

ICCS200: Assignment 3  
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**1: Tail sum of Squares**

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Given Program :

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int sumHelper(int n , int a) {  
    if (n==0) return a;  
    else return sumHelper(n-1, a + n*n); }  
  
int sumSqr(int n) {  
    return sumHelper(n, 0); }
```

Prove for  $n \geq 1$ ,  $\text{sumSqr}(n) \hookrightarrow 1^2 + 2^2 + 3^2 + \dots + n^2$ .

Predicate:

$$P(n) \equiv \forall a, \text{sumHelper}(n,a) \hookrightarrow a + \frac{n(2n+1)(n+1)}{6}$$

Base Case:

$$P(1) \equiv \text{sumHelper}(1,0) \hookrightarrow 1 = 1^2$$

Inductive Step:

Assume  $\forall k < n$   $P(k)$  is true.

$$\text{sumHelper}(n-1,a) \hookrightarrow a + 1^2 + 2^2 + 3^2 + \dots + (n-1)^2.$$

$$\text{sumHelper}(n,a) \hookrightarrow a + 1^2 + 2^2 + 3^2 + \dots + (n-1)^2 + n^2.$$

$$\text{We know that } a + 1^2 + 2^2 + 3^2 + \dots + (n-1)^2 = a + \frac{n(n-1)(2n-1)}{6}.$$

Therefore  $a + 1^2 + 2^2 + 3^2 + \dots + (n-1)^2 + n^2 = a + \frac{n(n-1)(2n-1)}{6} + n^2$ . By Inductive Hypothesis.

$$\Rightarrow a + 1^2 + 2^2 + 3^2 + \dots + (n-1)^2 + n^2$$

$$= a + \frac{n(n-1)(2n-1) + 6n^2}{6}$$

$$= a + \frac{2n^3 + 3n^2 + n}{6}$$

$$= a + \frac{n(2n^2 + 3n + 1)}{6}$$

$$= a + \frac{n(2n+1)(n+1)}{6}$$

Therefore proved

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**2: Mysterious Function**

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Predicate :

$$P(n) \equiv \forall n \geq 1, foo(n) \hookrightarrow (p, q), \frac{p}{q} = 1 - \frac{1}{n+1}$$

Base case :

$$P(1) \equiv foo(1) \hookrightarrow (1, 2) \text{ such that } \frac{1}{2} = 1 - \frac{1}{2} \text{ which is true}$$

Inductive Step:

Assume  $P(n-1)$  is true.

$$P(n-1) \equiv foo(n-1) \hookrightarrow (p, q), \frac{p}{q} = 1 - \frac{1}{n} = \frac{n-1}{n}$$

Therefore  $p = n-1$  and  $q = n$

By Inductive Hypothesis

$$foo(n) \hookrightarrow (q + p * n * (n+1), q * n * (n+1))$$

$$= (n + (n-1) * n * (n+1), n * n * (n+1))$$

$$= n + (n^2 - n) * (n+1), n^3 + n^2$$

$$= n + n^3 - n^2 + n^2 - n, n^3 + n^2$$

$$= n^3, n^3 + n^2$$

$$\text{for } P(n), p = n^3 \text{ and } q = n^3 + n^2$$

$$foo(n) \hookrightarrow \frac{n^3}{n^3 + n^2}$$

$$= \frac{n^2(n)}{n^2(n+1)}$$

$$= \frac{n}{n+1}$$

$$= \frac{n+1-1}{n+1}$$

$$= \frac{n+1}{n+1} - \frac{1}{n+1}$$

$$= 1 - \frac{1}{n+1}$$

$$\text{since } foo(n) \hookrightarrow 1 - \frac{1}{n+1}$$

Therefore  $P(n)$  is true.

### 3: Missing Tile

(task 1) proof by Induction

Predicate:

$$P(i) \equiv \text{grid of size } 2^n \times 2^n \text{ with a painted cell is L-tilable.}$$

Base Case:

$$P(1) \equiv \text{grid of size } 2^1 \times 2^1 \text{ is L-tilable with a painted cell is true.}$$

Inductive Step:

Assume  $P(k) \equiv$  grid size  $2^k \times 2^k$  with a painted cell is L-tilable.

We break the grid into four sections of size  $2^k \times 2^k$  each.

If the painted cell is in the middle and we add an L tile facing away from the painted cell, we will be left with 4 grids of size  $2^k \times 2^k$  each with with an already occupied cell from the painted cell and the L-tile we placed.

By Inductive Hypothesis,

each of the  $2^k \times 2^k$  grid is L-tilable, therefore the grid with a painted cell is tilable.

### 5: Midway Tower of Hanoi

#### (task 1)

$$T(n) = 2T(n-1) + 1$$

$$\text{Predicate: } P(i) \equiv T(i) = 2^i - 1$$

Base Case:

$$T(i) = 2^1 - 1 = 1$$

Inductive Step:

$$\text{Assume } T(n) = 2^n - 1$$

$$T(n+1) = 2T(n) + 1 \text{ (By Inductive Hypothesis)}$$

$$= 2(2^n - 1) + 1$$

$$= (2^{n+1}) - 2 + 1$$

$$= (2^{n+1}) - 1$$

$$\text{Therefore } T(n+1) = (2^{n+1}) - 1$$

Which is what we have to show

Therefore proved.

#### (task 3)

##### (1)

$$g(n) = 2g(n-1) + n$$

$$g(n) = af(n) + bn + c$$

$$f(n) = 2T(n-1) + 1$$

$$g(0) = af(0) + b \cdot 0 + c = 0$$

Therefore  $c=0$

**(2)**

$$g(n) = 2g(n-1) + n$$

$$g(n) = af(n) + bn$$

$$af(n) + bn = 2(af(n-1) + b(n-1)) + n$$

$$a(f(n) + 2f(n-1)) + bn = 2b(n-1) + n$$

$$a(1) + bn + 2b - 2bn - n = 0$$

$$a(1) + 2b - n(b+1) = 0$$

$$(P)_n + (Q)$$

By this P and Q are both 0

Solving for a and b:

$$b + 1 = 0$$

$$b = -1$$

$$a + 2b = 0$$

$$a = -2b$$

$$a = 2$$

Therefore  $g(n) = 2f(n) - n$

**(3)**

Since  $f(n)=2^n - 1$

and  $g(n)=2f(n)-n$

Therefore closed form of  $g(n) = 2(2^n - 1) - n$

**(4)**

Predicate:

$$P(i) \equiv g(i) = 2f(i) - i$$

Base Case:

$$g(0) \equiv 2f(0) - 0 = 0 \text{ is true}$$

Inductive Step:

Assume  $g(k) = 2f(k) - k \quad \forall k > 0$

Given:

$$g(n) = 2g(n-1) + n$$

$$g(k+1) = 2g(k) + k + 1 \text{ By IH:}$$

$$g(k+1)=2(2f(k)-k)+(n+1)$$

$$= 4f(k)-k+1+1-1$$

$$=2(f(k)+1)-(k-1)$$

$$=2(f(k+1))-(k-1)$$