

COL 351:

Analysis and Design of Algorithms

Lecture 26

Polynomial Multiplication

Given: Two polynomials $A(x) = a_0 + a_1x + \dots + a_nx^n$ and $B(x) = b_0 + b_1x + \dots + b_nx^n$, with degree less than equal to 'n' and integer coefficients.

Find: Product $A(x) \cdot B(x) = c_0 + c_1x + c_2x^2 + \dots + c_{2n}x^{2n}$ (Say, $C(x)$)

Example:

If $A(x) = 1 + x + x^2$ and

$B(x) = 1 + 2x + x^3$. Then,

$C(x) = 1 + 3x + 3x^2 + 3x^3 + x^4 + x^5$

$$c_i = \sum_{j=0}^i a_j b_{i-j}$$

Trivial: $O(n^2)$

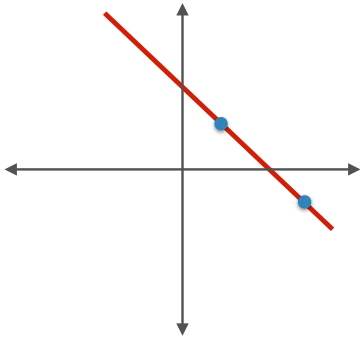
GOAL : $O(n \log n)$

Representation of a polynomial

- An alternate way to represent polynomial $A(x) = a_0 + a_1x + \cdots + a_nx^n$.

$$n = 1$$

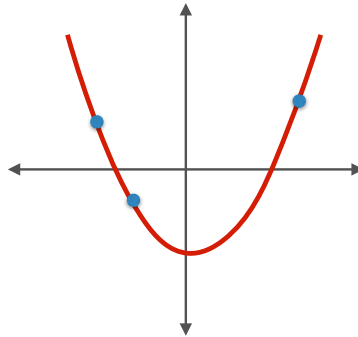
$$A(x) = a_0 + a_1x$$



Evaluation at 2
point suffices

$$n = 2$$

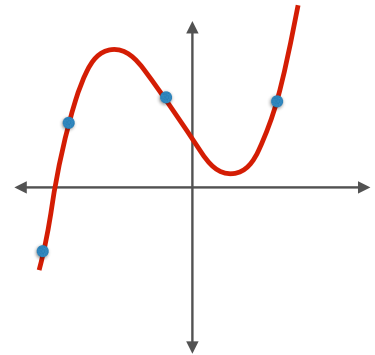
$$A(x) = a_0 + a_1x + a_2x^2$$



Evaluation at 3
point suffices

$$n = 3$$

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3$$



Evaluation at 4
point suffices

Representation of a polynomial

- An alternate way to represent polynomial $A(x) = a_0 + a_1x + \cdots + a_nx^n$.

Lemma: Given $n + 1$ pairs $(x_0, y_0), (x_1, y_1), \dots, (x_n, y_n)$, there exists a unique polynomial (say P) with degree at most n such that $y_i = P(x_i)$, for $i = 0, 1, \dots, n$.

Proof:

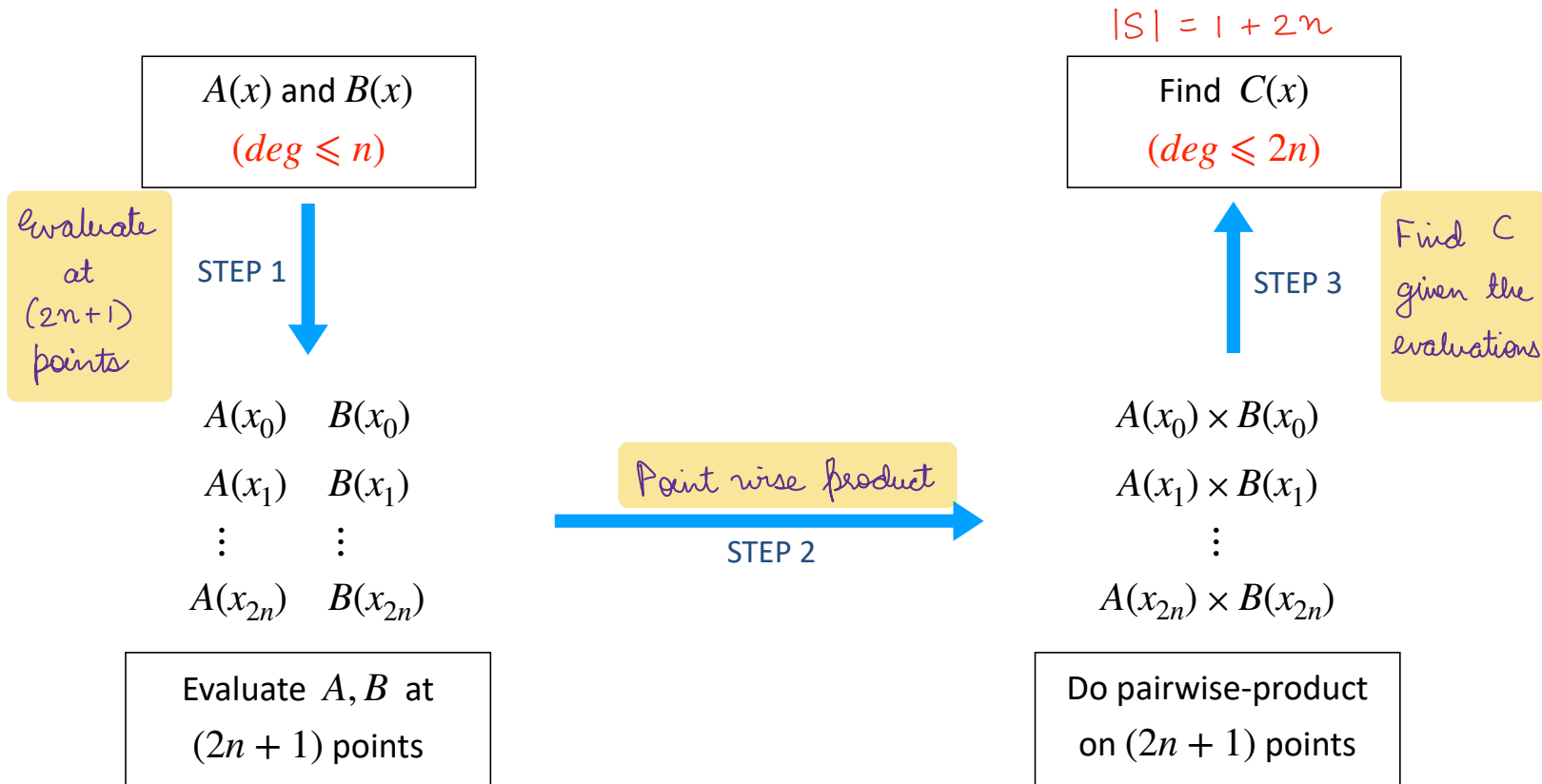
(Hint: A polynomial can be represented as product of monomials in Complex numbers)

1. Suppose P_1, P_2 have same evaluations on x_0, x_1, \dots, x_n .
2. Define $Q := P_1 - P_2$.
3. On the $(n + 1)$ points Q will evaluate to 0, but Q is not identically 0.
4. This is not possible as $\deg(Q) \leq n$.

Why are we looking at alternate representation?

- Answer: Efficient way to compute product.

Take a set $S = \{x_0, x_1, x_2, \dots, x_{2n}\}$



Step 1: Pointwise evaluation

Given: Polynomial 'A' of degree $\leq n$, find its evaluation on a set $S = \{x_0, x_1, \dots, x_n\}$.

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 \dots + a_nx^n$$

$$= (a_0 + a_2x^2 + a_4x^4 + \dots) + x(a_1 + a_3x^2 + a_5x^4 + \dots)$$

$$= A_{\text{even}}(x^2) + x \cdot A_{\text{odd}}(x^2)$$

Assume $N = (n + 1)$
is a power of 2

$|S| \leq N$
 $\text{deg} < N$

Problem
of size N

Remark 1: Degree of polynomials $A_{\text{even}}, A_{\text{odd}} \leq (n - 1)/2 < N/2$.

Remark 2: If $|S^2| \leq N/2$, then we get two subproblems of size $N/2$.

$$S^i = \{x^i \mid x \in S\}$$

Can we say $T(N) = 2T(N/2) + O(N)$?

Only if..

$$|S| = N$$

$$|S^2| = N/2$$

$$|S^4| = N/4$$

$$|S^N| = 1$$

\Rightarrow Set S should be N roots
of $x^N = 1$

$$S^i = \{x^i \mid x \in S\}$$

N-th Roots of Unity

$N = 8$

$$S = \{e^{\frac{2\pi i}{N}} \mid i \in [1, N]\}$$

$$S = \{1, \omega, \omega^2, \omega^3, \dots, \omega^8\}$$

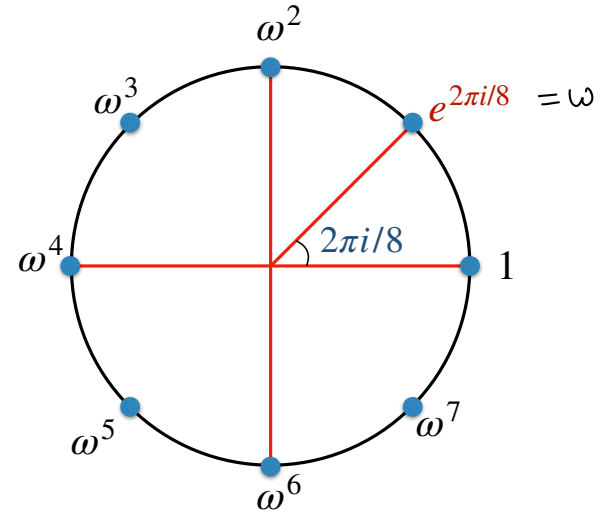
What is:

$$S = \{1, \omega, \omega^2, \dots, \omega^7\}$$

$$S^2 = \{1, \omega^2, \omega^4, \omega^6\} = \{1, i, -1, -i\}$$

$$S^4 = \{1, -1\}$$

$$S^8 = \{1\}$$



$$\omega = e^{2\pi i/8}$$

$$= \cos\left(\frac{2\pi}{8}\right) + i \sin\left(\frac{2\pi}{8}\right)$$

$$S^i = \{x^i \mid x \in S\}$$

How to generate all roots from a single root?

$N=8$

Primitive Root:

An N^{th} root of unity that can generate all other N^{th} roots.

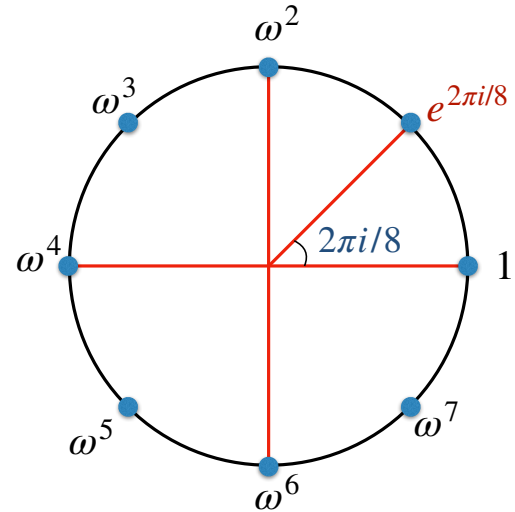
N-th root of unity:

ω such that $\omega^N = 1$

N-th primitive root of unity:

ω such that

- $\omega^N = 1$, and
- $\omega^i \neq 1$, for $0 < i < N$



Homework

Ques 1:

Suppose $\omega = e^{2\pi i/N}$, then list all i for which ω^i is an N^{th} primitive root of unity.

Ques 2:

If ω is N^{th} root of unity other than 1, then show that $1 + \omega + \cdots + \omega^{N-1} = 0$.

Ques 3:

If ω is N^{th} primitive root of unity and $i \in [1, N-1]$, then show that

$$1 + \omega^i + \cdots + \omega^{i(N-1)} = 0.$$