

- It is the foundation of generic programming which involves writing a code that is independent of particular type.
- It can be viewed as variable that can be instantiated to any data type, irrespective of whether this data type is fundamental or user defined.
- It enforces code reusability feature of C++

```
#include <iostream>
using namespace std;
template <class T>
void sort (T a[], int n)
{
    for (int i=0; i<n-1; i++)
    {
        for (int j=0; j<n-1-i; j++)
        {
            if (a[j] > a[j+1])
            {
                T temp = a[j];
                a[j] = a[j+1];
                a[j+1] = temp;
            }
        }
    }
}
template <class T>
void print (T a[], int n)
{
    cout << "Array is \t";
    for (int i=0; i<n; i++)
    {
        cout << a[i] << "\t";
    }
}
int main ()
{
    int a;
```

```

cout << "enter the size of the int array " << endl;
cin >> a;
int arr [a];
cout << "enter the elements " << endl;
for( int i=0; i<a; i++)
{
    cin >> arr[i];
}
cout << "original \n";
Print (arr,a);
cout << "sorted " << endl;
Print (arr,a);

cout << "enter the size for float array \n";
cin >> a ; float arr [a];
cout << "enter the elements " << endl;
for( int i=0; i<a; i++)
{
    cin >> arr[i];
}
cout << "original " << endl ;
Print (arr,a);
cout << " sorted \n ";
Print << (arr,a);
}

```

```

cout << "enter the size of the int array " << endl;
cin >> a;
int arr [a];
cout << "enter the elements " << endl;
for( int i=0; i<a; i++)
{
    cin >> arr[i];
}
cout << "original \n";
Print (arr,a);
cout << "sorted " << endl;
Print (arr,a);

cout << "enter the size for float array \n";
cin >> a ; float arr [a];
cout << "enter the elements " << endl;
for( int i=0; i<a; i++)
{
    cin >> arr[i];
}
cout << "original " << endl;
Print (arr,a);
cout << "sorted \n";
Print << (arr,a);
}

```

```

template <class T>
void secondlargest(T a[], int n)
{
    for (int k=0; k<n; k++)
    {
        for (int l=0; l<n-i; l++)
        {
            if (a[l] > a[l+1])
            {
                T temp = a[i];
                a[i] = a[i+1];
                a[i+1] = temp;
            }
        }
    }
    cout << "the second largest element is " << a[n-2];
}

int main ()
{
    int s;
    cout << "enter size of array \n";
    cin >> s;
    int arr[s];
    cout << "enter the elements of integer array " << endl;
    for (int i=0; i < s; i++)
    {
        cin >> arr[i];
    }
    float a;
    secondlargest(arr, s);
    cout << "enter the size for float array \n";
    cin >> a;
    float ar[a];
    cout << "enter the elements for float array " << endl;
    float for (int j=0; j < a; j++)
}

```

Scanned by CamScanner

```

    cin >> arr[a];
}

second_largest-(arr,a);
}

```

3) The steps are -

- 1) start
- 2) input a prefix expression in the 'prefix' string.
- 3) for each last character 'ch' of 'prefix' string
 - if ('ch' is an operand)
 - push 'ch' in 'operandstk'
 - else
 - $op2 = \text{pop an operand from 'operandstk'}$
 - $op1 = \text{pop an operand from 'operandstk'}$
 - value = result of applying 'ch' in ' $op2 \rightarrow op1$ '
 - push 'value' to 'operandstk'
- 4) Pop the top element of the stack 'operandstk' which is the required result.
- 5) Exit.

4) Given expression $- * + A B C D$ where
 $A = 4, B = 3, C = 2, D = 5$

Symbol	op1	op2	value	opnd stack
5				5
2				5,2
3				5,2,3
4				5,2,3,4
+	4	3	7	
*	7	2	14	5,2,7
-	14	5	9	5,14
				9

The steps are

- 1) if ch is an operand , immediately output.
- 2) close parenthesis : pop stack symbol until an open parenthesis appears.
- 3) operator, pop all stack symbols until a symbol of lower precedence or a right-associative symbol of equal precedence appears . Then push the operator.
- 4) End of input, pop all remaining stack symbols.

convert the infix exp $((a+b)*c-(d-e))^\wedge(f+g)$ to postfix using stack.

Symbol	stack Postfix string	string operator
((
c	cc	
a	cc	a
+	cc	a
b	cc+	ab
*	cc*	ab+
c	cc*	ab+c
-	cc*-	ab+c-
(cc*-	ab+c*-
d	cc*-	ab+c*-d
-	cc*-	ab+c*-d
e	cc*-	ab+c*-de
)	cc*-	ab+c*-de-
)	-	ab+c*-de--
\wedge	n	ab+c*-de--
(nc	ab+c*-de--
f	nc	ab+c*-de-f
+	nc+	ab+c*-de-f+
g	nc+	ab+c*-de-fg
)	n+	ab+c*-de-fg+

```

void addQ(int item)
{
    // add an item to the queue
    int k = (rear+1)%max-size;
    if(front == k)
    {
        cout << "queue full";
        return;
    }
    rear = k;
    queue[rear] = item;
}

```

Steps:

- 1) pass a value to the function,
- 2) assign $[(\text{rear}+1) \% \text{maxsize}]$ to some variable
- 3) check if k equal to front.
- 4) if true , show message queue is full.
- 5) else assign the value to queue [rear]

The steps are

- 1) Read in the expression
- 2) Process each character of the expression to a single digit number (0 to 9) then push it into the stack.
If the character corresponds to one of the arithmetic operations (+, -, *, /), then.
 - pop a number off the stack, call it op2.
 - pop a number off the stack, call it op1.
 - combine the operand using the arithmetic operator as follows

$$\text{Result} = \text{op1} \text{ operator } \text{op2}$$

- push result onto the stack.
- 3) When the end is reached, pop the remaining number from the stack. This number is the value of the expression.

Postfix expression $ab+c*de--fg+^n$

$12+3445--67+^n$

Scanned symbol	op1	op2	result	stack
1				1
2				1, 2
+	2	1	3	3
3				3, 3
*	3	3	9	9
4				4, 9
5	-	-	-	4, 5, 4, 9
-	4	5	-1	-1, 9
-	9	-1	10	10
6	-	-	-	6, 10
7	-	-	-	7, 6, 10
+	6	7	13	13, 10
n	10	13	10^{13}	10^{13}

The problems associated with linear queue are
By the definition of the queue ,when we add an element in queue ,rear pointer is increased by 1
element in the queue . when we remove front pointer is increased by 1.

Consider the operations performed on queue with size 5 as follows .

i) Initially empty queue is there so, front = 0,
so rear = -1

ii) when we add 5 elements to the queue ,the data of the queue becomes as follows front = 0 & rear = 4.

iii) Now we have deleted 2 elements from queue so there should be space for another 2 elements, but as rear pointer is pointing at last position & queue overflow condition is true, we can't insert any new element in the queue even if it has empty space. To overcome this problem there is another variation of queue known as circular queue.

Big Oh notation is a mathematical notation that describes the limiting behaviour of function when argument tends towards a particular value or infinity. In computer science, big O notation is used to classify algorithm by how they respond to change in input size such as how the processing time of an algorithm changes as the problem size becomes extremely large.

$f(n) = O(g(n))$ if there exists positive constants and n_0 such that $f(n) \leq cg(n)$ for all $n, n > n_0$.

Proof:

By the big oh notation $f(n)$ is $O(n^3)$ if $f(n) \leq cn^3$ for some $n > n_0$ let us check this condition.

$$\text{if } n^3 + 20n + 1 \leq cn^3$$

$$\text{then } 1 + \frac{20}{n^2} + \frac{1}{n^3} \leq c$$

then Big Oh notation holds for $n > n_0 = 1$ and $c > 22$.

Performance of an algorithm can be measured in two ways

- i) Space complexity
- ii) Time complexity.

space complexity refers to the amount of space needed by a program to run to completion. It has 2 parts.

- i) Fixed Part - it's independent of input and outputs, constants, fixed size component variable, instruction spaces, simple variable.
- ii) Variable part is space needed by component variable whose size is dependent on the particular problem instance being solved and recursion stack space.
 $S(P)$ space requirement of any program is
 $S(P) = c + S_p(I)$ where c is a constant and I is the instance characteristics.

```
float sum(float a[], const int n)
```

```
{ float sum = 0;
```

```
for(int i = 0; i < n; i++)
```

```
{ sum += a[i]; }
```

```
}
```

```
3.
```

Since n is passed by value, one word must be allocated to it. a is the base address, one word for it. Space needed by the fun is independent of n , hence

$$S_p(n) = 0.$$

Time complexity of any program $T(P)$ is given by compile time and run time, compile time doesn't depend upon instance characteristics. We focus on the run time of the program denoted by $T(P)$.

Program step is syntactically and semantically meaningful segment of the program that has run time and independent of the instance characteristics.

float sum (float a[], int n)

{ int s=0; // onetime

for (int i=0; i<n; i++) // for n+1 times

{ s+=a[i]; // ntimes

}

return s; //onetime

}

Total $2n+3$.

- 13) Recursion - It is the name given for expressing anything in terms of itself.

Recursion function is a function which calls itself until a particular condition is met.

Properties.

- Recursive algorithm should terminate at some point otherwise recursion will never end.
- hence recursive algorithm should have stopping condition to terminate along with recursive calls.

Advantages -

- (easier and simpler versions of algorithm can be created using recursion)
- Recursive definition of a problem can be easily translated into recursive functions.
- Lot of book keeping activities such as initialisation etc required in iterative solutions is avoided.

Disadvantage -

- When a function is called, the function saves form parameters, local variables and return addresses and hence consumes a lot of memory.
- lot of time in pushing and popping and hence consumes more time to compute the result

```

#include <iostream>
#include <string.h>
using namespace std;
bool isPal( const string &str, int start, int end )
{
    if( start >= end )
        return true ;
    if( str[ start ] != str[ end ] )
        return false ;
    return isPal( str, ++start, --end );
}

```

```

int main ()
{
    char str[50];
    cout << " enter string \n ";
    cin >> str;
    int i = 0;
    int O = strlen(str);
    O = O - 1;
    if( int check = isPal( str, i, s ) )
        cout << " is Palindrome ";
    else
        cout << " not Palindrome ";
}

```

```

#include <iostream>
using namespace std;
int fib( int n )
{
    if( n == 0 )
        return 0;
    if( n == 1 )
        return 1;
    else
        { return fib( n - 1 ) + fib( n - 2 ), } }

```

```
int main
{
    int f;
    f = fib(4);
    cout << f;
}
```

Fibonacci sequence

$$0, 1, 1, 2, 3, 5, 8 \dots$$

Each element is the sum of two preceding elements and fibonacci of a number is the value at that position in sequence.

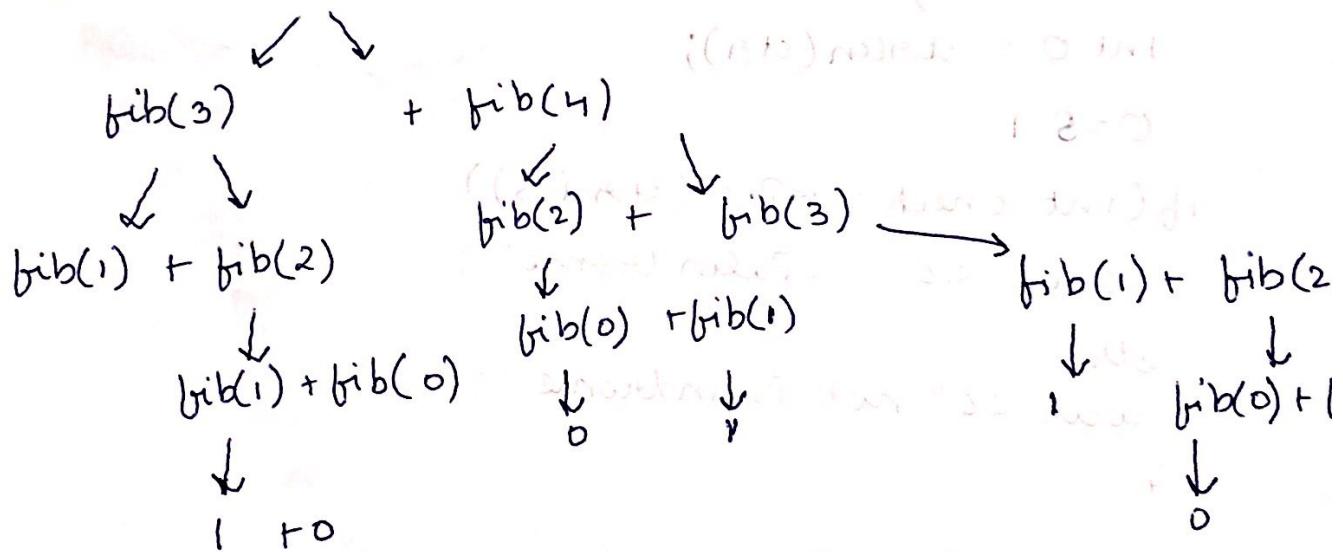
$$\text{ie } \text{fib}(0) = 0$$

$$fib(1) = 1$$

$$\text{fib}(2) = 1$$

$$\text{fib}(3) = 2$$

For fib(5)



- 16) A stack is an ordered list in which insertions and deletions are made at one end called the top. It is also called first in first out list

```
template <class T, int Maxsize>
```

class stack

{ T arr [maxsize];

int top;

```
stack() { top = -1; }
```

```
bool isEmpty() { return top == -1; }
```

```
bool isFull() { return top == maxsize - 1; }
```

```
void push(Titem)
```

```
{ if (Isfull()) {
```

```
cout << " stack overflow"; }
```

```
arr[++top] = item;
```

```
}
```

```
T pop {
```

```
if (IsEmpty()) {
```

```
cout << " stack underflow";
```

```
return -99; }
```

```
return arr[top--];
```

```
}
```

```
};
```

```
#include <iostream>
```

```
using namespace std;
```

```
int counter = 0;
```

```
void tower(int n, char source, char temp, char dest)
```

```
{ if (n == 1)
```

```
{
```

```
cout << " move disk1 from " << source << " to " << dest <<
```

```
endl;
```

```
counter++;
```

```
return ;
```

```
}
```

```
tower(n-1, source, dest, temp);
```

```
cout << " move disk" << n << " from " << source <<
```

```
" " << dest << endl;
```

```

        counter++;
    tower(n-1, temp, source, dest);
}
int main()
{
    char a, b, c;
    int n;
    cout << "enter the no of disk";
    cin >> n;
    tower(n, 'a', 'b', 'c');
    cout << "the steps are ";
}

```

```

#include <iostream.h>
#include <stdlib.h>
using namespace std;
class twoStack
{
    int *arr;
    int size;
    int top1, top2;
public:
    twoStack(int n)
    {
        size = n;
        arr = new int[n];
        top1 = -1;
        top2 = size;
    }
    void push(int x)
    {
        if (top1 < top2 - 1)
        {
            top1++;
            arr[top1] = x;
        }
    }
}

```

```
else
{
    cout << "stack overflow" << endl;
    exit(1);
}
```

```
void push2(int x)
```

```
{ if (top1 < top2 - 1)
{
    top2--;
    arr[top2] = x;
}
```

```
else
```

```
{ cout << "stack overflow \n";
    exit(0);
}
```

```
int pop1()
```

```
{ if (top1 >= 0)
{
    int x = arr[top1];
    top1--;
    return x;
}
```

```
else
```

```
{ cout << "stack underflow" << endl;
    exit(0);
}
```

```
}
```

```
int pop2()
```

```
{ if (top2 < size)
{
    int x = arr[top2];
    top2++;
    return x;
}
```

```

else
{
    cout << "stack underflow" << endl;
    exit(1);
}
};

int main()
{
    twostacks ts(5);
    ts.push1(5);
    ts.push2(10);
    ts.push2(15);
    ts.push1(11);
    ts.push(7);

    cout << "popped element from stack1 is " << ts.pop1();
    cout << endl;

    ts.push(40);

    cout << "popped element from stack2 is " << ts.pop2();
    cout << endl;

    return 0;
}

```

- 19) There is a special type of ADT called the priority queue. This is a special type of table whose items are assigned priority values as they are added into the structure. This priority value is used to determine the order in which the items in the table are to be processed.

Methods:

create

This is the priority queue's class constructor. It will create a new, empty priority queue.

isEmpty()

returns true if the

insert ()

Adds a new item to the priority queue.

delete ()

Removes and returns the item in the priority queue with the highest priority.

Priority queue can be implemented in 4 ways

a) Using arrays

The sort order will be from smallest to largest (ie the item with highest priority will always be the last item in the array).

Limitations :

- Only a limited number of items may be added to the priority queue. If their limit is set to a large number, then this structure may waste a lot of memory.
- Although the delete() method is very efficient, the insert () method is very inefficient.

b) using linked list

The sort order will be from largest to smallest (ie the item with highest priority is always the first item in the list).

Limitations :

- Although the delete() method is very efficient, the insert () method is inefficient.

c) using a sorting algorithm to make a priority queue

A sorting algorithm can be used to implement a priority queue.

Limitations .

- Although the insert () method is efficient, the delete() method is inefficient.

```

20) #include <iostream>
using namespace std;
#include "stack.h";
char choice;
void menu()
{
    cout << "\n\n Menu";
    cout << "\n 1. Arrival (A)";
    cout << "\n 2. Departure (D)";
    cout << "\n 3. Exit (E)";
    cout << "\nEnter your choice";
    cin >> choice;
}

int main()
{
    stack<int> cars(10);
    int cno;
    int pos;
    int c;
    int temp[10] = {0};
    int count[10] = {0};
    menu();
    while (choice != 'E')
    {
        switch (choice)
        {
            case 'A':
                cout << "Enter the car no:";
                cin >> cno;
                if (cars.push(cno))
                {
                    cout << "\nSorry garage is full";
                    break;
                }
            else
                cout << "\nCar no ";
        }
    }
}

```

```
break;  
}
```

case 'D':

```
cout << "\n Enter the car number to depart : ";
```

```
cin >> cno;
```

```
for (int i = 9; i >= 0; --i)
```

```
{
```

```
    c = cars.pop();
```

```
    if (c == cno)
```

```
{
```

```
    cout << "\n Car successfully departed.";
```

```
    cout << "\n It was moved " << count[i] << " times to
```

```
        let other cars depart.";
```

```
    break;
```

```
}
```

```
else
```

```
{
```

```
    temp[i] = c;
```

```
    cout << [i]++;
```

```
}
```

```
}
```

```
for (int i = 0; i < 10; ++i)
```

```
{ if (temp[i] == 0)
```

```
    continue;
```

```
}
```

```
else
```

```
{ cars.push(temp[i]);
```

```
}
```

```
}
```

```
g
```

```
g
```

```
g
```