Quantum Error Correcting Codes FQC Online Classes 2020

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Introduction

- A quantum error-correcting code (QECC) can be viewed as a mapping of k qubits into n qubits, where n > k.
- The k qubits are the "logical qubits" or "encoded qubits" that we wish to protect from errors.
- The additional n k qubits allow us to store the k logical qubits in a redundant fashion so that the encoded information is not easily damaged.

How does the classical error correction works?

- The simplest example of a classical error-correcting code is a repetition code: the bit we wish to protect by 3 copies of the bit.
- We encode logical 0 by 000 and 1 by 111

- Now an error may occur that causes one of the three bits to flip.
- Each bit is assumed to flip $(0 \leftarrow \rightarrow 1)$ independently with probability p.
- Thus a bit 0 sent through the channel will be received as 0 with probability 1-p and will be received as 1 with probability p.
- To reduce channel errors, we may invoke the majority vote.

- When 000 is sent through this channel, it will be received as 000 with probability $(1 p)^3$, as 100, 010 or 001 with probability 3p $(1 p)^2$, as 011, 101, or 110 with probability 3p²(1 p) and finally as 111 with probability p³.
- Note that summation of all the probabilities is 1.
- By taking the majority vote, we correctly reproduce the desired result 0 with probability $p_0 = (1 p)^3 + 3p (1 p)^2 = (1 p)^2 (1 + 2p)$ and fail with probability $p_1 = 3p^2(1 p) + p^3 = (3 2p) p^2$. We obtain $p_0 \gg p_1$ for sufficiently small $p \ge 0$. In fact, we find $p_0 = 0.972$ and $p_1 = 0.028$ for p = 0.1. The success probability p_0 increases as $p_0 = 0.972$ and $p_1 = 0.028$ for $p_0 = 0.1$.

Three Challenges in QECCs

- No cloning: Duplicating quantum states to get repetition code is impossible.
- Errors are continuous: A continuum of different errors may occur on a single qubit. Therefore, determination of which error occurred requires infinite precision and infinite resources.
- Measurement destroys quantum information: Measurement destroys quantum information and makes recovery impossible if quantum state is destroyed.

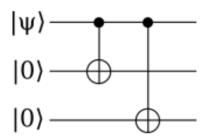
• The single bit flip error is described by:

$$|0\rangle \to |1\rangle$$
$$|1\rangle \to |0\rangle$$

The Bit flip error is equivalent to action of the X-gate.

- Suppose Alice wants to send qubit $\Psi = a |0\rangle + b |1\rangle$ to Bob through a noisy quantum channel
- Encoding: Alice uses 3 qubits to encode one qubit. Similar to the repetition code, we will append two ancillary bits $|00\rangle$ to Ψ , giving us Ψ_1 .

• A circuit performing this encoding is shown below:



$$\Psi_0 \equiv (a |0\rangle + b |1\rangle) |0\rangle |0\rangle$$
$$\equiv a |000\rangle + b |100\rangle$$
$$\Psi_1 \equiv a |000\rangle + b |111\rangle$$

• Error detection: Assuming that only a single qubit flip can occur during transmission, Bob will receive one of these 4 possible states

$$\Phi_0 \equiv a |000\rangle + b |111\rangle$$
 // No error $\Phi_1 \equiv a |001\rangle + b |110\rangle$ // 1^{st} qubit flip $\Phi_2 \equiv a |010\rangle + b |101\rangle$ // 2^{nd} qubit flip $\Phi_3 \equiv a |100\rangle + b |011\rangle$ // 3^{rd} qubit flip

• Now Bob has to extract from the received state which error occurred during the qubit transmission. For this purpose, Bob introduces two ancillary qubits in the state |00| as shown in Fig. 10.1.

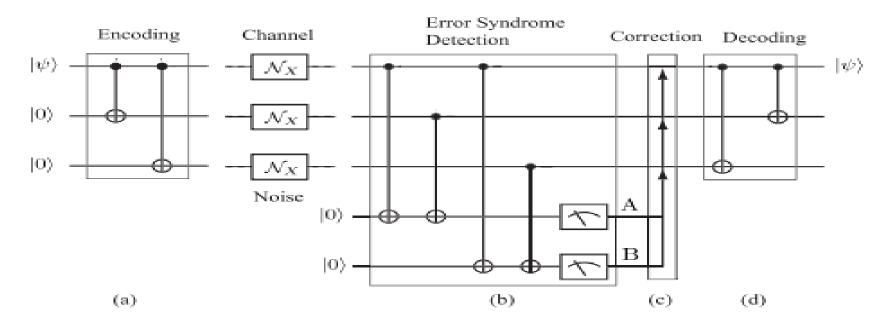


FIGURE 10.1

Quantum circuits to (a) encode, (b) detect bit-flip error syndrome, (c) make correction to a relevant qubit and (d) decode. The gate \mathcal{N}_X stands for the bit-flip noise. The circuit (a) belongs to Alice, while the circuits (b), (c) and (d) belong to Bob.

- Bob applies four CNOT operations whose control bits are the encoded qubits while the target qubits are Bob's two ancillary qubits. Let $|x_1x_2x_3\rangle$ be a basis vector Bob has received and let A (B) be the output state of the first (second) ancillary qubit. It is seen from Fig. 10.1 (b) that $A = x_1 \oplus x_2$ and $B = x_1 \oplus x_3$.
- Suppose $a|100\rangle + b|011\rangle$ is the received logical qubit. Then error extracting sequence transforms the ancillary qubits as

$$(a|100\rangle + b|011\rangle)|00\rangle \to a|10011\rangle + b|01111\rangle = (a|100\rangle + b|011\rangle)|11\rangle.$$

• Both of the ancillary qubits are flipped since $x_1 \oplus x_2 = x_1 \oplus x_3 = 1$ for both $|100\rangle$ and $|011\rangle$. The set of two bits is called error syndrome and it tells Bob in which physical qubit the error occurred during transmission

• The states after error extraction is made is given below:

State after error syndrome extraction
$(a 000\rangle + b 111\rangle) 00\rangle$
$(a 100\rangle + b 011\rangle) 11\rangle$
$(a 010\rangle + b 101\rangle) 10\rangle$
$(a 001\rangle + b 110\rangle) 01\rangle$

- Error recovery: Once the *error syndrome* is detected, error recovery is very simple.
- Suppose $(a|100) + b|011\rangle |11\rangle$ is the state after error syndrome extraction.
- Now Bob measures his two ancillary qubits and obtains two bits 11 of classical information.
- In this case it tells us that a bit-flip error occurred on qubit number 3 (or 11 in binary). So, Bob corrects the error by applying X-gate to qubit number 3.

•
$$a|100\rangle + b|011\rangle \longrightarrow a|000\rangle + b|111\rangle$$

Now, by applying the inverse of the encoding procedure he recovers $a|0\rangle + b|1\rangle$

$$a|000\rangle + b|111\rangle \longrightarrow a|0\rangle + b|1\rangle |00\rangle$$

• The same procedure works in the case that the second or third qubit experiences a bit-flip as well.

Error	Correction to be made
Syndrome	
(00)	Do Nothing
(01)	Apply X gate to 1st qubit (X ₁)
(10)	Apply X gate to 2 nd qubit (X ₂)
(11)	Apply X gate to 3 rd qubit (X ₃)

Phase-Flip QECC

 Now we will see, how to handle only a single qubit phase flip which might creep in during transmission. Phase flip error is equivalent to action of Pauli matrix Z

$$a|0\rangle + b|1\rangle \rightarrow a|0\rangle - b|1\rangle$$

• Here Hadamard gate comes to our rescue again. The idea is that phase flip error in the $|+\rangle$, $|-\rangle$ basis is same as bit flip error

$$|+\rangle \rightarrow |-\rangle$$

$$|-\rangle \rightarrow |+\rangle$$

This simple idea gives the whole algorithm for phase flip errors:

Phase-Flip QECC

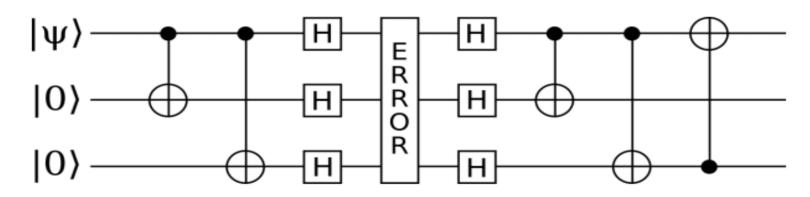


Figure 4: Phase Flip Encoding and Decoding Wik17

- Append ancilla bits: $a\ket{0} + b\ket{1} \rightarrow a\ket{000} + b\ket{111}$
- Apply Hadamard: $a |000\rangle + b |111\rangle \rightarrow a |+++\rangle + b |---\rangle$
- ullet Apply bit flip algorithm to detect and correct a bit flip error in the $|+\rangle\,,|-\rangle$ basis.
- ullet Apply Hadamard again to recover original message in $|0\rangle\,, |1\rangle$ basis.

Shor's Nine-qubit Code

• The Shor code is such a code which uses 9 qubits to encode 1 qubit. Essentially, the Shor code is a combination of bit flip and phase flip codes. We first encode a qubit using the encoding procedure of the phase flip code:

$$|0\rangle \rightarrow |+++\rangle, |1\rangle \rightarrow |---\rangle$$

 This is followed by encoding each of these 3 qubits according to the bit flip encoding procedure:

$$|+\rangle \rightarrow (|000\rangle + |111\rangle)/\sqrt{2}$$

 $|-\rangle \rightarrow (|000\rangle - |111\rangle)/\sqrt{2}$

Thus the final codewords are:

$$|0\rangle \rightarrow |0_L\rangle \equiv \frac{(|000\rangle + |111\rangle)(|000\rangle + |111\rangle)(|000\rangle + |111\rangle)}{2\sqrt{2}}$$

$$|1\rangle \rightarrow |1_L\rangle \equiv \frac{(|000\rangle - |111\rangle)(|000\rangle - |111\rangle)(|000\rangle - |111\rangle)}{2\sqrt{2}}$$

Shor's Nine-qubit Code

• The Fig. 5 below shows Shor's nine-qubit encoding

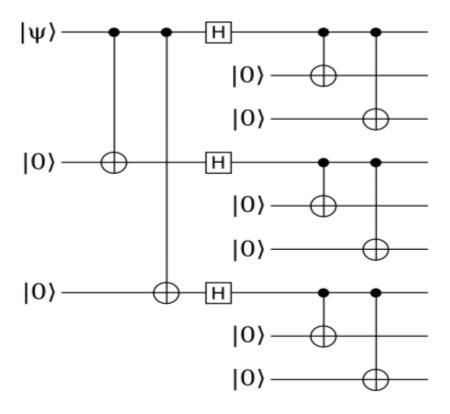


Figure 5: Shor Code Encoding [Wik17]

Shor's Nine-qubit Code

- One can show that shor's code detects and corrects the following errors:
 - Bit flip error on a single qubit.
 - Phase flip error on a single qubit.
 - Combined bit and phase flip error on a single qubit.