

## Tutorial - 4

- ①.1 4-bit 2's complement numbers  $\rightarrow (-8, 7)$   
 8-bit " " "  $\rightarrow (-2^7, 2^7 - 1)$   
 16-bit " " "  $\rightarrow (-2^{15}, 2^{15} - 1)$   
 $(-32768, 32767)$

①.2 Let's first think about the cases when the overflow affects the results. We write down the truth table corresponding to the addition of the last bit  $x_{n-1}$  &  $y_{n-1}$  given the carry-in  $c_{n-1}$  resulting in the sum  $s_{n-1}$  and carry-out  $c_n$ .

$x_{n-1}$	$y_{n-1}$	$c_{n-1}$	$c_n$	$s_{n-1}$
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

← Suspicious cases  
 $c_n \oplus c_{n-1} = 1$   
 ←

Now recall that the MSB indicates the sign of the number. Hence, adding two positive numbers should result in a +ve number, while adding two -ve numbers should result in a -ve number. With these underlying conditions, there are two suspicious cases where the overflow results into unexpected results.

Now, for the three cases of 4-bit, 8-bit & 16-bit numbers, it is straightforward to verify with some examples.



2.1 We have.

$$X = x_7 x_6 x_5 x_4 x_3 x_2 x_1 x_0$$

$$Y = y_7 y_6 y_5 y_4 y_3 y_2 y_1 y_0$$

$$C = c_8 c_7 c_6 c_5 c_4 c_3 c_2 c_1 c_0$$

Since  $c_i = g_{i-1} + p_{i-1} \cdot c_{i-1}$ , ~~the~~ given the inputs  $X$  and  $Y$ ,  $c_i$  will take the value of  $c_{i-1}$  after time  $\Delta t$  which is also equivalent to the delay corresponding an AND gate followed by an OR gate.

Now since ~~and~~ we want to make sure that  $c_8$  toggles after every  $\Delta t$ , we can have

$$C = 010101010$$

~~If this~~ we would ensure that the carry value propagate, we can allow  $c_8$  to get the value of  $c_7$  after  $\Delta t$ ,  $c_6$  after  $2\Delta t$ ,  $c_5$  after  $3\Delta t$  and so on.

Hence, the set of selected inputs should ensure that

①  $C = 010101010$

② the carry values propagate, but ~~the~~ <sup>are not</sup> generated.

To satisfy the first condition, we ~~need to~~ can have identical  $X$  and  $Y$  ~~as identical~~ with alternate bits as 0 and 1.

$$X = 01010101$$

$$Y = 01010101$$

Now at  $T=0$ , we have to set <sup>each</sup>  $x_i$  and  $y_i$ , such that

$$g_i = x_i \cdot y_i = 0 \quad \& \quad p_i = x_i + y_i = 1$$

Therefore, to satisfy second condition, we can have one of  $x_i$  and  $y_i$  as 1 and <sup>we</sup> set the other as 0.



~~That is,~~

For example,

$$X = 00000000$$

$$Y = 11111111$$

Now,  $c_8$  will toggle after every  $\Delta t$ .

		$c_8 c_7 c_6 c_5 c_4 c_3 c_2 c_1 c_0$
Time ↓	$T=0$	$Q = 010101010$
	$\Delta t$	<del>010101010</del>
	$2\Delta t$	010101000
	$3\Delta t$	010100000
	$4\Delta t$	010100000
	$5\Delta t$	101000000
	$6\Delta t$	010000000
	$7\Delta t$	100000000
	$8\Delta t$	000000000