

## Homework 2

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- 1) (10) Given the grammar below, identify which sentences are in the language (which are valid sentence).

- a. baab
- b. bbbab
- c. bbaaaaaa
- d. bbaab

$$\langle S \rangle \rightarrow \langle A \rangle a \langle B \rangle b$$
$$\langle A \rangle \rightarrow \langle A \rangle b \mid b$$
$$\langle B \rangle \rightarrow a \langle B \rangle \mid a$$

From the first line,  $\langle S \rangle \rightarrow \langle A \rangle a \langle B \rangle b$ , we know that the ending to any sentence must have the letter b, meaning c is invalid. Continuing through the grammar we have  $\langle A \rangle \rightarrow \langle A \rangle b \mid b$ , which means that the first letter must be the letter b. Plugging these values in we get two possible strings:

$ba\langle B \rangle b$  or  $bba\langle B \rangle b \rightarrow baab$  or  $bbaab \rightarrow$  Meaning only **a** and **d** are valid

- 2) (10) Identify all of the tokens (categories of lexemes) in the grammar below, and which lexemes they categorize. Put them in a table.

$$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$$
$$\langle \text{id} \rangle \rightarrow A \mid B \mid C$$
$$\langle \text{expr} \rangle \rightarrow \langle \text{id} \rangle + \langle \text{expr} \rangle$$
$$\mid \langle \text{id} \rangle * \langle \text{expr} \rangle$$
$$\mid ( \langle \text{expr} \rangle )$$
$$\mid \langle \text{id} \rangle$$

Tokens	Lexemes
Variable	A,B,C
Add_op	+
Mult_op	*
Assign_op	=
Parenthesis	(,)

- 3) (10) Given the grammar from question 2, show a left-most derivation and draw the parse tree for the following statement.

a.  $B = B + (C + (A * A))$

$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$

$\rightarrow B = \langle \text{expr} \rangle$

$\rightarrow B = \langle \text{id} \rangle + \langle \text{expr} \rangle$

$\rightarrow B = B + \langle \text{expr} \rangle$

$\rightarrow B = B + (\langle \text{expr} \rangle)$

$\rightarrow B = B + (\langle \text{id} \rangle + \langle \text{expr} \rangle)$

$\rightarrow B = B + (C + \langle \text{expr} \rangle)$

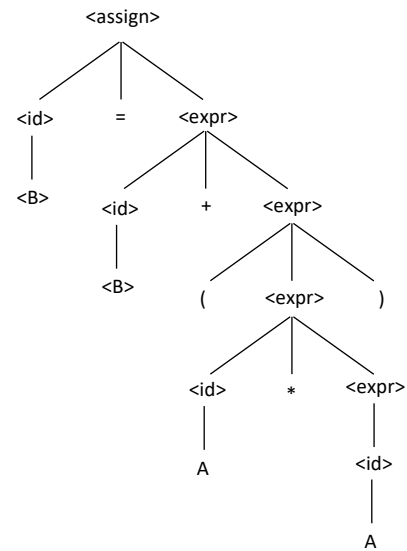
$\rightarrow B = B + (C + (\langle \text{expr} \rangle))$

$\rightarrow B = B + (C + (\langle \text{id} \rangle * \langle \text{expr} \rangle))$

$\rightarrow B = B + (C + (A * \langle \text{expr} \rangle))$

$\rightarrow B = B + (C + (A * \langle \text{id} \rangle))$

$\rightarrow B = B + (C + (A * A))$



- 4) (10) Remove all of the recursion from the following grammar:

$S \rightarrow Aa \mid Bb$

$A \rightarrow Aa \mid AbC \mid C$

$B \rightarrow S \mid bb$

$C \rightarrow c$



$S \rightarrow Aa \mid Bb$

$A \rightarrow AA'$

$A \rightarrow C$

$A' \rightarrow a \mid bC$

$B \rightarrow S \mid bb$

$C \rightarrow c$

- 5) (10) Use left factoring to resolve the pairwise disjointness problems in the following grammar:

$A \rightarrow aBc \mid ac \mid a$

$B \rightarrow b \mid aB$



$A \rightarrow aA'$

$A' \rightarrow Bc \mid c$

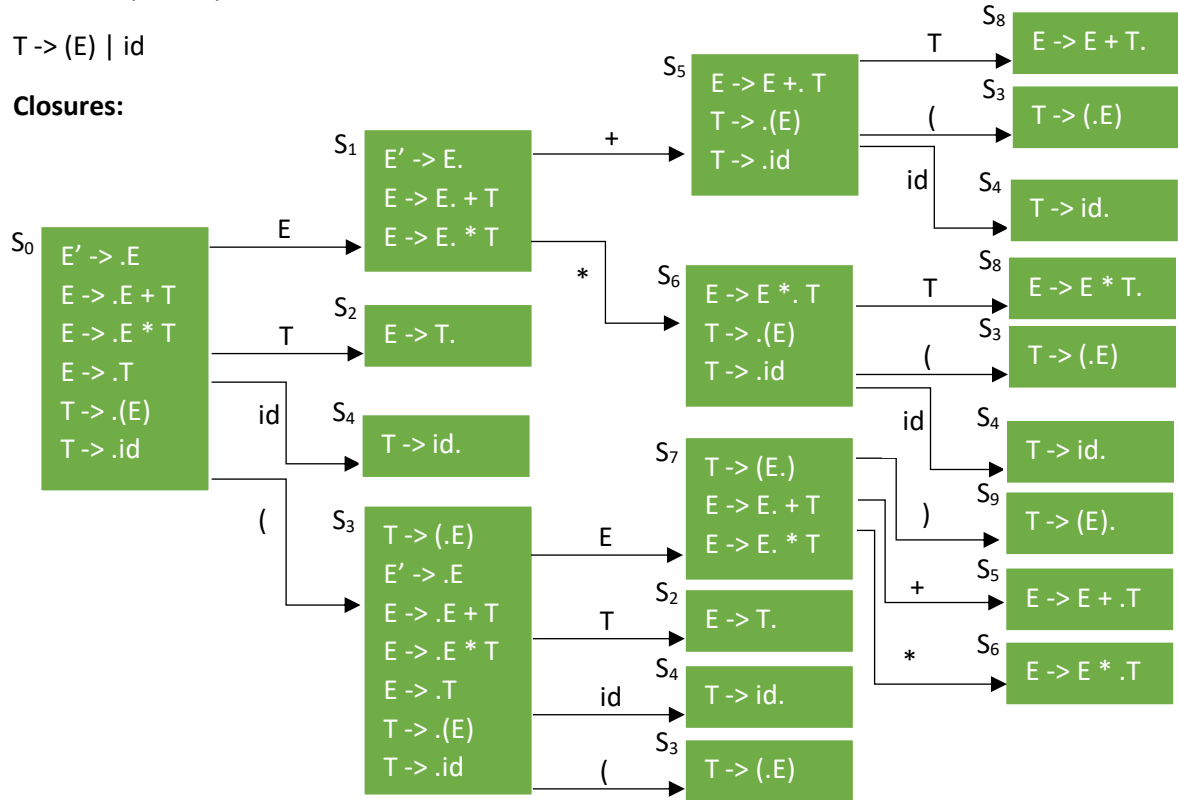
$B \rightarrow b \mid aB$

- 6) (20 pts) Create an LR(0) parse table for the following grammar. Show all steps (creating closures, the DFA, the transition table, and finally the parse table):

$E \rightarrow E + T \mid E * T \mid T$

$T \rightarrow (E) \mid id$

**Closures:**



State	Action						Goto	
	+	*	(	)	id	\$	E	T
0			S3		S4	acc	1	2
1	S5	S6						
2	R3	R3		R3				
3			S3		S4		7	2
4	R5	R5		R5				
5			S3		S4			8
6			S3		S4			8
7	S5	S6		S9				
8	R2	R2		R2				
9	R4	R4		R4				

- 7) (20 pts) Show a complete bottom-up parse, including the parse stack contents, input string, and action for the string below using the parse table you created in step 6. Think about how I went through this in class.

**Parse stack:**

**Input String:**

$(id + id) * id$

- 8) (10 pts) Show a rightmost derivation for the string above, and show how the bottom-up parse you completed in step 7 correctly finds all of the handles for the input string above.

$(id + id) * id$