Bit manipulations.

Bitwise operators → `&`, `|`, `^`, `~`, `<<`, `>>`

same same puppy shome

A	B	A&B	A|B	A^B	~A	~B
0	0	0	0	0	1	1
0	1	0	1	1	1	0
1	0	0	1	1	0	1
1	1	1	1	0	0	0

A = 20
B = 45

A = 0 0 0 1 0 1 0 0
B = 0 0 1 0 1 1 0 1

print (A & B) = 4 — 0 0 0 0 0 1 0 0 (4)

print (A | B) = 61 — 0 0 1 1 1 1 0 1 (61)

print (A ^ B) = 57 — 0 0 1 1 1 0 0 1 (57)

Commutative

a & b = b & a

a | b = b | a

a ^ b = b ^ a

Associative

a & b & c = b & a & c = c & a & b

"

"

a & 0 = 0	a | 0 = a	a ^ 0 = a
a & a = a	a | a = a	a ^ a = 0

Right shift operator

a = 50 : 0 0 1 1 0 0 1 0 — 50 $(2^1+2^4+2^5)$

a >> 1 : 0 0 0 1 1 0 0 1 — 25 $(2^0+2^3+2^4)$

a >> 2 : 0 0 0 0 1 1 0 0 — 12

50 >> 1 → $\frac{50}{2^1}$ → 50 >> 2 → $\frac{50}{2^2}$

$a >> n = \frac{a}{2^n}$

Left shift operator

a = 5 : 0 0 0 0 0 1 0 1 — (2^0+2^2)

a << 1 : 0 0 0 0 1 0 1 0 — 10 (2^1+2^3)

$20 = (2^2+2^4) =$

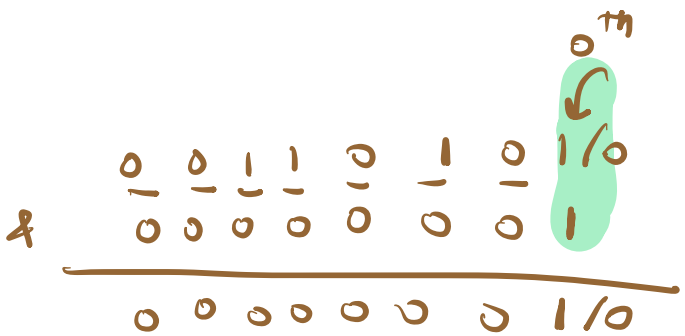
$a << N = a \times 2^n$

Q. Given a number N and i. Check if i^{th} bit is set or not.

$N = 53$ (bit positions 6 5 4 3 2 1 0): 0 1 1 0 1 0 1

$i = 3$ ans = False

$i = 2$ ans = True



$i = 0$ $N \& 1 == 1$

$i = 2^{th}$ $(N >> 2) \& 1 == 1$

```
if ( (N>>i)&1 == 1 )   return True
else                   return False
```

$N = N >> i$

Tc: O(1)
Sc: O(1)

Q. Given N elements. Every element repeats twice. Except one. Find unique element

arr[] : {3, 2, 3, 7, 2, 8, 7}

ans : 8

ans = xor of all elements

$a \wedge a = 0$

$ans = 3 \wedge 2 \wedge 3 \wedge 7 \wedge 2 \wedge 8 \wedge 7$

$= 3 \wedge 3 \wedge 2 \wedge 2 \wedge 7 \wedge 7 \wedge 8$

$= 0 \wedge 0 \wedge 0 \wedge 8$

$= 0 \wedge 0 \wedge 8$

$= 0 \wedge 8$

$= 8$

```
int val = 0
for (i=0; i<N; i++)
{
    val = val ^ arr[i]
}
return val
```

T.C : O(N)
S.C : O(1)

Q. Every element repeats thrice. Except one [unique]
Find unique element.

arr[] = {5, 7, 5, 4, 7, 11, 11, 9, 11, 7, 5, 4, 4}

ans = 9

BF : For each ele, go & count freq
if freq == 1
ans

TC : O(N²)
SC : O(1)

Idea2 : Use freq hash map

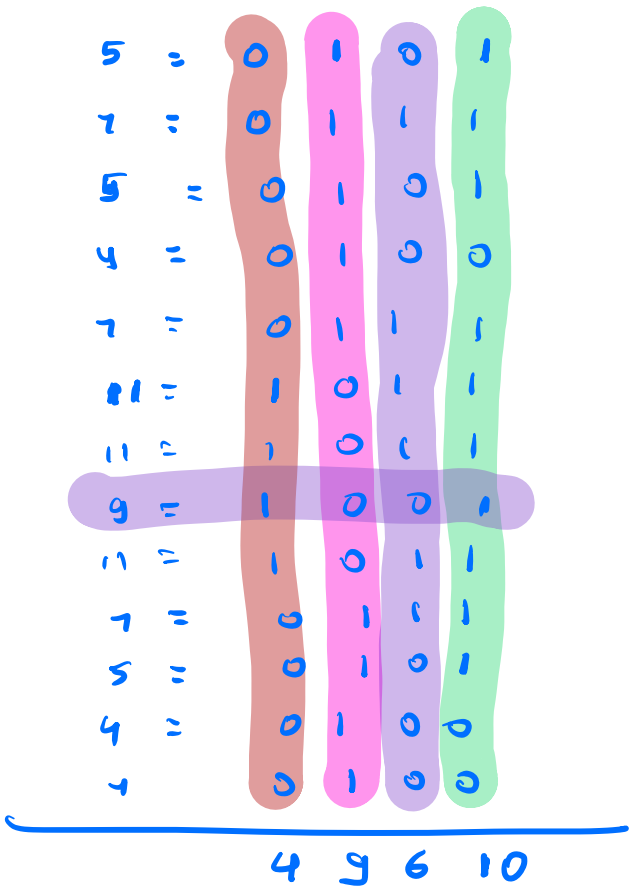
→ store freq of each element in hm

→ Iterate array again
if freq == 1
return ans

TC : O(N)
SC : O(N)

Idea 3.

{5, 7, 5, 4, 7, 11, 11, 9, 11, 7, 5, 4, 4}



10 % 3 = 1

1 0 0 1

$2^3 + 2^0 = 9$

At each bit, Calculate no of Element having this bit set (K)

If (K % 3 == 1) this bit is set in ans

else this bit is unset in ans

```
int ans = 0

for ( i=0; i<32; i++ )
{
    // ith bit

    cnt = 0
    for ( j=0; j<N; j++)
    { if ( checkbit ( j, i ) == True)
        cnt++

    if ( cnt % 3 == 1 )
    { ans = ans + 2^i
}

return ans
```

TC: $O(32N) \approx O(N)$

SC: $O(1)$

Q. Every element repeats thrice. Except one [twice]

Find unique element.

cnt % 3 == 2

Q. Every element repeats 4 times except one [once]

Find unique element.

`cnt % 4 == 1`

and

ans = XOR of all elements

Q. Given N elements. Every element repeats twice except 2 unique elements.

Find those 2 unique numbers.

arr[6] = 3, 4, 6, 4, 3, 8

ans = 6, 8

BF For every, check freq.

If freq == 1

return ans1, ans2

$TC: O(N^2)$

$SC: O(1)$

Use hashmap

$TC: O(N)$

$SC: O(N)$

arr[6] = 3, 4, 6, 4, 3, 8

3 ^ 4 ^ 6 ^ 4 ^ 3 ^ 8

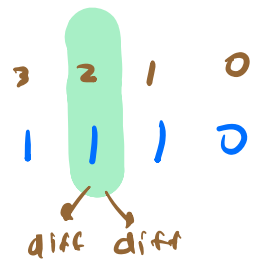
= 3 ^ 3 ^ 4 ^ 4 ^ 6 ^ 8

= 6 ^ 8

= 14

0 1 1 0
0 0 0 0
---
1 1 1 0

XOR (A ^ B) = 14



3, 4, 6, 4, 3, 8

2th bit is set

4, 6, 4

ans1 = xor

2th bit is unset

3, 3, 8

ans2 = xor

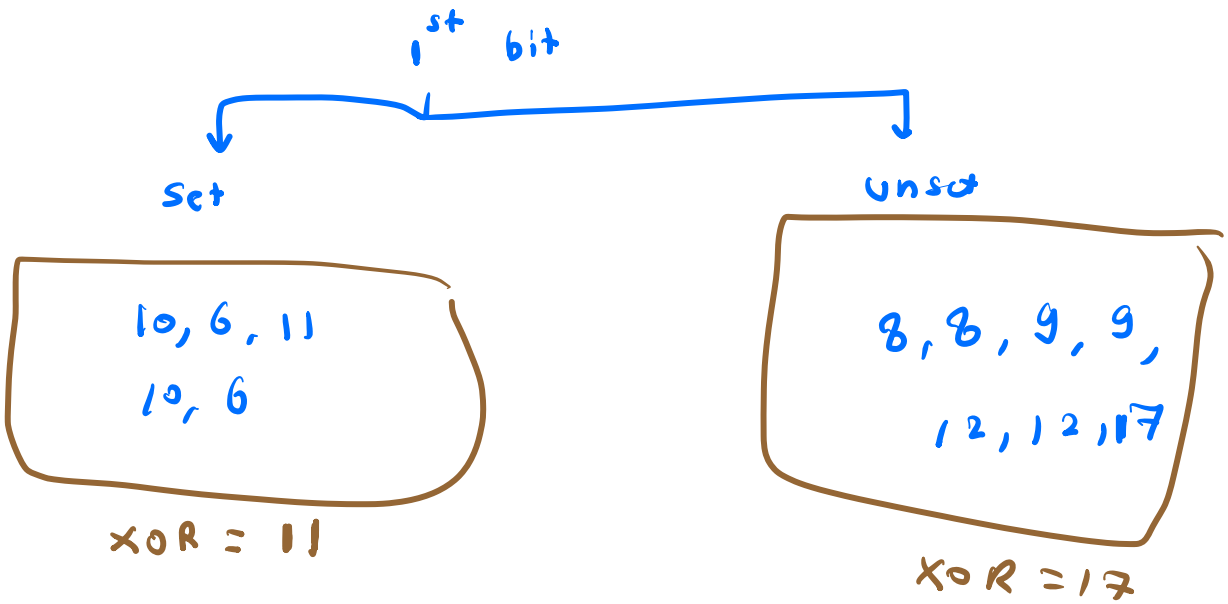
10, 8, 8, 9, 12, 9, 6, 11, 10, 6, 12, 17

XOR (array) = 26

4 3 2 1 0

1 1 0 1 0

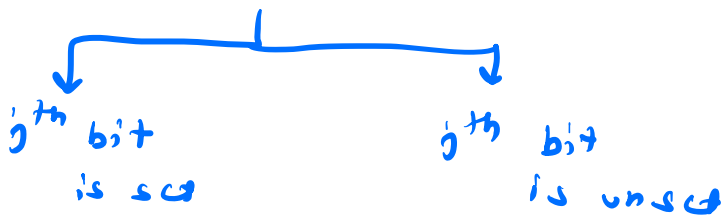
16+8+2



1. Calculate XOR of the array. = K

Todo 2. Take any set bit (K) = j^{th} bit

3. 2 Groups



ans1 = XOR(G1) ans2 = XOR(G2)

```
xor1 = 0
xor2 = 0

for( i=0; i<N; i++)
{
    if( checkbit(Arr[i], j) == True)
        xor1 = xor1 ^ arr[i]
    else
        xor2 = xor2 ^ arr[i]
```

ans1 = XOR1
ans2 = XOR2

TC: O(N)
SC: O(1)

Q. Given N array elements. Choose i, j
s.t arr[i] & arr[j] is maximised

Find max AND

arr[3] = {27, 18, 20}

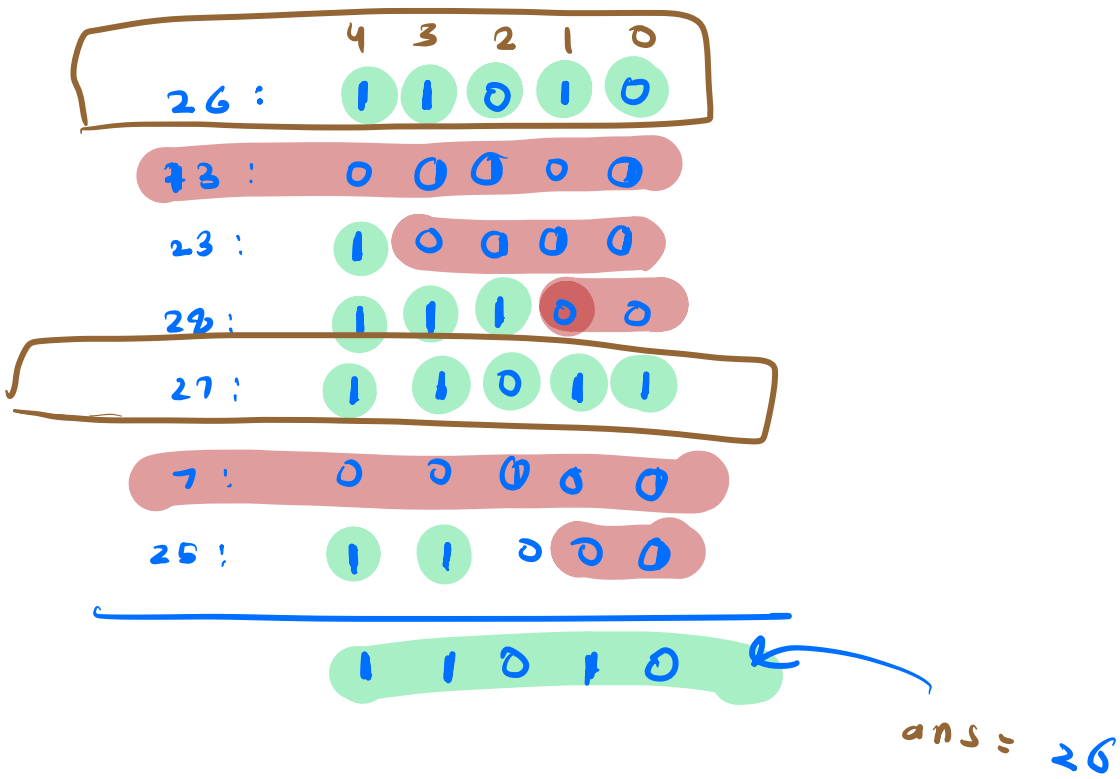
27 & 18 = 10010 = 18
27 & 20 = 10000 = 16
18 & 20 = 10000 = 18

27 = 11011
18 = 10010
20 = 10100

BF: Consider each pair, pick max AND

TC: O(N^2)
SC: O(1)

arr[] = { 26, 13, 23, 28, 27, 7, 25 }



ans = 26

```
ans=0
for( j=31 ; j≥0 ; j--)
{
    // jth bit
    cnt=0
    for( j=0 ; j<N ; j++)
    {   if( checkbit( a[j], j) == T)
            cnt++
    }
    if( cnt ≥ 2)
    {
        ans = ans + 2^j
        for( j=0 ; j<N ; j++)
        {   if(checkbit(a[j], j) == False)
            {   arr[j] = 0
            }
        }
    }
}
return ans
```

TC : O(32N) ~ O(N)

SC : O(1)

Break

10:40 - Wednesday

9PM

→ Bitwise operators &, |, ^, ~, <<, >>

→ Properties a & a = a

a & 0 = 0

→ Checkbit(N, i) (N >> i) & 1 == 1

→ All twice, except unique XOR (Array)

→ All thrice except unique

for each bit

{ cnt no of ele set

if(cnt % 3 == 1)

ans = ans + 2^i

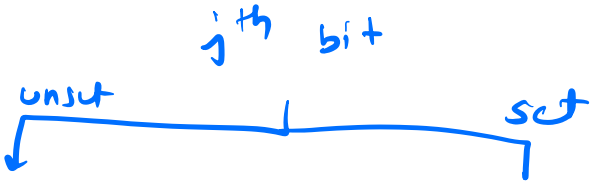
→ All twice except 2 unique

XOR(Array)

= XOR(A ∩ B) = K

Any set bit from XOR

j^{th} bit



D_1

P_1
P_1
P_3
P_3
P_4
P_4

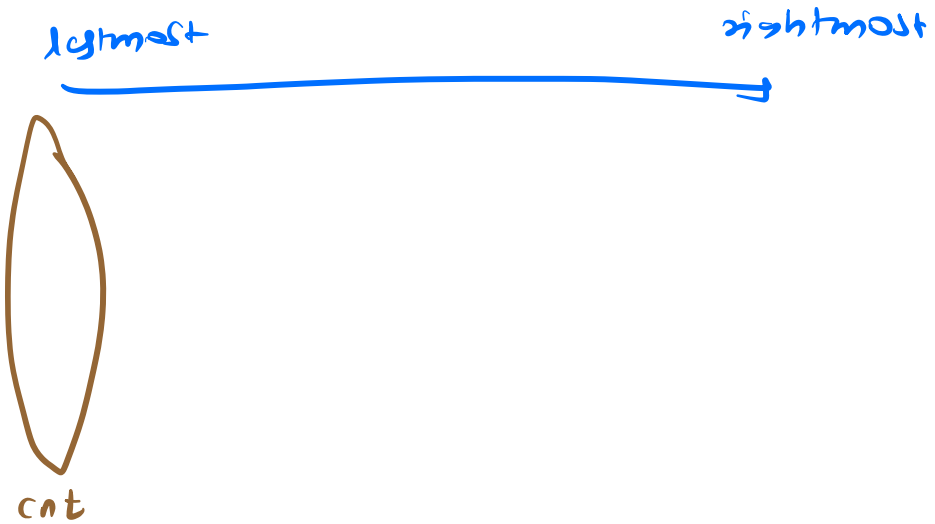
D_2

P_2
P_2
P_5
P_5

XOR(h_1) = ans1

XOR(h_2) = ans2

Q → Max AND of 2 numbers.



cnt ≥ 2

take bit in one ans

discard numbers with unset bit

ans