# Knowledge and Retrieval Modeling Hierarchies

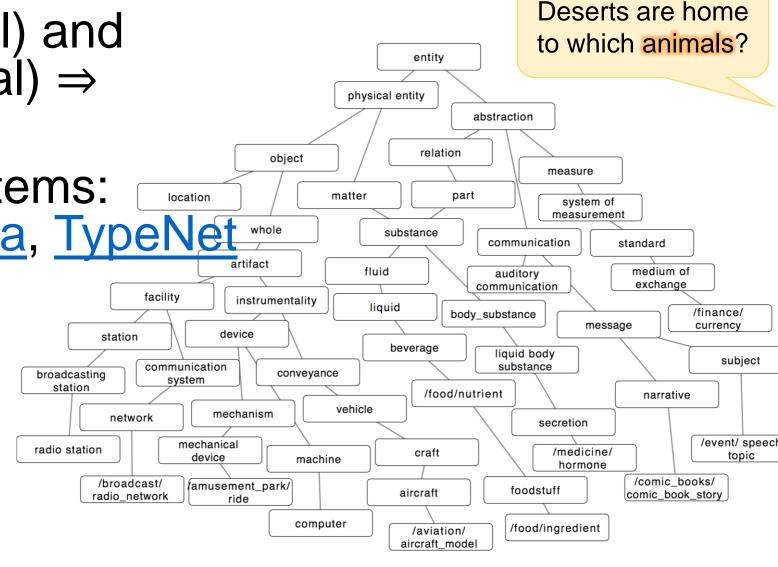
#### Hierarchies in KGs

Anti/symmetry ... how about transitivity?

 (camel, is-a, mammal) and (mammal, is-a, animal) ⇒ (camel, is-a, animal)

 Popular KG type systems: FIGER, GFT, DbPedia, TypeNet

 None of the KG embedding methods discussed thus far handle hierarchies



## Two views of embedding hierarchies

- Embed a hierarchy (DAG) in a space with a distance such that distances on DAG are approximately preserved by distance in the embedding space
  - Euclidean → Poincaré balls, hyperbolic embeddings
- Encode DAG nodes so that ancestor-descendant queries can be answered efficiently
  - Gaussian, order and box embeddings
- Should work with incomplete supervision
- Should play well with other embeddings

## Low-distortion graph embeddings

- Each node v in graph G = (V, E) embedded to  $x(v) \in \mathbb{R}^D$
- Graph distance between  $u, v \in V$  is  $d_G(u, v)$
- Distortion of embedding g is given by

$$\operatorname{distor}(x,G) = \frac{\max_{u,v} \frac{\|x(u) - x(v)\|}{d_G(u,v)}}{\min_{u,v} \frac{\|x(u) - x(v)\|}{d_G(u,v)}}$$
Maximum stretch
Minimum stretch

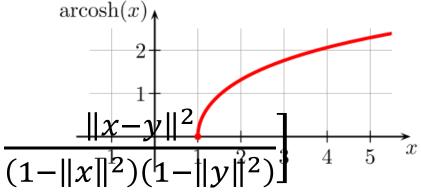
• Distortion of G is  $\inf_{x}$  distor(g, G)

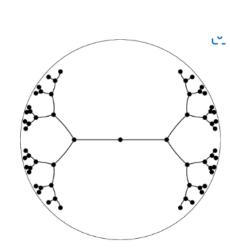
### **Euclidean distortion facts**

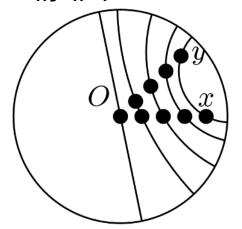
- With large enough D, any graph with n nodes can be embedded in  $\mathbb{R}^D$  with  $O(\log n)$  distortion
- Any connected planar graph can be embedded in  $\mathbb{R}^2$  with O(n) distortion; trees with  $O(\sqrt{n})$
- With large enough D, any tree can be embedded in  $\mathbb{R}^D$  with  $O(\log \log n)$  distortion
- Binary trees can be embedded in a line with  $O\left(\frac{n}{\log n}\right)$  distortion
- Binary trees can be embedded in  $\mathbb{R}^D$  with  $O\left(\frac{n^{1/D}}{\log n}\right)$  distortion
- Distortion of  $(1 + \epsilon)$  is possible in hyperbolic space

# Poincaré disk, hyperbolic space

- $\operatorname{acosh}(a) = \ln(a + \sqrt{a^2 + 1})$
- Points *x*, *y* strictly inside unit circle
- Hyperbolic distance  $d_H(x,y) = \operatorname{acosh} \left[1 + 2 \frac{\|x-y\|}{(1-\|x\|^2)(1)}\right]$
- As x, y approach perimeter  $d_H \to \infty$
- Natural tree embedding
- Precision bits not free!

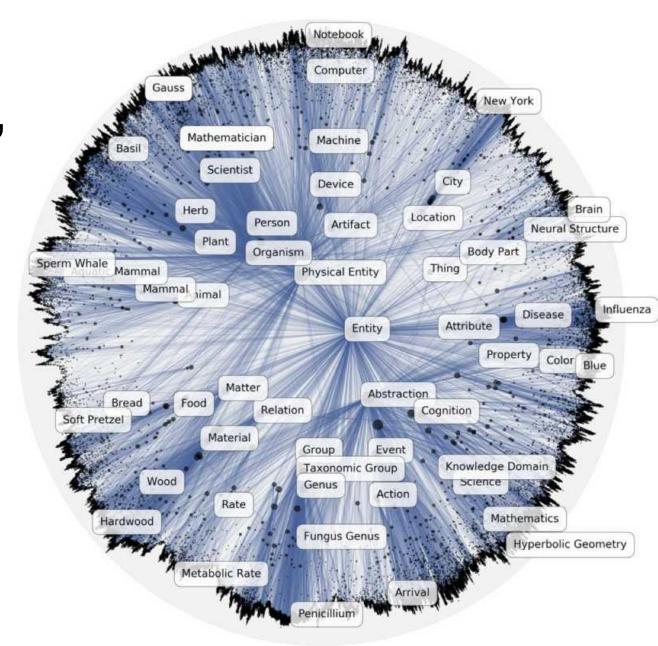






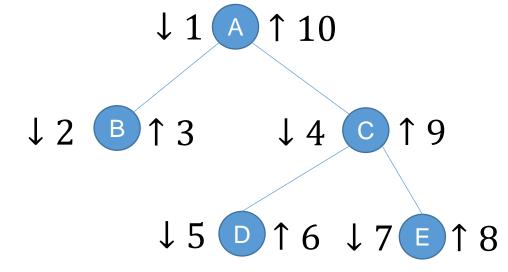
#### WordNet nouns in Poincaré disk

- As expected, generic synsets near disk center, specific near periphery
- What do applications (e.g., QA) need?
  - "scientists who played musical instruments"
  - "mammals living in the desert"
- word2vec ↔ Poincaré?



## Toward order embeddings

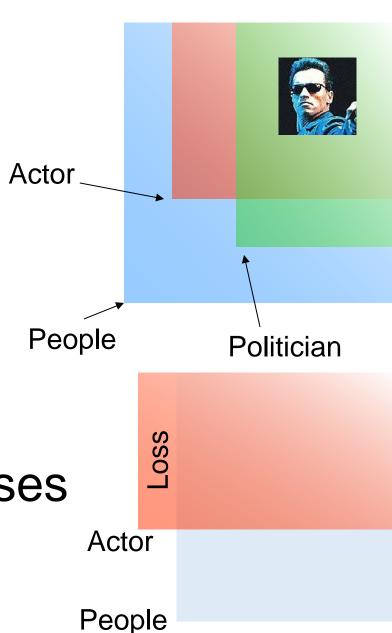
- Computing Euclidean or hyperbolic distance between items not the only choice
- Denote partial order by "x < y" meaning x is a descendant of y
- In case of a tree, associate with each item x the in-order traversal interval I(x)
  - I(B) = [2,3], I(A) = [1,10],I(D) = [5,6]
  - $I(x) \subset I(y) \Leftrightarrow x < y$



Nodes of a tree can be embedded in one dimension to answer ancestordescendant queries in constant time.

## Apex of axis-aligned open cones

- Item x represented by  $u_x \in \mathbb{R}^D$
- $u_x$  is the apex of an open cone
- $x \prec y \Leftrightarrow u_x \geq u_y$ , elementwise
- Design training loss function
  - Notation:  $ReLU(a) = [a]_+ = max\{0, a\}$
  - If x < y,  $\| \text{ReLU}(u_y u_x) \|$ , i.e. all D dims must satisfy constraint
  - If  $x \not\prec y$ , ReLU $[\alpha \|\text{ReLU}(u_y u_x)\|]$
- Does not recognize asymmetry in losses



## Addressing loss asymmetry

- If x < y, then for all dim d, want  $u_x[d] \ge u_y[d]$ 
  - E.g.,  $\ell_+(x,y) = \max_{d \in [D]} \text{ReLU}(u_y[d] u_x[d])$
- If  $x \not\prec y$ , then for some d, want  $u_x[d] < u_y[d]$ 
  - $\ell_{-}(x,y) = \min_{d \in [D]} \operatorname{ReLU}\left(\alpha \operatorname{ReLU}\left(u_{y}[d] u_{x}[d]\right)\right)$
- All open cones and their intersections have same measure of volume (unlike in-order intervals)
  - ... even though Politicians \( \beta \) People
- Hard to model negative correlation
  - X is-a fruit, or X is-not-a scientist

# (Hyper)rectangle/box embeddings

- Each type/item x characterized by interval  $I_x[d] = [b_{x,d}, h_{x,d}]$  for each dimension d
- Want  $I_x[d] \subseteq I_y[d] \ \forall d$ , iff x < y
- Learning to lay out Venn diagrams

