

## Project 1

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<https://github.com/harvanchik/335-project-1>

After analyzing my own code for the project, it's clear that the algorithm's efficiency class is  $O(n^2)$ , where  $n$  represents the number of meetings or schedule intervals for each person in the group. Here's a breakdown of my analysis:

The ***updateSchedule*** function copies the person's schedule and daily activities, both of which have an  $O(n)$  complexity since they iterate over the schedule once. So, the total complexity here is  $O(n + n) = O(n)$ .

The ***mergeSchedules*** function combines the schedules of two persons. It iterates through both schedules once, resulting in a nested loop structure. Hence, the complexity of this function is  $O(n^2)$  in the worst case.

The ***sortedSchedules*** function identifies possible availabilities in the merged schedule by iterating over it once, giving it a complexity of  $O(n)$ .

The ***matchedAvailabilities*** function filters availabilities based on the provided duration, iterating over the list of possible availabilities once, resulting in a complexity of  $O(n)$ .

Overall, the most significant factor contributing to the algorithm's time complexity is the ***mergedSchedules*** function with its  $O(n^2)$  complexity. This means that the overall time complexity of the algorithm is  $O(n^2)$ .

To improve efficiency, I could optimize the ***mergedSchedules*** function. One approach would be to sort both people's schedules based on their start times, reducing the complexity to  $O(n * \log(n))$ . However, it's important to note that this would require additional memory for sorting. Importantly, increasing the value of  $n$  wouldn't change the complexity class in this case.

In conclusion, my current algorithm has a time complexity of  $O(n^2)$ . While I recognize the potential for optimization by sorting schedules, I also see that this would introduce some trade-offs.