

# Discrete-Time Fourier Series

Lecture 29

DTFS (represented by N complex exponentials)

$$x[n] = \sum_{k=0}^{N-1} a_k e^{jk\Omega_o n}$$

# DTFS by Matrix

$$\begin{bmatrix} x[0] \\ x[1] \\ x[2] \\ x[3] \end{bmatrix} = \begin{bmatrix} e^{\frac{j2\pi}{N}0.0} & e^{\frac{j2\pi}{N}1.0} & e^{\frac{j2\pi}{N}2.0} & e^{\frac{j2\pi}{N}3.0} \\ e^{\frac{j2\pi}{N}0.1} & e^{\frac{j2\pi}{N}1.1} & e^{\frac{j2\pi}{N}2.1} & e^{\frac{j2\pi}{N}3.1} \\ e^{\frac{j2\pi}{N}0.2} & e^{\frac{j2\pi}{N}1.2} & e^{\frac{j2\pi}{N}2.2} & e^{\frac{j2\pi}{N}3.2} \\ e^{\frac{j2\pi}{N}0.3} & e^{\frac{j2\pi}{N}1.3} & e^{\frac{j2\pi}{N}2.3} & e^{\frac{j2\pi}{N}3.3} \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ a_2 \\ a_3 \end{bmatrix}$$

$$\begin{bmatrix} x[0] \\ x[1] \\ x[2] \\ x[3] \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & j & -1 & -j \\ 1 & -1 & 1 & -1 \\ 1 & -j & -1 & j \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ a_2 \\ a_3 \end{bmatrix}$$

# DTFS by Matrix

$$\begin{bmatrix} x[0] \\ x[1] \\ x[2] \\ x[3] \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & j & -1 & -j \\ 1 & -1 & 1 & -1 \\ 1 & -j & -1 & j \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ a_2 \\ a_3 \end{bmatrix}$$

Matrices are inverse of each other

$$\begin{bmatrix} a_0 \\ a_1 \\ a_2 \\ a_3 \end{bmatrix} = \frac{1}{4} \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & -j & -1 & j \\ 1 & -1 & 1 & -1 \\ 1 & j & -1 & -j \end{bmatrix} \begin{bmatrix} x[0] \\ x[1] \\ x[2] \\ x[3] \end{bmatrix}$$

Number of operations increases as  $N^2$

# Fast Fourier “Transform”

Divide FS of length  $2N$  into two of length  $N$  (divide and conquer)

$$\begin{bmatrix} a_0 \\ a_1 \\ a_2 \\ a_3 \end{bmatrix} = \frac{1}{4} \begin{bmatrix} W_4^0 & W_4^0 & W_4^0 & W_4^0 \\ W_4^0 & W_4^1 & W_4^2 & W_4^3 \\ W_4^0 & W_4^2 & W_4^0 & W_4^2 \\ W_4^0 & W_4^3 & W_4^2 & W_4^1 \end{bmatrix} \begin{bmatrix} x[0] \\ x[1] \\ x[2] \\ x[3] \end{bmatrix}$$

where  $W_4^m = e^{-j\frac{2\pi}{4}m}$

Number of Multiplications = 16

Number of Additions =  $4 \times 3 = 12$

Total=28

# Fast Fourier Transform

Divide it into two 2-point series (divide and conquer)

Even-numbered entries in  $x[n]$ :

$$\begin{bmatrix} b_0 \\ b_1 \end{bmatrix} = \begin{bmatrix} W_2^0 & W_2^0 \\ W_2^0 & W_2^1 \end{bmatrix} \begin{bmatrix} x[0] \\ x[2] \end{bmatrix}$$

Odd-numbered entries in  $x[n]$ :

$$\begin{bmatrix} c_0 \\ c_1 \end{bmatrix} = \begin{bmatrix} W_2^0 & W_2^0 \\ W_2^0 & W_2^1 \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix}$$

Number of Multiplications =  $2 \times 4$

Number of Additions =  $2 \times 2 = 4$

Total=12

# Fast Fourier Transform

Break the original 4-point DTFS coefficients  $a_k$  into two parts:

$$a_k = d_k + e_k$$

where  $d_k$  comes from the even-numbered  $x[n]$  and  $e_k$  comes from the odd-numbered  $x[n]$

# Fast Fourier Transform

The 2-point DTFS coefficients  $b_k$  of the even-numbered  $x[n]$

$$\begin{bmatrix} b_0 \\ b_1 \end{bmatrix} = \begin{bmatrix} W_2^0 & W_2^0 \\ W_2^0 & W_2^1 \end{bmatrix} \begin{bmatrix} x[0] \\ x[2] \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \end{bmatrix} \begin{bmatrix} x[0] \\ x[2] \end{bmatrix}$$

$$\begin{bmatrix} d_0 \\ d_1 \\ d_2 \\ d_3 \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 & W_4^0 & W_4^0 \\ W_4^0 & W_4^1 & W_4^2 & W_4^3 \\ W_4^0 & W_4^2 & W_4^0 & W_4^2 \\ W_4^0 & W_4^3 & W_4^2 & W_4^1 \end{bmatrix} \begin{bmatrix} x[0] \\ x[1] \\ x[2] \\ x[3] \end{bmatrix}$$



# Fast Fourier Transform

The 2-point DTFS coefficients  $b_k$  of the even-numbered  $x[n]$

$$\begin{bmatrix} b_0 \\ b_1 \end{bmatrix} = \begin{bmatrix} W_2^0 & W_2^0 \\ W_2^0 & W_2^1 \end{bmatrix} \begin{bmatrix} x[0] \\ x[2] \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \end{bmatrix} \begin{bmatrix} x[0] \\ x[2] \end{bmatrix}$$

$$\begin{bmatrix} d_0 \\ d_1 \\ d_2 \\ d_3 \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \\ W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \end{bmatrix} \begin{bmatrix} x[0] \\ x[2] \end{bmatrix}$$

# Fast Fourier Transform

The 2-point DTFS coefficients  $b_k$  of the even-numbered  $x[n]$

$$\begin{bmatrix} b_0 \\ b_1 \end{bmatrix} = \begin{bmatrix} W_2^0 & W_2^0 \\ W_2^0 & W_2^1 \end{bmatrix} \begin{bmatrix} x[0] \\ x[2] \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \end{bmatrix} \begin{bmatrix} x[0] \\ x[2] \end{bmatrix}$$

$$\begin{bmatrix} d_0 \\ d_1 \\ d_2 \\ d_3 \end{bmatrix} = \begin{bmatrix} b_0 \\ b_1 \\ b_0 \\ b_1 \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \\ W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \end{bmatrix} \begin{bmatrix} x[0] \\ x[2] \end{bmatrix}$$

# Fast Fourier Transform

The 2-point DTFS coefficients  $b_k$  of the even-numbered  $x[n]$

$$\begin{bmatrix} b_0 \\ b_1 \end{bmatrix} = \begin{bmatrix} W_2^0 & W_2^0 \\ W_2^0 & W_2^1 \end{bmatrix} \begin{bmatrix} x[0] \\ x[2] \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \end{bmatrix} \begin{bmatrix} x[0] \\ x[2] \end{bmatrix}$$

$$\boxed{\begin{bmatrix} d_0 \\ d_1 \\ d_2 \\ d_3 \end{bmatrix}} = \begin{bmatrix} b_0 \\ b_1 \\ b_0 \\ b_1 \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \\ W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \end{bmatrix} \begin{bmatrix} x[0] \\ x[2] \end{bmatrix}$$

# Fast Fourier Transform

The 2-point DTFS coefficients  $c_k$  of the odd-numbered  $x[n]$

$$\begin{bmatrix} c_0 \\ c_1 \end{bmatrix} = \begin{bmatrix} W_2^0 & W_2^0 \\ W_2^0 & W_2^1 \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix}$$

$$\begin{bmatrix} e_0 \\ e_1 \\ e_2 \\ e_3 \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 & W_4^0 & W_4^0 \\ W_4^0 & W_4^1 & W_4^2 & W_4^3 \\ W_4^0 & W_4^2 & W_4^0 & W_4^2 \\ W_4^0 & W_4^3 & W_4^2 & W_4^1 \end{bmatrix} \begin{bmatrix} x[0] \\ x[1] \\ x[2] \\ x[3] \end{bmatrix}$$

# Fast Fourier Transform

The 2-point DTFS coefficients  $c_k$  of the odd-numbered  $x[n]$

$$\begin{bmatrix} c_0 \\ c_1 \end{bmatrix} = \begin{bmatrix} W_2^0 & W_2^0 \\ W_2^0 & W_2^1 \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix}$$

$$\begin{bmatrix} e_0 \\ e_1 \\ e_2 \\ e_3 \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^1 & W_4^3 \\ W_4^2 & W_4^2 \\ W_4^3 & W_4^1 \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix}$$

# Fast Fourier Transform

The 2-point DTFS coefficients  $c_k$  of the odd-numbered  $x[n]$

$$\begin{bmatrix} c_0 \\ c_1 \end{bmatrix} = \begin{bmatrix} W_2^0 & W_2^0 \\ W_2^0 & W_2^1 \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix}$$

$$\begin{bmatrix} e_0 \\ e_1 \\ e_2 \\ e_3 \end{bmatrix} = \begin{bmatrix} W_4^0 c_0 \\ W_4^1 c_1 \\ W_4^2 c_0 \\ W_4^3 c_1 \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^1 & W_4^3 \\ W_4^2 & W_4^2 \\ W_4^3 & W_4^1 \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix}$$

# Fast Fourier Transform

The 2-point DTFS coefficients  $c_k$  of the odd-numbered  $x[n]$

$$\begin{bmatrix} c_0 \\ c_1 \end{bmatrix} = \begin{bmatrix} W_2^0 & W_2^0 \\ W_2^0 & W_2^1 \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^0 & W_4^2 \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix}$$

$$\boxed{\begin{bmatrix} e_0 \\ e_1 \\ e_2 \\ e_3 \end{bmatrix} = \begin{bmatrix} W_4^0 c_0 \\ W_4^1 c_1 \\ W_4^2 c_0 \\ W_4^3 c_1 \end{bmatrix}} = \begin{bmatrix} W_4^0 & W_4^0 \\ W_4^1 & W_4^3 \\ W_4^2 & W_4^2 \\ W_4^3 & W_4^1 \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix}$$

# Fast Fourier Transform

Combine  $d_k$  and  $e_k$  to get  $a_k$

$$\begin{bmatrix} a_0 \\ a_1 \\ a_2 \\ a_3 \end{bmatrix} = \begin{bmatrix} d_0 + e_0 \\ d_1 + e_1 \\ d_2 + e_2 \\ d_3 + e_3 \end{bmatrix} = \begin{bmatrix} b_0 \\ b_1 \\ b_0 \\ b_1 \end{bmatrix} + \begin{bmatrix} W_4^0 c_0 \\ W_4^1 c_1 \\ W_4^2 c_0 \\ W_4^3 c_1 \end{bmatrix}$$

FFT procedure:

- Compute  $b_k$  and  $c_k$ :  $2 \times (2 \times 2) = 8$  multiples
- Combine  $a_k = b_k + W_4^k c_k$ : 4 multiples
- Total 12 multiples: fewer than the original 16 multiples
- Total 8 additions: fewer than the original 12 additions



# Scaling of the FFT algorithm

Let  $M(n)$  be the number of multiples to compute n-point FFT

$$M(1) = 0$$

$$M(2) = 2M(1) + 2 = 2$$

$$M(4) = 2M(2) + 4 = 8$$

$$M(8) = 2M(4) + 8 = 24$$

$$M(N) = N \log_2 N$$

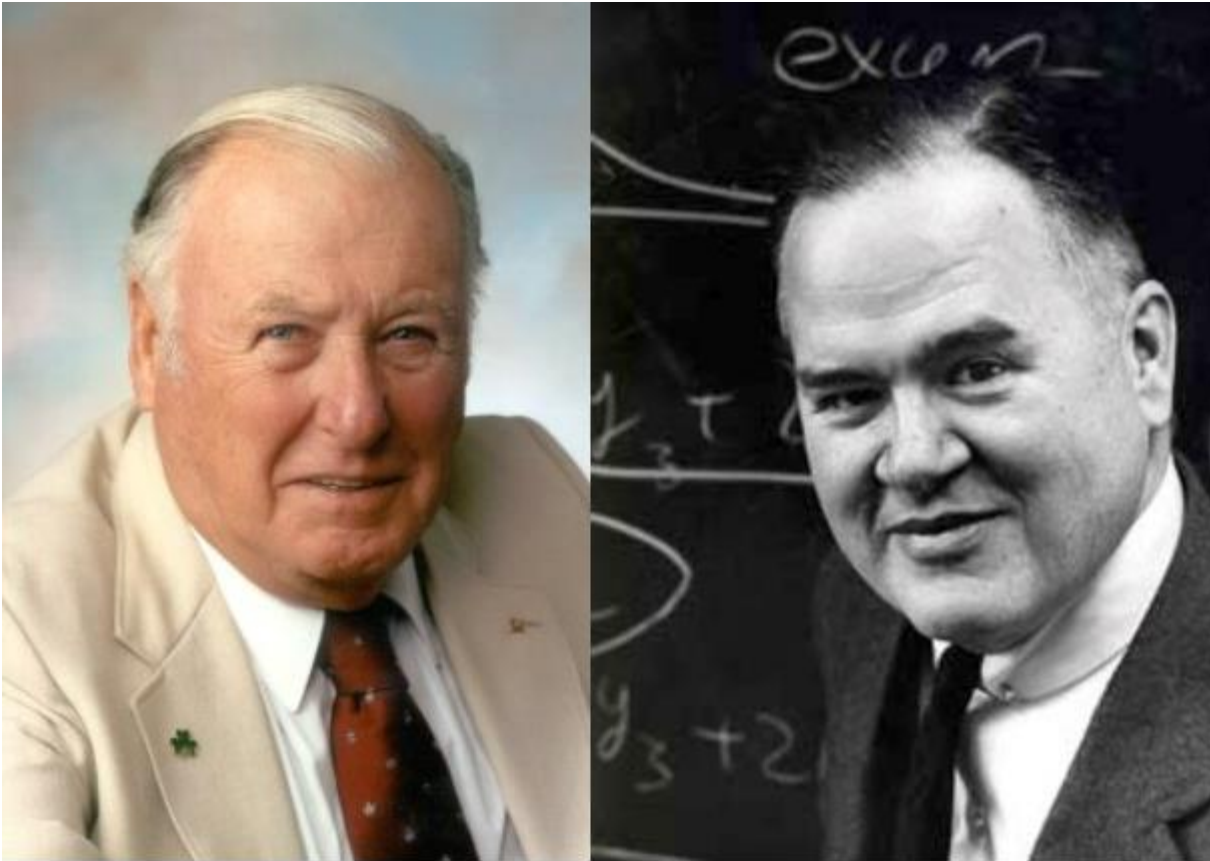
$$N = 1024$$

$$\text{Original} = 1048576$$

$$\text{FFT} = 10240$$

ONE MILLION OPERATIONS REDUCES  
TO TEN THOUSAND

# Cooley and Tukey



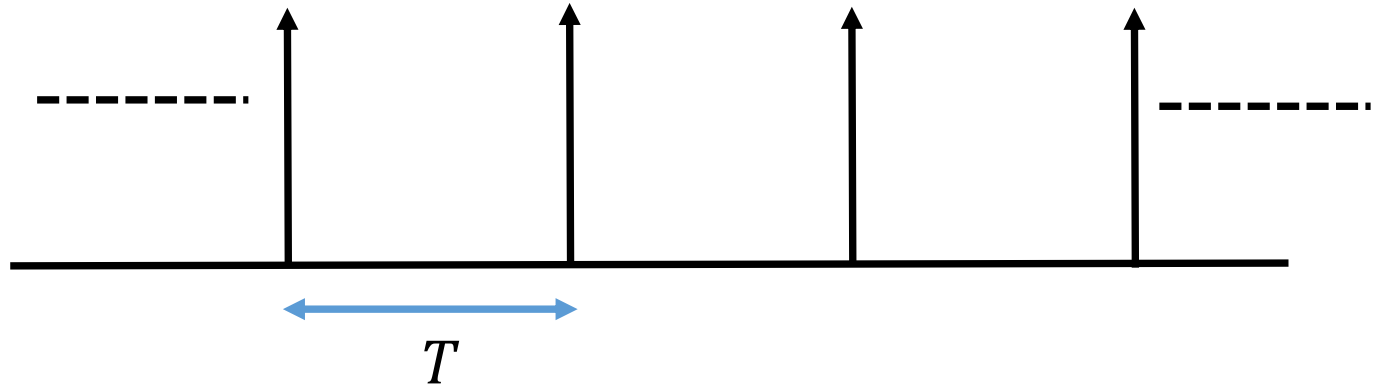
James William Cooley  
(1926-)

John Wilder Tukey  
(1915-2000)

- James Cooley and John Tukey were research scientists at IBM.
- They published their work in 1965 (read it on Google).
- He developed fast Fourier transforms to understand the data from sensors which were planted to detect nuclear explosions.

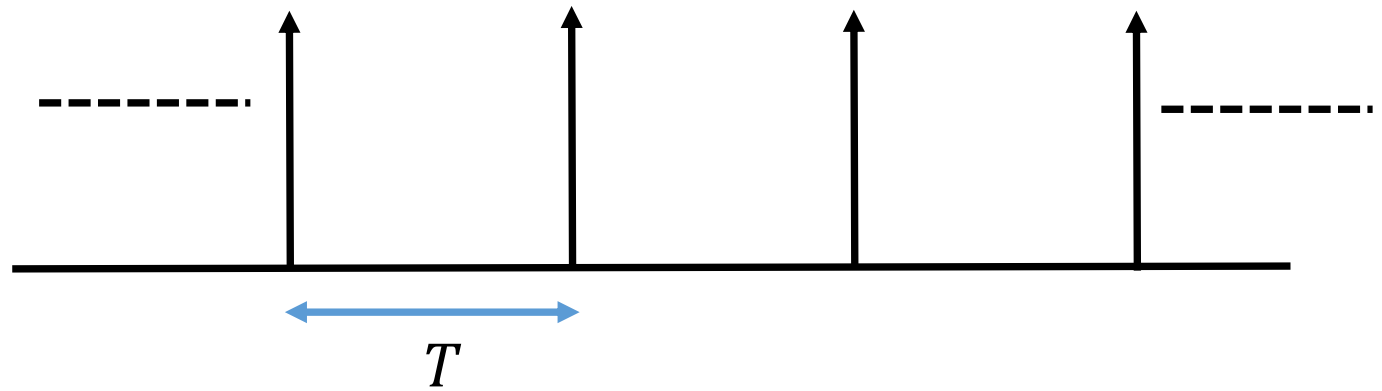
# Picket Fence

$$\sum_{k=-\infty}^{\infty} \delta(t - kT)$$

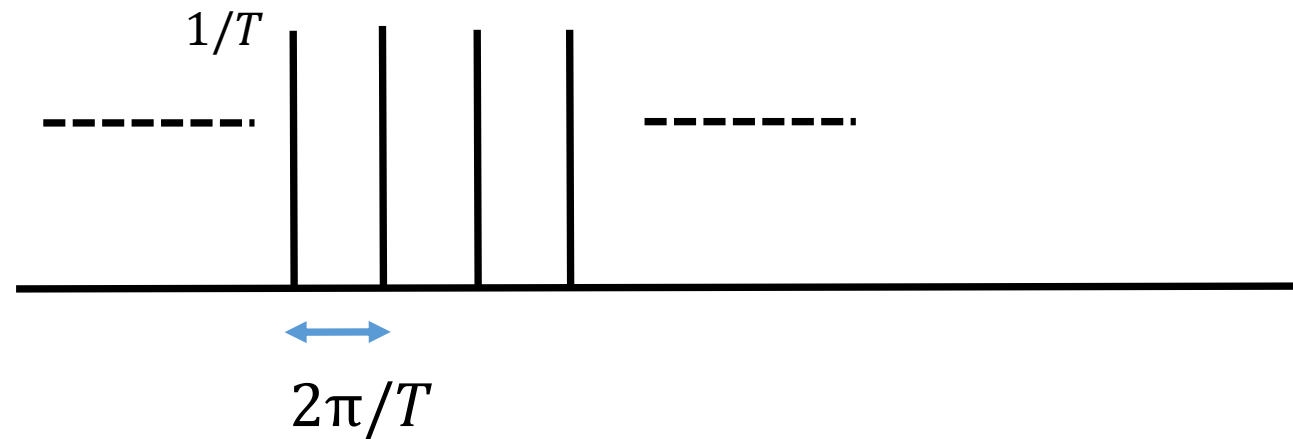


# Picket Fence

$$\sum_{k=-\infty}^{\infty} \delta(t - kT)$$

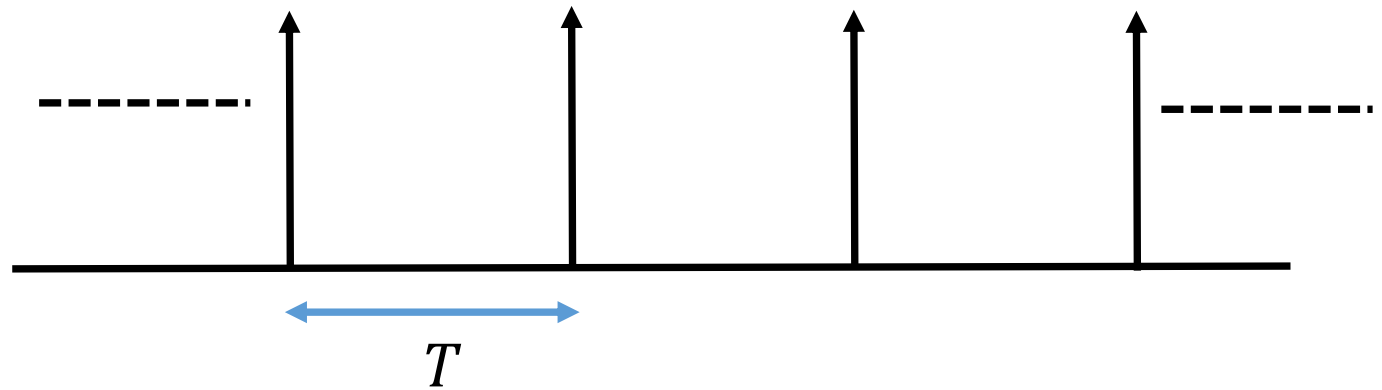


Line  
Spectrum

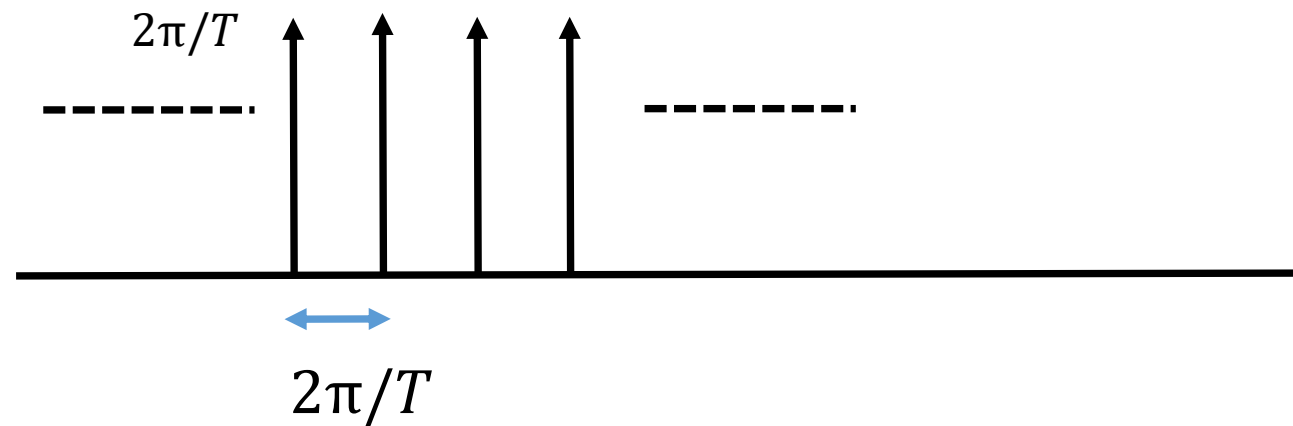


# Picket Fence

$$\sum_{k=-\infty}^{\infty} \delta(t - kT)$$

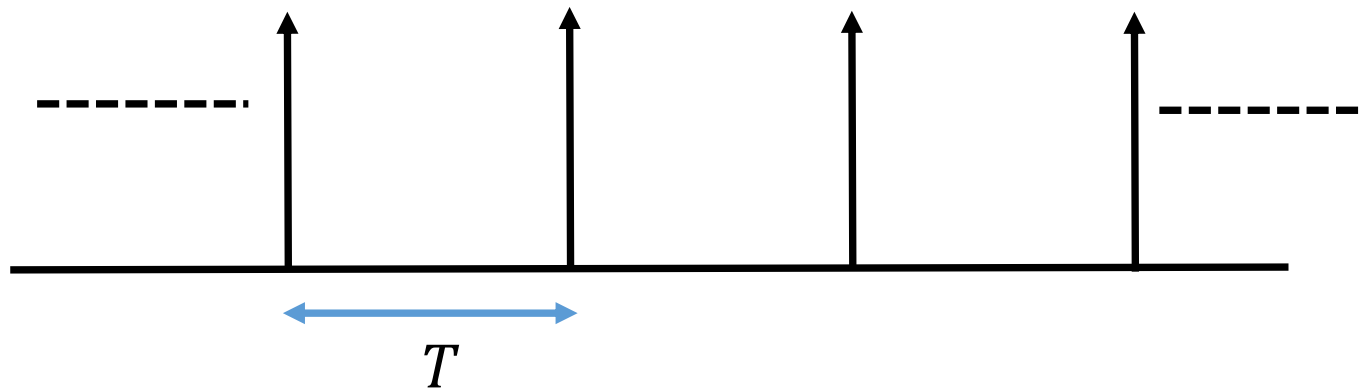


Fourier  
Spectrum

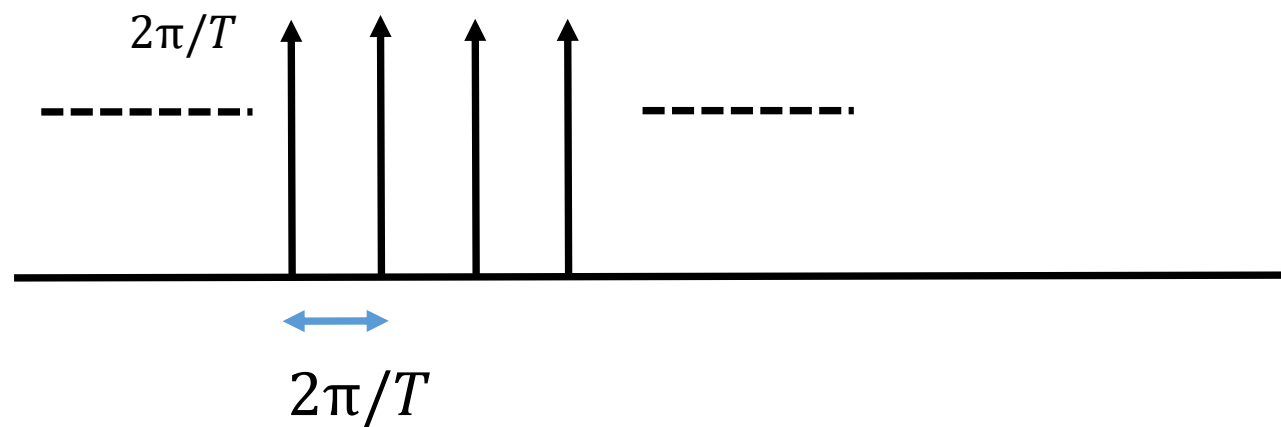


# Picket Fence

$$\sum_{k=-\infty}^{\infty} \delta(t - kT)$$



$$\frac{2\pi}{T} \sum_{k=-\infty}^{\infty} \delta\left(\omega - k \frac{2\pi}{T}\right)$$



# Picket Fence

$$\sum_{k=-\infty}^{\infty} \delta(t - kT) \quad \xleftrightarrow{\mathbb{F}} \quad \frac{2\pi}{T} \sum_{k=-\infty}^{\infty} \delta\left(\omega - k \frac{2\pi}{T}\right)$$

**Picket Fence**

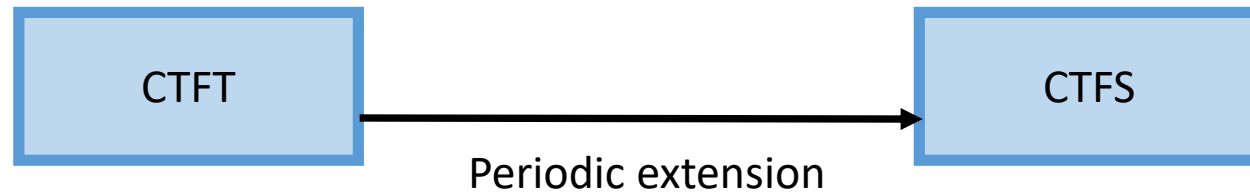
**Picket Fence**

# Relations between Fourier Transforms

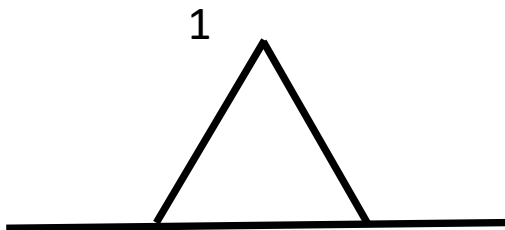




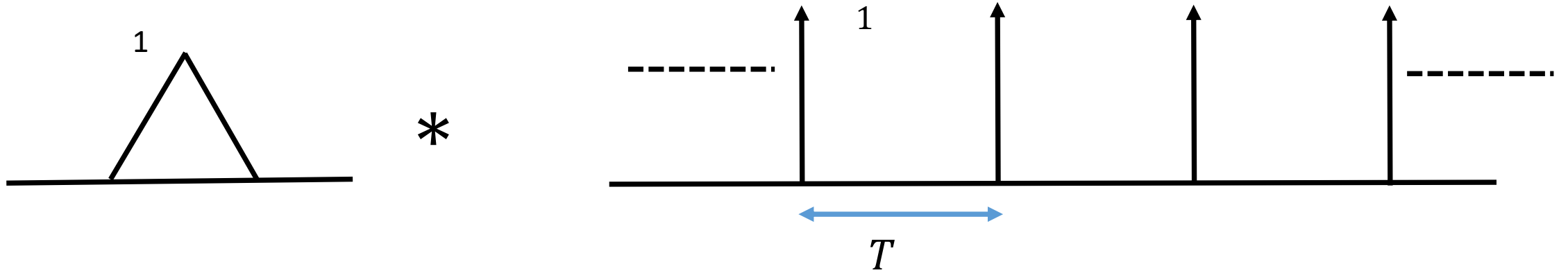
# Relations between Fourier Transforms



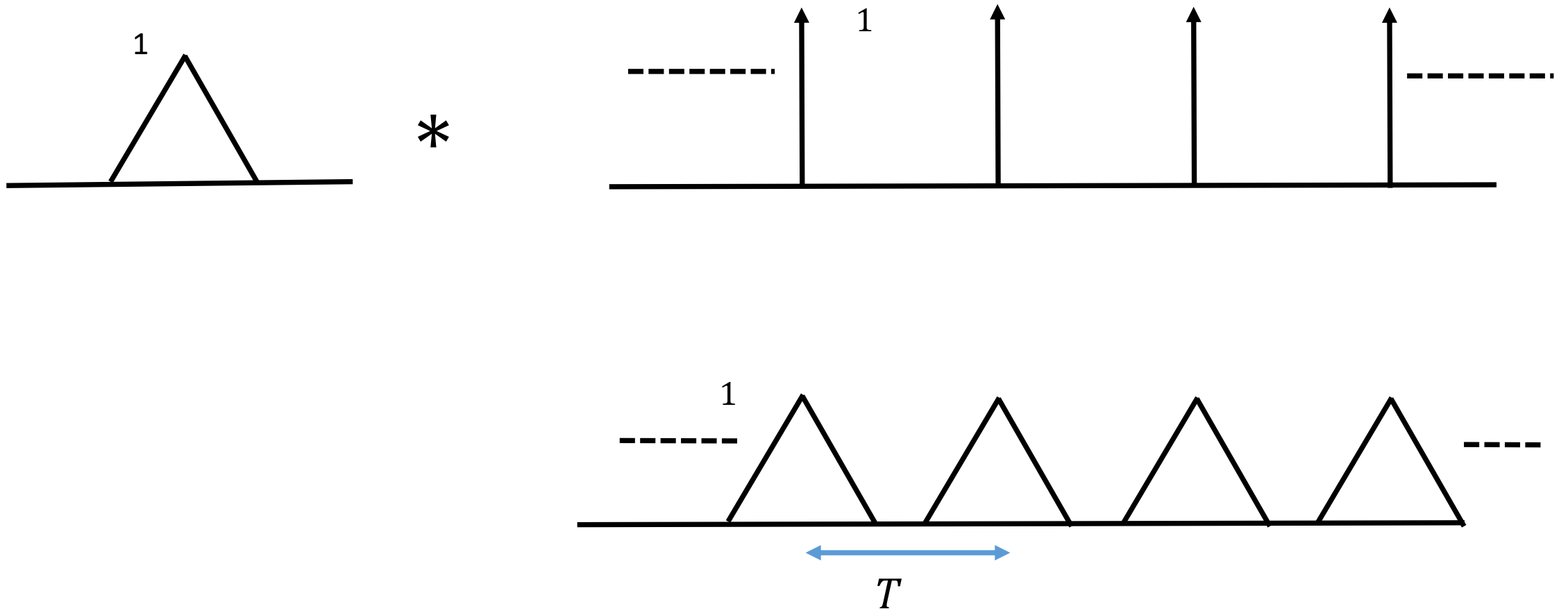
# Periodic extension



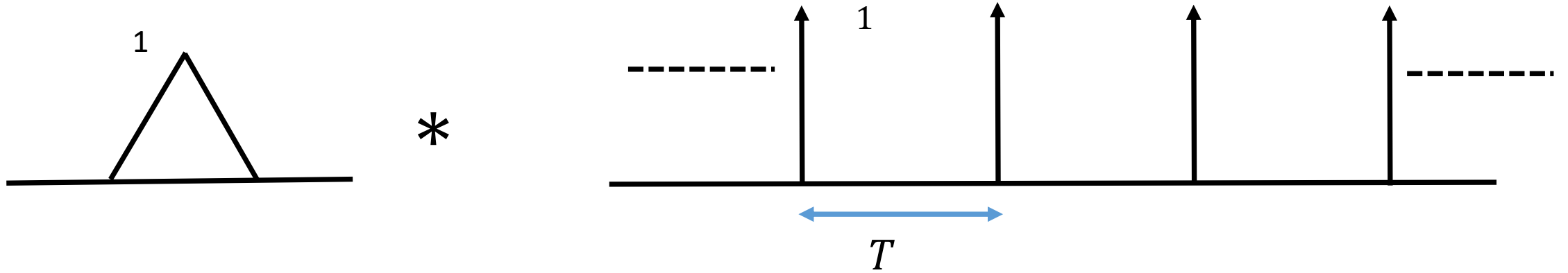
# Periodic extension



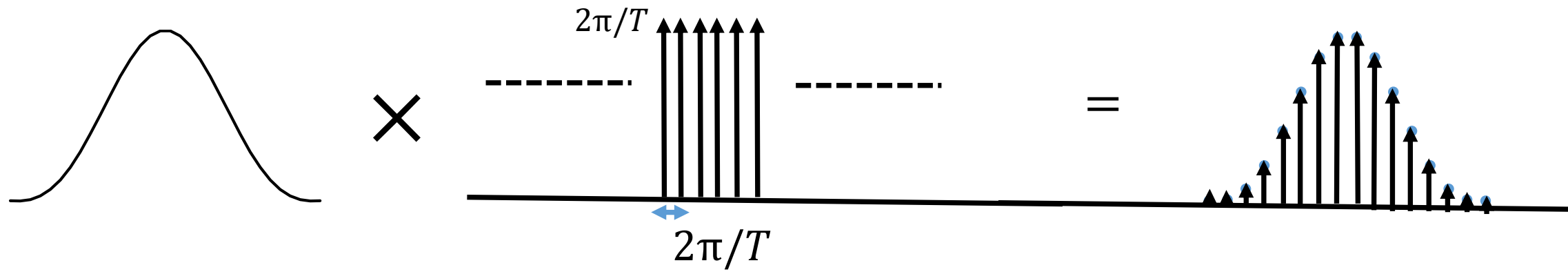
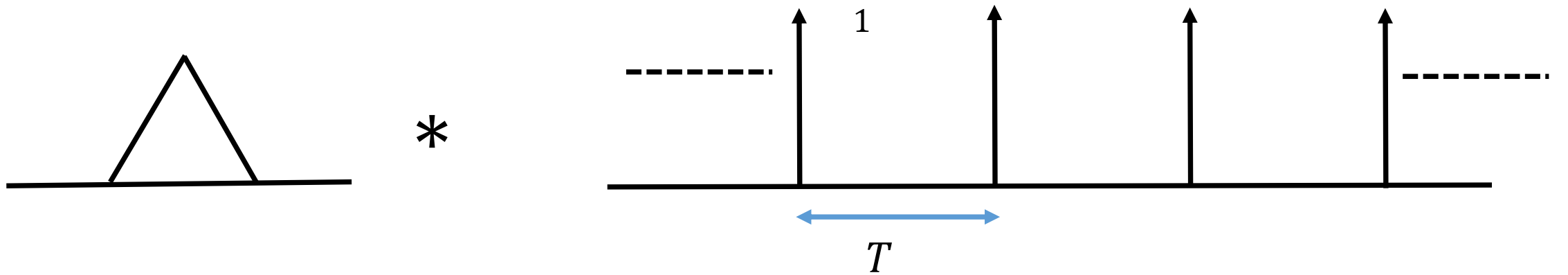
# Periodic extension



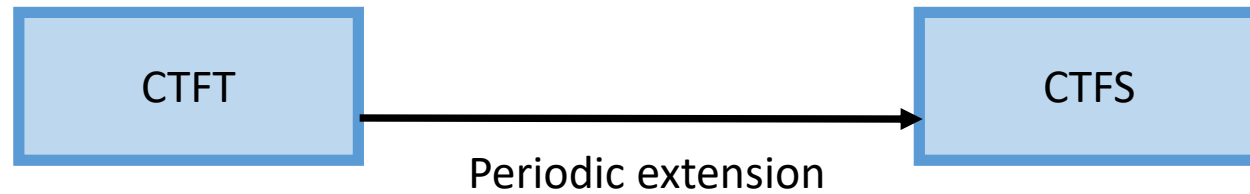
# Periodic extension



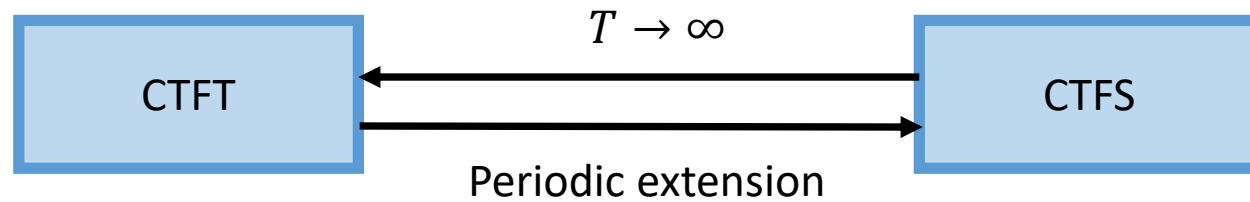
# Periodic extension



# Relations between Fourier Transforms

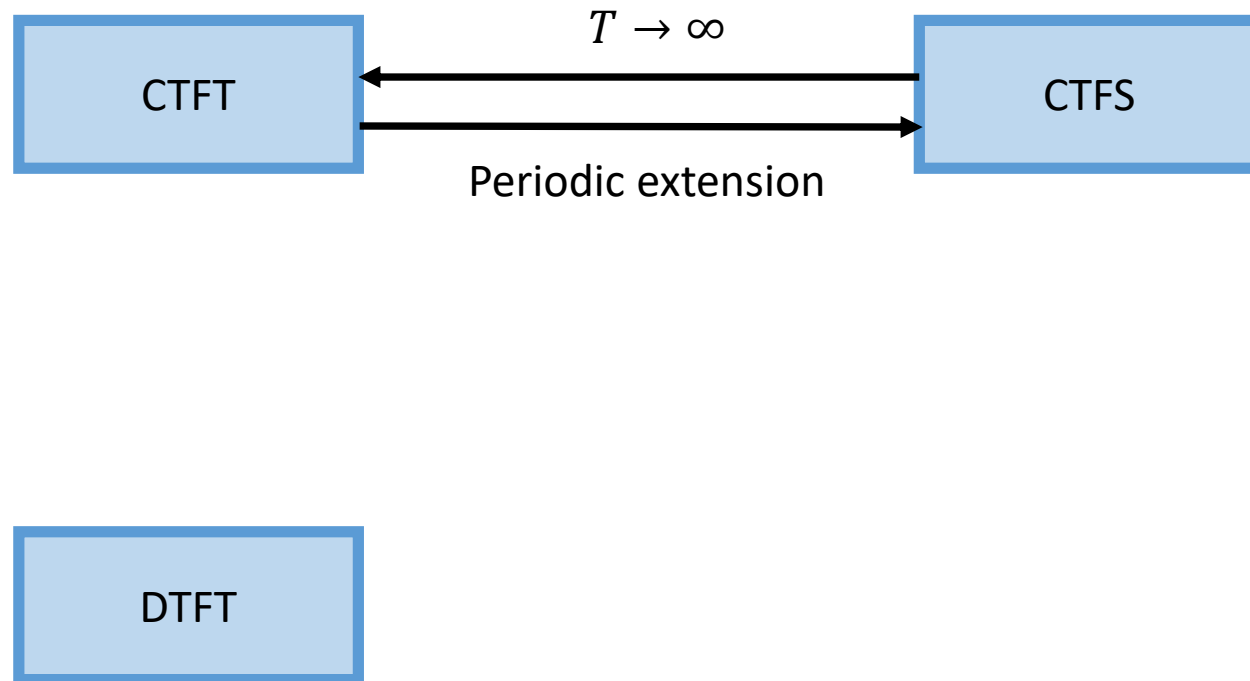


# Relations between Fourier Transforms

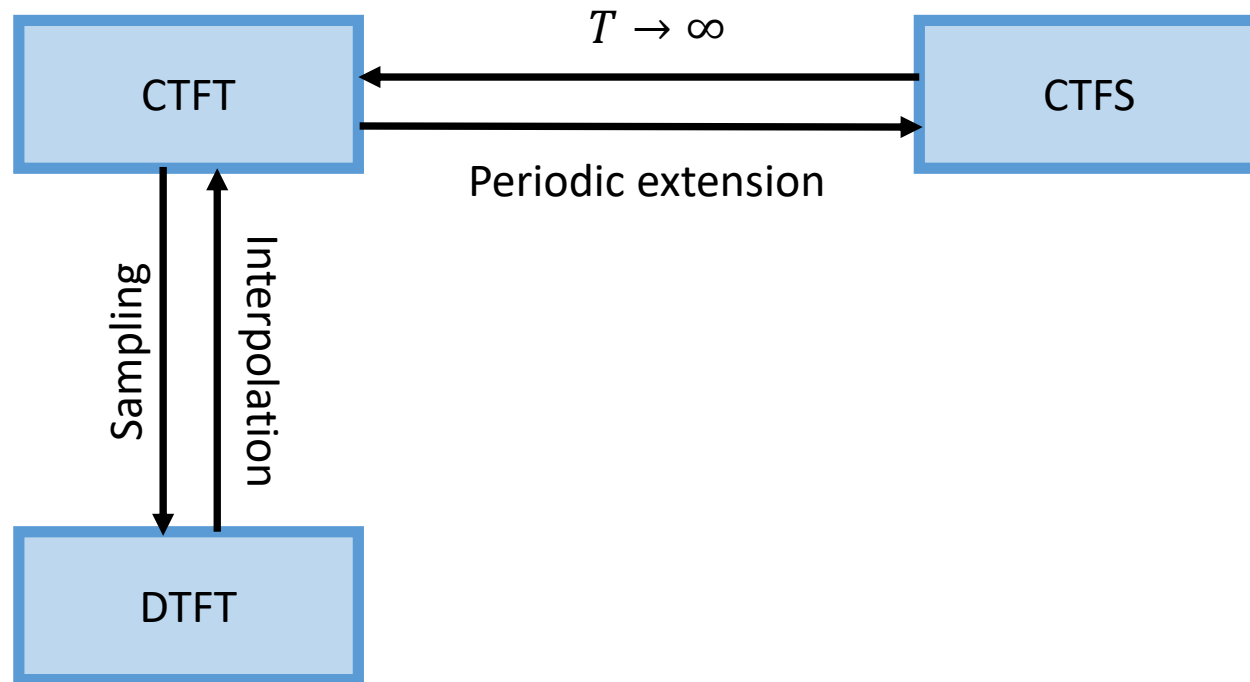




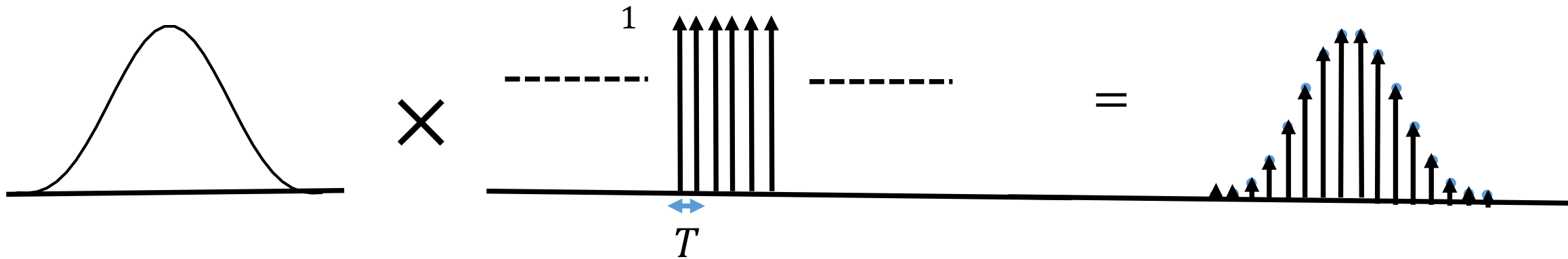
# Relations between Fourier Transforms



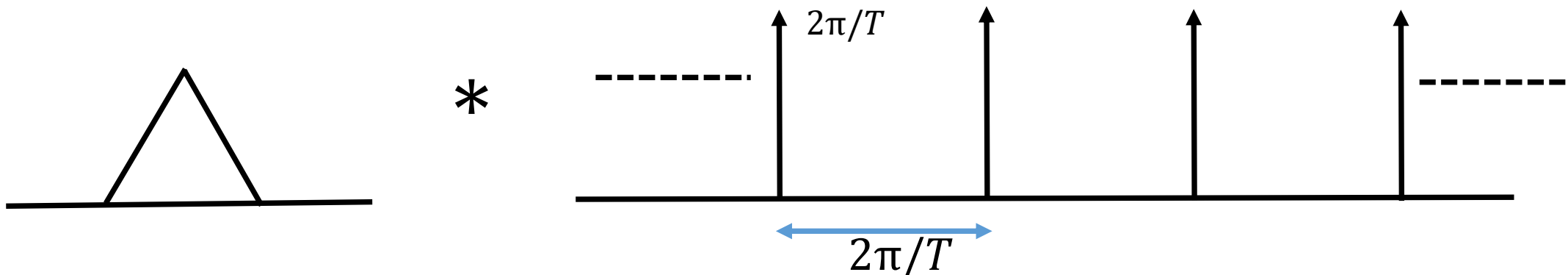
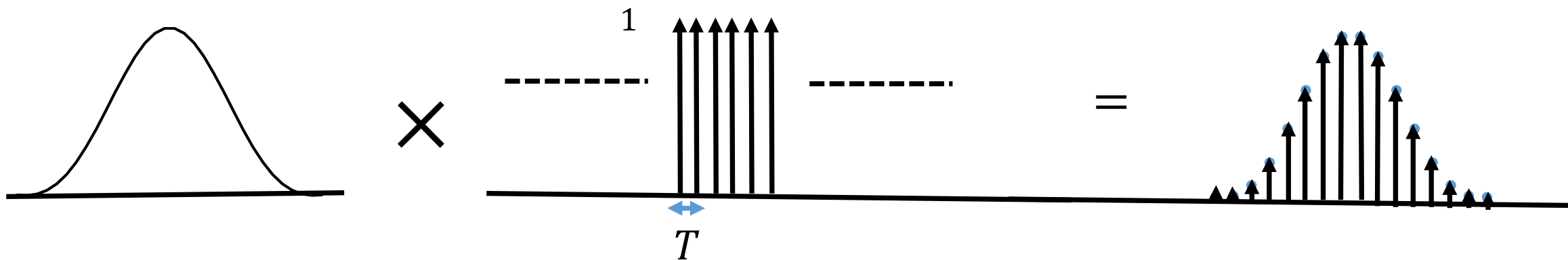
# Relations between Fourier Transforms



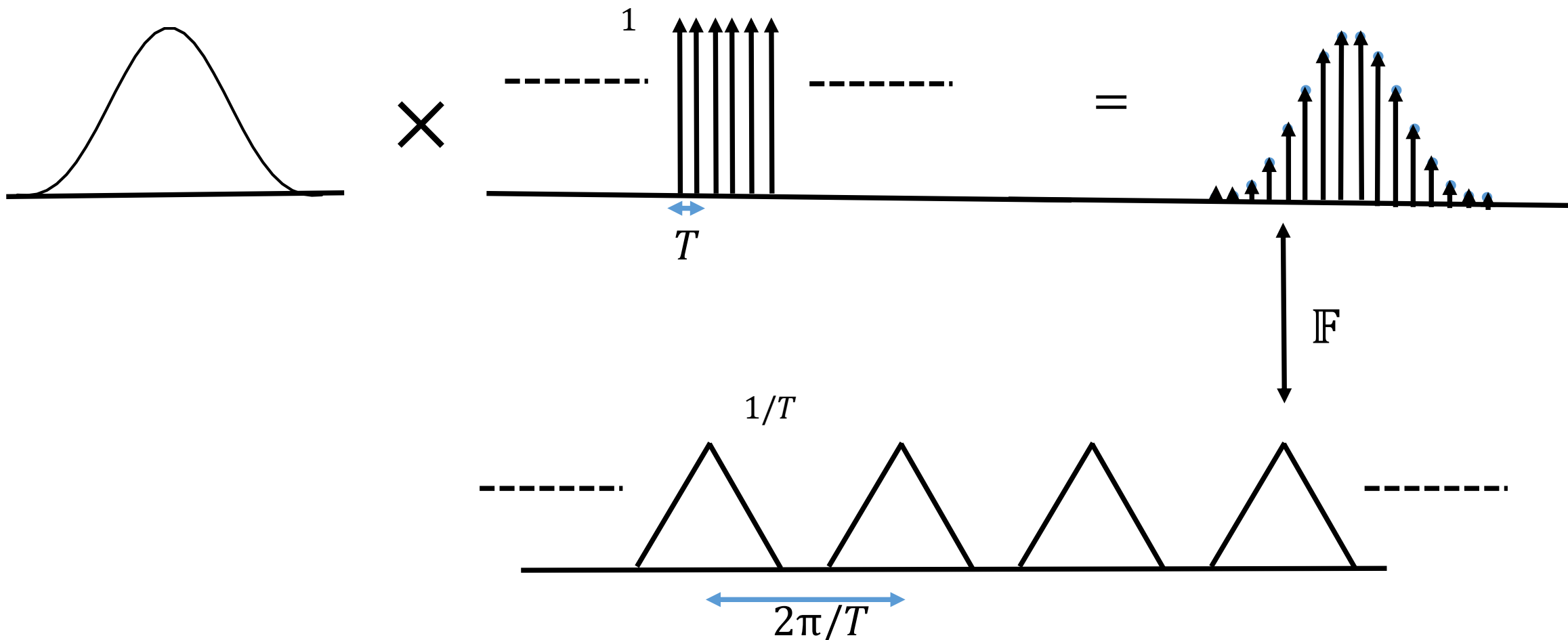
# Sampling



# Sampling

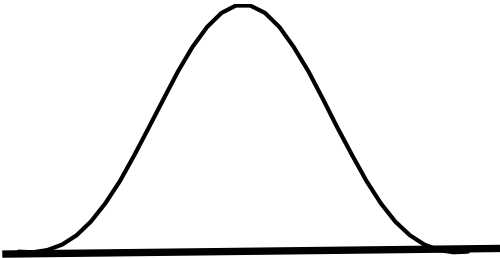


# Sampling



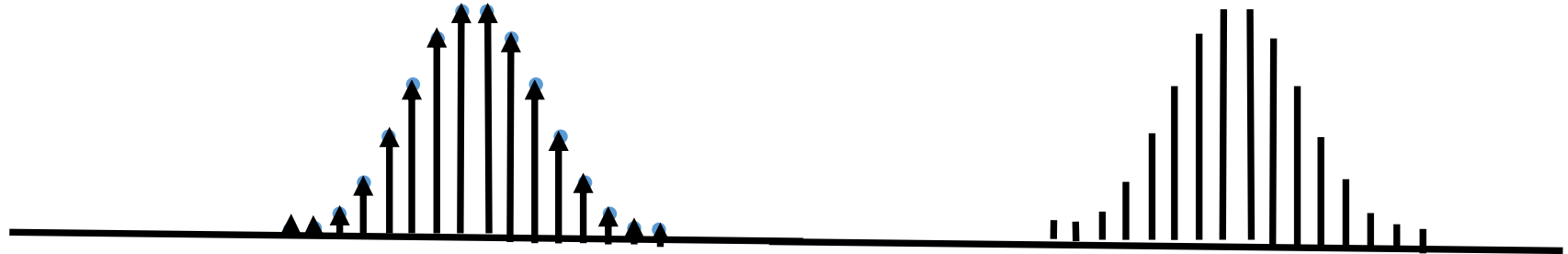
# Sampling

$x(t)$



$\hat{X}(\omega)$

$$\sum_k x(kT)\delta(t - kT)$$



$x[n] = x(nT)$

$$X(\omega) = \sum_k \frac{1}{T} \hat{X}\left(\omega - k \frac{2\pi}{T}\right)$$

$X(e^{j\Omega})$

# Relationship between Fourier Transforms

$$X(\omega) = \int_{-\infty}^{\infty} x_d(t) e^{-j\omega t} dt$$

$$X(\omega) = \sum_n x[n] e^{-j\omega nT} \int_{-\infty}^{\infty} \delta(t - nT) dt$$

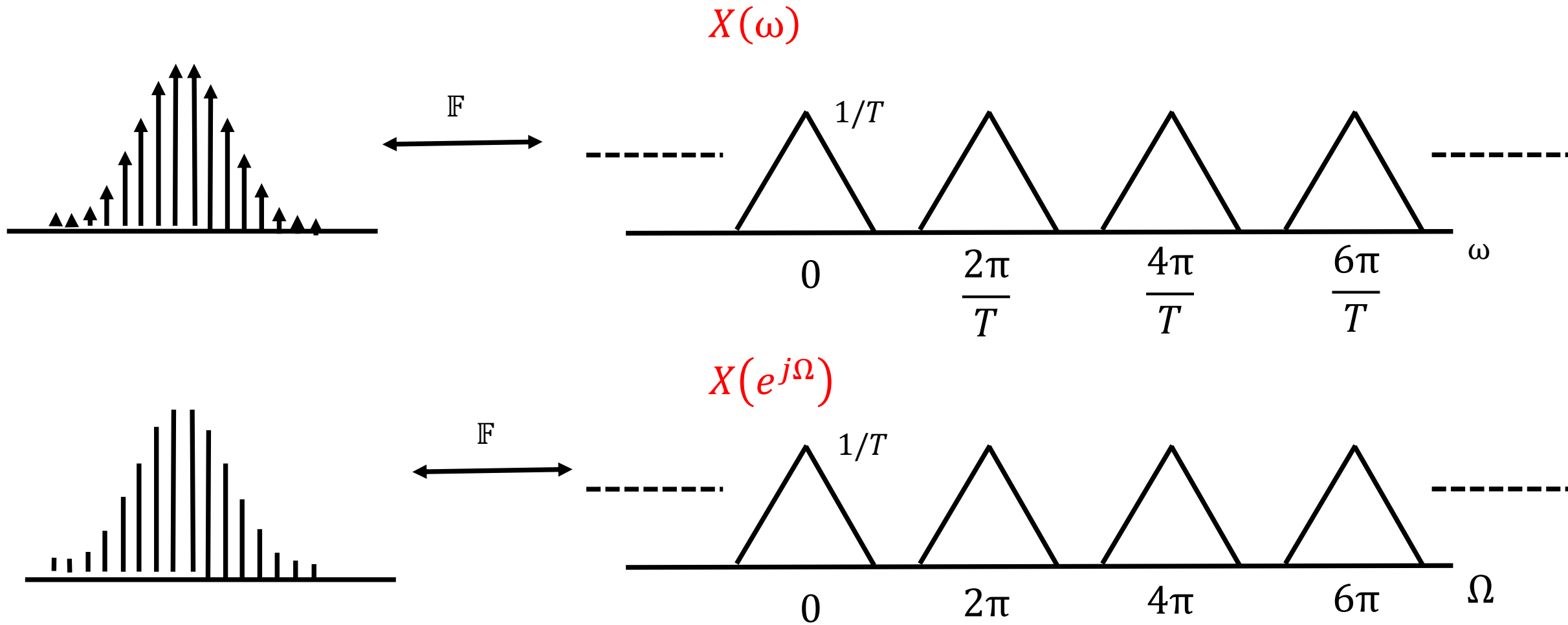
$$X(\omega) = \int_{-\infty}^{\infty} \sum_k x(kT) \delta(t - kT) e^{-j\omega t} dt$$

$$X(\omega) = \sum_n x[n] e^{-j\omega nT} = X(e^{j\Omega}) \Big|_{\Omega=\omega T}$$

$$X(\omega) = \int_{-\infty}^{\infty} \sum_n x[n] \delta(t - nT) e^{-j\omega t} dt$$

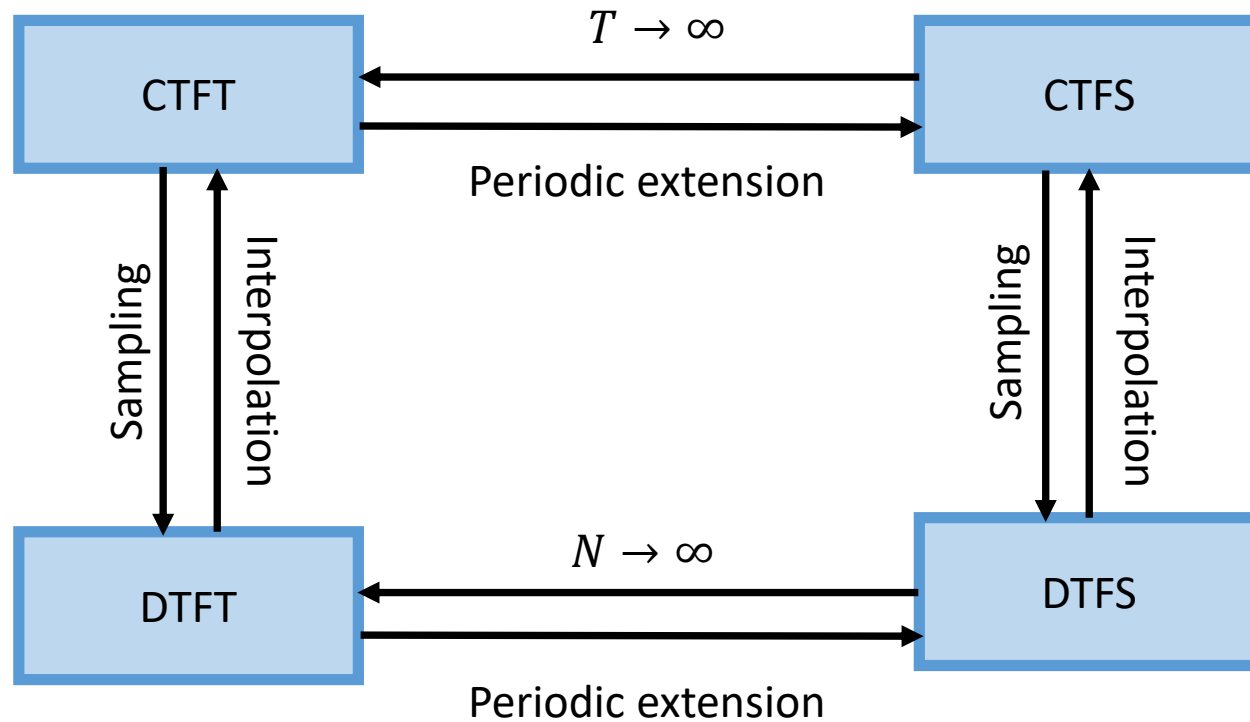
$$X(\omega) = \sum_n x[n] \int_{-\infty}^{\infty} \delta(t - nT) e^{-j\omega t} dt$$

# Relationship between Fourier Transforms

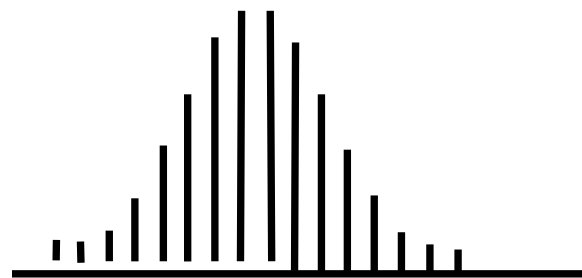




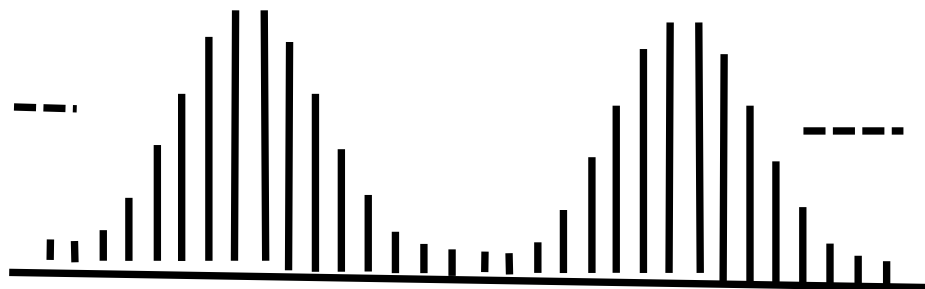
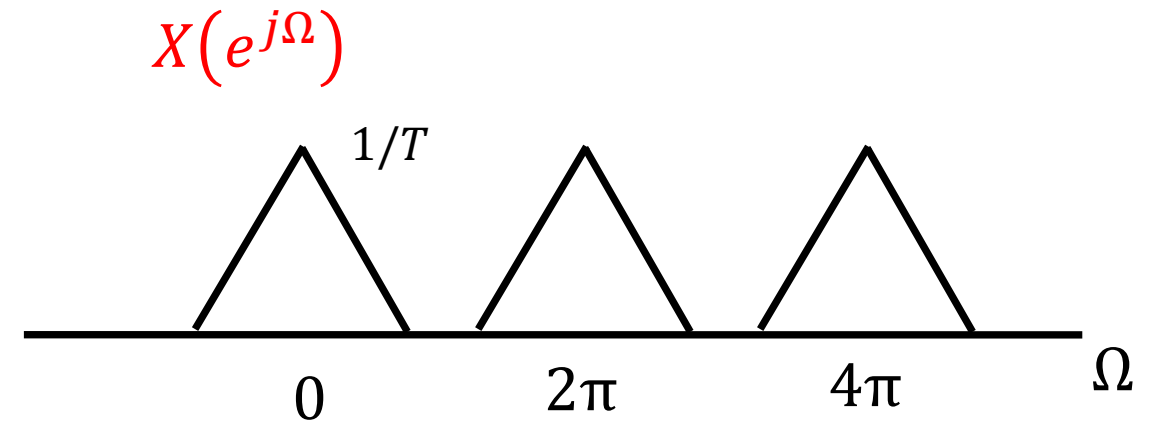
# Relations between Fourier Transforms



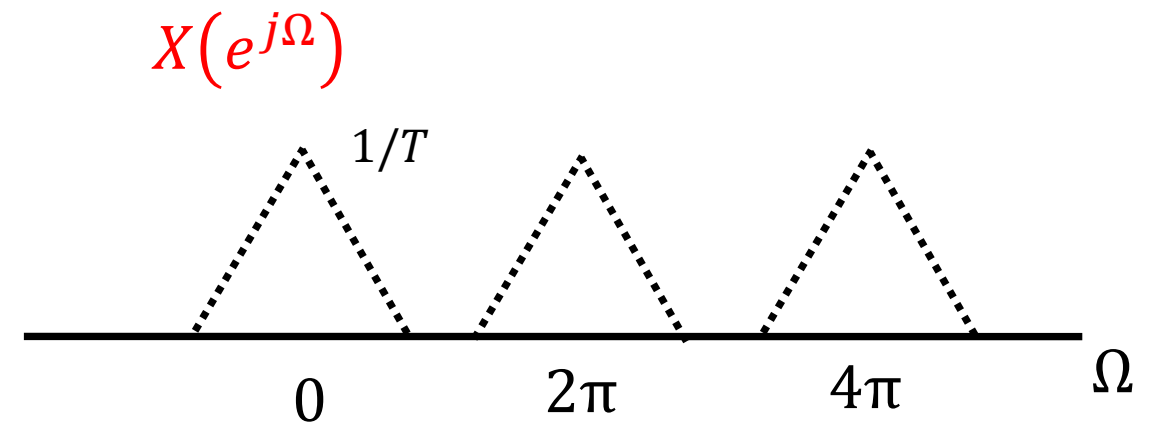
# Relationship between Fourier Transforms



$\mathbb{F}$



$\mathbb{F}$



# Summary

- DFT is most useful
- FFT is the most used computer algorithm