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**Problem 0**

<b>Points:</b>
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**Acknowledgements**

- (a) I did not work in a group.
- (b) I did not consult with anyone in my group members.
- (c) I did not consult any non-class materials.

## Problem 1

Points:

**Algorithm 1:** GREEDY-HORN

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**Input** : set of Horn clauses  
**Output**: either the assignment or "unsatisfiable"

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1 Set all variables to 0;
2 while  $\exists$  an " $\implies$ " that is not satisfied do
3   | Set its RHS to 1;
4 end while
5 if all pure negative clauses are 1 then
6   | return the assignment
7 end if
8 else
9   | return "unsatisfiable"
10 end if

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- (a) According to the algorithm, first set all variables to 0  $\implies$

$$w = 0, x = 0, y = 0, z = 0$$

Now there are 5 clauses having " $\implies$ " out of which 4<sup>th</sup> clause ( $\implies x$ ) is not satisfied. So, we will set RHS of this clause to 1, that is, **x=1**.

This will lead to reconsidering the assignment of y because according to 3<sup>rd</sup> clause ( $x \implies y$ ), if LHS is True then RHS should be set to 1, that is, **y=1**.

This will lead to reconsidering the assignment of w because according to 5<sup>th</sup> clause ( $x \wedge y \implies w$ ), if LHS is True then RHS should be set to 1, that is, **w=1**.

This will lead to reconsidering the assignment of z because according to 1<sup>st</sup> clause ( $w \wedge y \wedge z \implies x$ ), if RHS is True then LHS should be resolved to 1, that is, **z=1**.

Now, pure negative clauses are failed to satisfy, so there is no satisfying assignment, hence algorithm will return "**unsatisfiable**".

- (b) According to the algorithm, first set all variables to 0  $\implies$

$$w = 0, x = 0, y = 0, z = 0$$

Now there are 4 clauses having " $\implies$ " out of which 4<sup>th</sup> clause ( $\implies z$ ) is not satisfied. So, we will set RHS of this clause to 1, that is, **z=1**.

This will lead to reconsidering the assignment of w because according to 2<sup>nd</sup> clause ( $z \implies w$ ), if LHS is True then RHS should be set to 1, that is, **w=1**.

Here, x and y need not to be changed since the implications are still satisfied with having **x=0, y=0**.

Now, pure negative clauses are still 1 and hence, satisfied. So, the algorithm will return the assignment

$$\mathbf{w=1, x=0, y=0, z=1}$$

**Problem 2**

<b>Points:</b>
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**Problem 3**

<b>Points:</b>
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