

* Simplex Method

$$1) \quad b \geq 0 \quad 2) \quad x \geq 0$$

3) convert all inequalities into equal

Slack Variables ≥ 0

with cost coefficients set 0 with each of them in objective funcⁿ

$$Ax = b, \quad x \geq 0.$$

Identity submatrix should be present in A .

$$\begin{array}{c} Ax \\ \downarrow \\ m \times n \end{array} \quad \left[\begin{array}{c} I_{m \times m} \end{array} \right]$$

This ease out the choice of initial BFS to start the algo.

\therefore We can take basis matrix $B = I_{m \times m}$

$$A = [B : R] \quad \text{set } x_R = 0$$

$$B x_B = b$$

$$\Rightarrow x_B = b \geq 0 \quad \text{used the assumption } D$$

$\therefore x_B$ is the BFS.

If I is already a submatrix of A , then we can start the algorithm

Else if not, we find out which columns of I are missing. We force these columns into the system by using additional > 0 variables called artificial variables.

In order to get rid of these variables quickly, we modify the objective funcⁿ

Suppose we aim

$$\max z = c^T x - m \alpha$$

$$\text{s.t. } Ax + \alpha a = b$$

$$x, \alpha \geq 0$$

α : artificial variable

of $\alpha \leq m$

by attaining a high $-ve$ value in the obj. funcⁿ, $M \gg 0$.

Result: $\max z = c^T x$ (we have converted it into eqⁿs = equation forms)
s.t. $x \in S$
(added slack + artificial variable).

$$S = \{Ax = b, x \geq 0\}$$

equivalent to $x \in S = \{x \in \mathbb{R}^n, Ax \leq b, x \geq 0\}$

$$c = (\underbrace{c_1, \dots, c_k}_{\text{original}}, \underbrace{0, \dots, 0}_{\text{slack}}, \underbrace{-m, \dots, -m}_{\text{artificial}})$$

let $x = \begin{pmatrix} x_B \\ 0 \end{pmatrix}$ is the current BFS @ some of iterⁿ of algo.

let $B = [d_1, d_2, \dots, d_m]_{m \times m}$ where each d_i is some column of A matrix.

Let $a_j \in A$ (column of A) which is not in B
i.e. $a_j \notin B$.

$$a_j: m \times 1 \Rightarrow a_j \in \mathbb{R}^m.$$

* $\{d_1, d_2, \dots, d_m\}$ is a basis of \mathbb{R}^m .
 $m \times 1$

We can write

$$a_j = \sum_{i=1}^m y_{ij} d_i, \text{ for some scalars} \\ (\text{unique combination})$$

Assumption: At least one $y_{rj} > 0$ in the combination $1 \leq r \leq n$.

$$\Rightarrow d_r = \sum_{\substack{i=1 \\ i \neq r}}^n \frac{-y_{ij}}{y_{rj}} d_i + \frac{1}{y_{rj}} a_j$$

Substitute this d_r in the eqⁿ

$$B x_B = b.$$

$$\Rightarrow d_1 x_{B_1} + \dots + d_m x_{B_m} = b.$$

$$\sum_{i=1}^m d_i x_{B_i} = b.$$

$$\Rightarrow \sum_{\substack{i=1 \\ i \neq r}}^m d_i x_{B_i} + \left(\sum_{i=1}^m \frac{-y_{ij}}{y_{rj}} d_i + \frac{a_j}{y_{rj}} \right) x_{B_i} = b.$$

$$\Rightarrow \sum_{\substack{i=1 \\ i \neq r}}^m \left(x_{B_i} - \frac{y_{ij}}{y_{rj}} x_{B_r} \right) + \frac{x_{B_r}}{y_{rj}} a_j = b.$$

$$\Rightarrow \sum_{\substack{i=1 \\ i \neq r}}^m \hat{n}_{Bi} d_i + \hat{n}_{Br} a_j = b.$$

\Rightarrow B is changed as a_j enters into B and d_r leaves B.

So, new matrix,

$$B = (d_1, \dots, d_{r-1}, a_j, d_{r+1}, \dots, d_m)_{m \times m}$$