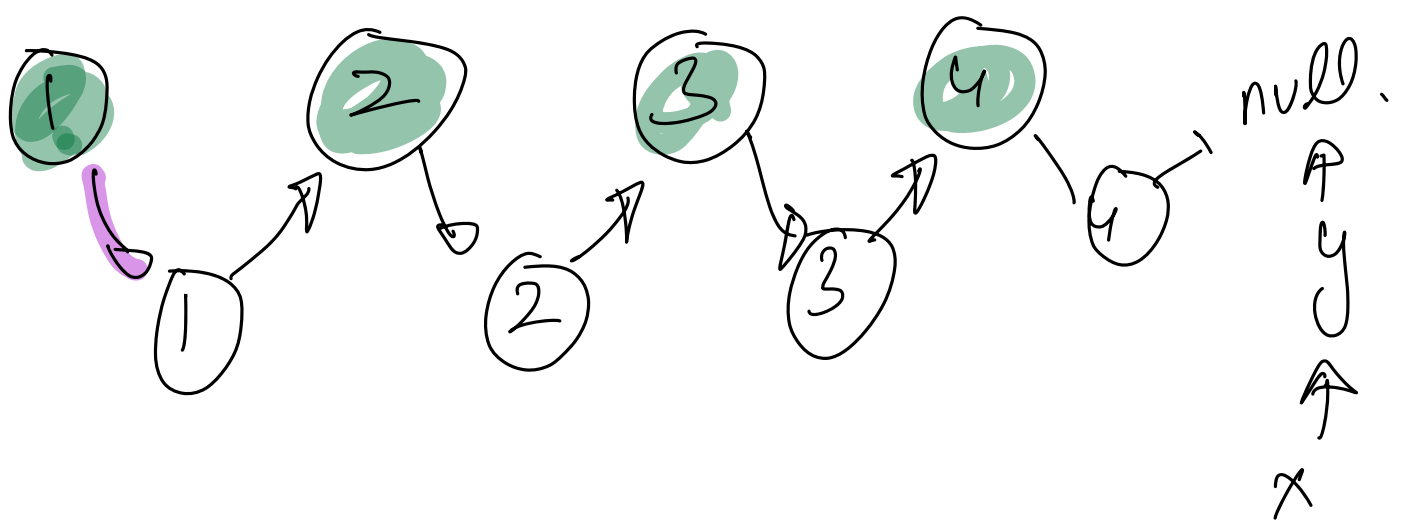
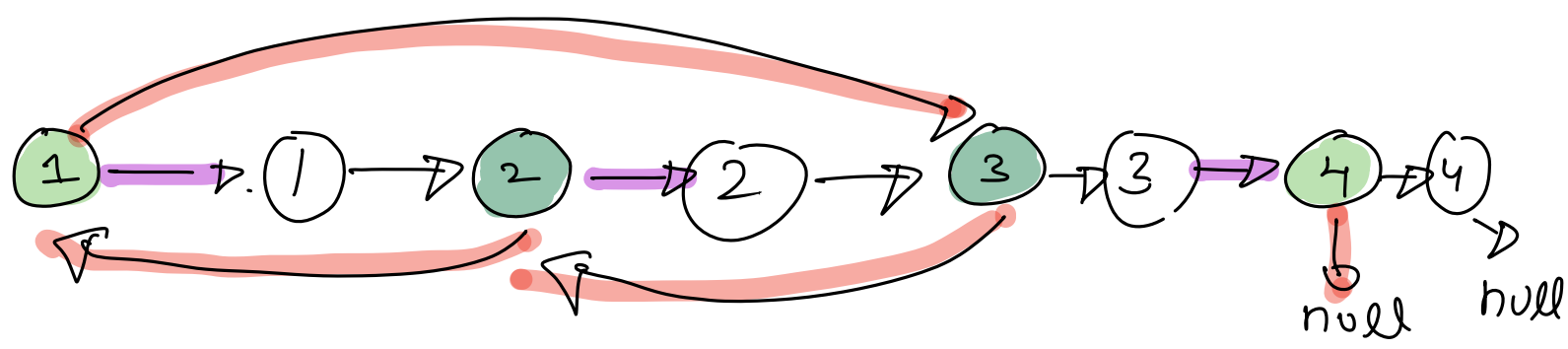
$T_c : O(n)$
 $Sc : O(n)$

$Map < Node, Node > mp$

Approach 2 [Reduce Space]



Step 1 : Add new nodes after old nodes in the same linked list



$x \Rightarrow \text{head};$

$y \Rightarrow \text{head.next};$

while ($x \neq \text{null}$) \hookrightarrow

$x.\text{next} = \text{new Node}(x.\text{val})$

$x.\text{next}.\text{next} = y;$

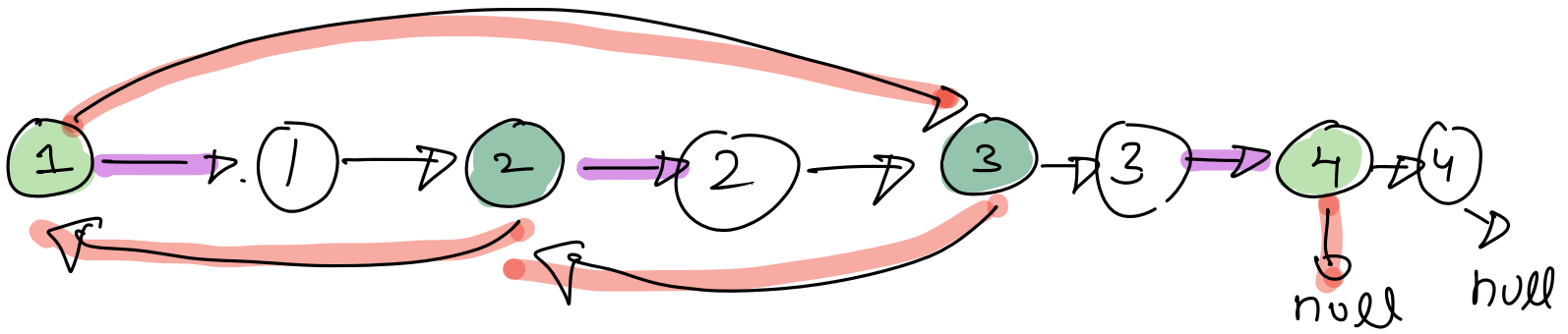
$x = y;$

if ($y \neq \text{null}$) \hookrightarrow

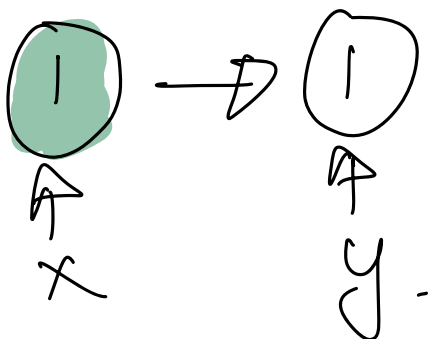
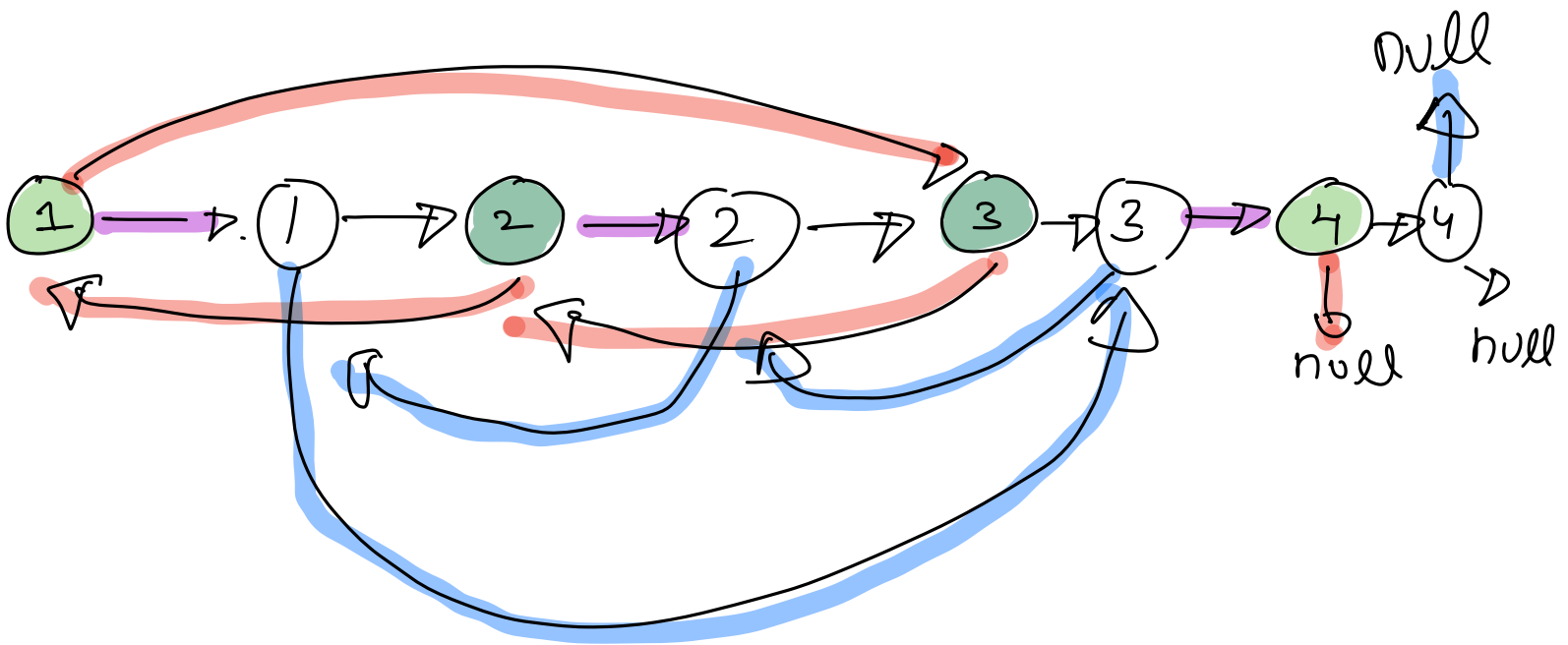
$y = y.\text{next};$

3

3

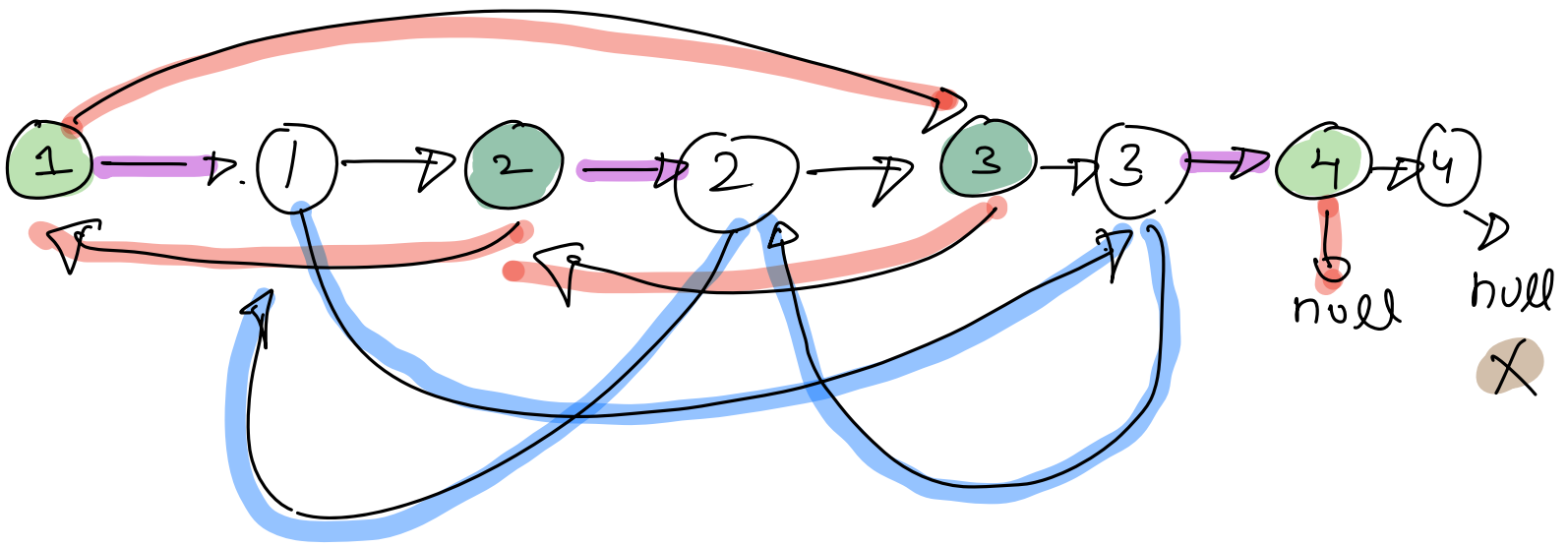


Step 2 : Point random pointer of new nodes correctly.



$y \cdot \text{rand} \Rightarrow x \cdot \text{rand} \cdot \text{next}$

$x \cdot \text{next} \cdot \text{rand} = x \cdot \text{rand} \cdot \text{next}$



Node. $x \Rightarrow \text{head};$

while ($x \neq \text{null}$) {

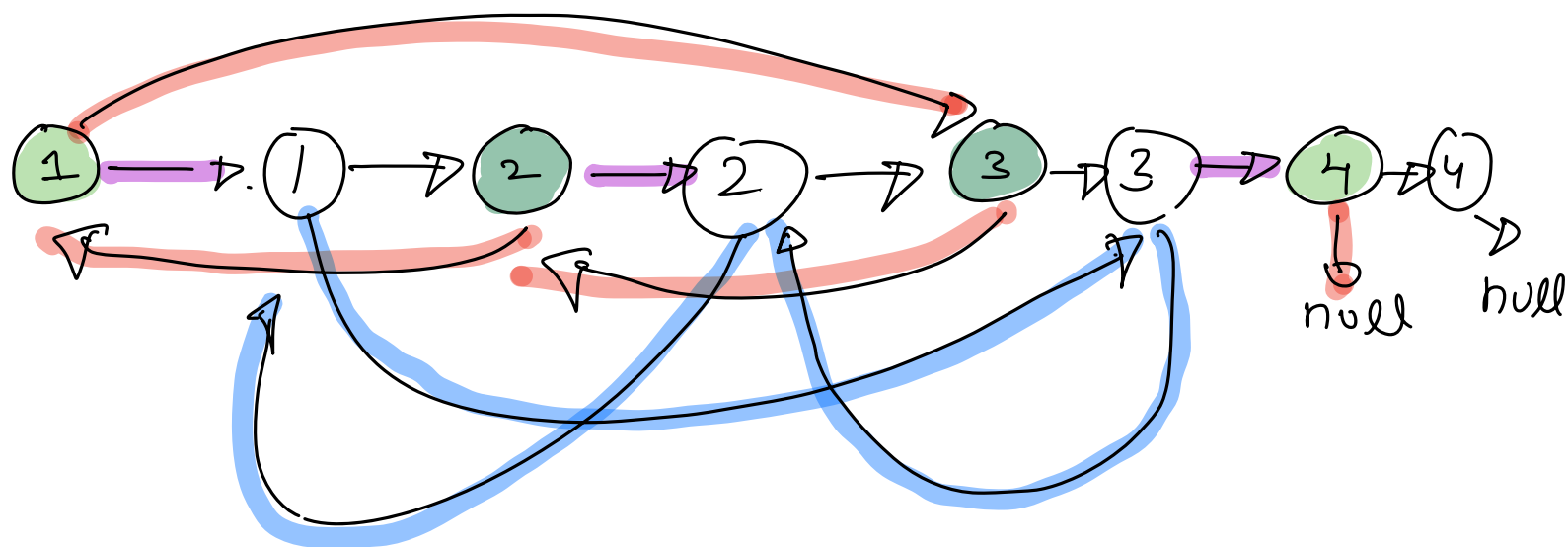
if ($x.\text{rand} \neq \text{null}$) {

$x.\text{next}.\text{rand} = x.\text{rand}.\text{next}$

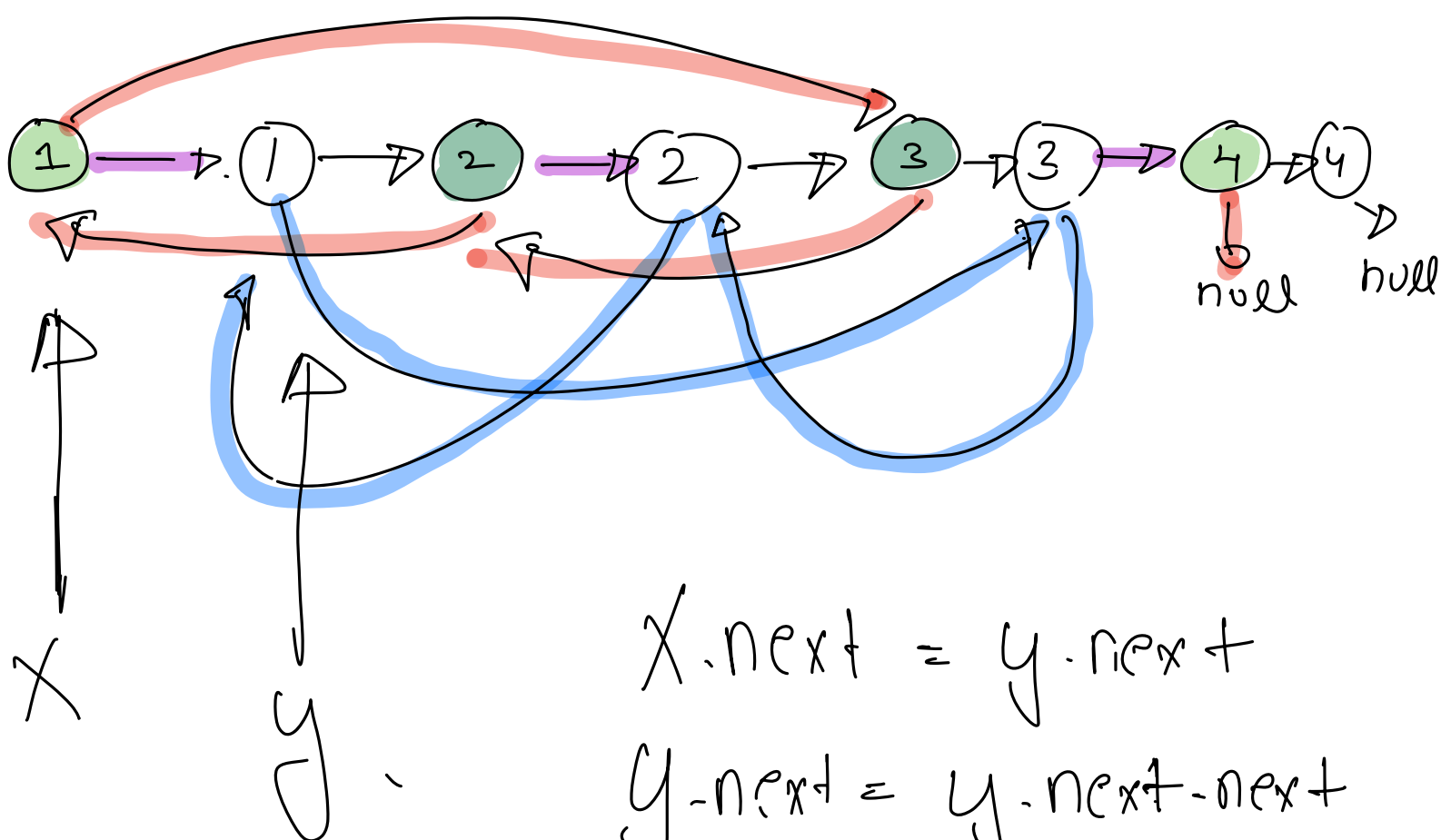
}

$x = x.\text{next}.\text{next};$

}



Step 3: Decouple the linked list

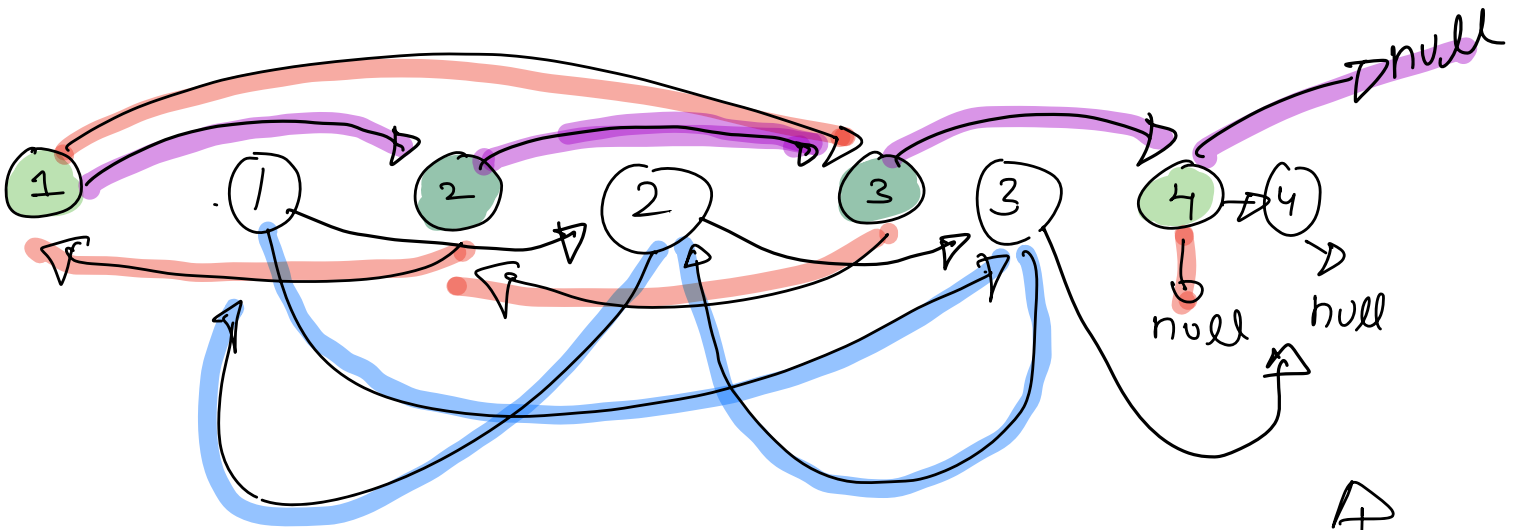


$$X.next = Y.next$$

$$Y.next = Y.next.next$$

$$X = X.next$$

$$Y = Y.next$$



$x = \text{head};$

$y = \text{head.next}$

while ($x \neq \text{null}$) $\{$

$x.\text{next} = y.\text{next}$

if ($y.\text{next} \neq \text{null}$) $\{$

$y.\text{next} = y.\text{next}.\text{next}$

$x = x.\text{next}$

$y = y.\text{next}$

$\}$

