

**CS-201**  
**Fundamental Structures of Computer  
Science I**

**HOMEWORK-2**  
**Mehmet Hasat Serinkan**  
**21901649**

## Question 1

*Algorithm 1:* The time complexity of the first algorithm is  $O(N)$ . There are 2 statements before the loop. The for loop will be executed for  $N+1$  times and the statements in the for loop will be executed for  $N$  times. Lastly, the return statement will be executed for only one time. Therefore, when this function is called, the total amount of execution is  $2+(N+1)+N+N+1 = 3N+4$ . Hence, its upper bound is  $O(3N+4)$ , which is simply  $O(N)$ .

*Algorithm 2:* In all cases, 2 statements will be executed before the loop. There are 3 different cases for algorithm 2:

- **Worst Case:** If the condition is satisfied on the last step of the loop, it would be the worst case. Loop will be executed for  $N$  times, the statements in the loop will run for  $N-1$  times. After that, the “else” part will be executed. There are 3 statements before the loop in the “else” part. Then, the loop will execute for  $N$  times the condition of “ $n \% i$ ” equals  $n$ . The statements in the loop will execute for  $N-1$  times. In the end, it will return the function which will be executed only one time. Therefore, the entire time of execution equals to  $2+N+2(N-1)+3+N+2(N-1)+1 = 6N+2$  which is equal to  $O(N)$ . So, it can be easily seen that the time complexity of the worst case of algorithm 2 is  $O(N)$ .
- **Best Case:** If the loop executes only one time, then it will be the best case. Therefore, the time complexity of the best case of algorithm 2 is  $O(1)$ .
- **Average Case:** Average case equals to all possible case times divided by the number of cases. Since the best case of algorithm 2 is  $O(1)$  and the worst case of algorithm 2 is  $O(N)$ , all possible case times can be simplified as  $1+2+3+4+\dots+N = \frac{N(N+1)}{2}$ , then the total number of cases is  $N$ . Therefore, the average case will be equal to  $O(\frac{(N+1)}{2})$ , which is simplified as  $O(N)$ .

*Algorithm 3:* The base case is  $n=1$  which means when  $n=1$ , function will return. The recurrence relation when  $n$  is even is  $T(N) = T(\frac{N}{2}) + 1$ , and when  $n$  is odd  $T(N) = T(\frac{N-1}{2}) + 1$  which can be simplified as  $T(N) = T(\frac{N}{2}) + 1$ . Therefore, the general recurrence relation is  $T(N) =$

$T(\frac{N}{2}) + 1$ . Now, assume that the function will return when it makes k times execution. Then, the recurrence relation becomes  $T(N) = T(\frac{N}{2^k}) + k$ . To reach the base case,  $2^k$  should be N which means  $k = \log N$ . Then equation becomes  $T(N) = T(1) + \log N = 1 + \log N$ . Its time complexity  $O(1 + \log N)$  can be simplified as  $O(\log N)$ . Therefore, the time complexity of the third algorithm is  $O(\log N)$ .

## Question 2

Specifications of my computer:

- System Manufacturer and Model: MONSTER-TULPAR T7 V19.5
- Processor: Intel® Core™ i7-9750H CPU @ 2.60 GHz
- Installed RAM: 32.0 GB
- System type: 64-bit operating system
- Windows specification: Windows 10 Home Single Language
- Graphic Card: NVIDIA GeForce RTX 2070

## Question 3

Running times of algorithm 1,2, and 3 with different n and p values. Execution times are reported as milliseconds. Value of parameter “a” is 3.

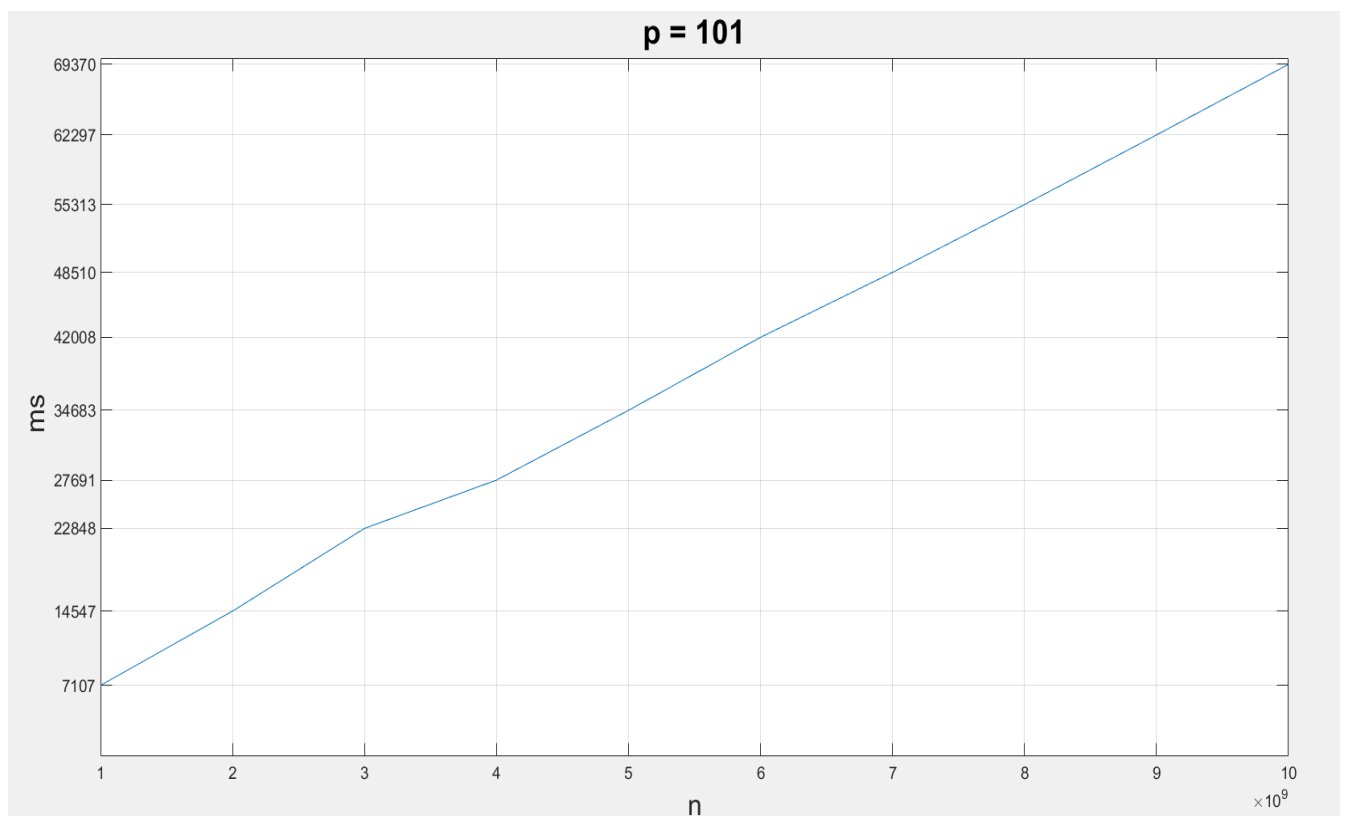
N	Algorithm 1			Algorithm 2			Algorithm 3		
	P = 101	P=1009	P=10007	P = 101	P=1009	P=10007	P = 101	P=1009	P=10007
$10^9$	7107	7041	7170	0.000094	0.001354	0.001026	0.0000356	0.0000351	0.0000380
$2 \times 10^9$	14547	13836	14705	0.000085	0.001549	0.000967	0.0000372	0.0000371	0.0000433
$3 \times 10^9$	22848	20824	21467	0.000099	0.001264	0.000929	0.0000386	0.0000412	0.0000399
$4 \times 10^9$	27691	27840	27761	0.000087	0.001425	0.000828	0.0000399	0.0000418	0.0000389
$5 \times 10^9$	34683	34879	34590	0.000085	0.001578	0.001031	0.0000418	0.0000434	0.0000450
$6 \times 10^9$	42008	41730	41360	0.000119	0.001380	0.000972	0.0000435	0.0000455	0.0000456
$7 \times 10^9$	48510	48360	48515	0.000084	0.001483	0.000902	0.0000450	0.0000458	0.0000461
$8 \times 10^9$	55313	55451	55548	0.000108	0.001282	0.000858	0.0000489	0.0000470	0.0000476
$9 \times 10^9$	62297	62284	62291	0.000093	0.001400	0.001046	0.0000500	0.0000495	0.0000499
$10^{10}$	69370	68724	69230	0.000092	0.001559	0.001041	0.0000505	0.0000515	0.0000519

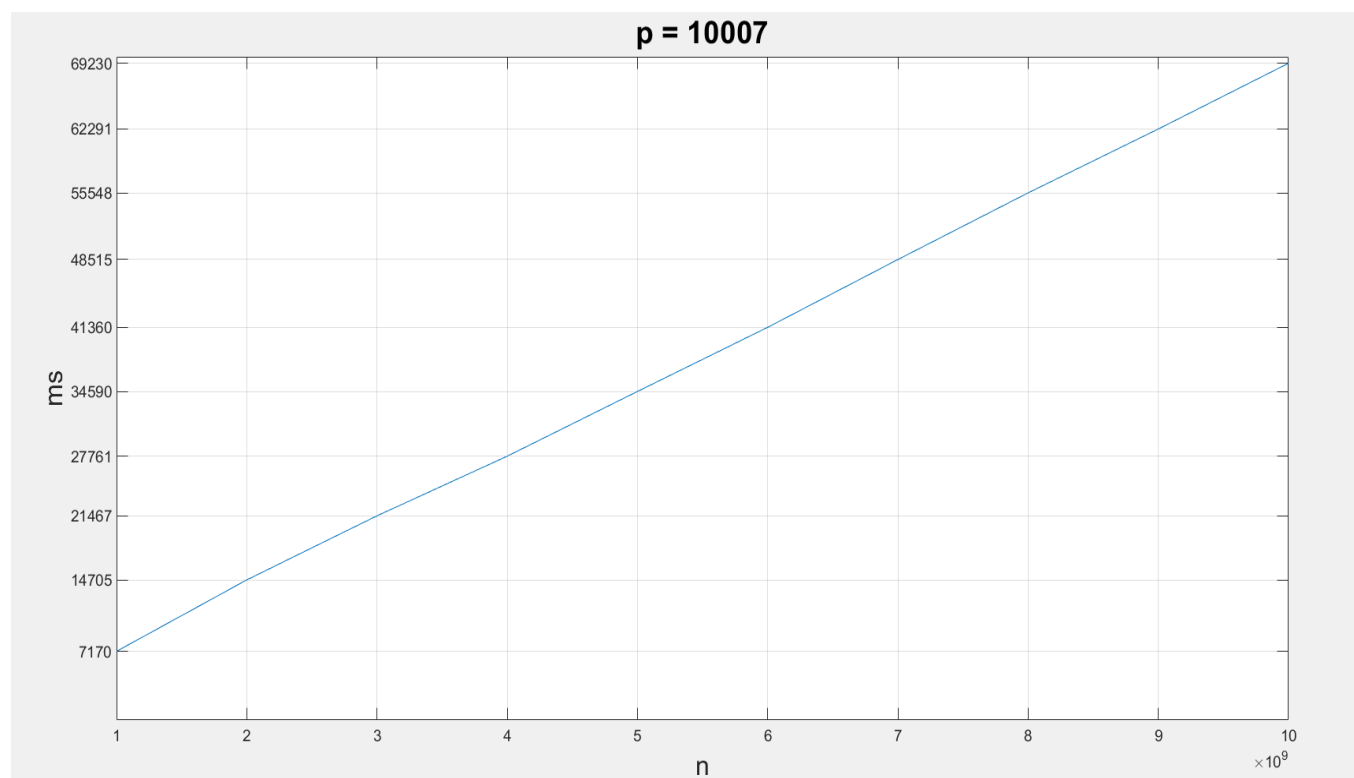
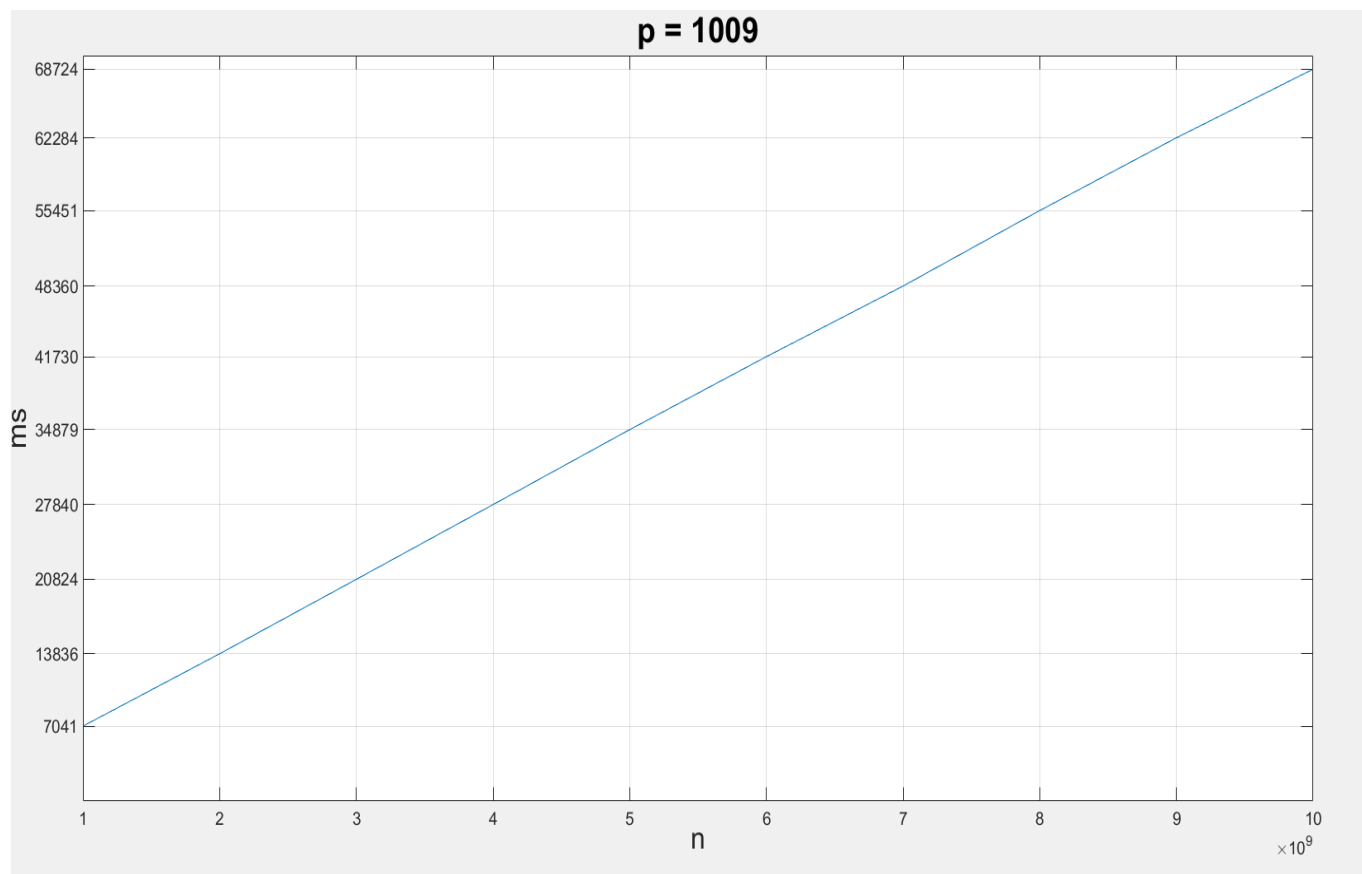
### Comments:

- It can be easily seen that Algorithm 1 is not efficient. It waits almost a minute when  $p$  is a huge number.
- Algorithm 2-3 is much more efficient than Algorithm 1. They execute the function very rapidly.
- The  $p$  values do not crucially affect run times in all algorithms.
- When  $p$  is a small value, it is best to use Algorithm 2.
- The time complexity of Algorithm 2 is almost always  $O(1)$ . As we mentioned before, it is best case for Algorithm 2. The main reason for this is the numbers are huge; hence, it reaches the else part of algorithm 2 almost everytime.
- These values are not very accurate because of non-ideal experimental conditions. The conditions of observations are not the same. For instance, my computer gets heater in Algorithm 1; therefore, the execution time got much larger.

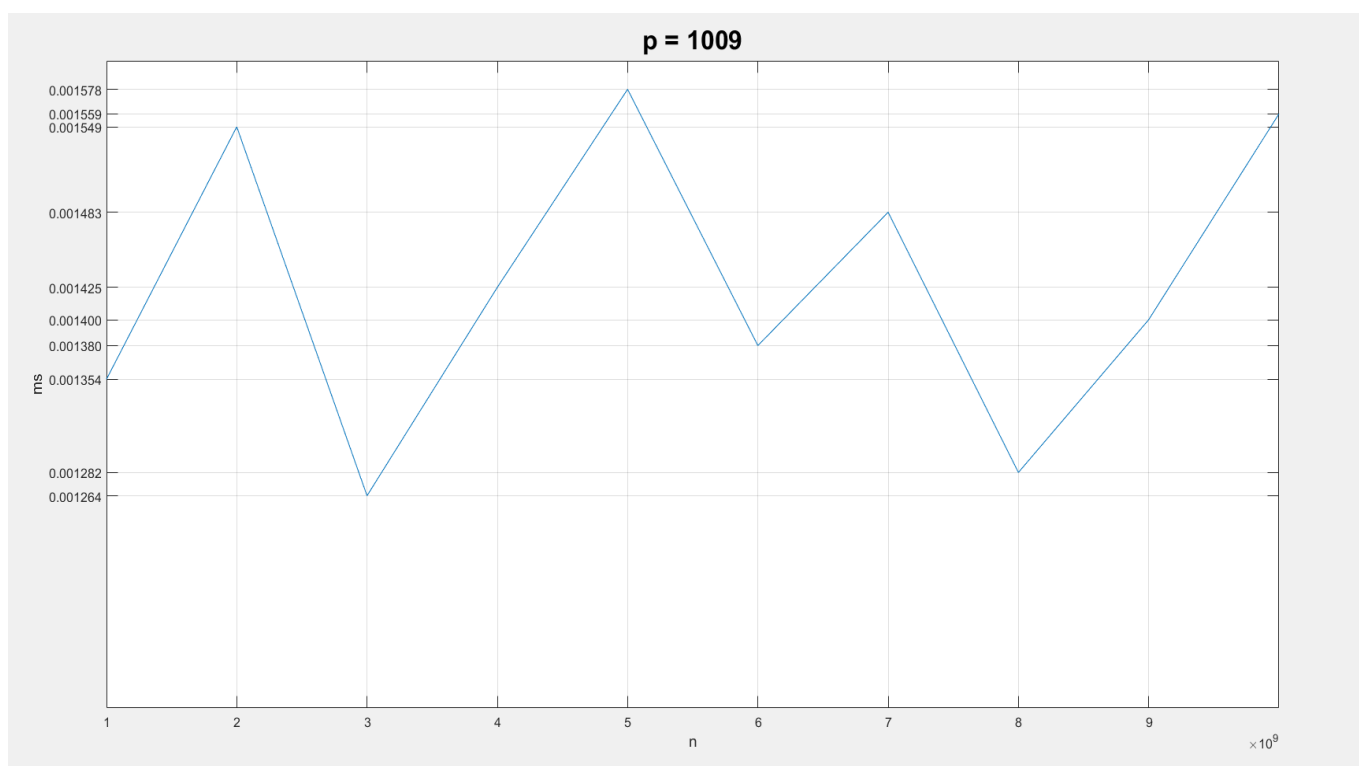
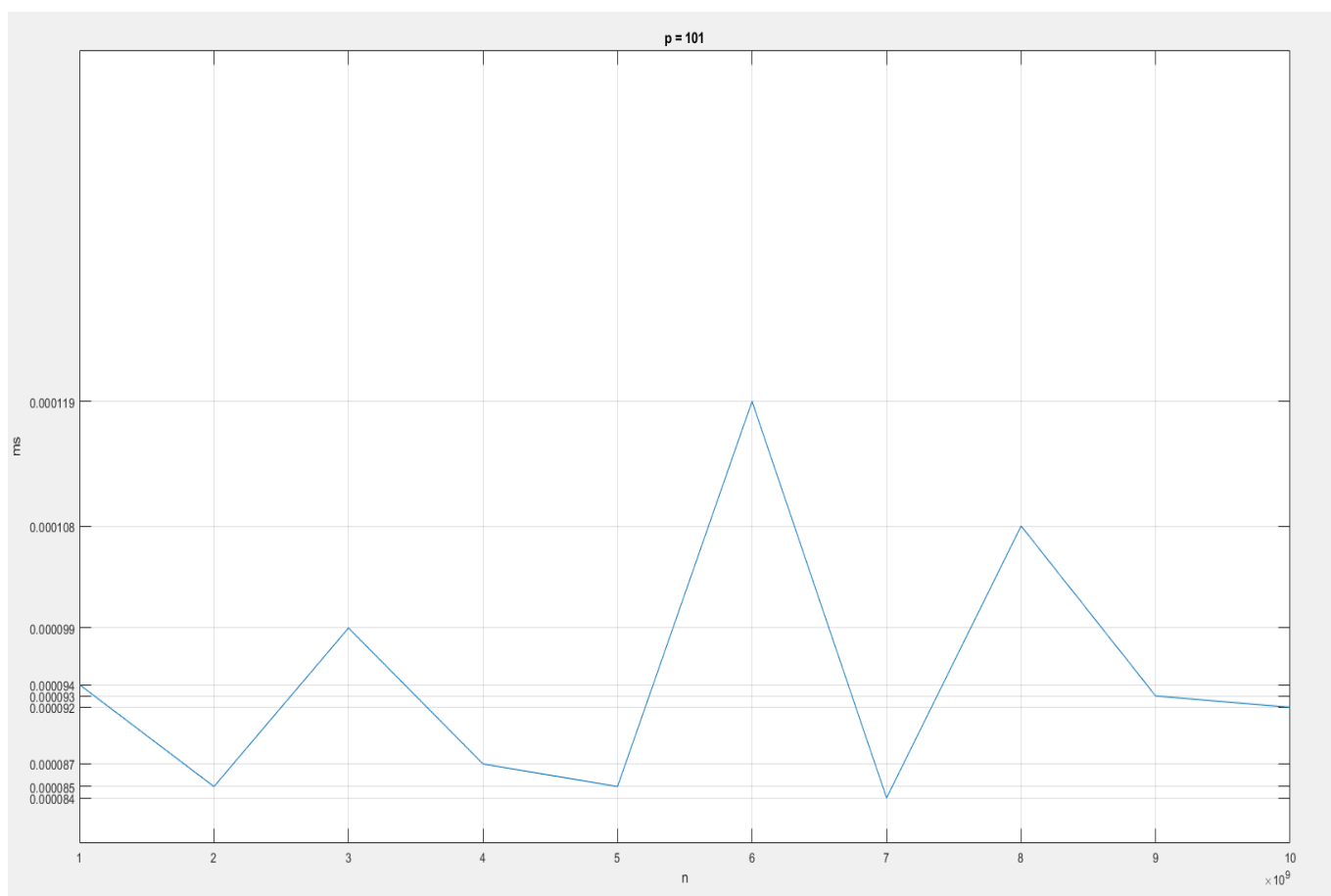
## Question 4

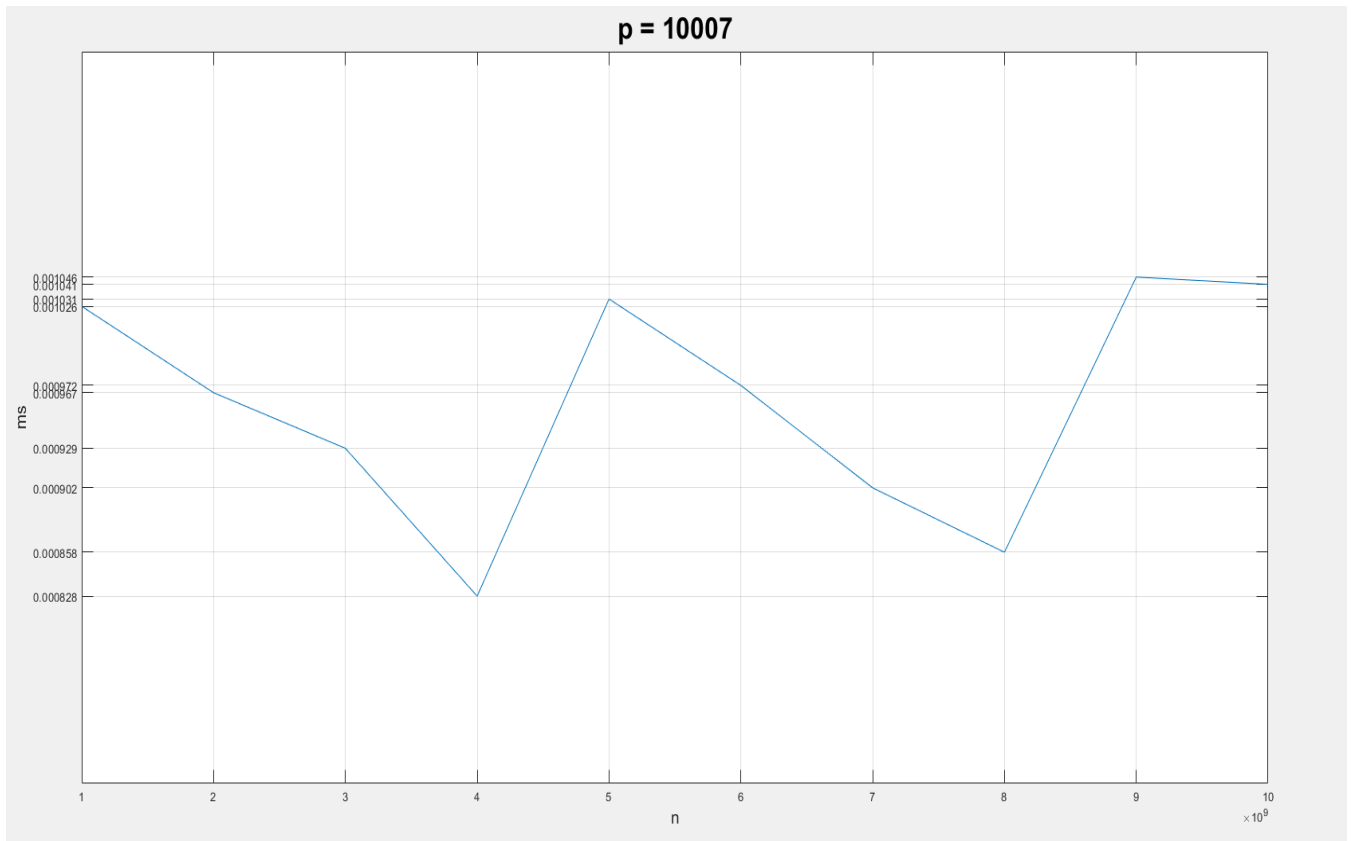
### Plots of Algorithm 1:



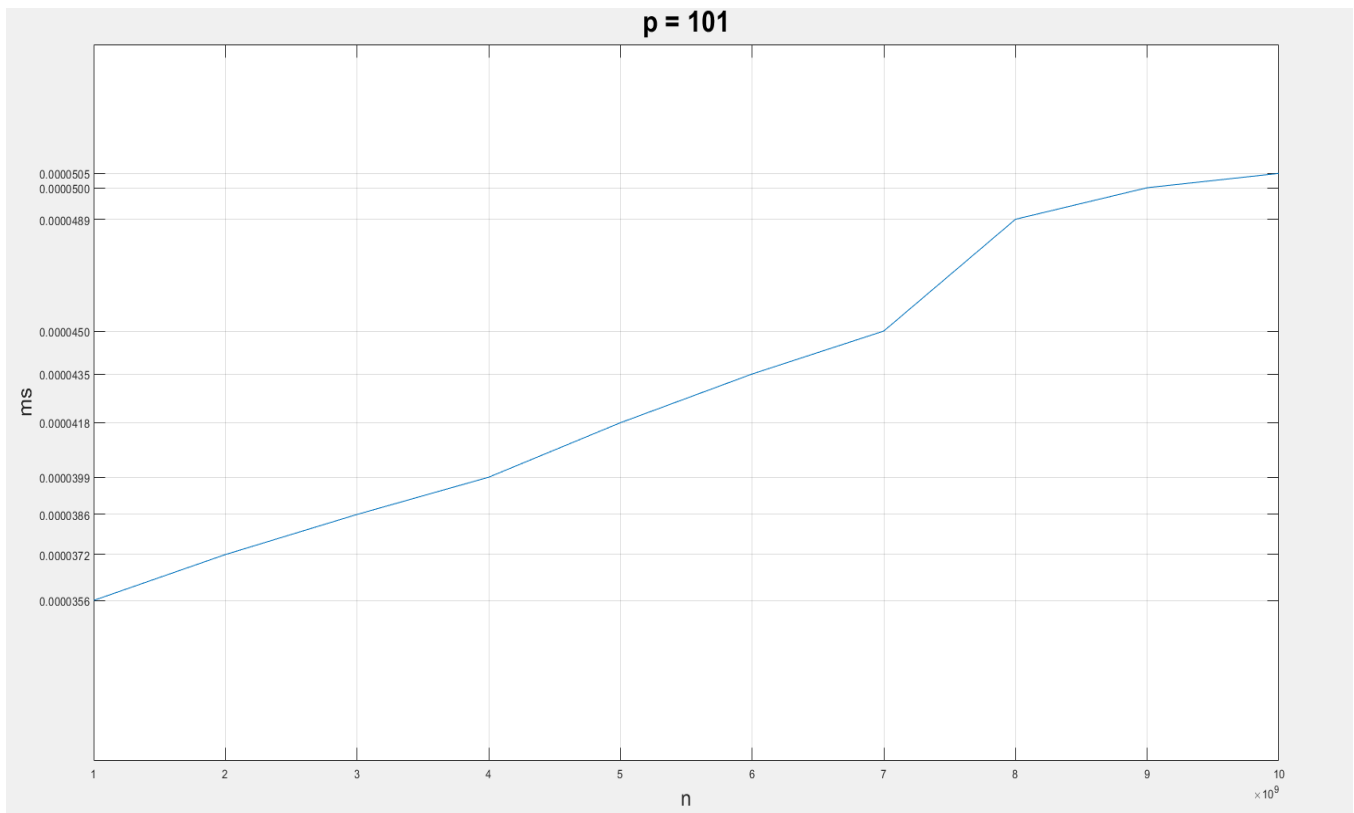


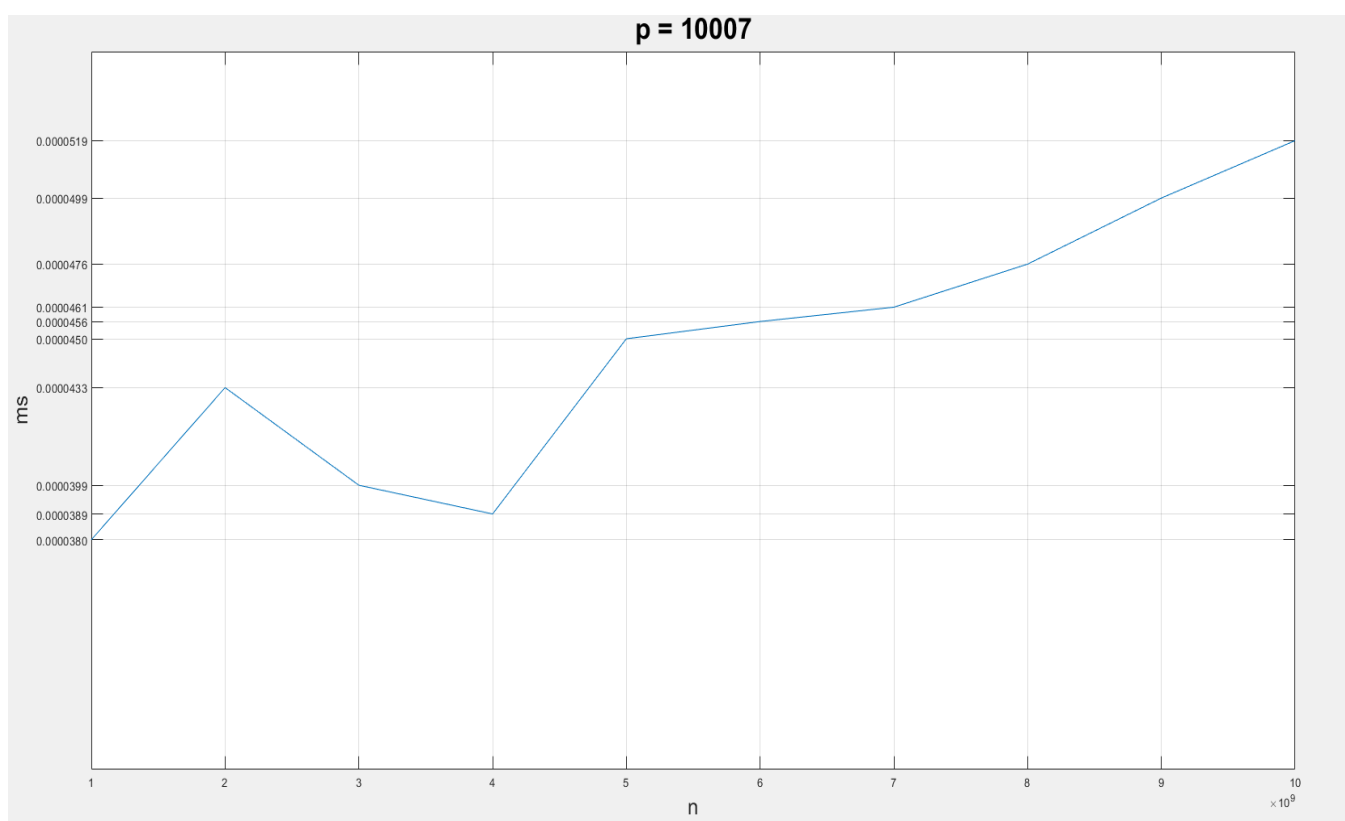
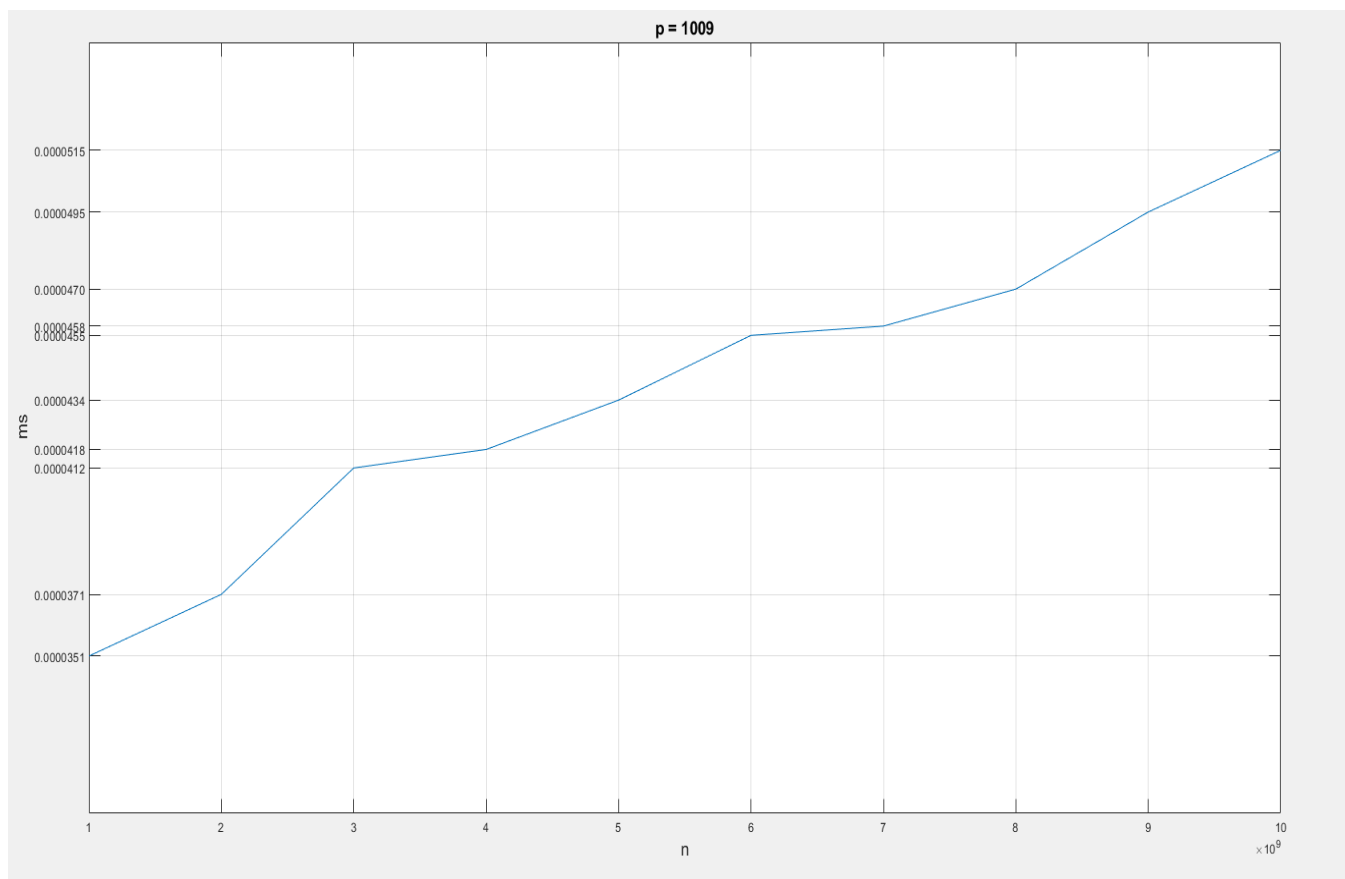
## Plots of Algorithm 2:





Plots of Algorithm 3:







### Discussion:

- Since the numbers are large, it can be easily seen from the graph that Algorithm 1's time complexity is linear, which is  $O(N)$ . If the numbers were smaller, it would be not easy to see the linear relation.
- If there were much more  $n$  values in the plots of Algorithm 2, it would easier to see it has  $O(1)$  time complexity. Moreover, if the range of  $y$ -axis were larger, it would make easier as well. However, to have better observation, less range of  $y$ -axis and fewer  $n$  values were used. Therefore, the plots of Algorithm 2 are not very accurate. By looking only at the plots, it is very hard to see Algorithm 2 has  $O(1)$  time complexity.
- Approximately every plot of Algorithm 3 has increased when  $n$  gets greater. However, the third plot is not defining the time complexity of Algorithm 3 explicitly. It gets less run time when  $n$  is converted from  $2 \times 10^9$  to  $3 \times 10^9$ . Nevertheless, if the  $n$  values were much more, then plots would be more accurate. It would be easier to see Algorithm 3 has  $O(\log N)$  time complexity.
- In brief, the plots would show the time complexities of algorithms more accurately if we had used much more  $n$  values. Moreover, with a better computer and more ideal conditions, the time complexities would be easier to see from the plots.