Lossy Compression and Mixed Precision Strategies for Memory-Bound Linear Algebra

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Running iterative methods in different precision formats

Linear System Ax=b with cond(A) ≈ 10^7

(apache2 from SuiteSparse) **NVIDIA V100 GPU**

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Double precision GMRES
Initial residual norm
sqrt(r^t r): 9670.36
Final residual norm
sqrt(r^T r): 9.6639e-09
GMRES iteration count: 23271
GMRES execution time: 43801 ms
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Relative residual ~10^{-12}

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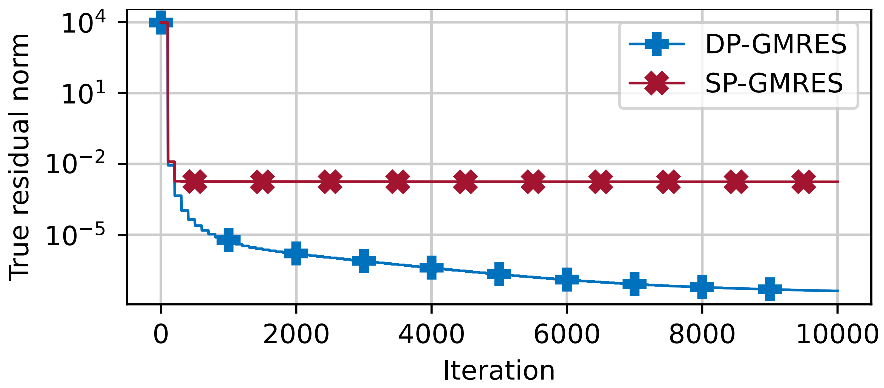

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Relative residual ~10^{-7}

~2x faster!

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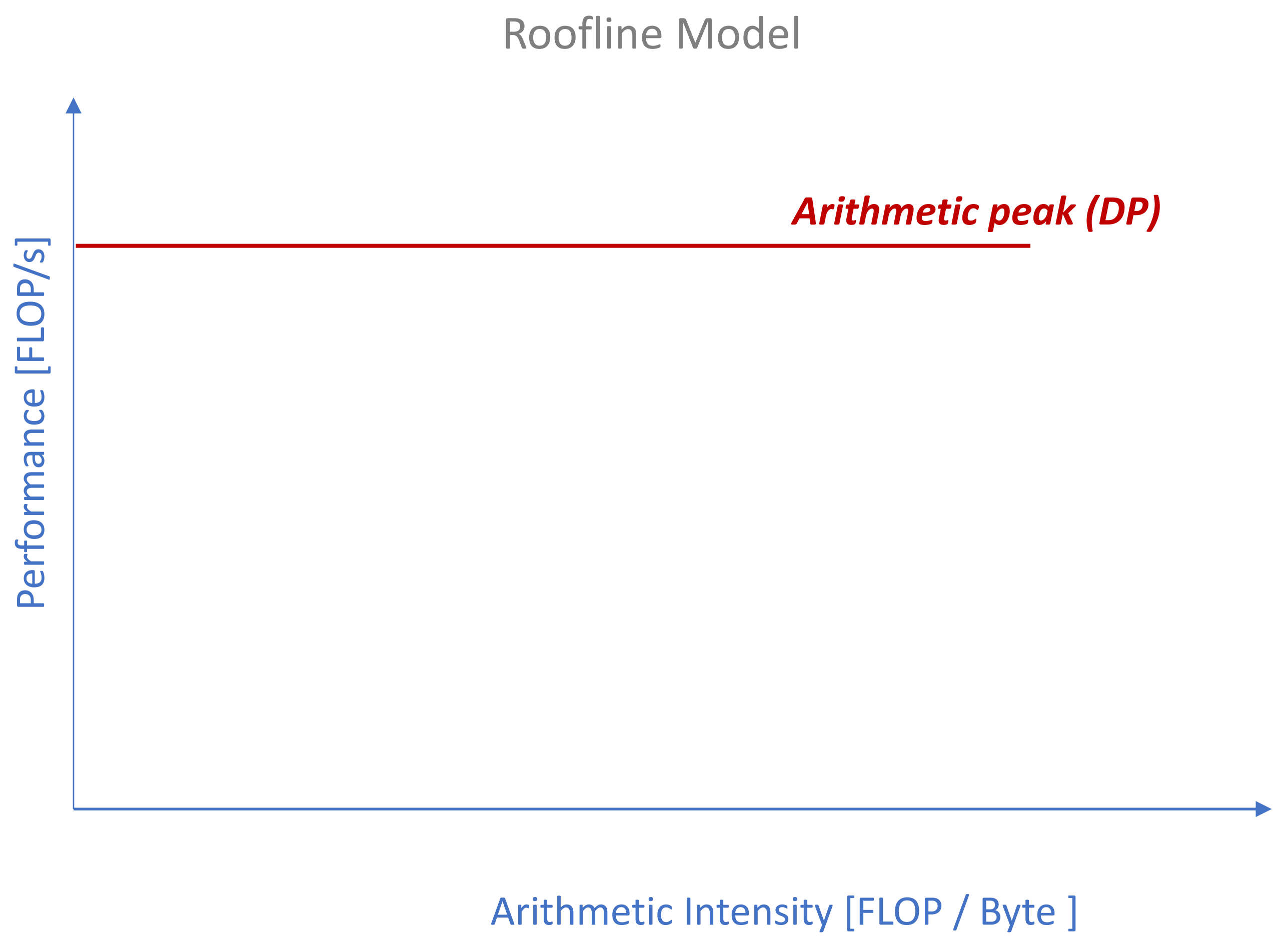

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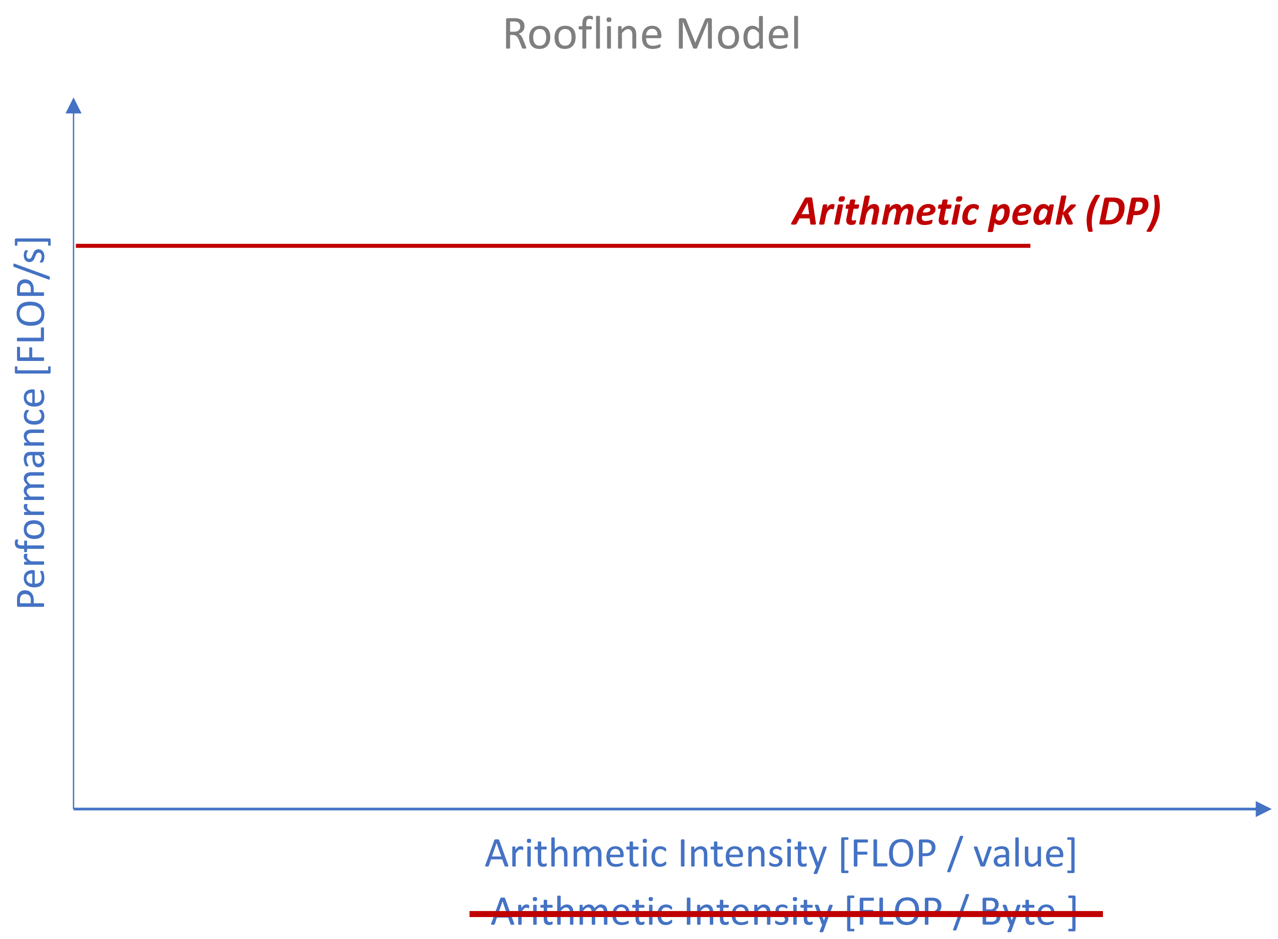

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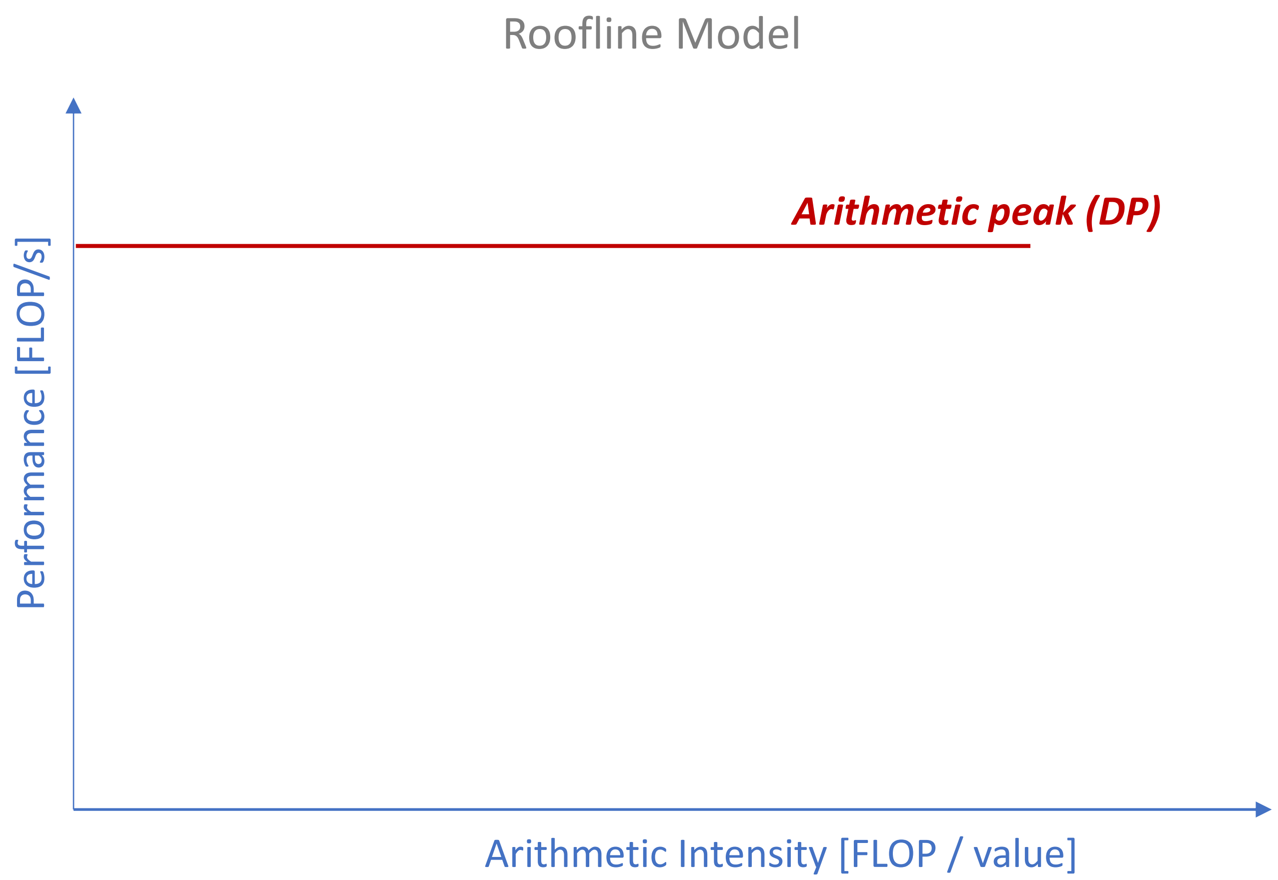

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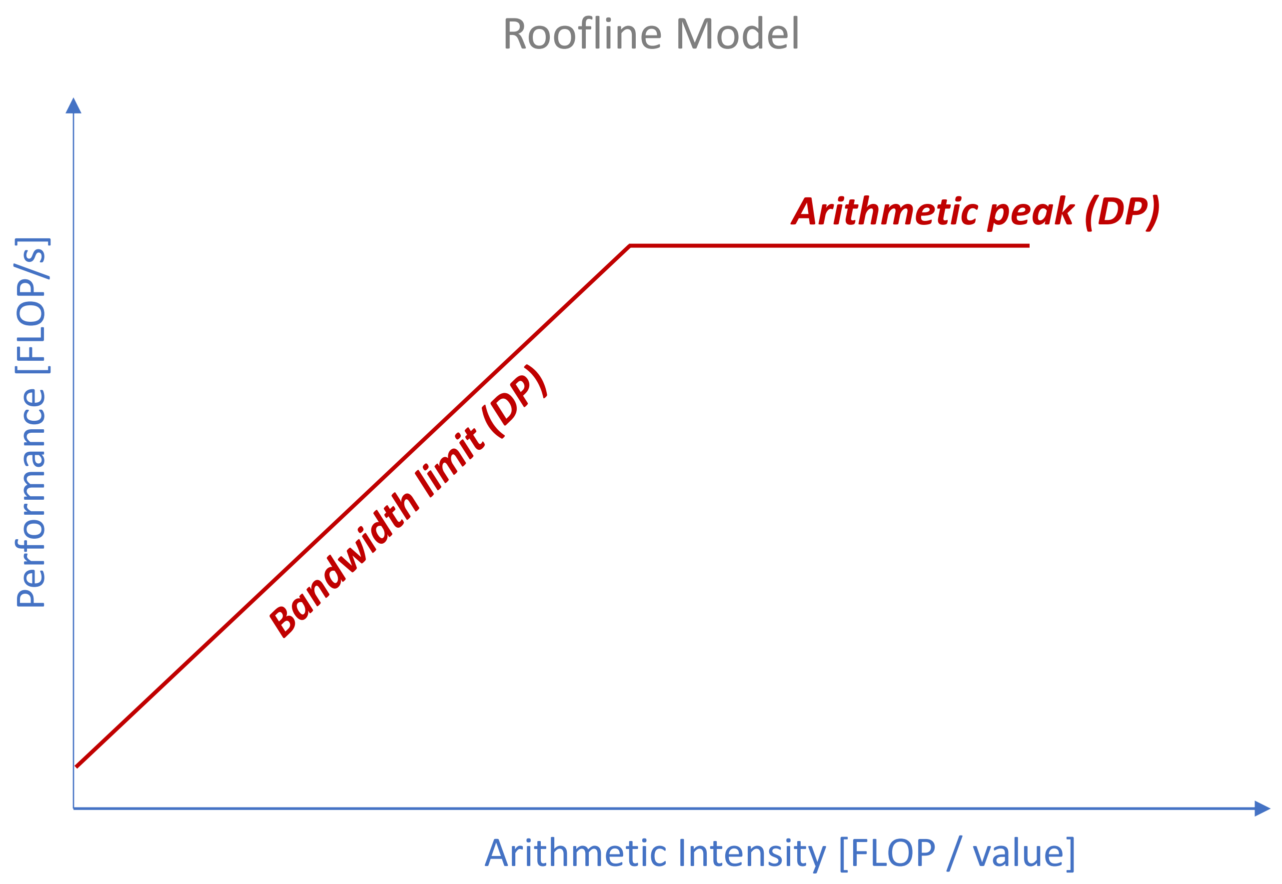

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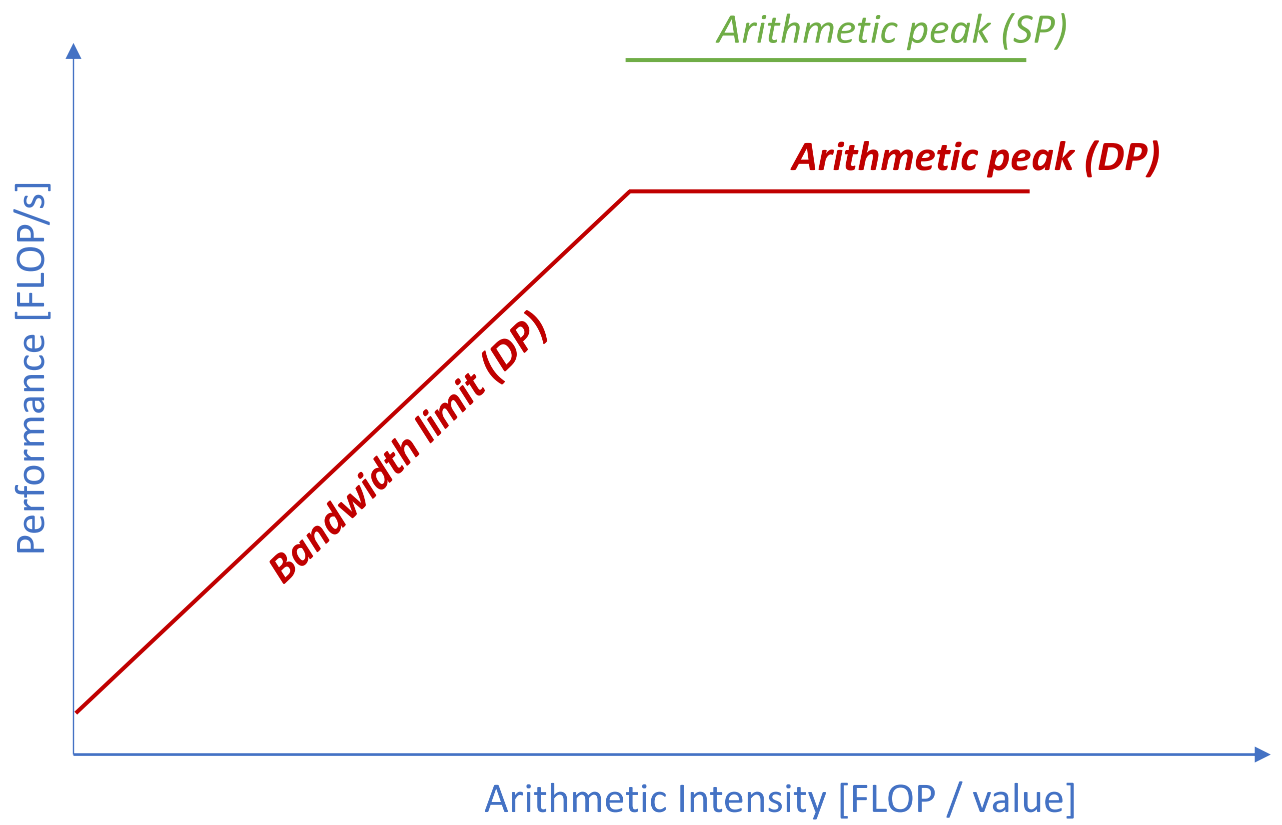

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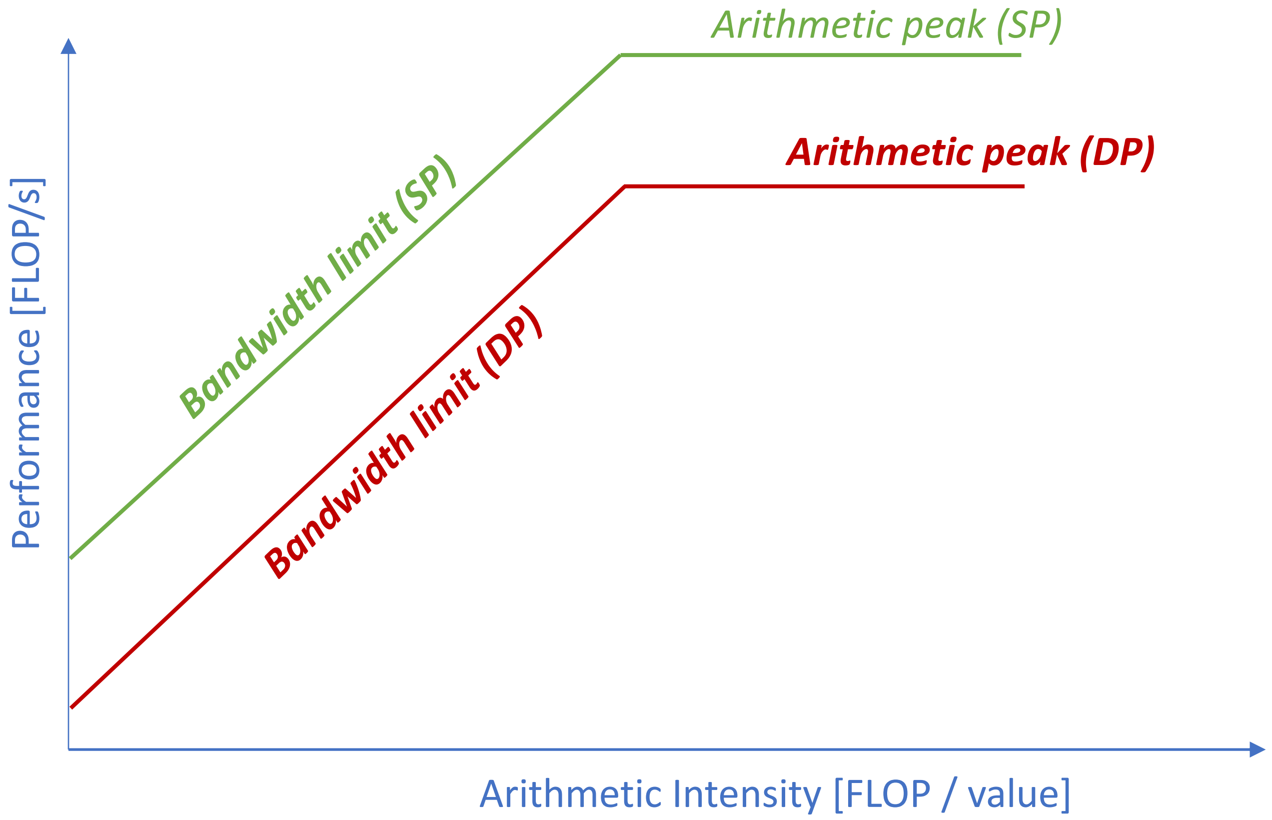

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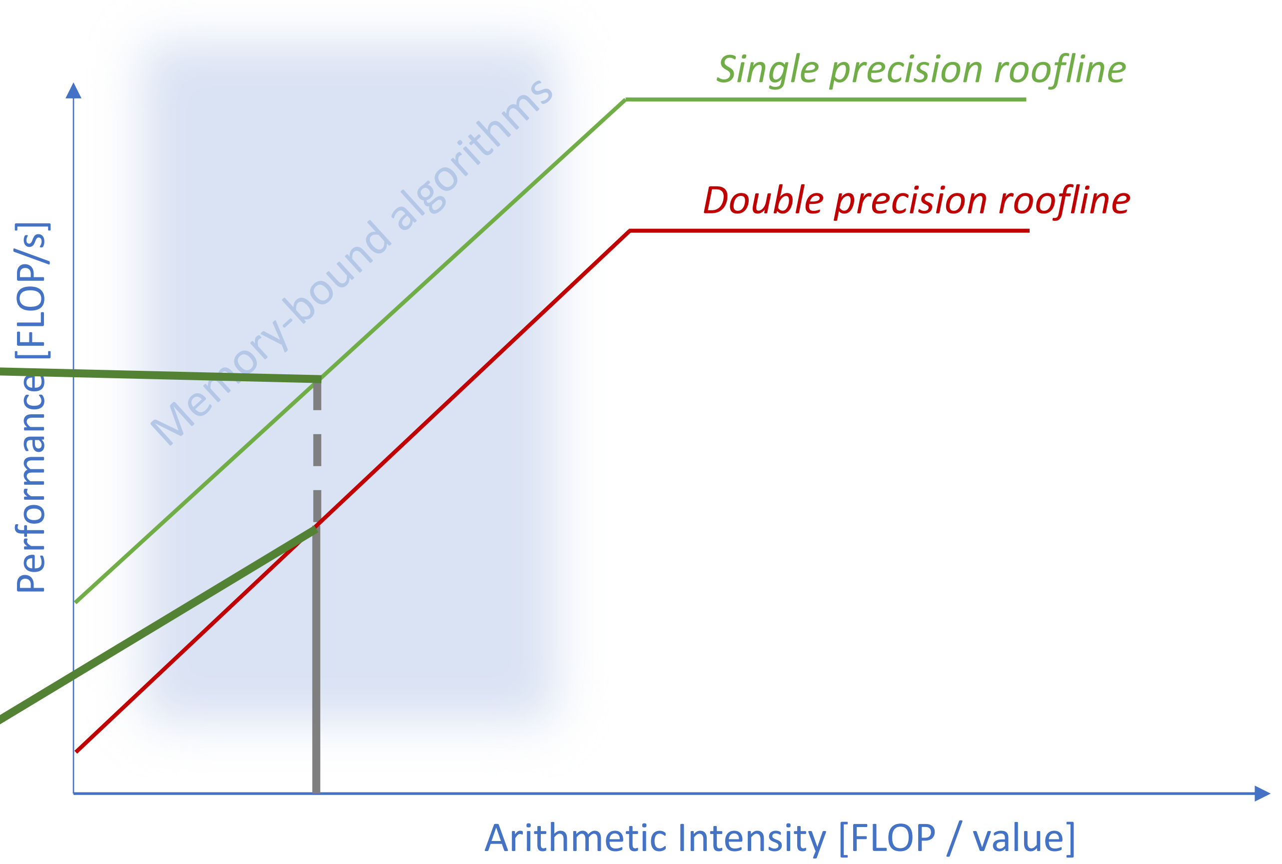
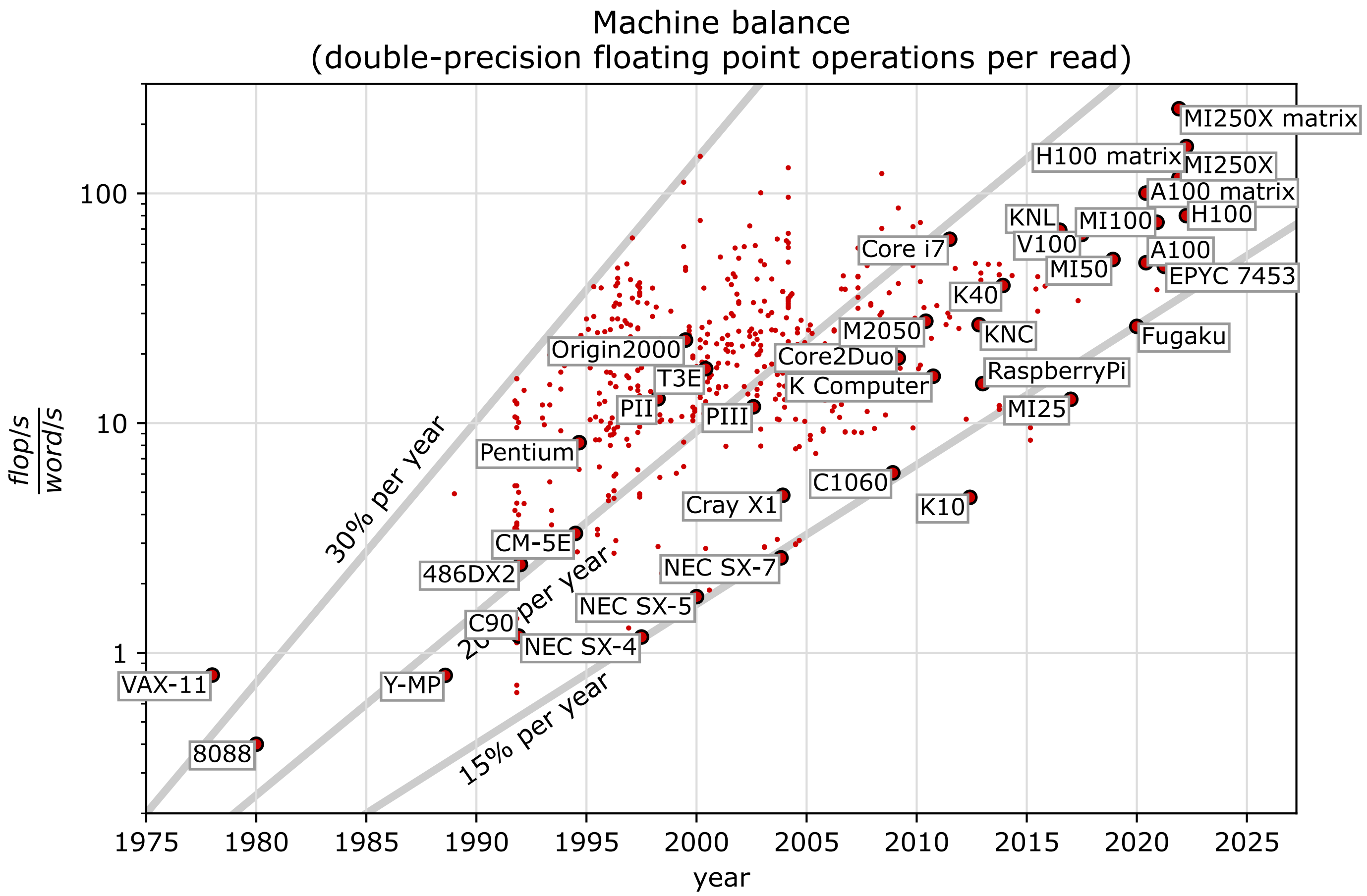


The memory wall



M. Gates

Compute performance grows faster than memory bandwidth.

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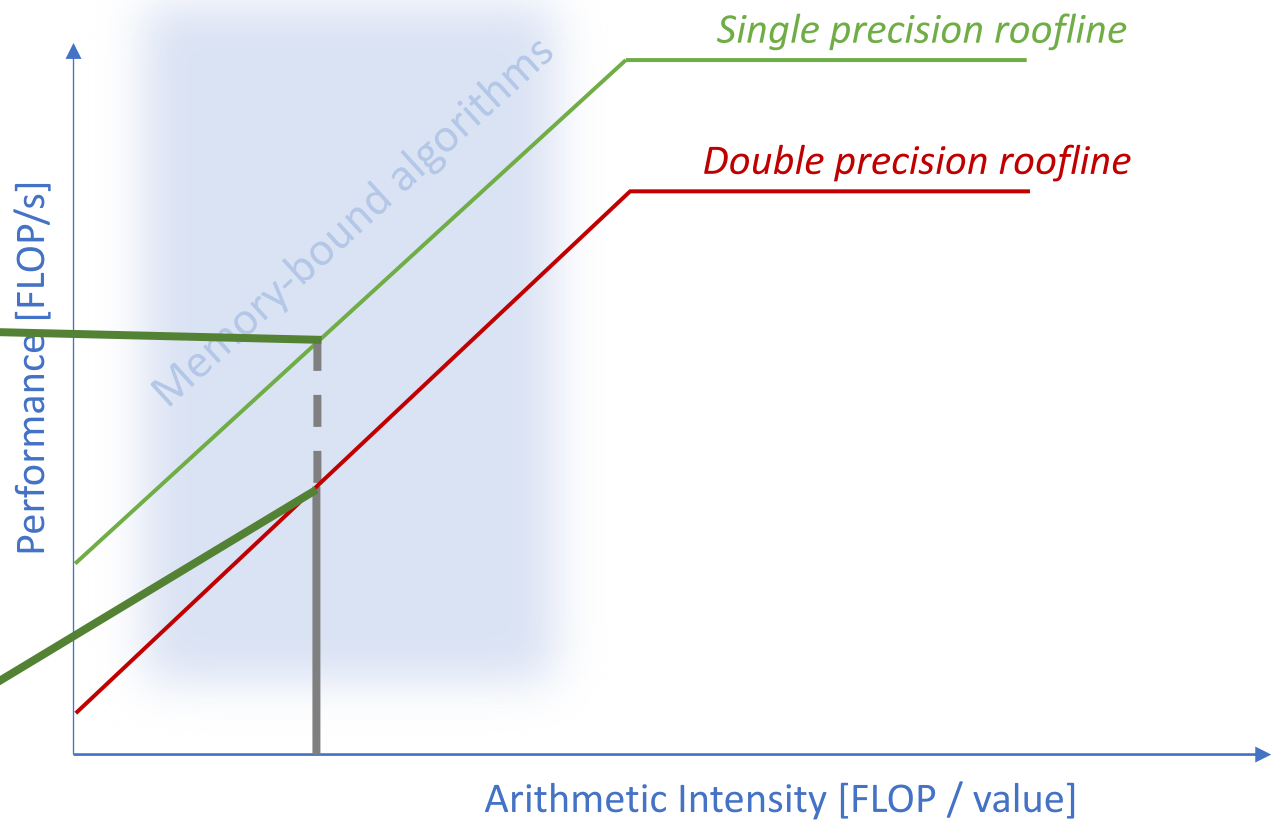
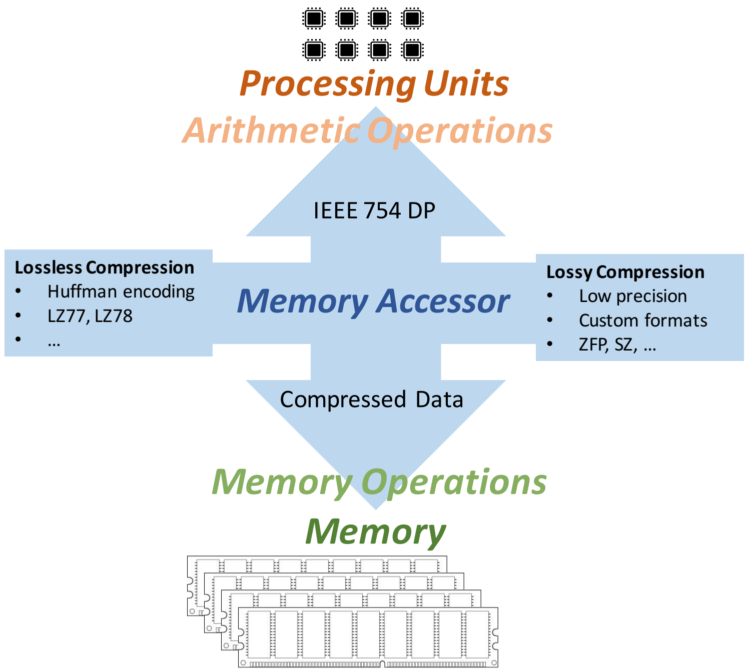
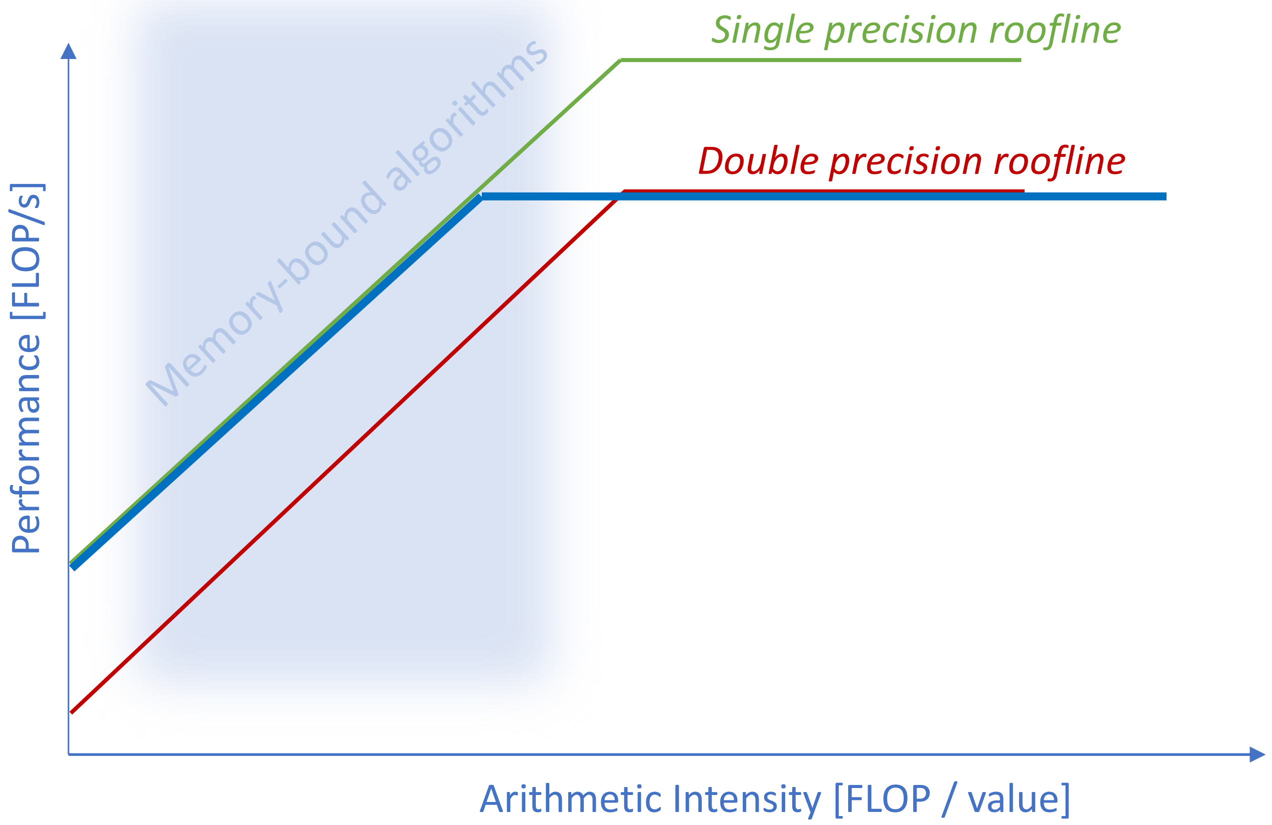


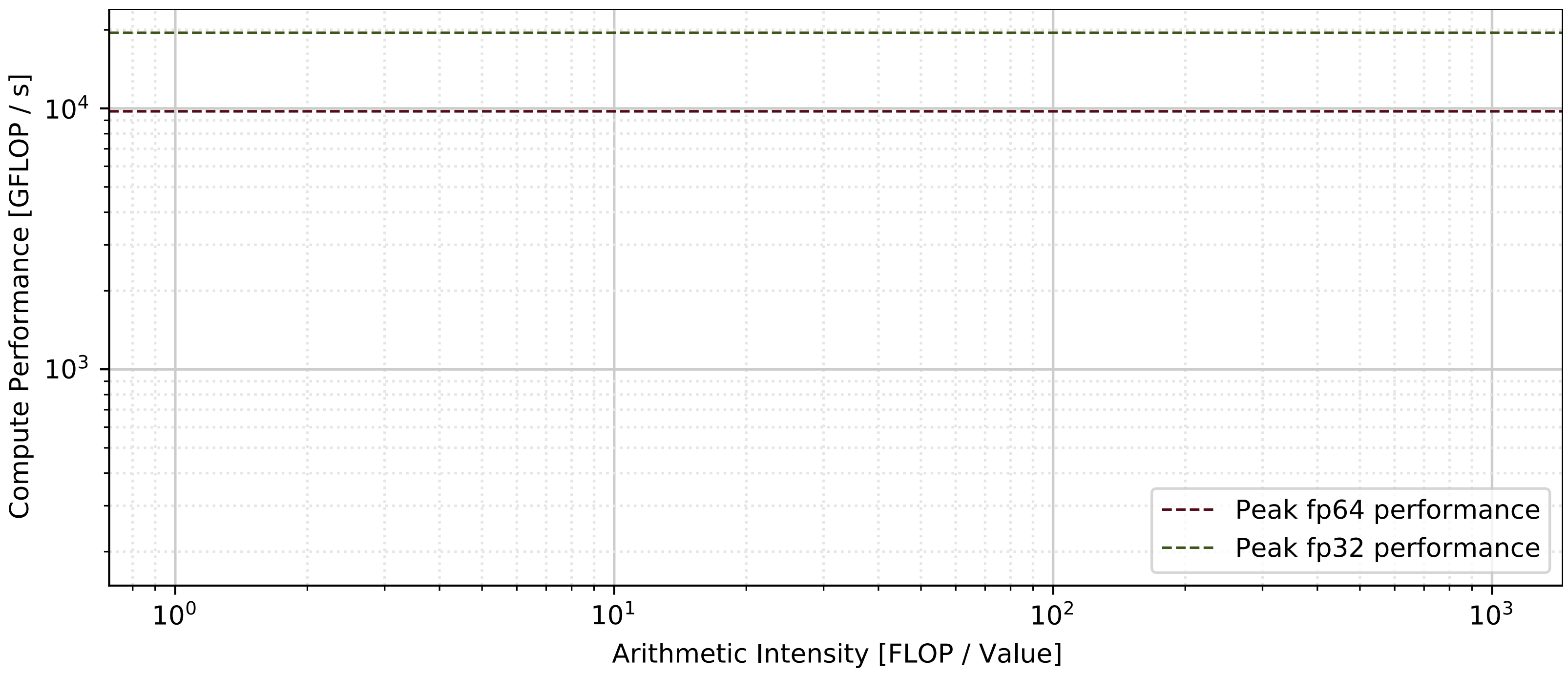
Can we get the best of both worlds?

- For memory-bound algorithms, the arithmetic operations are free, can use high precision formats.
- Data access should be as cheap as possible, use reduced precision.
- In-Register Compression





Memory Accessor for NVIDIA A100 GPU



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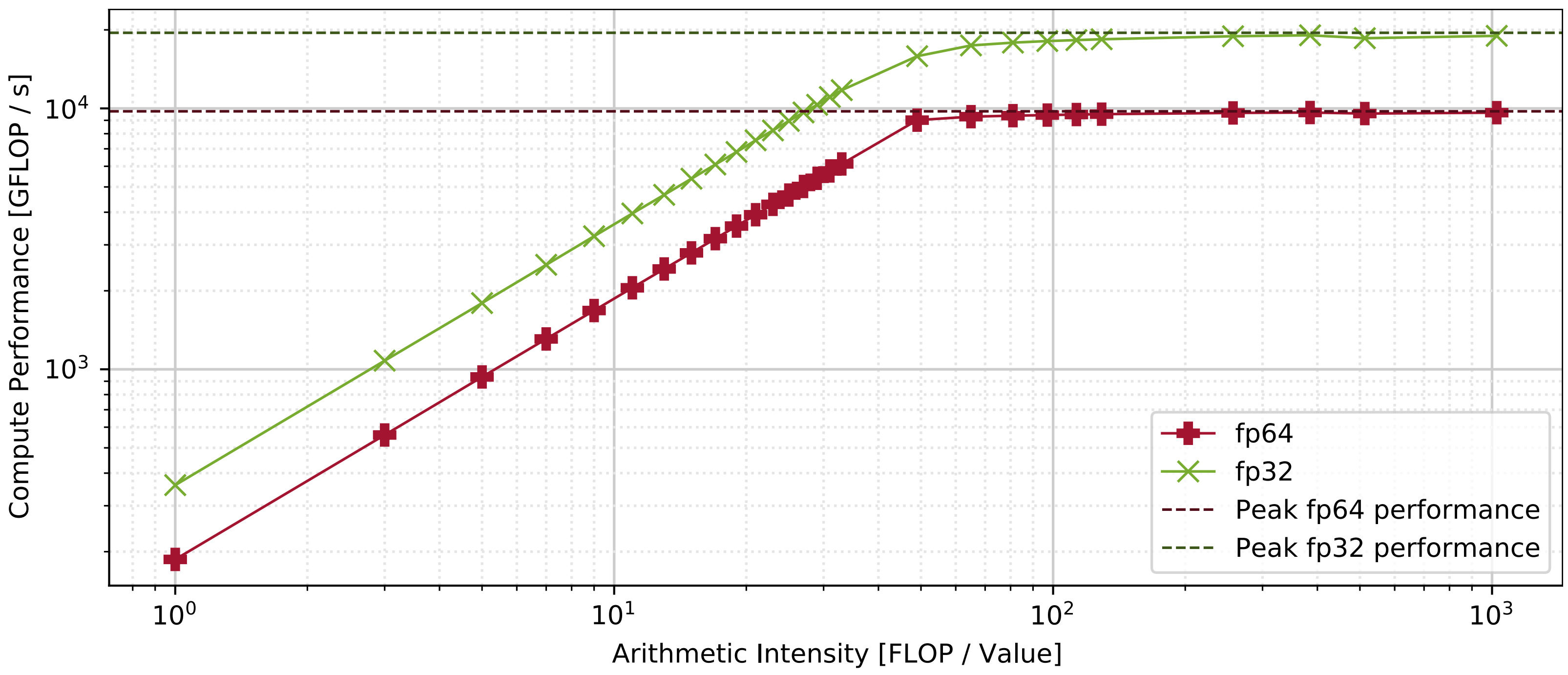

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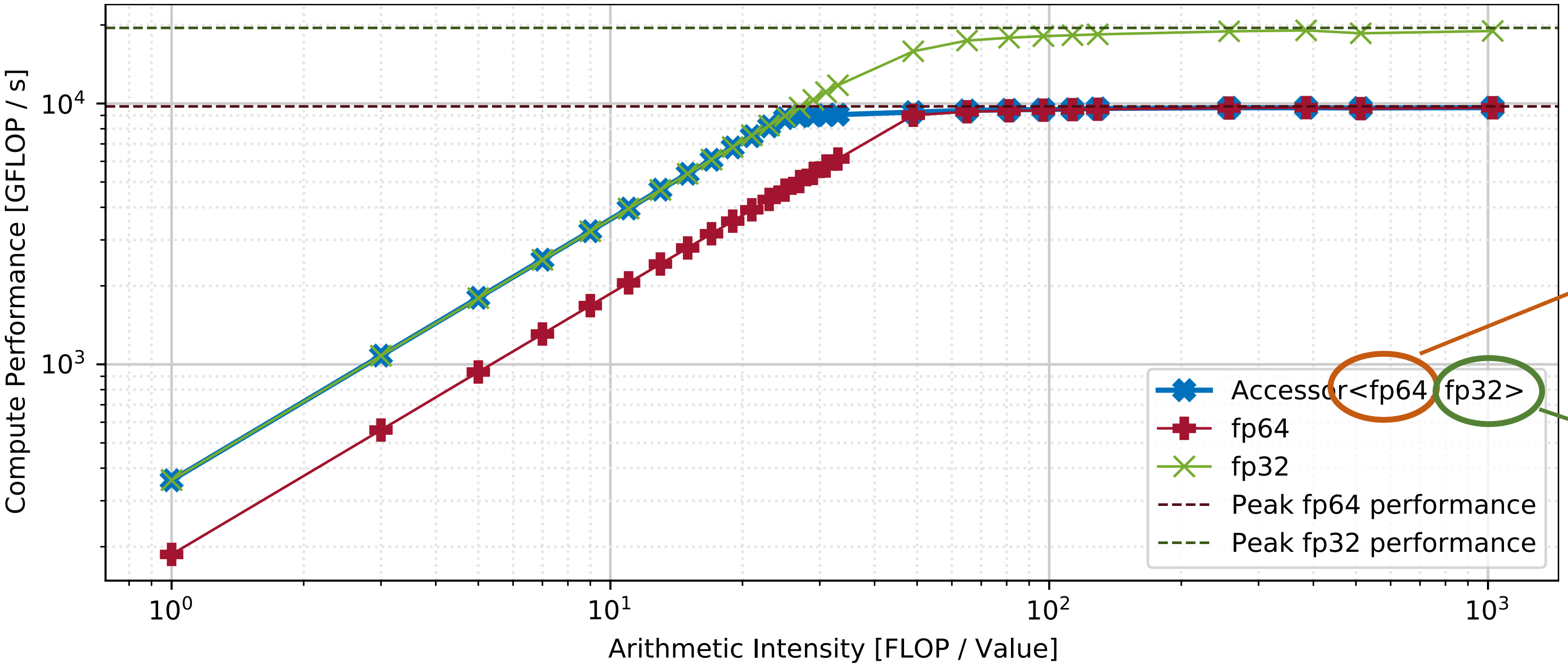

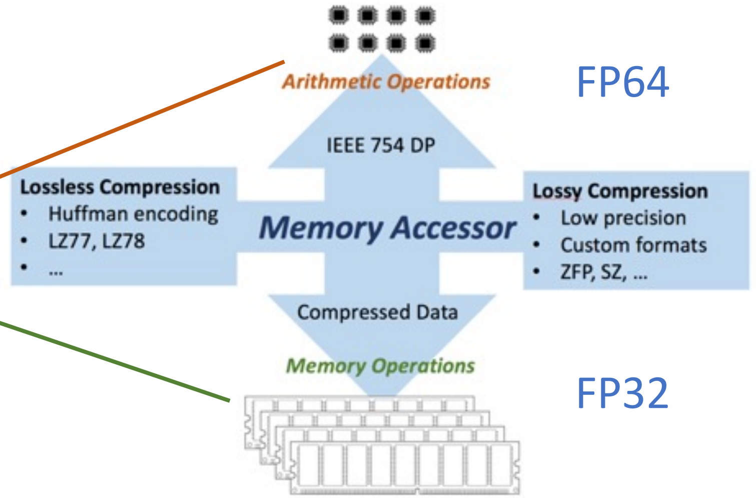

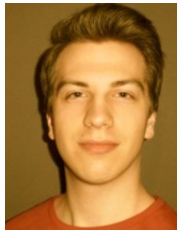

T. Grützmacher

Use the memory accessor to boost performance

- *Start from double precision algorithm*
- *Use memory accessor to store intermediate data in compressed form*
- *Require double precision output accuracy*

Use the memory accessor to boost performance

Goran Flegar

E.S. Quintana-Orti

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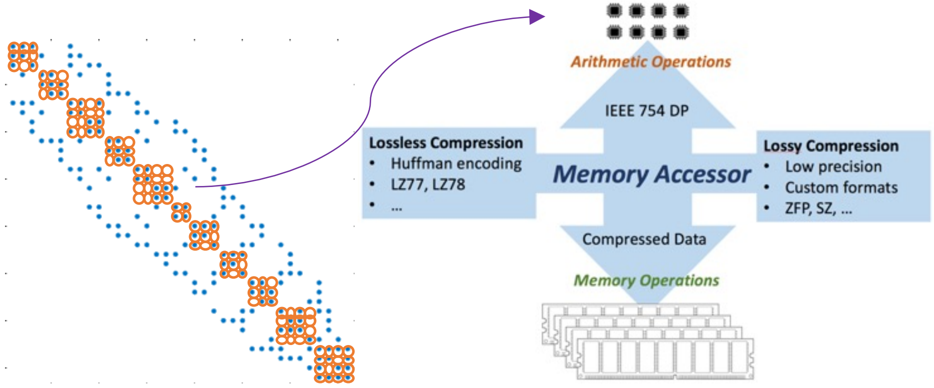


- *Preconditioning*

[1] G Flegar, H Anzt, T Cojean, ES Quintana-Ortí, "Adaptive precision block-Jacobi for high performance preconditioning in the Ginkgo linear algebra software," ACM Transactions on Mathematical Software (TOMS) 47 (2), 1-28.

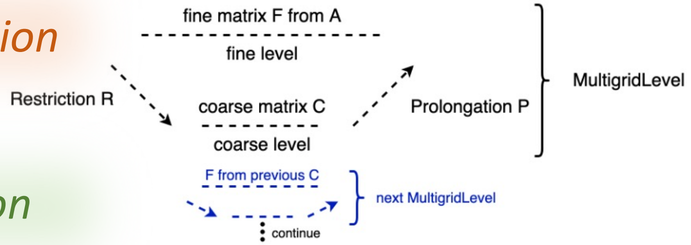
[2] F Göbel, T Grützmacher, T Ribizel, H Anzt, "Mixed precision incomplete and factorized sparse approximate inverse preconditioning on GPUs," European Conference on Parallel Processing, 550-564.

- *Multigrid*

[3] Stephen F. McCormick, Joseph Benzaken, Rasmus Tamstorf "Algebraic error analysis for mixed-precision multigrid solvers", https://arxiv.org/abs/2007.06614
[4] M Tsai, N Beams, H Anzt, "Mixed precision algebraic multigrid on GPUs," PPAM 2022.

high precision

low precision



- *Krylov solvers*

[5] J Aliaga, H Anzt, T Grützmacher, E S Quintana-Ortí, A Tomás, "Compressed basis GMRES on high-performance graphics processing units", IJHPCA 2022

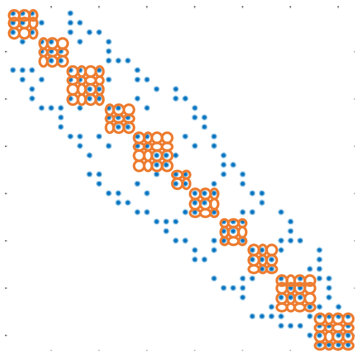
Mike Tsai

Using the memory accessor for mixed precision preconditioning

- **Preconditioning iterative solvers.**
  - Idea: Approximate inverse of system matrix to make the system “easier to solve”: $P^{-1} \approx A^{-1}$

    and solve $Ax = b \quad \Leftrightarrow \quad P^{-1}Ax = P^{-1}b \quad \Leftrightarrow \quad \tilde{A}x = \tilde{b}\,.$

- **Block-Jacobi preconditioner** is based on **block-diagonal scaling**: $P = diag_B(A)$
  - Each block corresponds to one (small) linear system.
    - *Larger* blocks typically improve convergence.
    - *Larger* blocks make block-Jacobi more expensive.



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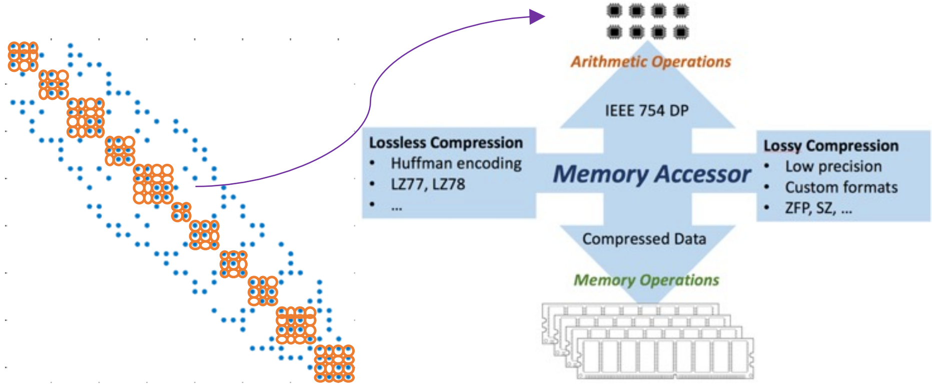
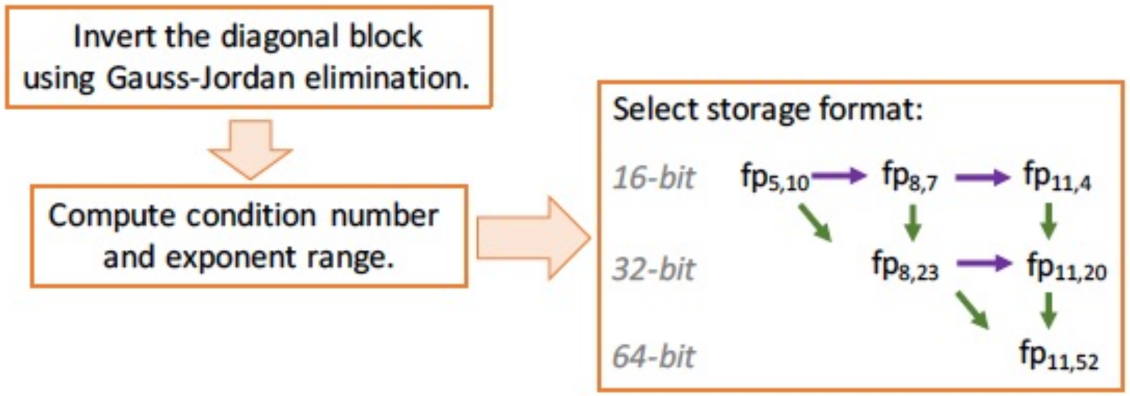


- *Why should we store the preconditioner matrix P^{-1} in full (high) precision?*

- Use the accessor to store the inverted diagonal blocks in lower precision.
  - *Be careful to preserve the regularity of each inverted diagonal block!*

Using the memory accessor for mixed precision preconditioning

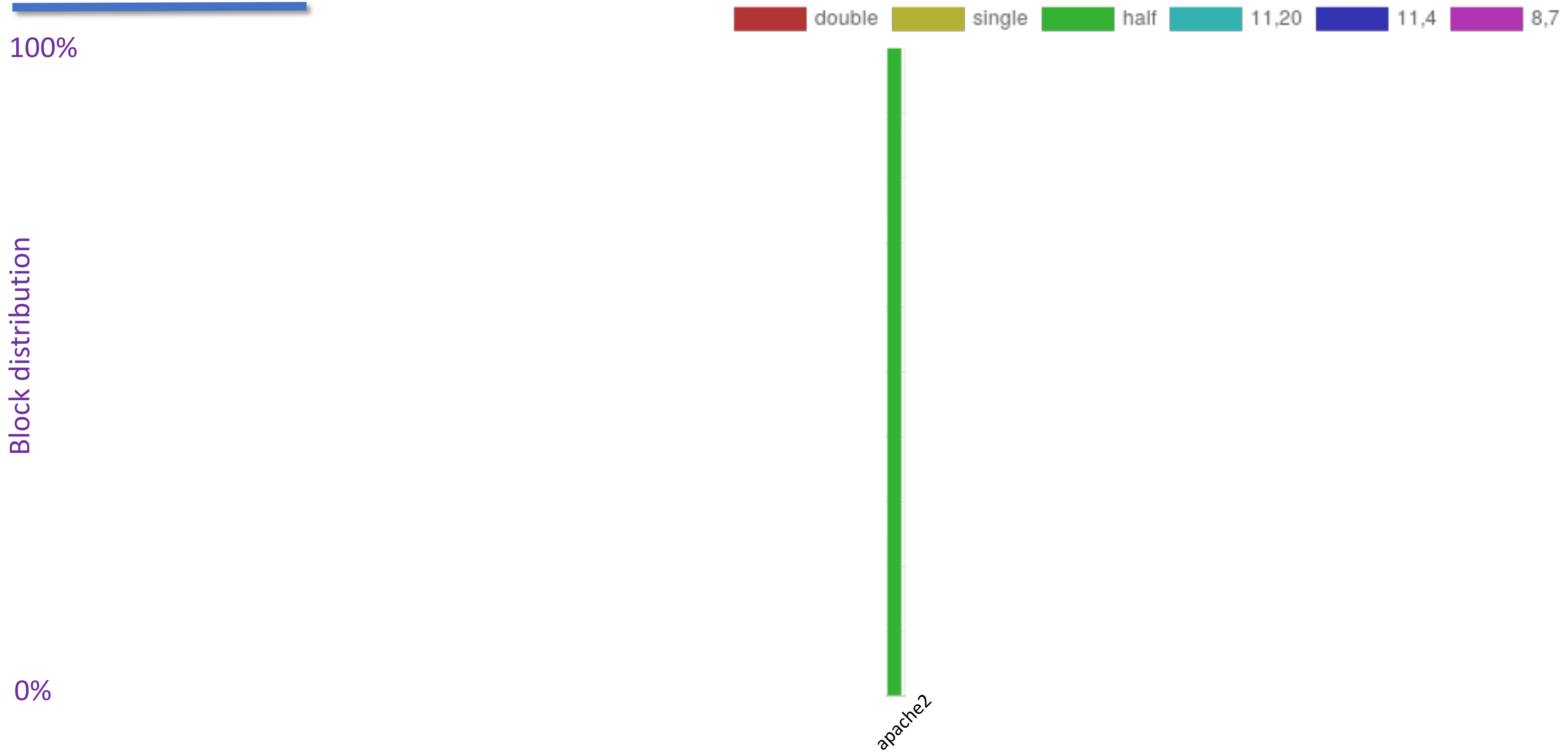
- Choose how much accuracy of the preconditioner should be preserved in the selection of the storage format.
- All computations use double precision, but store blocks in lower precision.



+ **Regularity preserved;**
+ Flexibility in the accuracy;
+ ”Not a low precision preconditioner”
    + Preconditioner is a constant operator;
    + No flexible Krylov solver needed ;

- **Overhead** of the **precision detection**
  (condition number calculation);
- **Overhead** from storing **precision information**
  (need to additionally store/retrieve flag);
- Speedups / preconditioner quality **problem-dependent**;

Using the memory accessor for mixed precision preconditioning



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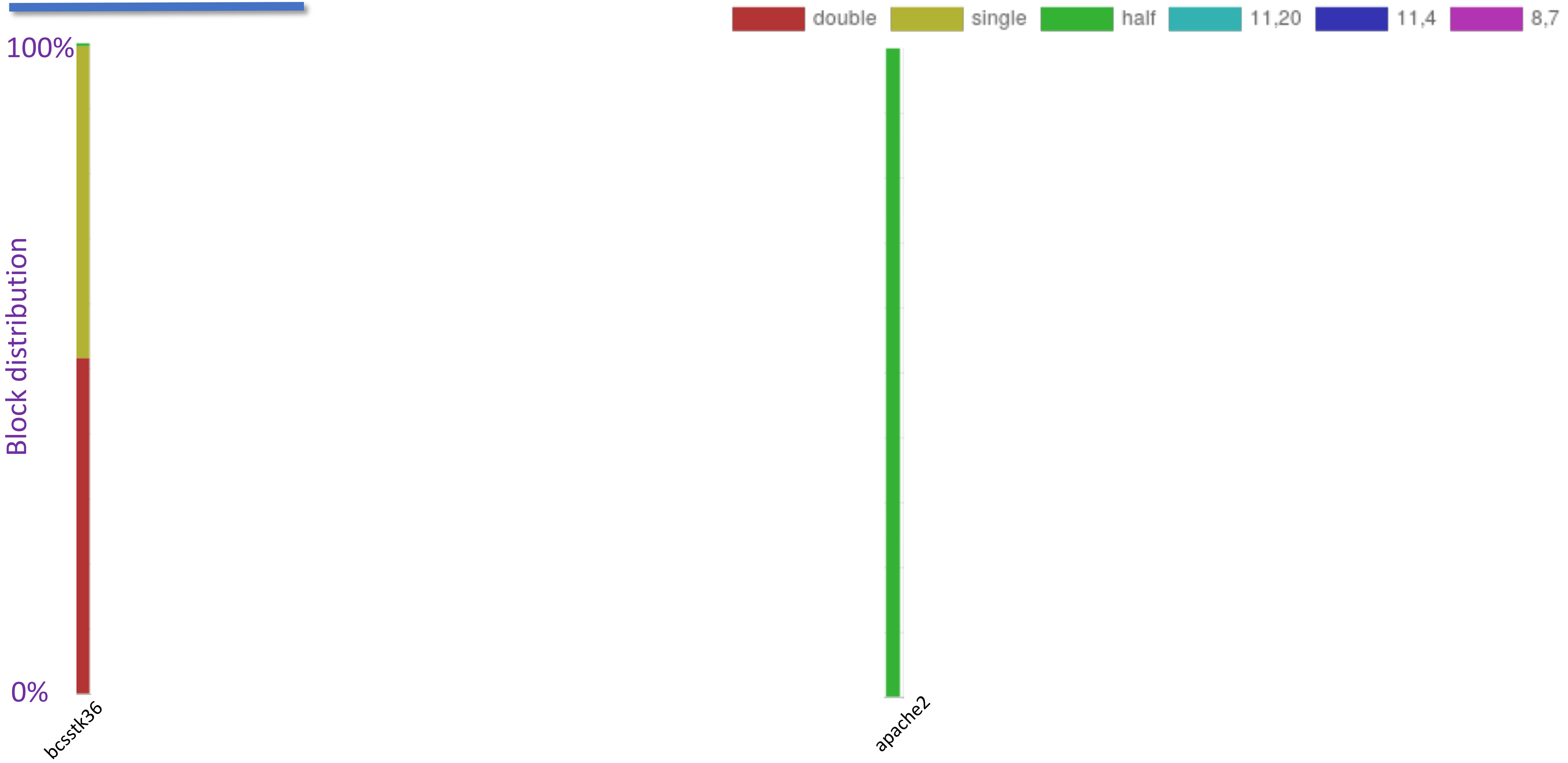

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Linear System Ax=b with cond(A) ≈ 10^7 *(apache2 from SuiteSparse)* **NVIDIA A100 GPU**

Double Precision CG + Double Precision Preconditioner

```
Initial residual norm sqrt(r^T r):
%%MatrixMarket matrix array real general
1 1
1390.67
Final residual norm sqrt(r^T r):
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1 1
3.97985e-06
CG iteration count:     4797
CG execution time [ms]: 2971.18
```

Accuracy improvement ~10^9

Single Precision CG + Single Precision Preconditioner

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Initial residual norm sqrt(r^T r):
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1390.67
Final residual norm sqrt(r^T r):
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1 1
1588.77
CG iteration count:     8887
CG execution time [ms]: 2972.46
```

No improvement

Experiments based on the Ginkgo library https://ginkgo-project.github.io/
ginkgo/examples/adaptiveprecision-blockjacobi/adaptiveprecision-blockjacobi.cpp

Using the memory accessor for mixed precision preconditioning

Ginkgo

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Double Precision CG + **Mixed Precision Preconditioner**

- *Preconditioner remains a constant operator*
- *Attainable accuracy of CG unaffected*
- *Faster because of less data movement*

Experiments based on the Ginkgo library https://ginkgo-project.github.io/
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3.98574e-06
CG iteration count:     4794
CG execution time [ms]: 2568.1
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16% runtime improvement

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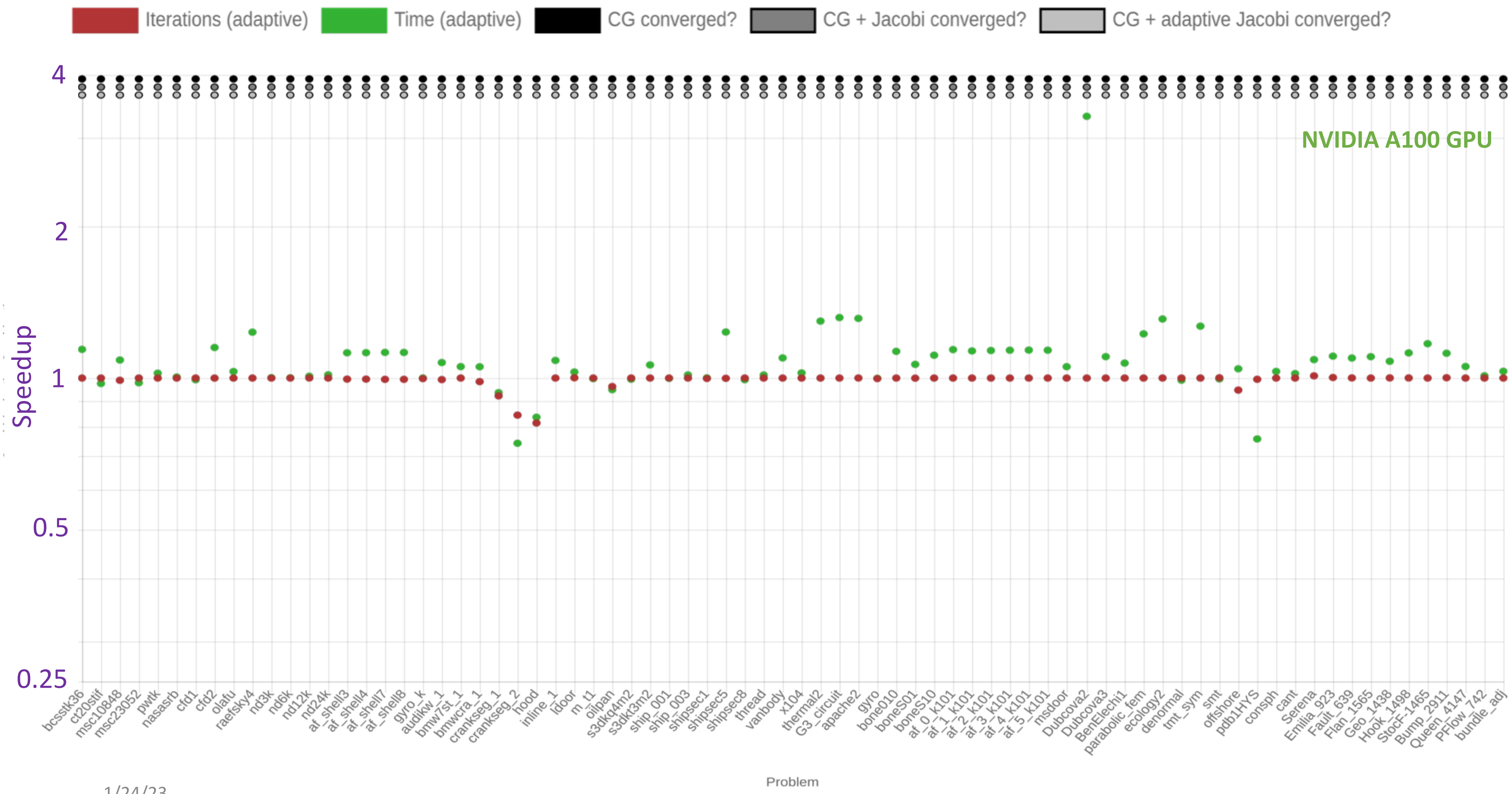
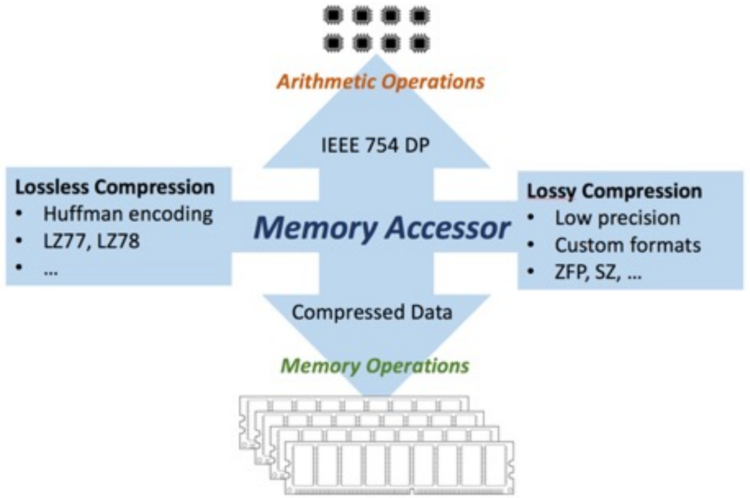


Iterations (adaptive)
Time (adaptive)
CG converged?
CG + Jacobi converged?
CG + adaptive Jacobi converged?
NVIDIA A100 GPU
Speedup
4
2
1
0.5
0.25
bcsstk36 ct20stif msc10848 msc23052 pwtk nasasrb cfd1 cfd2 olafu raefsky4 nd3k nd6k nd12k nd24k af_shell3 af_shell4 af_shell7 af_shell8 gyro_k audikw_1 bmw7st_1 bmwcra_1 crankseg_1 crankseg_2 hood inline_1 ldoor m_t1 oilpan s3dkq4m2 s3dkt3m2 ship_001 ship_003 shipsec1 shipsec5 shipsec8 thread vanbody x104 thermal2 G3_circuit apache2 gyro bone010 boneS01 boneS10 af_0_k101 af_1_k101 af_2_k101 af_3_k101 af_4_k101 af_5_k101 msdoor Dubcova2 Dubcova3 BenElechi1 parabolic_fem ecology2 denormal tmt_sym smt offshore pdb1HYS consph cant Serena Emilia_923 Fault_639 Flan_1565 Geo_1438 Hook_1498 StocF-1465 Bump_2911 Queen_4147 PFlow_742 bundle_adj
Problem

Use the memory accessor to boost performance

Can we use the memory accessor to accelerate the algorithm without changing the final result?

- Yes, if we can do all operations in registers and write the final result in high precision.
- Not in general, if we read/write intermediate date in low precision.
- We need to analyze the error propagation and adapt the algorithms to the application & data.



Possibilities in the context of solving linear systems:

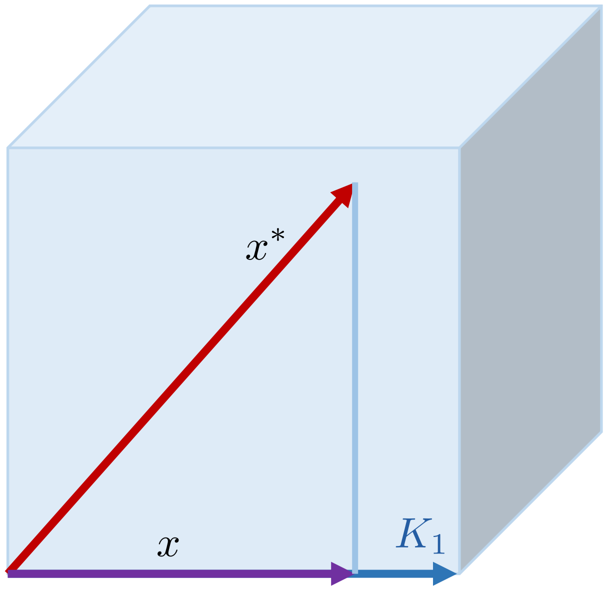
- *Approximate linear operators / Preconditioners / Inner solvers;*
- ***"Self-healing" iterative methods;***

Rethinking Algorithms: Self-Healing Iterative Methods

- **Krylov iterative solvers**
- Krylov methods aim at approximating the solution to a linear problem in a subspace.
- Over the iterations, a nested sequence of Krylov subspaces is generated, adding one basis vector in each iteration.
- Orthonormalization ensures a orthonormal basis is formed (Classical Gram-Schmidt, Modified Gram Schmidt…).

$$K_0 \subset K_1 \subset K_2 \subset \ldots$$

$$K_i(A, r) = span\{b, Ab, A^2b, \ldots A^{i-1}b\}$$



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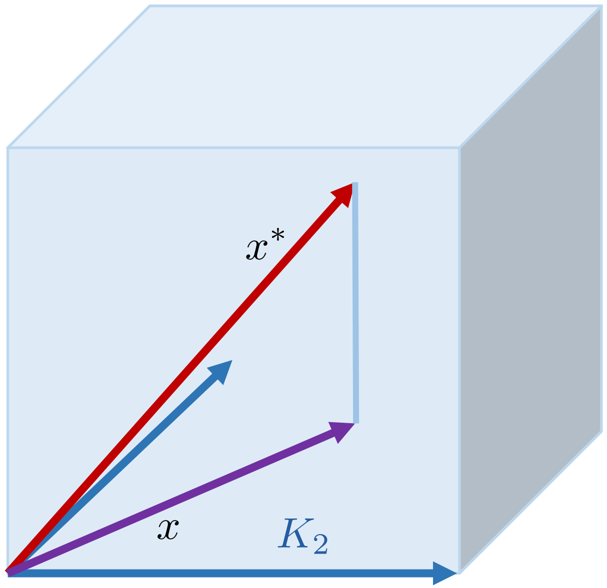

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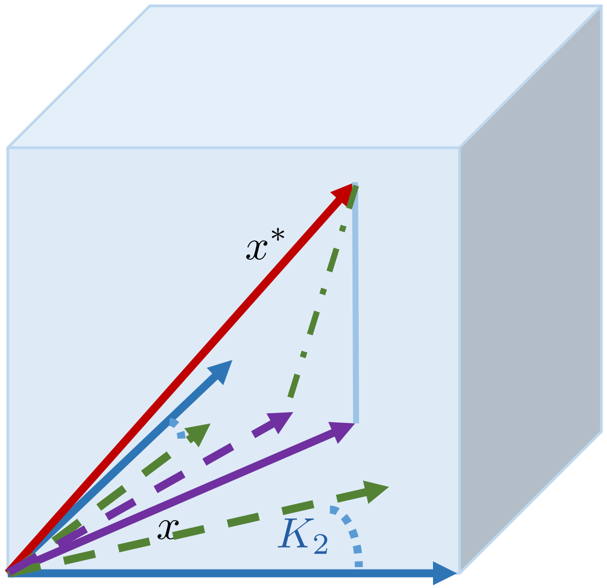
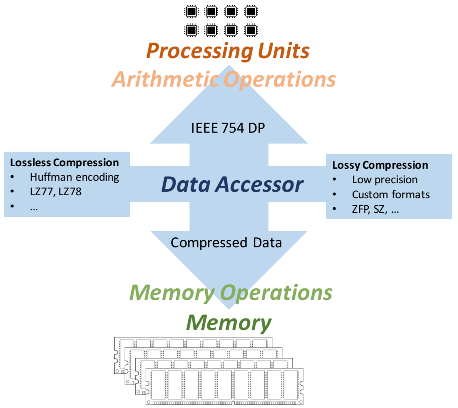


Compressed Basis (CB-) GMRES

- Use double precision in all arithmetic operations;
- Store Krylov basis vectors in lower precision;
  - Search directions are no longer DP-orthogonal;
  - Hessenberg system maps solution to “perturbed” Krylov subspace;
  - Additional iterations may be needed;
  - As long as the loss-of-orthogonality is moderate, we should see moderate convergence degradation;



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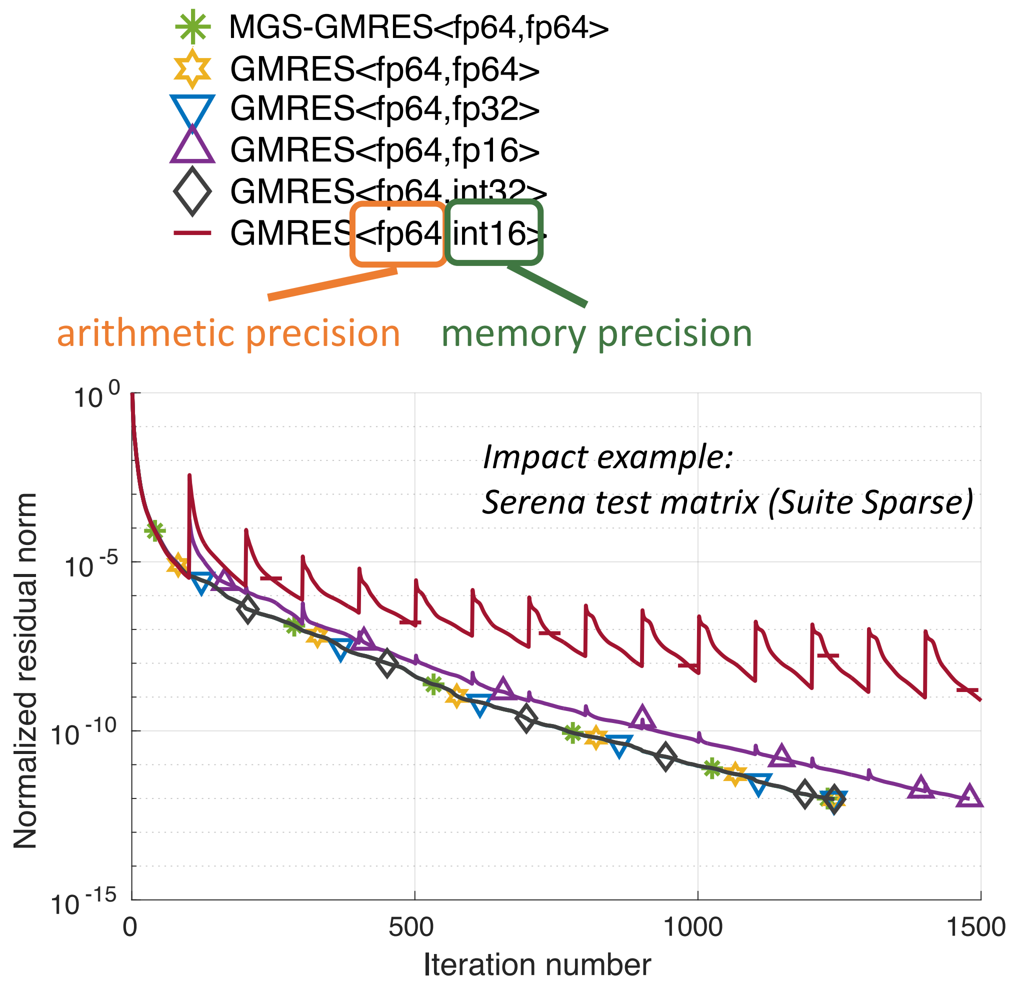

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sqrt(r^T r): 0.00175464
GMRES iteration count: 25000
GMRES execution time: 27376 ms
```
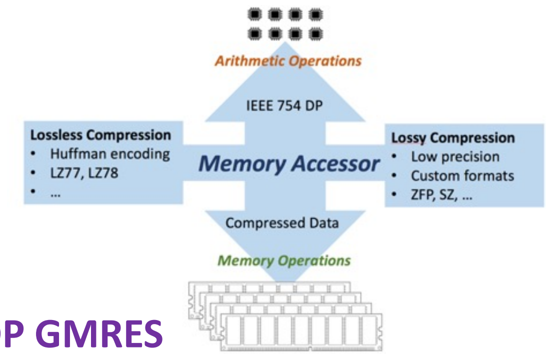

Relative residual ~10^{-7}

```
Compressed Basis GMRES
Initial residual norm
sqrt(r^t r): 9670.36
Final residual norm
sqrt(r^T r): 9.6591e-09
GMRES iteration count: 23271
GMRES execution time: 29369 ms
```

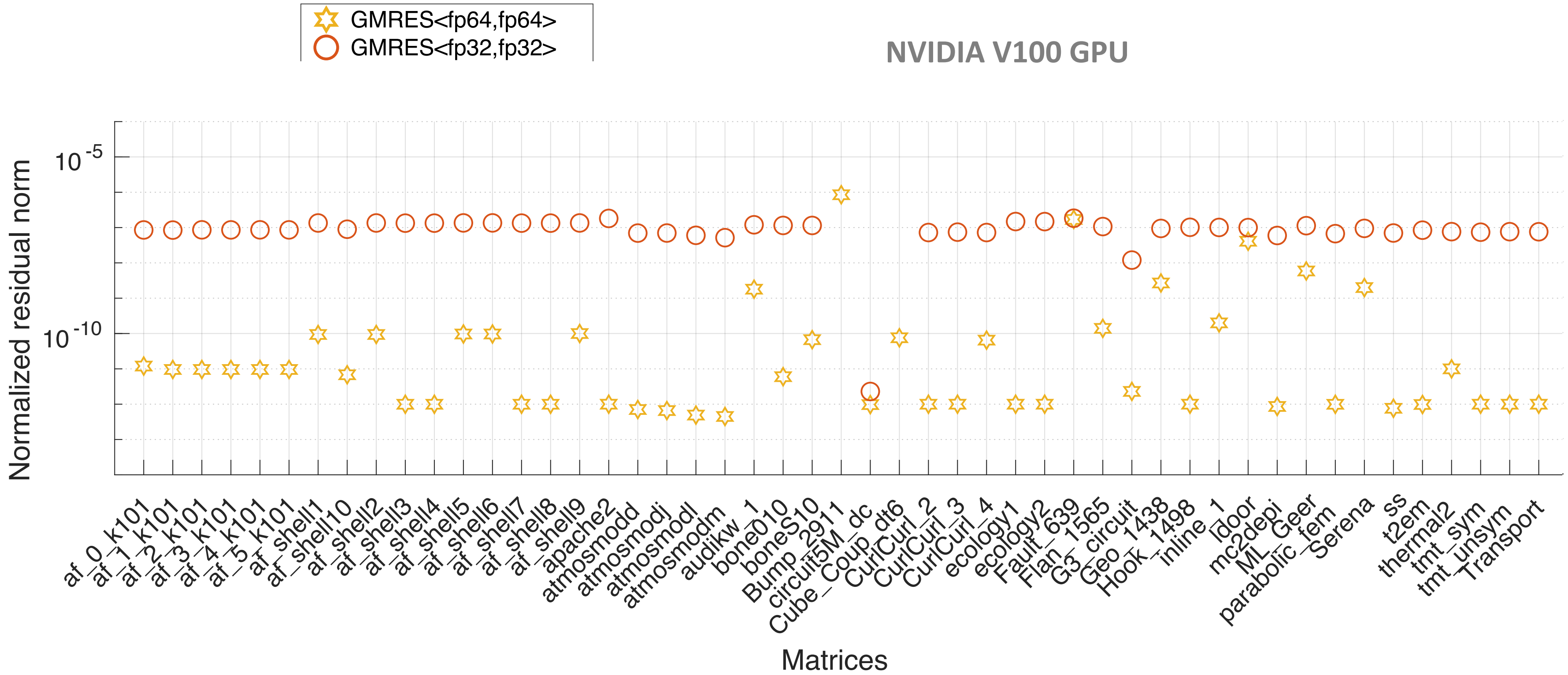
Relative residual ~10^{-12}



Accuracy of DP GMRES

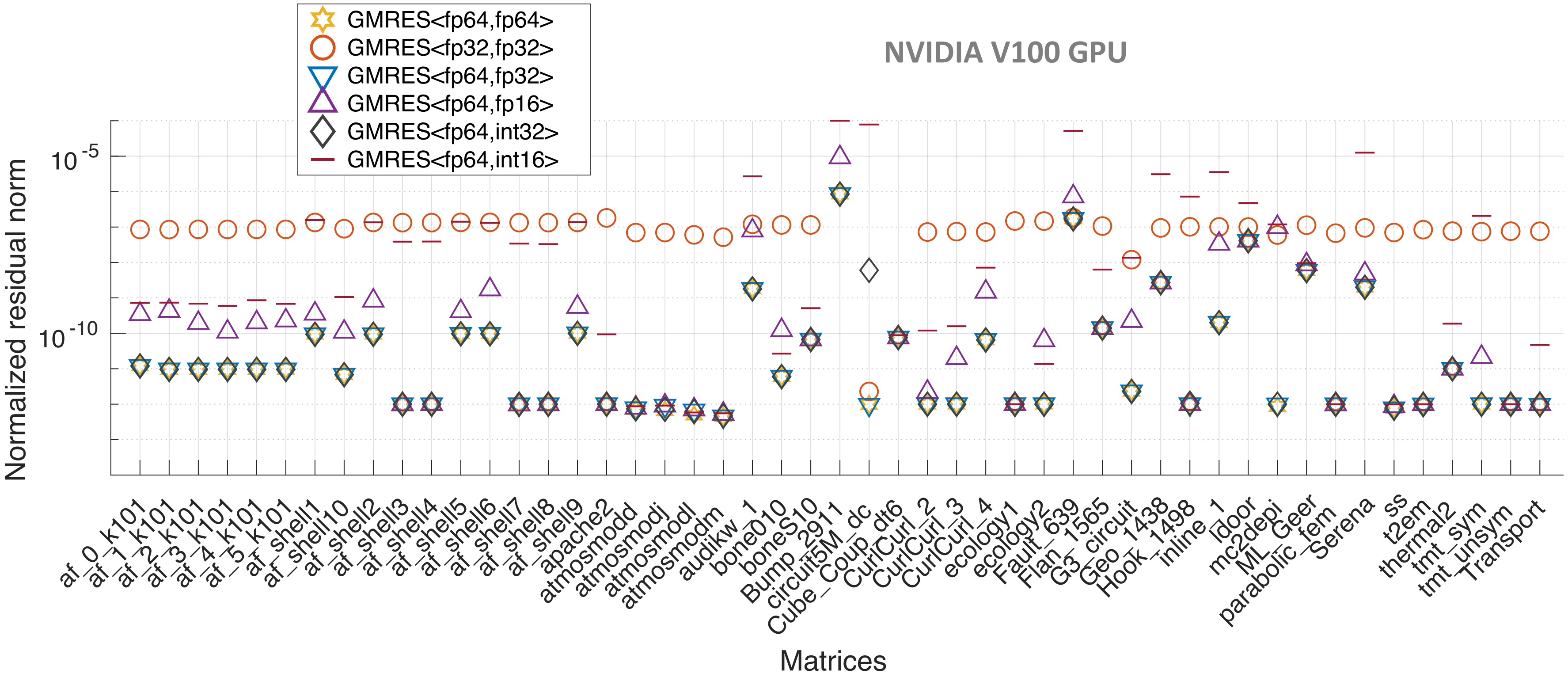
Performance similar to SP GMRES

Compressed Basis GMRES



Compressed Basis GMRES

- CB-GMRES using 32-bit storage preserves DP accuracy (SP-GMRES does not)

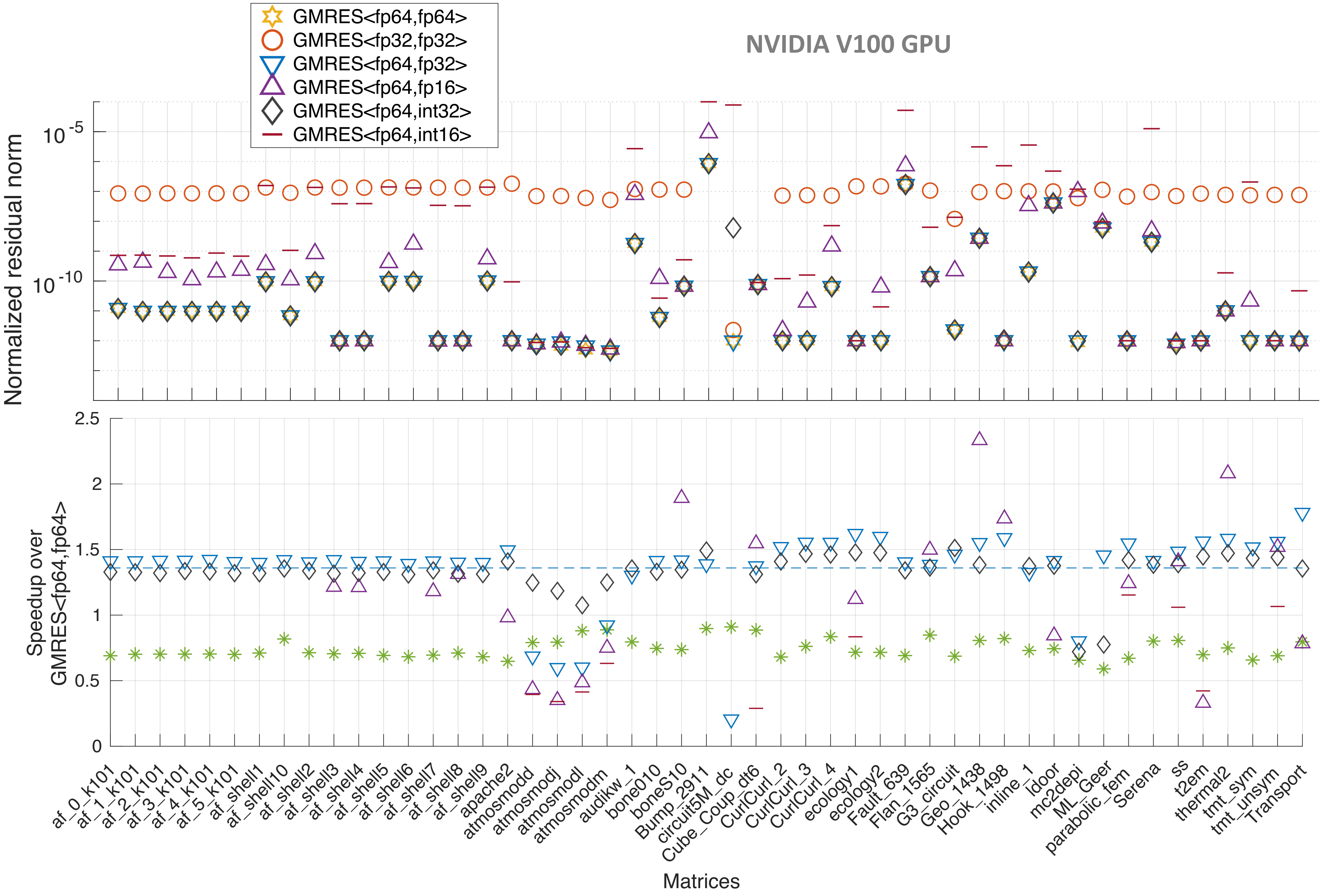


Compressed Basis GMRES

- CB-GMRES using 32-bit storage preserves DP accuracy (SP-GMRES does not)

- Speedups problem-dependent
- Speedup Ø1.4x (for restart 100)
- 16-bit storage mostly inefficient

Ginkgo



Integration into MFEM

Improve current Ginkgo-MFEM integration:

- √ MFEM and Ginkgo operate directly on same data without copies
- √ New `GinkgoExecutor` class automatically matches MFEM `Device` configuration - for CPU, CUDA, or HIP
- √ Ginkgo can use MFEM matrix-free operators in solvers

Add Ginkgo preconditioners to MFEM:

- √ Ginkgo preconditioners can be used with Ginkgo solvers, or used with MFEM solvers
- √ Includes Ginkgo's new ILU-ISAI/IC-ISAI preconditioners, which use the Incomplete Sparse Approximate Inverse to apply the ILU or IC factorization for improved GPU performance

Add new Ginkgo solver to MFEM:

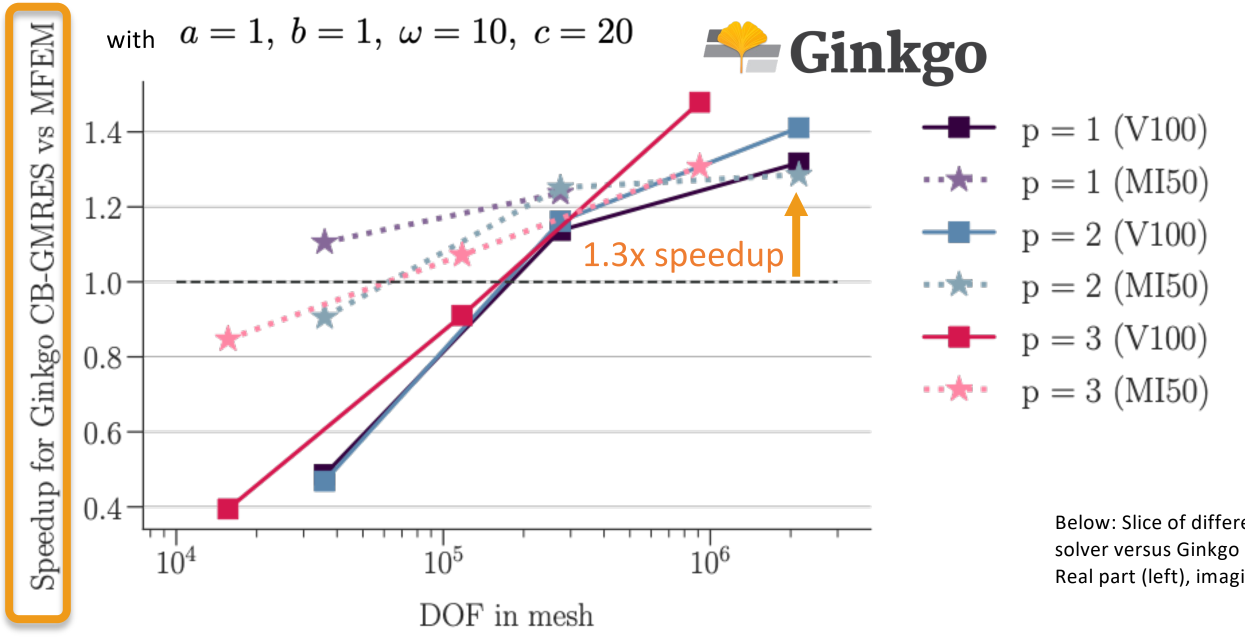
- √ Integration for Ginkgo's Compressed Basis GMRES solver, which uses mixed precision techniques for speedup (see example to right)

Example: Speeding up MFEM's "example 22" on NVIDIA and AMD GPUs

Example 22 solves harmonic oscillation problems, with a forced oscillation imposed at the boundary. For this test, we use variant 1:

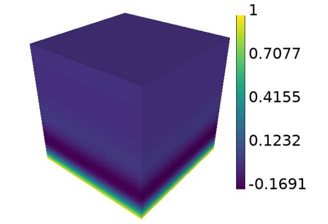
$$-\nabla \cdot (a\nabla u) - \omega^2 bu + i\omega cu = 0$$

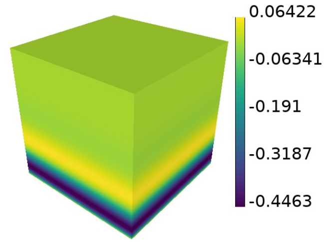
with $a = 1,\ b = 1,\ \omega = 10,\ c = 20$



Speedup of Ginkgo's Compressed Basis-GMRES solver vs MFEM's GMRES solver for three different orders of basis functions (p) for MFEM's example 22. The tests use the "partial assembly" type of MFEM matrix-free operators.

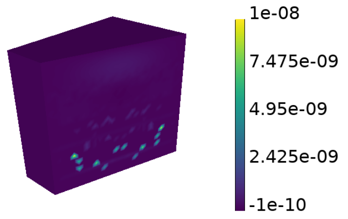
CUDA 10.1/NVIDIA V100 and ROCm 4.0/AMD MI50.
GMRES(50) used for both solvers. CB-GMRES used float/double.

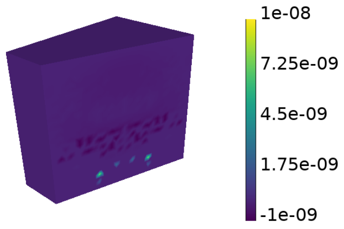




From top: Real part of solution, imaginary part of solution.

Below: Slice of difference in solution output using MFEM solver versus Ginkgo CB-GMRES.
Real part (left), imaginary part (right)





Natalie Beams (Univ. of Tennessee)

Using the memory accessor to boost accuracy

Instead of improving the performance of memory-bound high precision algorithms, the memory accessor can be used to increase the accuracy of memory-bound low precision algorithms – at no cost.

Design

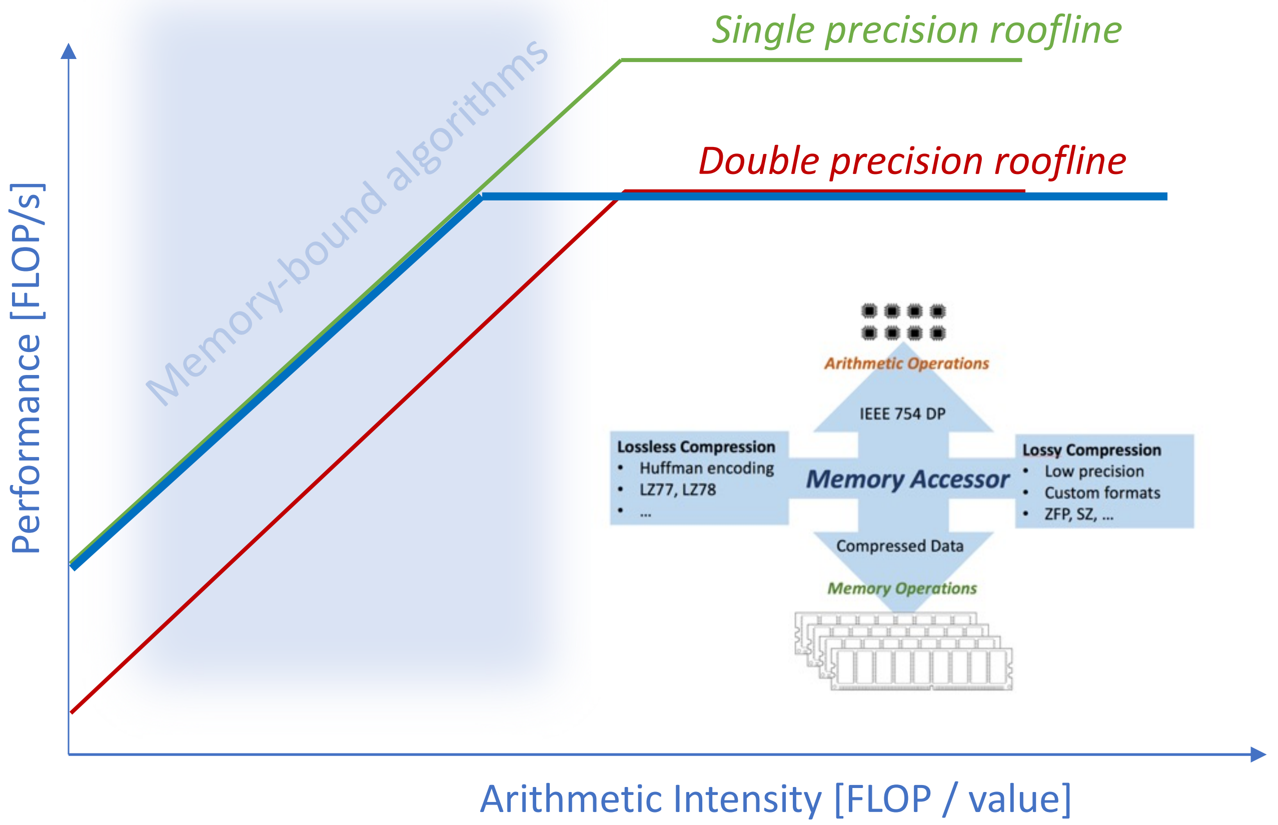
- Memory access in low precision (e.g. fp32);
- Computations in high precision (e.g. fp64);

Characteristics

- Performance of low precision BLAS;
- Higher accuracy than low precision BLAS;

Usage

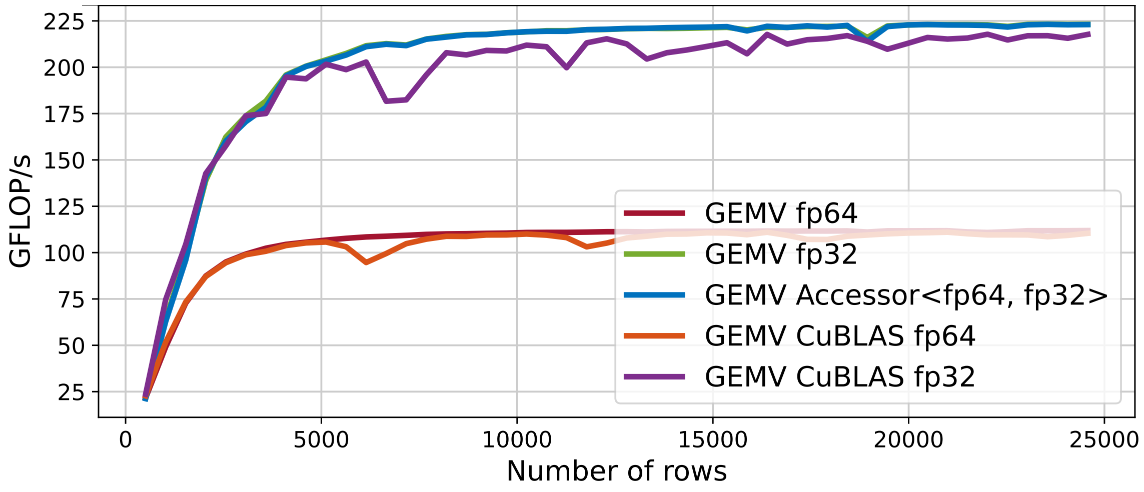
1. Can replace low precision BLAS to increase accuracy;
2. Can replace high precision BLAS if information loss is acceptable;
   (without having to deal with explicit mixed precision usage)



Accessor-BLAS: Replacing LP BLAS to improve accuracy

GEMV

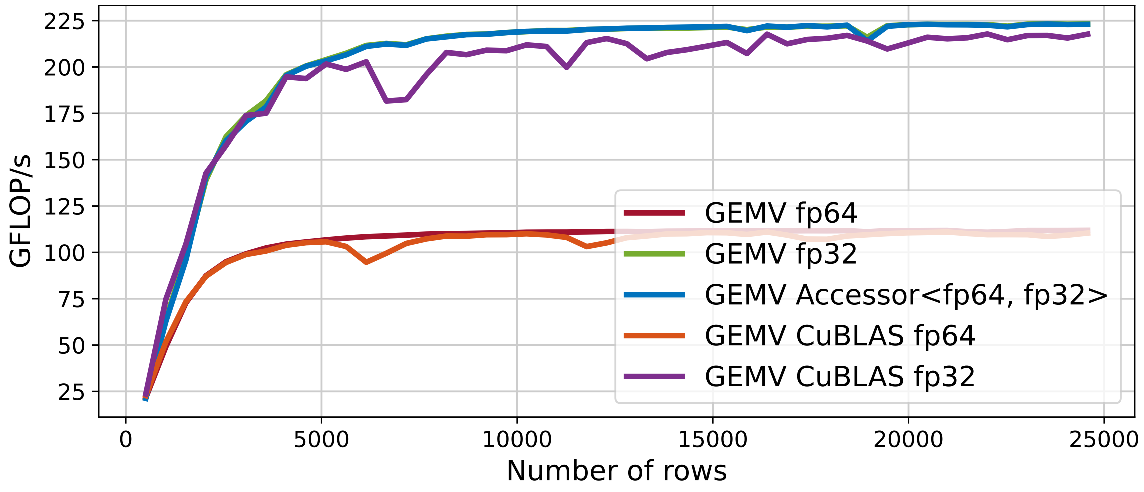
NVIDIA V100 GPU (Summit)

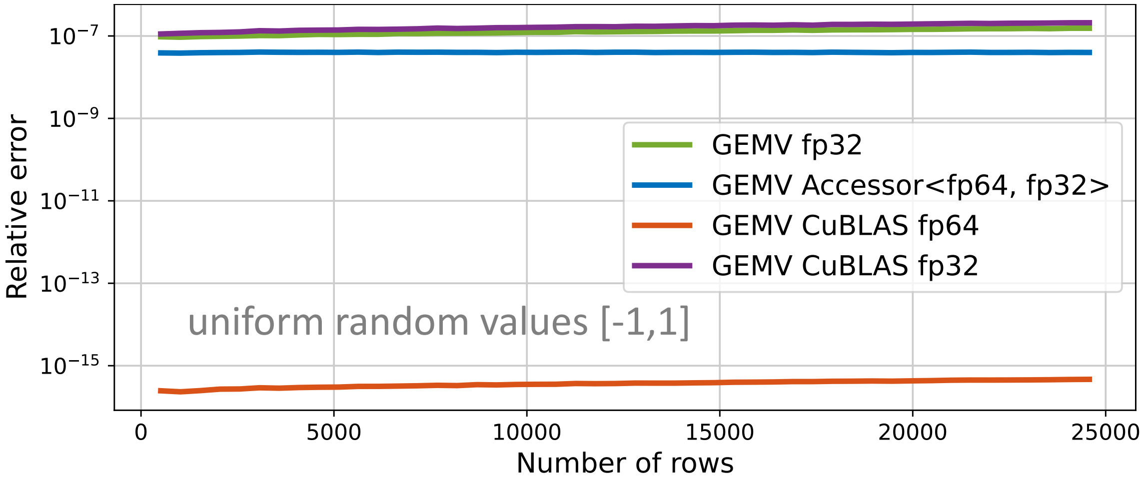


Accessor-BLAS: Replacing LP BLAS to improve accuracy

GEMV

NVIDIA V100 GPU (Summit)





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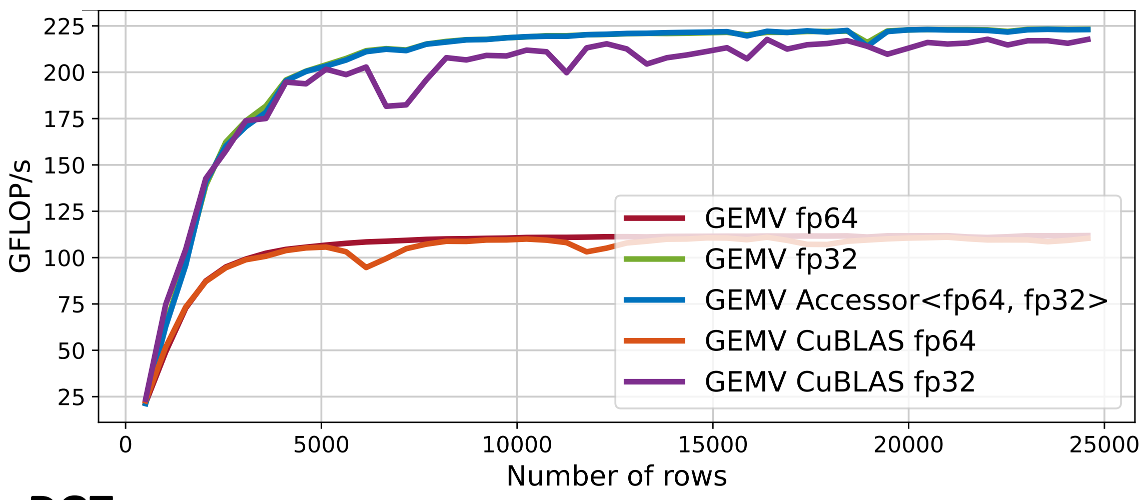
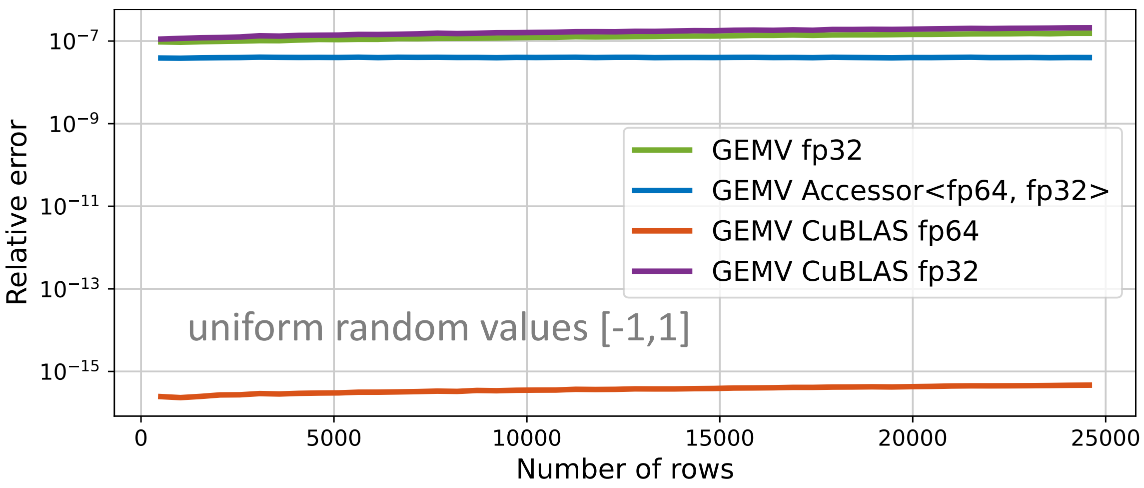
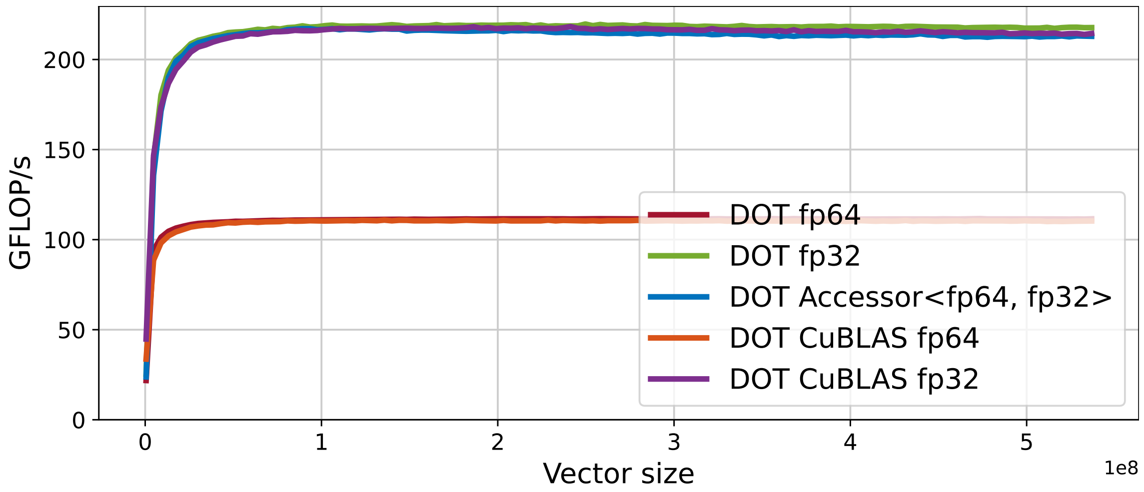
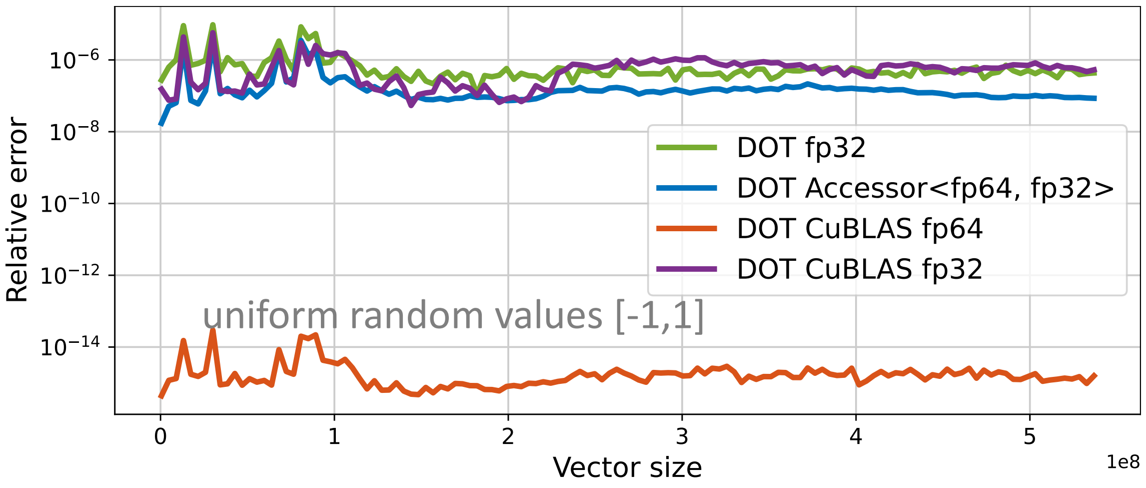




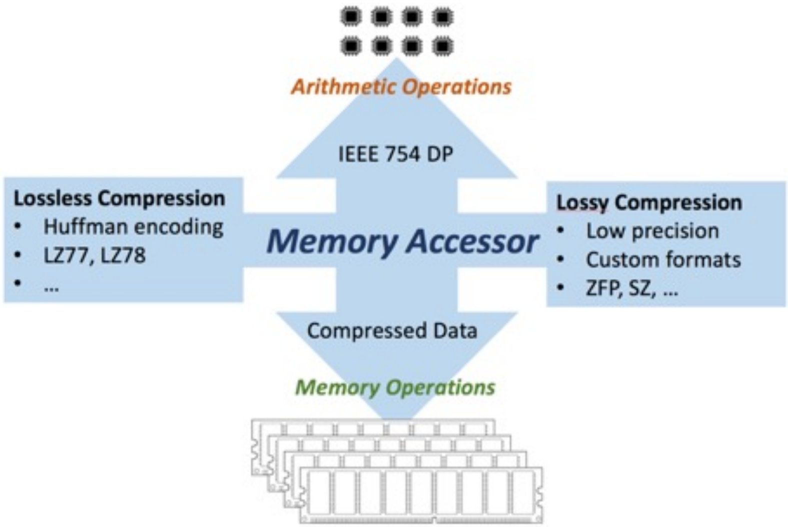
DOT





Summary and resources

- *For memory-bound algorithms, mixed precision can boost performance through reduced data movement.*
- **Memory accessor** allows to **compress data in main memory** but **do all arithmetic in high (double) precision**.
- **Approximate operators** (preconditioners, lower multigrid levels) and **self-healing iterative methods** can accept/compensate **information loss**.
- **Memory-bound low precision algorithms can increase accuracy at no cost.**



Ginkgo

Accessor-based GEMV and DOT available as open-source code:
https://github.com/ginkgo-project/accessor-BLAS

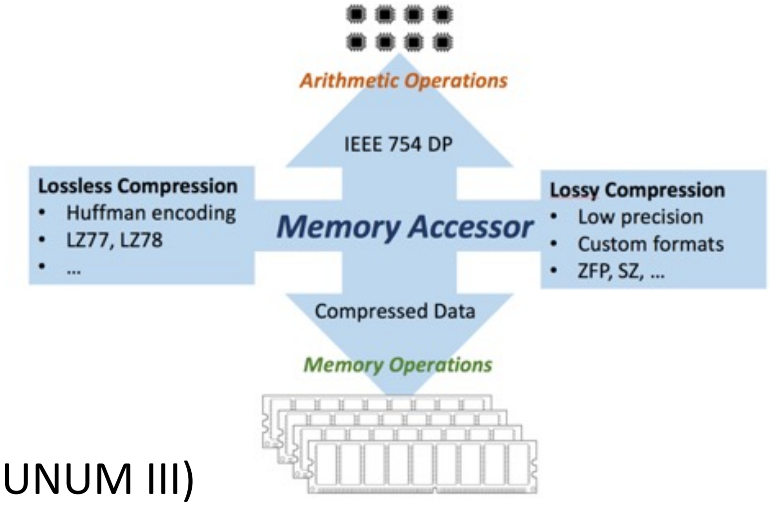
Mixed Precision block-Jacobi preconditioning:
https://github.com/ginkgo-project/ginkgo/tree/develop/examples/adaptiveprecision-blockjacobi

Mixed Precision Iterative Refinement:
https://github.com/ginkgo-project/ginkgo/tree/develop/examples/mixed-precision-ir

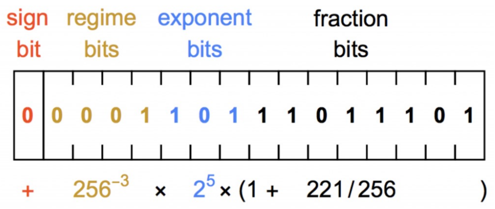
Compressed-Basis GMRES:
https://github.com/ginkgo-project/ginkgo/tree/develop/examples/cb-gmres

Let's try harder

- **IEEE 754 fp64 in arithmetic operations**
- **More sophisticated in-register compression**
  - Custom formats
  - Compression techniques (SZ, ZFP)
- **Store data in compressed format**



POSIT (UNUM III)



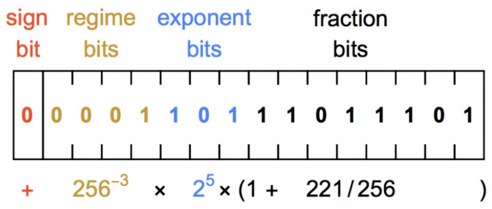
John L. Gustafson

Dynamic Range

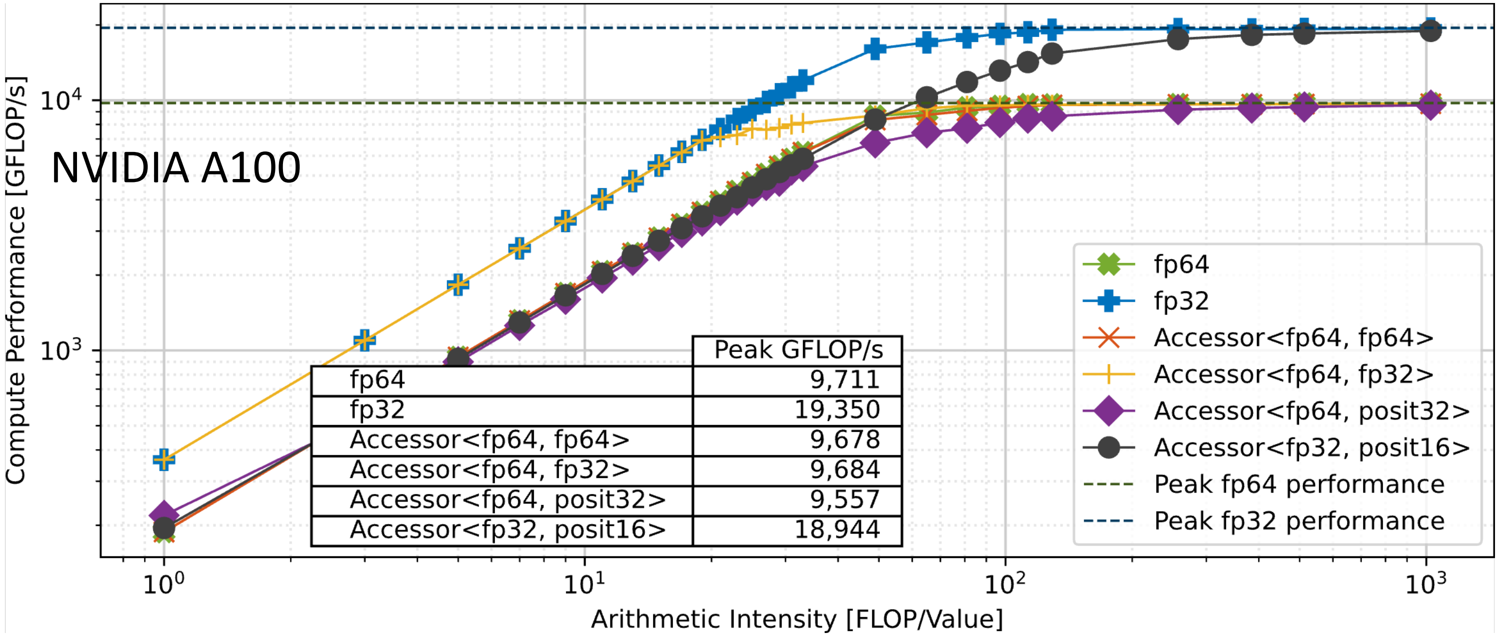
Size [bits]	IEEE exponent size [bits]	Approx. IEEE dynamic range	Approx. Posit dynamic range	Posit exp. bits
16	5	$[6 \cdot 10^{-8}, 7 \cdot 10^{4}]$	$[1 \cdot 10^{-17}, 7 \cdot 10^{16}]$	2
32	8	$[1 \cdot 10^{-45}, 3 \cdot 10^{38}]$	$[8 \cdot 10^{-37}, 1 \cdot 10^{36}]$	2
64	11	$[5 \cdot 10^{-324}, 2 \cdot 10^{308}]$	$[2 \cdot 10^{-75}, 5 \cdot 10^{75}]$	2

	IEEE	Posit
Special values	• IEEE defines ± 0, $\pm\infty$ and NaN (quiet and signaling) as special values • A lot of NaN representations (fp32 has $2^{24} - 1 \approx 10^{7}$ different NaNs)	• Posit only has 2: 0 and NaR • NaR (Not a Real) is used as an error-value (like NaN and $\pm\infty$)
Gradual over- and underflow	• IEEE supports gradual underflow with subnormal numbers (fraction has an implicit 0.) • No support for gradual overflow	• Posit supports both gradual over- and underflow through the regime • The farther away from 1.0, the fewer fraction bits

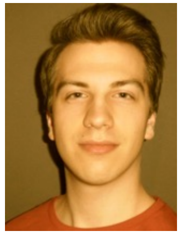
Using POSIT as memory format



Unum Type III
John L. Gustafson

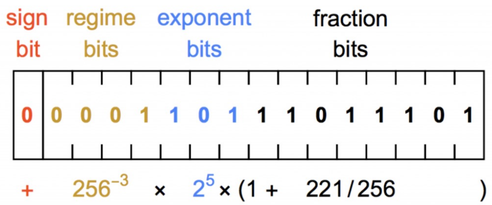


	Peak GFLOP/s
fp64	9,711
fp32	19,350
Accessor<fp64, fp64>	9,678
Accessor<fp64, fp32>	9,684
Accessor<fp64, posit32>	9,557
Accessor<fp32, posit16>	18,944



T. Grützmacher

Using POSIT as memory format



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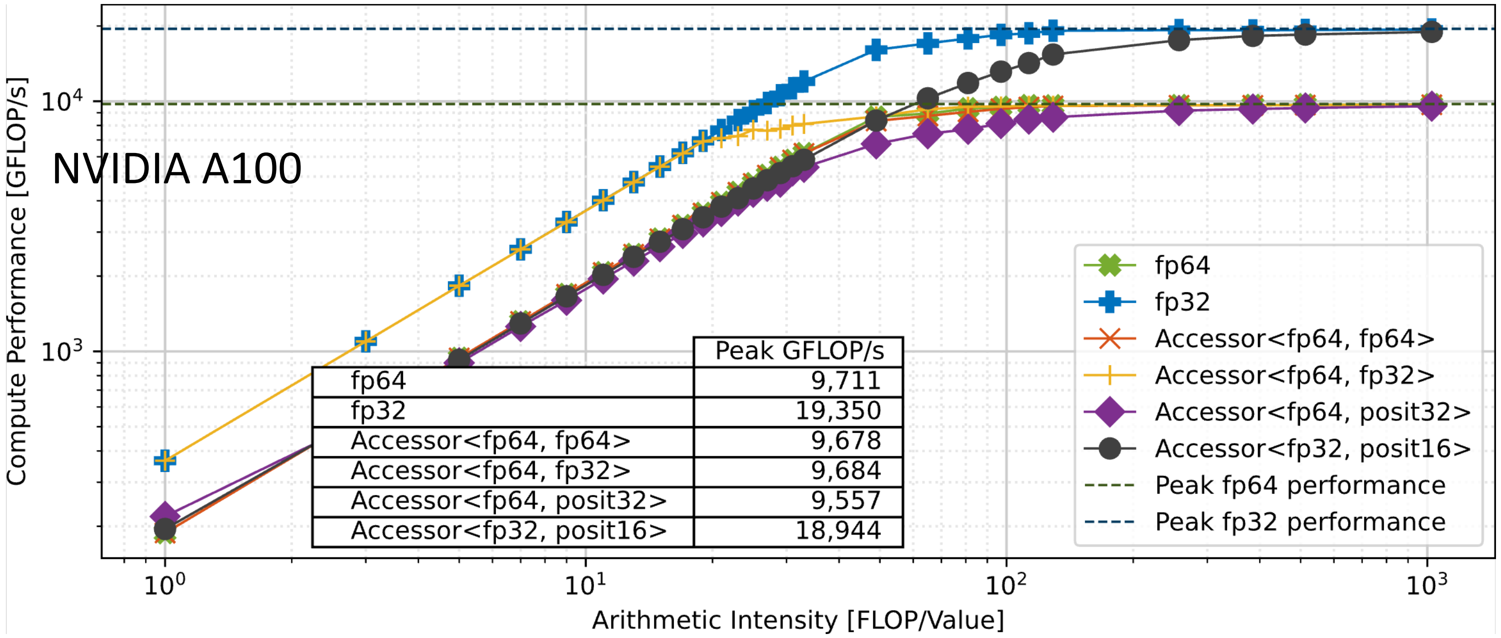
John L. Gustafson



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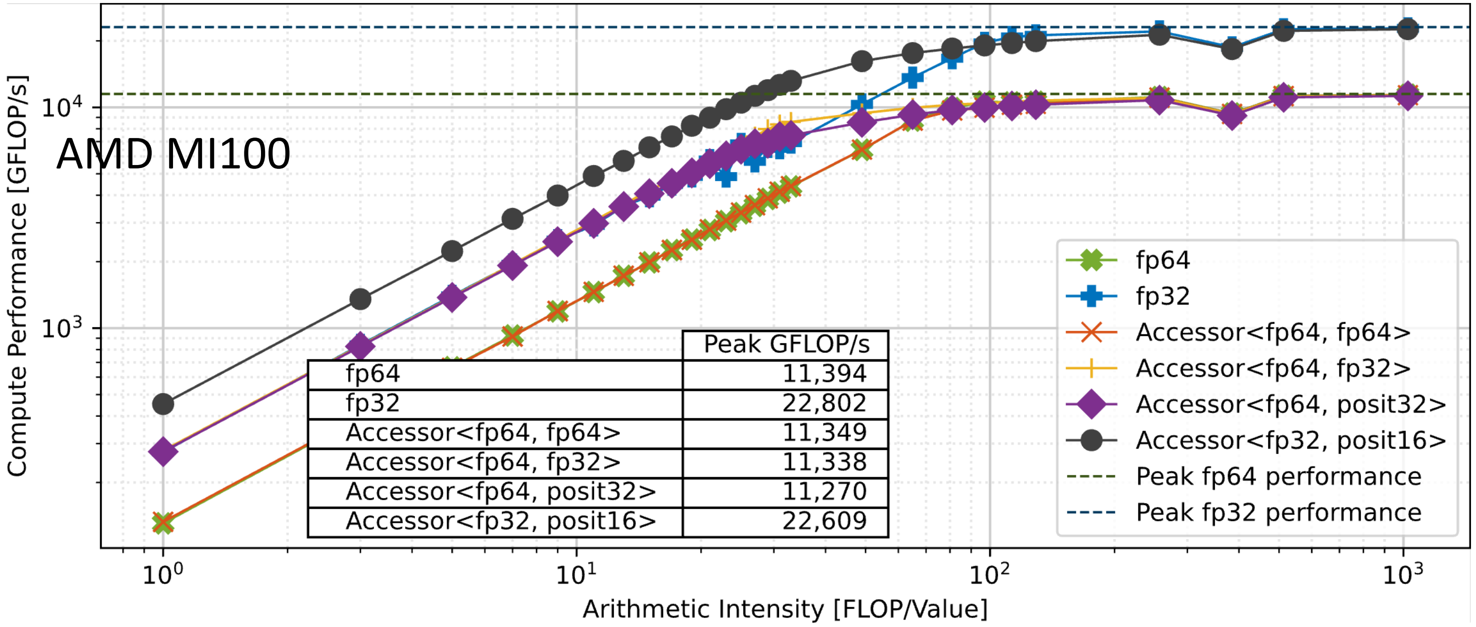
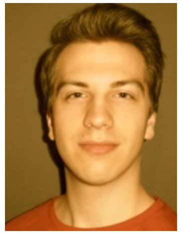


	Peak GFLOP/s
fp64	11,394
fp32	22,802
Accessor<fp64, fp64>	11,349
Accessor<fp64, fp32>	11,338
Accessor<fp64, posit32>	11,270
Accessor<fp32, posit16>	22,609



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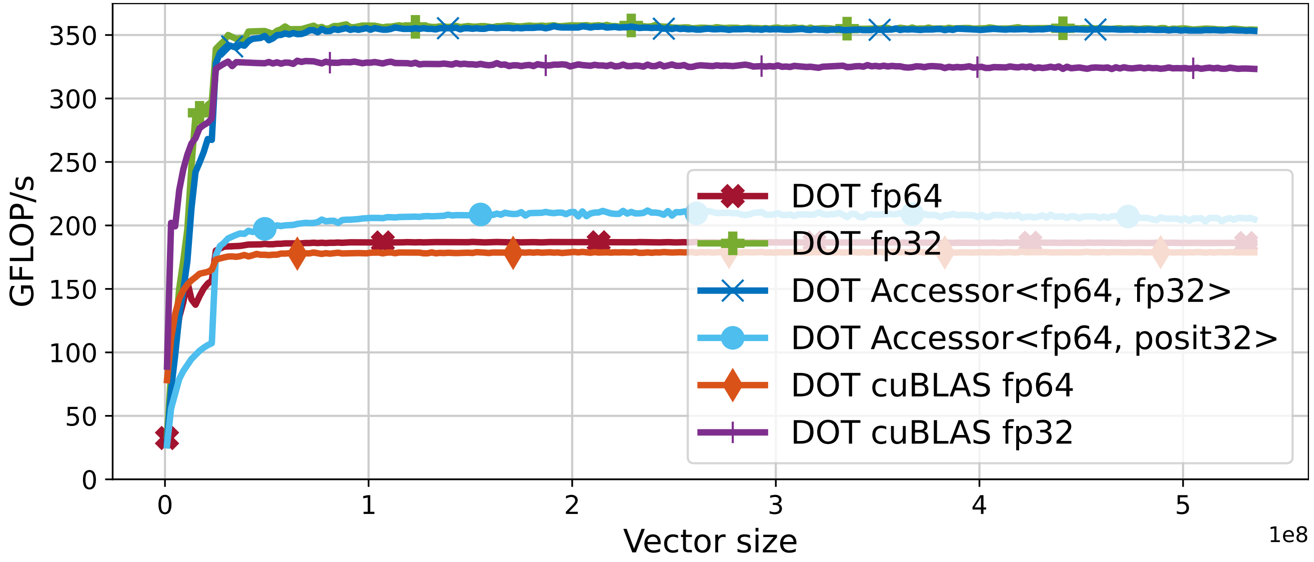
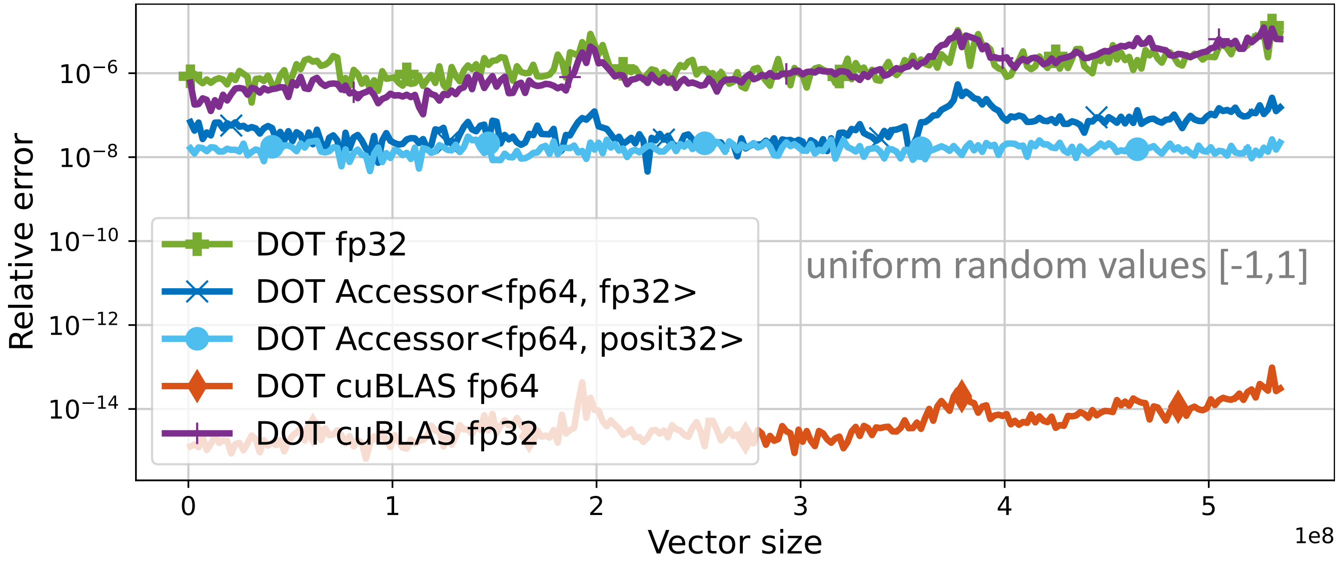
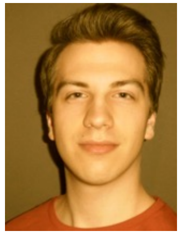
Using POSIT as memory format

NVIDIA A100





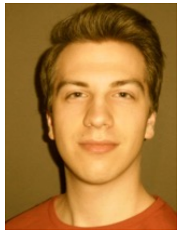


T. Grützmacher

Using ZFP / SZ compression as memory format


F. Capello


R. Underwood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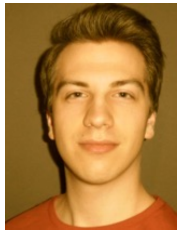
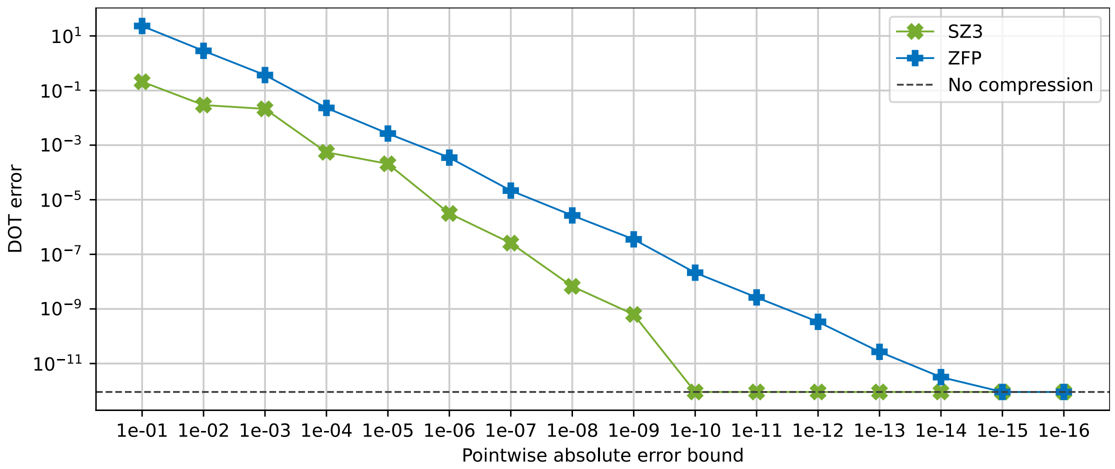
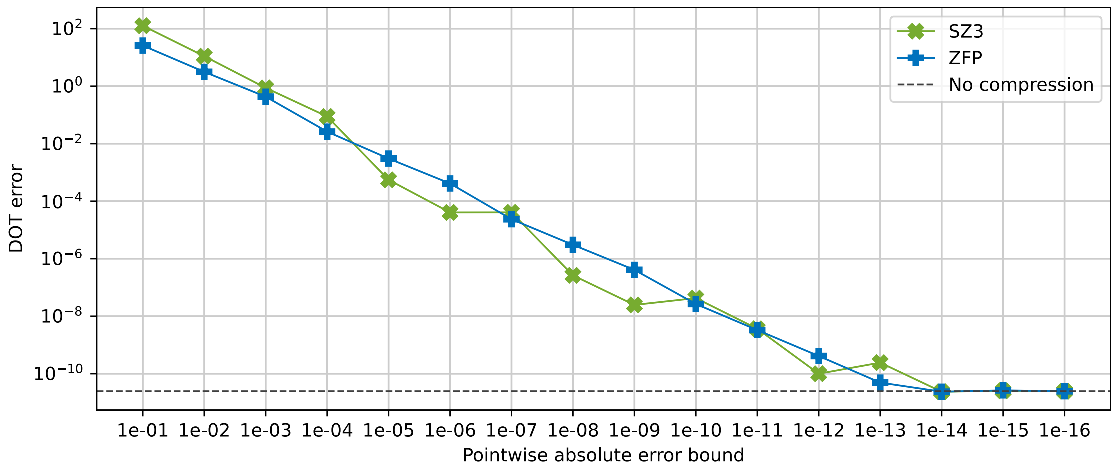

T. Grützmacher

Random data



1D sine function



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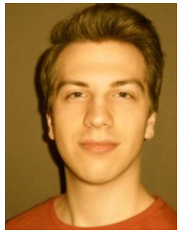
T. Grützmacher

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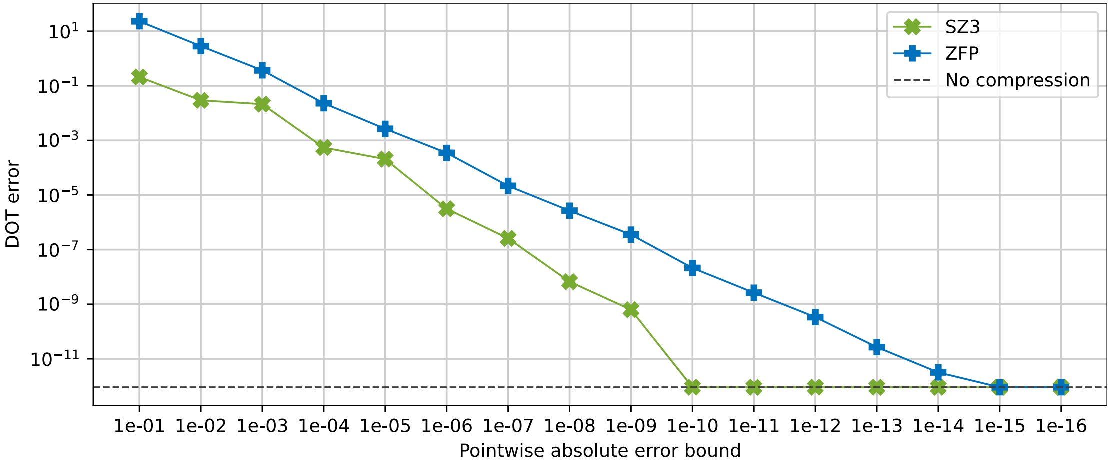
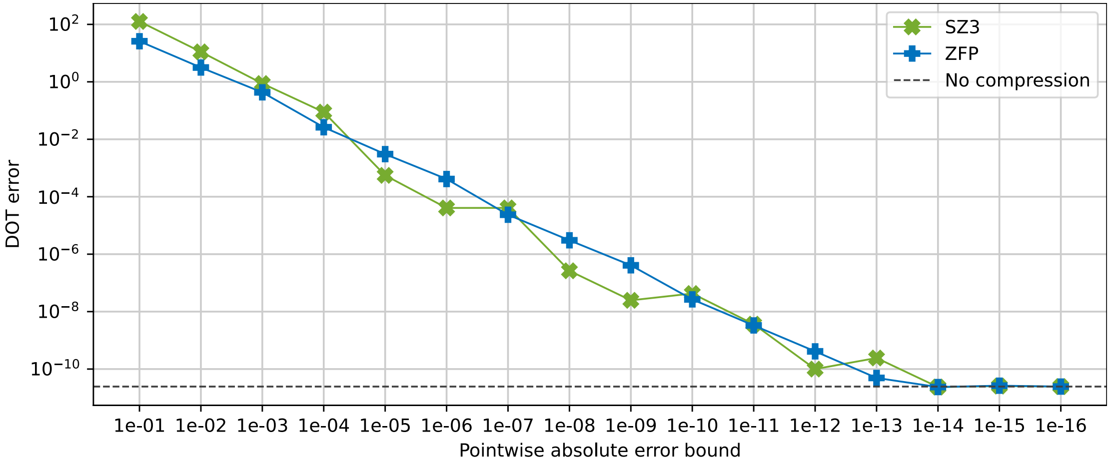
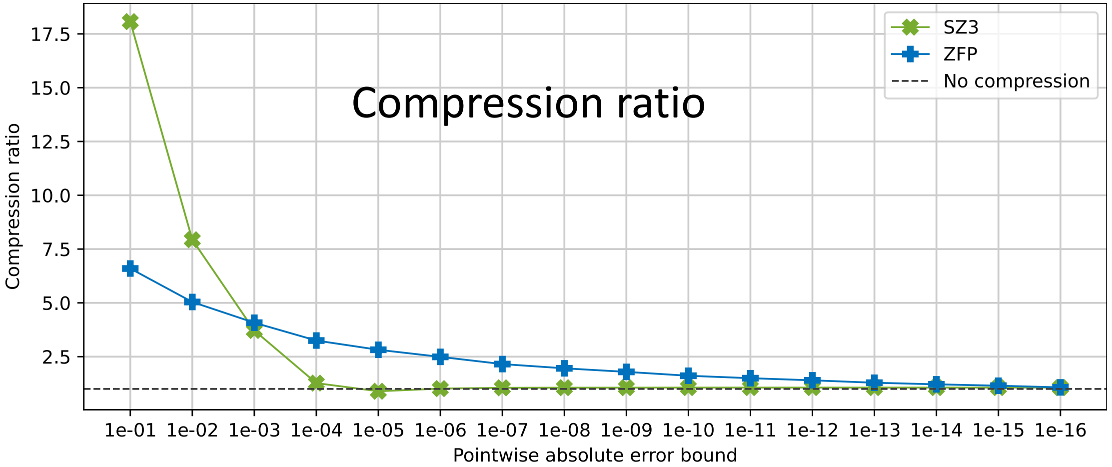
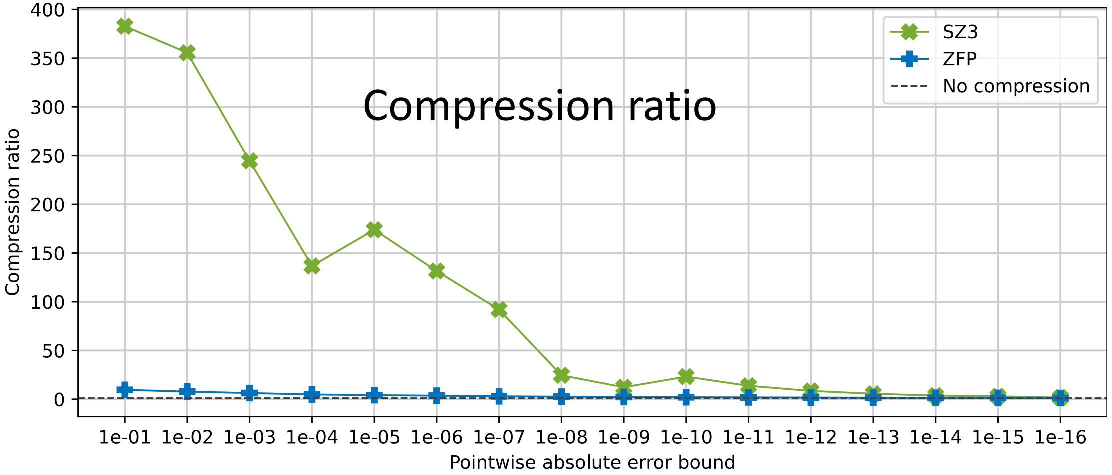


1D sine function

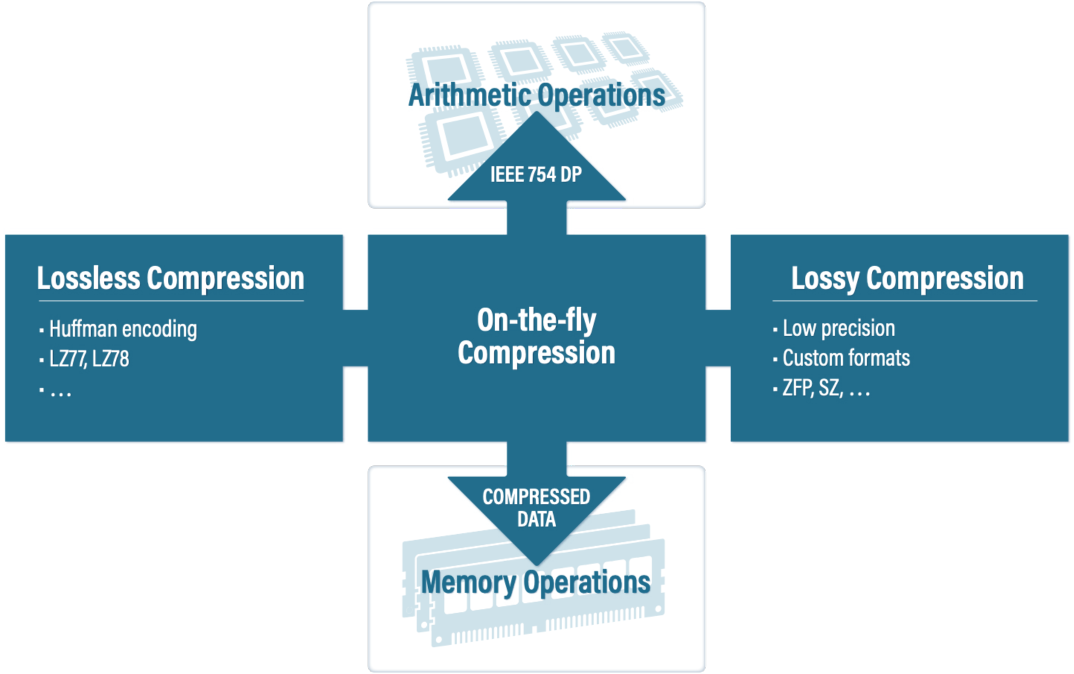






Let's try harder

- IEEE 754 fp64 in arithmetic operations
- More sophisticated in-register compression
    - Custom formats
    - Compression techniques (SZ, ZFP)
- Store data in compressed format



Trade-off:

- Aggressive compression comes with larger information loss
- Element-wise compression allows only for moderate compression ratios
- Block-wise compression makes random access difficult
- Register count limits the block size (hardware specific)
- Data-dependent compression efficiency