

The Role of Software in HPC – Lessons Learnt in the US Exascale Computing Project



Approved for public release

Hartwig Anzt, University of Tennessee



The US Exascale Computing Project



Addressing a National Imperative

The Exascale Computing Project is an aggressive research, development, and deployment project focused on delivery of mission-critical applications, an integrated software stack, and exascale hardware technology advances.

Application Development



Software Technology

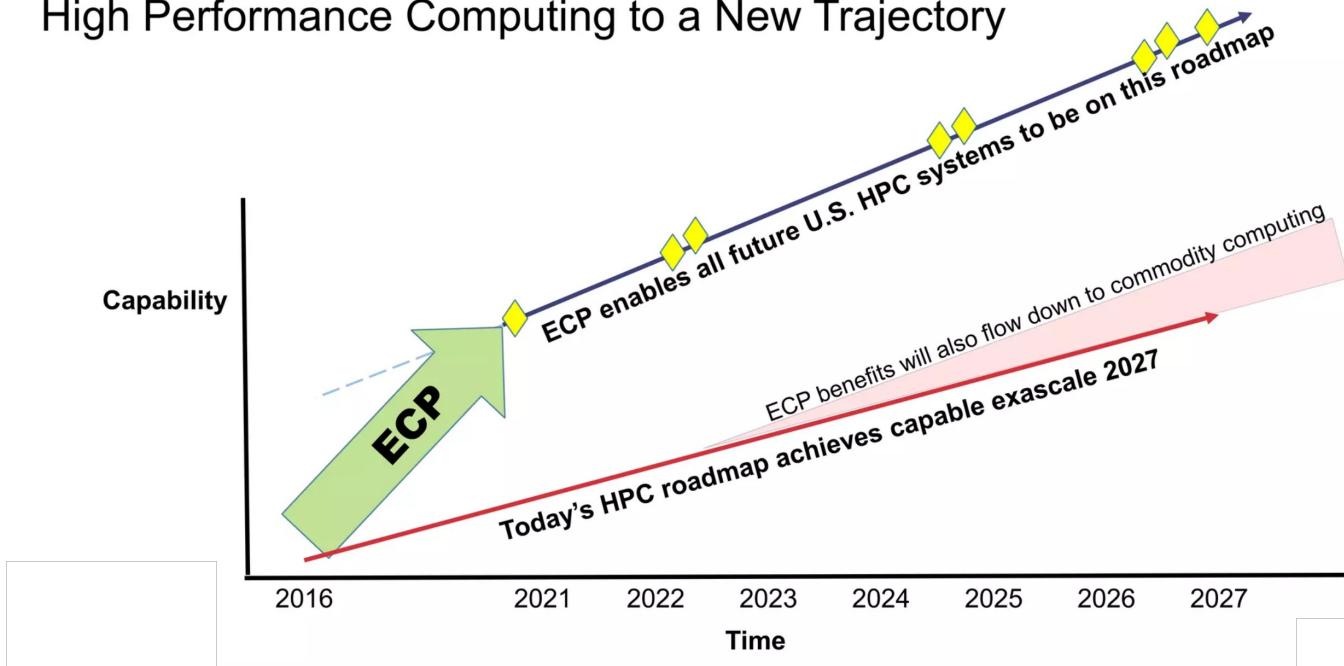


Hardware & Integration



Vision: Exascale Computing Project (ECP) Lifts all U.S. High Performance Computing to a New Trajectory

© Paul Messina



The US Exascale Computing Project



Advancing Scientific Discovery

The ECP aims to ensure availability of the exascale computing ecosystem necessary for developing clean energy systems, improving the resilience of our infrastructure, designing new materials that can perform in extreme environments, adapting to changes in the water cycle, developing smaller and more powerful accelerators for use in medicine and industry, and much more. Several projects focus on data-intensive problems to enable effective use of the data streams from powerful scientific facilities, complex environmental genomes, and cancer research (patient genetics, tumor genomes, molecular simulations, and clinical data).



Strengthening National Security

The ECP teams are also developing new applications for supporting the NNSA Stockpile Stewardship Program, which is responsible for maintaining the readiness and reliability of our nuclear weapons systems—without underground testing. Assessing the performance of weapons systems subject to hostile environments and potential threat scenarios exceeds the capabilities of current HPC systems and codes. NNSA application projects are focused on providing the sophisticated modeling and analysis tools needed to sustain the U.S. nuclear deterrence.



Improving Industrial Competitiveness

Exascale systems will be used to accelerate research that leads to innovative products and speeds commercialization, creating jobs and driving US competitiveness across industrial sectors, such as the emerging energy economy. To ensure alignment with US industry needs, the ECP is engaging senior technology decision makers from among the country's most prominent private sector companies.

The US Exascale Computing Project



US\$4B – what is it spent on?

- 3 computers
 - \$600M each
 - \$400M to vendors for Design, Path, Fast - Forward
- 21 Applications



AMD Based
(Up & running)



Intel Based
(Being installed)



AMD Based
(Planned)

Sustainable software development

Domain*	Base Challenge Problem	Domain*	Challenge Problem	PMR Core (17)	Compilers and Support (7)	Tools and Technology (11)	xSDK (16)	Visualization Analysis and Reduction (9)	Data mgmt, I/O Services, Checkpoint restart (12)	Ecosystem/E4S at-large (12)
Wind Energy	2x2 5 MW turbine array in 3x3x1 km ³ domain	Quantum Materials	Predict & control mats. @ quantum level	QUO	openarc	TAU	hypre	ParaView	SCR	mpiFileUtils
Nuclear Energy	Small Modular Reactor with complete in-vessel coolant loop	Astrophysics	Supernovae explosions, neutron star mergers	Papyrus	Kitsune	HPC Toolkit	FleCSI	Catalyst	FAODEL	TriBITS
Fossil Energy	Burn fossil fuels cleanly with CLRs	Cosmology	Extract "dark sector" physics from upcoming cosmological surveys	SICM	LLVM	Dyninst Binary Tools	MFEM	VTK-m	ROMIO	MarFS
Combustion	Reactivity controlled compression ignition	Earthquakes	Regional hazard and risk assessment	Legion	CHILL autotuning comp	Gotcha	Kokkoskernels	SZ	Mercury (Mochi suite)	GUFI
Accelerator Design	TeV-class 10 ²⁻³ times cheaper & smaller	Geoscience	Well-scale fracture propagation in wellbore cement due to attack of CO ₂ -saturated fluid	Kokkos (support)	LLVM openMP comp	Caliper	Trilinos	zfp	HDF5	Intel GEOPM
Magnetic Fusion	Coupled gyrokinetics for ITER in H-mode	Earth System	Assess regional impacts of climate change on the water cycle @ 5 SYPD	RAJA	OpenMP V & V	PAPI	SUNDIALS	Visit	Parallel netCDF	BEE
Nuclear Physics: QCD	Use correct light quark masses for first principles light nuclei properties	Power Grid	Large-scale planning under uncertainty; underfrequency response	PaRSEC*	Flang/LLVM Fortran comp	Program Database Toolkit	PETSc/TAO	ASCENT	ADIOS	FSEFI
Chemistry: GAMESS	Heterogeneous catalysis: MSN reactions	Cancer Research	Scalable machine learning for predictive preclinical models and targeted therapy	DARMA	GASNet-EX	Search (random forests)	libEnsemble	Cinema	Darshan	Kitten Lightweight Kernel
Chemistry: NWChemEx	Catalytic conversion of biomass	Metagenomics	Discover and characterize microbial communities through genomic and proteomic analysis	Qthreads	BOLT	Siboka	STRUMPACK	ROVER	UnifyCR	COOLR
Extreme Materials	Microstructure evolution in nuclear mats.	FEL Light Source	Protein and molecular structure determination using streaming light source data	UPC++	MPICH	C2C	SuperLU	Veloc	NRM	ArgoContainers
Additive Manufacturing	Born-qualified 3D printed metal alloys			Open MPI	Tasmanian	Sonar	ForTrilinos	IOSS	HXHIM	Spack
				Umpire			SLATE			PMR
				AML			MAGMA			Tools
							DTK			Math Libraries
							Tasmanian			Legend
							Ginkgo			Data and Vis
										Ecosystems and delivery



Designing an ECP library for sustaining simulation performance



Building Trusted Scientific Software

SHARE in f t o

PUBLISHED JUN 28, 2018 AUTHOR MIKE HER

Software Verification

SHARE in f t o

PUBLISHED AUG 15, 2018 AUTHOR ANSHU I

Think Locally, Act Globally: Outreach for Better Scientific Software

SHARE in f t o

I have worked in the scientific software field for more than 20 years. One phrase I often hear is “Verification is doing things right, and validation is checking to see if things are done right.” I have had to memorize this phrase to memory in order to avoid confusion when the

Pairing internal and external concerns

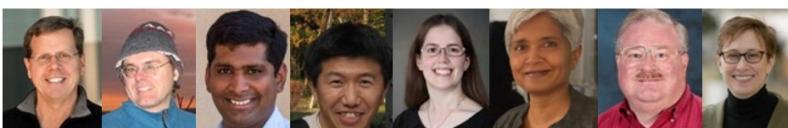
Verification focuses on internal concerns of a good software system. Validation focuses on external concerns of a system, such as the

Helping code teams improve their software development, productivity, and sustainability is no small challenge. In the IDEAS Productivity project, we have found that one of the keys to aiding the Exascale Computing Project (ECP) software development teams involves extensive outreach to the broader community of computational scientists and engineers (CSE) in high-performance computing (HPC).

PUBLISHED JUL 17, 2018 AUTHOR DAVID BERNHOLDT TOPICS BETTER SKILLS ↗ PERSONAL PRODUCTIVITY AND SUSTAINABILITY

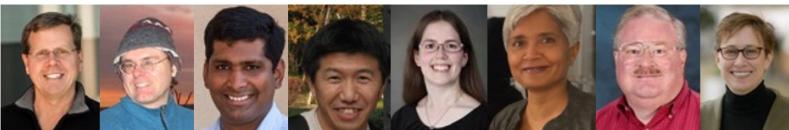
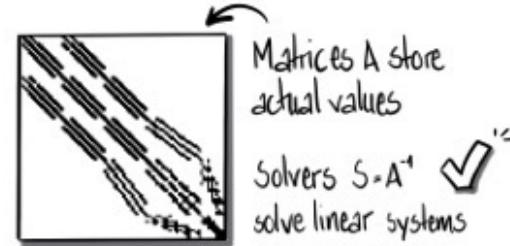
An ambitious goal

The ECP needs to deliver a software environment and applications ready to run on exascale computers, which are scheduled to be deployed starting in 2021. Achieving this goal entails a major, large-scale software development effort. Recognizing the challenges development teams will face, the ECP is supporting the IDEAS Productivity project to help scientific researchers improve their development practices.



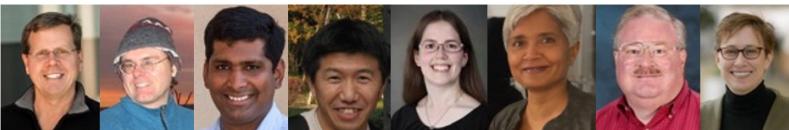
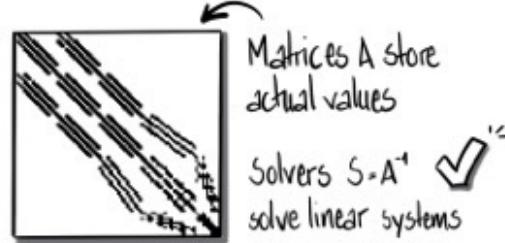
Designing an ECP library for sustaining simulation performance

Ginkgo - A sparse linear algebra library for HPC



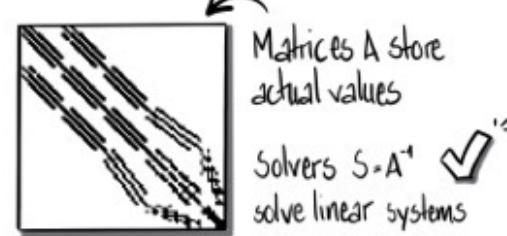
Designing an ECP library for sustaining simulation performance

written in C++ → Ginkgo - A sparse linear algebra library for HPC



Designing an ECP library for sustaining simulation performance

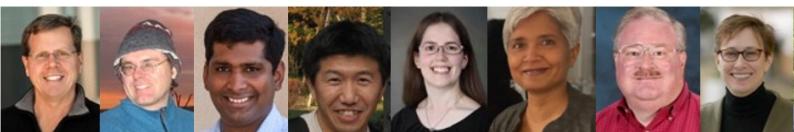
written in C++ → Ginkgo - A sparse linear algebra library for HPC



Contains architecture-
agnostic algorithm
implementation

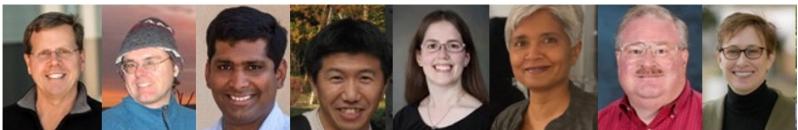
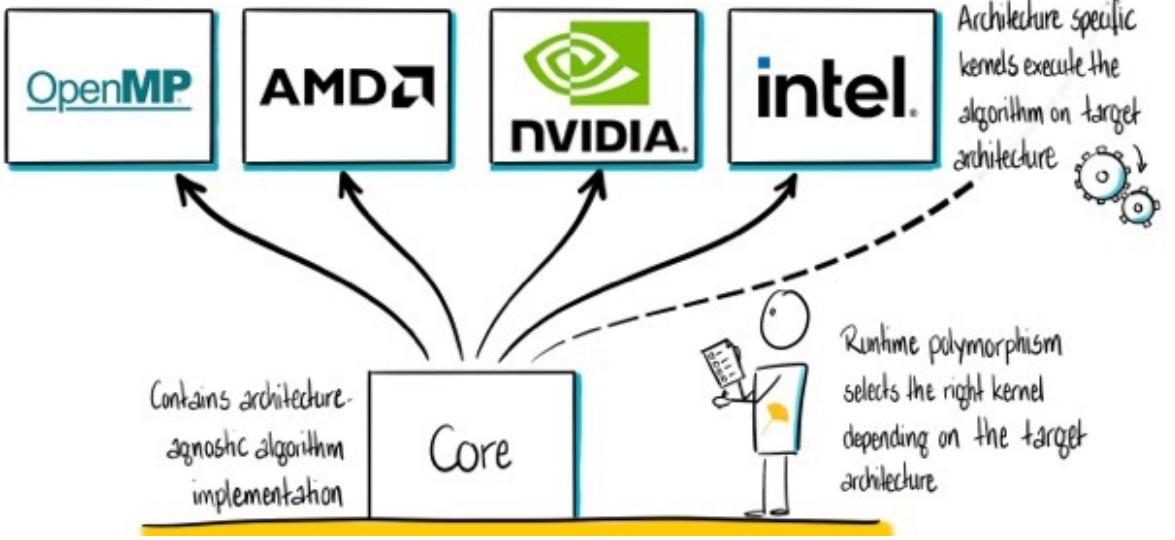
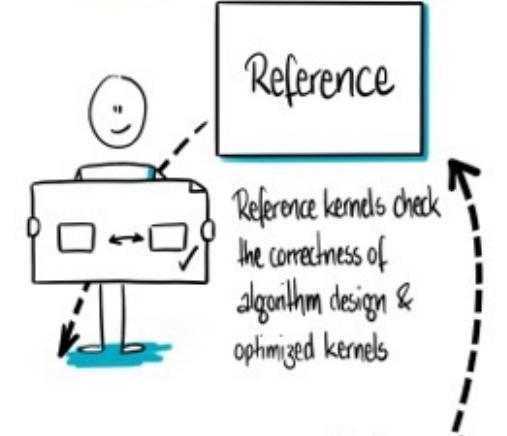
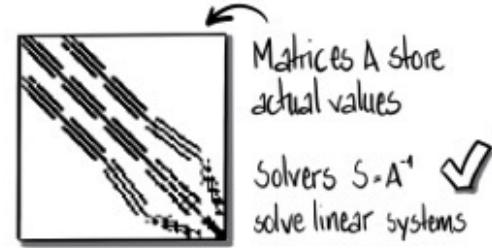
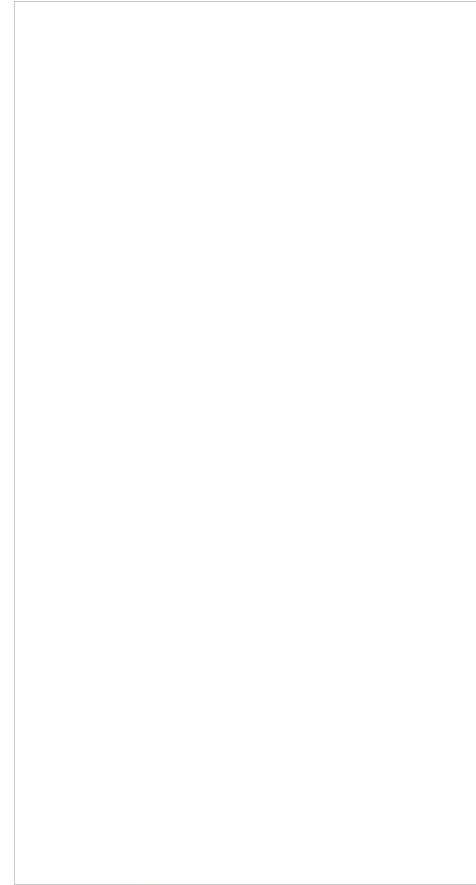
Core

Runtime polymorphism
selects the right kernel
depending on the target
architecture

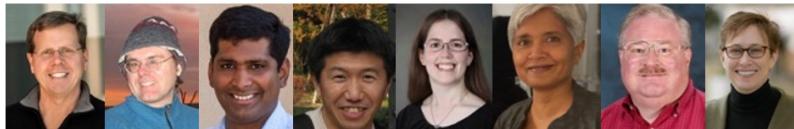
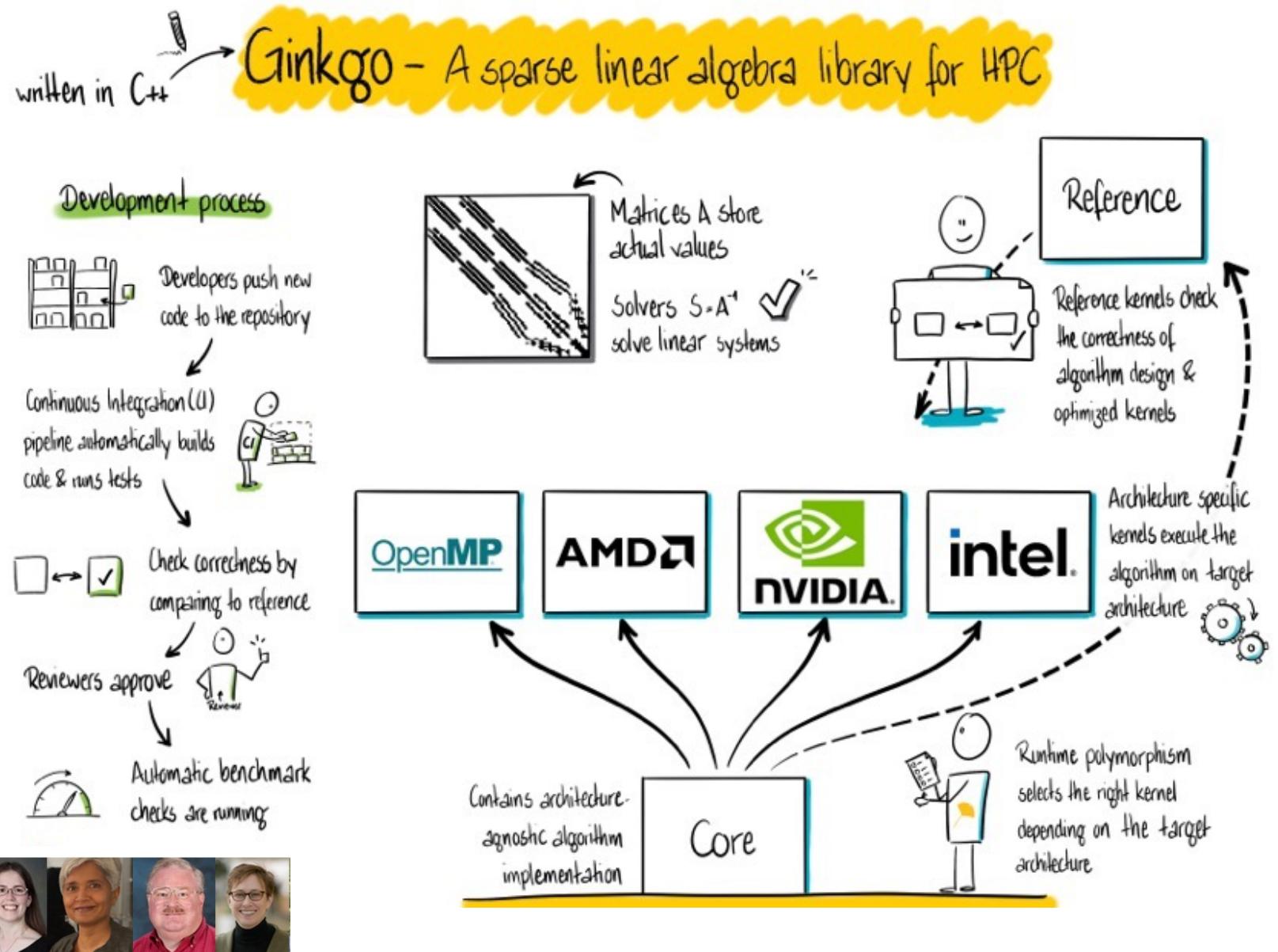


Designing an ECP library for sustaining simulation performance

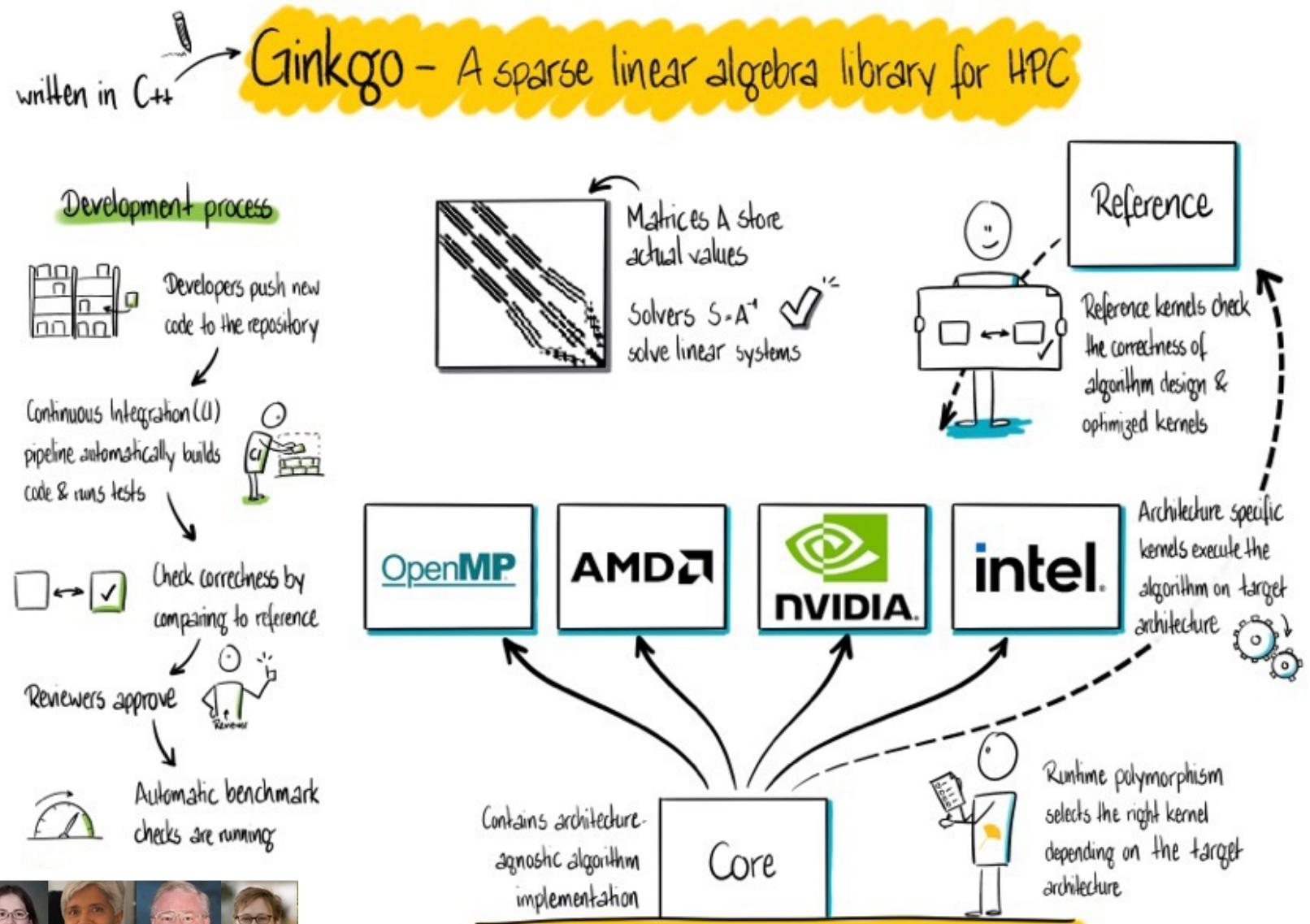
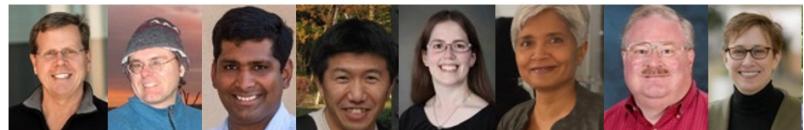
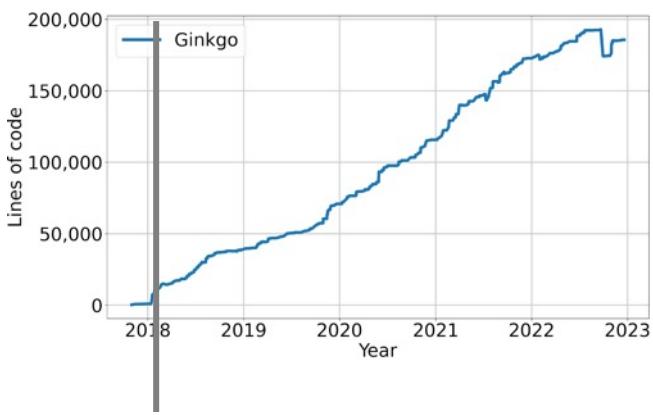
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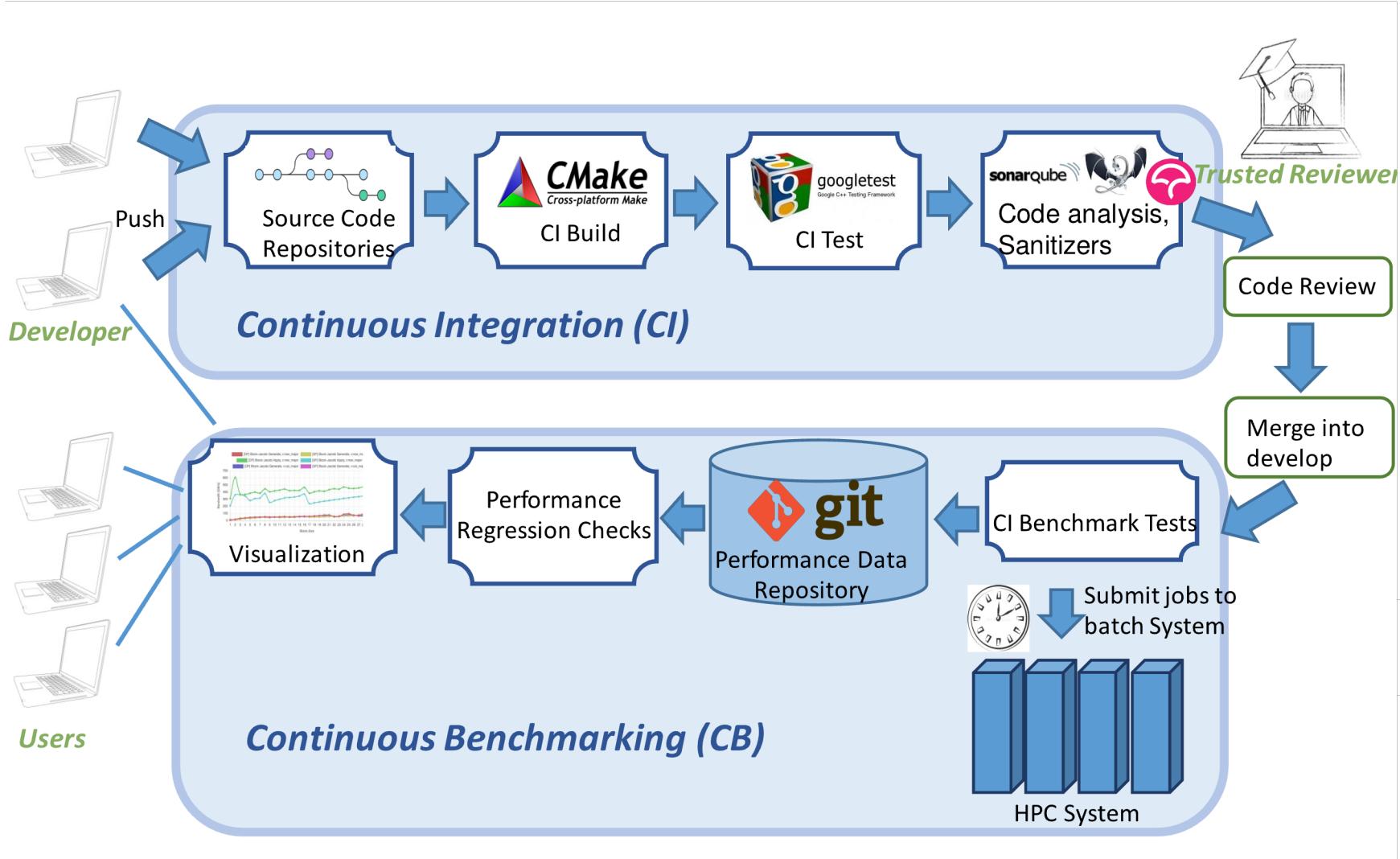
Designing an ECP library for sustaining simulation performance



Designing an ECP library for sustaining simulation performance



Sustainable software development & CI/CD



Starting with the CUDA backend

Linear Operator Interface

We express everything as Linear Operator.

- Internally, we leverage C++ class inheritance.
- Applications can apply any functionality as a linear operator.

Matrix-Vector Product

$$x := A \cdot b$$

Preconditioner (for matrix A)

$$\begin{aligned} x &:= M^{-1} \cdot b \\ M^{-1} &\approx A^{-1} \end{aligned}$$

Solver (for system $Ax = b$)

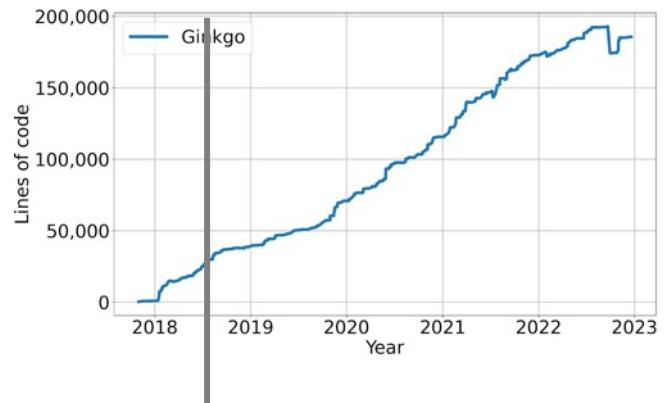
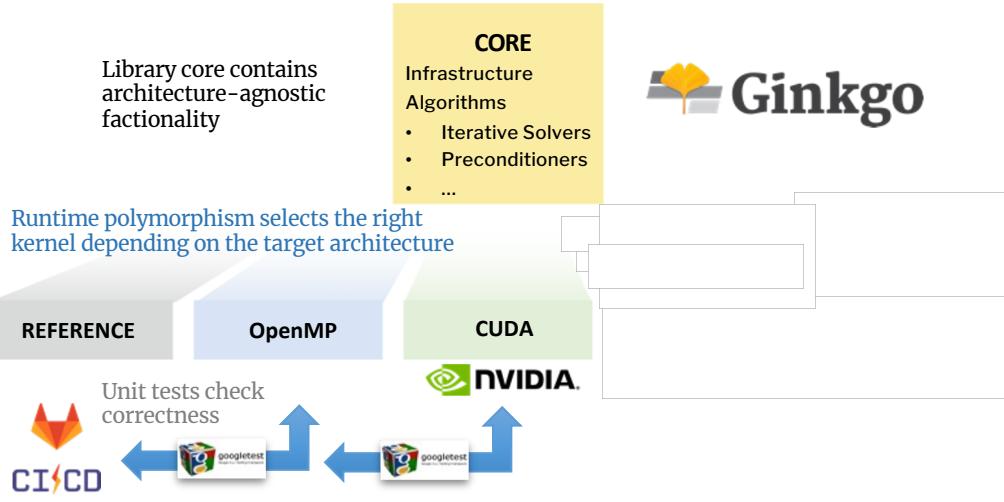
$$M^{-1} = \Pi(A)$$

$$S \approx A^{-1}$$

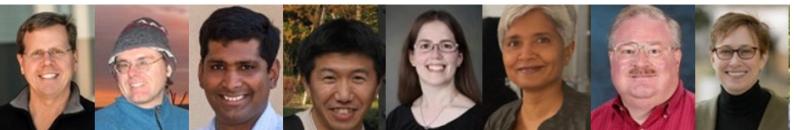
$$S = \Sigma(A)$$

All of them can be expressed as

Application of a linear operator* (LinOp) $L : \mathbb{F}^m \rightarrow \mathbb{F}^m$

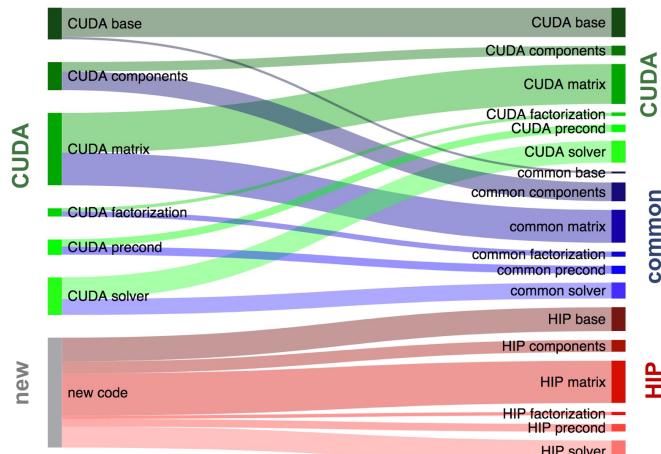


	OMP	CUDA
Basic		
SpMV	✓	✓
SpMM	✓	✓
SpGeMM	✓	✓
Krylov solvers		
BiCG	✓	✓
BICGSTAB	✓	✓
CG	✓	✓
CGS	✓	✓
GMRES	✓	✓
IDR	✓	✓
(Block-)Jacobi	✓	✓
ILU/IC		✓
Parallel ILU/IC	✓	✓
Parallel ILUT/ICT	✓	✓
Sparse Approximate Inverse	✓	✓
Preconditioners		



Extending to AMD GPUs

The screenshot shows a GitHub blog post from the Ginkgo repository. The title is "Porting the Ginkgo Package to AMD's HIP Ecosystem". The post discusses the experience of porting CUDA code to AMD's Heterogeneous-compute Interface for Portability (HIP). It includes a section on hardware diversification, portability, and sustainability. The post was published on June 25, 2020, by Hartwig Anzt. Topics covered include better reliability, better planning, testing, and design. The post also features a diagram illustrating the porting process from CUDA to HIP.



Library core contains architecture-agnostic functionality

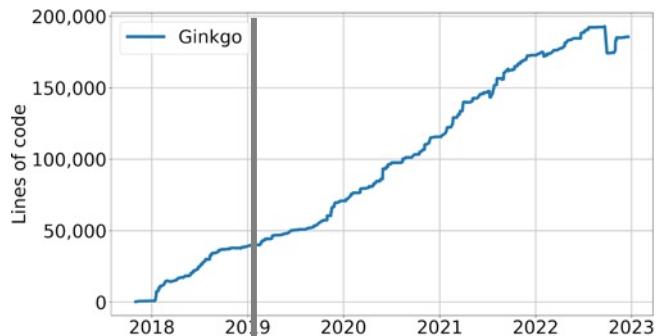
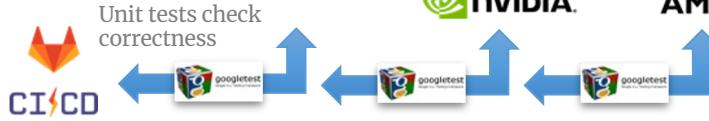
CORE
Infrastructure Algorithms
• Iterative Solvers
• Preconditioners
• ...

Runtime polymorphism selects the right kernel depending on the target architecture

REFERENCE OpenMP

CUDA NVIDIA

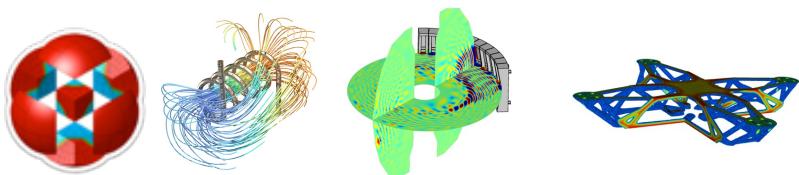
HIP AMD



	OMP	CUDA	HIP
Basic	✓	✓	✓
SpMV	✓	✓	✓
SpMM	✓	✓	✓
SpGeMM	✓	✓	✓
Krylov solvers			
BiCG	✓	✓	✓
BICGSTAB	✓	✓	✓
CG	✓	✓	✓
CGS	✓	✓	✓
GMRES	✓	✓	✓
IDR	✓	✓	✓
(Block-)Jacobi	✓	✓	✓
ILU/IC		✓	✓
Parallel ILU/IC	✓	✓	✓
Parallel ILUT/ICT	✓	✓	✓
Sparse Approximate Inverse	✓	✓	✓
Preconditioners			



Input from the “first customer”



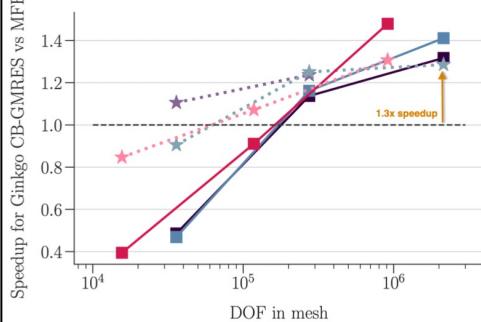
MFEM is a *free, lightweight, scalable C++ library* for finite element methods.

Speeding up MFEM’s “example 22” on GPUs

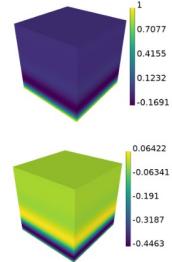
Example 22 of the MFEM finite element library solves harmonic oscillation problems, with a forced oscillation imposed at the boundary. In this test, we use variant 1:

$$-\nabla \cdot (a \nabla u) - \omega^2 b u + i \omega c u = 0$$

with $a = 1$, $b = 1$, $\omega = 10$, $c = 20$



- p = 1 (V100)
- ★··· p = 1 (MI50)
- p = 2 (V100)
- ★··· p = 2 (MI50)
- p = 3 (V100)
- ★··· p = 3 (MI50)



Real part of solution (top),
imaginary part of solution

Speedup of Ginkgo’s Compressed Basis-GMRES solver vs MFEM’s GMRES solver for three different orders of basis functions (p), using MFEM matrix-free operators and the Ginkgo-MFEM integration wrappers in MFEM. CUDA 10.1/V100 and ROCm 4.0/MI50.



Library core contains architecture-agnostic functionality

CORE

- Infrastructure Algorithms
- Iterative Solvers
 - Preconditioners
 - ...



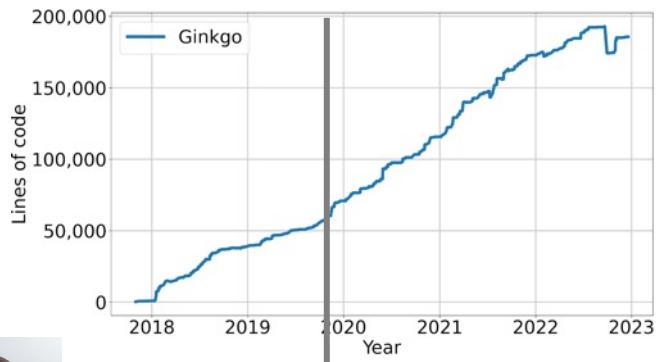
Runtime polymorphism selects the right kernel depending on the target architecture

Architecture-optimized kernels

REFERENCE OpenMP



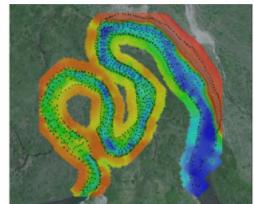
CI/CD



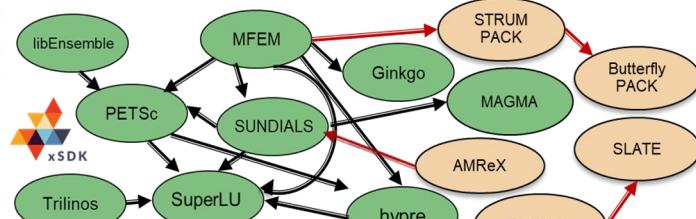
	OMP	CUDA	HIP
Basic			
SpMV	✓	✓	✓
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Krylov solvers			
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BICGSTAB	✓	✓	✓
CG	✓	✓	✓
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IDR	✓	✓	✓
(Block-)Jacobi	✓	✓	✓
ILU/IC		✓	✓
Parallel ILU/IC	✓	✓	✓
Parallel ILUT/ICT	✓	✓	✓
Sparse Approximate Inverse	✓	✓	✓
Preconditioners			
On-Device Matrix Assembly	✓	✓	✓
MC64/RCM reordering	✓		
Wrapping user data		✓	
Utilities			

Part of the xSDK effort

xSDK: Extreme-scale Scientific Software Development Kit



Integrated surface-subsurface hydrology simulations of river meanders require the combined use of xSDK packages.



xSDK-examples v.0.3.0

The xSDK provides infrastructure for and interoperability of a **collection of related and complementary software elements**—developed by diverse, independent teams throughout the high-performance computing (HPC) community—that provide the building blocks, tools, models, processes, and related artifacts for rapid and efficient development of high-quality applications.

November 2022

- 26 math libraries
- 2 domain components
- 16 mandatory xSDK community policies
- Spack xSDK installer

xSDK community policies:

- 16 mandatory policies,
- 8 recommended policies,
- 4 Spack variant guidelines
- Available on Github
<https://xSDK.info/policies/>



Library core contains architecture-agnostic functionality

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- Iterative Solvers
- Preconditioners
- ...

Runtime polymorphism selects the right kernel depending on the target architecture

REFERENCE

OpenMP

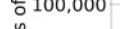
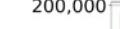
CUDA

HIP



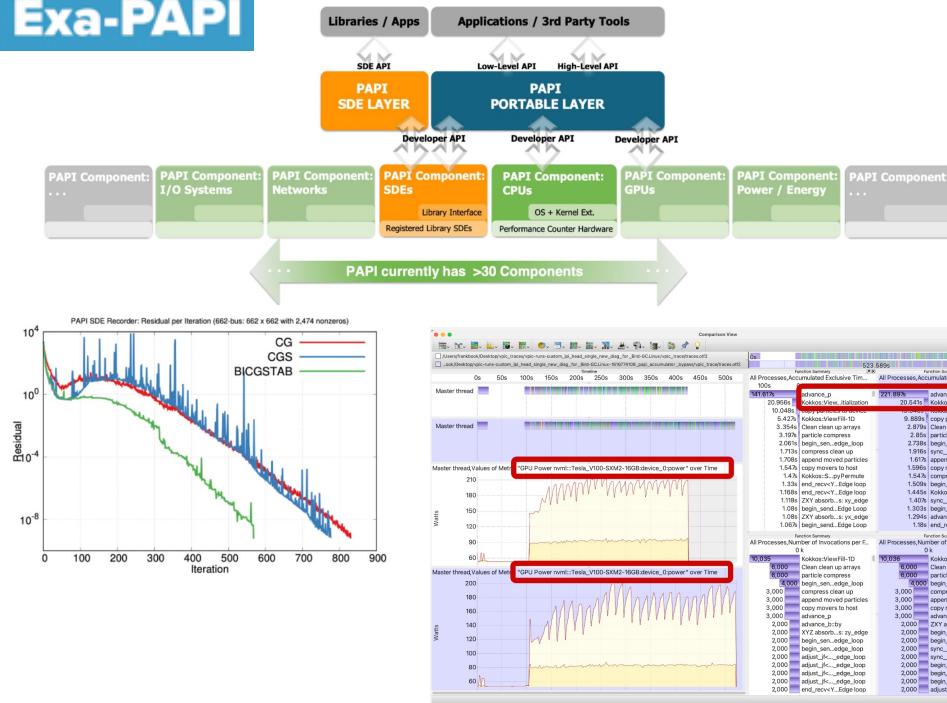
CI/CD

Unit tests check correctness



Adding profiling functionality

Exa-PAPI

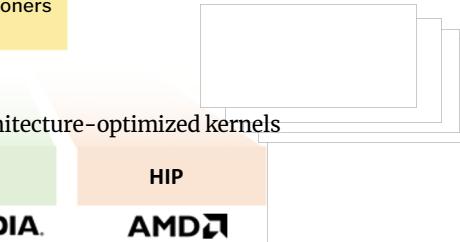


Library core contains architecture-agnostic functionality

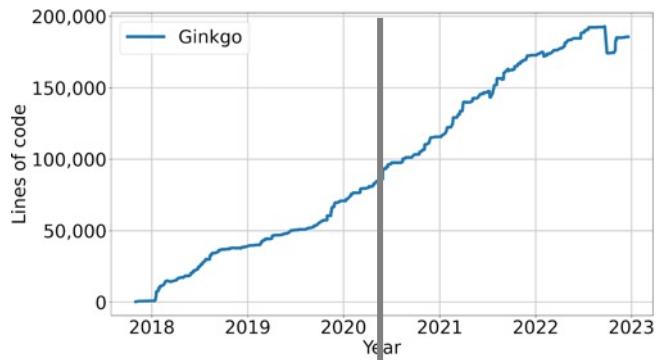
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Sparse Approximate Inverse	✓	✓	✓
Preconditioners			
On-Device Matrix Assembly	✓	✓	✓
MC64/RCM reordering	✓		
Wrapping user data		✓	
Logging		✓	
PAPI counters		✓	



Extending to Intel GPUs

Preparing for the Arrival of Intel's Discrete High-Performance GPUs

By Hartwig Anzt

March 23, 2021

[yhmtsa1/try_oneapi](#) Private

Code Issues Pull requests Actions Projects Security Insights

master 1 branch 0 tags Go to file Add file ▾

yhmtsa1 format update	69d2173 on Aug 7, 2022	70 commits
arg_struct	WIP	2 years ago
atomic	atomic and get_in_template	2 years ago
check_uninit	some checker	last year
classical_csr	cpu barrier issue in classical csr spmv	last year
clinfo	clinfo	2 years ago
coop_cuda	keep some history but I do not check them detail	last year
coop_draft	keep some history but I do not check them detail	last year

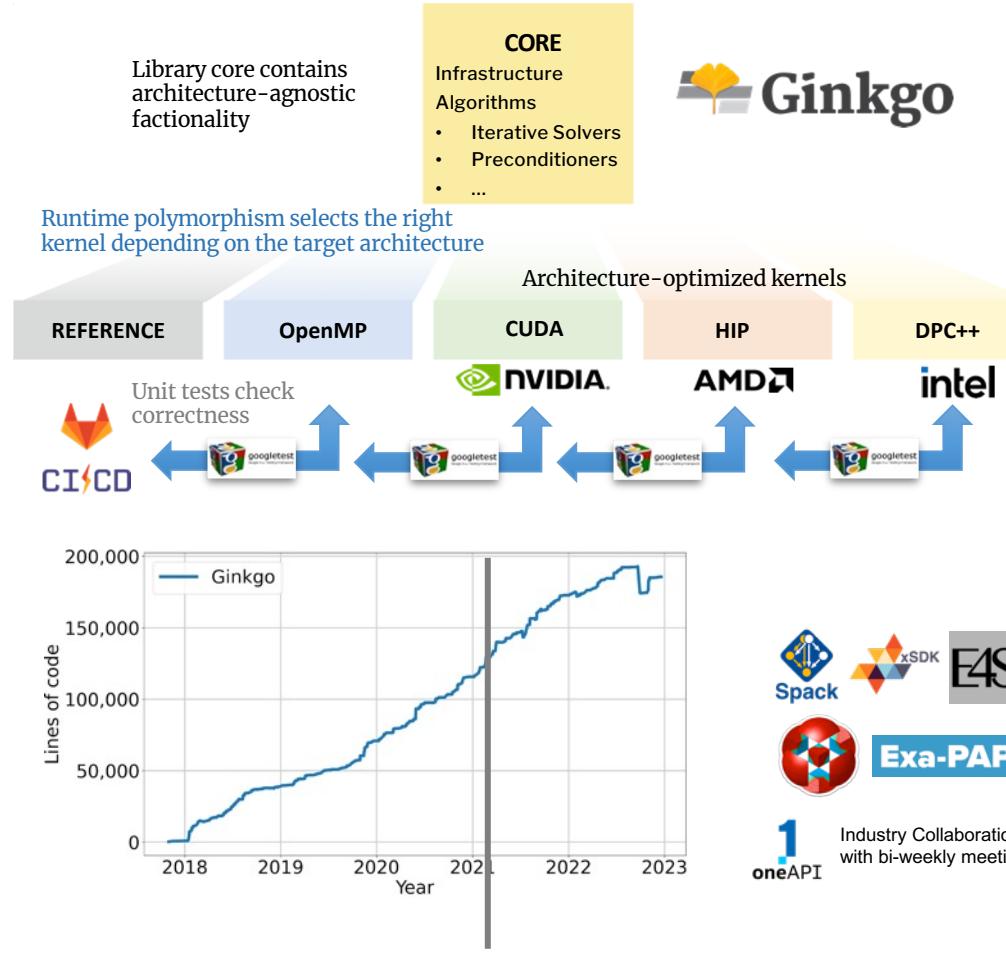
About

No description, website, or topics provided.

Readme 0 stars 2 watching 0 forks

Releases

No releases published Create a new release



	Functionality	OMP	CUDA	HIP	DPC++
Basic	SpMV	✓	✓	✓	✓
	SpMM	✓	✓	✓	✓
	SpGeMM	✓	✓	✓	✓
Krylov solvers	BICG	✓	✓	✓	✓
	BICGSTAB	✓	✓	✓	✓
	CG	✓	✓	✓	✓
Preconditioners	CGS	✓	✓	✓	✓
	GMRES	✓	✓	✓	✓
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Utilities	(Block-)Jacobi	✓	✓	✓	✓
	ILU/IC		✓	✓	✓
	Parallel ILU/IC	✓	✓	✓	✓
	Parallel ILUT/ICT	✓	✓	✓	✓
	Sparse Approximate Inverse	✓	✓	✓	✓
<hr/>					
Utilities	On-Device Matrix Assembly	✓	✓	✓	✓
	MC64/RCM reordering	✓			
	Wrapping user data		✓		
	Logging		✓		
Utilities	PAPI counters		✓		



Extending to Intel GPUs



- Bi-Weekly technical meetings with Intel
- Long list of bug reports, feature requests, performance data discussions, documentation improvements ...

cuBLAS backend (and potentially other domains) fails with latest LLVM builds
#223

Summary
As first observed in #219 many tests in cuBLAS backend is failing with latest LLVM builds.

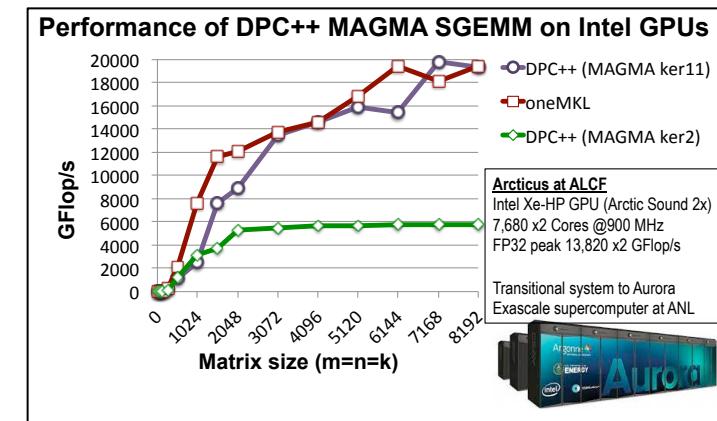
Version
I have tried LLVM commit: 66361038b63caaae566fc9648f5da50b74222b83 and got the below tests failing (showing only a few of them)

```
1 - BLAS/RT/Nrm2TestSuite/Nrm2Tests.RealSinglePrecision/Column_Major_TITAN_RTX (Failed)
3 - BLAS/RT/Nrm2TestSuite/Nrm2Tests.RealDoublePrecision/Column_Major_TITAN_RTX (Failed)
7 - BLAS/RT/Nrm2TestSuite/Nrm2Tests.ComplexDoublePrecision/Column_Major_TITAN_RTX (Failed)
17 - BLAS/RT/LamaXTestSuite/LamaTests.RealSinglePrecision/Column_Major_TITAN_RTX (Failed)
19 - BLAS/RT/LamaXTestSuite/LamaTests.RealDoublePrecision/Column_Major_TITAN_RTX (Failed)
23 - BLAS/RT/LamaXTestSuite/LamaTests.ComplexDoublePrecision/Column_Major_TITAN_RTX (Failed)
27 - BLAS/RT/DotxTestSuite/DotxTests.ComplexSinglePrecision/Column_Major_TITAN_RTX (Failed)
35 - BLAS/RT/DotxTestSuite/DotxTests.ComplexSinglePrecision/Column_Major_TITAN_RTX (Failed)
67 - BLAS/RT/AsmTestSuite/AsmTests.ComplexSinglePrecision/Column_Major_TITAN_RTX (Failed)
81 - BLAS/RT/ScalTestSuite/ScalTests.RealSinglePrecision/Column_Major_TITAN_RTX (Failed)
85 - BLAS/RT/ScalTestSuite/ScalTests.ComplexSinglePrecision/Column_Major_TITAN_RTX (Failed)
```

From [DPCPP AoT documentation](#), not clear:

- The options are also required at linking time? Unused in files without kernels?
- Any example of other projects integrating AoT in a CMake setup?

[Intel Compiler \(Fortran/C/C++/L0\) - Intel Discrete GPU Accelerator - Joint Laboratory for System Evaluation \(anl.gov\)](#)
hang_atomic_on_local
Ticket number: CMPLRLLVM-36572 (works in PVC, but still fails on ATS node)
related to driver not compiler self



... but also docker image contributions and bug fixes!

tid % subgroup size >= 4 gives wrong division

(double) 1/a gives wrong result when the tid % subgroup size >= 4. For example, when a = 1.07338829563753890
1/a should be 0.9316293125835232
if (local_id == assign_id) { a = double(1/a); }
when assign_id < 4, Gen9 GPU still give the correct result
when assign_id >= 4, Gen9 GPU gives wrong 0.9316293591.0000000506
CPU has more worse result

It is connected to optimizations (not reproducible with O0). If fp-speculation=off do not improve results.
Ticket number: XDEPS-4031 ()

Devcloud node issue

- sycl-ls/clinfo does not give any output for nodes s001-n225, s011-n006
- no gpu on the nodes s001-n232, s011-n233, s011-n008
- aithub.com is not accessible on login-node

Fix cuda/hip backend location #219

Merged mkrainuk merged 2 commits into oneapi-src:develop from yhmtsaifix_cuda_backend_location 20 days ago

Description
From intel/llvm#6407, it moves almost all headers from CL/sycl to sycl
I followed #199 way
make the header can use sycl/* if they exist and allow the old intel llvm.
I also update the CL/sycl.hpp which are not changed before.

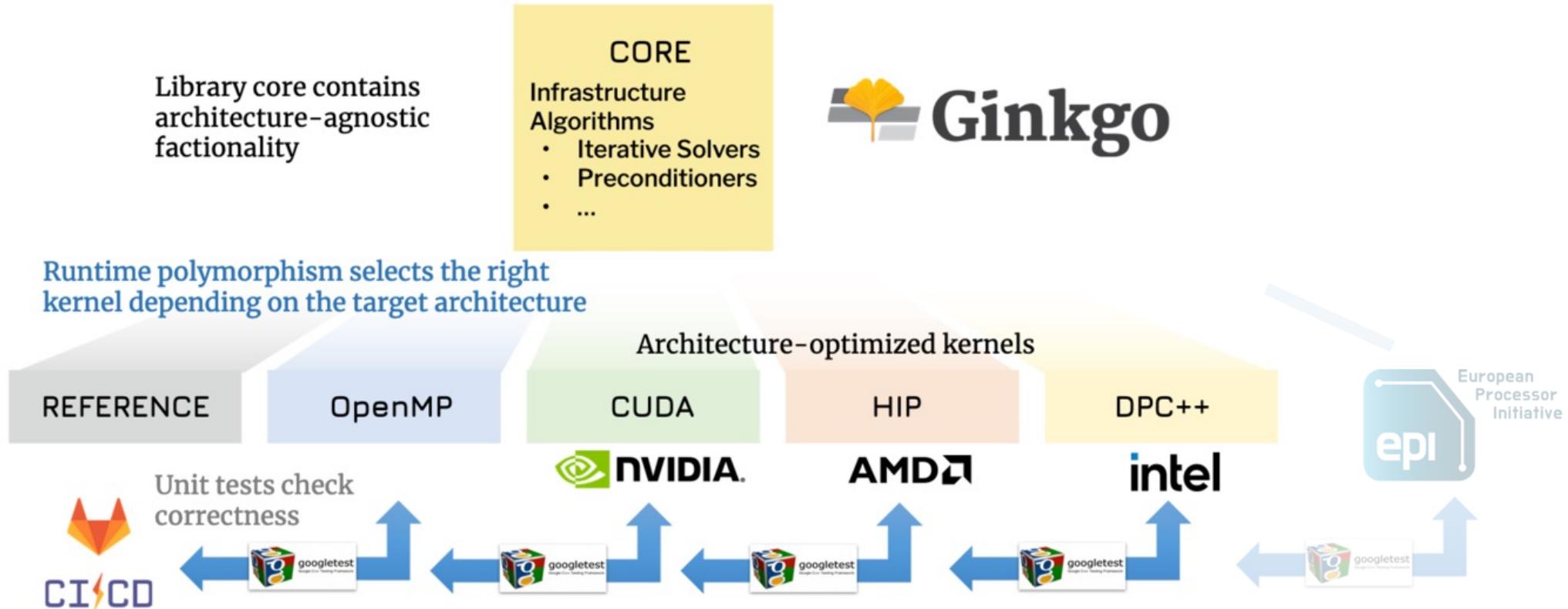
All Submissions
Do all unit tests pass locally? Attach a log. A: It is a compiling issue.
Have you formatted the code using clang-format?

Bug fixes
Have you added relevant regression tests? A: It is a compiling issue.
Have you included information on how to reproduce the issue (either in a GitHub issue or in this PR)?
Reproduce:
compile the latest intel llvm and this repo, it will not be able to compile due to missing headers.

yhmtsaifixed 2 commits 2 months ago
use the correct sycl path after intel/llvm#6407
fix the missing sycl/sycl.hpp

mkmrainerel commented on Aug 1 edited
@yhmtsaif Thanks for the PR. Is the description from sycl/CL to CL correct? My understanding is all header files moved

Portability as central design principle



Focus efforts as lightweight tool in ECP to address challenges



Focus efforts

- Mixed precision
- batched
- *Address recent hardware trends (tensor cores, etc.)*
- *Address hardware requirements*



Library core contains architecture-agnostic functionality

CORE
Infrastructure Algorithms

- Iterative Solvers
- Preconditioners
- ...



Runtime polymorphism selects the right kernel depending on the target architecture

REFERENCE

OpenMP



Unit tests check correctness



CUDA



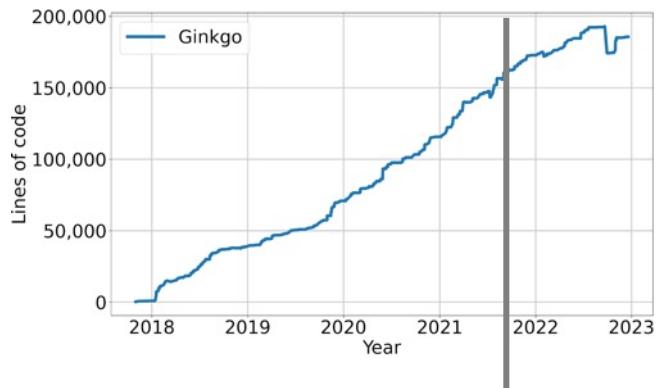
Architecture-optimized kernels



HIP



DPC++



	OMP	CUDA	HIP	DPC++
Basic				
SpMV	✓	✓	✓	✓
SpMM	✓	✓	✓	✓
SpGeMM	✓	✓	✓	✓
Krylov solvers				
BiCG	✓	✓	✓	✓
BICGSTAB	✓	✓	✓	✓
CG	✓	✓		✓
CGS	✓	✓	✓	✓
GMRES	✓	✓	✓	✓
IDR	✓	✓	✓	✓
(Block-)Jacobi	✓	✓	✓	✓
ILU/IC		✓	✓	✓
Parallel ILU/IC	✓	✓	✓	✓
Parallel ILUT/ICT	✓	✓	✓	✓
Sparse Approximate Inverse	✓	✓	✓	✓
Preconditioners				
On-Device Matrix Assembly	✓	✓	✓	✓
MC64/RCM reordering	✓			
Wrapping user data		✓		
Logging		✓		
PAPI counters		✓		



Mixed precision focus effort



Focus efforts

- Mixed precision
- batched
- *Address recent hardware trends (tensor cores, etc.)*
- *Address hardware requirements*



Advances in Mixed Precision Algorithms: 2021 Edition

by the ECP Multiprecision Effort Team (Lead: Hartwig Anzt)

Ahmad Abdelfattah, Hartwig Anzt, Alan Ayala, Erik G. Boman, Erin Carson, Sebastien Cayrols, Terry Cojean, Jack Dongarra, Rob Falgout, Mark Gates, Thomas Grützmacher, Nicholas J. Higham, Scott E. Kruger, Sherry Li, Neil Lindquist, Yang Liu, Jennifer Loe, Piotr Luszczek, Pratik Nayak, Daniel Osei-Kuffuor, Sri Pranesh, Sivasankaran Rajamanickam, Tobias Ribizel, Barry Smith, Kasia Swirydowicz, Stephen Thomas, Stanimire Tomov, Yaohung M. Tsai, Ichi Yamazaki, Urike Meier Yang

Available access | Research article | First published online March 19, 2021

A survey of numerical linear algebra methods utilizing mixed-precision arithmetic

Ahmad Abdelfattah, Hartwig Anzt and Ulrike Meier Yang View all authors and affiliations

Volume 35, Issue 4 | <https://doi.org/10.1177/10943420211003313>



Library core contains architecture-agnostic functionality

CORE

- Infrastructure Algorithms
- Iterative Solvers
 - Preconditioners
 - ...

Runtime polymorphism selects the right kernel depending on the target architecture

REFERENCE

OpenMP



Unit tests check correctness



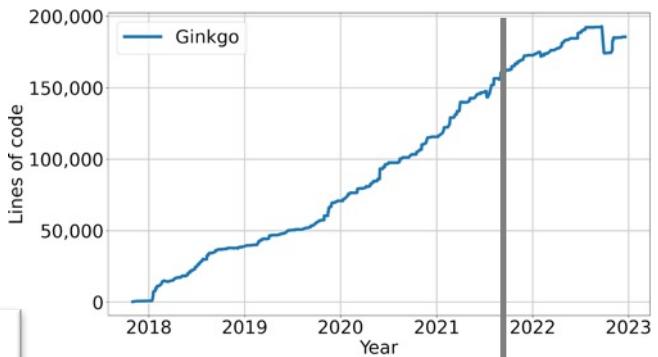
CUDA



HIP



DPC++



	OMP	CUDA	HIP	DPC++
Basic				
SpMV	✓	✓	✓	✓
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BiCG	✓	✓	✓	✓
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CGS	✓	✓	✓	✓
GMRES	✓	✓	✓	✓
IDR	✓	✓	✓	✓
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ILU/IC		✓	✓	✓
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Preconditioners				
On-Device Matrix Assembly	✓	✓	✓	✓
MC64/RCM reordering	✓			
Wrapping user data		✓		
Logging		✓		
PAPI counters		✓		

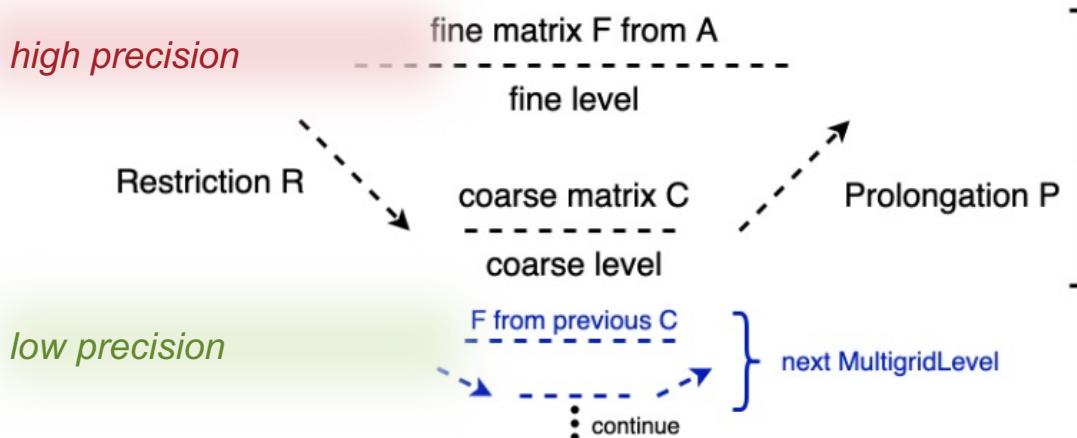
Mixed precision AMG on GPUs

- Preconditioning iterative solvers

- Idea: Approximate inverse of system matrix to make the system “easier to solve”: $P^{-1} \approx A^{-1}$

$$\text{and solve } Ax = b \Leftrightarrow P^{-1}Ax = P^{-1}b \Leftrightarrow \tilde{A}x = \tilde{b}$$

- Mixed Precision Multigrid Preconditioner



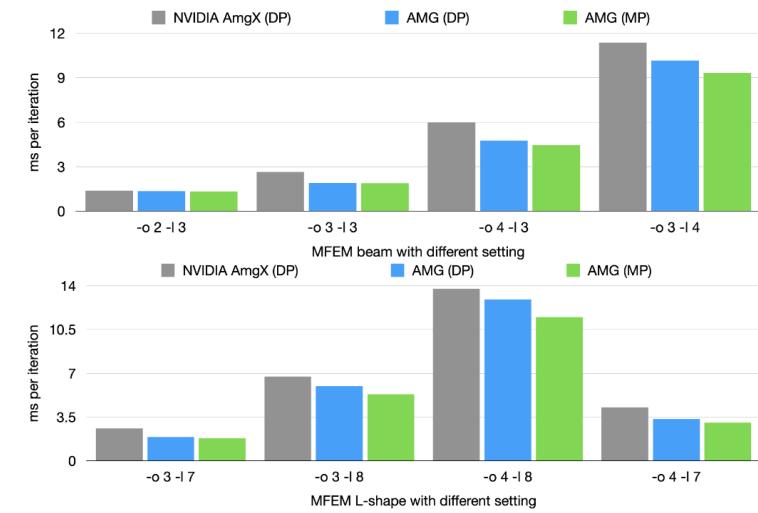
```
1 multigrid::build()
2   .with_max_levels(10u) // equal to NVIDIA/AMGX 11 max levels
3   .with_min_coarse_row(64u)
4   .with_pre_smoothen(sm, sm_f)
5   .with_mg_level(pgm, pgm_f)
6   .with_level_selector(
7     [](const size_type level, const LinOp*) -> size_type {
8       // Only the first level is generated by MultigridLevel(double).
9       // The subsequent levels are generated by MultigridLevel(float).
10      return level >= 1 ? 1 : 0;
11    })
12   .with_coarsest_solver(coarsest_solver_f)
```



MultigridLevel



Mike Tsai



Mixed precision AMG on GPUs

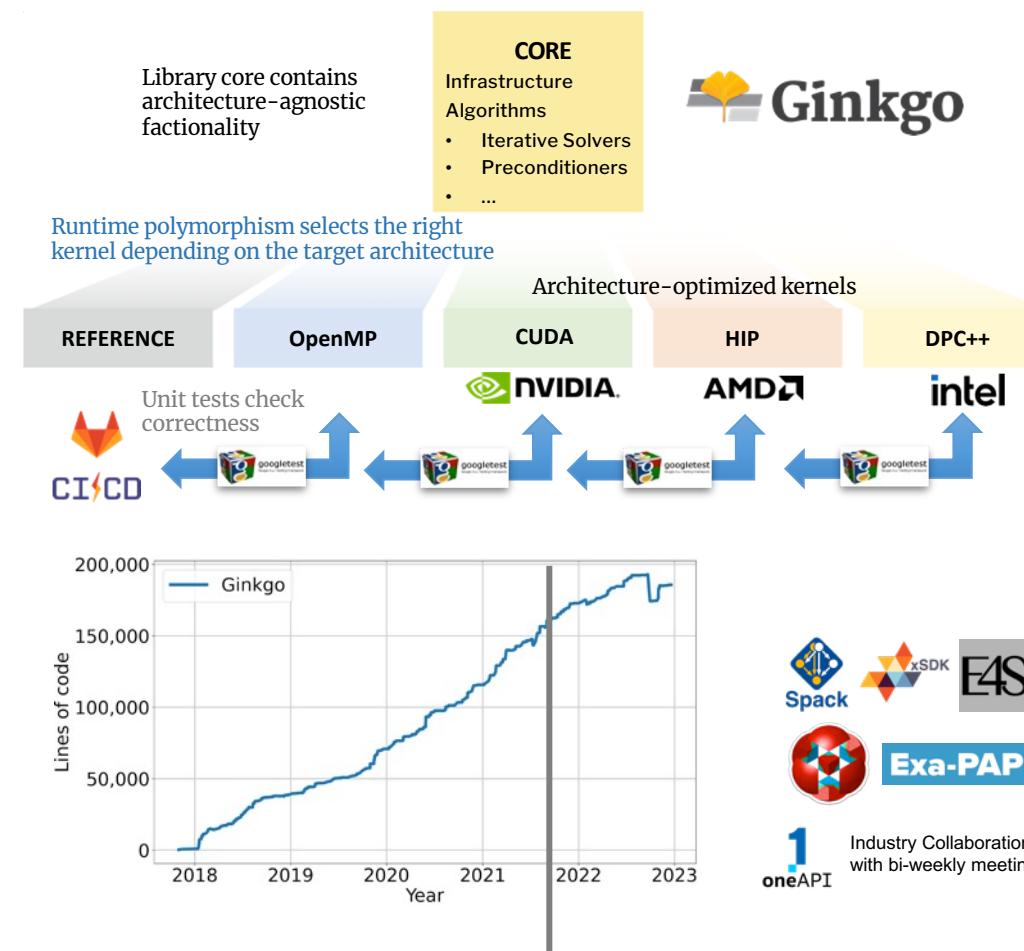


- Focus efforts
- Mixed precision
 - batched



ICL UTK @ICL_UTK · Sep 13

Congratulations to Yu-Hsiang Mike Tsai from [@KITKarlsruhe](#), in collaboration with ICL's Natalie Beams and [@HartwigAnzt](#)! Their paper "Mixed Precision Algebraic Multigrid on GPUs" took home a best paper award at PPAM2022. ppam.edu.pl



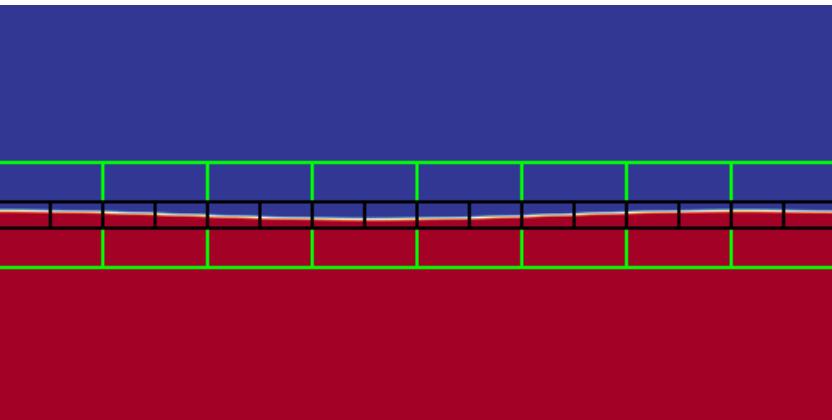
	OMP	CUDA	HIP	DPC++
Basic				
SpMV	✓	✓	✓	✓
SpMM	✓	✓	✓	✓
SpGeMM	✓	✓	✓	✓
Krylov solvers				
BiCG	✓	✓	✓	✓
BICGSTAB	✓	✓	✓	✓
CG	✓	✓	✓	✓
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GMRES	✓	✓	✓	✓
IDR	✓	✓	✓	✓
(Block-)Jacobi	✓	✓	✓	✓
ILU/IC		✓	✓	✓
Parallel ILU/IC	✓	✓	✓	✓
Parallel ILUT/ICT	✓	✓	✓	✓
Sparse Approximate Inverse	✓	✓	✓	✓
Preconditioners				
AMG preconditioner	✓	✓	✓	✓
AMG solver	✓	✓	✓	✓
Parallel Graph Match	✓	✓	✓	✓
Utilities				
On-Device Matrix Assembly	✓	✓	✓	✓
MC64/RCM reordering	✓			
Wrapping user data		✓		
Logging		✓		
PAPI counters		✓		

Batched focus effort – Combustion Simulations

Batched iterative solvers for SUNDIALS / PeleLM

PeleLM is a parallel, adaptive mesh refinement (AMR) code that solves the reacting Navier-Stokes equations in the low Mach number regime. The core libraries for managing the subcycling AMR grids and communication are found in the [AMReX source code](#).

<https://amrex-combustion.github.io/PeleLM/overview.html>

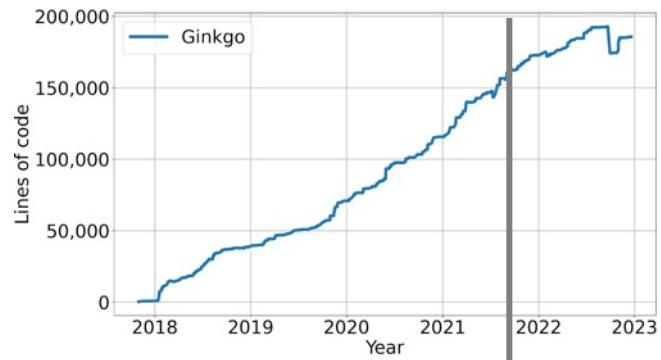
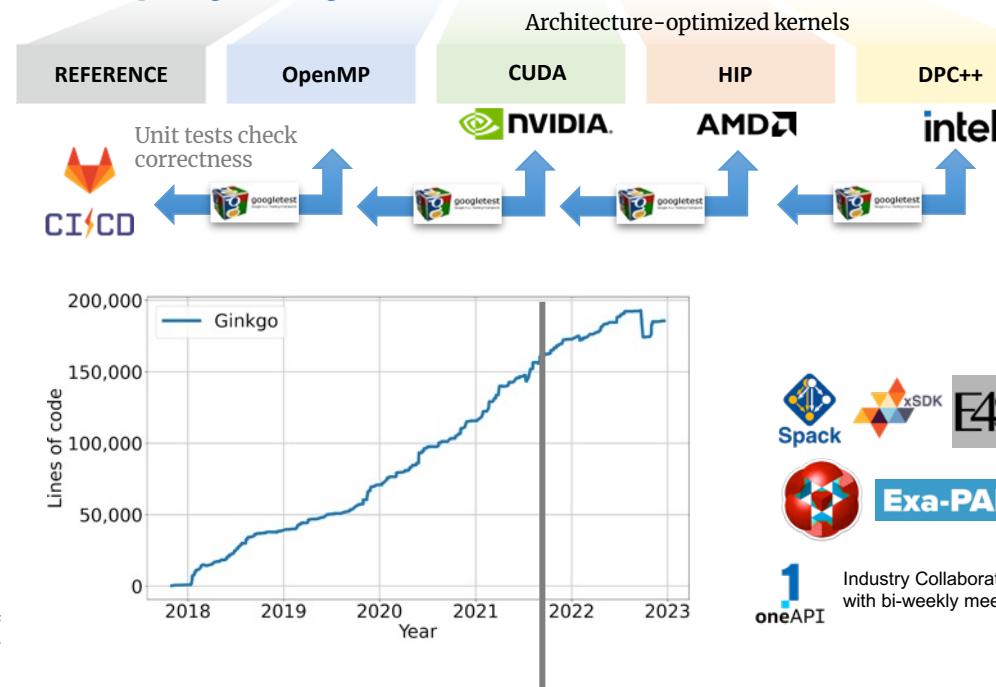


Problem	Size	Non-zeros (A)	Non-zeros (L+U)
dodecane_lu	54	2,332 (80%)	2,754 (94%)
drm19	22	438 (90%)	442 (91%)
gri12	33	978 (90%)	1,018 (93%)
gri30	54	2,560 (88%)	2,860 (98%)
isoctane	144	6,135 (30%)	20,307 (98%)
lidryer	10	91 (91%)	91 (91%)

Library core contains architecture-agnostic functionality

CORE
Infrastructure Algorithms
• Iterative Solvers
• Preconditioners
• ...

Runtime polymorphism selects the right kernel depending on the target architecture

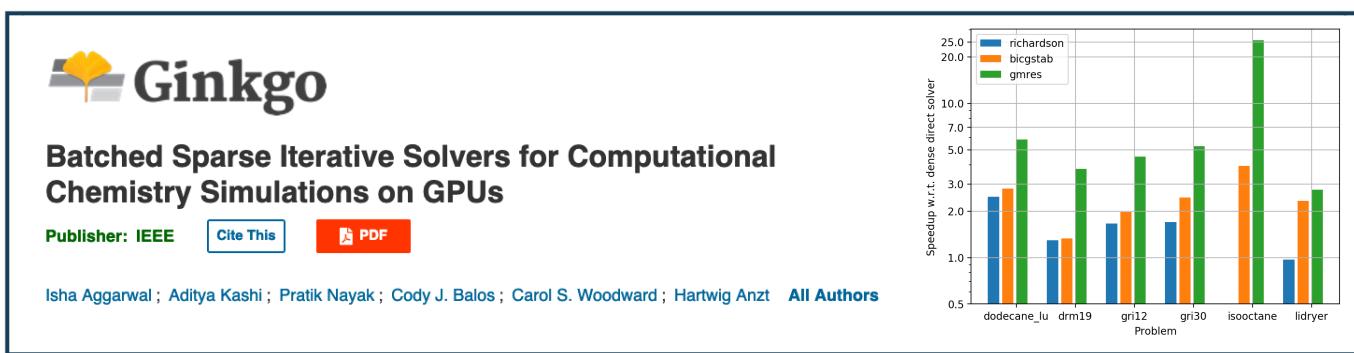
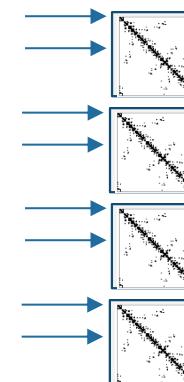
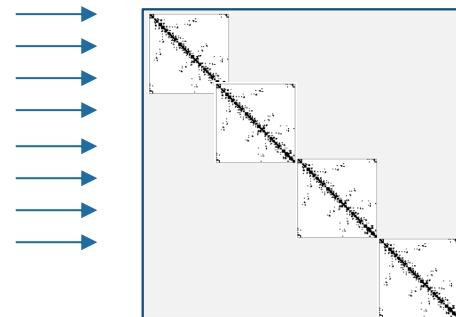
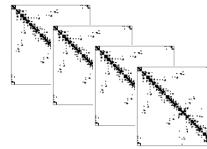


	OMP	CUDA	HIP	DPC++
Basic				
SpMV	✓	✓	✓	✓
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SpGeMM	✓	✓	✓	✓
Krylov solvers				
BiCG	✓	✓	✓	✓
BICGSTAB	✓	✓	✓	✓
CG	✓	✓	✓	✓
CGS	✓	✓	✓	✓
GMRES	✓	✓	✓	✓
IDR	✓	✓	✓	✓
(Block-)Jacobi	✓	✓	✓	✓
ILU/IC		✓	✓	✓
Parallel ILU/IC	✓	✓	✓	✓
Parallel ILUT/ICT	✓	✓	✓	✓
Sparse Approximate Inverse	✓	✓	✓	✓
Preconditioners				
AMG preconditioner	✓	✓	✓	✓
AMG solver	✓	✓	✓	✓
Parallel Graph Match	✓	✓	✓	✓
Utilities				
On-Device Matrix Assembly	✓	✓	✓	✓
MC64/RCM reordering	✓			
Wrapping user data		✓		
Logging		✓		
PAPI counters		✓		



Batched focus effort – Combustion Simulations

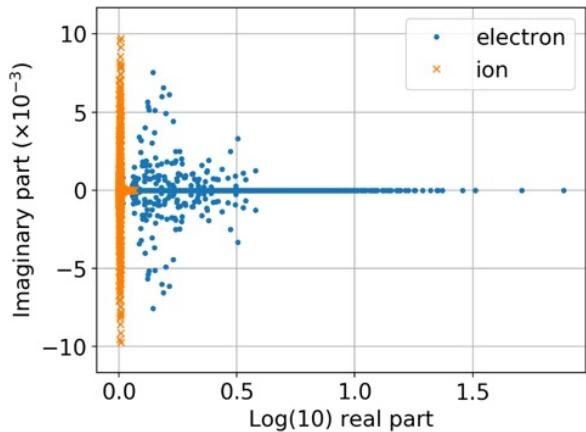
- Many sparse problems of medium size have to be solved concurrently.
 - ~ 50 – 2,000 unknowns, < 50% dense;
 - All sparse systems may share the same sparsity pattern;
 - An approximate solution may be acceptable (e.g., inside a non-linear solver);
- One solution is to arrange the individual systems on the main diagonal of one large system.
 - Convergence determined by the “hardest” problem;
 - No reuse of sparsity pattern information;
 - Global synchronization points;
- Better approach: design batched iterative solve functionality that solves all problems concurrently.
 - Problem-dependent convergence accounted for;
 - No global synchronization;
 - Reuse of sparsity pattern information;



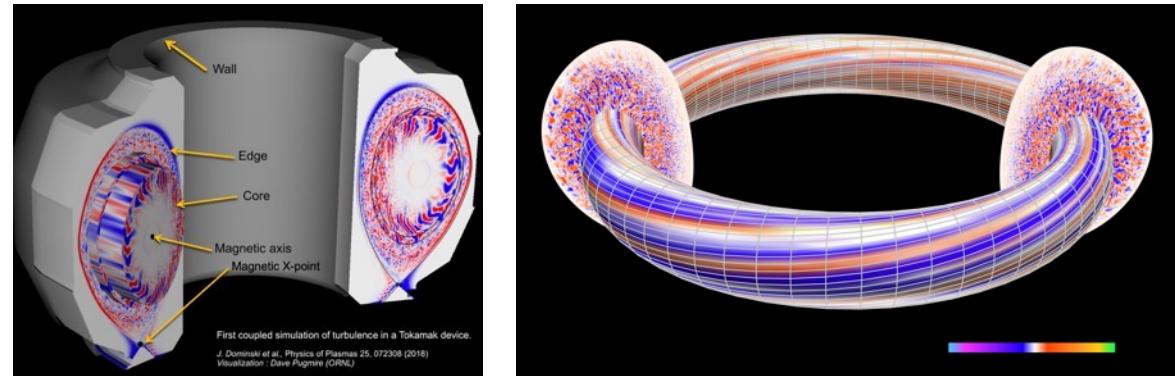
Batched focus effort – Fusion Plasma Simulations

XGC is a gyrokinetic particle-in-cell code, which specializes in the simulation of the edge region of magnetically confined thermonuclear fusion plasma. The simulation domain can include the magnetic separatrix, magnetic axis and the biased material wall. XGC can run in total-delta-f, and conventional delta-f mode. The ion species are always gyrokinetic except for ETG simulation. Electrons can be adiabatic, massless fluid, driftkinetic, or gyrokinetic.

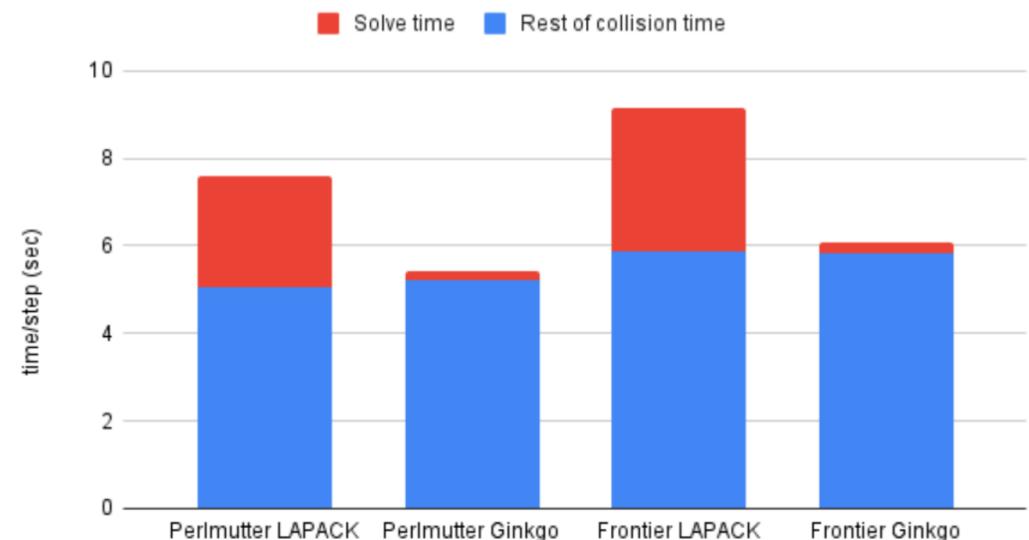
Source: https://xgc.pppl.gov/html/general_info.html



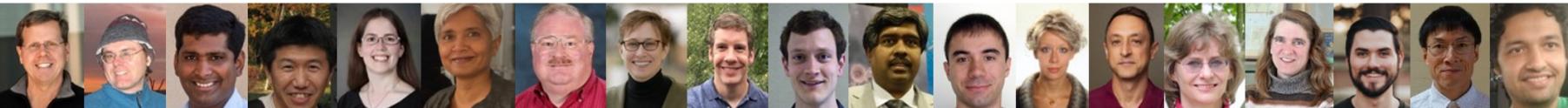
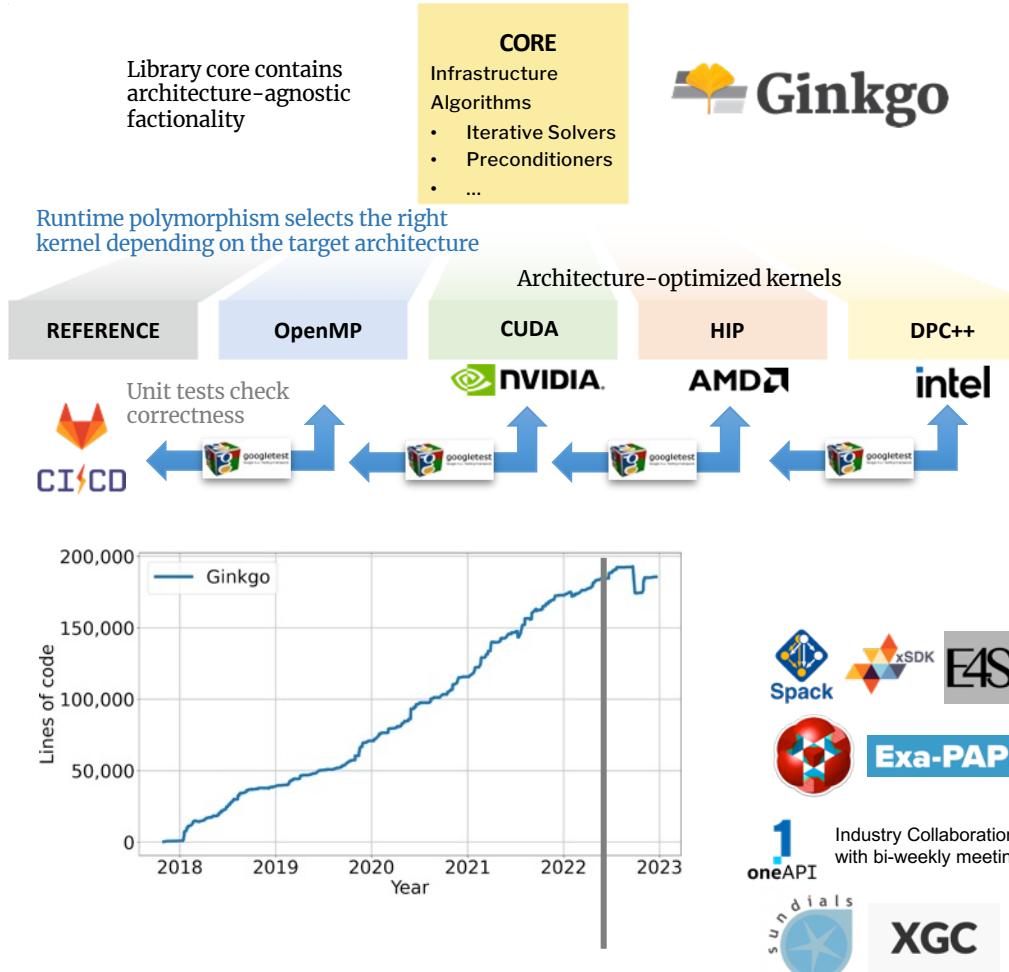
- Two species
- Ions easy to solve
- Electrons hard to solve
- Banded matrix structure
- Non-symmetric, need BiCGSTAB
- n = ~1,000
- nz = ~9,000



XGC collision time reduction (64 nodes)



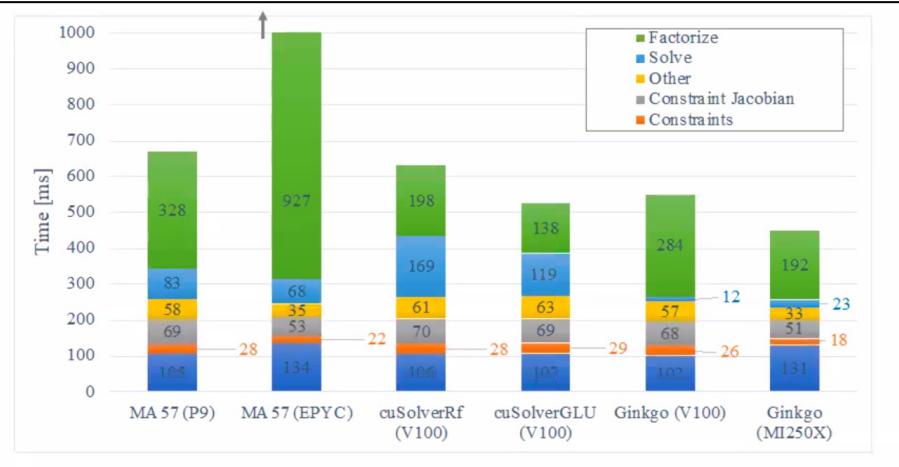
Adding Batched Functionality



Sparse direct solvers for power grid simulations

EXASGD

© Slaven Peles



- Power Grid Simulations
- All GPU solvers outperform CPU solvers
- Ginkgo first GPU-resident solver

Library core contains architecture-agnostic functionality

CORE
Infrastructure Algorithms

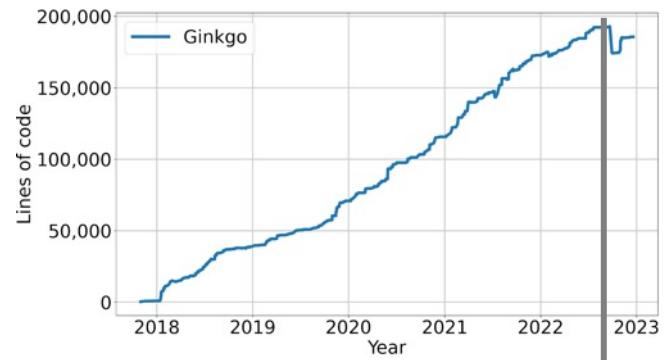
- Iterative Solvers
- Preconditioners
- ...

Runtime polymorphism selects the right kernel depending on the target architecture

 Ginkgo

REFERENCE OpenMP CUDA HIP DPC++

Unit tests check correctness
 CI/CD



EXASGD

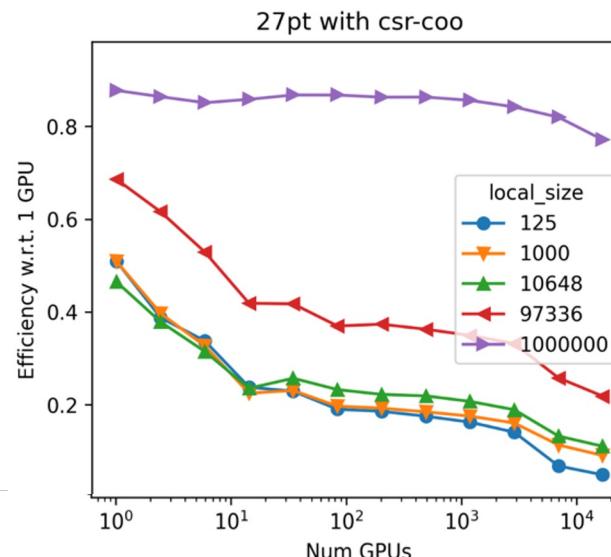
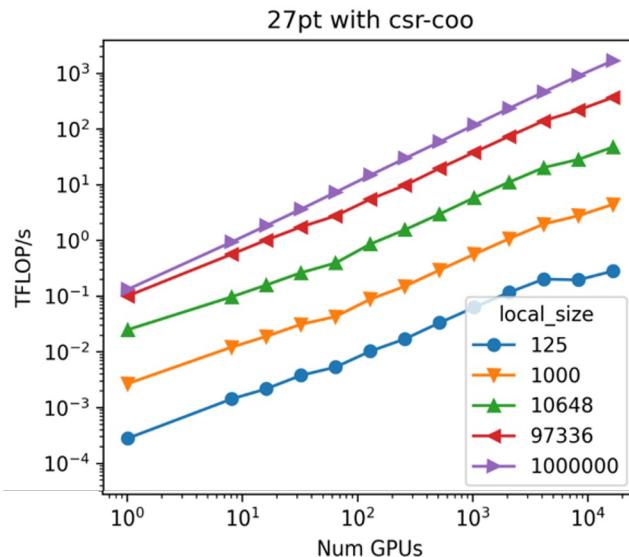
	OMP	CUDA	HIP	DPC++
Basic				
SpMV	✓	✓	✓	✓
SpMM	✓	✓	✓	✓
SpGeMM	✓	✓	✓	✓
Krylov solvers				
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BICGSTAB	✓	✓	✓	✓
CG	✓	✓	✓	✓
CGS	✓	✓	✓	✓
GMRES	✓	✓	✓	✓
IDR	✓	✓	✓	✓
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ILU/IC		✓	✓	✓
Parallel ILU/IC	✓	✓	✓	✓
Parallel ILUT/ICT	✓	✓	✓	✓
Preconditioners				
Sparse Approximate Inverse	✓	✓	✓	✓
Batched BiCGSTAB	✓	✓	✓	
Batched CG	✓	✓	✓	
Batched GMRES	✓	✓	✓	
Batched ILU	✓	✓	✓	
Batched ISAI	✓	✓	✓	
Batched Jacobi	✓	✓	✓	
AMG				
AMG preconditioner	✓	✓	✓	✓
AMG solver	✓	✓	✓	✓
Parallel Graph Match	✓	✓	✓	✓
Sparse direct				
Symbolic Cholesky	✓	✓	✓	✓
Numeric Cholesky				UNDER DEVELOPMENT
Symbolic LU	✓	✓	✓	✓
Numeric LU	✓	✓	✓	
Sparse TRSV	✓	✓	✓	
Utilities				
On-Device Matrix Assembly	✓	✓	✓	✓
MC64/RCM reordering	✓			
Wrapping user data		✓		
Logging		✓		
PAPI counters		✓		



Distributed runs on Frontier (Cray + AMD MI250 GPUs)

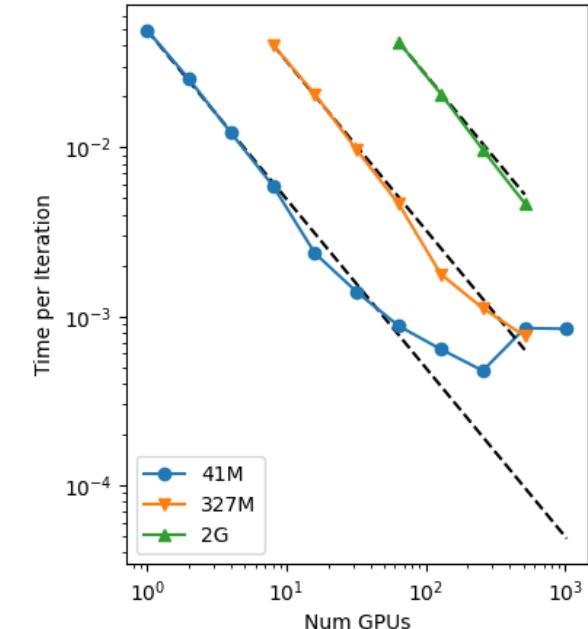
Weak scaling: problem size increases with parallel resources

Weak scaling up to 16k GCDs (8k GPUs)

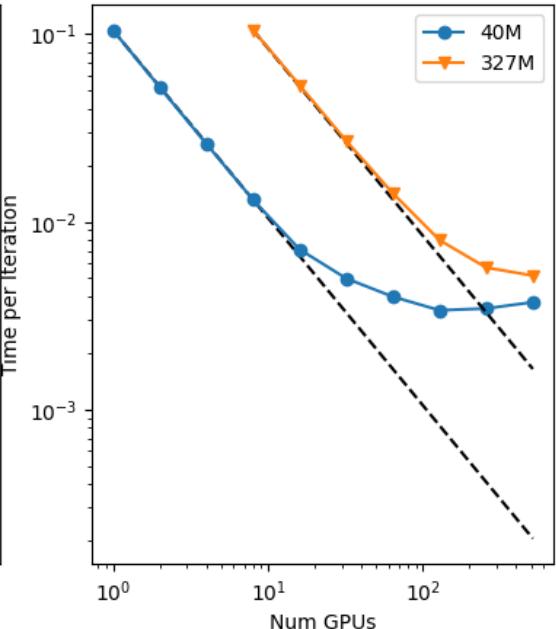


Strong scaling: problem size constant

Strong Scaling - Cg Local Jacobi

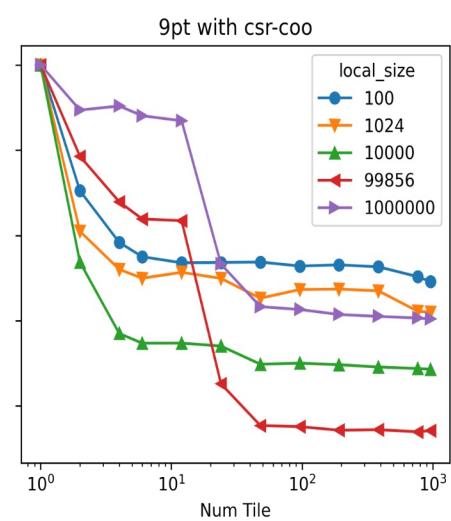
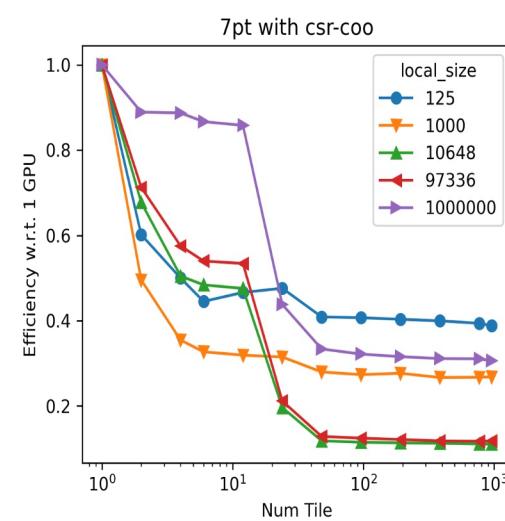
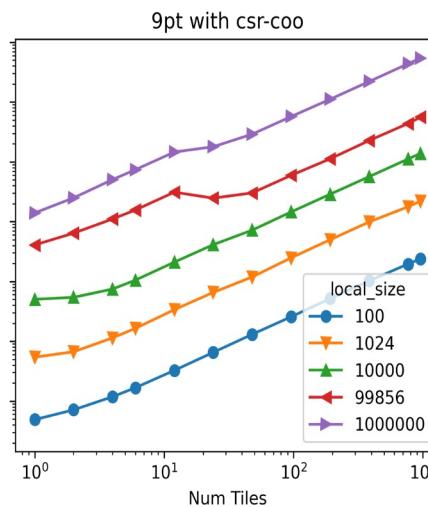
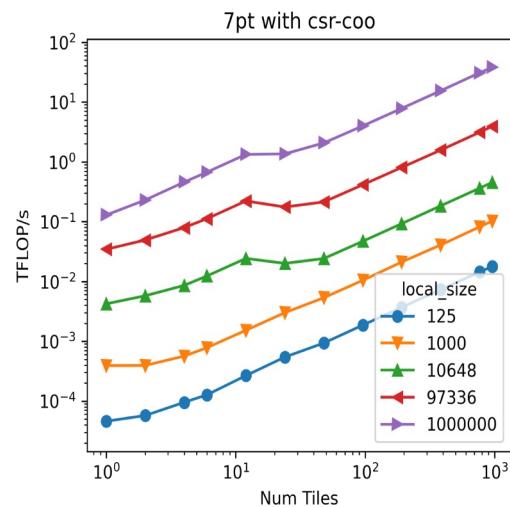
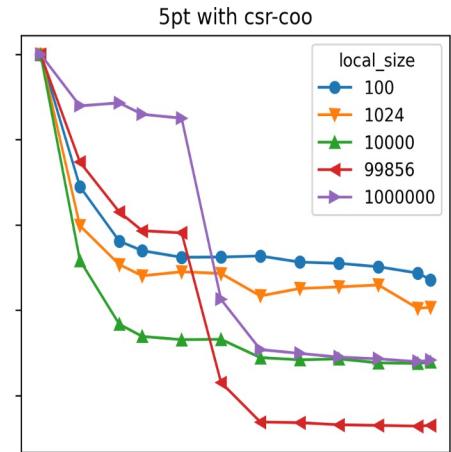
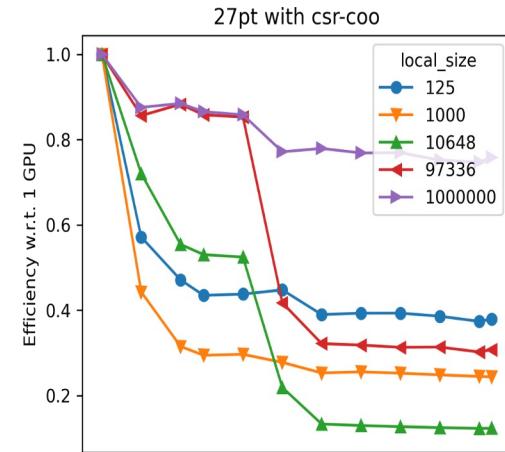
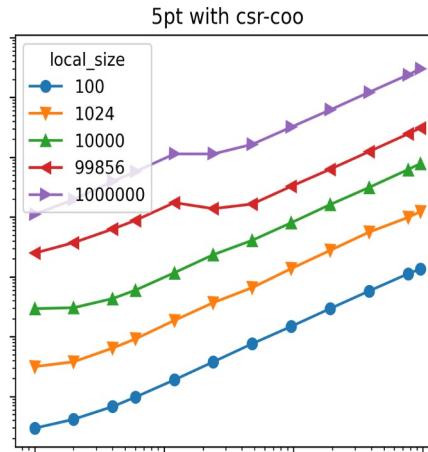
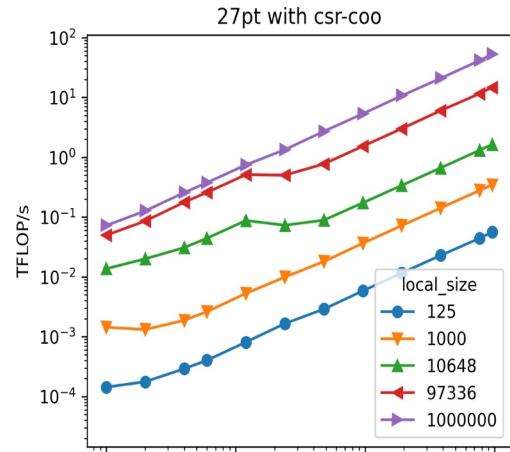


Strong Scaling - Gmres Local Jacobi



Distributed runs on Sunspot (Intel PVCA GPUs)

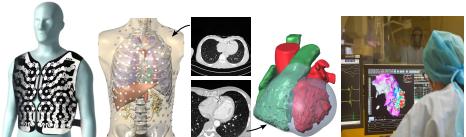
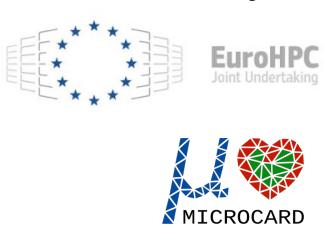
Weak scaling: problem size increases with parallel resources



“Now” – Near completion of ECP

- Sustainable software design ready for the addition of new backends.

- EuroHPC Project MICROCARD uses Ginkgo



<https://www.microcard.eu>

- BMBF PDExa project uses Ginkgo

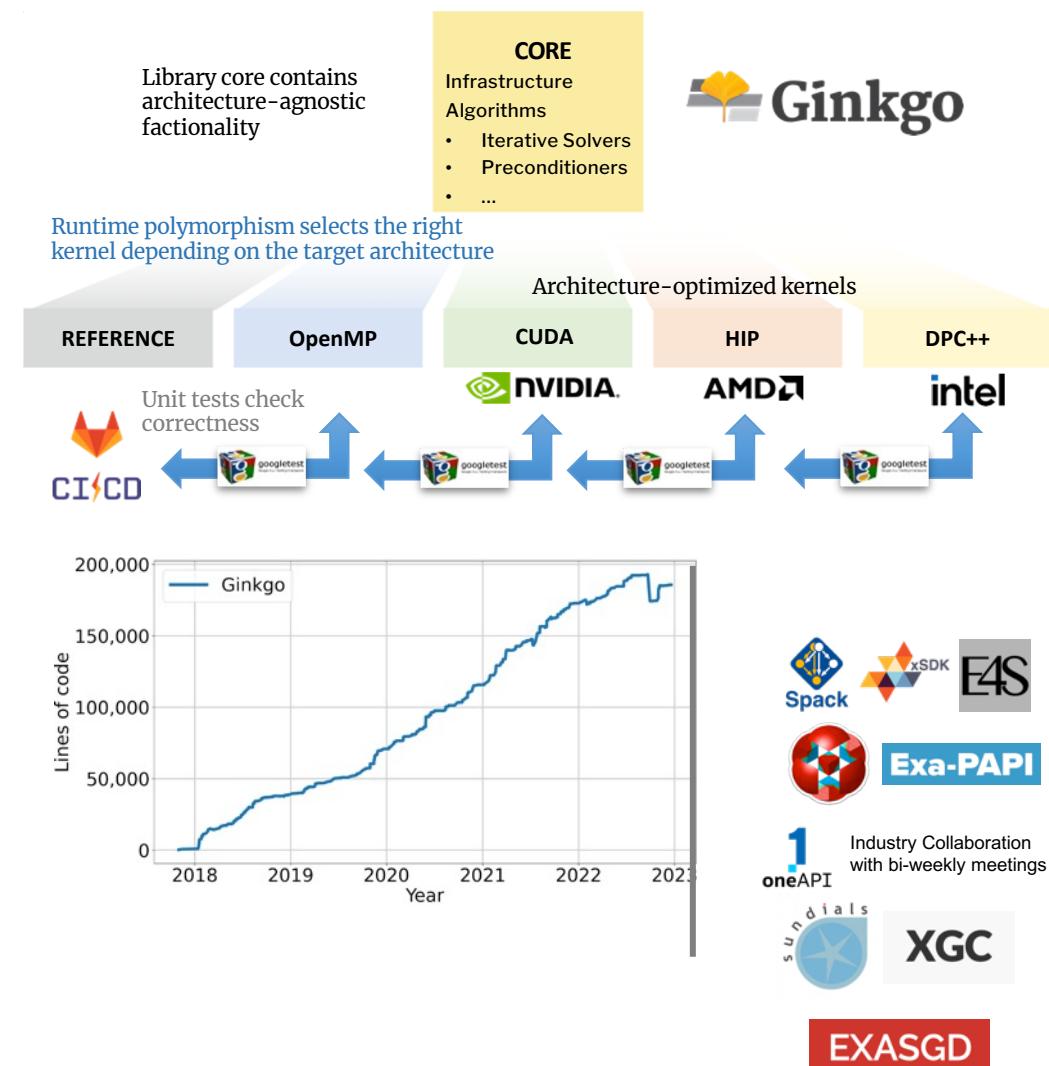


- BMBF ExaSIM project uses Ginkgo



The Open Source CFD Toolbox

<https://exasim-project.com>



Lessons learnt from the Ginkgo development process

- ECP earmarking roughly half the budget to Software & App development is a game changer.
 - Central component for the success of ECP.
 - This concept needs to – and does become - the blueprint for other nations and projects.
- Workforce recruitment and workforce retention are the key to success in software development.
 - Money does not write software. RSEs do. We need to create attractive career plans.
 - We need to make research software development attractive to students. Academic recognition.
- Anticipating the future in hardware development accelerates the porting process.
 - Blueprints and early access systems both useful.
 - Interaction with industry is mutually beneficial.
- Management, tools, and strategic initiatives, interaction and collegial behavior are important.
 - Jira/Notion/[...] milestones and deliverables give projects and collaborative interactions a structure and timeline.
 - Strategic focus groups, conferences, and meetings bring experts together and create collaboration.
 - Listen to the application needs. Value input and acknowledge collaborators.