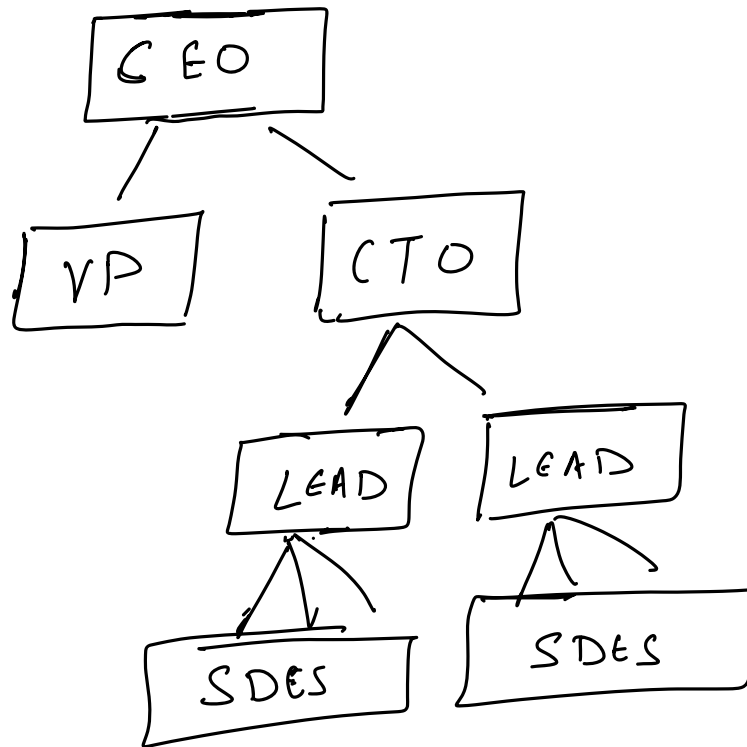
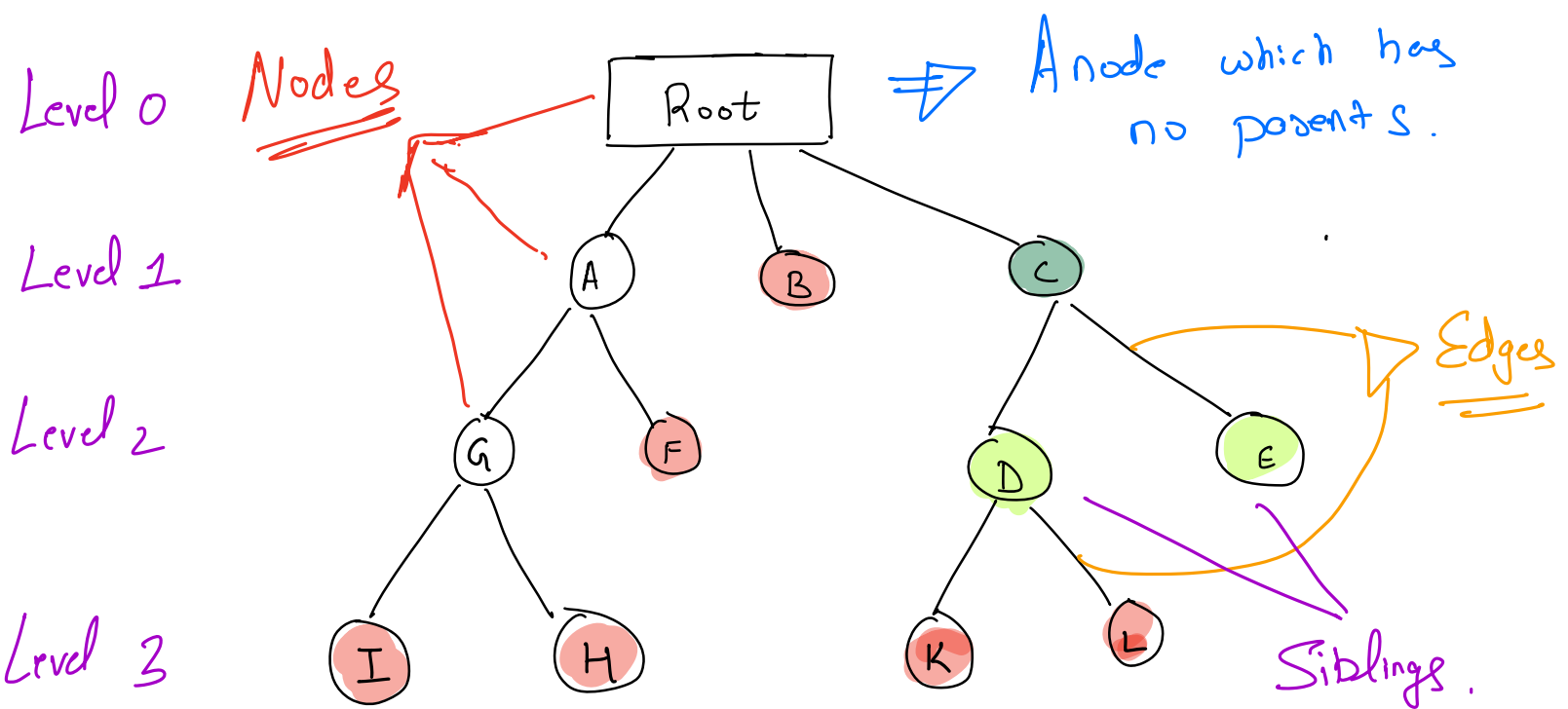


# Hierarchical Data



# TREES



1) A, B, C are children of Root.

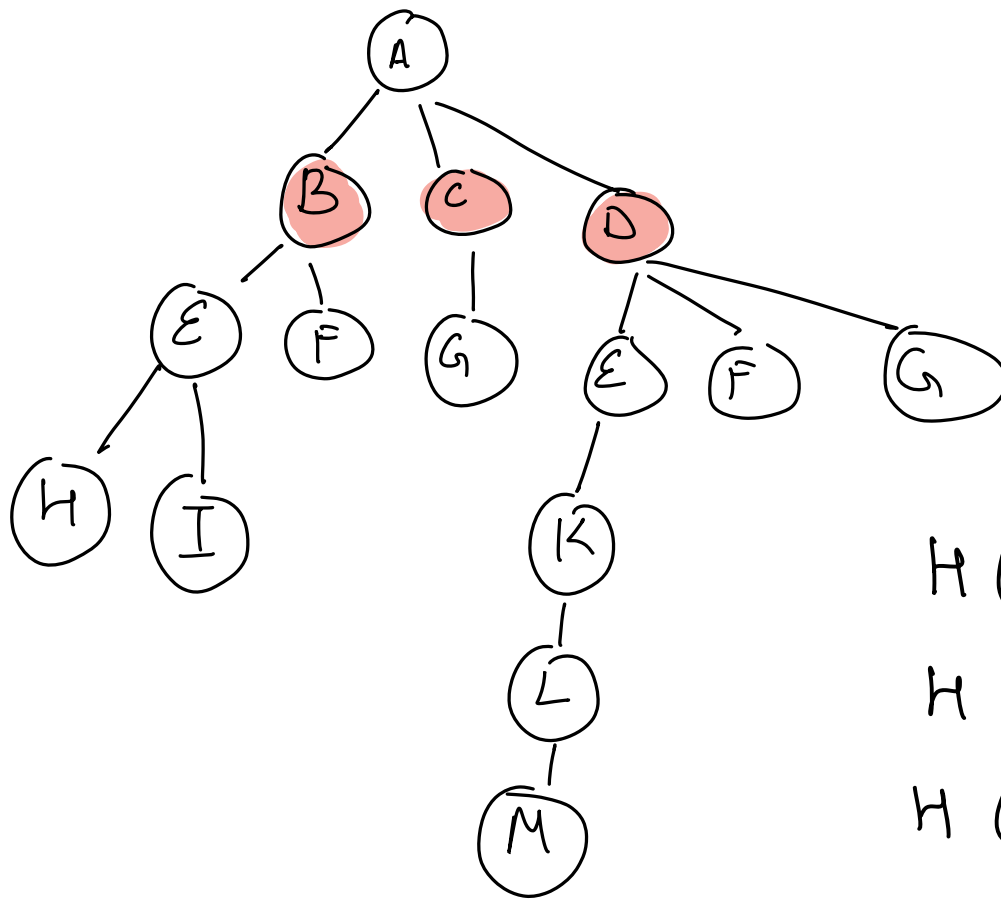
2) C is the parent of D & E.

3) D & E are siblings as they are children of the same parent.

4) G, F, D, E are at same level.

5) Nodes having no child are leaf nodes.

Height of a node : length of the longest path starting from that node to any leaf node. We will go only to higher levels. Length is calculated using edges traversed.



$$H(D) = 4$$

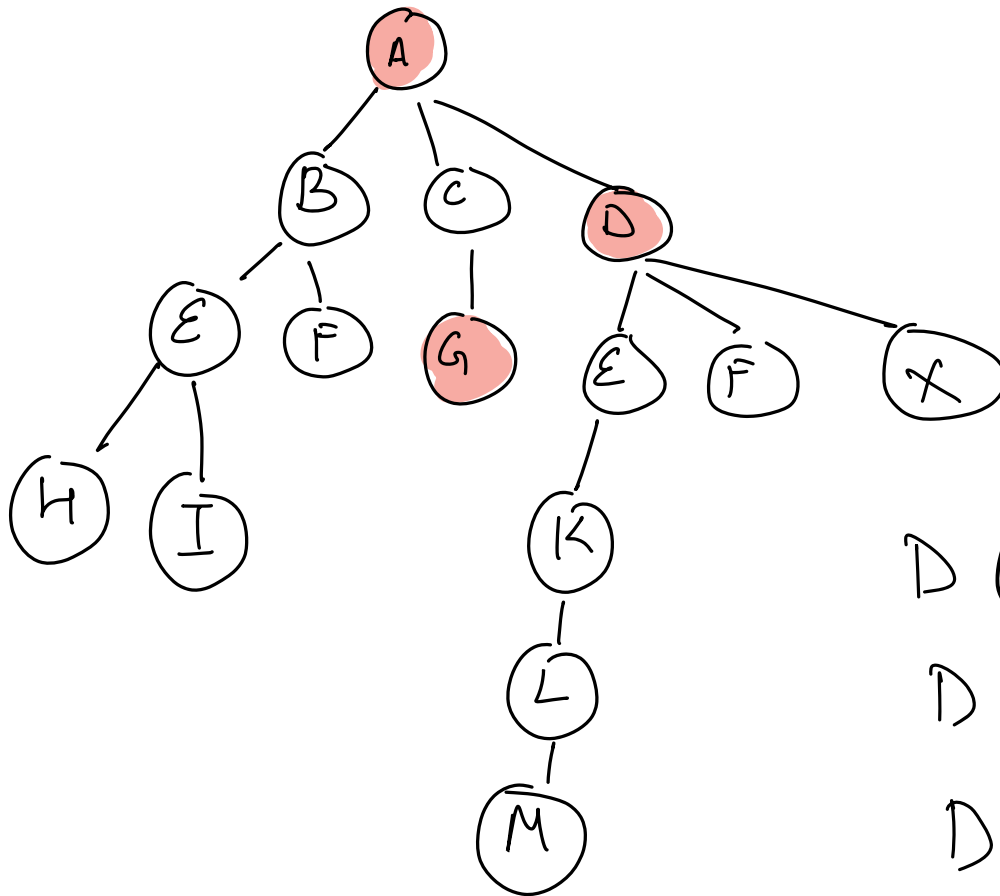
$$H(C) = 1$$

$$H(B) = 2$$

$$h(\text{node}) \Rightarrow \text{max-height (children)} + 1$$

Depth of node : The length of the path from that node to the root.

Level = Depth.



$$D(D) = 1$$

$$D(G) = 2$$

$$D(A) = 0$$

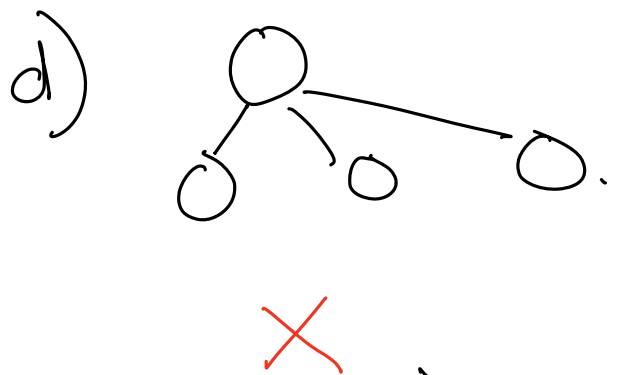
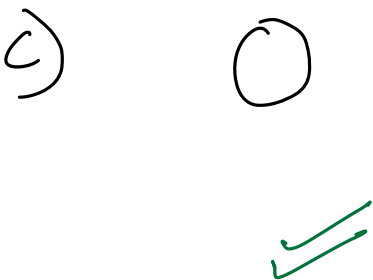
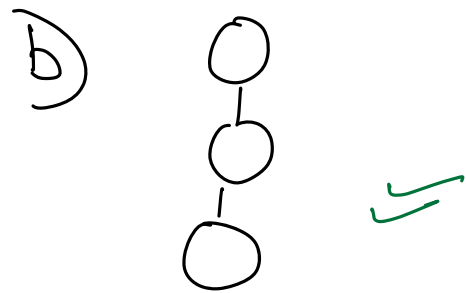
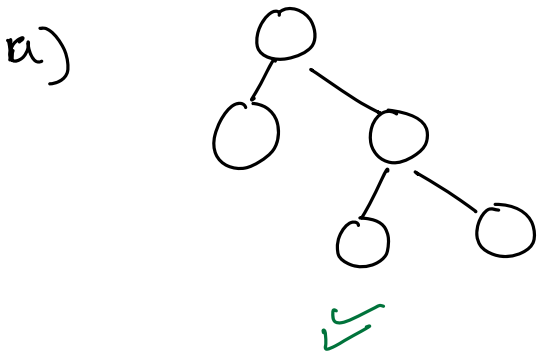
# Types of Trees:

1) n-ary tree: nodes can have any number of children.

2) Binary Tree: nodes can have at max 2 children.

A node can have 0, 1 or 2 children.

Which one is a Binary tree



class TreeNode {

int val;

TreeNode left;

TreeNode right;

TreeNode (int val) {

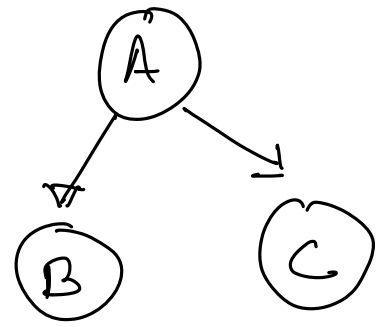
this.val = val;

this.left = null;

this.right = null;

}

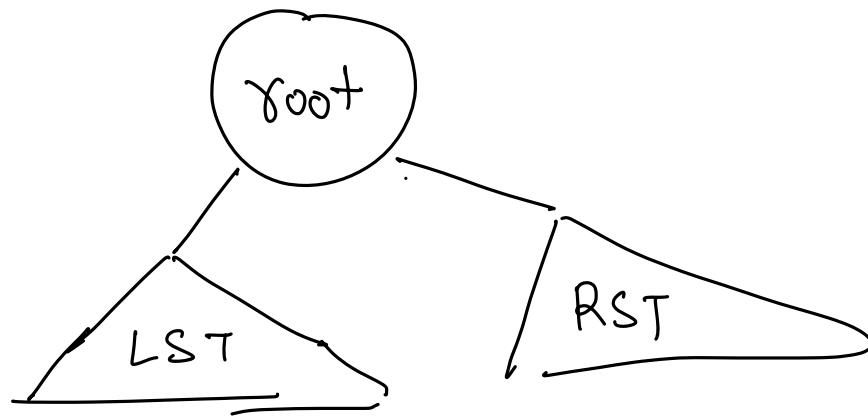
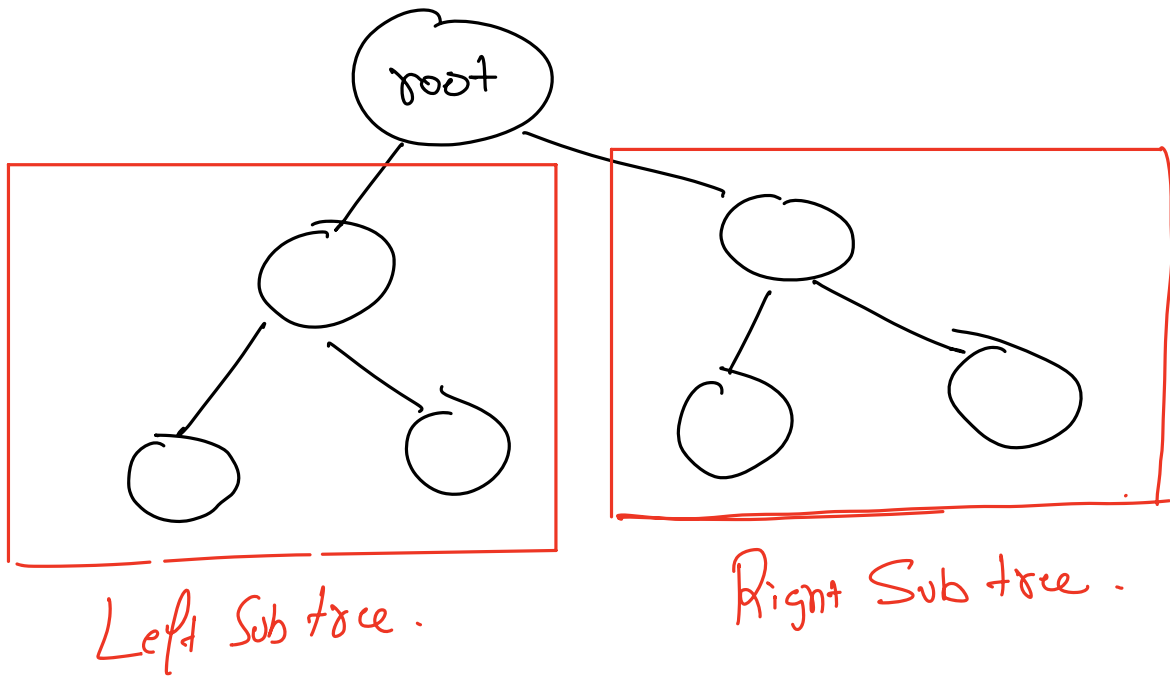
}



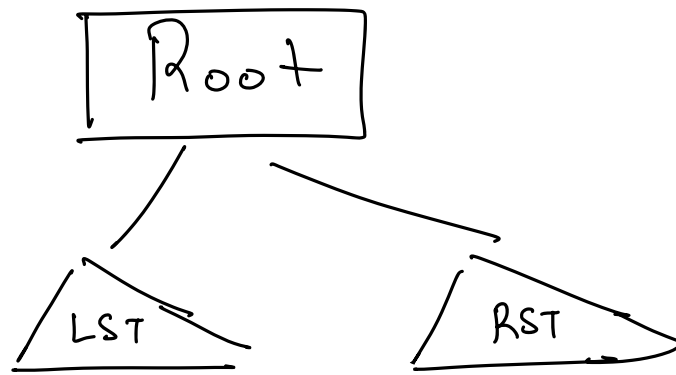
B: left child  
of A

C: right child  
of A.

# Binary Tree



Recursion will be used a lot to solve problems of trees because of the structure.



Possible ways of traversing.

Root LST RST	Root RST LST	LST Root RST	LST RST Root	RST LST Root	RST Root LST
--------------------	--------------------	--------------------	--------------------	--------------------	--------------------

# LST will always be traversed before RST!

Root LST RST	<del>Root</del> <del>RST</del> <del>LST</del>	LST Root RST	LST RST Root	<del>RST</del> <del>LST</del> <del>Root</del>	<del>RST</del> <del>Root</del> <del>LST</del>
--------------------	---	--------------------	--------------------	---	---

↓

Preorder

↓

Inorder

↓

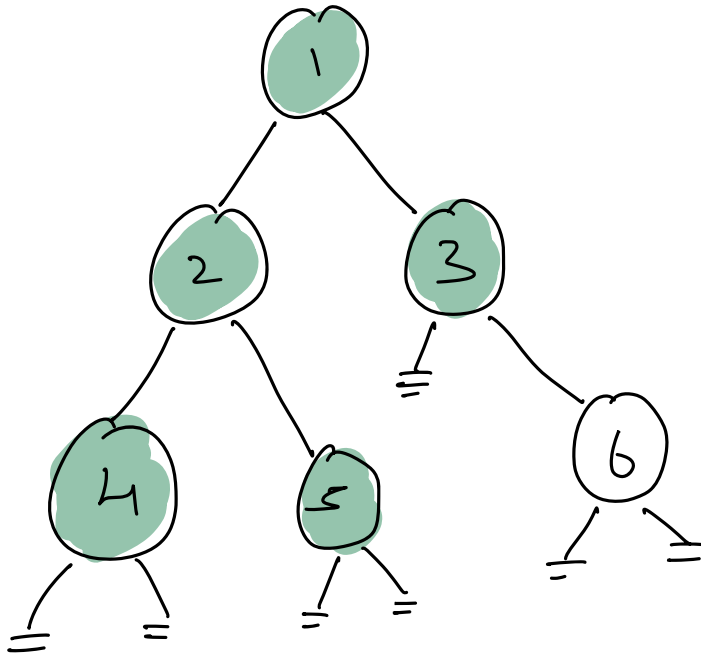
Postorder



Pre-order

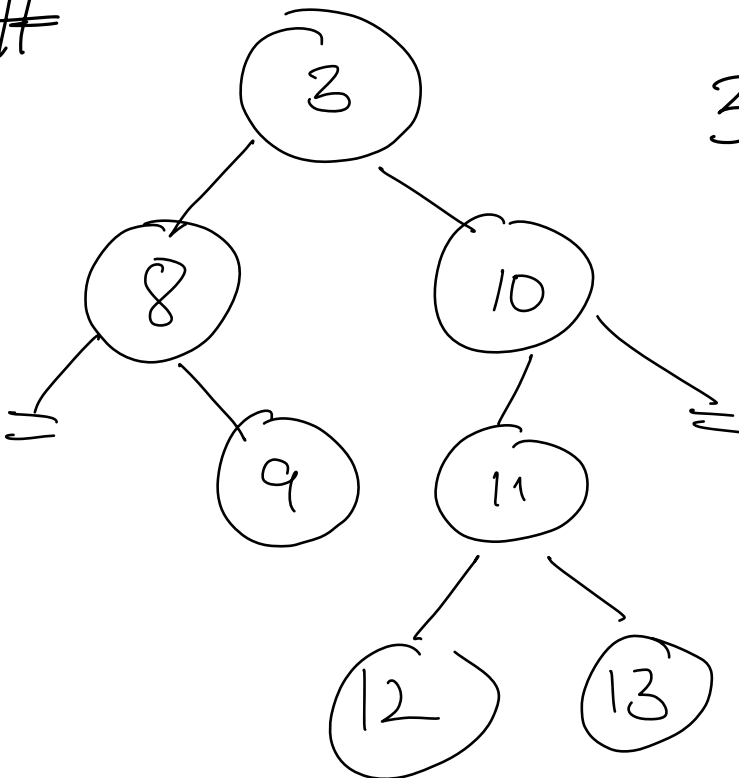


Root LST RST



1, 2, 4, 5, 3, 6

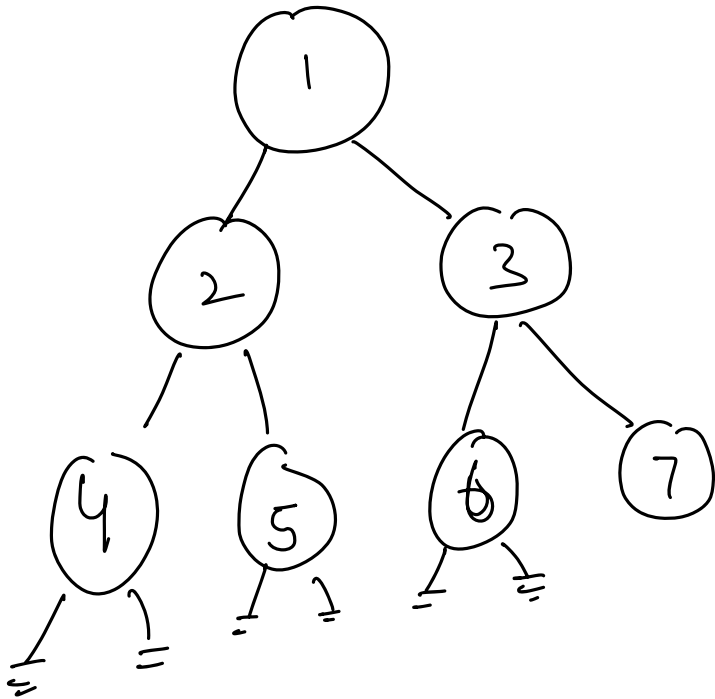
#



3, 8, 9, 10, 11, 12, 13

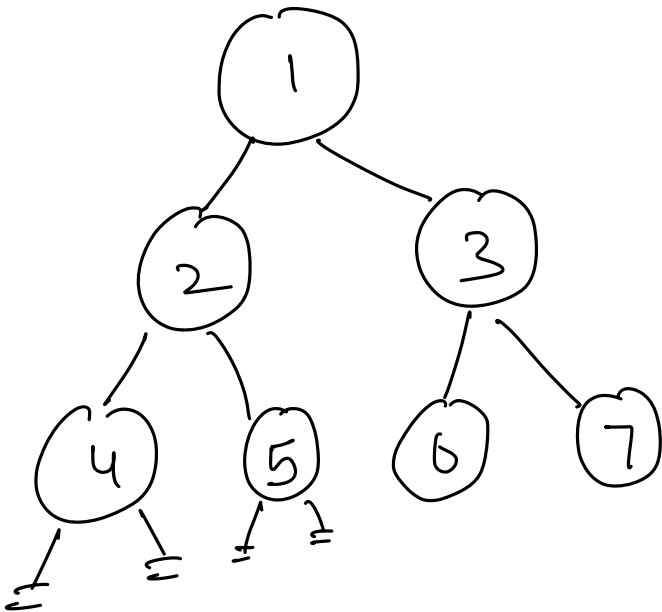
Inorder

$\Rightarrow$  LST, Root, RST



4, 2, 5, 1, 6, 3, 7

Post order : LST, RST, Root



4, 5, 2, 6, 7, 3, 1

void pre\_order (Node root) {

if (root == null) 1  
return;

print (root.val); 2

pre\_order (root.left); 3

pre\_order (root.right); 4.

}

1) Pre order = 1, 2, 3, 4

2) In order = 1, 3, 2, 4

3) Postorder = 1, 3, 4, 2

$T_c: O(n)$

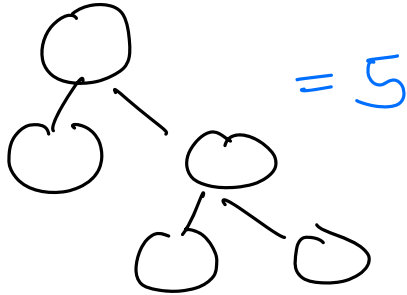
$SC: O(h)$

↳ height of tree.

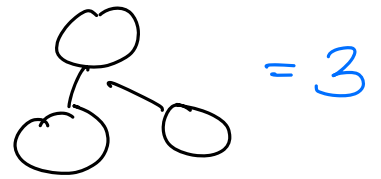
$SC: O(n)$ : Worst case.

Q1 Find the number of nodes in a tree.  
You will be given the root.

Ex1



Ex2



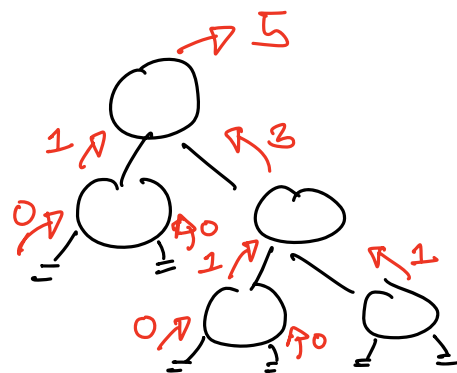
```
int count_nodes (Node root) {
    if (root == null)
        return 0
```

```
    return count_nodes (root.left)
        +
        count_nodes (root.right)
        +
        1
```

3

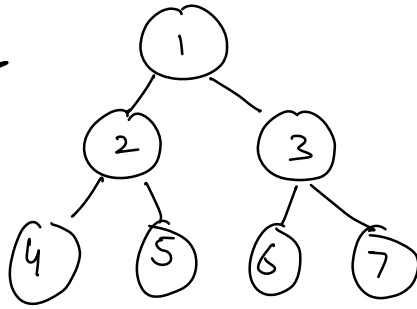
Tc:  $O(n)$

Sc:  $O(n)$



Q<sub>2</sub> Find the sum of all nodes in a tree.

Ex 1



⇒ 28

```
int count_nodes (Node root) {  
    if (root == null)  
        return 0
```

```
    return count_nodes (root.left) +  
           count_nodes (root.right) +  
           root.val;
```

3

Q3 Find height of a tree ToDo

↳ height (root)!

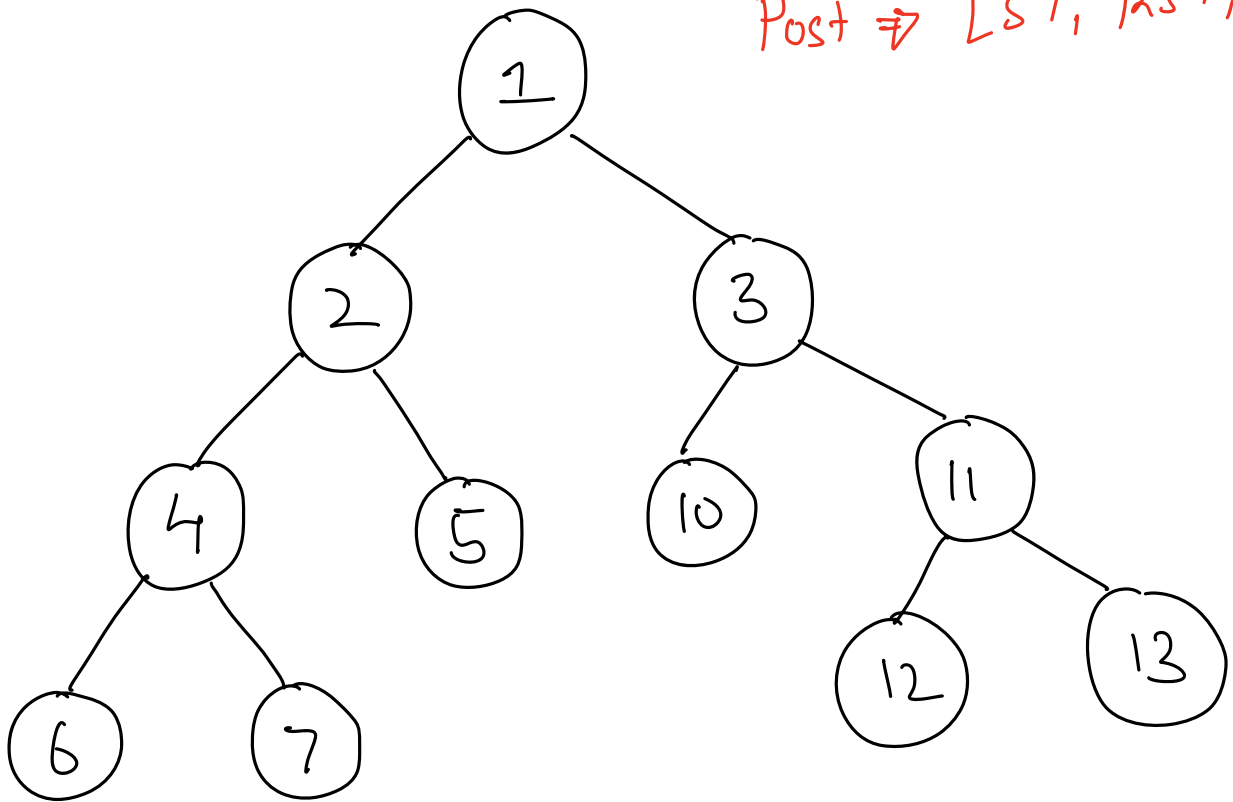
# Handle Base Case properly!

Q3 Given the inorder and postorder traversal of a tree,  
Construct the Binary tree. All values are distinct.

in[] = [ 6 4 7 2 5 1 10 3 12 11 13 ]

post[] = [ 6 7 4 5 2 10 12 13 11 3 1 ]

In  $\Rightarrow$  LST, Root, RST  
Post  $\Rightarrow$  LST, RST, Root



$\rightarrow$  Last Element of Post-order traversal  
will be root.

$in[] = [$ 

0	1	2	3	4
6	4	7	2	5

 $]$

$\uparrow$  index  
 $\nearrow$  index+1  
 $\nearrow e$

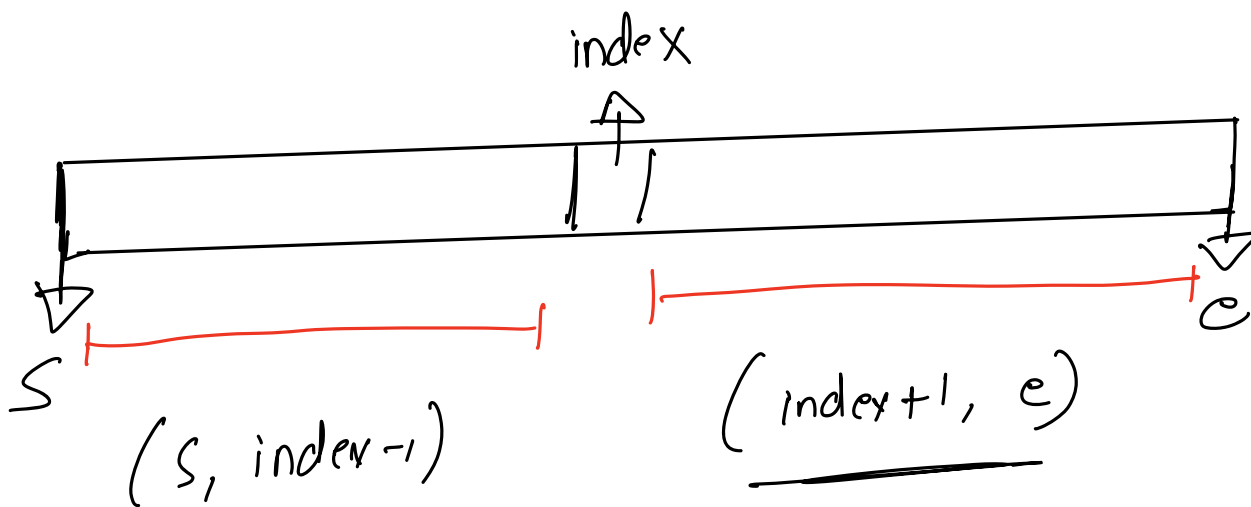
$post[] = [$ 

0	1	2	3	4
6	7	4	5	2

 $]$

$\uparrow$  index  
 $\nearrow$  index+1  
 $\nearrow e$

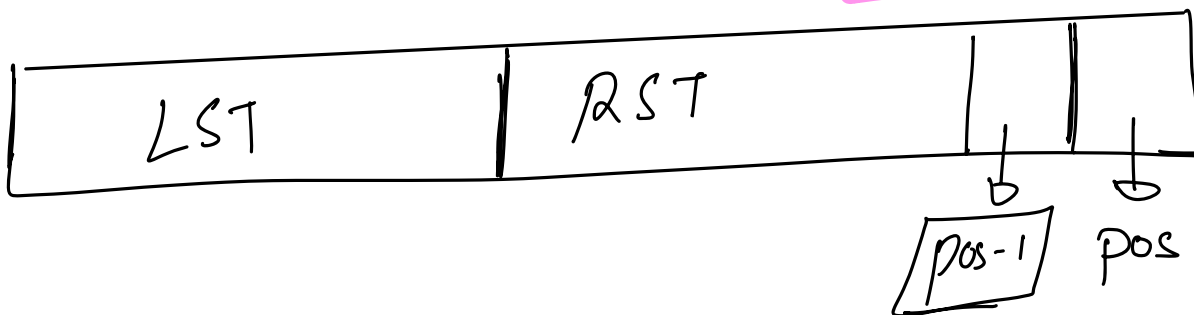
$\uparrow$  pos



$$[index+1, e]$$

$$= e - (index+1) + 1$$

$$= e - index$$



initial call : constructTree (0, n-1, n-1)



# Pseudo Code !

in[], post[]; map <int, int> mp;

Node constructTree (int s, int e, int pos) {

int root-val = post[pos]

Node root = new Node (root-val);

int index = mp[root-val];

root->left = constructTree (s, index-1, pos-(e-index)-1);

root->right = constructTree (index+1, e, pos-1);

return root;

}

Base Case: if (s == e) {

Node root = new Node (in[s])

return root;

}

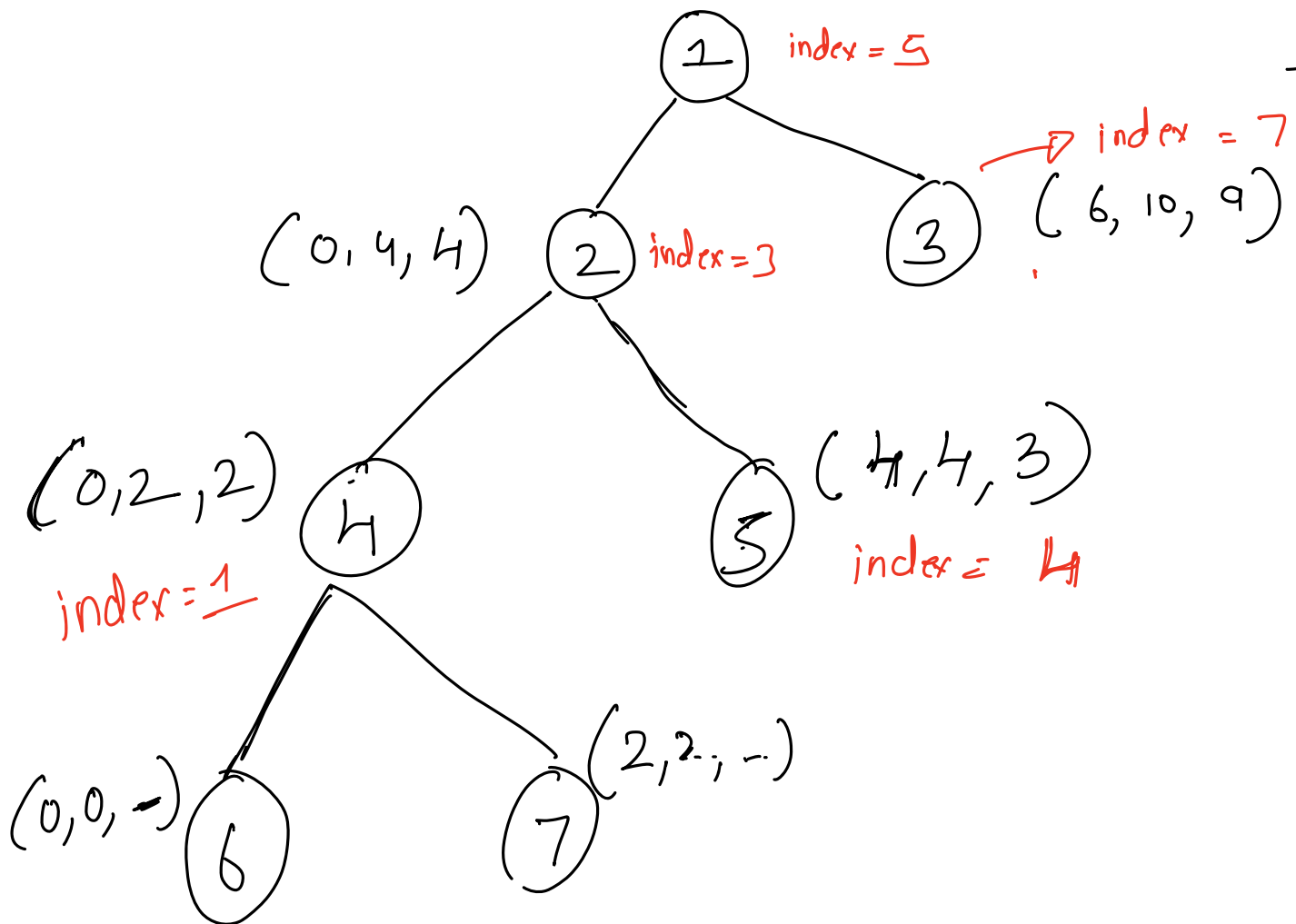
if (s > e)

return null;

in[] = [ 0 1 2 3 4 5 6 7 8 9 10  
6 4 7 2 5 1 10 3 12 11 13 ]

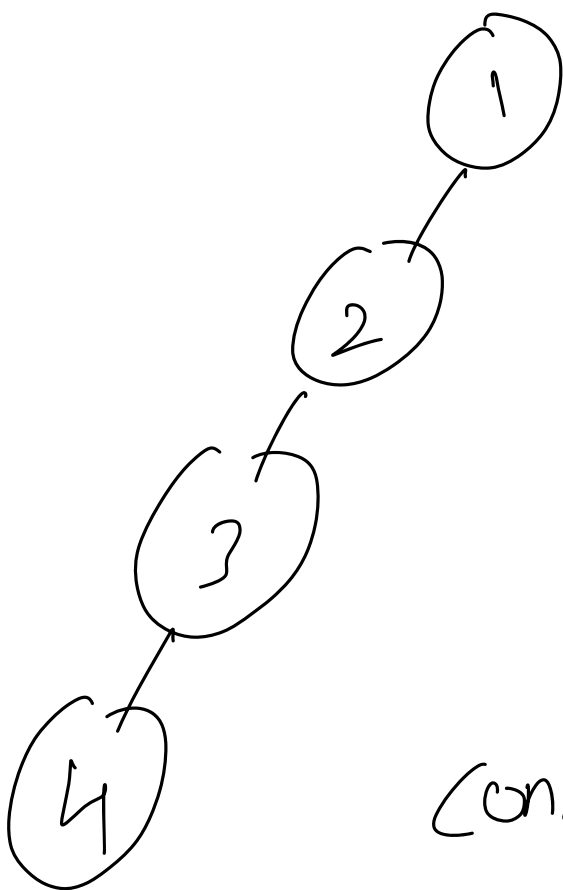
post[] = [ 0 1 2 3 4 5 6 7 8 9 10  
6 7 4 5 2 10 12 13 11 3 1 ]

Construct (0, 10, 10)



$T.C: O(n)$

$S.C: O(n)$  — Hashmap  
 — Stack  
 Space.



in  $\Rightarrow$  4 3 2 1

post  $\Rightarrow$  4 3 2 1

construct (0, 3, 3)

post (0, 2)

post (4, 3)

S7e