

Ingredients of TQC

Physics : Topological phases of Matter (2D anyons)
(TPMs)

Defn: A system is in topological phase if at low energies, long distances (length), the effective field theory is a Topological Quantum Field Theory.
(due to Nayak, ---)

Bosonic (2+1) TQFTs \longleftrightarrow Modular \otimes categories

State spaces? $\mathcal{H}(M, \ell)$

where $\ell \subseteq \mathbb{Z} = \{0, 1, \dots, k-1\}$

we have an involution $*: \ell \rightarrow \ell$ s.t. $0^* = 0$.

- \mathcal{H} is a 2D topological modular functor.

Recall, we had defined

$$N_{ab}^c = \dim(\mathcal{H} \left(\begin{array}{c} a \\ b \\ \backslash \\ c \end{array} \right))$$
 these have symmetries

It helps: if $N_{ab}^c \in \{0, 1\}$

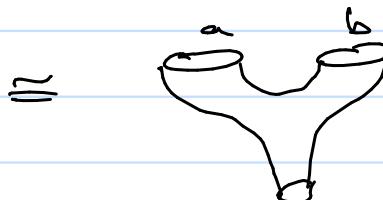
: also if $a^* = a$

(these don't always happen)

Input data: $\{N_{ab}^c\}_{(a,b,c) \in \ell^3}$

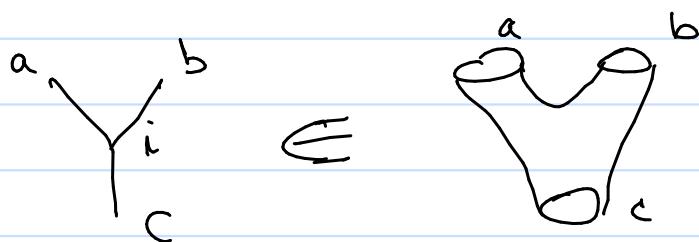
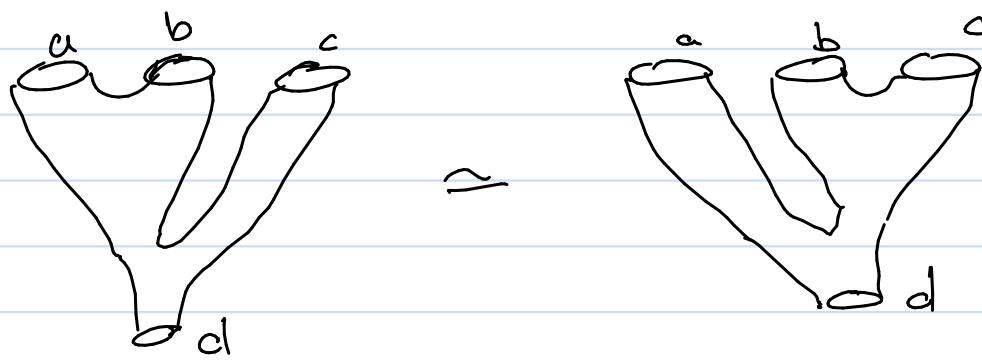
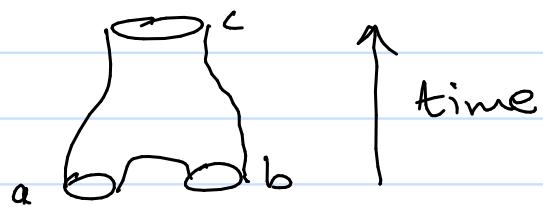


\simeq



Flat surface at low energy
a, b anyons. They can swim

Things of anyon
< breaking up to
give 2 anyons



↳ think of it
as an element in the Hilbert space $\ell^2(\mathbb{R}^4)$

So,



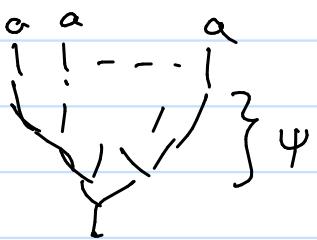
$$\Rightarrow \begin{array}{c} a \\ b \\ c \\ \diagdown \\ x \end{array} = \sum_y F_{\dots} \begin{array}{c} a \\ b \\ c \\ \diagdown \\ y \end{array}$$

associativity

$$fl \left(\text{Diagram} \right) = fl^{\circ}_{a,a,\dots,a}$$

B_n acts by particle exchange

a vector in $fl^{\circ}_{a,\dots,a} =$



$$\rho(\sigma_i) \begin{array}{|c|} \hline \psi \\ \hline \end{array} = \begin{array}{|c|} \hline \psi \\ \hline \end{array} \quad \begin{array}{l} \text{a a} \\ \text{---} \\ \text{a} \end{array} \quad \begin{array}{l} \text{a a} \\ \text{---} \\ \text{a} \end{array}$$

There are unitary operators on $fl(\dots)$
 \Rightarrow These are gates
(Braiding)

$\{ \rho(\sigma_i) \} \rightarrow$ gate set

Idea: Construct a MTC from explicit algebras

Temperley - Lieb algebras:

$TL_n(A)$: u_1, u_2, \dots, u_{n-1} generators

such that

$$1) u_i^2 = d u_i \quad \text{where } A = - (A^2 + A^{-2})$$

$$2) u_i u_{i+1} u_i = u_i$$

$$3) [u_i, u_j] = 0 \quad \text{if } |i-j| \neq 1$$

This is an algebra over $\mathbb{C}(A)$

But working over $\mathbb{Q}[A, A^{-1}]$ suffices.

Group algebra of $B_n = \mathbb{Q}[A, A^{-1}] B_n$

$$\varphi: \mathbb{Q}[A, A^{-1}] B_n \longrightarrow \text{TL}_n(A)$$

$$\varphi(\sigma_i) = \beta_i := A^{-1}u_i + A \mathbf{1}$$

Exercise: Check that this is a surjective homomorphism.

(For homo: show β_i satisfy braid relation)
 (For surj: show β_i generate TL_n)

Theorem: TL_n is f.d.-g.s. (for generic A)

Diagrammatic version : $\mathcal{TL}_n(A)$

Define $U_i : | \cdots | \stackrel{U}{\underset{i \text{ int}}{\sim}} | \cdots |$

We Make a monoid with generators U_i

$$\text{Unit} : 1_n \rightarrow ||\cdots|$$

product operation: stacking pictures

elements are up to "d-isotopy"

- loops: remove loops & multiply by $d^{\# \text{loops}}$
- $+$: formal sums, action of $\mathbb{Q}[A, A^{-1}]$ as coeff.

$$\text{Example} : U_i^2 = \left\{ \begin{array}{c} | \cdots | \stackrel{U}{\sim} | \cdots | \\ | \cdots | \underset{i \text{ int}}{\sim} | \cdots | \end{array} \right\} = d \left\{ \begin{array}{c} | \cdots | \stackrel{U}{\sim} | \cdots | \\ | \cdots | \underset{n \text{ int}}{\sim} | \cdots | \end{array} \right\} = d U_i$$

Thm : $\text{TL}_n(A) \simeq \mathcal{TL}_n(A)$
 $U_i \mapsto U_i$

A basis for $\mathcal{TL}_n(A)$ is noncrossing perfect matchings on A .

$\mathcal{TL}_3(A)$ has basis

$$\{ \text{|||}, \text{~\textbackslash~\textbackslash~}, \text{\textbackslash~\textbackslash~}, \text{\textbackslash~\textbackslash~}, \text{\textbackslash~\textbackslash~} \}$$

Thm: $\dim(\mathcal{TL}_n(A)) = C_n$ catalan number

$$\mathcal{TL}_3(A) \cong \mathbb{C} \oplus M_2$$

We have Kauffman bracket

$$o_i \mapsto A^{-1} u_i + A$$

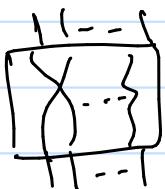
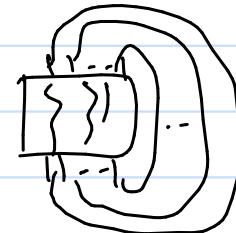
$$\begin{aligned} (\text{Skien relations}) \quad & \langle \times \rangle = A^{-1} \langle \text{\textbackslash~\textbackslash~} \rangle + A \langle \text{||} \rangle \\ & \langle 0 \rangle = d \end{aligned}$$

On $\bigcup_{n \geq 1} \mathcal{TL}_n(A)$, we have a Markov trace

$$\text{Tr}(D) = \begin{array}{c} \text{Diagram of } D \\ \text{with loops} \end{array} = d^{\# \text{loops}}$$

$$\text{Example: } \text{Tr}(U_1) = \text{Tr} \left(\begin{array}{|c|c|} \hline \text{D} & \text{D} \\ \hline \text{D} & \text{D} \\ \hline \end{array} \right) = \begin{array}{c} \text{Diagram of } U_1 \\ \text{with loops} \end{array} = d$$

$$\text{Tr}(U_2) = \text{Tr} \left(\begin{array}{|c|c|} \hline \text{D} & \text{D} \\ \hline \text{D} & \text{D} \\ \hline \end{array} \right) = \begin{array}{c} \text{Diagram of } U_2 \\ \text{with loops} \end{array} = d^2$$

$\beta \in \mathcal{B}_n$  $\rightarrow \hat{\beta} = L$  $L = \hat{\beta}$ a linkWe can write $\beta = \sigma_{i_1}^{\alpha_1} \cdots \sigma_{i_k}^{\alpha_k}$

$$e(\beta) := \sum_{j=1}^k \alpha_j$$

The Jones polynomial

$$J(L, q) = \frac{(-A^{-3})^{e(\beta)} \operatorname{Tr}(q(\beta))}{d} \quad \left| \begin{array}{l} A = q^{-\frac{1}{4}} \\ d \in \mathbb{Q}[q^{\pm \frac{1}{4}}, i^{\pm \frac{1}{2}}] \end{array} \right.$$

- Jones poly. is defined for oriented links.
- We can orient all lines upwards to get oriented links.

Alexander's Thm: Any oriented link is the closure of some oriented braid.

Markov's Thm: Two braids β and γ have $\hat{\beta} = \hat{\gamma}$ iff \exists a sequence of Markov moves

$$\beta \xrightarrow{M_{i_1}} \dots \xrightarrow{M_{i_m}} \gamma$$

$$(M1) \quad \beta \mapsto \begin{array}{c} \text{---} \\ | \quad | \\ \alpha \quad \beta \\ | \quad | \\ \text{---} \end{array}$$

$$(M2) \quad \beta \mapsto \begin{array}{c} \text{---} \\ | \quad | \\ \beta \quad \alpha \\ | \quad | \\ \text{---} \end{array} \in \mathcal{B}_{n+1}$$

$$\text{Example: } J\left(\frac{\sigma_1^3}{n}, q\right) = q + q^3 - q^4$$

Theorem [Vertigan] Exact computation of $J(L, q)|_{q=e^{2\pi i/\ell}}$ is $\text{FP}^{\#P}$ -hard in # of crossings of L for $\ell \neq 1, 2, 3, 4, 6$

$\text{FP}^{\#P} \rightarrow$ function version of NP-hard

Not known: Fully Polynomial Randomized Approximation Scheme.

Q Can we approximate $J(L, q)$ on Q.C. in polynomial time?
(with high accuracy & high confidence)

YES (Technically $|J(L, q)|^2$)

JONES - WENZL PROJECTIONS

$$\Delta_0(d) = 1, \Delta_1(d) = d, \Delta_{n+1} = d\Delta_n - \Delta_{n-1}$$

$$P_n \in \text{TL}_n(A)$$

$$P_1 = \boxed{1} := \begin{smallmatrix} 1 \\ \vdots \\ 1 \end{smallmatrix}, P_2 = \boxed{1 - \frac{1}{d}} \cup \begin{smallmatrix} 1 \\ \vdots \\ 1 \end{smallmatrix} =: \boxed{\frac{1}{2}}$$

$$P_{n+1} = \boxed{n} \Big] - \frac{\Delta_{n-1}}{\Delta_n} \boxed{n} \Big] \quad \begin{smallmatrix} 1 \\ \vdots \\ n \\ \vdots \\ 1 \end{smallmatrix}$$

Exercise: Draw P_3

- Lemma:
- 1) $P_n^2 = P_n$
 - 2) $P_n U_i = U_i P_n = 0, \quad 1 \leq i \leq n-1$
 - 3) $\text{Tr}(P_n) = \Delta_n$

Exercise: Prove this

Theorem: Let $r \geq 3$ integer

Set $A = \begin{cases} ie^{-2\pi i / 4r} & r \text{ even} \\ ie^{\pm 2\pi i / 2r} & r \text{ odd} \\ \pm \rightarrow r \bmod 4 \end{cases}$

$\mathcal{J}_{L_n}(A)$ bad things happen
(we can have zero divisors)

Thm: $\frac{\mathcal{J}_{L_n}(A)}{\langle P_{n-1} \rangle}$ is s.s. and
composing with φ , we get
a unitary B_m representation
 $\neq n$

need sesquilinear form for unitary

Use $\langle \cdot \rangle$ on $\mathcal{J}_{L_n}(A)$ defined by

$$\langle P, Q \rangle = \text{Tr} \left(\begin{array}{c|cc|c} & \bar{P} & & \\ \hline & & \bar{Q} & \\ & & & I \end{array} \right) \quad \begin{matrix} \bar{P} \text{ is } P \\ \text{upside down} \end{matrix}$$

This is a sesquilinear form &

$$\langle P(\beta)|\psi\rangle, \varphi(\beta)|\pi\rangle \rangle = \langle |\psi\rangle, |\pi\rangle \rangle$$

on $\frac{\mathcal{J}_{L_n}(A)}{\langle P_{n-1} \rangle}$

is positive definite

We want to come up with a category

Next time: construct a category

Objects:



intervals with
points on them

Morphisms: linear combinations of
TL-diagrams