## Artificial Intelligence: Adversarial Search

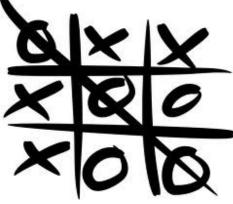
#### Motivation







chess

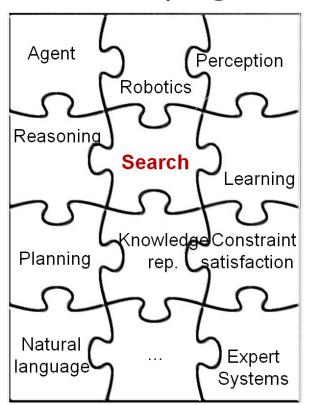


tic-tac-toe

## Today



- State Space Search for Game Playing
  - MiniMax
  - Alpha-beta pruning
  - Stochastic Games
- Where we are today



#### Adversarial Search

- Classical application for heuristic search
  - simple games: exhaustibly searchable
  - complex games: only partial search possible
  - additional problem: playing against opponent
- Here, we look at 2-player adversarial games
  - □ win, lose, or tie

#### Types of Games

#### Perfect Information

- A game with the perfect information is that in which agents can look into the complete board. Agents have all the information about the game, and they can see each other moves also.
- Examples: Chess, Checkers, Go, etc.

#### Imperfect Information

- Game state only partially observable, choices by opponent are not visible (hidden)
- Example: Battleship, Stratego, many card games, etc.

#### Types of Games (II)

#### Deterministic games

- No games of chance (e.g., rolling dice)
- Examples: Chess, Tic-Tac-Toe, Go, etc.

#### Non-deterministic games

- Games with unpredictable (random) events (involving chance or luck)
- Example: Backgammon, Monopoly, Poker, etc.

#### Types of Games (III)

#### Zero-Sum Game

- If the total gains of one player are added up, and the total losses are subtracted, they will sum to zero (example: cutting a cake)
- A gain by one player must be matched by a loss by the other player
- One player tries to maximize a single value, the other player tries to minimize it
- Examples: Checkers, Chess, etc.

#### Non-Zero-Sum Game

- Win-Win or Lose-Lose type games
- Famous example: The Prisoner's Dilemma

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  MiniMax
  - □ MiniMax
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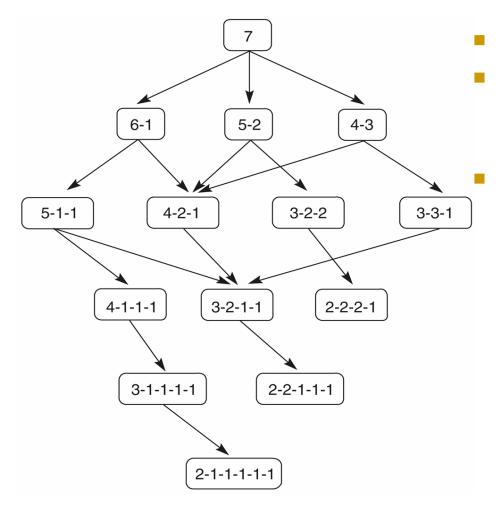
#### Example: Game of Nim

#### Rules

- 2 players start with a pile of tokens
- move: split (any) existing pile into two non-empty differently-sized piles
- game ends when no pile can be unevenly split
- player who cannot make his move loses

→ Worksheet #2 ("Game of Nim")

#### State Space of Game Nim



start with one pile of tokens each step has to divide one pile of tokens into 2 non-empty piles of different size

player without a move left loses game

source: G. Luger (2005)

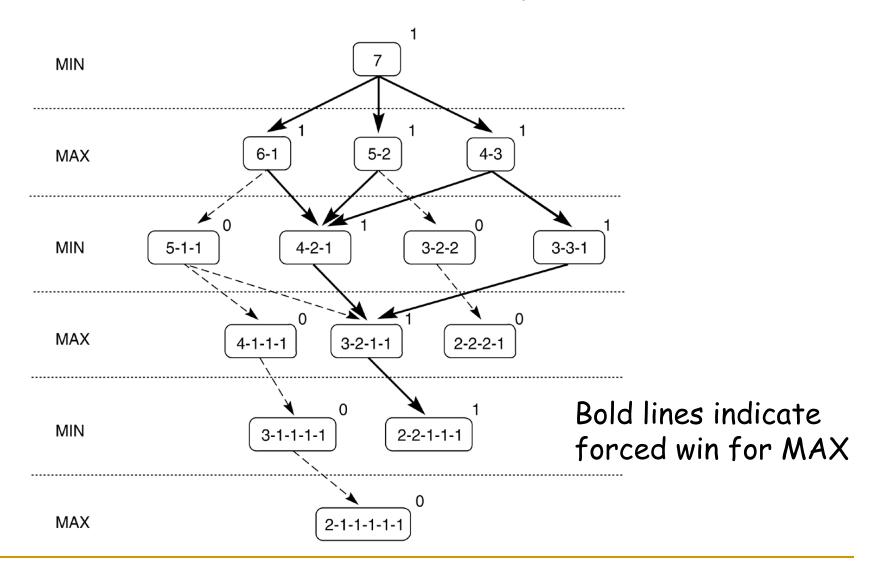
#### MiniMax Search

- Game between two opponents, MIN and MAX
  - MAX tries to win, and
  - MIN tries to minimize MAX's score
  - Existing heuristic search methods do not work
    - would require a helpful opponent
    - Need to incorporate "hostile" moves into search strategy

#### Exhaustive MiniMax Search

- For small games where exhaustive search is feasible
- Procedure:
  - build complete game tree
  - 2. label each level according to player's turn (MAX or MIN)
  - 3. label leaves with a utility function to determine the outcome of the game
    - e.g., (0, 1) or (-1, 0, 1)
  - 4. propagate this value up:
    - if parent=MAX, give it max value of children
    - if parent=MIN, give it min value of children
    - 5. Select best next move for player at root as the move leading to the child with the highest value (for MAX) or lowest values (for MIN)

#### Exhaustive MiniMax for Nim



source: G. Luger (2005)

#### n-ply MiniMax with Heuristic

- Exhaustive search for interesting games is rarely feasible
- Search only to predefined level
  - called n-ply look-ahead
  - n is number of levels
- No exhaustive search
  - nodes evaluated with heuristics and not win/loss
  - indicates best state that can be reached
  - horizon effect
- Games with opponent
  - simple strategy: try to maximize difference between players using a heuristic function e(n)

#### Heuristic Function for 2-player games

- simple strategy:
  - try to maximize difference between MAX's game and MIN's game
- typically called e(n)
- $\bullet$  e(n) is a heuristic that estimates how favorable a node n is for MAX
  - $\neg e(n) > 0 \longrightarrow n$  is favorable to MAX
  - $\neg e(n) < 0 \longrightarrow n$  is favorable to MIN
  - = e(n) = 0 --> n is neutral

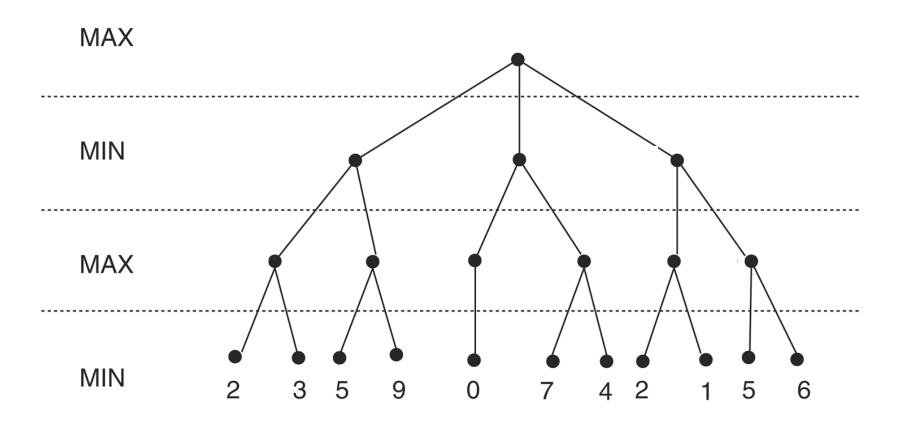
#### Choosing a Heuristic Function e(n)

Usually e(n) is a weighted sum of various features:

$$e(n) = \sum w_i f_i(n)$$

- E.g. of features:
  - $\neg$   $f_1$  = number of pieces left on the game for MAX
  - $\neg f_2 = \text{number of possible moves left for MAX}$
  - $\neg$   $f_3 = -(number of pieces left on the game for MIN)$
  - $\neg$   $f_4 = -(number of possible moves left for MIN)$
- E.g. of weights:
  - $w_1 = 0.5 // f_1$  is a very important feature
  - $w_2 = 0.2 // f_2$  is not very important
  - $w_3 = 0.2 // f_3$  is not very important
  - $w_4 = 0.1 // f_4$  is really not important

#### MiniMax with Fixed Ply Depth



Leaf nodes show the actual heuristic value e(n)

→ Worksheet #2 ("MiniMax")

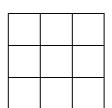
### Example: e(n) for Tic-Tac-Toe

#### Possible e(n)

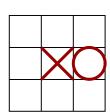
number of rows, columns, and diagonals open for MAX

- number of rows, columns, and diagonals open for MIN

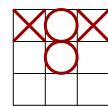
$$e(n) = \begin{cases} +\infty, & \text{if n is a forced win for MAX} \\ -\infty, & \text{if n is a forced win for MIN} \end{cases}$$



$$e(n) = 8-8 = 0$$



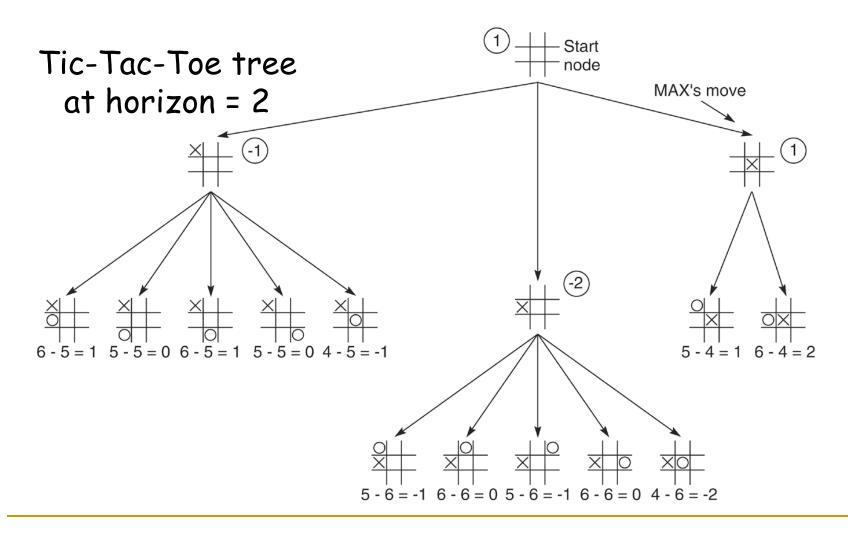
$$e(n) = 6-4 = 2$$



$$e(n) = 3-3 = 0$$

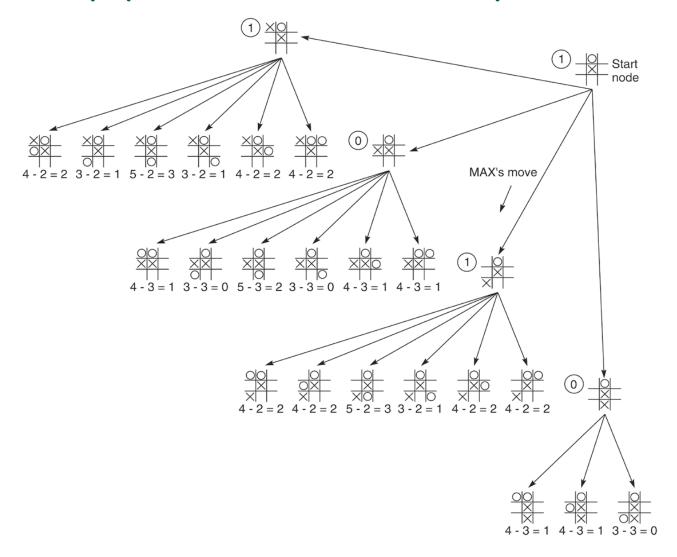
→ Worksheet #2 ("MiniMax Heuristic for Tic-Tac-Toe")

#### Two-ply MiniMax for Opening Move

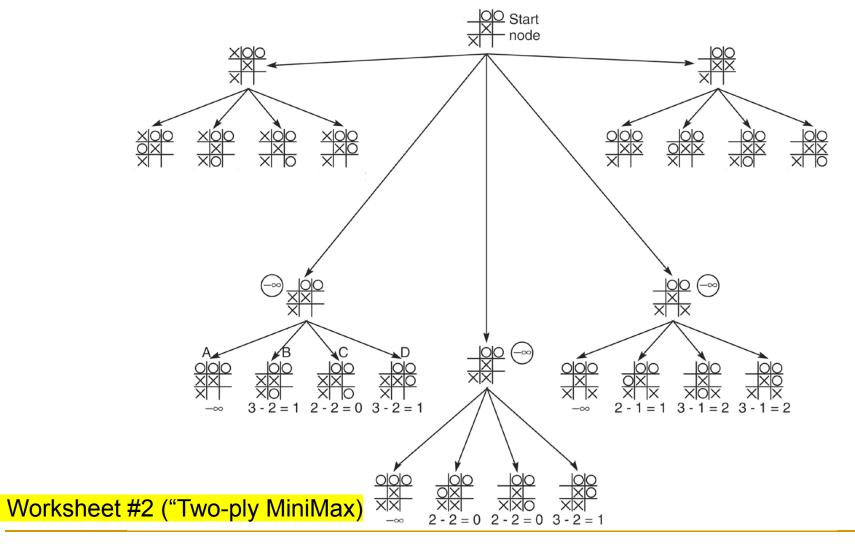


source: G. Luger (2005)

#### Two-ply MiniMax: MAX's possible 2nd moves



#### Two-ply minimax: MAX's move at end?



source: G. Luger (2005)

### Today

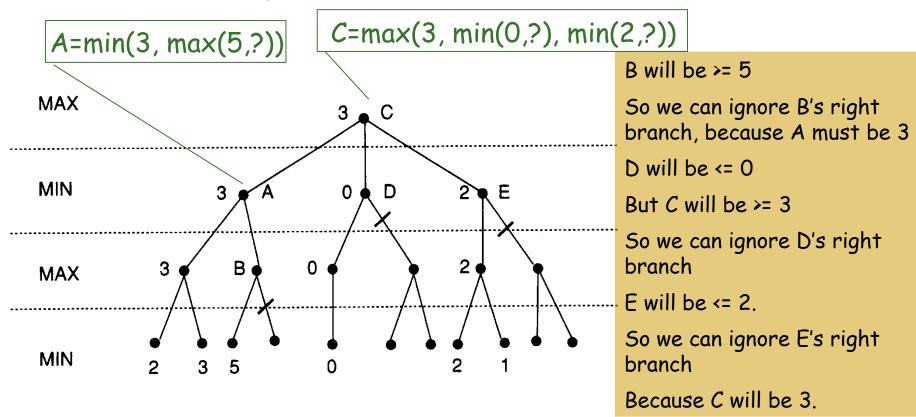
- State Space Search for Game Playing

  - Alpha-beta pruning
     Stock
  - Stochastic games
- Where we are today

#### Alpha-Beta Pruning

- Optimization over MiniMax, that:
  - ignores (cuts off, prunes) branches of the tree
     that cannot possibly lead to a better solution
  - reduces branching factor
  - allows deeper search with same effort

- With minimax, we look at all possible nodes at the n-ply depth
- With a-B pruning, we ignore branches that could not possibly contribute to the final decision



source: G. Luger (2005)

# Alpha-Beta Pruning Algorithm

- $\alpha$ : lower bound on the final backed-up value.
- $\beta$ : upper bound on the final backed-up value.
- Alpha pruning:

eg. if MAX node's  $\alpha$  = 6, then the search can prune branches from a MIN

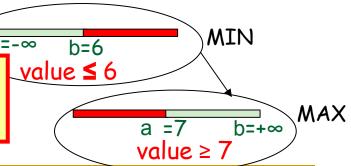
descendant that has a  $\beta \leftarrow 6$ .

□ if child  $\beta$  <= ancestor  $\alpha \rightarrow$  prune

incompatible...
so stop searching the right branch;
the value cannot come from there!

- Beta pruning:
  - eg. if a MIN node's  $\beta$  = 6, then the search can prune branches from a MAX descendant that has an  $\alpha >= 6$ .
  - □ if ancestor  $\beta$  <= child  $\alpha \rightarrow$  prune

incompatible...
so stop searching the right branch;
the value cannot come from there!



b=5

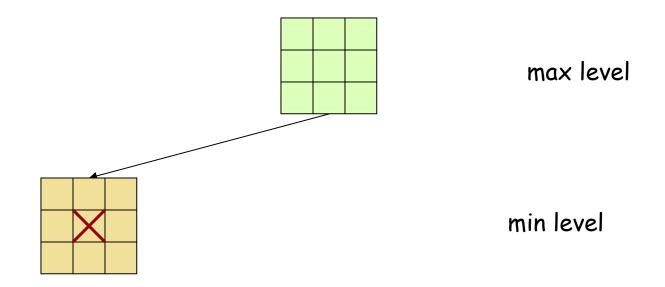
value ≤ 5

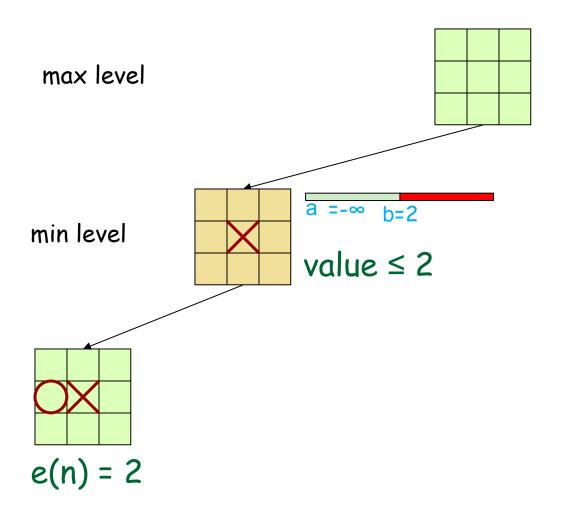
value ≥ 6

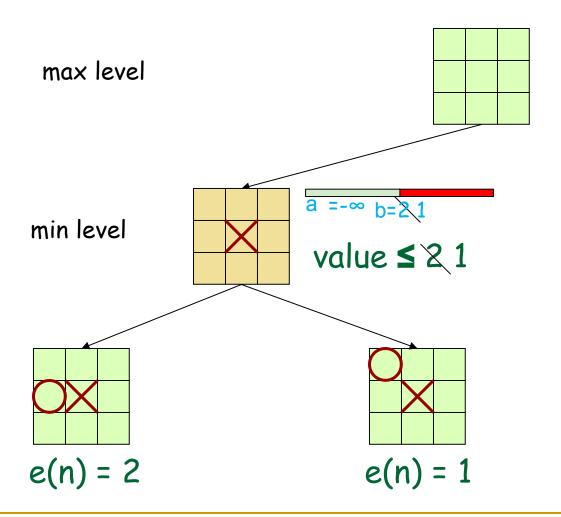
NIM

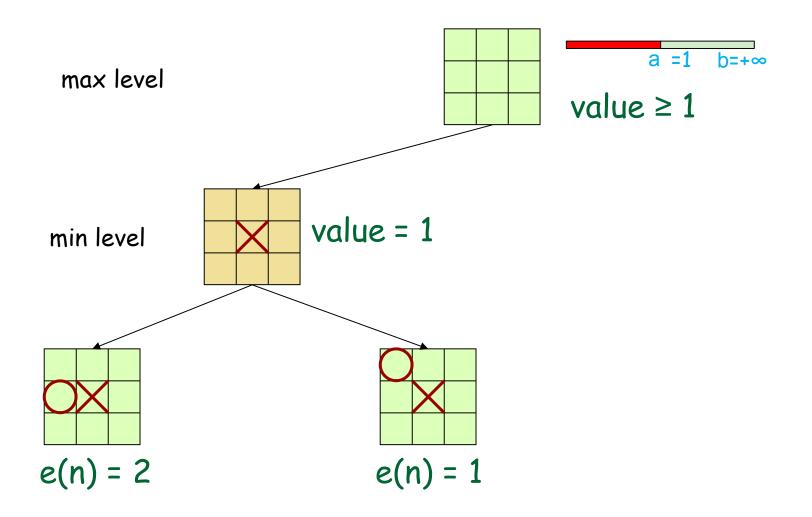
## Alpha-Beta Pruning Algorithm

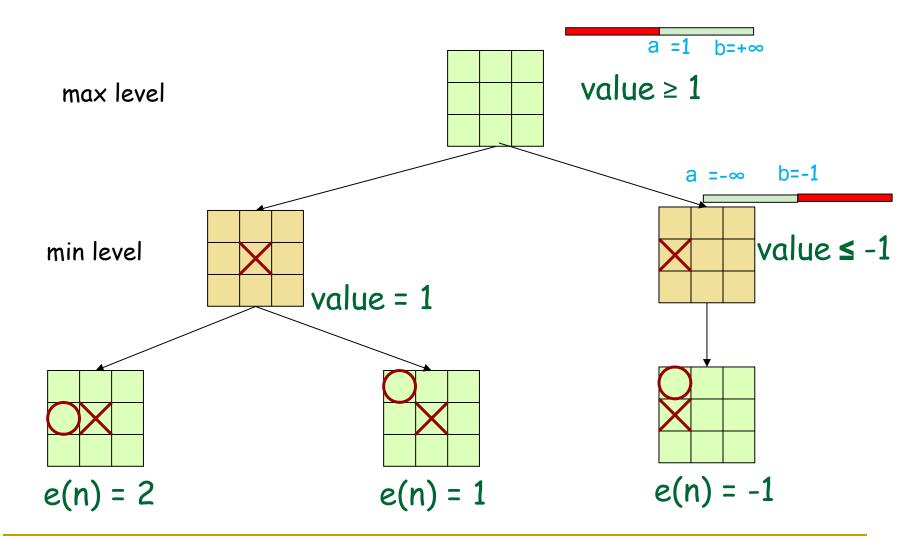
```
01 function alphabeta (node, depth, \alpha, \beta, maximizingPlayer)
02
         if depth = 0 or node is a terminal node
              return the heuristic value of node
03
0.4
         if maximizingPlayer
                                                             Initial call:
0.5
              v := -∞
                                                             alphabeta (origin, depth, -\infty, +\infty, TRUE)
06
              for each child of node
07
                   v := max(v, alphabeta(child, depth - 1, \alpha, \beta, FALSE))
08
                   \alpha := \max(\alpha, v)
                   if \beta \leq \alpha
09
                        break (* ß cut-off *)
10
11
              return v
12
         else
1.3
              ∨ := ∞
              for each child of node
14
15
                   v := min(v, alphabeta(child, depth - 1, \alpha, \beta, TRUE))
                   \beta := \min(\beta, v)
16
                   if \beta \leq \alpha
17
                        break (* α cut-off *)
18
19
              return v
```

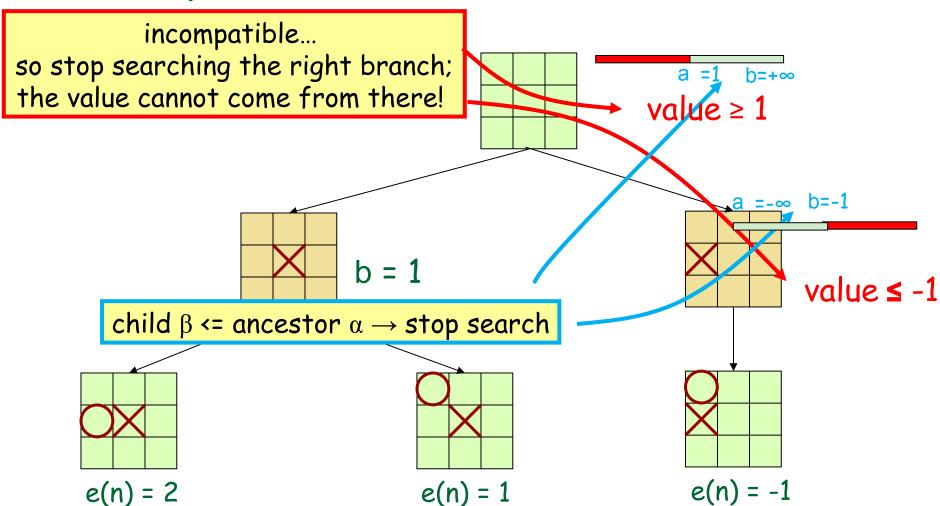


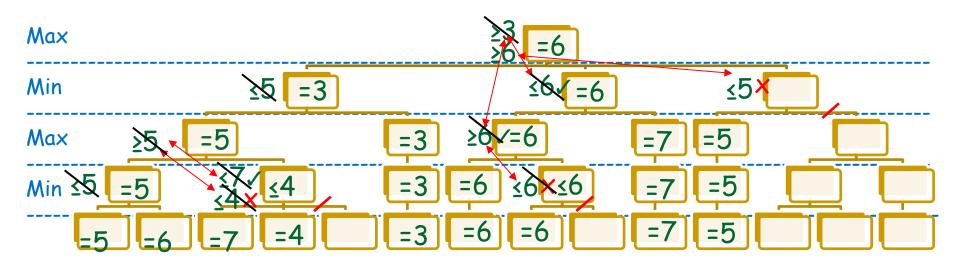


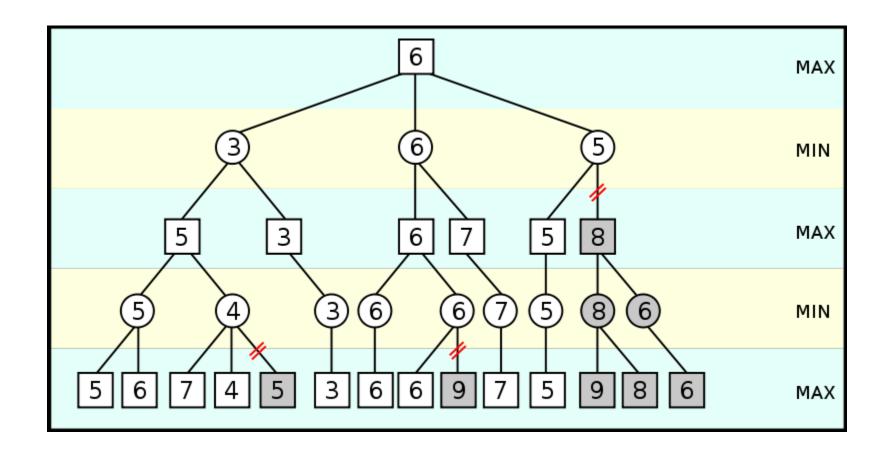


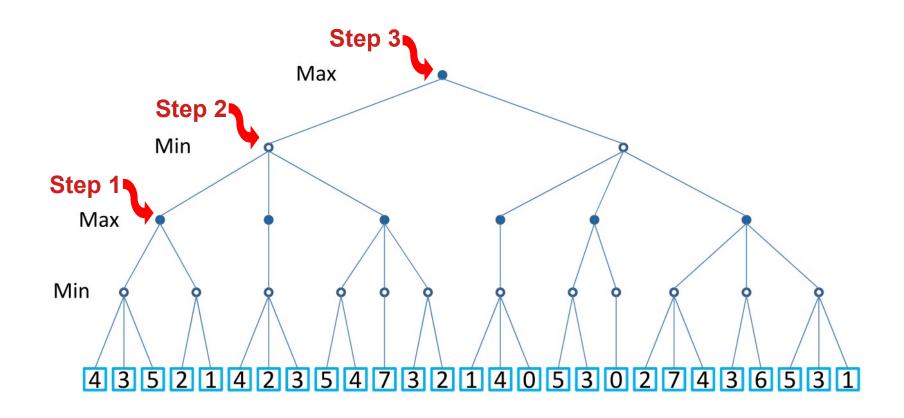








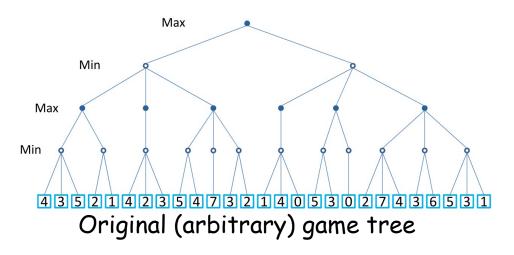


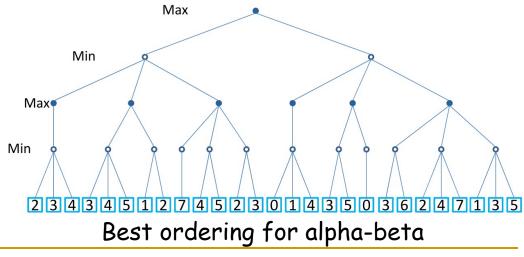


→ Worksheet #2 ("Alpha-Beta Pruning")

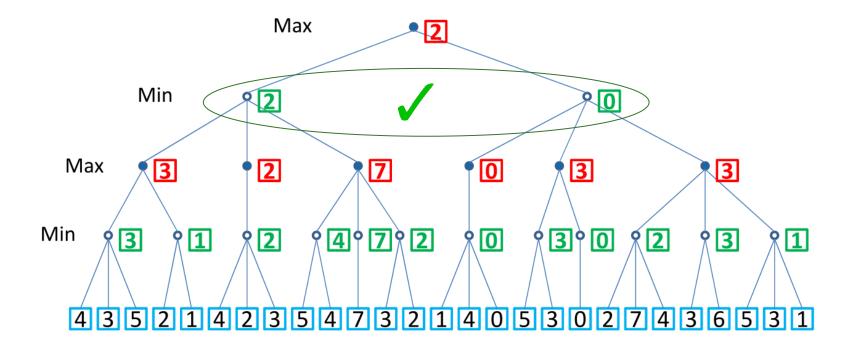
## Efficiency of Alpha-Beta Pruning

- Depends on the order of the siblings
- In worst case:
  - alpha-beta provides no pruning
- In best case:
  - branching factor is reduced to its square root

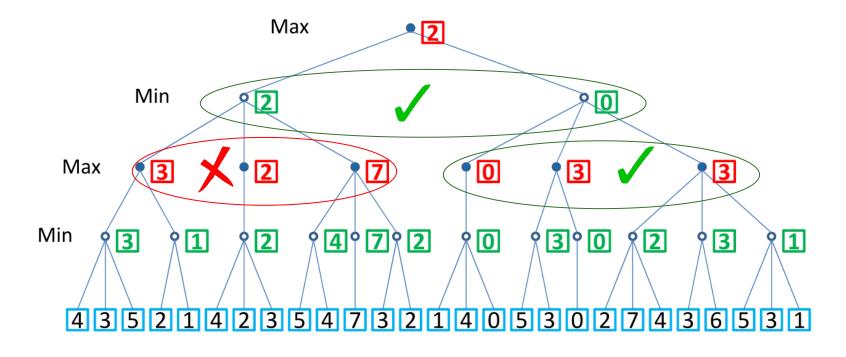




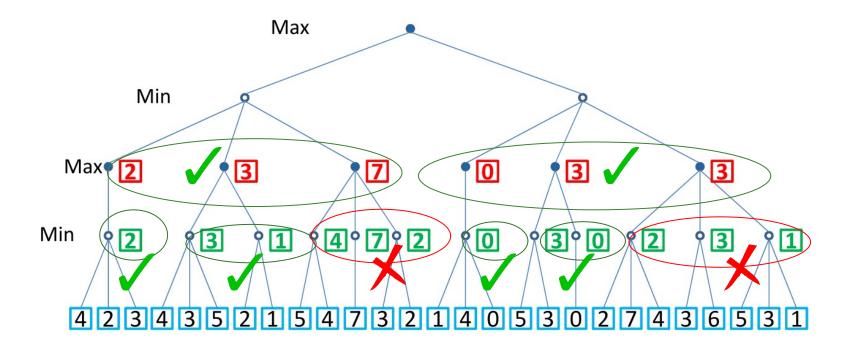
- best ordering:
  - children of MIN : smallest node first
  - children of MAX: largest node first

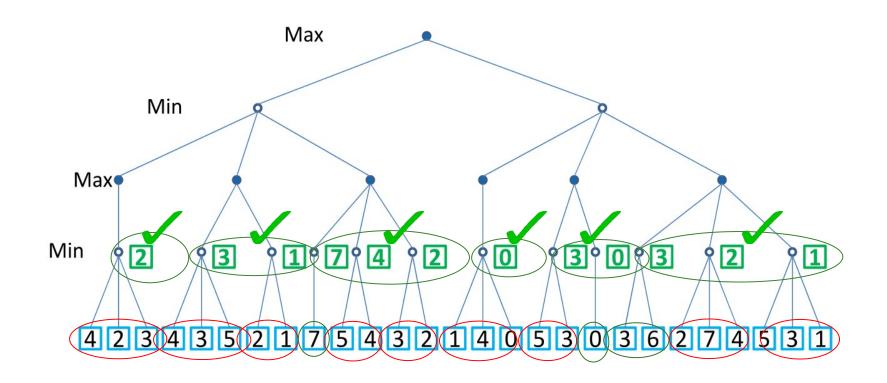


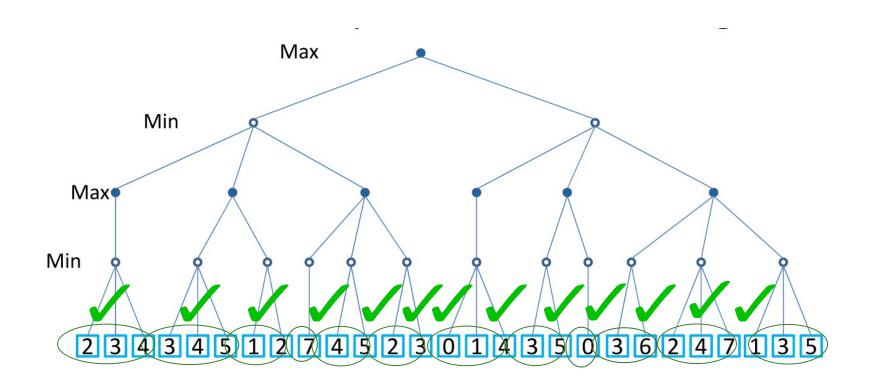
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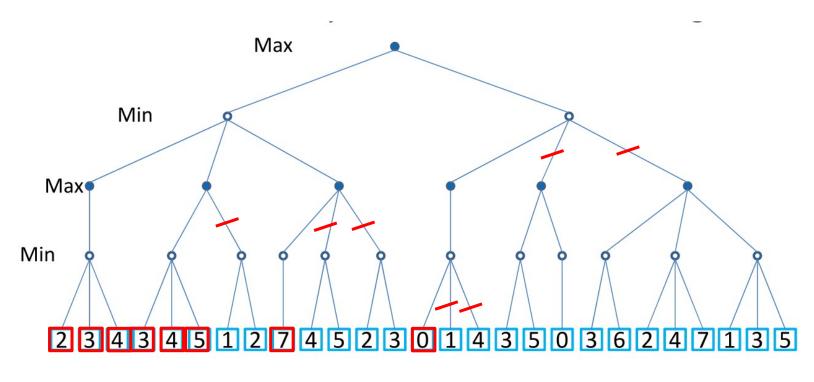


- best ordering:
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  - children of MAX: largest node first









8 nodes explored out of 27

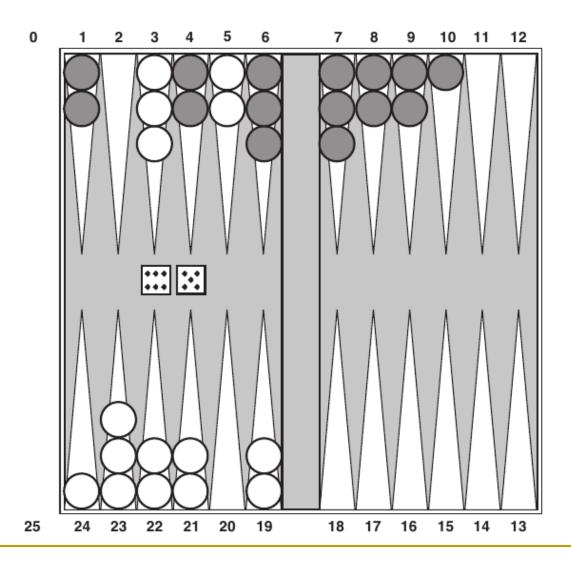
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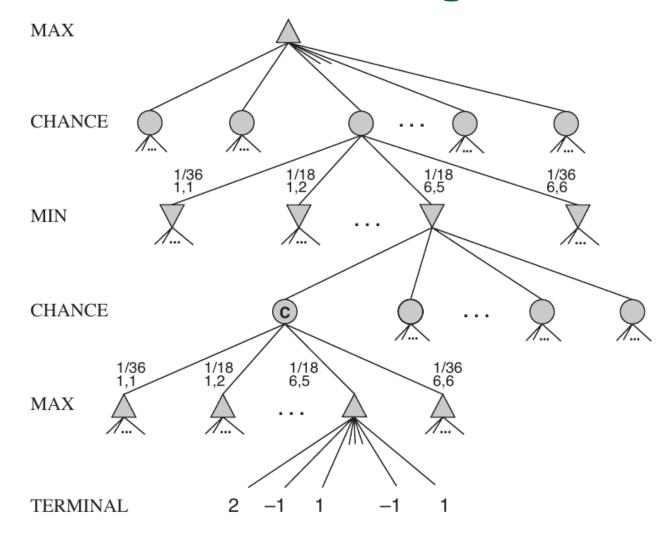
## Backgammon



## Stochastic (Non-Deterministic) Games

- Search tree for games of chance
  - white can calculate its own legal moves
  - but it does not know what black will roll...
- Idea: add chance nodes to the search tree
  - branches indicate possible dice rolls
  - each branch labeled with the roll and its probability (e.g., 1/6 for a single dice roll)

#### Search Tree for Backgammon



### EXPECTIMINIMAX Algorithm

- Calculating Expectiminimax
  - Like MiniMax, but using the sum of the weighted sum for Chance nodes:

$$\sum P(r)Expectiminimax(Result(s,r))$$

- $\neg$  r is a possible dice roll (or other random event)
- P(r) the probability of the event
- Result(s, r) is the same state s with dice roll result r
- Note: very expensive due to the high branching factor!
- See <a href="https://en.wikipedia.org/wiki/Expectiminimax">https://en.wikipedia.org/wiki/Expectiminimax</a>
   for the whole algorithm

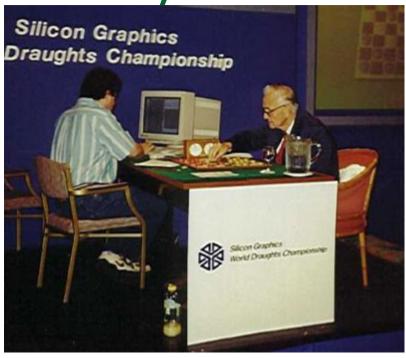
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Where we are today

#### 1992-1994 - Checkers: Tinsley vs. Chinook



Marion Tinsley

World champion for over 40 years

**VS** 

Chinook

Developed by Jonathan Schaeffer, professor at the U. of Alberta

1992: Tinsley beat Chinook in 4 games to 2,

with 33 draws.

1994: 6 draws

In 2007, Schaeffer announced that checkers was solved, and anyone playing against Chinook would only be able to draw, never win.

Play against Chinook: http://games.cs.ualberta.ca/cgi-bin/player.cgi?nodemo

#### 1997 - Othello: Murakami vs. Logistello



Takeshi Murakami World Othello (aka Reversi) champion

**VS** 

Logistello

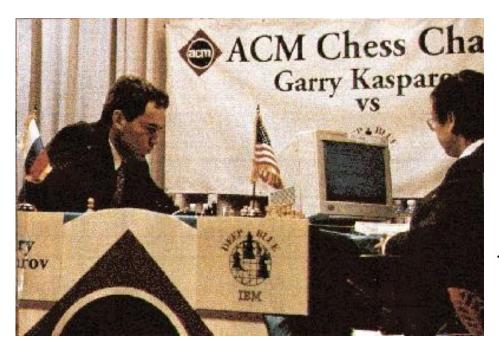
developed by Michael Buro runs on a standard PC

https://skatgame.net/mburo/log.html

(including source code)

Logistello beat Murakami by 6 games to 0

#### 1997- Chess: Kasparov vs. Deep Blue



Garry Kasparov
50 billion neurons
2 positions/sec
VS
Deep Blue

32 RISC processors + 256 VLSI chess engines 200,000,000 pos/sec

Deep Blue wins by 3 wins, 1 loss, and 2 draws

#### 2003 - Chess: Kasparov vs. Deep Junior



Garry Kasparov

still 50 billion neurons still 2 positions/sec

VS

Deep Junior

8 CPU, 8 GB RAM, Win 2000 2,000,000 pos/sec Available at \$100

Match ends in a 3/3 tie!

#### 2016 - Go: AlphaGo vs Lee Se-dol

- GO was always considered a much harder game to automate than chess because of its very high branching factor (35 for chess vs 250 for Go!)
- In 2016, AlphaGo beat Lee Sedol in a five-game match of GO.
- In 2017 AlphaGo beat Ke Jie, the world No.1 ranked player at the time
- uses a Monte Carlo tree search algorithm to find its moves based on knowledge previously "learned" by deep learning



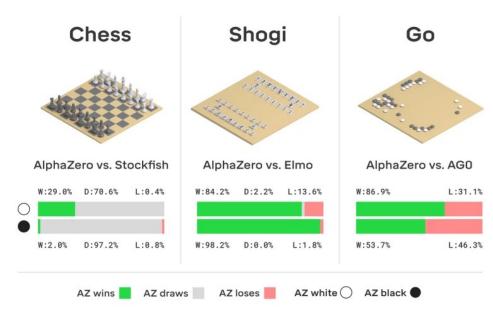
#### 2017 - AlphaGo Zero & AlphaZero

AlphaGo Zero learned the Game by itself, without input of human games

- Became better than all old versions after 40 days of training
- In the first three days, AlphaGo Zero played 4.9 million games against itself using reinforcement learning

AlphaZero can learn other games, like Chess and Shogi

- In 2018, it beat the thenbest chess program, Stockfish 8 in a 100-game tournament
- Trained using 5,000 tensor processing units (TPUs), run on four TPUs and a 44-core CPU during matches



#### 2018 - AlphaZero vs Stockfish 8



# 2019 - Deep learning to answer math questions

Ongoing work on solving other problems with a general AI

 In 2019, Google engineers published on their work training a neural network system to answer math questions, like

What is the sum of 1+1+1+1+1+1?

- The system's answer? 6
- But it did solve 14 out of 40 questions on a standard test correctly

See <a href="https://arxiv.org/abs/1904.01557">https://arxiv.org/abs/1904.01557</a>

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