

The identifiability equation follows by a change of variable $t \mapsto 2t/\beta$. \square

3. Regularization

3.1. Preliminaries

Definition 158 (Operator Fourier transform). *Given a Hamiltonian H and an operator A , define the **operator Fourier transform (FT)** \hat{A}_H by*

$$\hat{A}_H[\omega] = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} A_H(t) e^{-i\omega t} f(t) dt,$$

where $f(t) = e^{-\sigma^2 t^2} \sqrt{\sigma \sqrt{2/\pi}}$ is a Gaussian filter. The “regularizing” role of $f(t)$ will be become clearer in the sequel. Its Fourier transform $\hat{f}[\omega] = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} e^{-i\omega t} f(t) dt$ satisfies $\hat{f}[\omega] = \frac{1}{\sqrt{\sigma \sqrt{2\pi}}} \exp(-\omega^2/4\sigma^2)$.

Note that the operator FT commutes with imaginary time evolution:

$$e^{\beta H} \hat{A}_H[\omega] e^{-\beta H} = (\widehat{e^{\beta H} A e^{-\beta H}})_H[\omega]$$

Taking the operator FT of both sides of Lemma 150 results in the following useful identity:

$$\hat{A}_H[\omega] = \sum_{\nu} A_{\nu} \hat{f}[\omega - \nu].$$

In other words, the operator FT gives “soft” access to the components in the Bohr decomposition of A . We have a corresponding “soft” Bohr decomposition, by Fourier duality.

Lemma 159. *For any operator A and Hermitian H ,*

$$A = C_{\sigma} \int_{-\infty}^{\infty} \hat{A}_H[\omega] d\omega$$

$$\text{for } C_{\sigma} := \frac{1}{\sqrt{2\sigma\sqrt{2\pi}}}.$$

Importantly, a straightforward calculation shows that the Gaussian filter ensures the operator FT decays exponentially in the frequency ω :

Lemma 160. *For any frequency ω and operator A satisfying $\|A\|_{\text{op}} \leq 1$,*

$$\hat{A}_H[\omega] = e^{-\beta\omega + \sigma^2 \beta^2} e^{\beta H} \hat{A}_H[\omega - 2\sigma^2 \beta] e^{-\beta H}$$

To see why this is useful, note that because $\|\hat{A}_H[\omega]\|_{\text{op}} \leq \hat{f}(0) = O(\sigma^{-1/2})$, this ensures that $\|e^{\beta H} \hat{A}_H[\omega'] e^{-\beta H}\|_{\text{op}} \lesssim e^{\sigma^2 \beta^2 + \beta \omega' \sigma^{-1/2}}$. Crucially, the right-hand scales *exponentially in the frequency ω'* , rather than exponentially in the system size! In contrast, norm of the imaginary time-evolved observable $\|e^{\beta H} A e^{-\beta H}\|_{\text{op}}$ can scale exponentially in the system size.

3.2. Truncating the identifiability observable

Using Lemma 159, we can decompose A in $\langle O, [A, H - H'] \rangle_{\rho}$ into low-frequency and high-frequency terms under operator FT with respect to H' and apply the

identifiability equation in Theorem 153 to obtain

$$\begin{aligned} \frac{\beta}{2C_\sigma} \langle O, [A, H - H'] \rangle_\rho &= \\ \int_{|\omega'| \leq \Omega'} \text{tr}(\rho \bar{\Delta} [[H'; O, \hat{A}_{H'}[\omega']]]) \, d\omega' + \frac{\beta}{2} \int_{|\omega'| \geq \Omega'} \langle O, [\hat{A}_{H'}[\omega'], H - H'] \rangle_\rho \, d\omega' . \end{aligned}$$

Let us try to write down a slightly more palatable expression for the first integral that doesn't involve the operator FT. Note that

$$\hat{A}_{H'}[\omega']_{H'}(t - i\beta/2) = (\sqrt{\rho'} \hat{A}_{H'}[\omega'] \sqrt{\rho'^{-1}})_{H'}(t),$$

and

$$\begin{aligned} &\int_{|\omega'| \leq \Omega'} \sqrt{\rho'} \hat{A}_{H'}[\omega'] \sqrt{\rho'^{-1}} \, d\omega' \\ &= \int_{|\omega'| \leq \Omega'} \hat{A}_{H'}[\omega' - \sigma^2 \beta] e^{-\beta\omega'/2 + \sigma^2 \beta^2/4} \, d\omega' \\ &= \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} A_{H'}(t') \underbrace{\int_{|\omega'| \leq \Omega'} e^{-i(\omega' - \sigma^2 \beta)t'} e^{-\beta\omega'/2 + \sigma^2 \beta^2/4} \, d\omega' f(t')}_{h_+(t')} \, dt' , \end{aligned}$$

where in the first step we used Lemma 160, and similarly

$$\begin{aligned} &\int_{|\omega'| \leq \Omega'} \sqrt{\rho'^{-1}} \hat{A}_{H'}[\omega'] \sqrt{\rho'} \, d\omega' \\ &= \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} A_{H'}(t') \underbrace{\int_{|\omega'| \leq \Omega'} e^{-i(\omega' + \sigma^2 \beta)t'} e^{\beta\omega'/2 + \sigma^2 \beta^2/4} \, D\omega' f(t')}_{h_-(t')} \, dt' . \end{aligned}$$

Observe that

$$|h_+(t)|, |h_-(t)| \leq O\left(\frac{\sqrt{\sigma}}{\beta} e^{-\sigma^2 t^2 + \beta\Omega'/2 + \sigma^2 \beta^2/4}\right),$$

i.e. these functions are rapidly decaying in t .

Summarizing, we have the following:

Lemma 161. *Let $\Omega' > 0$ and define the truncated observable*

$$\begin{aligned} &\bar{\Delta}^{\leq \Omega'} [[H'; O, A]] \\ &:= \frac{1}{\sqrt{2\pi}} \iint_{-\infty}^{\infty} (h_+(t') O_H^\dagger(t) A_{H'}(t' + t) - h_-(t') A_{H'}(t' + t) O_H^\dagger(t)) g_\beta(t) \, dt' \, dt . \end{aligned} \tag{59}$$

Then

$$\frac{\beta}{2C_\sigma} \langle O, [A, H - H'] \rangle_\rho = \text{tr}(\rho \bar{\Delta}^{\leq \Omega'} [[H'; O, A]]) + \frac{\beta}{2} \int_{|\omega'| \geq \Omega'} \langle O, [\hat{A}_{H'}[\omega'], H - H'] \rangle_\rho \, d\omega' .$$

Let's take stock of what this buys us. First, because g_β , h_+ , and h_- are rapidly decaying, the bulk of the double integral in the truncated observable $\bar{\Delta}^{\leq \Omega'} [[H'; O, A]]$ is coming from short-time evolutions of O and A , which are local by the aforementioned Lieb-Robinson bounds. In short, if we only look at the “low-degree” term

in Lemma 161, we now have an observable which is entirely local which captures the discrepancy between H and H' .

It still remains to control the truncation error term $\int_{|\omega'| \geq \Omega'} \dots$. For this, we can use Lemma 160 in conjunction with locality of $H - H'$ and A to show that for $\Omega' \geq \Omega(\sigma^2/\mathfrak{d})$, where \mathfrak{d} is the degree of the dual interaction graph, the truncation error is negligible. We defer the details to [CAN25, Lemma III.5].

As discussed above, there is still one important missing piece before we can turn the above into a learning algorithm. The issue is that the truncated observable in Eq. (59) ultimately still depends on H through $O_H^\dagger(t)$. We explain the workaround for this next.

4. Learning Algorithm

In this section we describe how to exploit the ingredients from the preceding sections, deferring a complete proof of correctness to [CAN25].

To sidestep the issue that the truncated observable defined in Eq. (59) depends on H , we first define a broader class of observables that contains this observable.

Definition 162 (General truncated observables). *Fix $\Omega' > 0$. Given operators K, O, A, G , with K and G Hermitian, define*

$$\Delta^* \llbracket G, K; O, A \rrbracket := \frac{1}{\sqrt{2\pi}} \iint_{-\infty}^{\infty} \left(h_+(t') O_G^\dagger(t) A_K(t' + t) - h_-(t') A_K(t' + t) O_G^\dagger(t) \right) g_\beta(t) dt' dt. \quad (60)$$

Note that $\Delta^* \llbracket H, H'; O, A \rrbracket = \overline{\Delta}^{\leq \Omega'} \llbracket H'; O, A \rrbracket$.

By design, we have the following:

Proposition 163. *When $K = H$, then $\text{tr}(\rho \Delta^* \llbracket G, K; O, A \rrbracket) = 0$ for all O, G, A .*

PROOF. In the proof of Lemma 161, instead of passing to h_+, h_- , we can directly express the “low-degree” term $\text{tr}(\rho \overline{\Delta}^{\text{trunc}}(H'; O, A))$ as

$$\frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} \text{tr} \left[\rho \int_{|\omega'| \leq \Omega'} \left(O_H^\dagger(t) \hat{A}_{H'}[\omega']_{H'}(t+i\beta/2) - \hat{A}_{H'}[\omega']_{H'}(t-i\beta/2) O_H^\dagger(t) \right) d\omega' \right] dt.$$

In the definition of $\Delta^* \llbracket G, K; O, A \rrbracket$, H' and H above are replaced by K and G respectively, yielding

$$\frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} \text{tr} \left[\rho \int_{|\omega'| \leq \Omega'} \left(O_G^\dagger(t) \hat{A}_K[\omega']_K(t+i\beta/2) - \hat{A}_K[\omega']_K(t-i\beta/2) O_H^\dagger(t) \right) d\omega' \right] dt.$$

If $K = H$ however, then mirroring the proof of the KMS condition, we have that $\text{tr}(\rho O_G^\dagger(t) \hat{A}_H[\omega']_H(t+i\beta/2)) = \text{tr}(\sqrt{\rho} \hat{A}_K[\omega']_H(t) \sqrt{\rho}) = \text{tr}(\rho \hat{A}_H[\omega']_H(t-i\beta/2) O_H^\dagger(t))$,

so $\Delta^* \llbracket G, H; O, A \rrbracket = 0$ as claimed. \square

This suggests that we can simply brute-force enumerate over a net of different K ’s, and for each one we check whether $\text{tr}(\rho \Delta^* \llbracket G, K; O, A \rrbracket) \approx 0$ for all 1-local Paulis A , and O, G in a suitable net. Previously we considered taking $O = [A, H - H']$, but given that this depends on H , we can instead use the fact that

$$\|[A, H - H']\|_\rho^2 \leq 2\mathfrak{d} \max_a |\langle [A, P_a], [A, H - H'] \rangle_\rho| \quad (61)$$

to restrict to $O = [A, P_a]$ for all 1-local Paulis A and terms a in the support of the Hamiltonian.

The (regularized) identifiability equation in Lemma 161, combined with Lemma 154 and the inequality in Eq. (61), ensures that if $\text{tr}(\rho \Delta^* \llbracket G, K; O, A \rrbracket) \approx 0$ for all 1-local A , $O = [A, P_a]$, and G in a suitable net, then $K \approx H$.

Only one step remains: how do we enumerate over G, K ? Naively, if the Hamiltonian has m terms, this would require enumerating over a net over $O(m)$ -dimensional parameter space and incurring a runtime scaling exponentially in $O(m)$. Fortunately, there is a workaround that again exploits locality. The intuition is that in the definition of $\Delta^* \llbracket G, K; O, A \rrbracket$ in Eq. (60), if A is a 1-local Pauli acting on site i , then $A_K(t'+t)$ and $O_G^\dagger(t) = [A, P_a]_G^\dagger(t)$ are roughly supported on a small neighborhood around i (because t', t are not too large because of the exponential damping of g_β, h_+, h_-). Moreover, Lieb-Robinson bounds ensure that these operators do not change much when G and K are replaced by their truncations to a suitable neighborhood around the i -th site. Formally, we have the following estimate:

Lemma 164 (Lieb-Robinson bound). *If Hamiltonian $H = \sum_a \lambda_a P_a$ with coefficients satisfying $|\lambda_a| \leq 1$ has interaction degree \mathfrak{d} , then for any operator A acting on subsystem $S \subseteq [n]$ and satisfying $\|A\|_{\text{op}} \leq 1$, if H_ℓ is given by removing all terms from H at distance at least ℓ from S , then*

$$\|A_{H_\ell}(t) - A_H(t)\|_{\text{op}} \leq O\left(|S| \cdot \frac{(2\mathfrak{d}|t|)^\ell}{\ell!}\right).$$

The proof of this will be the subject of one of the homework exercises.

With this in hand, we essentially have a complete, albeit informal, description of the algorithm:

- For each qubit $i \in [n]$:
 - (1) Enumerate over a net of local Hamiltonians K_ℓ acting on the neighborhood $V(\ell, i)$ of radius ℓ around the i -th site
 - (2) For each such K_ℓ , use $O(\log n)$ copies of ρ to estimate the observable values $\text{tr}(\rho \Delta^* \llbracket G_\ell, K_\ell; [A, P_a], A \rrbracket)$ for all local Hamiltonians G_ℓ acting on $V(\ell, i)$ and all terms P_a and 1-local Paulis A .
 - (3) If for any such K_ℓ all of these observable values are small, then we will take our estimate of H over the local patch $V(\ell, i)$ to be K_ℓ .

The quantitative details are somewhat dense and do not provide much additional insight beyond the intuition outlined above, so we defer these to [CAN25].