COMP 416: Computer Networks - Project 3

Network Layer Analysis and DV Routing Simulator Hasan Can Aslan - 60453

Part 1. Network Layer Analysis

Part 1A. Network Interface Analysis

1. We provide 2 parameters to netstat command. First -a parameter displays all active TCP connections and the TCP and UDP ports on which the computer is listening, -i parameter displays network interface status. So netstat -ai command shows the state of the network interfaces and their TCP connections.

• •																	
Name		Mtu Network	Address	Ipkt	ts Ier	rs 0	pkts Oer	s Coll									
lo0	16384	<link#1></link#1>		1042147	0 1	.042147											
lo0	16384	127	localhost	1042147	- 1	.042147			llw0	1500	<link#13></link#13>	1a:43:3d:51:25:33	0	0	0	0	0
			224.0.0.251						llw0	1500	fe80::1843:	fe80:d::1843:3dff					
			all-systems.mcast.	net								ff01::1	(refs: 1)				
lo0	16384			1042147	- 1	.042147						ff02::1	(refs: 1)				
			ff02::fb	(refs: 1)								ff02::1:ff51:2533	(refs: 1)				
			ff02::2:ff25:e58e	(refs: 1)								ff02::fb	(refs: 1)				
			ff01::1	(refs: 1)								ff02::1:ffd7:2354	(refs: 3)				
			ff02::1	(refs: 1)								ff02::1:ff2f:6d75	(refs: 4)				
			ff02::1:ff00:1	(refs: 1)								ff02::2:ff8a:dd10	(refs: 11)				
Lo0	16384	hasan-macbo	fe80:1::1	1042147	- 1	.042147						ff02::1:ffbd:2c9a	(refs: 3)				
			ff02::fb	(refs: 1)					utun0	1380	<link#14></link#14>		· 0	Θ	20	Θ	Θ
			ff02::2:ff25:e58e	(refs: 1)					utun0	1380	hasan-macbo	fe80:e::e457:d006	0		20		
			ff01::1	(refs: 1)								ff01::1	(refs: 1)				
			ff02::1	(refs: 1)								ff02::2:ff8a:dd10	(refs: 1)				
			ff02::1:ff00:1	(refs: 1)								ff02::1	(refs: 1)				
rif0∗	1280	<link#2></link#2>		0	Θ	Θ	0	Θ				ff02::1:ffda:fd79	(refs: 1)				
tf0*		<link#3></link#3>			ē	0	9	ō	utun1	2000	<link#15></link#15>		0	0	20	Θ	Θ
p1*		<link#5></link#5>	f2:18:98:62:13:01	ő	0	0	o o	0				fe80:f::60e1:e539	ō		20		
		<link#6></link#6>	f0:18:98:62:13:01	14531107	0 20	312842	75214	0				ff01::1	(refs: 1)				
			fe80:6::851:25c0:			312842						ff02::2:ff8a:dd10	(refs: 1)				
			ff02::fb	(refs: 1)								ff02::1	(refs: 1)				
			ff01::1	(refs: 1)								ff02::1:ff69:f22e					
			ff02::2:ff8a:dd10	(refs: 1)					utun2	1380	<link#16></link#16>		102	0	100	Θ	Θ
			ff02::1	(refs: 1)								fe80:10::4e2b:2ad	102		100		
			ff02::1:ff3a:78a4	(refs: 1)					0.00112	1000	nasan maces	ff01::1	(refs: 1)		200		
en0	1500	192.168.1		14531107	- 26	312842						ff02::2:ff8a:dd10	(refs: 1)				
	1000		239.255.255.250	1.001101		.0120.2						ff02::1	(refs: 1)				
			224.0.0.251									ff02::1:ffbc:8a45	(refs: 1)				
			all-systems.mcast.	net					utun3	1380	<link#17></link#17>	1102	0	0	20	е	ю
n1	1500		82:dd:b9:a2:4c:01	0		Θ	Θ	Θ				fe80:11::97c5:f07	9		20		
			82:dd:b9:a2:4c:00	0	0	0	0	0	acaris	1300	nasan macbo	ff01::1	(refs: 1)				
	1500		82:dd:b9:a2:4c:05		0	0	0	0				ff02::2:ff8a:dd10	(refs: 1)				
			82:dd:b9:a2:4c:04	0	0	0	0	0				ff02::1	(refs: 1)				
oridg			82:dd:b9:a2:4c:01	0		0	0	0				ff02::1:ff27:e9b7	(refs: 1)				
awdl0		<link#11></link#11>	1a:43:3d:51:25:33	14761		14426	9	0	en5	1500	<link#4></link#4>	ac:de:48:00:11:22	41518		41198	66	Θ
awd to			fe80:c::1843:3dff	14761		14426			en5			fe80:4::aede:48ff	41518		41198	-	
ma co	1000		ff01::1	(refs: 1)		14420			CIIJ	1300	resoaede:	ff01::1	(refs: 1)		41130		
			ff02::1	(refs: 1)								ff02::1	(refs: 1)				
			ff02::1:ff51:2533	(refs: 1)								ff02::2:ff8a:dd10	(refs: 1)				
			ff02::2:ff8a:dd10	(refs: 1)								ff02::1:ff00:1122					
			ff02::fb	(refs: 1)								11021.1100:1122	(reis: 1)				

2. <u>Destination column:</u> Shows the destination network.

Gateway column: Indicates the router through which packets are forwarded.

<u>Flags column:</u> The U flag indicates that the route is up route is valid.

The G flag indicates that the route is to a gateway.

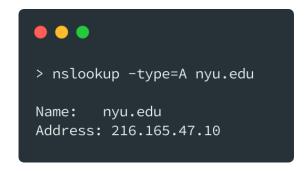
The S flag indicates route added with the route command.

The H flag indicates route is to a host rather than to a network, where the destination address is a complete address.

<u>Netif column:</u> Indicates the interface on the local host that is the source endpoint of the transmission.

Internet:				
Destination	Gateway	Flags		Expire
default	192.168.1.1	UGScg	en0	
127	localhost	UCS	lo0	
localhost	localhost	UH	lo0	
169.254	link#6	UCS	en0	
169.254.91.108	link#6	UHLSW	en0	
169.254.103.19	link#6	UHLSW	en0	
169.254.196.189	link#6	UHLSW	en0	
192.168.1	link#6	UCS	en0	
192.168.1.1/32	link#6	UCS	en0	
192.168.1.1	c8:3a:35:dc:2d:58	UHLWIir	en0	1166
192.168.1.2	7c:1:91:65:88:b6	UHLWIi	en0	293
192.168.1.3	40:b8:9a:63:fa:d	UHLWI	en0	616
192.168.1.5	6c:96:cf:89:b1:95	UHLWI	en0	733
192.168.1.12	f2:4b:73:fe:12:19	UHLWIi	en0	967
192.168.1.13/32	link#6	UCS	en0	
192.168.1.255	ff:ff:ff:ff:ff	UHLWbI	en0	
224.0.0/4	link#6	UmCS	en0	
224.0.0.251	1:0:5e:0:fb	UHmLWI	en0	
239.255.255.250		UHmLWI	en0	
255.255.255.255/32	link#6	UCS	en0	1

3. My client using ports 64532 and 64531 to connect nyu.edu. First, I get the IP address of the host, using nslookup.



Then using netstat -P tcp command, we can find nyu.edu's IP address in the list.

```
> netstat -an

Active Internet connections (including servers)
Proto Recv-Q Send-Q Local Address Foreign Address (state)
.
.
.
.
.
tcp4 0 0 192.168.1.13.64532 216.165.47.10.80 ESTABLISHED tcp4 0 0 192.168.1.13.64531 216.165.47.10.80 ESTABLISHED
.
.
.
```

Part 1B. ICMP Analysis

4. TTL, time to live is the number of hops or amount of time that a packet exists inside a network. When TTL exceeded, packet is discarded. Time to live is important for caching information. It determines how long the data can be cached and when the information should be updated. We can find TTL under the network layer, for example under the Internet Protocol.

5. ICMP protocol is designed to communicate between network layers. However, port numbers are used to communicate between application layers, so ICMP has neither destination nor source port number.

6. Minimum TTL is 1.

```
Frame 126: 182 bytes on wire (1456 bits), 182 bytes captured (1456 bits) on interface en0, id 0 Ethernet II, Src: TendaTec_dc:2d:58 (c8:3a:35:dc:2d:58), Dst: Apple_62:13:01 (f0:18:98:62:13:01)
Internet Protocol Version 4, Src: 64.125.26.172, Dst: 192.168.1.13
Internet Control Message Protocol
   Type: 11 (Time-to-live exceeded)
  Code: 0 (Time to live exceeded in transit)
   Checksum: 0xbe96 [correct]
   [Checksum Status: Good]
   Unused: 00000000
  Internet Protocol Version 4, Src: 192.168.1.13, Dst: 216.165.47.10
      0100 .... = Version: 4
      .... 0101 = Header Length: 20 bytes (5)
     Differentiated Services Field: 0x00 (DSCP: CS0, ECN: Not-ECT)
      Total Length: 52
      Identification: 0x9507 (38151)
   > Flags: 0x00
      Fragment Offset: 0
   > Time to Live: 1
      Protocol: UDP (17)
     Header Checksum: 0x5b4d [validation disabled]
      [Header checksum status: Unverified]
      Source Address: 192.168.1.13
     Destination Address: 216.165.47.10
  User Datagram Protocol, Src Port: 38120, Dst Port: 33465
> ICMP Multi-Part Extensions
Data (24 bytes)
```

7. My computer sends first a packet with TTL = 1, the second packet has TTL = 2 so and so on. A router will decrement the packet's TTL when the packet passes through. If the packet arrives the router with TTL value equals to 1, then router sends an ICMP packet to the source. In this way we can count the router along the path to the destination.

8. For each router along the path, traceroute sends 3 probes.

```
traceroute to nyu.edu (216.165.47.18), 64 hops max, 52 byte packets

1 192.168.1.1 (192.168.1.1) 9.176 ms 2.175 ms 3.090 ms

2 212.156.201.180.static.turktelekom.com.tr (212.156.201.180) 21.905 ms 19.413 ms

19.555 ms

3 81.212.71.241.static.turktelekom.com.tr (81.212.71.241) 13.059 ms 13.178 ms 13.232
ms

4 01-adana-xrs-t2-1----31-hatay-t3-1.statik.turktelekom.com.tr (195.175.167.56) 19.527
ms 17.382 ms 17.085 ms

3 35-izmir-xrs-t2-1---01-adana-xrs-t2-1.statik.turktelekom.com.tr (195.175.166.0)
28.367 ms 30.671 ms 31.680 ms

6 * 35-ebgp-izmir-sr12e-k---35-izmir-xrs-t2-1.statik.turktelekom.com.tr (81.212.30.5)
31.320 ms 29.615 ms

7 301-fra-col-1---06-ebgp-ulus-sr12e-k.statik.turktelekom.com.tr (212.156.101.32)
73.037 ms 69.714 ms 68.996 ms

8 62.157.248.1 (62.157.248.1) 868.168 ms * *

9 pd900cb02.dip0.t-ipconnect.de (217.0.203.2) 85.037 ms 79.513 ms *

10 new-york-un.ear3.newyork1.level3.net (4.28.130.118) 160.253 ms

80.156.170.214 (88.150.170.214) 91.782 ms

new-york-un.ear3.newyork1.level3.net (4.28.130.118) 153.203 ms

11 ael2.csl.fra6.de.eth.zayo.com (64.125.26.172) 150.142 ms 146.613 ms

dmzgwb-p2p-extgwc.net.nyu.edu (128.122.254.110) 157.920 ms

ae2.csl.ams17.nl.eth.zayo.com (64.125.29.19) 167.329 ms 167.260 ms

12 nyugwa-ptp-dmzgwb-vl3082.net.nyu.edu (128.122.254.116) 157.550 ms 160.002 ms *

14 * * *

15 * * * wsqdcgwa-vl902.net.nyu.edu (128.122.138) 164.298 ms

16 ae5.cs3.lga5.us.eth.zayo.com (64.125.29.126) 148.268 ms * *

17 * * *

18 * 209.66.118.177.idia-282827-zyo.zip.zayo.com (209.66.118.177) 173.936 ms 167.051 ms

19 * * *

20 * * *

21 nyufw-outside-ngfw-vl3080.net.nyu.edu (128.122.254.116) 163.018 ms * 154.461 ms

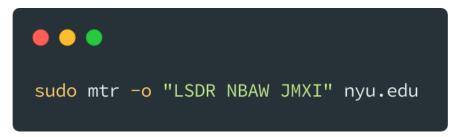
22 * * *

23 wsqdcgwa-vl902.net.nyu.edu (128.122.1.38) 157.344 ms * *
```

9. I reach to the destination much faster, at 11 hops, when tracing 18.31.0.200.

```
traceroute to 18.31.0.200 (18.31.0.200), 64 hops max, 52 byte packets 1 192.168.1.1 (192.168.1.1) 7.753 ms 21.956 ms 13.085 ms
  212.156.201.180.static.turktelekom.com.tr (212.156.201.180) 240.803 ms 148.522 ms
ms
4 01-adana-xrs-t2-2---31-hatay-t3-1.statik.turktelekom.com.tr (212.156.119.103) 21.969
ms 28.330 ms 58.632 ms
5 06-ulus-xrs-t2-2---01-adana-xrs-t2-2.statik.turktelekom.com.tr (81.212.216.58) 34.451
ms 103.286 ms 115.133 ms
   06-ebgp-ulus-sr12e-k---06-ulus-xrs-t2-2.statik.turktelekom.com.tr (81.212.217.121)
154.806 ms 149.470 ms 129.796 ms
 7 305-vie-col-3---06-ebgp-ulus-sr12e-k.statik.turktelekom.com.tr (212.156.140.170)
217.891 ms 162.312 ms
 8 80.157.206.133 (80.157.206.133) 275.066 ms
    4.68.63.201 (4.68.63.201) 180.306 ms
    pd9ef317d.dip0.t-ipconnect.de (217.239.49.125) 147.381 ms
10 80.156.160.134 (80.156.160.134) 198.457 ms
    ae-2-3.bear1.boston1.level3.net (4.69.159.249) 222.319 ms
    80.156.160.134 (80.156.160.134) 145.162 ms
   massachuset.bear1.boston1.level3.net (4.53.48.98) 236.803 ms 283.898 ms
```

- 10. Routing Blackhole is the routing the all network traffic into blackhole and lost it.
 We can use such a system to prevent attacks to our networks like DDoS attack.
 However, normal traffic is also affected such prevention system.
- 11. We can arrange the output fields by using the -o flag. I use following command:



	Packets				Pings							
Host	Loss%	Snt	Drop	Rcv	Last	Best	Avg	Wrst	Jttr .	Javg .	Jmax	Jint
1. 192.168.1.1	3.4%				12.5	2.0	23.7	317.8		17.6	295.	156.
2. 212.156.201.180.static.turktelekom.com.tr				87	19.6	15.2	50.2		21.2	23.7	183.	307.
3. 81.212.71.241.static.turktelekom.com.tr	0.0%			88	13.3	13.1	37.5	301.5	4.6	17.7	144.	188.
4. 01-adana-xrs-t2-131-hatay-t3-1.statik.turktelekom.com.tr	0.0%	88		88	18.4	15.3	37.5	276.3		18.7	164.	162.
5. 35-izmir-xrs-t2-101-adana-xrs-t2-1.statik.turktelekom.com.tr	0.0%	88		88	33.0	29.7		253.2		18.2	183.	129.
6. 35-ebgp-izmir-sr12e-k35-izmir-xrs-t2-1.statik.turktelekom.com.tr	2.3%	88		86		28.9	46.9	210.7	16.6	16.4	167.	218.
7. 301-fra-col-106-ebgp-ulus-sr12e-k.statik.turktelekom.com.tr		88		87	73.4	68.9	85.8	403.2		20.0	330.	179.
8. 62.67.19.245	96.6%		84		80.8	77.3	79.2	80.8				
9. ae-1-3501.ear3.newyork1.level3.net	84.1%			14	164.3	156.6	176.4	354.0		33.1	197.	252.
10. new-york-un.ear3.newyork1.level3.net		88		87	160.1	153.3		549.2	22.5	25.2	302.	245.
11. dmzgwb-p2p-extgwc.net.nyu.edu	81.6%	88	71		157.9	157.0	215.1	479.1	22.9	52.3	322.	443.
12. nyugwa-ptp-dmzgwb-vl3082.net.nyu.edu	1.1%	88		87	159.5	155.5	194.7	487.4		34.8	250.	460.
13. nyufw-outside-ngfw-vl3080.net.nyu.edu	79.3%	88			158.3	155.7	169.2	208.5	50.2	17.3 !	50.2	208.
14. (waiting for reply)												
15. wsqdcgwa-vl902.net.nyu.edu	95.4%		83			159.3	271.1	468.2	86.6	148.	308.	542.
16. (waiting for reply)												

These fields are described in man page of mtr as follows:



12. mtr continuously updates times by polling a remote server and allows us to see how the latency and performance changes over time. On the other hand, traceroute allows us to discover the pathway to a host. When we compare two Wireshark record, we can see that mtr continuously sends probes whilst traceroute sends predetermined number of probes.

Part 2. DV Routing Simulator

Part 2A. Implementation

When constructing Node class, I initialize id from static counter, get cost and neighbors from static arrays of DVSimulator class. myDV is initially equals to cost of that node. I initialize bestPath array based on following conditions: If nodes id or neighbor's id is equal to destination id, then use that id, otherwise get a random node from randomNeighbor method.

```
public Node() {
    this.id = count++;
    this.cost = DVSimulator.cost[this.id];
    this.neighbors = DVSimulator.neighbors[this.id];
    this.myDV = this.cost;

for (int i = 0; i < DVSimulator.NUMNODES; i++) {
        if (i == this.id) {
            this.bestPath[i] = this.id;
        } else if (Node.contains(neighbors, i)) {
            this.bestPath[i] = i;
        } else {
            this.bestPath[i] = randomNeighbor();
        }
    }

notifyNeighbors();
}</pre>
```

For the notifyNeighbors method, I create a new packet for each neighbor that created with current node id as source, neighbor id as destination and current node's myDV array as the dv. Then, I send the packet.

```
public void notifyNeighbors() {
    for (int neighbor : neighbors) {
        Packet packet = new Packet(this.id, neighbor, this.myDV);
        DVSimulator.sendPacket(packet);
    }
}
```

Citation: I get the following method from https://stackoverflow.com/a/34541755. It is simple method for ease of use when searching a value in array.

```
public static boolean contains(final int[] arr, final int key) {
   return Arrays.stream(arr).anyMatch(i -> i == key);
}
```

Part 2B. Optional (Bonus)

I compare for each value in the DV received from neighbor as new_cost and current DV as current_cost. If there is a cheaper path, I update myDV with corresponding cost and update bestPath with corresponding neighbor's id. If there is such a change in current Node's DV I update hasDVChanged value as true. If hasDVChanged value is true, then I notifyNeighbors and increment numUpdates. By doing that, I increment and notify neighbors just once for each packet.

```
public void updateDV(Packet p) {
      int neighbor_id = p.getSource();
      neighborDV[neighbor_id] = p.getDV();
      boolean hasDVChanged = false;
      for (int i = 0; i < myDV.length; i++) {</pre>
          int new_cost = cost[neighbor_id] + neighborDV[neighbor_id][i];
          int current_cost = myDV[i];
          if (new_cost < current_cost) {</pre>
              myDV[i] = new_cost;
              bestPath[i] = neighbor_id;
              hasDVChanged = true;
          }
      if (hasDVChanged) {
          notifyNeighbors();
          numUpdates++;
```