

CS698D, Assignment 1

Harshit Maheshwari

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1 PROBLEM 1

The question is done according to the algorithm described in notes.

Data: A sequence of words and numbers in an array: A

Result: Decompressed sequence of words: in *output*

while ($A[i] \neq \phi$) **do**

if $A(i+1).type == word$ **then**

$output += A(i+1);$

$arr.first = A(i+1);$

$i += 2;$

else

if $i > arr.length()$ **then**

$return Error$

end

$word = arr(i);$

$output += word;$

$arr.remove(i);$

$arr.begin = word;$

$i ++;$

end

end

2 PROBLEM 2

Expected Length of first run:

$$\begin{aligned}
 E_l(2) &= \lim_{n \rightarrow \infty} \left(\sum_{l=1}^n p(l)l \right) \\
 &= 1 \left(\frac{1}{2^2} + \frac{1}{2^2} \right) + 2 \left(\frac{1}{2^3} + \frac{1}{2^3} \right) + \dots \infty \\
 &= 1 \cdot \left(\frac{1}{2} \right) + 2 \left(\frac{1}{2^2} \right) + \dots + i \left(\frac{1}{2^i} \right) \dots \infty \\
 \frac{E_1(2)}{2} &= 1 \cdot \left(\frac{1}{2^2} \right) + 2 \left(\frac{1}{2^3} \right) + \dots + i \left(\frac{1}{2^{(i+1)}} \right) \dots \infty
 \end{aligned}$$

Subtracting the above 2 equations:

$$\begin{aligned}
 \frac{E_1(2)}{2} &= \frac{1}{2} + \frac{1}{2^2} + \frac{1}{2^3} \dots \\
 &= 1 \\
 E_1(2) &= 2
 \end{aligned}$$

For alphabet os size: s

$$\begin{aligned}
 E_l(s) &= 1 \left(\frac{s(s-1)}{s^2} \right) + 2 \left(\frac{s(s-1)}{s^3} \right) + 3 \left(\frac{s(s-1)}{s^4} \right) + \dots \infty \\
 \frac{E_l(s)}{s} &= 1 \left(\frac{s(s-1)}{s^3} \right) + 2 \left(\frac{s(s-1)}{s^4} \right) + 3 \left(\frac{s(s-1)}{s^5} \right) + \dots \infty
 \end{aligned}$$

Subtracting the above 2 equations:

$$\begin{aligned}
 E_l(s) \left(1 - \frac{1}{s} \right) &= s(s-1) \left(\frac{1}{s^2} + \frac{1}{s^3} + \dots \right) \\
 &= \frac{s(s-1)}{s^2} \left(\frac{1}{1 - \frac{1}{s}} \right) \\
 E_l(s) &= \frac{s}{s-1}
 \end{aligned}$$

3 PROBLEM 3

We have a sequence of words: $w_1, w_2, w_3 \dots w_n$

To show:

$$l_p(w_{i-1}, w_i) \geq l_p(w_i, w_j) \forall j < (i-1) \quad (3.1)$$

Proof by contradiction:

Suppose, $\exists j < i-1$ s.t. $l_p(w_{i-1}, w_i) [= l_1] \geq l_p(w_i, w_j) [= l_2]$

Then, w_{i-1}, w_j match till l_1 and $w_{i-1}(l_1 + 1) \neq w_j(l_1 + 1)$

$$w_{i-1}(l_1 + 1) \leq w_i(l_1 + 1) \text{ [Lexicographically sorted]}$$

$$w_{i-1}(l_1 + 1) \leq w_j(l_1 + 1)$$

$$w_{i-1}(l_1 + 1) < w_j(l_1 + 1)$$

But then $j > i - 1$

$\Rightarrow \Leftarrow$

Therefore, $l_p(w_{i-1}, w_i) \geq l_p(w_i, w_j) \forall j < (i - 1)$

4 PROBLEM 4

To prove

$$H(Y|X) = 0 \leftrightarrow Y(X) \text{ or}$$

$$\forall x_0 \in X \exists y_0 \in Y \text{ s.t. } P(y_0|x_0) = 1 \text{ and}$$

$$P(y_1|x_0) = 0 \forall y_1 \neq y_0$$

Proof

\Rightarrow

$$\sum_y \sum_x P(x, y) \log \frac{1}{P(y|x)} = 0 \quad (4.1)$$

Since, each of the term in the summation is greater than 0, each term has to individually 0.

Now, suppose,

$$\exists y_1 \text{ s.t. } P(y_1|x) \in (0, 1)$$

$$\exists y_2 \text{ s.t. } P(y_2|x) \in (0, 1)$$

Then, $P(x, y_1) \log \frac{1}{P(y_1|x)} \neq 0$.

Also, $P(x, y_2) \log \frac{1}{P(y_2|x)} \neq 0$

$\Rightarrow \Leftarrow$

\Leftarrow

Suppose, $\exists y_0$ s.t. $Pr(y_0|x_0) = 1$ and $Pr(y_1|x_0) = 0 \forall y \neq y_0$

Then, it is easy to see that all the terms in the equation 4.2 are 0.

Hence, proved.

5 PROBLEM 5

5.1

Case 1: Sequence is unbounded.

Then trivially limit is ∞ because the sequence is monotonically decreasing.

Case 2: Sequence is bounded.

Let $x = \inf x_n$

$\forall \epsilon > 0 \exists x_n$ s.t. $(x + \epsilon) > x_n$ [By inf definition]

$\forall m > n \ x_m < x_n$

$\forall m > n \ x_m < x + \epsilon$

$\forall \epsilon \forall m > N \ |x_m - x| < \epsilon$

So, $x_m \rightarrow x$

5.2

Let the sequence be x_1, x_2, x_3, \dots

$$\limsup x = \lim_{n \rightarrow \infty} \sup_{m > n} x_m$$

$$\sup_{m > i} x_m \geq \sup_{m > (i+1)} x_m$$

So, $\sup_{m > i} x_m, \sup_{m > (i+1)} x_m, \dots$ form a decreasing sequence. By analogy of previous result in-creasing sequence of real numbers has a limit.

Similarly,

$$\liminf x = \lim_{n \rightarrow \infty} \inf_{m > n} x_m$$

$$\inf_{m > i} x_m \leq \inf_{m > (i+1)} x_m$$

So, $\inf_{m > i} x_m, \inf_{m > (i+1)} x_m, \dots$ form an increasing sequence. By previous result this forms a decreasing sequence of real numbers has a limit.

6 PROBLEM 6

We will show,

$$f = O(g) \leftrightarrow \limsup_{n \rightarrow \infty} \frac{|f(n)|}{|g(n)|} < \infty \quad (6.1)$$

\Rightarrow

If $g(n) = 0$ then trivially $f(n) = 0$. If $g(n) \neq 0$:

$\exists C \exists n > N \ |f(n)| < C|g(n)|$

$$\frac{|f(n)|}{|g(n)|} \leq C \quad (6.2)$$

If a sequence: x is bounded above by C then $\limsup x \leq C$

Since, the sequence is bounded above by C . So,

$$\limsup_{n \rightarrow \infty} \frac{|f(n)|}{|g(n)|} \leq C \quad (6.3)$$

←

We say that $\alpha = \limsup_{n \rightarrow \infty} f(n)$ if following conditions hold:

$$(\epsilon > 0)(\exists N)(\forall n > N)(f(n) < \alpha + \epsilon)$$

$$(\epsilon > 0)(\forall N)(\exists n > N)(f(n) > \alpha - \epsilon)$$

Putting $\epsilon = 1$ in the above definition.

$$\forall n > N \frac{|f(n)|}{|g(n)|} < C + 1 \quad (6.4)$$

$$\forall n > N |f(n)| \leq (C + 1)|g(n)| \quad (6.5)$$

So, $f(n) = O(g(n))$. Hence, proved.

Reference: parc.im.pwr.wroc.pl/cichon/Math/BigO.pdf

7 PROBLEM 7

The function $f(x) = \frac{1}{x}$ is undefined on $x = 0$ so it is discontinuous. [

$$\lim_{x \rightarrow 0^-} f(x) = -\infty$$

$$\lim_{x \rightarrow 0^+} f(x) = +\infty$$

Since both the limits are different we cannot define $f(x)$ at $x = 0$ s.t. the function becomes continuous.]

$f: \mathbb{R} \rightarrow \mathbb{R}$ is continuous $\leftrightarrow \forall \epsilon, x \exists \delta \forall y (|x - y| < \delta) \rightarrow (|f(x) - f(y)| < \epsilon)$

Consider some point $x \neq 0$ at which we will show $f(x)$ is continuous. If $y > \frac{x}{2}$

Choose $\delta < \min\left(\frac{|x|}{2}, \frac{\epsilon x^2}{2}\right)$

$$\left|\frac{1}{x} - \frac{1}{y}\right| = \left|\frac{y - x}{xy}\right| \leq \frac{\delta}{|x|(|x| - \delta)} \leq \frac{2\delta}{x^2} \leq \epsilon$$

If $x \neq 0$ then $\forall \epsilon$ we can find a suitable value of δ s.t. the continuity condition holds.

At $x = 0$ the function is discontinuous.

8 PROBLEM 8

Given:

Continuous function $f: [0, 1] \rightarrow \mathbb{R}$ Since, f is continuous on $[0, 1]$,

$\forall \epsilon, x \exists \delta \forall y \ |x - y| < \delta \rightarrow |f(x) - f(y)| < \epsilon, \ x, y \in [0, 1]$

Let the partitions of the interval $[0, 1]$ be denoted by x_1, x_2, \dots, x_n s.t. $|x_{i+1} - x_i| < \delta \ \forall i \in [n]$

For partition P_i

$$\begin{aligned}
 U_i &= f(x_{i+1})(x_{i+1} - x_i) \\
 L_i &= f(x_i)(x_{i+1} - x_i) \\
 U_i - L_i &= (f(x_{i+1}) - f(x_i))(x_{i+1} - x_i) \\
 &\leq \epsilon(x_{i+1} - x_i) \\
 \sum_{i=1}^{i=n} (U_i - L_i) &= \epsilon(1 - 0)
 \end{aligned}$$

$\Rightarrow f$ is integrable.

9 PROBLEM 9

$$\begin{aligned}
 L^1 &= 0 \\
 L^2 &= \frac{1}{2^3}(0^2 + 1^2) \\
 L^3 &= \frac{1}{3^3}(0^2 + 1^2 + 2^2) \\
 L^4 &= \frac{1}{4^3}(0^2 + 1^2 + 2^2 + 3^2) \\
 L^5 &= \frac{1}{5^3}(0^2 + 1^2 + 2^2 + 3^2 + 4^2) \\
 L^n &= \frac{1}{n^3}(0^2 + 1^2 + \dots + (n-1)^2) \\
 &= \frac{(2n-1)n(n-1)}{6n^3} \\
 L_{n \rightarrow \infty}^n &= \frac{1}{3}
 \end{aligned}$$

Now, checking for upper limit:

$$\begin{aligned}
 U^1 &= 1 \\
 U^2 &= \frac{(1^2 + 2^2)}{2^3} \\
 U^3 &= \frac{(1^2 + 2^2 + 3^2)}{3^3} \\
 U^n &= \frac{(2n+1)(n+1)n}{6n^3} \\
 U_{n \rightarrow \infty}^n &= \frac{1}{3}
 \end{aligned}$$

Since, $L_{n \rightarrow \infty}^n = U_{n \rightarrow \infty}^n$ the function is Riemann integrable and the value of the integral is: $\frac{1}{3}$

10 PROBLEM 10

To prove:

$$\mu(A_1 \triangle A_2) = 0 \rightarrow \mu(A_1) = \mu(A_2) \quad (10.1)$$

Proof:

$$\begin{aligned} \mu(A_1 \triangle A_2) &= 0 \\ \Rightarrow \mu(A_1 \setminus A_2 \cup A_2 \setminus A_1) &= 0 \\ \text{Since both the sets are disjoint} \\ \Rightarrow \mu(A_1 \setminus A_2) + \mu(A_2 \setminus A_1) &= 0 \end{aligned}$$

Since the range of μ is $\mathbb{R}^+ \cup 0$, both the terms have to be individually 0.

$$\begin{aligned} \Rightarrow \mu(A_1 \setminus A_2) &= 0 \\ \text{Also, } \mu(A_1) &= \mu(A_1 \setminus A_2 \cup A_1 \cap A_2) \\ \mu(A_1) &= \mu(A_1 \setminus A_2) + \mu(A_1 \cap A_2) \\ \text{Since both the sets are disjoint} \\ \mu(A_1) &= \mu(A_1 \cap A_2) \\ \text{Similarly, } \mu(A_2) &= \mu(A_1 \cap A_2) \end{aligned}$$

Therefore, $\mu(A_1) = \mu(A_2)$