

## Recursion

20 August 2023 13:02

- function calls itself.
- When sol. of a problem depends on sol. of smaller instances of same problem.
  - (i) Base Case
  - (ii) Recursive Case

Trees  
Graphs  
DP

$$5! = 120 \quad (5 \times 4 \times 3 \times 2 \times 1)$$
$$(5 \times 4!)$$

for( i=n; i>=1; i-1 )  
} = loop

$$n! = n \times (n-1)!$$

$$5! = 5 \times 4!$$

4x3!

$$3 \times 21$$

2 x 1!

$$1 \times 6!$$

$$n = 5$$

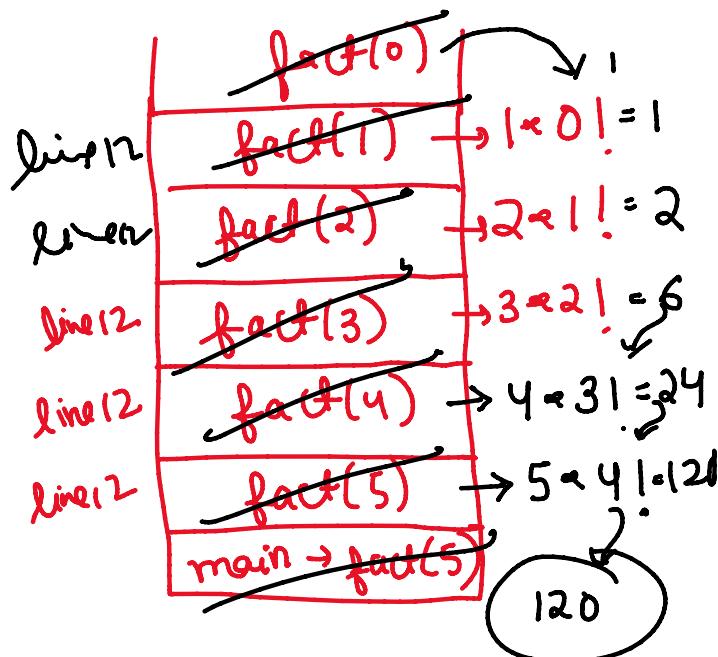
Base Case

if( $n == 0$ )  
return 1;

Base Case  
 $r = 0$

① Figure out the base off.

① figure out:  
② assume that subproblems will be solved automatically



.....

- ② Assume that subproblems will be solved automatically
- ③ Using subproblems write ans for current problem



Time  $\rightarrow O(N)$

Space  $\rightarrow O(N)$

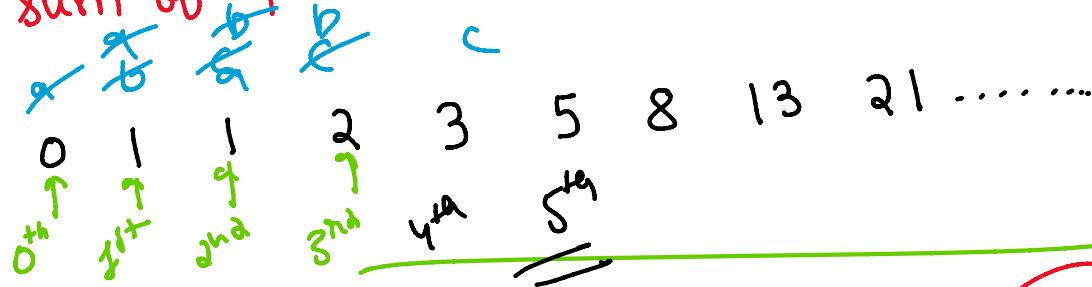
$O(1)$

## Fibonacci series

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Find  $n^{\text{th}}$  Fibonacci number.

→ sum of previous two numbers.



$$\begin{aligned} n &= 6 \\ \underline{\underline{\text{ans}}} &= 8 \end{aligned}$$

Base Case

$$0 = a + b$$

$\left. \begin{array}{l} \text{if } (n == 0 \text{ || } n == 1) \\ \text{return } n; \end{array} \right\}$

$\text{if } (n < 1)$

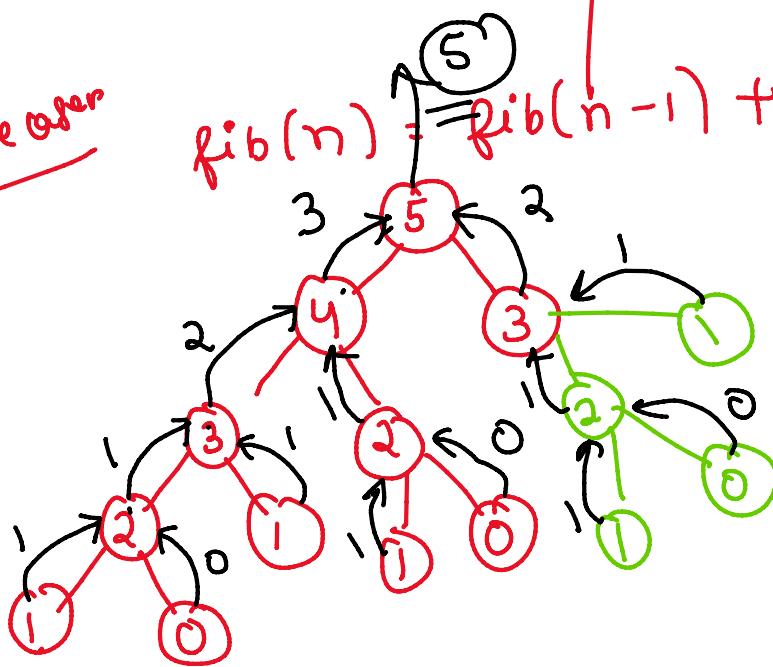
$$n = 5$$

$$\underline{\underline{\text{ans}}} = 5$$

return n;  $\xrightarrow{\text{array}} O(n) \rightarrow \text{space}$

Recursive Case

$$\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$$



$$\begin{aligned} 1 \\ 2 \\ 4 = 8 \\ 16 \end{aligned}$$

Time:  $O(2^N)$   
Space:  $O(N)$

Left to right direction.

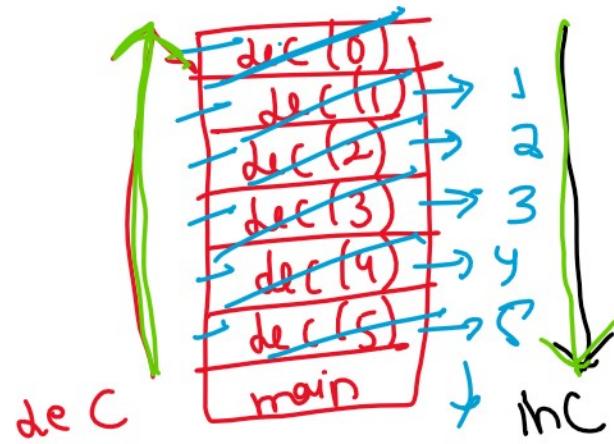
## Inc dec order

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n = 5

dec↑ 5, 4, 3, 2, 1

5, 4, 3, 2, 1

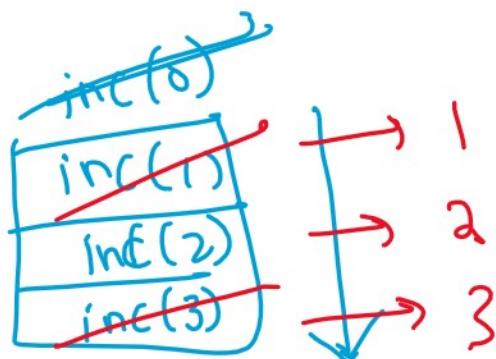


Incr 1, 2, 3, 4, 5

Time  $O(N)$   
Space  $O(N)$

Bad code  
if ( $n == 0$ )  
return;

```
void dec (int n)
{
    cout << n; // dec
    dec(n-1);
    cout << endl; // inc
}
```



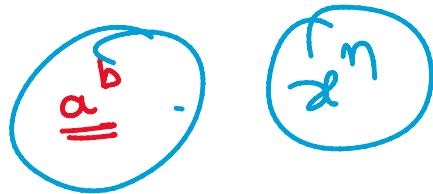
1, 2, 3

$a, b$ 

$$a = 5, b = 3$$

$$\text{ans} = 5^3 = 125$$

Base case    ~~if ( $b == 0$ )  
return 1;~~



$$5^0 = 1$$

$$10^0 = 1$$

if ( $b == 1$ )  
return a;

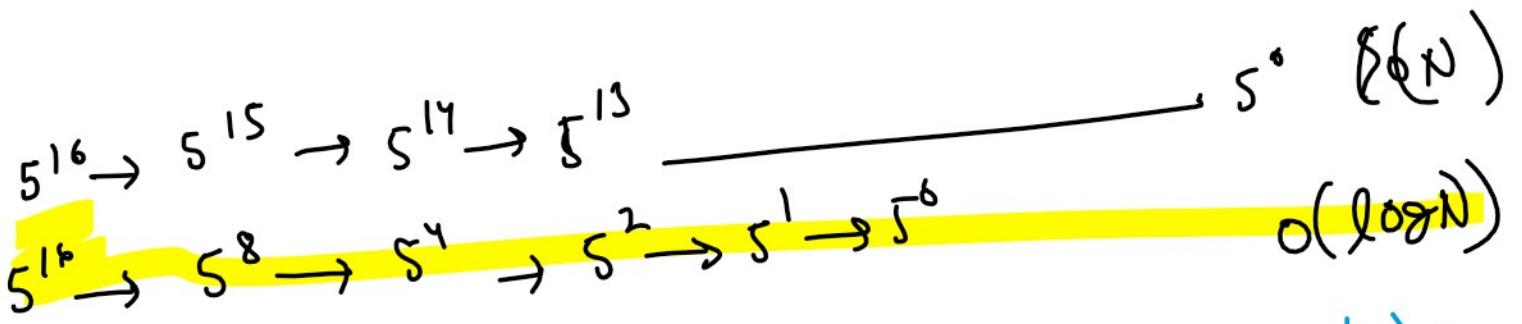
recursive case     $\text{power}(a, b) = a * \text{power}(a, b-1);$

$$\begin{aligned}
 5^3 &= 5 * 5^2 \\
 125 &\quad \downarrow \\
 5 * 5^2 &= 5 * 5 * 5^1 \\
 125 &\quad \downarrow \\
 5 * 5 * 5^1 &= 5 * 5 * 1 \\
 125 &\quad \downarrow \\
 5 * 5 * 1 &= 1
 \end{aligned}$$

Time =  $O(N)$   
Space =  $O(N)$

odd/even

power (2, 5)

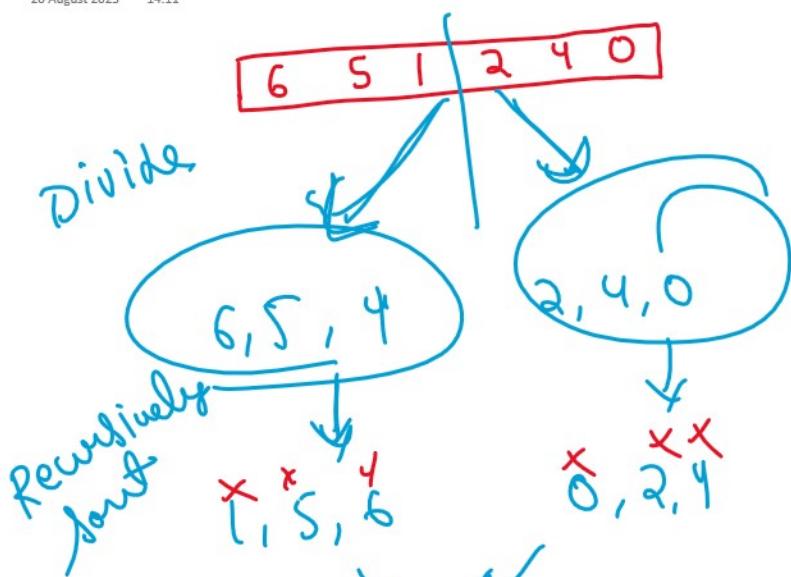


even  $b \cdot j \cdot 2 = 0$   
 $\text{return power}(a, b/2) * \text{power}(a, b/2);$

odd  $\text{return } a * \text{power}(a, b/2) * \text{power}(a, b/2);$

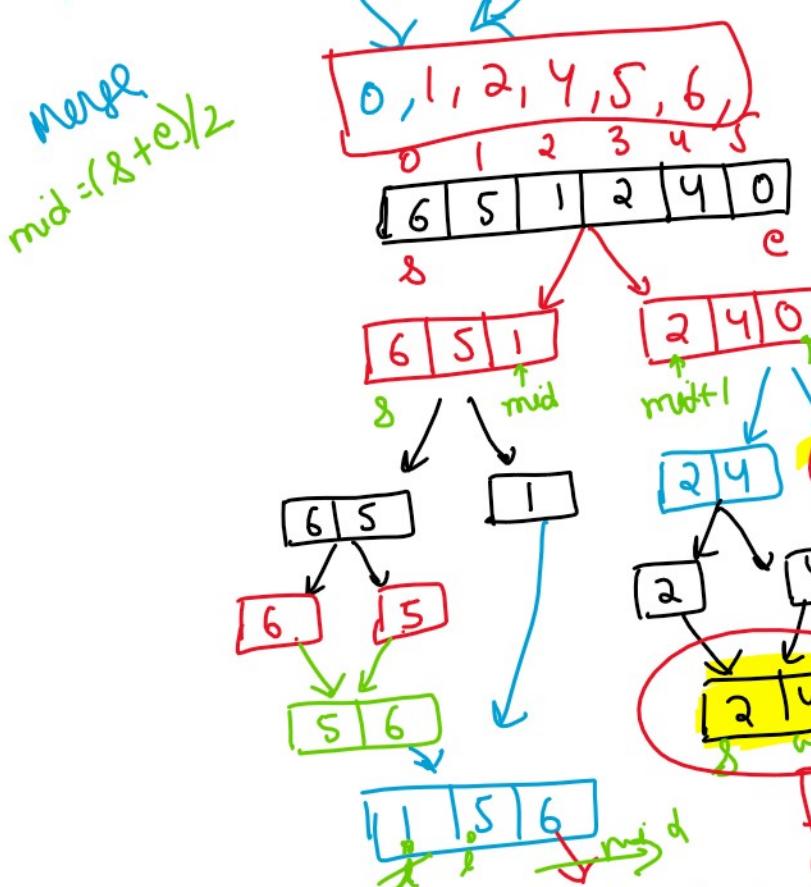
$5^{16} \rightarrow 5^8 \rightarrow 5^4 \rightarrow 5^2 \rightarrow 5^1 \rightarrow 5^0$

ans = 0, 1, 2, 4, 5, 6

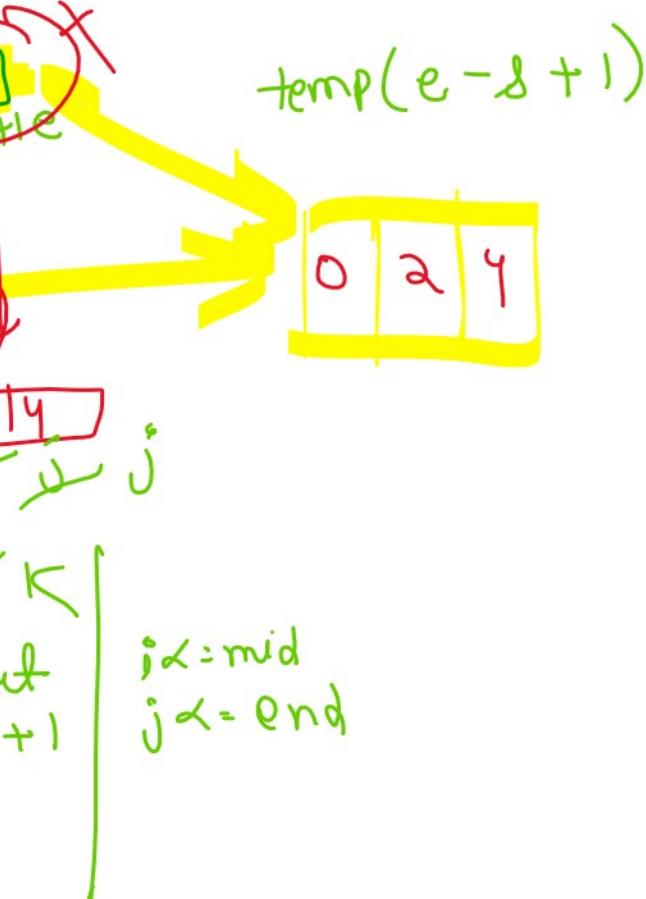


✓ merge Sort  
✗ quick Sort

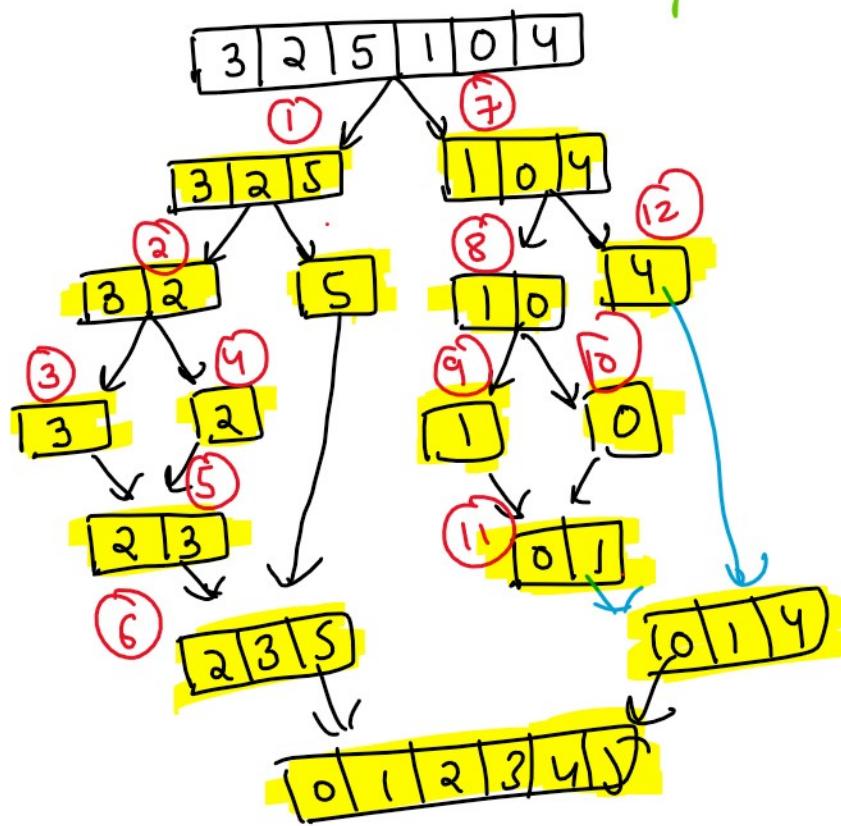
1. Divide
2. Recursively sort
3. Merge.



mid:  $(\delta + e)/2$   
left half  
right half

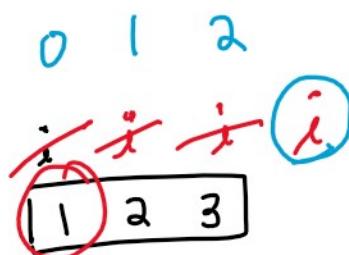
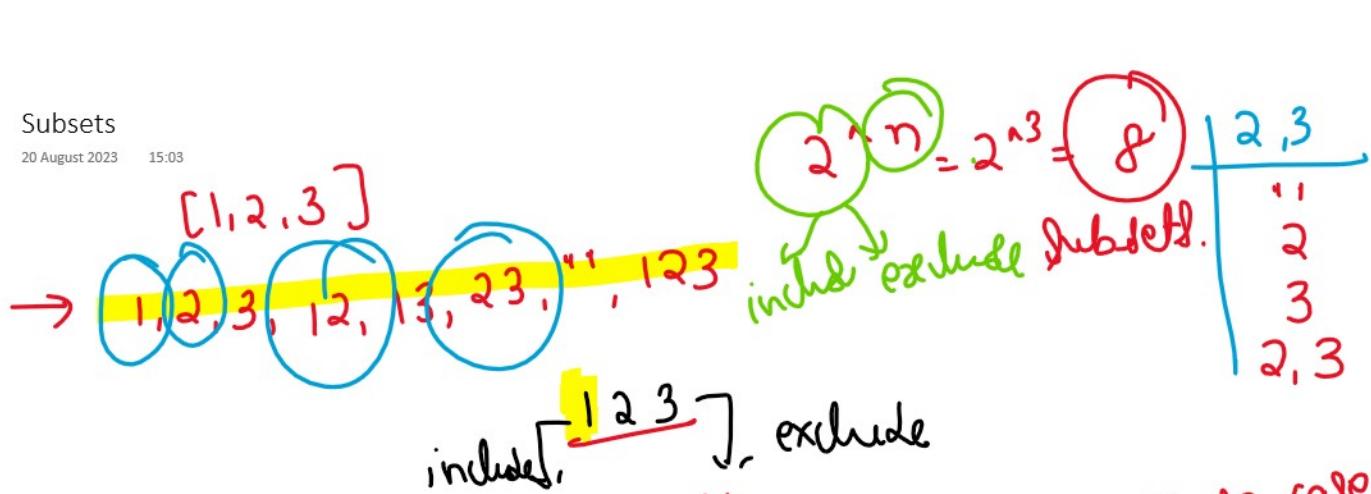


9.8



# Subsets

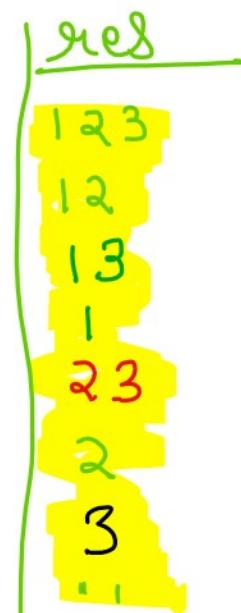
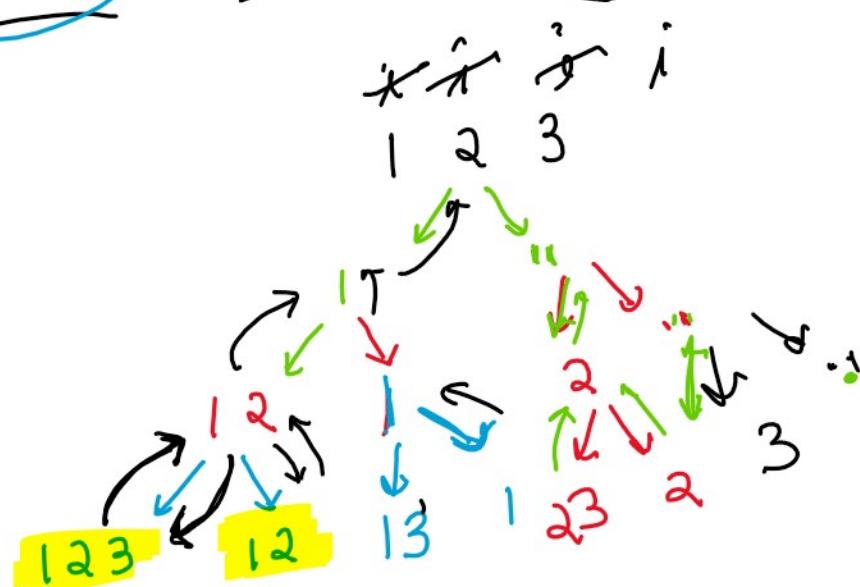
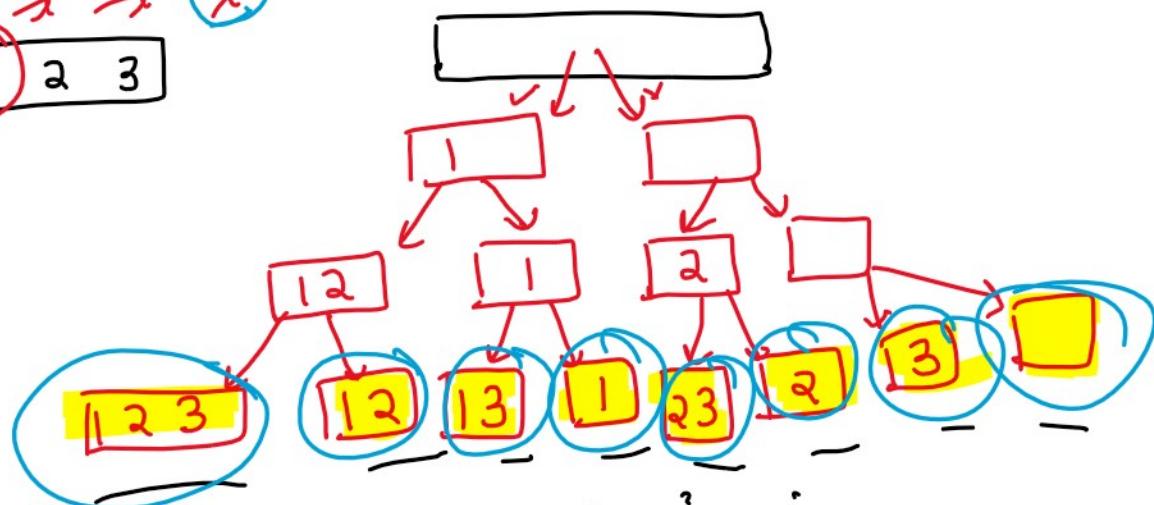
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1	1
12	2
13	3
123	23

Base case  
index =  $n$   
add in res

$j = -3$

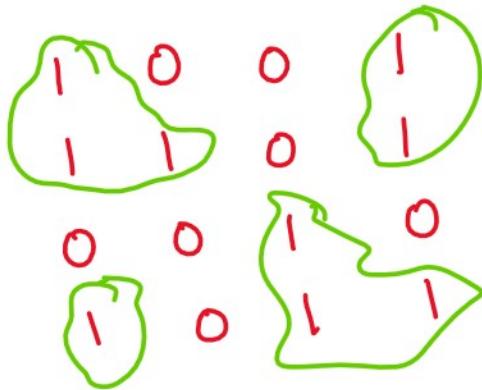


✓ ✓ ✓ ✓ ✓ ✓ ✓ ✓



## Number of islands

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ans = 4

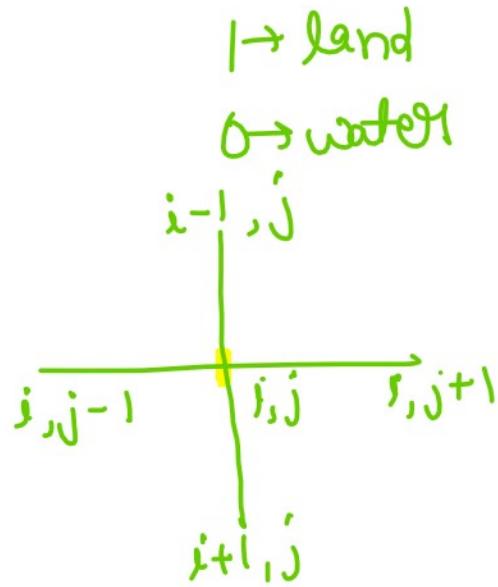
```
for(i=0; i<N) → rows
    for(j=0; j<M) → cols.
        if(grid[i][j] == 1)
            // count++
```

$i=0, j=\emptyset$  |

ans = 4

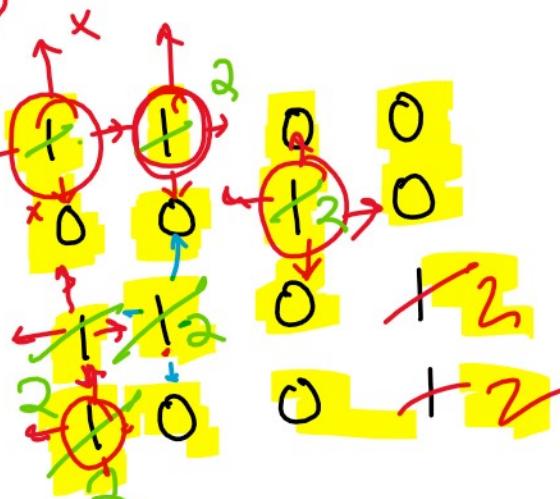
count = 0 / 1 / 2 / 3 / 4

Boundary check / base case  
i & j should be within grid



N x M

dfs( )



1 → land

0 → water

2 → visited

Time =  $O(N \times M)$   
Space =  $O(N \times M)$

Time : -  
Space :  $O(\overline{N \times M})$

- ✓ → flood fill algo
- ✓ → max area island
- ✓ → surrounded regions
- ✓ → word search(i, i)

Rotten orange  
mn step by knkt