

# COL215L: Digital Logic & System Design

## Lecture 10: Binary Arithmetic



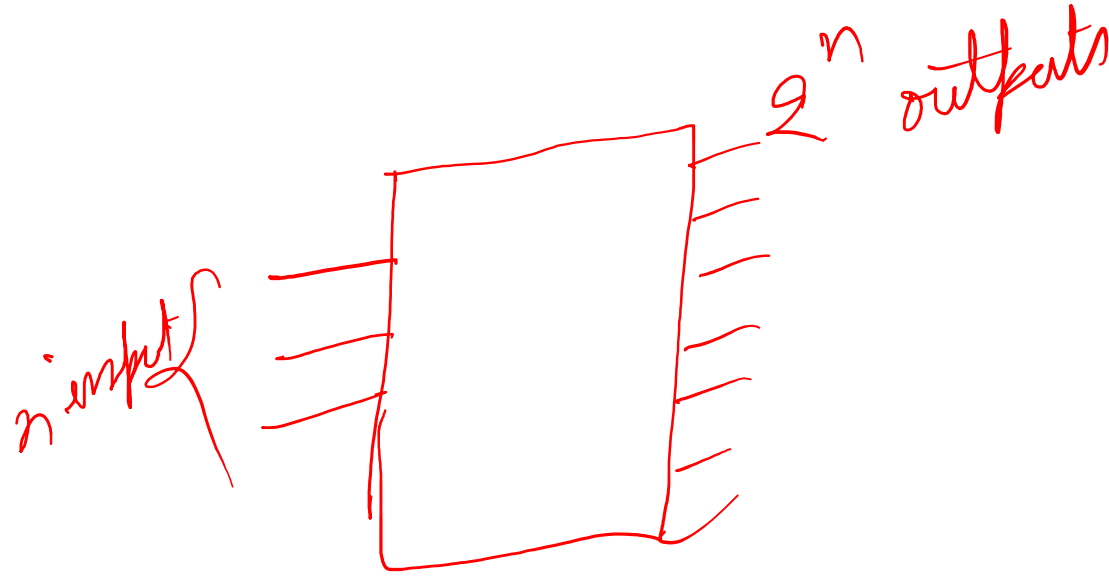
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# Decoder

- $n-2^n$



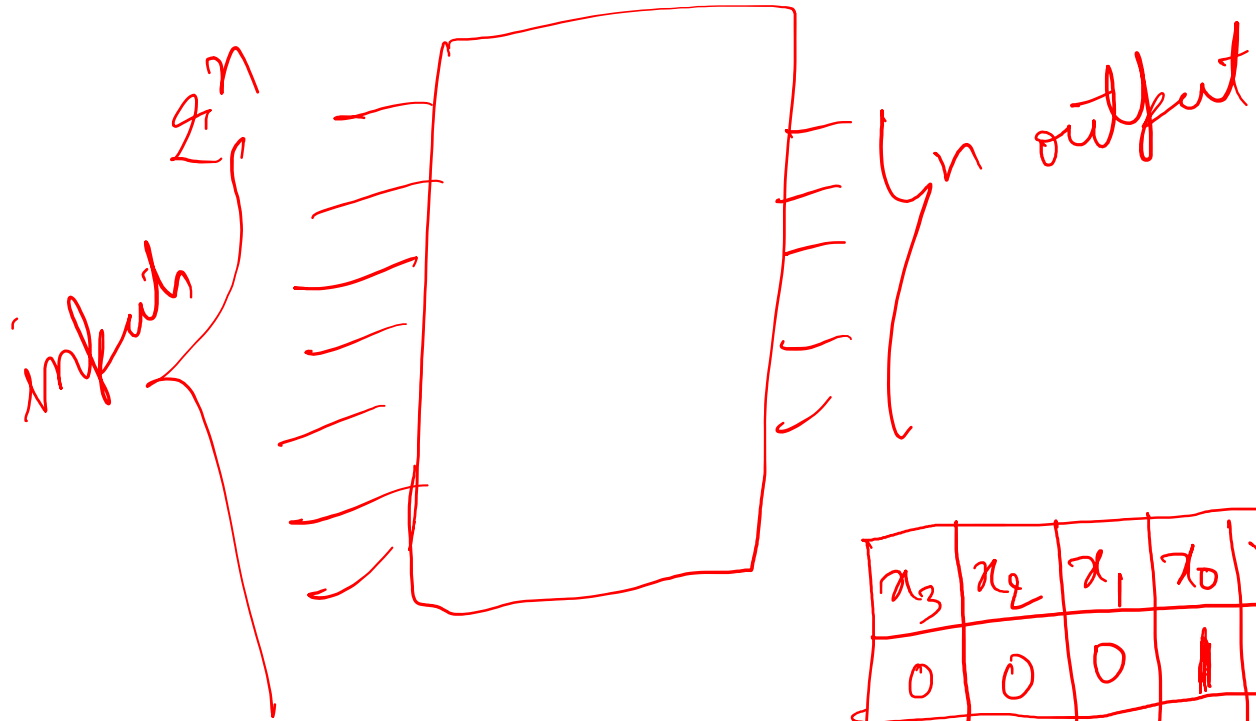
# Encoder

- $2^n - n$

• 4-to-2 Encoder

$$y_1 = ($$

$$y_0 = ($$



$x_3$	$x_2$	$x_1$	$x_0$	$y_1$	$y_0$
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1

# Priority Encoder

- $2^n - n$
- Inputs have priorities

4 inputs  $x_3 > x_2 > x_1 > x_0$

$$Z = (x_3 + x_2 + x_1 + x_0)$$

$$y_1 = \left( \begin{array}{c} 1 \\ 1 \\ 1 \\ 1 \end{array} \right)$$

$$y_0 = \left( \begin{array}{c} 1 \\ 1 \\ 1 \\ 1 \end{array} \right)$$

inputs				outputs		
$x_3$	$x_2$	$x_1$	$x_0$	$y_1$	$y_0$	$Z$
0	0	0	0	X	X	0
0	0	0	1	0	0	1
0	0	1	X	0	1	1
0	1	X	X	1	0	1
1	X	X	X	1	1	1

# Half Adder

$(x), (y)$		$(s)$ <sup>sum</sup>	$(c)$ <sup>carry</sup>
0	0	1	1
0	1	0	1
1	0	1	0
1	1	0	1

x	y	s	c
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

$$s = x \oplus y$$

$$c = xy$$

$$s = x \oplus y$$

→ Unsigned integer

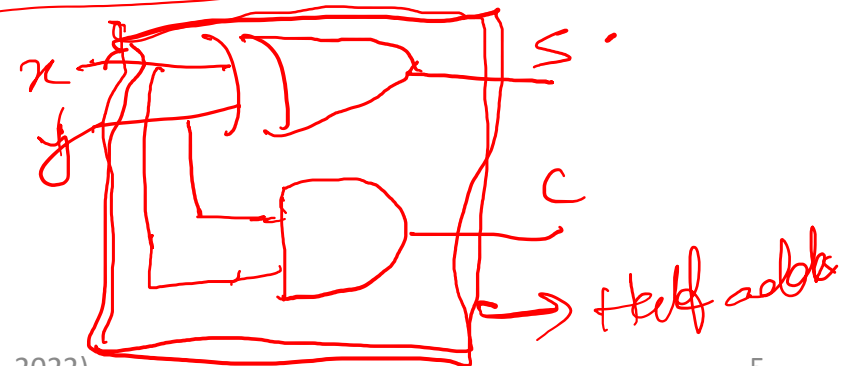
Signed integer

$$(X)_B = x_{n-1} x_{n-2} \dots x_0$$

$$(X)_S = x_{n-1} \cdot 2^{n-1} + x_{n-2} \cdot 2^{n-2} + \dots + x_0 \cdot 2^0$$

$$(X)_S = 15$$

$$\rightarrow (X)_B = 1111$$



# Full Adder

$C_i$	$x_i$	$y_i$	$C_{i+1}$	$S_i$
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	1	0
1	0	1	1	1
1	1	0	1	1
1	1	1	1	0

$$C_{i+1} = x_i y_i + x_i C_i + y_i C_i$$


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$$S_i = x_i \oplus y_i \oplus C_i$$

$$x_{n-1} x_{n-2} \dots x_{i+1} x_i \dots x_0$$

$$y_{n-1} y_{n-2} \dots y_{i+1} y_i \dots y_0$$

$$S_{n-1} S_{n-2} \dots S_{i+1} S_i \dots S_0$$

$$C_{n-2} \dots C_{i+1} C_i \dots C_0$$

