

CS2023 - Data Structures and Algorithms

In-class Lab Exercise

Week 11

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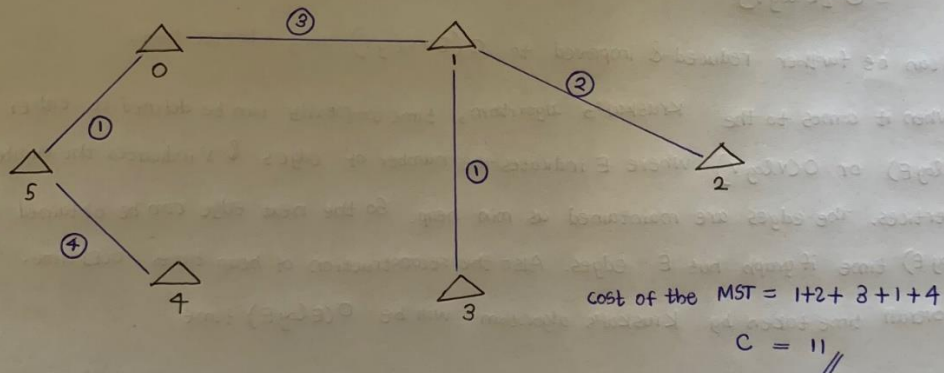
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Question 1

Q1) Adjacency matrix for the given graph.

	0	1	2	3	4	5
0	0	3	0	0	0	1
1	3	0	2	1	10	0
2	0	2	0	3	0	5
3	0	1	3	0	5	0
4	0	10	0	5	0	4
5	1	0	5	0	4	0

Q2) Starting from Node 3



Q3) Yes, both MSTs have the same edges with same cost. If each edge of the tree has a distinct weight then there will only be one, unique minimum spanning tree.

3)

```
"D:\Semester 4\CS2023 - Data Structures and Algorithms\In Class labs\In-class Lab 11\MST_with_Prim's_Algorithm\bin\Debug\MST_with_Prim's_Algorithm.exe"
Edge  Weight
0 - 1   3
1 - 2   2
1 - 3   1
5 - 4   4
0 - 5   1

Process returned 0 (0x0)   execution time : 0.014 s
Press any key to continue.
```

Q5) when it comes to the prim's Algorithm, if adjacency list is used to represent the graph, using BFS all the vertices can be traversed in $O(V+E)$ time. Then we can use a min heap for storing the vertices not yet included in the MST.

This min heap can be used as a priority queue to get the minimum weight edge. In min heap, there are operations like extracting minimum element & decreasing the key values which takes $O(\log V)$ time.

So overall time complexity can be given as,

$$\begin{aligned} &= O(E+V) \cdot O(\log V) \\ &= O[(E+V) \log V] \\ &= O[E \log V] \end{aligned}$$

This can be further reduced & improved to $O(E+V \log V)$ using Fibonacci heap.

But when it comes to the Kruskal's algorithm, time complexity can be defined as either $O(E \log E)$ or $O(V \log V)$ where E indicates the number of edges & V indicates the number of vertices. The edges are maintained as min heap. So the next edge can be obtained in $O(\log E)$ time if graph has E edges. Also the reconstruction of heap takes $O(E)$ time.

So overall time taken by Kruskal's algorithm will be $O(E \log E)$ time.