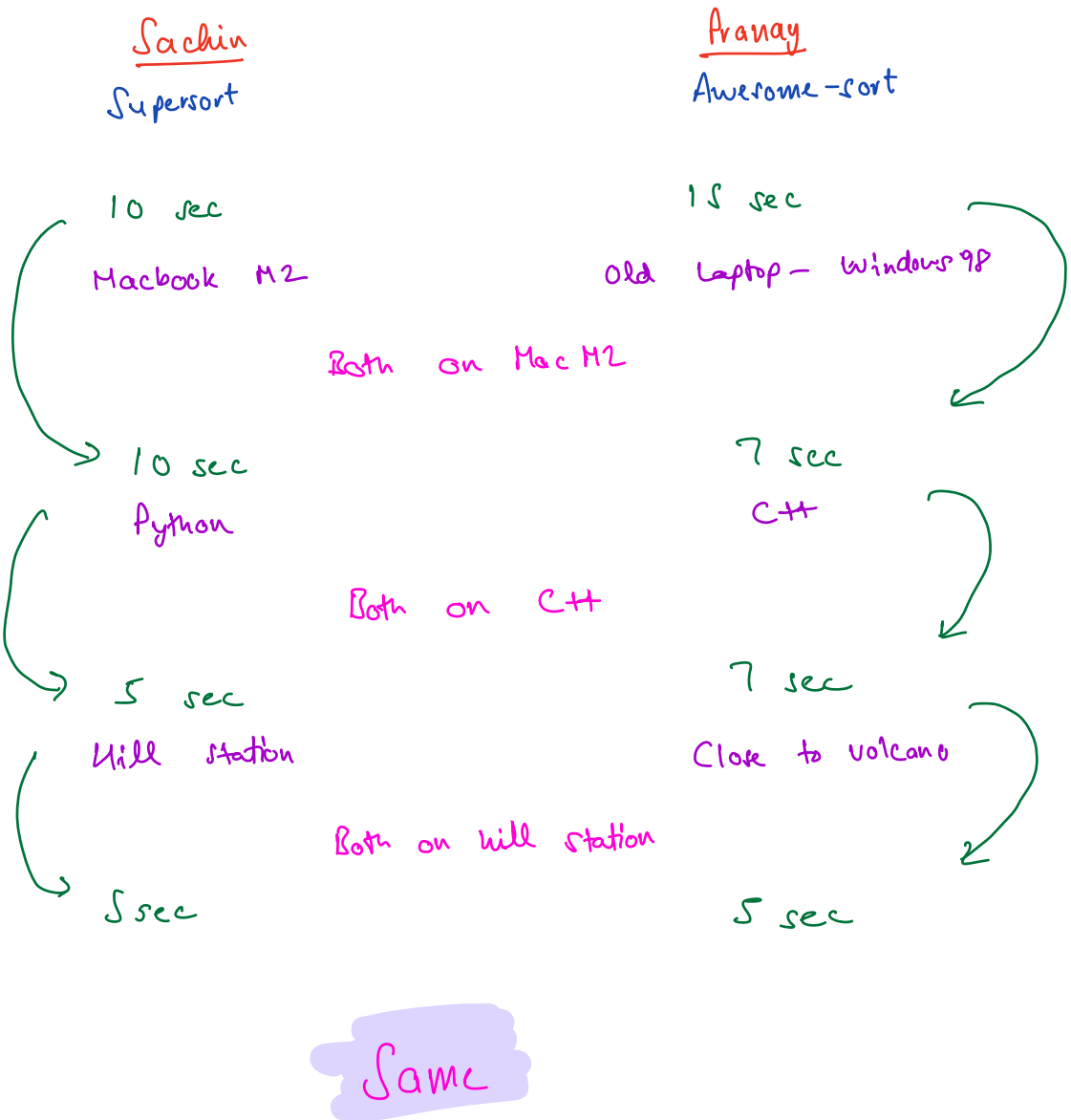
Time Complexity 2

AGENDA:

- Asymptotic Analysis
- Big O
- Issues with Big O
- Space complexity
- TLE

Comparing Algorithms

Given an array, sort it.



Execution Time – Not a reliable measure of performance for algo's

→ It depends on external factors

→ SW / HW / Environment

```
for (i=0; i<N; i++) {
    =
    =
}
```

→ N iterations

No of iterations is independant of external factors

Angeline: $100 \log_2 N$

Shubhankur: $N/10$

Before 3550: Shubhankur preferred

After 3550: Angeline preferred

Overall – Angeline's algo is more efficient

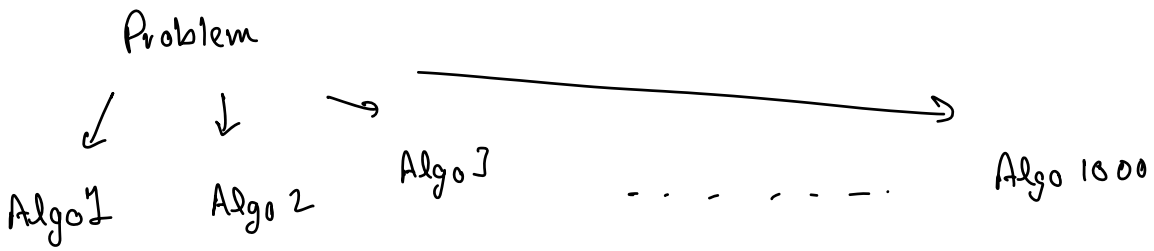
Hotstar – 2×10^8 people

Google – 10^8 search results in 0.012 s

Despacito – 7B+

Huge data

Increasing



Graph method is not possible to compare.

Asymptotic Analysis of Algorithms

Observing performance of algorithms for very large inputs.

↳ Big O

Comparisons using Big O notation

Angeline

$100 \log_2 N$

↓

$O(\log_2 N)$

Better

Snubhankur

$N/10$

↓

$O(N)$

1) Calculate the no of iterations
2) Neglect all lower order terms
3) Neglect constant coefficients

Why neglect lower order terms ?

Algo : $N^2 + 10N$ ← Lower order

$N = 100 \rightarrow 10^4 + 10^3$

% contribution of lower order term

$\frac{10^3}{10^4 + 10^3} \times 100 = 9.09\ \%$

$N = 10^4 \rightarrow 10^8 + 10^5$

$\frac{10^5}{10^8 + 10^5} \times 100 = 0.1\ \%$

Contribution of lower order terms is significantly small for larger inputs

Negligible

Why neglect constant coefficient ?

Manaswini	Sufyan
$10 \log_2 N$	N
$10^2 \log_2 N$	N
$10^3 \log_2 N$	N
$10^4 \log N$	$10N - 10^3$

Issues with Big-O

Issue 1

	Ajith	AbinDas	More efficient
	100N	N^2	
N = 50	100 x 50	50 x 50	AbinDas
N = 80	100 x 80	80 x 80	AbinDas
N = 100	100 x 100	100 x 100	Same
N = 120	100 x 120	120 x 120	Ajith
N = 150	100 x 150	150 x 150	Ajith

↓ $O(N)$

↓ $O(N^2)$

Big O doesn't give you complete picture.

Big O comparison is not true for all input sizes.

Issue 2

Vignesh

$N^2 + 10N$

↓

$O(N^2)$

Manikanta

$2N^2 + 5N$

↓

$O(N^2)$

Issue : Comparison says → Both are same

$N^2 + 10N$

$2N^2 + 5N$

Note: Big O will solve 99% of your problems

Linear Search

```
for (i=0; i<N; i++)
    if (arr[i] == x)
        return true
return false
```

Best - O(1)

Worst - O(N)

Default - Worst

Space Complexity

Break till 10:14 AM

Amount of extra space taken by your algorithm.

Extra space

Q1

```
func (int N) {
    int x = N
    int y = x + x
    int z = x + y
}
```

int - 4 bytes

3 x 4 = 12 bytes

Constant

↓

O(1)

Q2

Quiz 1

```
func (int N) {
    int x = N
    int y = x^2
    int z = x+y
    arr[] = int [N]      <- Array of integers
                 ^
                 Size of array
}
```

4B + 4B + 4B + 4N bytes

= (12 + 4N) bytes

↓

O(N)

Q3

Quiz 2

```
func (int N) {
    int x = N          - 4 bytes
    int y = x^2        - 4 bytes
    int z = x+y        - 4 bytes
    arr[] = int [N]    <- 1D Array of size N   <- 4N bytes
    l[][] = int [N][N] <- 2D Array of NxN      <- 4N^2 bytes
}
```

Total = (4 + 4 + 4 + 4N + 4N^2) bytes

↓

$O(N^2)$

Q4

Array of int ↓ Size of array ↓

```
sumOfArray ( int arr[], int N ) {
    sum = 0                                  - 4B
    for ( i=0 ; i<N ; i++ ) {                - 4B
        sum = sum + arr[i]
    }
    return sum
}
```

Total extra space
= (4 + 4) bytes
= 8 bytes
↓
O(1)

Q5

Quiz 3

What is the space complexity of the abc() function of the following code snippet, assuming the input matrix is of size N*M?

```java
class Main {

  static int[] abc(int[][] a) {
    int rows = a.length;
    int cols = a[0].length;
    int[] copy = new int[rows * cols];
    int k = 0;
    for (int i = 0; i < rows; i++) {
      for (int j = 0; j < cols; j++) {
        copy[k] = a[i][j];
        k++;
      }
    }
    return copy;
  }

  public static void main(String args[]) {
    int[][] a = { { 2, 7, 9 }, { 0, 1, 5 } };
    int[] b = abc(a);
    for (int x : b) {
      System.out.print(x + " ");
    }
  }
}
```

← N*M

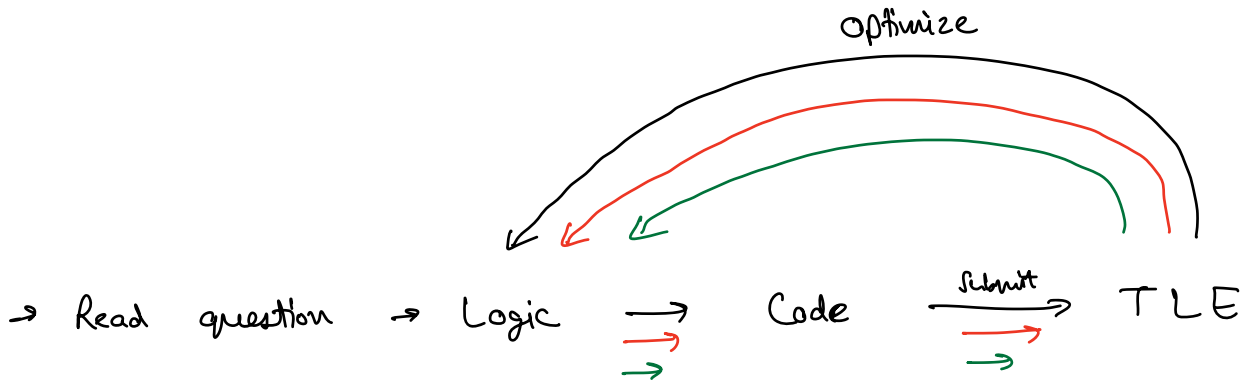
```python
def abc(a):
    rows = len(a)
    cols = len(a[0])

    copy = [0] * (rows * cols)
    k = 0
    for i in range(rows):
        for j in range(cols):
            copy[k] = a[i][j]
            k += 1

    return copy


def main():
    a = [[2, 7, 9], [0, 1, 5]]
    b = abc(a)
    print(b)


main()
```

← N*M

Extra space O(N*M)

TLE - Time Limit Exceeded

Google Contest → 2 hrs
3 Questions



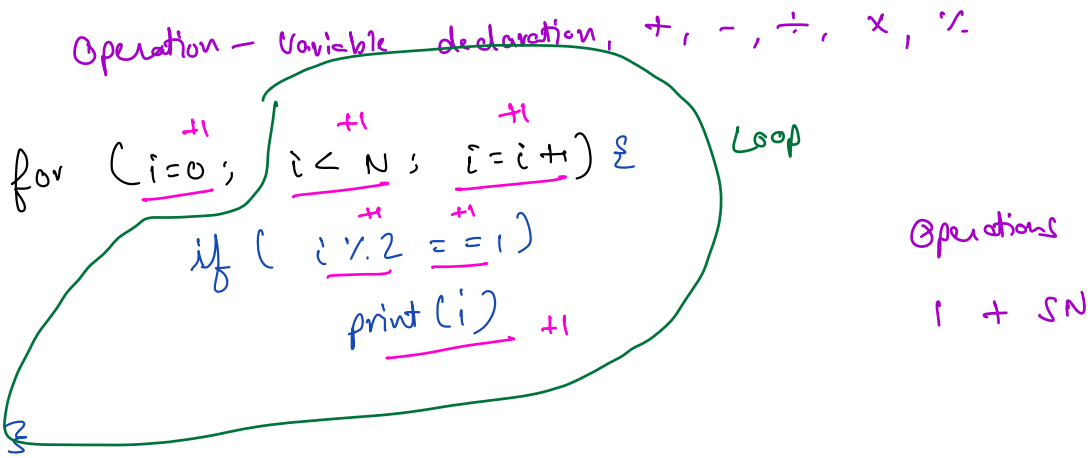
// Check whether TLE
without even writing a single line of code

Online Judges → Their servers → 1 GHz

Time Limit - 1 sec

10^9 operations / second

At max, our code can have
10^9 operations



Assumption

1 iteration → 10 operations

10^8 iterations → 10^9 operations

At max, our code can have

10^8 iterations

(Approx)

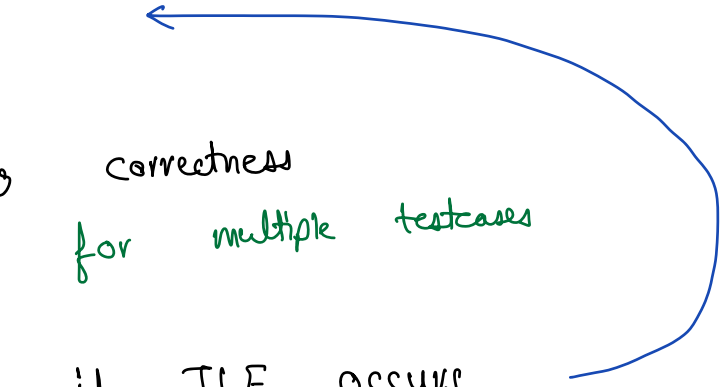
Process to solve

1) Read & understand the question

2) Logic

3) Checking correctness
   Try it for multiple testcases

4) Check if TLE occurs



↓ If no TLE

5) Write code & execute

Importance of constraints

Eg

$1 <= N <= 10^6$

$1 <= arr[i] <= 10^8$

Algo - $O(N^2)$ time

$N = 10^6 \rightarrow 10^{12}$ iterations

TLE

1 sec → 10^8 iterations

Eg

$1 <= N <= 5 \times 10^3$

Algo - $O(N^2)$ ← Not optimized

$N = 5 \times 10^3$

$N^2 \rightarrow (5 \times 10^3)^2 = 25 \times 10^6 < 10^8$

Will work

No TLE

Eg $1 <= N <= 5 \times 10^2$

Algo: $O(N^3)$ time

Rare

$N = 5 \times 10^2$

$N^3 = 125 \times 10^6$

$= 1.25 \times 10^8$ iterations

$\lesssim 10^8$ iterations

Can't be said

↙ It might work

→ It might not work

Corner Case

Doubts

Thank you

C++ → Compiled language

Python → Interpreted language

Python ⟶ C++ →

10^8 N

↑ Cannot be ignored

$1 <= N <= 10^8$

```
for i → [1,n]
    for j → [1, 3^i]
```

i	j	Iterations
1	$1 \to 3^1$	3^1 +
2	$1 \to 3^2$	3^2 +
3	$1 \to 3^3$	3^3 +
...		...
N	$1 \to 3^N$	3^N

Total = $3 + 3^2 + 3^3 + \cdots 3^N$

GP formula

10 operations → 1 iteration

1 operation → $\frac{1}{10}$ iteration

10^9 operations → $\frac{1}{10} \times 10^9$ iterations

$= 10^8$ iterations

$10000 \log x$ $\qquad$ $10x - 10000$

```
for (i=1; i<=n; i+=2)
```

$<=N$

i → 1, 3, 5, 7, 9, 11 .....

Approx - $\frac{N}{2}$ iterations $\rightarrow$ $O(N)$

$\frac{N}{2} = \frac{1}{2} \times N$

↑ Constant coefficient

```
for  i -> [0, n-1]
    for  j -> [0, i]
```

i	j	Iterations
0	0 → 0	1
1	0 → 1	2
2	0 → 2	3
3	0 → 3	4
⋮		
n-1		

Total =

$1 + 2 + 3 + 4 + \ldots + N$

$= \frac{N(N+1)}{2}$

$1 <= N <= 10^3$

$1 <= arr[i] <= 10^5$

$O(N^2)$

$O(\max(arr[i]))$

```
arr = []
for i → [0, n-1]
    arr.append(i)
```

} Space - O(N)

O, Θ, Ω, → Discrete mathematics

```
for (i=0; i<100; i++)
```

=

N = 100 → 100 iterations

N = 10^{25} → 100 iterations

Does no of iterations depend on N?

↳ No

Constant no of iterations = 100

Time - O(1)

main() {

Come from input

sum(N, M)

sum(a, b) {

return arr[a to b]

}

}

↓

Space - O(N+M)

Good Night

Thank You

Friday