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**Problem 0**

<b>Points:</b>
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**Acknowledgements**

- (a) I did not work in a group.
- (b) I did not consult without anyone my group members.
- (c) I did not consult any non-class materials.

## Problem 1

Points:

## DFS Basics

(a)  $A \rightarrow B \rightarrow D \rightarrow E \rightarrow G \rightarrow F \rightarrow C \rightarrow H \rightarrow I$ 

$A$	$(1, 12)$
$B$	$(2, 11)$
$C$	$(13, 18)$
$D$	$(3, 6)$
(b) $E$	$(4, 5)$
$F$	$(8, 9)$
$G$	$(7, 10)$
$H$	$(14, 17)$
$I$	$(15, 16)$

	Edge	Type
	$(A, B)$	Tree
	$(A, E)$	Forward
	$(B, D)$	Tree
	$(D, E)$	Tree
	$(E, D)$	Back
(c)	$(B, G)$	Tree
	$(G, F)$	Tree
	$(G, D)$	Cross
	$(C, H)$	Tree
	$(H, I)$	Tree
	$(C, I)$	Forward

**Problem 2**

<b>Points:</b>
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**Pre and Post Processing**

(a) For our DFS algorithm, if  $\text{post}(u) < \text{post}(v)$ , then,

- Case 1 :  $[\text{pre}(u), \text{post}(u)][\text{pre}(v), \text{post}(v)]$
- Case 2 :  $[\text{pre}(v), [\text{pre}(u), \text{post}(u)]\text{post}(v)]$

These are the only 2 Cases for an undirected graph where  $\text{post}(u) < \text{post}(v)$ .

Since we know there is an edge between these 2 nodes, Case 1 can not happen because we must visit all the neighbors of a node before marking it as visited. That means, in Case 1, node  $u$  is marked visited before exploring edge  $v$  as it has  $\text{post}(u) < \text{post}(v)$  which is violating the DFS rule of exploring all neighbor nodes before marking it visited. So, Case 2 is the only possible one, which yields  $v$  as the ancestor of  $u$ . The statement is True.

(b) First, Run Depth First Search Algorithm on the tree while also keeping the timestamps for the time when  $u$  is started to get explored (pre-number) and the time when  $u$  is finished getting explored (post-number). For every node  $u$ ,  $\text{pre}(u)$  denotes the time when we began exploring  $u$  and  $\text{post}(u)$  denote the time when we finished it. This process will take linear time since the *Explore()* function will never be called on a Node more than once.

Now, to check whether  $u$  an ancestor of  $v$ , we have to check whether:

$$\text{pre}(u) < \text{pre}(v) \text{ and } \text{post}(u) > \text{post}(v)$$

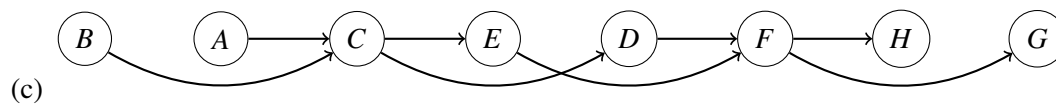
If the above condition becomes True, then  $u$  is the ancestor of  $v$ . This condition will itself take a constant time because it is just comparing two numbers.

## Problem 3

Points:

## Linearization Basics

(a)	$A$	$(1, 14)$
	$B$	$(15, 16)$
	$C$	$(2, 13)$
	$D$	$(3, 10)$
	$E$	$(11, 12)$
	$F$	$(4, 9)$
	$G$	$(5, 6)$
	$H$	$(7, 8)$

(b) Sources :  $A, B$  ( $A, B$  are having most elevated post number)Sinks :  $G, H$  ( $G, H$  are having least post number)(d) We have three times where we can choose between two vertices on a path so  $2^3 = 8$  linearization.