

TRANSACTIONS

Agenda

Transaction:

ACID

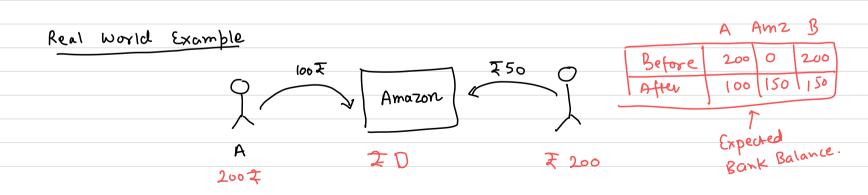
Commits & Rollbacks

Transaction Isolation Levels

RUC, RC, RR, Gerializable

Deadlocks

0

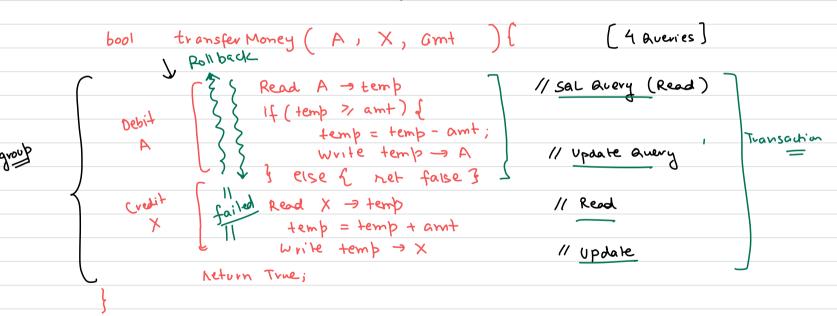


Next class

Bank Balance

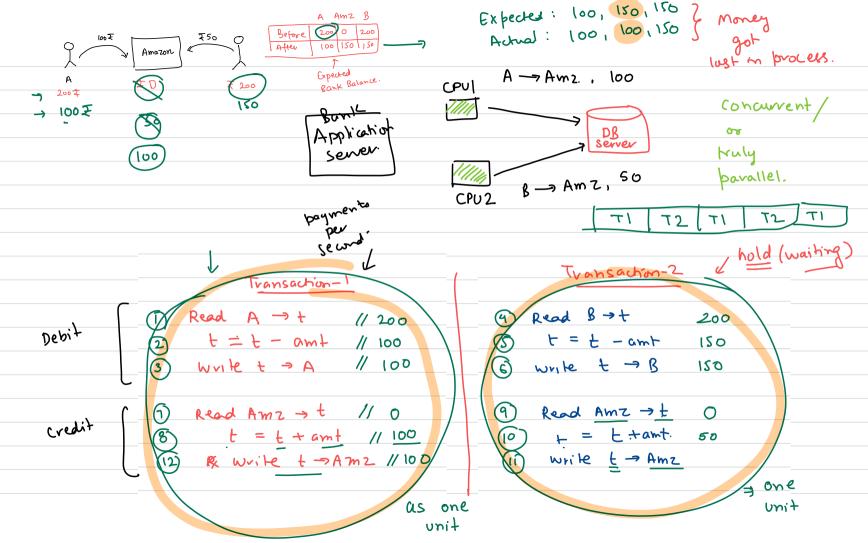
	uid	AIC	Bal.	
1	A	=	200 -	-100 L
	B	=	200	
	Am2	=	<u></u> (6)	

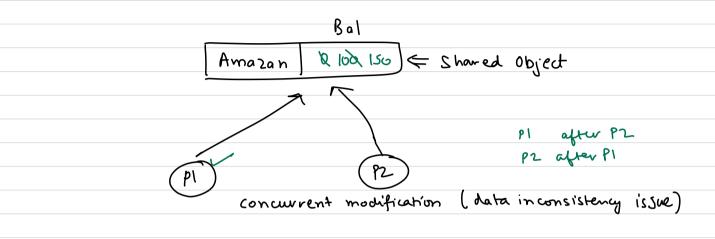
Table (Bank Balance)



Concurrency

Concurrent Behaviour	Notes	Zom	other	20 apps
When	app		/apps/	11
multiples	•			
ordosses				
make		1 core CPU	motes os	M W
progress		CPU	2 00ms	
α \-			Contex	+ Switching
Same			- takis - Cashes	Zoms ->
h'me:	2	core	Notes Zoom N	ster (amera)
	Core	s core?	OS Chro	m 05]
	1			
	h brocess has	a priority.	=) get more lless	CPU time
		, , , , , ,	it is decided	by OS Scheduler.
			• •	0





Transaction:	is	a set	of	DB operations	logically groupe	Ł
Turns action:		together	to	perform a	task.	

- -> money Transfer
- -> Railway Ticket Booking
 -> Movie ticket Booking

ACID Properties

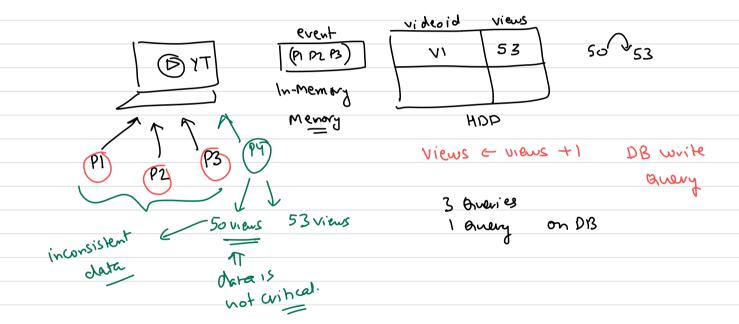
Not necessary for a transaction to have all 4 properties. depends upon use-case.

- 1) Atomicity
- 2) Consistent
- 3) Isolation
- 4) Durability
- (i) Atomicity: Opr/trans should appear as atomic/one entity to an end user.

UPI Payment -> single outcome & fail (nothing has happened)

2) Consistency - logically correct

Banking => Consistency is expected.



3	Isolation

one transaction shouldn't affect another transaction running at same on same DB in a wrong way.

"Lock" > mechanism to provide "isolation" Isolation

(4) Durable

Lo long lasting
Lo once change is made, it should pensist in DB forever.

- · Date HDD/SD cand + Backup
- · Cloud Storage (Moltiple Backups)

across different Regions of world.

Commit & Roll back

Saving

Crayes

to

DB

(Ourable)

Session-1

Session-2

Changes

Changes

Levels.

Session-2

What 'State'

of data

you will

end

Up reading.

4 Transaction Isolation Levels defined for every session. Lenient 1) Read Uncommitted (RU) = dirty read problem 2) Read Committed (RC) 3) Repeatable Read (RR) [Default in MySQL] 4) Sevralizable (i) 0 = 10 Session 2 (RUC Session (RR)

RUC

good: efficient in execution, hu locks or concurrency control are regd. divty Reads, inconsistent data. Init: 100, + 0, 100 finally: 50 100, 100, 50 => 2507 A: 100 50 100 B → Amz 8:100 Amz: 0 lead B → + (100) amt: So t = t -50 (50) write t-7 A (SD) write t-7 B (50) RUC Read Amz $\rightarrow \pm$ (0) $t = \pm +50$ (50) write $t \rightarrow Amz$ (50) failure Read Amz -> t (50) (aused (5) t = t + 50 (100)
(7) WYI'LE t → AMZ problem Commits, Dirty Read