Unix System Calls

PART-1

Process Management Model

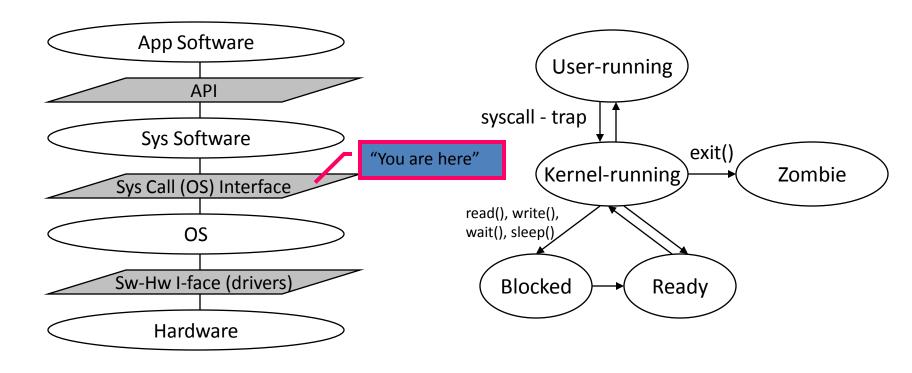
The Unix process management model is split into two distinct operations:

- 1. The creation of a process.
- 2. The running of a new program.

Overview

- System Call Interface
- Process Management with C
 - fork()
 - -exec()
 - wait()
 - -exit()

System Call Interface



The fork() System Call (1)

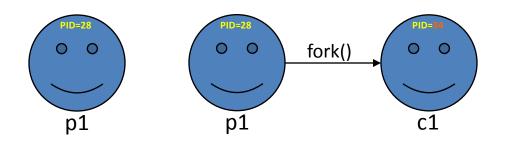
- A process calling fork() spawns a child process.
- The child is almost an identical clone of the parent:
 - Program Text (segment .text)
 - Stack (ss)
 - PCB (eg. registers)
 - Data (segment .data)

```
#include <sys/types.h>
#include <unistd.h>
pid_t fork(void);
```

f you only need the OS types, say for the prototypes of your functions, then just #include <sys/types.h>. However if you need the function definitions, then you #include <unistd.h> or any of the other system headers, as needed.

The fork() System Call (2)

- The fork() is one of the those system calls, which is called once, but returns twice!
- After fork() both the parent and the child are executing the same program.
- On error, fork() returns -1



```
Consider a piece of program
...
pid_t pid = fork();
printf("PID: %d\n", pid);
...

The parent will print:
PID: 34
And the child will always print:
PID: 0
```

The **pid_t** data type represents process IDs. You can get the process ID of a process by calling getpid. The function getppid returns the process ID of the parent of the current process (this is also known as the parent process) (this is also known as the parent process)

The following is a simple example of fork()

```
#include <stdio.h>
#include <sys/types.h>
#include <unistd.h>
int main(void)
 printf("Hello \n");
 fork();
 printf("bye\n");
 return 0;
```

Hello –is printed once by parent process

bye consequent and once by the child

The exec () System Call (1)

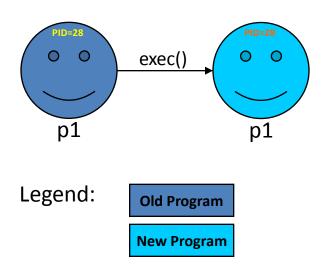
- The exec() call replaces a current process' image with a new one (i.e. loads a new program within current process).
- The new image is either regular executable binary file or a shell script.
- There's no a syscall under the name exec(). By exec() we usually refer to a family of calls:

```
- int execl(char *path, char *arg, ...);
- int execv(char *path, char *argv[]);
- int execle(char *path, char *arg, ..., char *envp[]);
- int execve(char *path, char *argv[], char *envp[]);
- int execlp(char *file, char *arg, ...);
- int execvp(char *file, char *argv[]);
```

- Here's what I, v, e, and p mean:
 - I means an argument list,
 - v means an argument vector,
 - e means an environment vector, and
 - p means a search path.

The exec () System Call (2)

- Upon success, exec() never returns to the caller. If it does return, it means the call failed. Typical reasons are: non-existent file (bad path) or bad permissions.
- Arguments passed via exec() appear in the argv[] of the main() function.

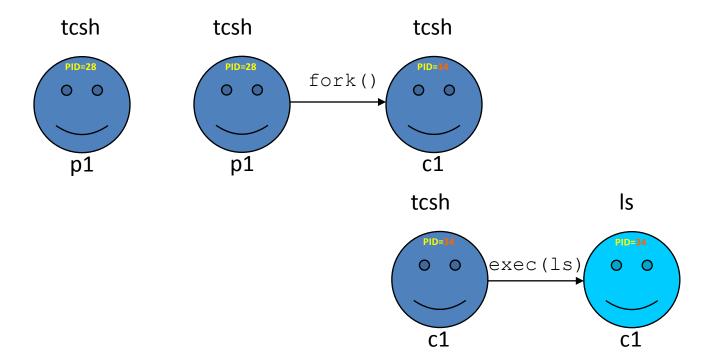


Environment

- The *e*-exec calls use the environment when attempt to invoke a new program.
- Name = Value
 - HOME
 - PATH
 - SHELL
 - USER
 - LOGNAME
 - **–** ...
- set or env will display current environment, which you can modify with:
 - the export command in a shell or a shell script (bash);
 - the setenv for tcsh
 - the getenv(), setenv(), putenv(), etc. in C

fork() and exec() Combined

• Often after doing fork() we want to load a new program into the child. *E.g.*: a shell.



The System wait () Call

Forces the parent to suspend execution, i.e.
 wait for its children or a specific child to die
 (terminate is more appropriate terminology,
 but a bit less common).

The System wait() Call (2)

```
#include <sys/types.h>
#include <sys/wait.h>

pid_t wait(int *status);
pid_t waitpid(pid_t pid, int *status, int options);
```

- The wait() causes the parent to wait for any child process.
- The waitpid() waits for the child with specific PID.
- The status, if not NULL, stores exit information of the child, which can be analyzed by the parent using the W^* () macros.
- The return value is:
 - PID of the exited process, if no error
 - (-1) if an error has happened

The exit() System Call

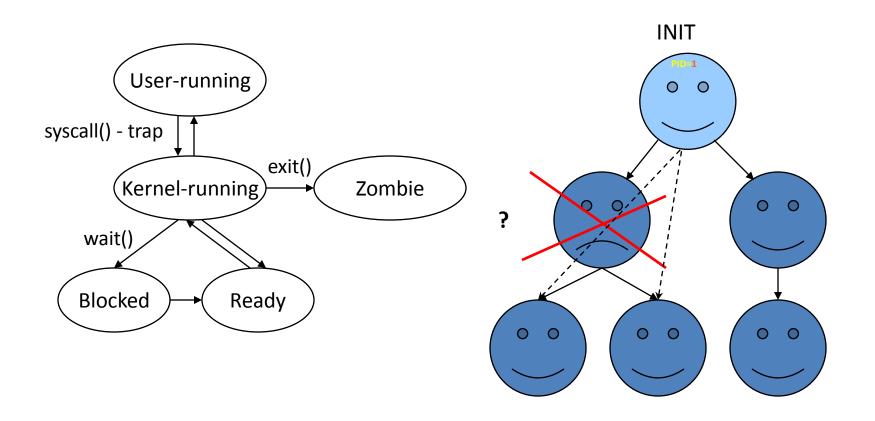
```
#include <stdlib.h>
void exit(int status);
```

- This call gracefully terminates process execution. Gracefully means it does clean up and release of resources, and puts the process into the zombie state.
- By calling wait (), the parent cleans up all its zombie children.
- exit() specifies a return value from the program, which a parent process might want to examine as well as status of the dead process.
- $\underline{-}$ exit() call is another possibility of quick death without cleanup.

exit()

- The exit operation typically performs clean-up operations within the process space before returning control back to the operating system. Some systems and programming languages allow user subroutines to be registered so that they are invoked at program termination before the process actually terminates for good. As the final step of termination, a primitive system exit call is invoked, informing the operating system that the process has terminated and allows it to reclaim the resources used by the process.
- It is sometimes possible to bypass the usual cleanup; C99 offers the _exit() function which terminates the current process without any extra program clean-up. This may be used, for example, in a forkexec routine when the exec call fails to replace the child process.

Process Overview



```
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     Example: to demonstrate fork() and exect() and system calls
 ************************
#include <stdio.h>
#include <stdlib.h>
#include <sys/types.h>
#include <unistd.h>
#include <sys/wait.h>
int main( int argc, char *argv[], char *env[] )
  pid t my pid, parent pid, child pid;
  int status:
/* get and print my pid and my parent's pid. */
  my pid = getpid(); parent pid = getppid();
  printf("\n Parent: my pid is %d\n\n", my pid);
  printf("Parent: my parent's pid is %d\n\n", parent pid);
/* print error message if fork() fails */
  if((child\ pid\ =\ fork())\ <\ 0\ )
  {
     perror("fork failure");
     EXIT | 1 | :
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```

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```
/* fork() == 0 for child process */
   if (child pid == 0)
   { printf("\nChild: I am a new-born process!\n\n");
     my pid = getpid(); parent pid = getppid();
     printf("Child: my pid is: %d\n\n", my pid);
      printf("Child: my parent's pid is: %d\n\n", parent pid);
      printf("Child: I will sleep 3 seconds and then execute - date - command \n\n");
      sleep(3);
      printf("Child: Now, I woke up and am executing date command \n\n");
      execl("/bin/date", "date", 0, 0);
      perror("execl() failure!\n\n");
      printf("This print is after execl() and should not have been executed if execl were successful! \n\n");
     exit(1);
/*
 * parent process
 */
   else
     printf("\nParent: I created a child process.\n\n");
      printf("Parent: my child's pid is: %d\n\n", child pid);
      system("ps -acefl | grep ercal"); printf("\n \n");
      wait(&status); /* can use wait(NULL) since exit status
                        from child is not used. */
      printf("\n Parent: my child is dead. I am going to leave.\n \n ");
  return 0;
}
```

C library function - perror()

- The C library function void perror(const char *str) prints a descriptive error message to stderr. First the string str is printed, followed by a colon then a space.
- Declaration
 - Following is the declaration for perror() function.
 - void perror(const char *str)Parameters
- **str** -- This is the C string containing a custom message to be printed before the error message itself.
 - Return Value
 - This function does not return any value.

Example

The following example shows the usage of perror() function.

```
#include <stdio.h>
int main ()
   FILE *fp;
   /* first rename if there is any file */
   rename("file.txt", "newfile.txt");
   /* now let's try to open same file */
   fp = fopen("file.txt", "r");
   if( fp == NULL )
      perror("Error: ");
      return(-1);
   fclose(fp);
   return(0);
  Course Instructor:
                                   Ram Chatterjee
```