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**Problem 0**

<b>Points:</b>
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**Acknowledgements**

- (a) I did not work in a group.
- (b) I did not consult without anyone my group members.
- (c) I did not consult any non-class materials.

**Problem 1****Points:****Compare Growth Rates**

- (a)  $n^{1.5} = \Omega(n^{1.3}) \quad \because n^{1.3} < n^{1.5}$
- (b)  $2^{n-1} = \Theta(2^n) \quad \because 2^{n-1} = \frac{2^n}{2}$  and constants (in this case,  $\frac{1}{2}$ ) does not matter
- (c)  $n^{1.3 \log n} = \Omega(n^{1.5}) \quad \because$  asymptotically  $n^{1.3 \log n} > n^{1.5}$
- (d)  $3^n = \Omega(n \cdot 2^n) \quad \because$  asymptotically  $3^n > n \cdot 2^n$
- (e)  $(\log n)^{100} = O(n^{0.1}) \quad \because$  power of  $\log n$  will only weaken the infinity of  $\log n$  and  $n^{0.1} > (\log n)^{100}$
- (f)  $n = \Omega((\log n)^{\log(\log n)}) \quad \because$  infinity of  $(\log n)^{\log(\log n)} >$  infinity of  $n$
- (g)  $2^n = \Omega(n!)$   $\because$  infinity of exponential function is greater than the infinity of factorial function
- (h)  $\log(e^n) = O(n \cdot \log n) \quad \because \log(e^n) = n$  and we know that  $n < n \log n$
- (i)  $n + \log n = \Theta(n + (\log n)^2) \quad \because$  dominating term is  $n$  and powers of  $\log n$  will not affect infinities
- (j)  $5n + \sqrt{n} = \Omega(\log n + n) \quad \because$  comparable functions are  $\sqrt{n}$  and  $\log n$ ;  $\sqrt{n} > \log n$

**Problem 2**

<b>Points:</b>
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**Problem 3**

<b>Points:</b>
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Your solution starts here ...