

Everything Everywhere All at Once - How Improving Software Sustainability, Productivity, and Performance Naturally Go Hand in Hand

A story from the trenches of the US Exascale Computing Project

Prof. Dr. Hartwig Anzt

Technical University of Munich

05/23/2024



The US Exascale Computing Project



Advancing Scientific Discovery

The ECP aims to ensure availability of the exascale computing ecosystem necessary for developing clean energy systems, improving the resilience of our infrastructure, designing new materials that can perform in extreme environments, adapting to changes in the water cycle, developing smaller and more powerful accelerators for use in medicine and industry, and much more. Several projects focus on data-intensive problems to enable effective use of the data streams from powerful scientific facilities, complex environmental genomes, and cancer research (patient genetics, tumor genomes, molecular simulations, and clinical data).



Strengthening National Security

The ECP teams are also developing new applications for supporting the NNSA Stockpile Stewardship Program, which is responsible for maintaining the readiness and reliability of our nuclear weapons systems—without underground testing. Assessing the performance of weapons systems subject to hostile environments and potential threat scenarios exceeds the capabilities of current HPC systems and codes. NNSA application projects are focused on providing the sophisticated modeling and analysis tools needed to sustain the U.S. nuclear deterrence.



Improving Industrial Competitiveness

Exascale systems will be used to accelerate research that leads to innovative products and speeds commercialization, creating jobs and driving US competitiveness across industrial sectors, such as the emerging energy economy. To ensure alignment with US industry needs, the ECP is engaging senior technology decision makers from among the country's most prominent private sector companies.

The US Exascale Computing Project

Addressing a National Imperative

The Exascale Computing Project is an aggressive research, development, and deployment project focused on delivery of mission-critical applications, an integrated software stack, and exascale hardware technology advances.

Application Development



Software Technology

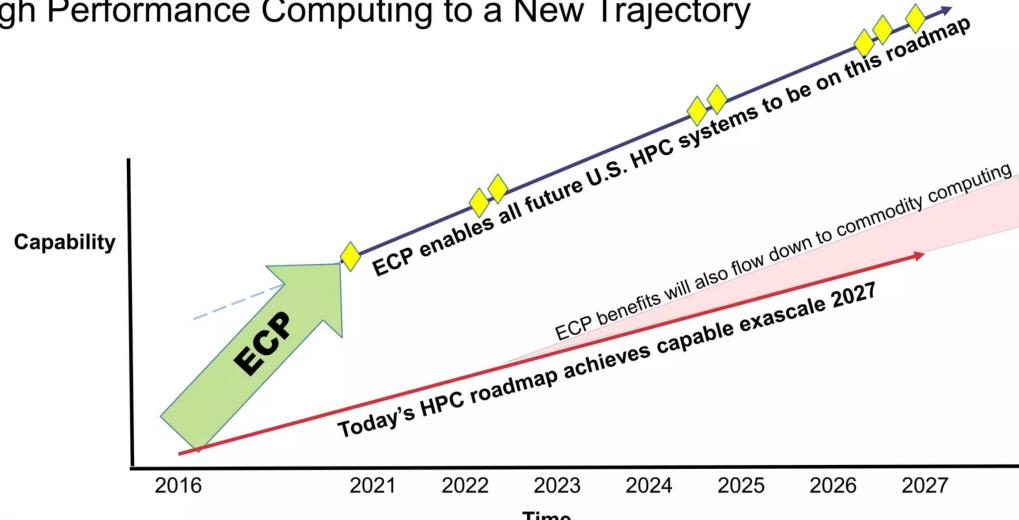


Hardware & Integration



© Paul Messina

Vision: Exascale Computing Project (ECP) Lifts all U.S. High Performance Computing to a New Trajectory



The US Exascale Computing Project



US\$4B – what is it spent on?

- 3 computers
 - \$600M each
 - \$400M to vendors for Design, Path, Fast - Forward
- 21 Applications



AMD Based
(Up & running)
20 MW



Intel Based
(Up & running)
40 MW



AMD APU Based
(panned)

The US Exascale Computing Project

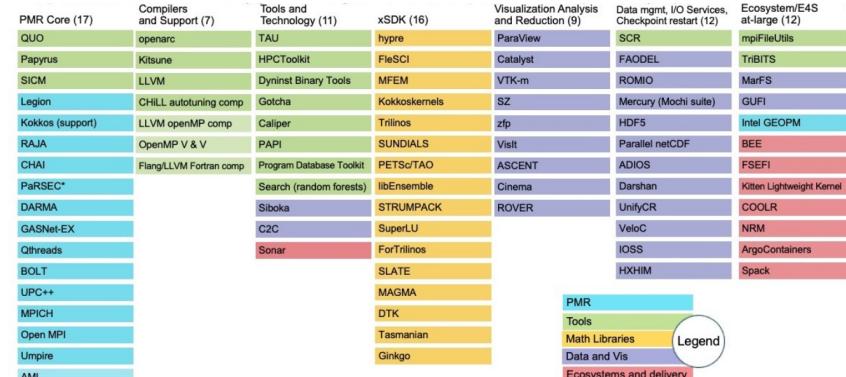
US\$4B – what is it spent on?

- 3 computers
 - \$600M each
 - \$400M to vendors for Design, Path, Fast - Forward
- 21 Applications



Domain*	Base Challenge Problem	Domain*	Challenge Problem
Wind Energy	2x2 5 MW turbine array in 3x3x1 km ³ domain	Quantum Materials	Predict & control mats @ quantum level
Nuclear Energy	Small Modular Reactor with complete in-vessel coolant loop	Astrophysics	Supernovae explosions, neutron star mergers
Fossil Energy	Burn fossil fuels cleanly with CLRs	Cosmology	Extract "dark sector" physics from upcoming cosmological surveys
Combustion	Reactivity controlled compression ignition	Earthquakes	Regional hazard and risk assessment
Accelerator Design	TeV-class 10 ²⁻³ times cheaper & smaller	Geoscience	Well-scale fracture propagation in wellbore cement due to attack of CO ₂ -saturated fluid
Magnetic Fusion	Coupled gyrokinetics for ITER in H-mode	Earth System	Assess regional impacts of climate change on the water cycle @ 5 SYPD
Nuclear Physics: QCD	Use correct light quark masses for first principles light nuclei properties	Power Grid	Large-scale planning under uncertainty; underfrequency response
Chemistry: GAMESS	Heterogeneous catalysis: MSN reactions	Cancer Research	Scalable machine learning for predictive preclinical models and targeted therapy
Chemistry: NWChemEx	Catalytic conversion of biomass	Metagenomics	Discover and characterize microbial communities through genomic and proteomic analysis
Extreme Materials	Microstructure evolution in nuclear mats	FEL Light Source	Protein and molecular structure determination using streaming light source data
Additive Manufacturing	Born-qualified 3D printed metal alloys		

Sustainable software development



A few words about myself

- Born and raised in Karlsruhe
- PhD in Numerical Mathematics from KIT
- Focus on computational linear algebra and high performance computing (HPC)
- Linear solvers, preconditioners, ...
- During my PostDoc at the University of Tennessee, I developed MAGMA sparse



MAGMA SPARSE

MAGMA-sparse as a “child” of MAGMA explores the development of sparse linear algebra functionality for NVIDIA GPUs.



Limitations:

- *C code with hand-written build system*
- *Sparse unit testing*
- *Focus on NVIDIA GPUs*
- *Design-specific limitations (flexibility/extensibility)*

Designing an ECP math library



Building Trusted Scientific Software

SHARE in f t d



Software Verification

SHARE in f t d



Think Locally, Act Globally: Outreach for Better Scientific Software

SHARE in f t d

I have worked in the scientific software field for more! phrase "Verification is doing things right, and validation phrase to memory in order to avoid confusion when the

PUBLISHED JUN 28, 2018 AUTHOR MIKE HER

Pairing internal and external concerns

Verification focuses on internal concerns of a good sol

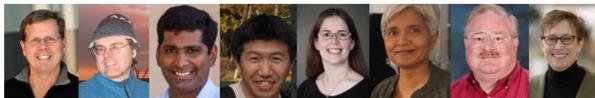
In the realm of software, verification is often erroneo proper subset of verification for gaining confidence in the holistic process by which the developers convinc it was designed to do. In scientific software this coul numerical stability, and efficacy of the method in the expected results. Note that verification is limited to e model specification, not that the model itself matche validation process.

PUBLISHED AUG 15, 2018 AUTHOR ANSHU I

An ambitious goal

The ECP needs to deliver a software environment and applications ready to run on exascale computers, which are scheduled to be deployed starting in 2021. Achieving this goal entails a major, large-scale software development effort. Recognizing the challenges development teams will face, the ECP is supporting the IDEAS Productivity project to help scientific researchers improve their development practices.

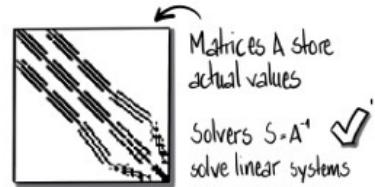
PUBLISHED JUL 17, 2018 AUTHOR DAVID BERNHOLDT TOPICS BETTER SKILLS PERSONAL PRODUCTIVITY AND SUSTAINABILITY



Designing software for performance, portability, & sustainability



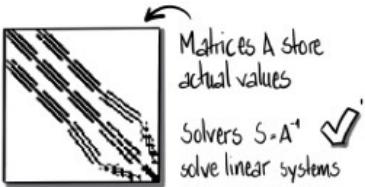
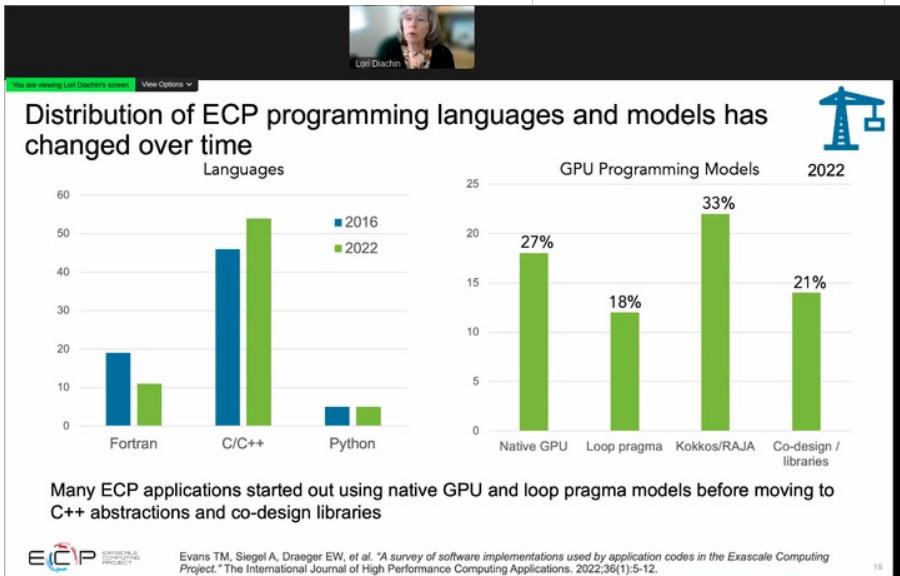
Ginkgo - A sparse linear algebra library for HPC



Designing software for performance, portability, & sustainability



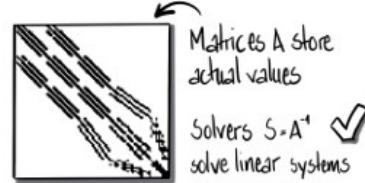
written in C++ → Ginkgo - A sparse linear algebra library for HPC



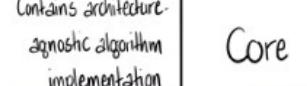
Designing software for performance, portability, & sustainability



written in C++ → Ginkgo - A sparse linear algebra library for HPC



Contains architecture-agnostic algorithm implementation

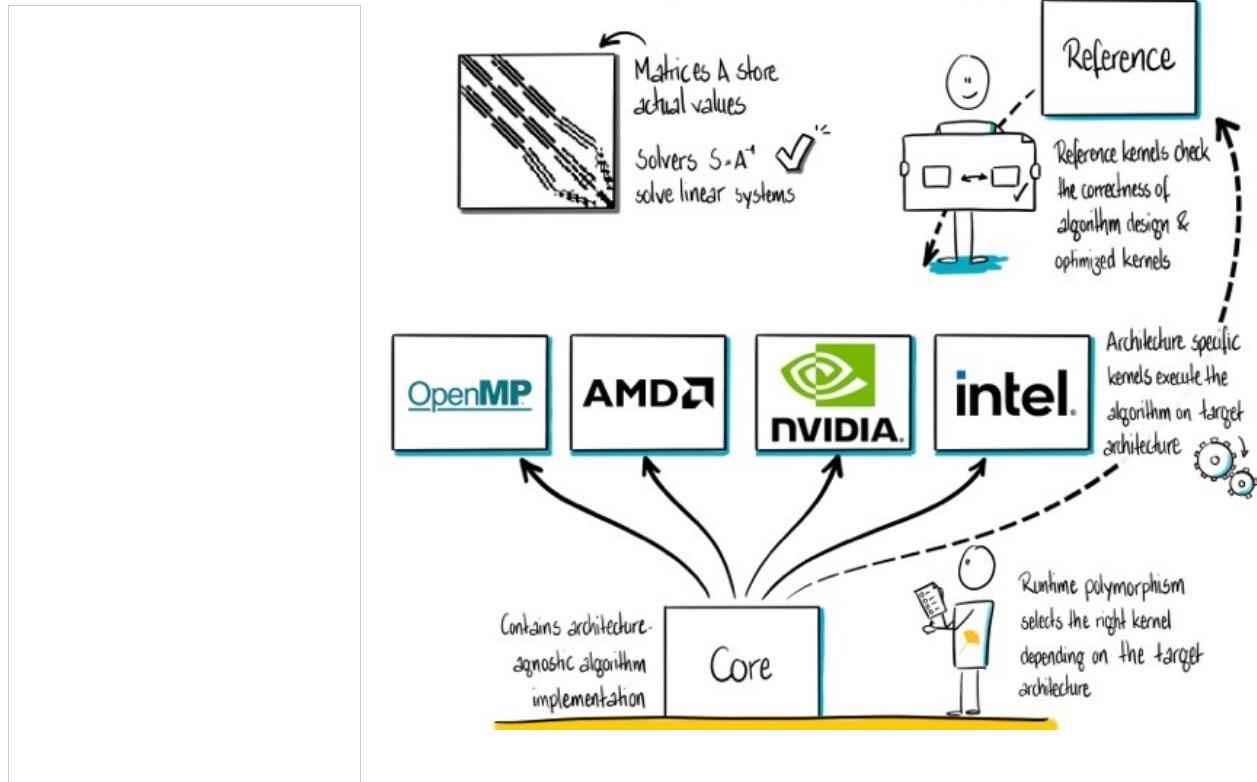


Runtime polymorphism selects the right kernel depending on the target architecture

Designing software for performance, portability, & sustainability



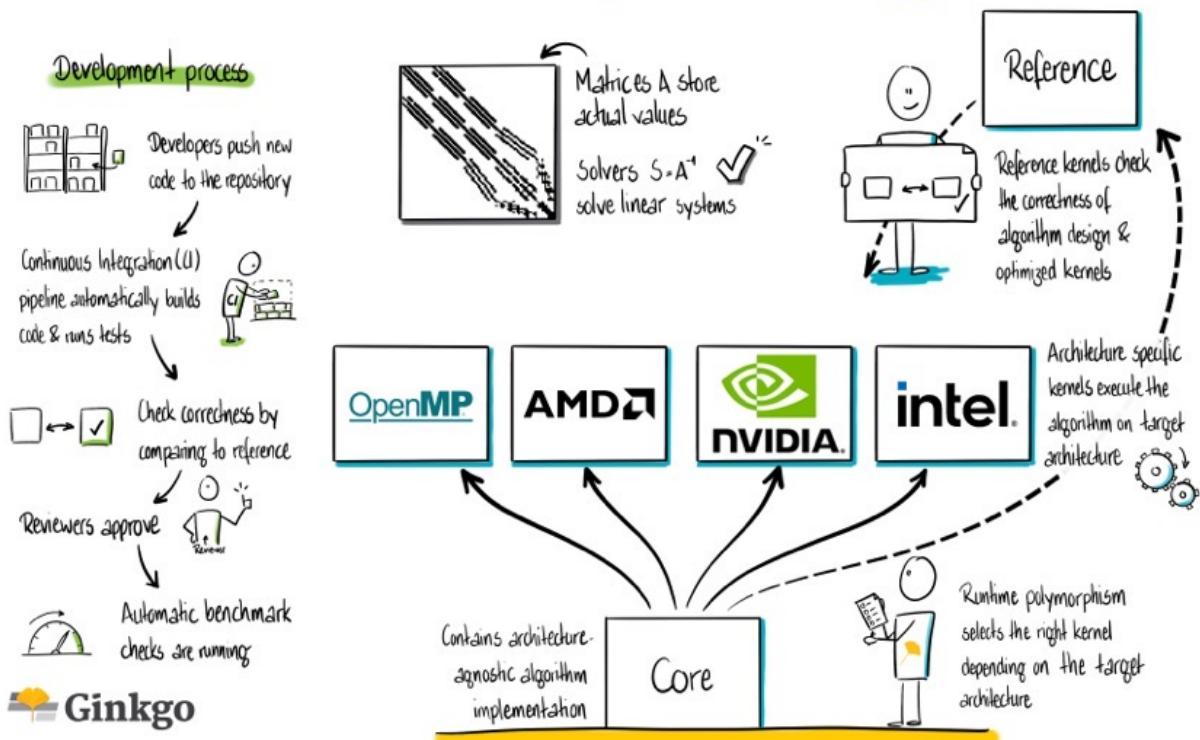
written in C++ → Ginkgo - A sparse linear algebra library for HPC



Designing software for performance, portability, & sustainability



written in C++ → Ginkgo - A sparse linear algebra library for HPC



Designing software for performance, portability, & sustainability



Ginkgo - A generic linear algebra library for HPC

Linear Operator Interface

We express everything as Linear Operator.

- Internally, we leverage C++ class inheritance.
- Applications can apply any functionality as a linear operator.

Matrix-Vector Product Preconditioner (for matrix A) Solver (for system $Ax = b$)

$x := A \cdot b$	$x := M^{-1} \cdot b$	$x := S \cdot b$
$M^{-1} \approx A^{-1}$	$S \approx A^{-1}$	
$M^{-1} = \Pi(A)$	$S = \Sigma(A)$	

All of them can be expressed as

Application of a linear operator* (LinOp) $L : \mathbb{F}^m \rightarrow \mathbb{F}^m$

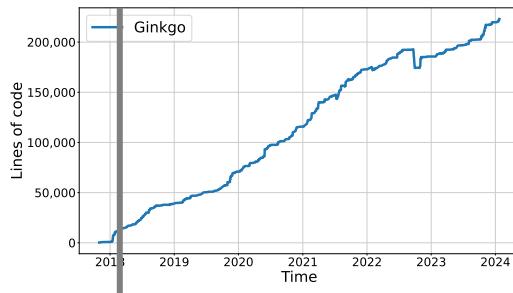


Runtime polymorphism
selects the right kernel
depending on the target
architecture

Designing software for performance, portability, & sustainability



written in C++ → Ginkgo - A sparse linear algebra library for HPC



Development process



Developers push new code to the repository

Continuous Integration (CI) pipeline automatically builds code & runs tests



Check correctness by comparing to reference

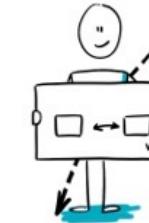
Reviewers approve

Automatic benchmark checks are running



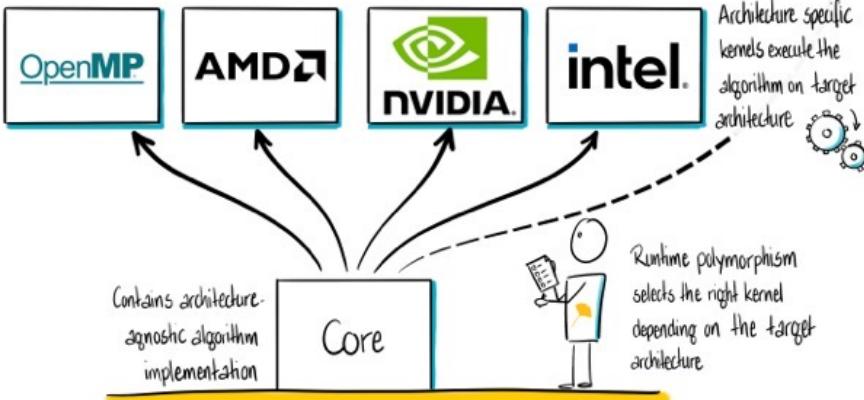
Matrices A store actual values

Solvers S=A⁻¹ solve linear systems

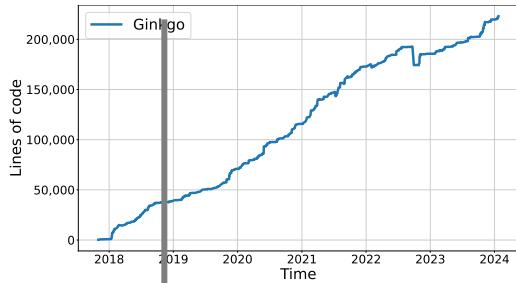
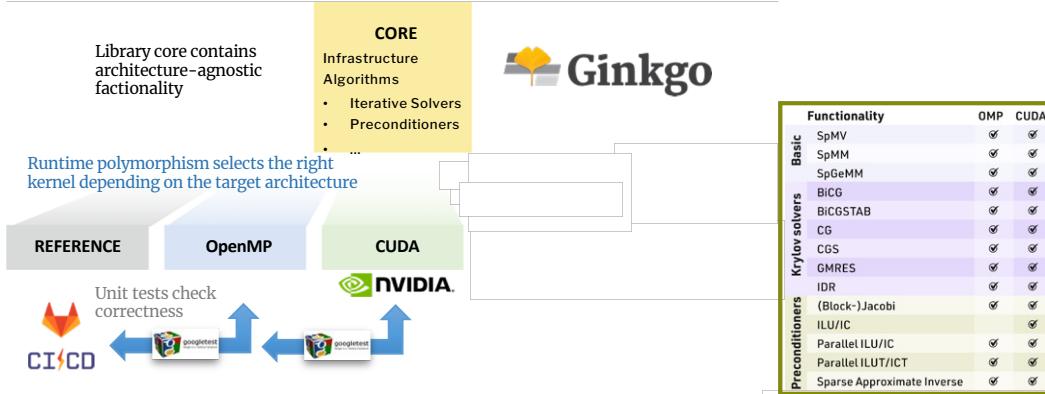


Reference

Reference kernels check the correctness of algorithm design & optimized kernels



Starting with the CUDA backend



	Functionality	OMP	CUDA
Basic	SpMV	✓	✓
	SpMM	✓	✓
	SpGeMM	✓	✓
	BICG	✓	✓
	BICGSTAB	✓	✓
	CG	✓	✓
	CGS	✓	✓
	GMRES	✓	✓
	IDR	✓	✓
	(Block-)Jacobi	✓	✓
Krylov solvers	ILU/IC	✓	✓
	Parallel ILU/IC	✓	✓
	Parallel ILUT/ICT	✓	✓
	Sparse Approximate Inverse	✓	✓

Extending to AMD GPUs

~2 months



better scientific software

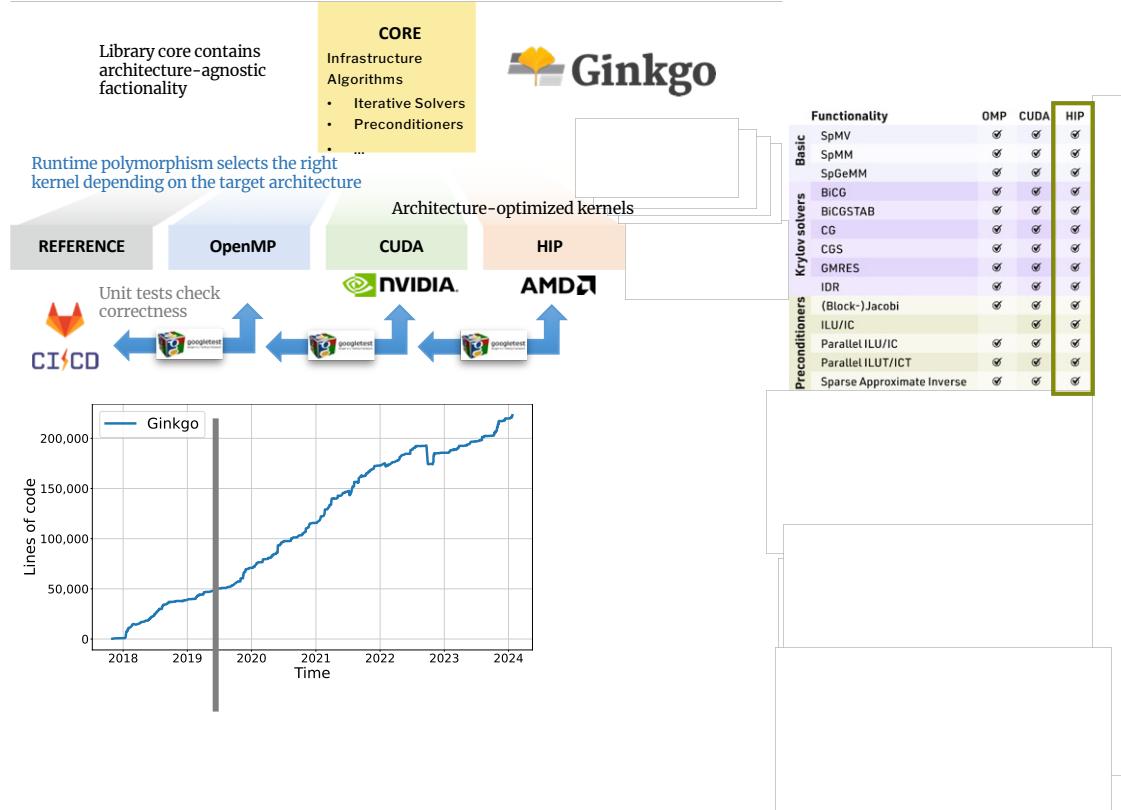
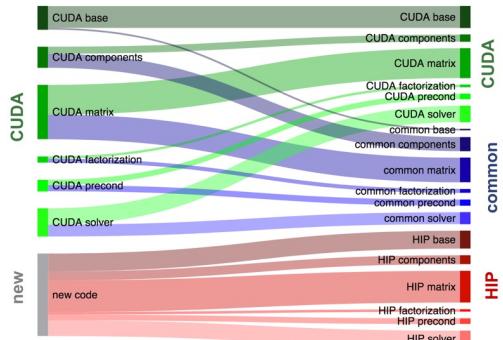
Resources ▾ Blog Events About ▾

HOME > BLOG > Porting the Ginkgo Package to AMD's HIP Ecosystem

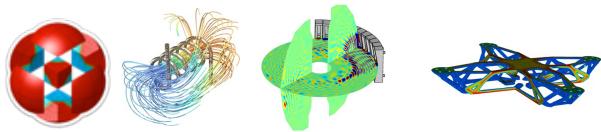
SHARE in f t p

In response to the explosion-like diversification in hardware architectures, hardware portability and the ability to adopt new processor designs have become a central priority in realizing software sustainability. In this blog article, we discuss the experience of porting CUDA code to AMD's Heterogeneous-compute Interface for Portability (HIP).

PUBLISHED JUN 25, 2020 AUTHOR HARTWIG ANZT TOPICS BETTER RELIABILITY TESTING BETTER PLANNING DESIGN



Input from the “first customer”



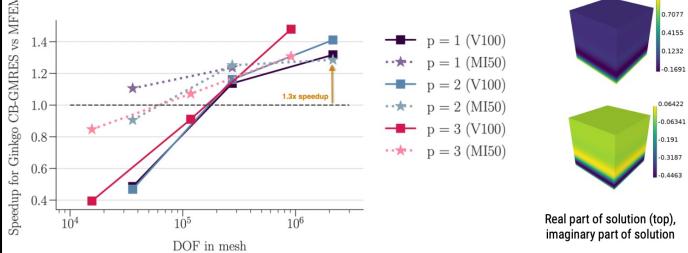
MFEM is a *free, lightweight, scalable C++ library* for finite element methods.

Speeding up MFEM’s “example 22” on GPUs

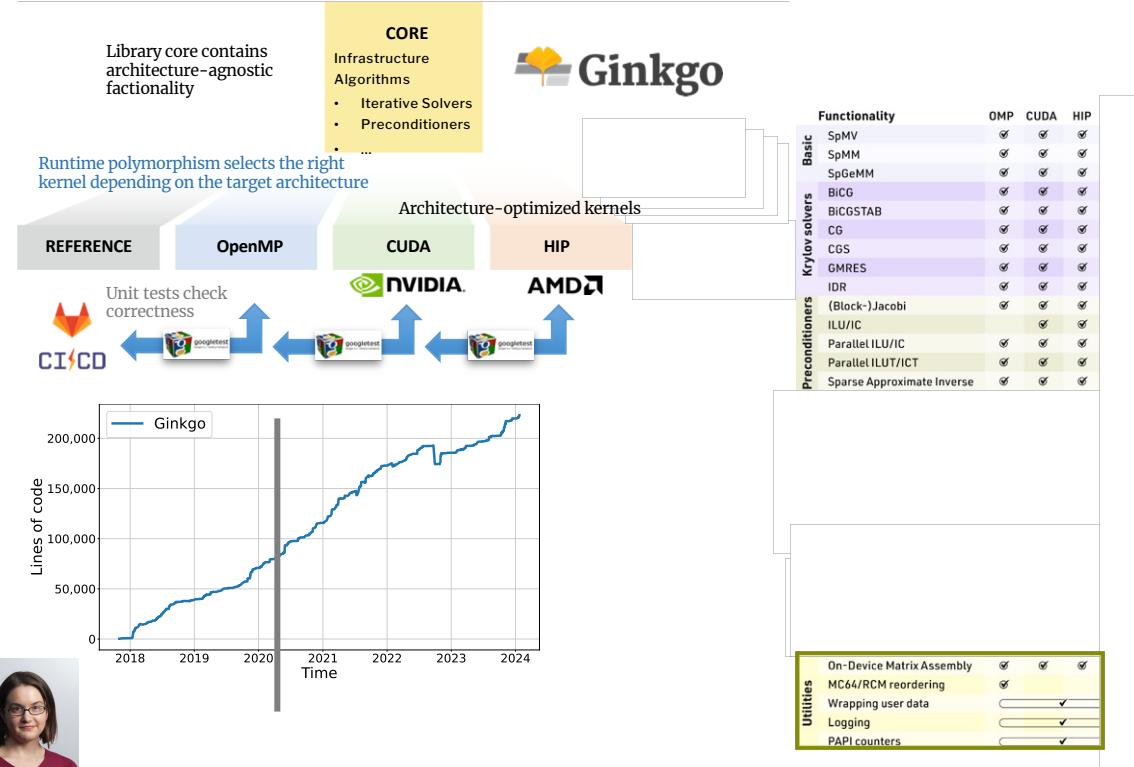
Example 22 of the MFEM finite element library solves harmonic oscillation problems, with a forced oscillation imposed at the boundary. In this test, we use variant 1:

$$-\nabla \cdot (a \nabla u) - \omega^2 bu + i\omega cu = 0$$

with $a = 1$, $b = 1$, $\omega = 10$, $c = 20$

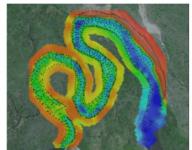


Speedup of Ginkgo's Compressed Basis-GMRES solver vs MFEM's GMRES solver for three different orders of basis functions (p), using MFEM matrix-free operators and the Ginkgo-MFEM integration wrappers in MFEM. CUDA 10.1/V100 and ROCm 4.0/MI50.

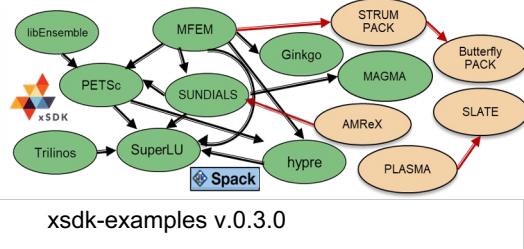


Part of the xSDK effort

xSDK: Extreme-scale Scientific Software Development Kit



Integrated surface-subsurface hydrology simulations of river meanders require the combined use of xSDK packages.



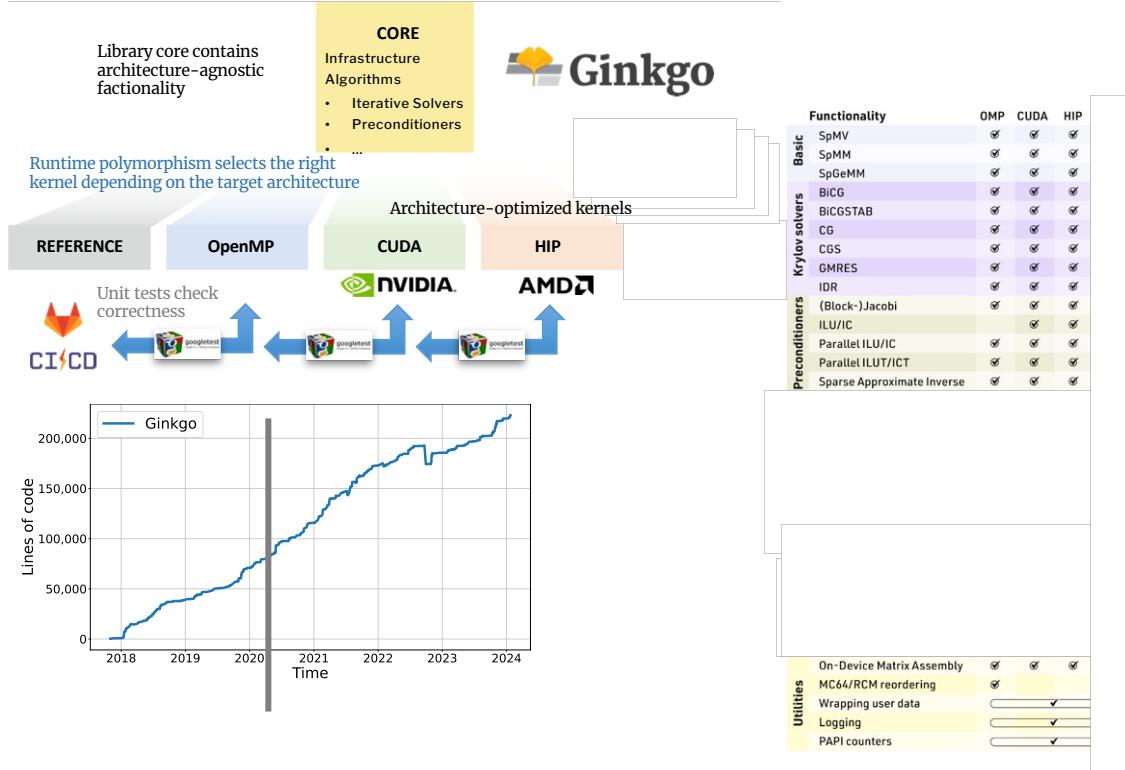
The xSDK provides infrastructure for and interoperability of a **collection of related and complementary software elements**—developed by diverse, independent teams throughout the high-performance computing (HPC) community—that provide the building blocks, tools, models, processes, and related artifacts for rapid and efficient development of high-quality applications.

November 2022

- 26 math libraries
- 2 domain components
- 16 mandatory xSDK community policies
- Available on Github
- Spack xSDK installer

xSDK community policies:

- 16 mandatory policies,
- 8 recommended policies,
- 4 Spack variant guidelines
- Available on Github
<https://x sdk.info/policies/>



Extending to Intel GPUs

~12 months



Since 1987 - Covering the Fastest Computers in the World and the People Who Run Them

- Home
- Technologies
- Sectors
- COVID-19
- AI/ML/DL

Preparing for the Arrival of Intel's Discrete High-Performance GPUs
By Hartwig Anzt

March 23, 2021

yhmtsai / try_oneapi · Private

Code Issues Pull requests Actions Projects Security Insights

1 master · 1 branch · 0 tags

Go to file Add file · Code · About

No description, website, or topics provided.

arg_struct WIP 2 years ago

atomic atomic and get_in_template 2 years ago

check_unit some checker last year

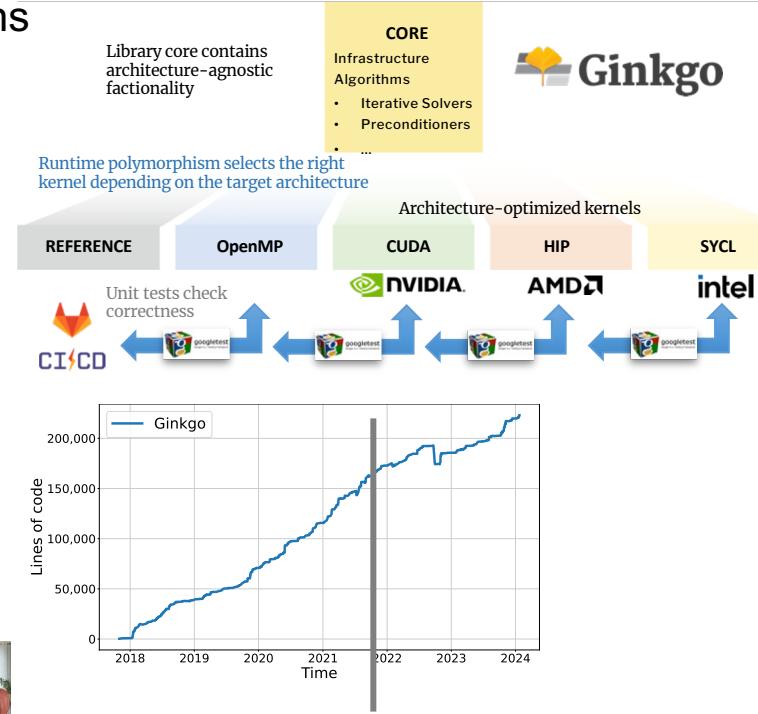
classical_csr cpu barrier issue in classical csr spmv last year

clinfo clinfo 2 years ago

coop_cuda keep some history but I do not check them detail last year

coop_draft keep some history but I do not check them detail last year

Releases No releases published Create a new release



Functionality	OMP	CUDA	HIP	DPC++
SpMV	✓	✓	✓	✓
SpMM	✓	✓	✓	✓
SpGeMM	✓	✓	✓	✓
BICG	✓	✓	✓	✓
BICGSTAB	✓	✓	✓	✓
CG	✓	✓	✓	✓
CGS	✓	✓	✓	✓
GMRES	✓	✓	✓	✓
IDR	✓	✓	✓	✓
(Block-)Jacobi	✓	✓	✓	✓
ILU/IC		✓	✓	✓
Parallel ILU/IC	✓	✓	✓	✓
Parallel ILUT/ICT	✓	✓	✓	✓
Sparse Approximate Inverse	✓	✓	✓	✓
Krylov solvers				
Preconditioners				
Utilities				

Extending to Intel GPUs



- Bi-Weekly technical meetings with Intel
- Long list of bug reports, feature requests, performance data discussions, documentation improvements ...

cuBLAS backend (and potentially other domains) fails with latest LLVM builds #223

[New issue](#)

[Closed](#) mmeterel opened this issue 22 days ago · 3 comments

Summary

As first observed in #219 many tests in cuBLAS backend is failing with latest LLVM builds.

Version

I have tried LLVM commit: 66361038b63caee566fc9648f5da50b7422b83 and got the below tests failing (showing only a few of them)

```
1 - BLAS/RT/NrmTestsSuite/NrmTests.RealSinglePrecision/Column_Major_TITAN RTX (Failed)
3 - BLAS/RT/NrmTestsSuite/NrmTests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
5 - BLAS/RT/NrmTestsSuite/NrmTests.ComplexDoublePrecision/Row_Major_TITAN RTX (Failed)
7 - BLAS/RT/NrmTestsSuite/NrmTests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
19 - BLAS/RT/ImaxTestsSuite/ImaxTests.RealDoublePrecision/Column_Major_TITAN RTX (Failed)
23 - BLAS/RT/ImaxTestsSuite/ImaxTests.ComplexDoublePrecision/Column_Major_TITAN RTX (Failed)
25 - BLAS/RT/ImaxTestsSuite/ImaxTests.ComplexDoublePrecision/Row_Major_TITAN RTX (Failed)
36 - BLAS/RT/DotTestsSuite/DotTests.ComplexSinglePrecision/Column_Major_TITAN RTX (Failed)
67 - BLAS/RT/AusmTestsSuite/AusmTests.ComplexSinglePrecision/Column_Major_TITAN RTX (Failed)
85 - BLAS/RT/DotTestsSuite/DotTests.ComplexSinglePrecision/Column_Major_TITAN RTX (Failed)
```

From [DPC++ AoT documentation](#), not clear:

- The options are also required at linking time? Unused in files without kernels?
- Any example of other projects integrating AoT in a CMake setup?

Intel Compiler (Fortran/C/C++/L0) - Intel Discrete GPU Accelerator - Joint Laboratory for System Evaluation (enl.gov)
hang_atomic_on_local
Ticket number: CMPLRLLVM-36572 (works in PVC, but still fails on ATS node)
related to driver not compiler self

tid % subgroup size >= 4 gives wrong division

(double) 1/a gives wrong result when the tid % subgroup size >= 4. For example, when a = 1.0733829563753890 1/a should be 0.9316293125835232 if (local_id == assign_id) { a = double(1)/a; } when assign_id >= 4, Gen9 GPU still give the correct result when assign_id >= 4, Gen9 GPU gives wrong 0.9316293125835232 CPU has more worse result

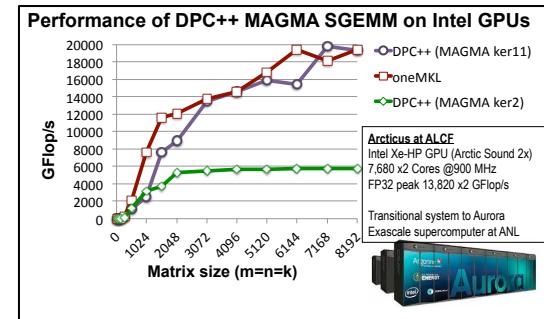
It is connected to optimizations (not reproducible with O0). fp-speculation=off do not improve results. Ticket number: XDEPS-4031 ()

IMAGE LAYERS

- 1 ADD file ... in /
- 2 OMD ["bosh"]
- 3 ENV INARCH=x86_64

Jevcloud node issue

- sycl-ls/clinfo does not give any of s001-n225, s011-n006
- no gpu on the nodes s001-n232, s001-n233, s011-n008
- github.com is not accessible on Jevcloud



... but also docker image contributions and bug fixes!

ginkgohub/oneapi:cuda11.6

DIGEST: sha256:8bc4c18d79e75b183ac1leafcd0753365c6e1a94edc3046d9a0eb8ba2d7bd934
OS/ARCH: linux/amd64
COMPRESSED SIZE: 0
LAST PUSHED: 22 days ago by [yhmtsai](#)

fix cuda/nip backend location #219

[I-Merged](#) mkrainiuk merged 2 commits into [oneapi-rc1-devel](#) from [yhmtsai/fix_cuda_backend_location](#) 20 days ago

Description

From Intel/lm#6407, it moves almost all headers from CL/sycl to sycl I followed #219 way to fix it. make the header can use sycl if they exist and allow the old intel lmv. I also update the CL/sycl.hpp which are not changed before.

All Submissions

Do all unit tests pass locally? Attach a log. A: It is a compiling issue.
✓ Have you formatted the code using clang-format?

Bug fixes

Have you added relevant regression tests? A: It is a compiling issue.
✓ Have you included information on how to reproduce the issue (either in a GitHub issue or in this PR)?
Reproduce:
compile the latest intel lmv and this repo, it will not be able to compile due to missing headers.

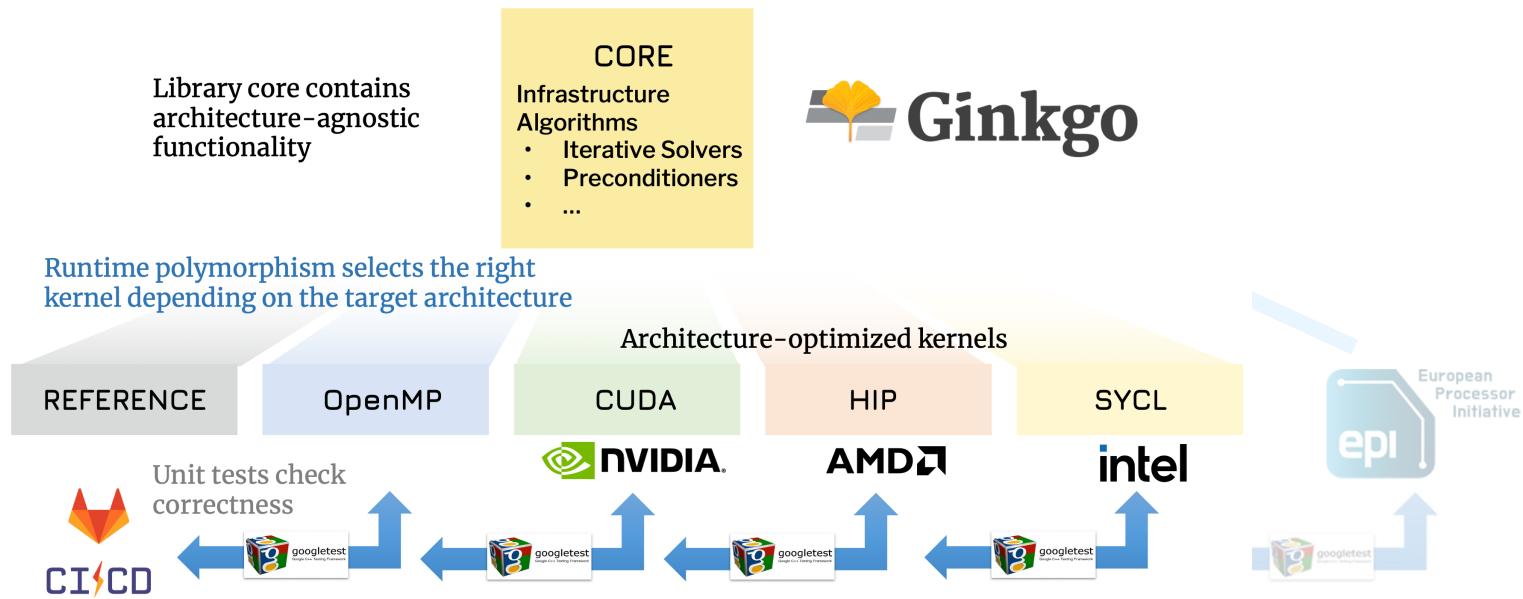
yhmtsai added 2 commits 2 months ago

- use the correct sycl path after Intel/lm#6407
- Fix the missing sycl/sycl.hpp

mmeterel commented on Aug 1

@yhmtsai Thanks for the PR. Is the description from sycl/CL to CL correct? My understanding is all header files moved

Portability as central design principle



This software design gives portability, performance, and sustainability.

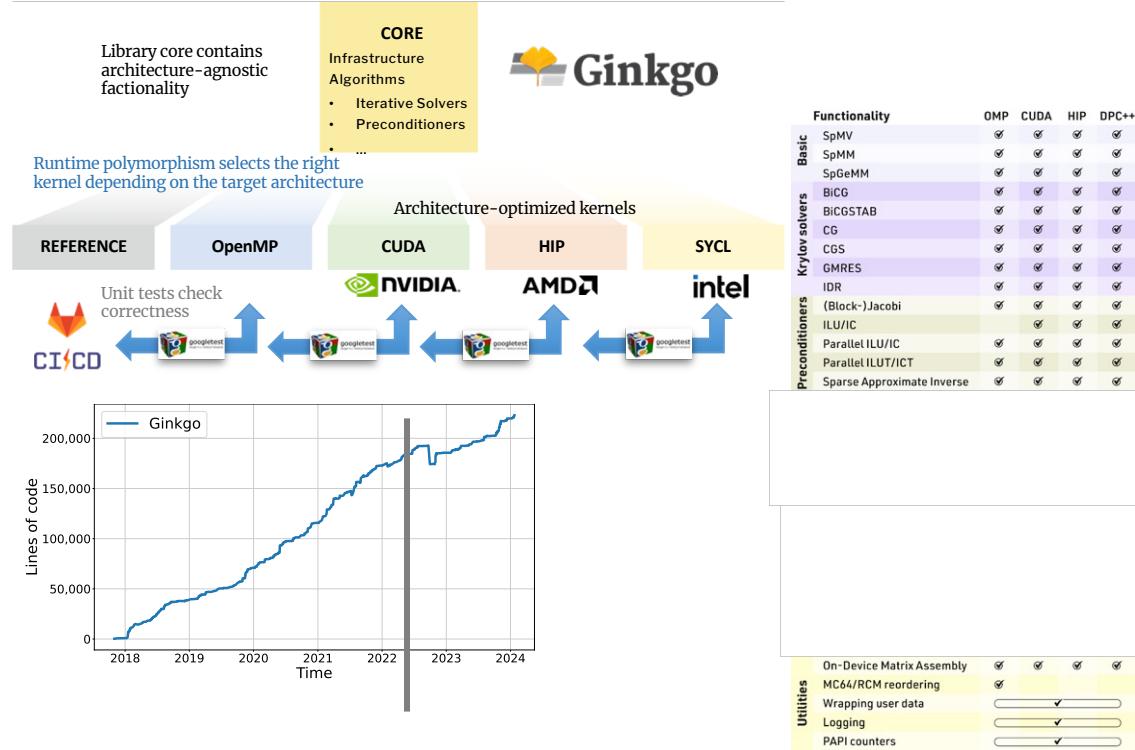
Focus efforts as lightweight tool in ECP to address challenges



Focus efforts



- Mixed precision
 - *Address recent hardware trends (tensor cores, etc.)*
- Batched Routines
 - *Address application requirements*



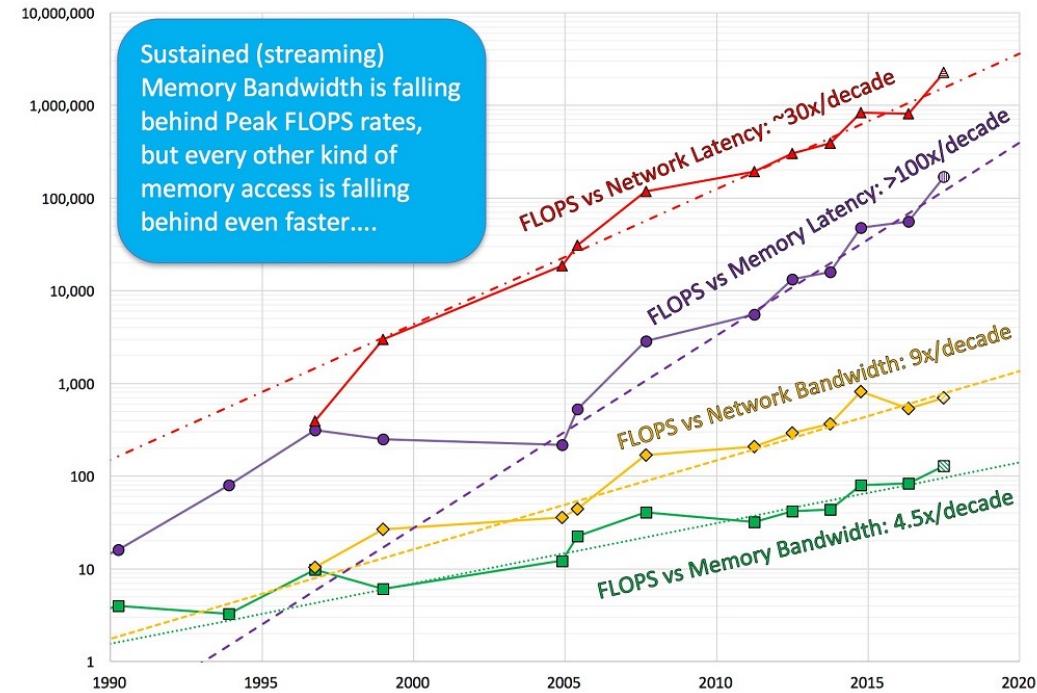
Mixed precision focus effort



Form Factor	H100 SXM
FP64	34 teraFLOPS
FP64 Tensor Core	67 teraFLOPS
FP32	67 teraFLOPS
TF32 Tensor Core	989 teraFLOPS ^z
BFLOAT16 Tensor Core	1,979 teraFLOPS ^z
FP16 Tensor Core	1,979 teraFLOPS ^z
FP8 Tensor Core	3,958 teraFLOPS ^z
INT8 Tensor Core	3,958 TOPS ^z
GPU memory	80GB
GPU memory bandwidth	3.35TB/s

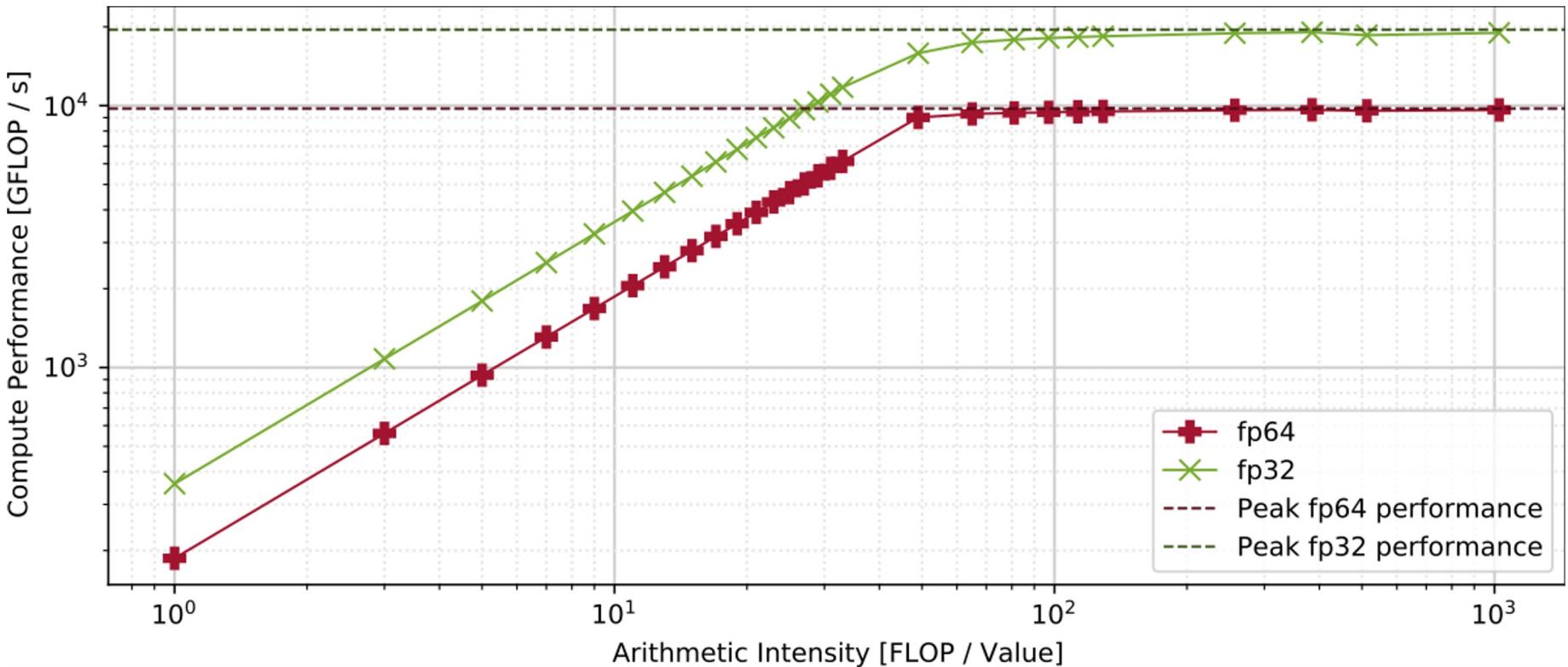
- (Dense) Matrix Performance
 - > Vector Operation Performance
- Low Precision Performance
 - > High Precision Performance

Balance: computation vs. communication



Trends in the relative performance of floating-point arithmetic and several classes of data access for select HPC servers over the past 25 years. Source: John McCalpin

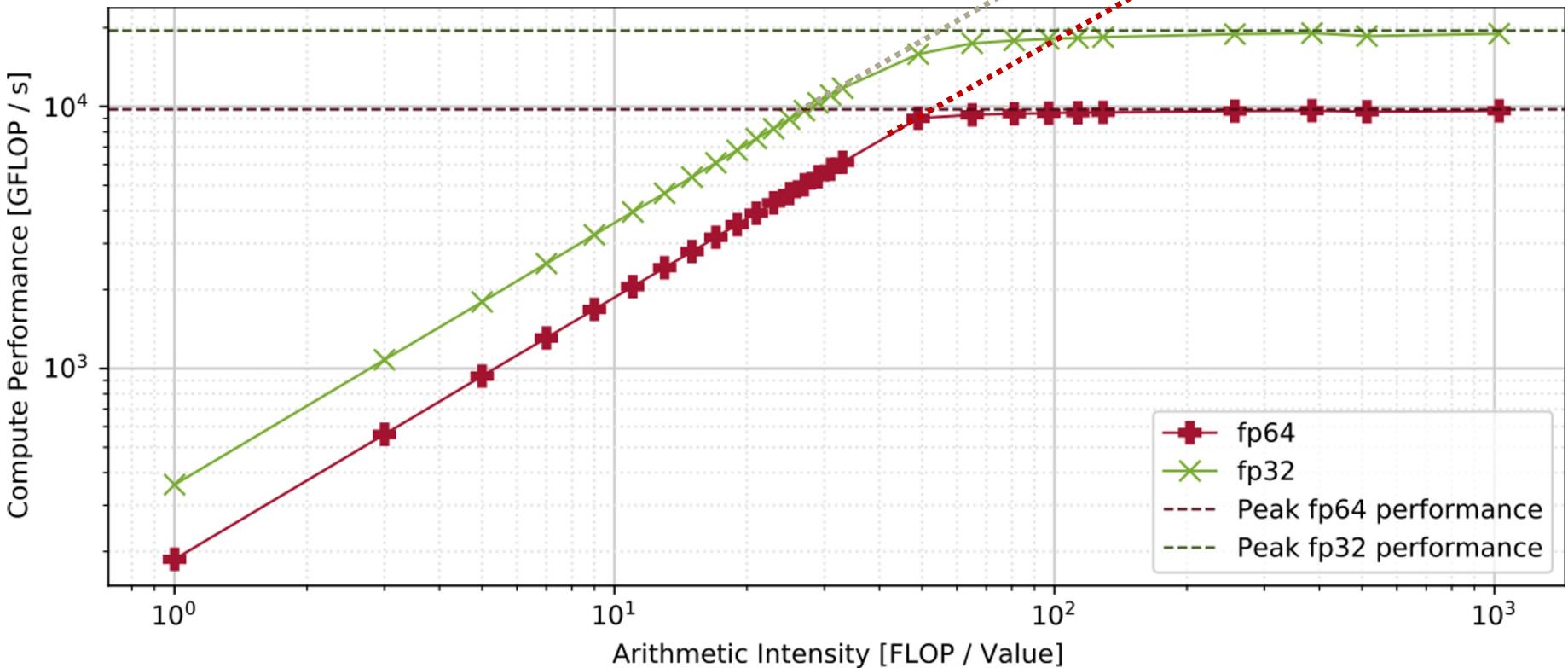
NVIDIA A100



Matrix fp32

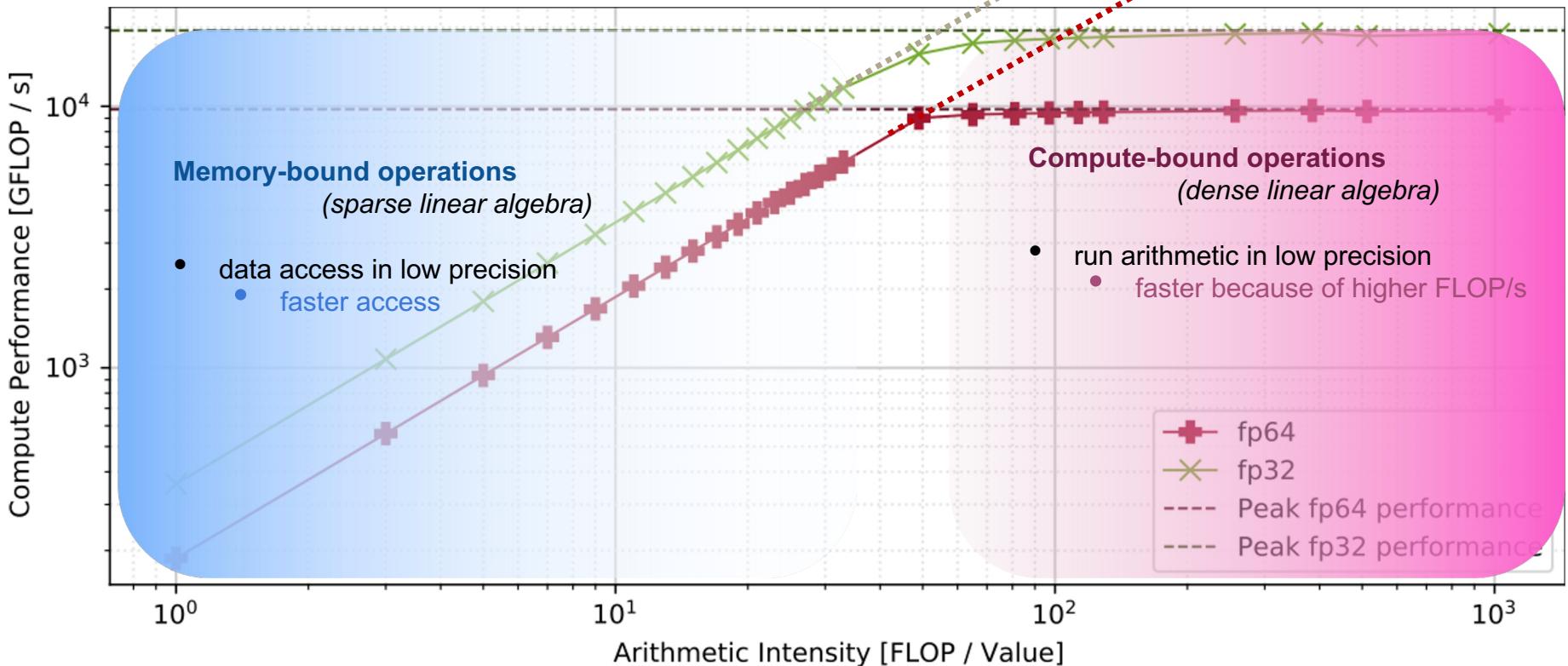


Matrix fp64



Matrix fp32

Matrix fp64



Linear System $Ax=b$ with $\text{cond}(A) \approx 10^7$
 (apache2 from SuiteSparse) NVIDIA V100 GPU

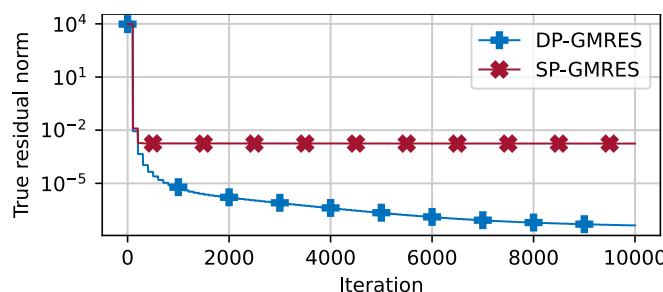
Double precision GMRES

Initial residual norm $\text{sqrt}(r^T r)$: 9670.36
 Final residual norm $\text{sqrt}(r^T r)$: $9.6639e-09$
 GMRES iteration count: 23271
 GMRES execution time: 43801 ms

Single precision GMRES

Initial residual norm $\text{sqrt}(r^T r)$: 9670.36
 Final residual norm $\text{sqrt}(r^T r)$: 0.00175464
 GMRES iteration count: 25000
 GMRES execution time: 27376 ms

~2x faster!

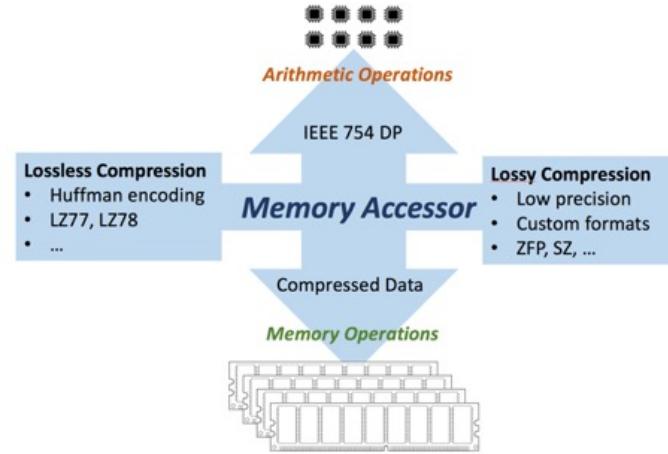


forward error \approx (unit round-off) * (linear system's condition number)

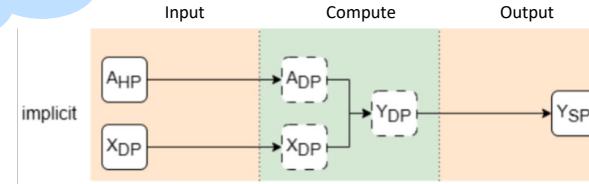
N. Higham: Accuracy and stability of numerical algorithms. SIAM, 2002.

Mixed precision focus effort

- Traditionally, we use a strong coupling between the precision formats used for **arithmetic operations** and **storing data**.
- *We should compute in fp64*
- *Data should be compressed for main memory access (low precision/compression)*
- *Compression / Conversion needs to happen on-the-fly*

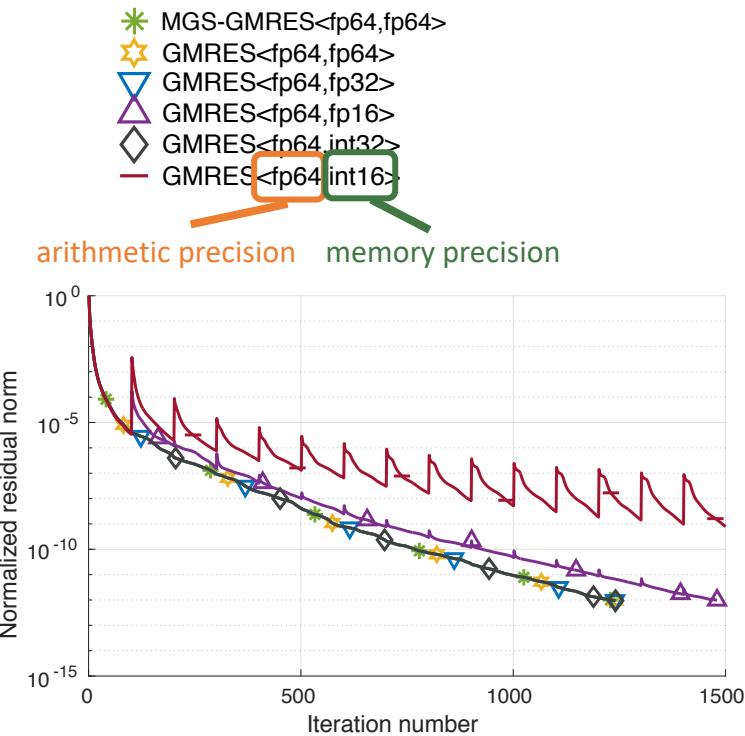


See Felix Liu's thesis



Compressed Basis (CB-) GMRES

- Use double precision in all arithmetic operations;
- Store Krylov basis vectors in lower precision;
 - Search directions are no longer DP-orthogonal;
 - Hessenberg system maps solution to “perturbed” Krylov subspace;
 - Additional iterations may be needed;
 - As long as the loss-of-orthogonality is moderate, we should see moderate convergence degradation;



Linear System $Ax=b$ with $\text{cond}(A) \approx 10^7$
(apache2 from SuiteSparse) NVIDIA V100 GPU

Double precision GMRES

Initial residual norm $\sqrt{r^T r}$: 9670.36
Final residual norm $\sqrt{r^T r}$: $9.6639e-09$
GMRES iteration count: 23271
GMRES execution time: 43801 ms

Single precision GMRES

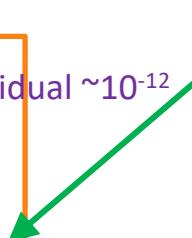
Initial residual norm $\sqrt{r^T r}$: 9670.36
Final residual norm $\sqrt{r^T r}$: 0.00175464
GMRES iteration count: 25000
GMRES execution time: 27376 ms

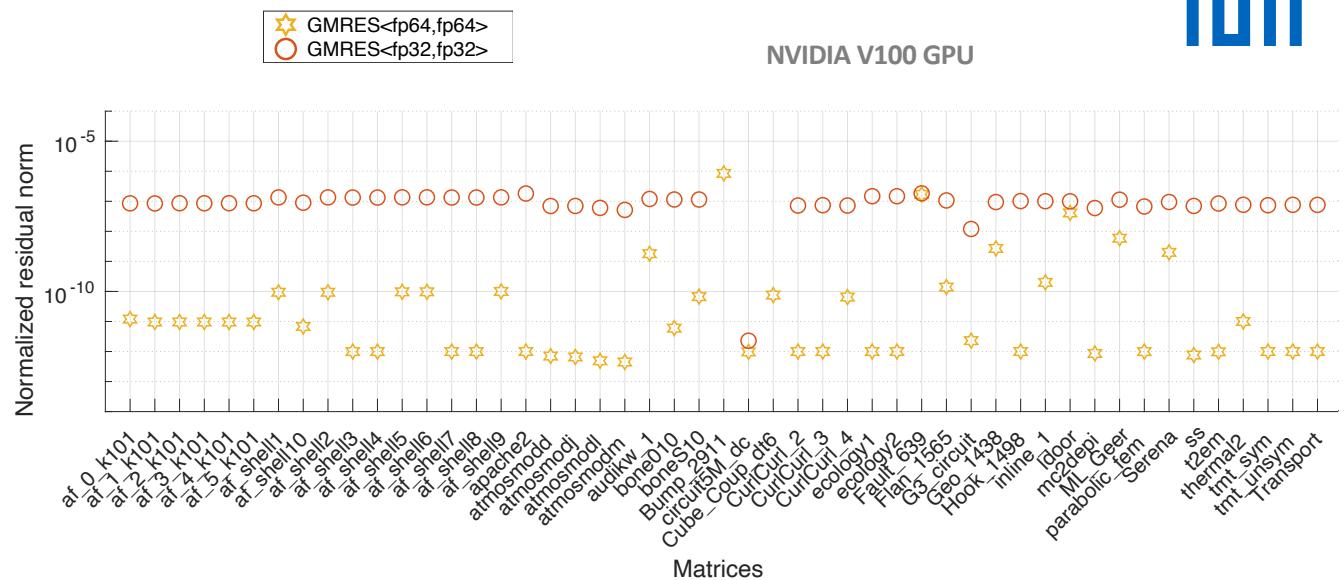
Compressed Basis GMRES

Initial residual norm $\sqrt{r^T r}$: 9670.36
Final residual norm $\sqrt{r^T r}$: $9.6591e-09$
GMRES iteration count: 23271
GMRES execution time: 29369 ms

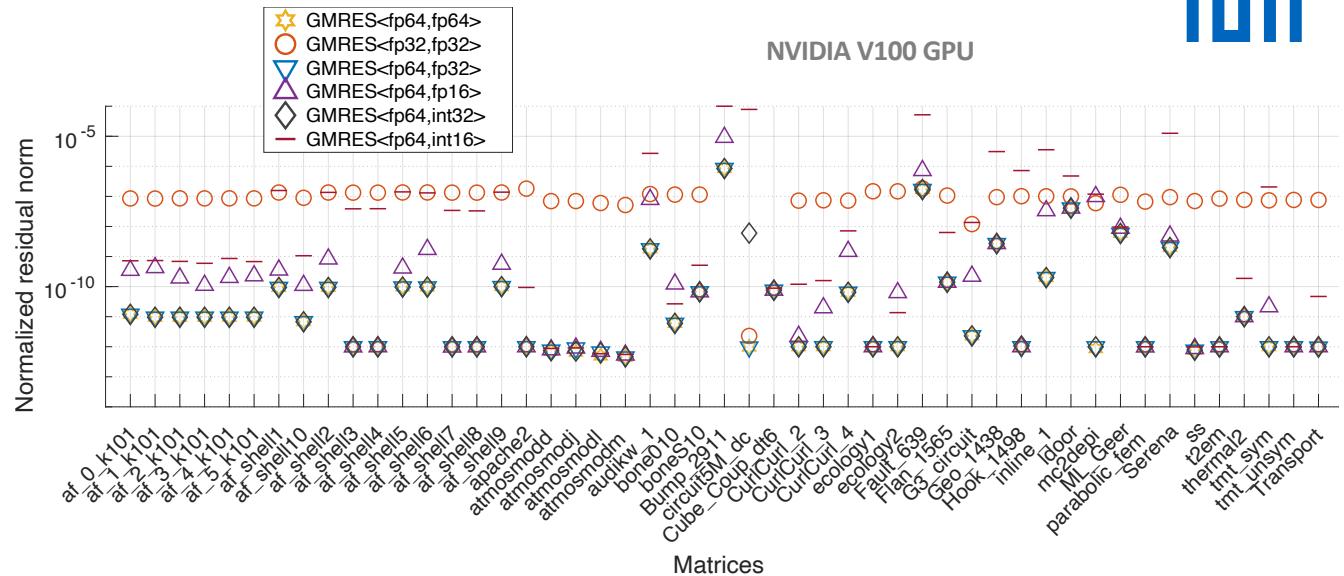
Relative residual $\sim 10^{-12}$

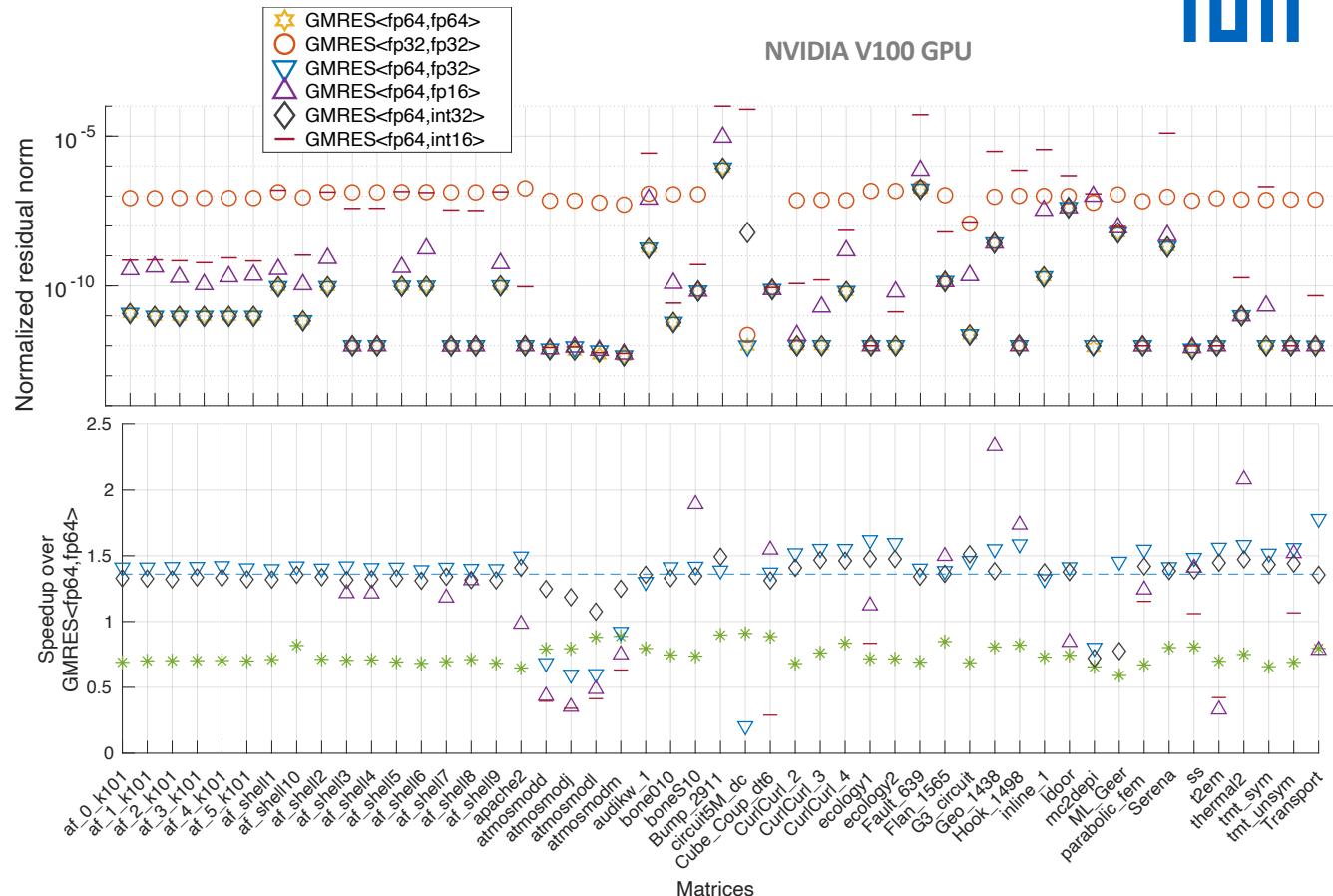
Accuracy of DP GMRES
Performance similar to SP GMRES





- CB-GMRES using 32-bit storage preserves DP accuracy
(SP-GMRES does not)





- CB-GMRES using 32-bit storage preserves DP accuracy (SP-GMRES does not)
- Speedups problem-dependent
- Speedup $\varnothing 1.4x$ (for restart 100)
- 16-bit storage mostly inefficient



Aliaga JI, Anzt H, Grützmacher T, Quintana-Ortí ES, Tomás AE. Compressed basis GMRES on high-performance graphics processing units. *The International Journal of High Performance Computing Applications*. 2022;0(0). doi:[10.1177/10943420221115140](https://doi.org/10.1177/10943420221115140)

Mixed precision AMG on GPUs

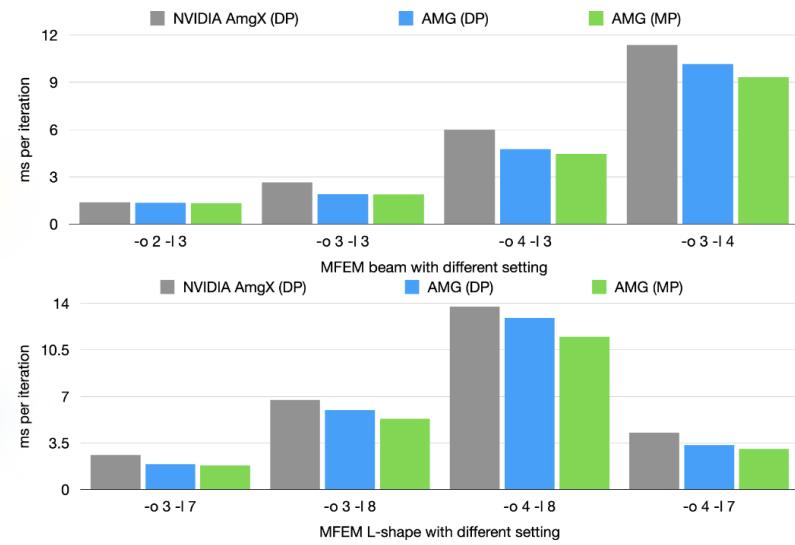
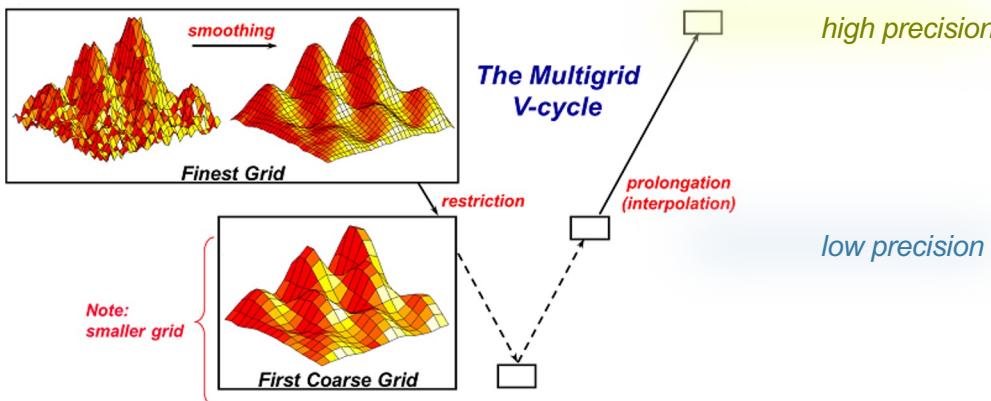


Mike Tsai

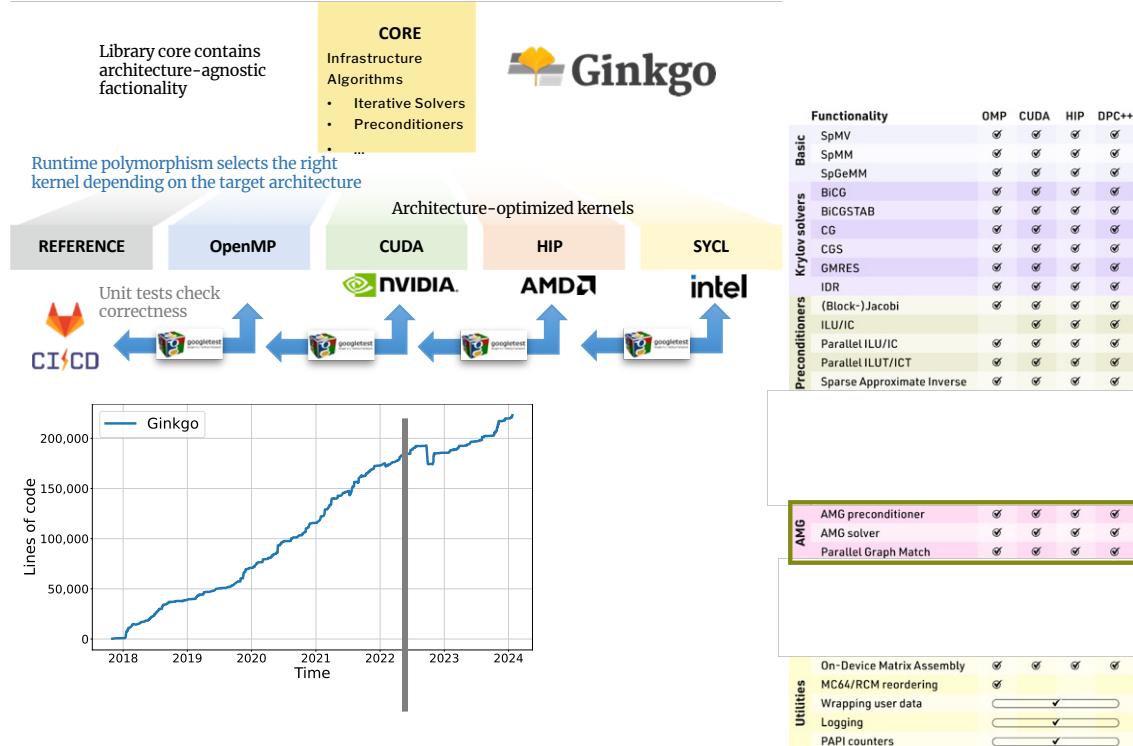
- Preconditioning iterative solvers

- Idea: Approximate inverse of system matrix to make the system “easier to solve”: $P^{-1} \approx A^{-1}$
and solve $Ax = b \Leftrightarrow P^{-1}Ax = P^{-1}b \Leftrightarrow \tilde{A}x = \tilde{b}$

- Mixed Precision Multigrid Preconditioner



Mixed precision focus effort

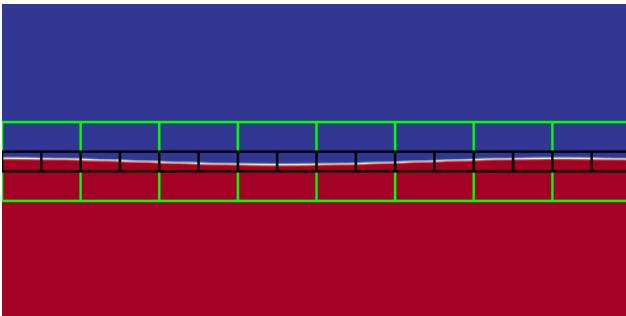


Batched focus effort – Combustion Simulations

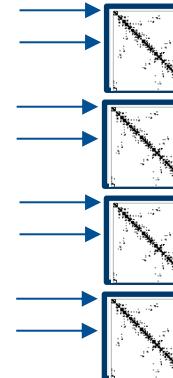
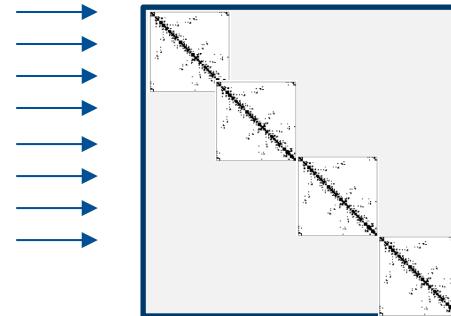
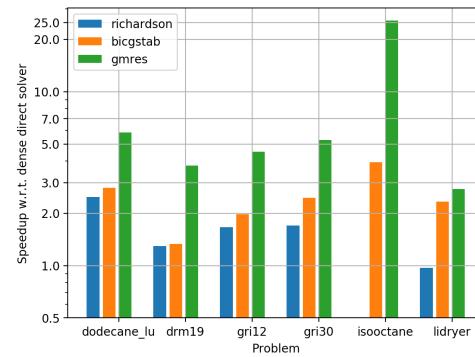
Batched iterative solvers for SUNDIALS / PeleLM

PeleLM is a parallel, adaptive mesh refinement (AMR) code that solves the reacting Navier-Stokes equations in the low Mach number regime. The core libraries for managing the subcycling AMR grids and communication are found in the [AMReX source code](#).

<https://amrex-combustion.github.io/PeleLM/overview.html>



Problem	Size	Non-zeros (\mathbf{A})	Non-zeros ($\mathbf{L+U}$)
dodecane_lu	54	2,332 (80%)	2,754 (94%)
drm19	22	438 (90%)	442 (91%)
gr12	33	978 (90%)	1,018 (93%)
gr30	54	2,560 (88%)	2,860 (98%)
isoctane	144	6,135 (30%)	20,307 (98%)
lidryer	10	91 (91%)	91 (91%)



Batched Sparse Iterative Solvers for Computational Chemistry Simulations on GPUs

Publisher: IEEE

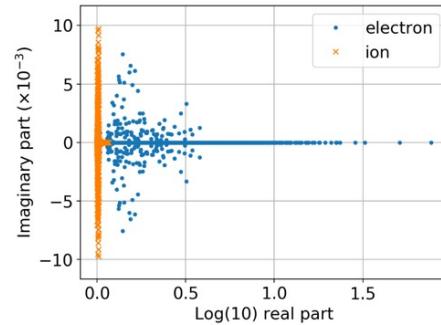
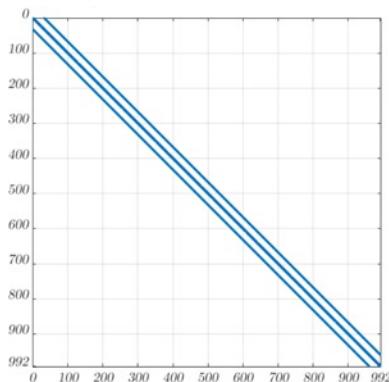
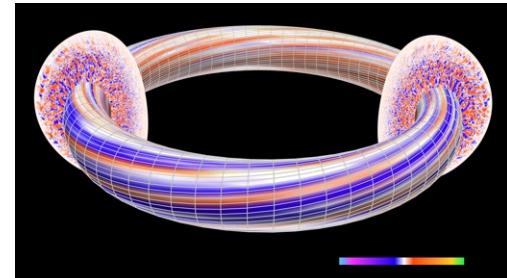
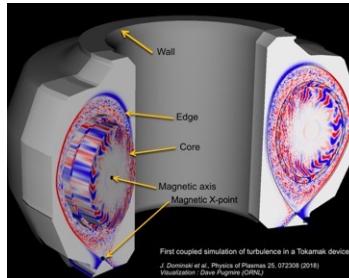
Cite This

PDF

Batched focus effort – Fusion Plasma Simulations

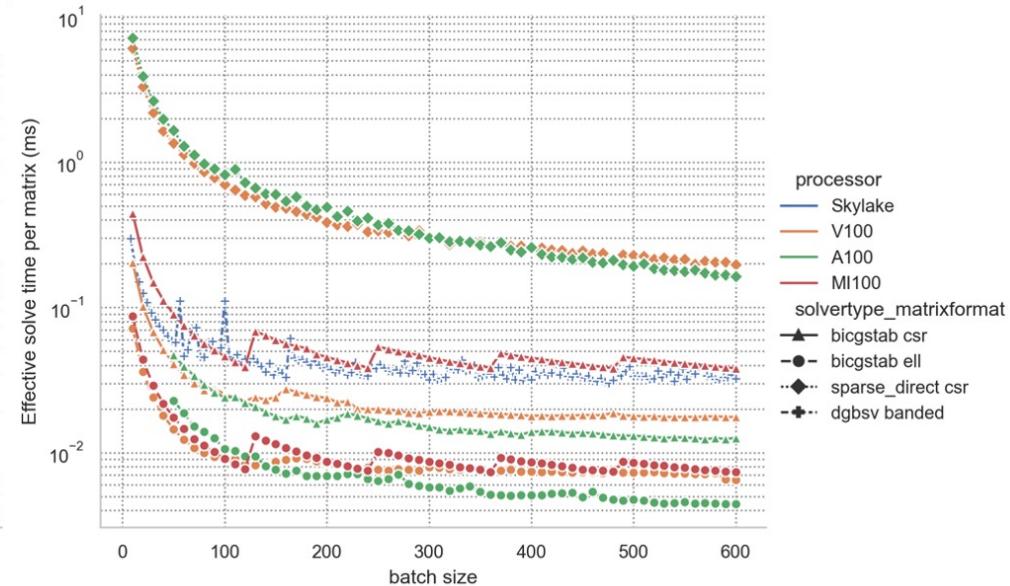
XGC is a gyrokinetic particle-in-cell code, which specializes in the simulation of the edge region of magnetically confined thermonuclear fusion plasma. The simulation domain can include the magnetic separatrix, magnetic axis and the biased material wall. XGC can run in total-delta-f, and conventional delta-f mode. The ion species are always gyrokinetic except for ETG simulation. Electrons can be adiabatic, massless fluid, driftkinetic, or gyrokinetic.

Source: https://xgc.pppl.gov/html/general_info.html

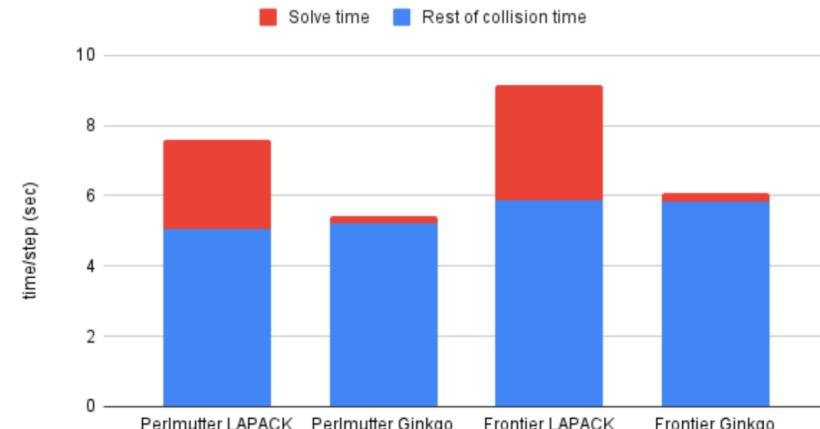


- Two species
- Ions easy to solve
- Electrons hard to solve
- Banded matrix structure
- Non-symmetric, need BiCGSTAB
- $n = \sim 1,000$
- $nz = \sim 9,000$

Batched focus effort – Fusion Plasma Simulations



XGC collision time reduction (64 nodes)



Sparse direct solvers for power grid simulations

Mathematical Formulation of the ExaSGD Core Challenge Security constrained multiperiod AC optimal power flow analysis

Posed as an optimization problem:

Find

$$\min_{x_t, y_{tsk}} (\sum_t F_t(x_t) + \sum_{tsk} G_{tsk}(x_t + y_{tsk}))$$

generator fuel cost

wind curtailment,
load shedding,
power imbalance, etc.

flow definitions,
power balance

bounds: generator power,
voltage, branch flow
generator ramping limit

Subject to:

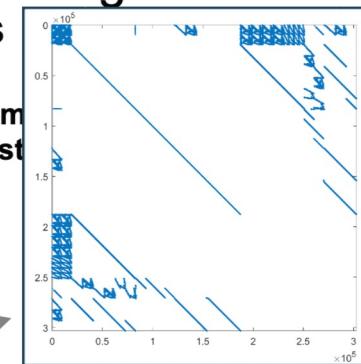
$$H_{tsk}(x_t, y_{tsk}) = 0$$

$$Q_{tsk}(x_t, y_{tsk}) \leq 0$$

$$R_t(x_t, x_{t+1}) \leq 0$$

The optimization problem
the underlying linear system

$$\begin{bmatrix} K_1 & & & \\ & K_2 & & \\ & & K_3 & \\ & & & \dots \end{bmatrix} \begin{bmatrix} B_1^T & B_2^T & B_3^T & \dots & B_N^T & K_0 \end{bmatrix} \begin{bmatrix} x \end{bmatrix} = \begin{bmatrix} v_N \\ \vdots \\ r_N \\ \vdots \\ r_0 \end{bmatrix}$$



- The characteristic block-arrow coupling structure can be exploited to decompose the optimization problem, nevertheless there is no solver that can tackle this on a GPU-based architecture.

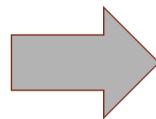
Sparse direct solvers for power grid simulations

Underlying KKT Linear System Properties

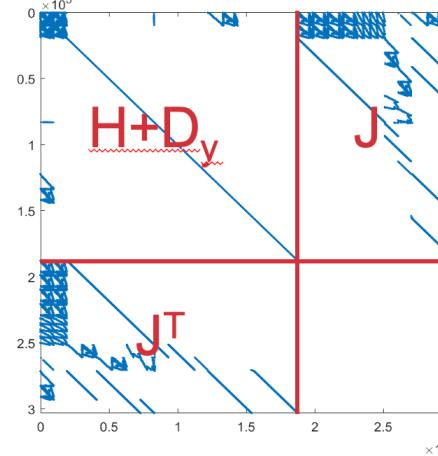
See Felix Liu's thesis

- Security constrained optimal power flow analysis.
- The interior method strategy leads to symmetric indefinite linear systems

$$\underbrace{\begin{bmatrix} H + D_y & J \\ J^T & 0 \end{bmatrix}}_{K_k} \underbrace{\begin{bmatrix} \Delta y \\ \Delta \lambda \end{bmatrix}}_{\Delta x_k} = \underbrace{\begin{bmatrix} r_y \\ r_\lambda \end{bmatrix}}_{r_k},$$



- J – sparse constraints Jacobian,
- H – sparse Hessian,
- D_y – arises from log-barrier function



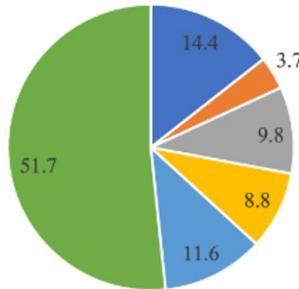
Typical sparsity pattern of optimal power flow matrices: No obvious structure that can be used by linear solver.

- The challenge: we need to solve a long sequences of such systems.

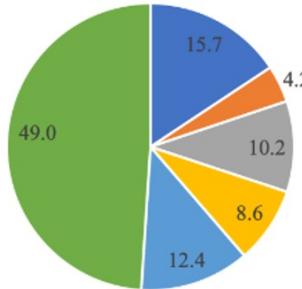
Sparse direct solvers for power grid simulations

Grid	Buses	Generators	Lines	$N(K_k)$	$\text{nnz}(K_k)$
Northeastern US	25 K	4.8 K	32.3 K	108 K	1.19 M
Eastern US	70 K	10.4 K	88.2 K	296 K	3.20 M
Western and Eastern US	82 K	13.4 K	104.1 K	340 K	3.73 M

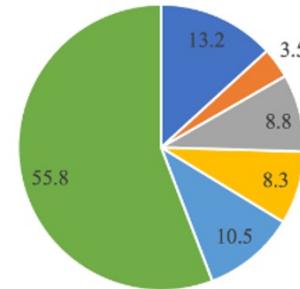
■ Hessian ■ Constraints ■ Constraint Jacobian ■ Other ■ Solve ■ Factorize



(a) Northeast U.S. grid



(b) Eastern U.S. grid



(c) Eastern and Western U.S. grids

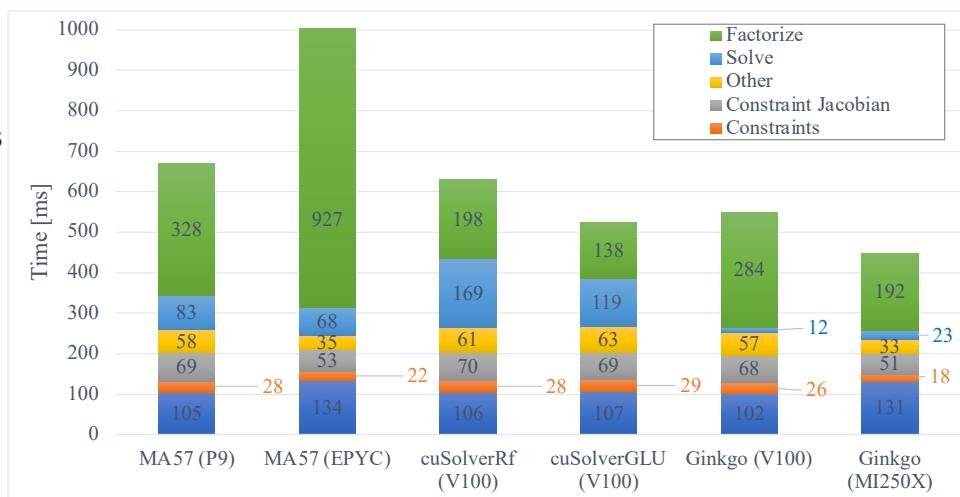
Sparse direct solvers for power grid simulations

You are viewing Presentation's screen [Switch](#) View Options ▾

Liner Solver Performance within Optimization Algorithm Average per iteration times (including first iteration on CPU)



- Each GPU solution outperforms all CPU baselines.
- Ginkgo performance improves on a better GPU.
- Iterative refinement configuration affects linear solver performance and optimization solver convergence.
- Ginkgo is the first GPU-resident sparse direct linear solver.



Solver	Factorize	Solve	Other	Constraint Jacobian	Constraints	Total
MA57 (P9)	328	83	58	69	105	675
MA57 (EPYC)	927	68	35	53	134	1129
cuSolverRF (V100)	198	169	61	70	106	533
cuSolverGLU (V100)	138	119	63	69	107	527
Ginkgo (V100)	284	57	68	29	102	512
Ginkgo (MI250X)	192	33	51	12	131	517

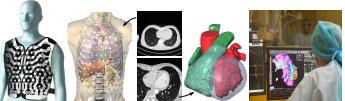
Multiple promising GPU-resident sparse linear solvers

29

Now, after the completion of ECP



- Sustainable software design ready for the addition of new backends.
- EuroHPC Project MICROCARD uses Ginkgo



<https://www.microcard.eu>



- BMBF PDExa and ExaSIM projects use Ginkgo



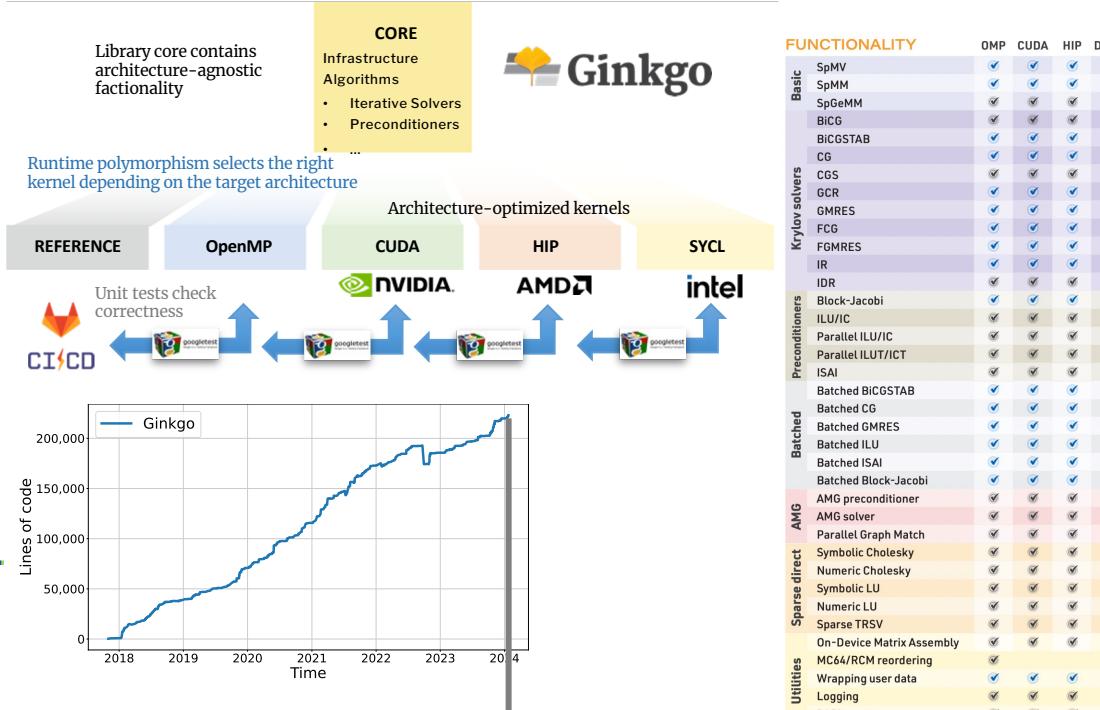
The Open Source CFD Toolbox

CEED/NekRS

Mirror of NekRS - GPU-oriented version of Nek5060. Please use the official repository. New git commits will be pushed to this mirror to create issues and pull requests.

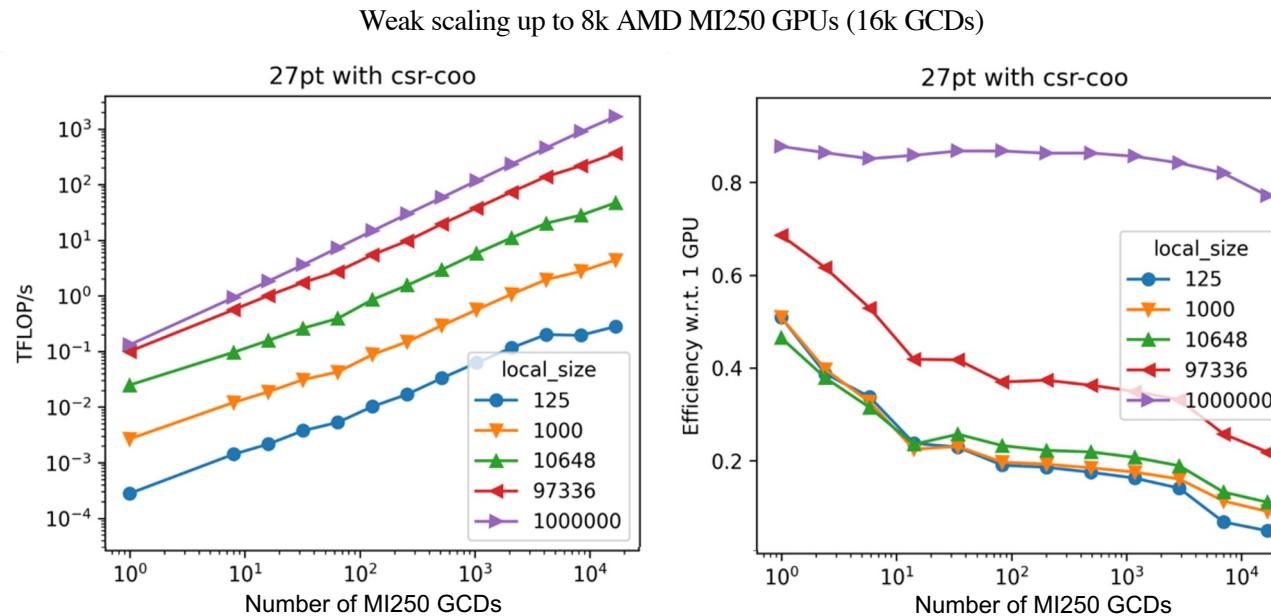


- Companies are evaluating Ginkgo



Scalability of Ginkgo on Frontier (#1 TOP500, AMD MI250)

Weak scaling: problem size increases with parallel resources



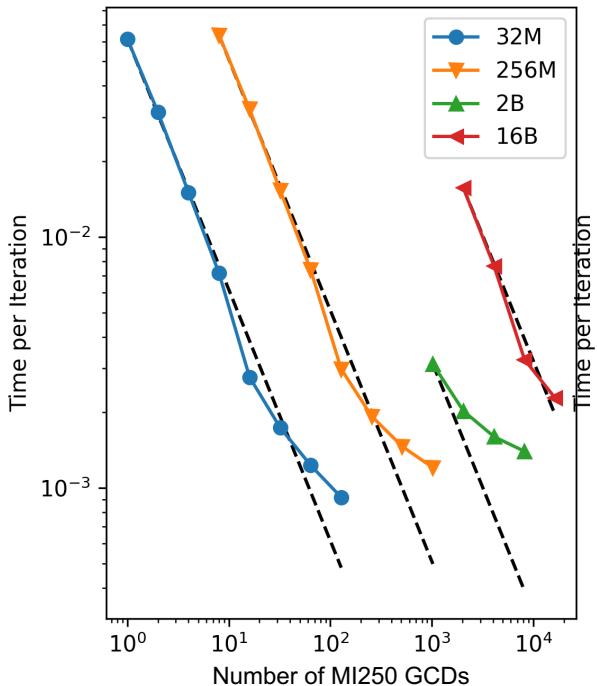
Significant Compute Waste!

Scalability of Ginkgo on Frontier (#1 TOP500, AMD MI250)

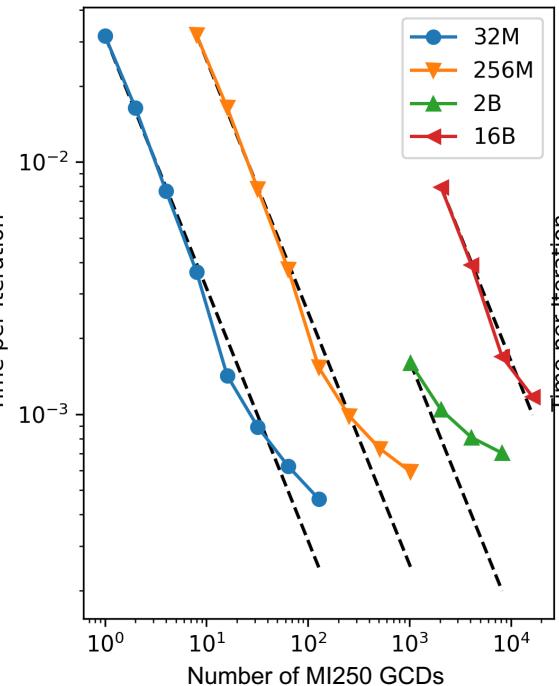
Strong scaling: problem size increases with parallel resources

Strong scaling up to 8k AMD MI250 GPUs (16k GCDs)

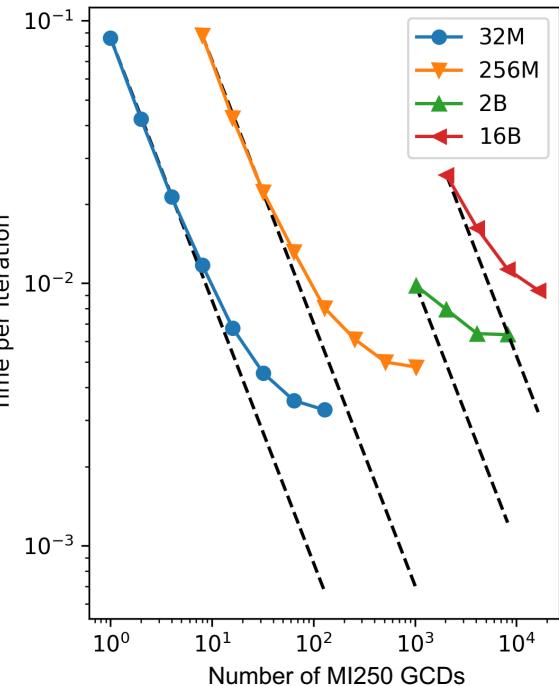
Strong Scaling - Bicgstab Local Jacobi



Strong Scaling - Cg Local Jacobi



Strong Scaling - Gmres Local Jacobi



Lessons learnt over the years

- ECP earmarking roughly half the budget to Software & App development is a game changer.
 - Central component for the success of ECP.
 - This concept needs to – and does become - the blueprint for other nations, companies, and projects.
- Workforce recruitment and workforce retention are the key to success in software development.
 - Money does not write software. RSEs do. We need to create attractive career plans.
 - We need to make research software development attractive to students. Academic recognition. Industry career paths.
- Anticipating the future in hardware development accelerates the porting process.
 - Blueprints and early access systems both useful.
 - Interaction with industry is mutually beneficial.
- Management, tools, and strategic initiatives, interaction and collegial behavior are important.
 - Jira/Notion/[...] milestones and deliverables give projects and collaborative interactions a structure and timeline.
 - Strategic focus groups, conferences, and meetings bring experts together and create collaboration.
 - Listen to the application needs. Value input and acknowledge collaborators.