

Problem 1 : Algorithm

Let the two stacks be S_1 and S_2 and the n members in the queue be a_1, a_2, \dots, a_n .

What we do is that first we insert a_1 into stack S_1 .

- 1) Then for ~~for each~~ each a_i where $i \geq 1$ first we check that if the topmost element in S_2 (let it be $a_k, k < i$) is larger than the a_i .
 if $a_k \geq a_i$ we add a_i to ~~top of~~ S_1 , else if $a_k < a_i$, we ~~pop~~ ^{pop} a_k from S_2 and insert it into S_1 .

Then again we check if the ~~next~~ ^{same} ~~element~~ for next element of S_2 with a_i .

- 2) Now if a_i is smaller than topmost element of S_2 , we check if it is larger than topmost element of S_1 . If it is, ~~smaller~~ then we add a_i to S_1 , else we pop elements from S_1 and keep adding them to S_2 until we find an element which is smaller than a_i .

3) finally, when we ~~are~~ have done this for all x_i $1 \leq i \leq n$, we pop elements from S_2 and add them to S_1 till S_2 is empty and we get a sorted stack S_1 with minimum element at bottom.

The time complexity for this algorithm would be $\underline{n^2}$ since for each step we might have to shift the entire i sized stack from 1 stack to another and we have to do this n times.

~~Problem~~

Narshit Mandaria

Problem 2.1

Marshall Newman

For function $\text{Int overflow}(x, y)$ where

we first use the function $\text{Int sum}(x, y)$

and then use the function $\text{Int less than equals to}$

$\text{Int less than equals to}(x, y)$
Let the result of Int sum be (z)

We now try $\text{Int less than equals to}(z, p)$

if this returns false (0) we return 1

if this returns true (1), we try

~~if~~ $\text{Int is equal}(z, p)$. If this returns 1
we return 1, else we return 0.

Problem 2.2.

Harshit Mawandha

For Int Base (p) - Create (a).

First we use the function from Problem 2.1

~~Int mod (p) - overflow(a, 0).~~

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If this returns 1, we use

~~Int mod (a, p)~~ and store the result
in ~~Int + list~~, then we use

~~Int - div (a, p)~~ and ~~store the~~ then restore the
~~step of first step~~ ~~Int mod~~ step 1.

now we use ~~Int list - add front~~ to add
this result to the front of the list.

2) Now when overflow returns 0, we just
add the remaining value to the front of
the list.

Problem 2.3

We first ~~use~~ ~~two~~ ~~Int.Base(p)~~
~~use~~ ~~make~~ ~~two~~ length of the lists to be n
 n

now for $i = 0, i \leq n$,

we iterate.

For any i^{th} element we ~~find~~ ^{have}
 a carry (say k) if $i \geq 1$, $k \geq 0$
 and element at position (i)

~~we~~ we use $\text{Int-sum}(k, a_i, b_i)$
 (a_i, b_i) be the i^{th} elements of 2 lists

~~Then~~ Then we use $\text{Int-mod}(\text{Int-sum}(k, a_i, b_i))$
 which will be our i^{th} element of (P)

~~and that answer~~
 and $\text{Int-div}(\text{Int-sum}(k, a_i, b_i), P)$
 which will be our k_{i+1}

Now when $i > n$ we just use
 k_{i+1} with a_i (the element of larger list)

and take rest b_i 's to be 0
 to get our final $\text{Int-mod}(P)$ list.