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Problem 0

Points:

Acknowledgements

- (a) I did not work in a group.
- (b) I did not consult without anyone my group members.
- (c) I did not consult any non-class materials.

Problem 1

Points:

DFS variations

- (a) **Algorithm:** It is given that the car can only hold enough gas to cover L miles while any stretch of highway (l_e) $e \in E$ may or may not be greater than L . Therefore, initially we can simply remove edges (e) of a greater value. Now, the next step is to see if the city t is still reachable from s . To check that, we will be using Breadth First Search (*BFS*) algorithm. So if vertex t is reachable, then the path is feasible otherwise, it is not feasible.

Runtime: Running *BFS* will take $O(|V| + |E|)$.

- (b) Here, we will modify Dijkstra's algorithm to find paths that minimize the maximum weight of any edge on the path (instead of the path length).

Here, we take minimum of $\max(\text{visited}, \text{not visited adjacent})$.

Algorithm:

ModifiedDijkstra($G = (V, E)$)

 for all $u \in V$ do

$\text{dist}(u) = \infty$

$\text{prev}(u) = \text{nil}$

 end

$\text{dist}(s) = 0$

 Let H be a priority queue constructed with all nodes in V using dist as the key

 while H is not empty do

$u = \text{deleteMin}(H)$

 for all edges $(u, v) \in E$ do

 if $\text{dist}(v) > \max(\text{dist}(u), l(u, v))$

$\text{dist}(v) = \max(\text{dist}(u), l(u, v))$

$\text{prev}(v) = u$

$\text{decreaseKey}(H, v)$

 end

 end

 end

end

Runtime: Time complexity is same as Dijkstra's algorithm i.e. $O((|V| + |E|)\log|E|)$.

Problem 2

Points:

Shortest bitonic paths

(a)

(b)

Problem 3

Points:

Dijkstra's on negative

(a)

(b)