

Midterm 2 Practice

Due No due date **Points** 30 **Questions** 22 **Available** Mar 11 at 10am - Mar 18 at 3pm 7 days
Time Limit None **Allowed Attempts** 7

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Attempt History

	Attempt	Time	Score
LATEST	Attempt 1	1,334 minutes	4.67 out of 30 *

* Some questions not yet graded

Submitted Mar 16 at 12:15pm

Question 1

0.33 / 1 pts

Select all of the **true** statements about the query processor/optimizer.

Correct!

☒ A SQL query is converted from a string to a parse tree by the parser component.

☐

A logical query tree has relational algebra operators and has selected algorithms to use to implement those operations.

Correct!

☒ The parser will generate a syntax error if the SQL query is not valid according to the SQL grammar.

You Answered

☒ The preprocessor validates table and field names and runs before the parser builds the parse tree.

Correct Answer

☐ A rule in a grammar may be applied more than once when parsing a query.

Question 2

1 / 1 pts

Select all of the **true** statements about heuristic optimization.

☐ Heuristic optimization rules always improve query plan performance.

Correct!

☒ Joins are commutative and associative.

Correct!

☒ Selection and cross-product operators can be combined into a join.

Correct!

☒ Conjunctive selections can be separated into a sequence of selection operations.

Correct!

☒ It is better to perform selections closer to the bottom (leaf nodes) of the query tree.

Question 3**0 / 1 pts**

Select the **true** statement about optimization.

Correct Answer**You Answered**

- ☐ A canonical logical query tree shows the join order in the query plan.
- ☐ Cost-based optimization uses data statistics to estimate operator costs.
- ☒ Heuristic and cost-based optimization are not used together.
- ☐ The selectivity is the fraction of tuples eliminated by a selection operator.
- ☐ The size of a join result is never larger than the size of the largest relation being joined.

Question 4**0 / 1 pts**

Select the **true** statement about transactions.

You Answered**Correct Answer**

- ☐ A transaction is an atomic program that always executes completely.
- ☒ A consistent database always accurately reflects real-world reality.
- ☐ A transaction can be in both aborted and committed states during its lifetime.
- ☐ A schedule may interleave (reorder) operations within a transaction.
- ☐ A schedule can be used to detect concurrency issues but does not prevent them.

Question 5**0 / 1 pts**

Select the **true** statement about serializability and schedule properties.

Correct Answer

- ☐ Two operations conflict if they use the same data item and one of them is a write.
- ☐ A schedule is conflict serializable if its precedence graph contains a cycle.
- ☐ A conflict serializable schedule is made equivalent to a serial schedule by swapping conflicting operations.

You Answered

- ☐ A recoverable schedule is always cascade-free.
- ☒ A strict schedule guarantees that no committed data item is overwritten.

Use the SQL statement below to answer the following questions.

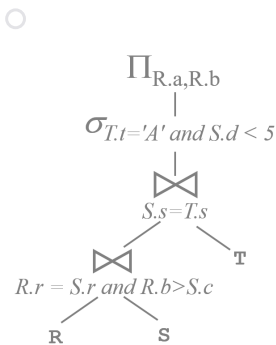
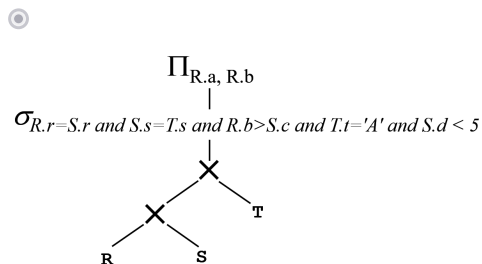
```
SELECT R.a, R.b
FROM R, S, T
WHERE R.r = S.r and S.s = T.s and R.b > S.c and T.t = 'A' and S.d < 5
```

Question 6

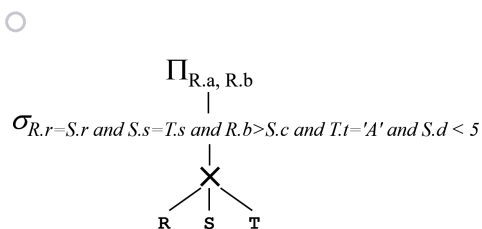
0 / 2 pts

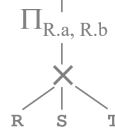
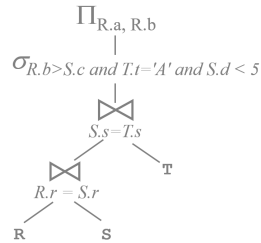
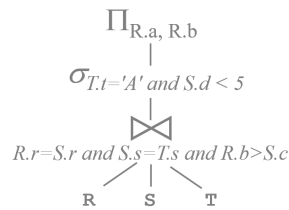
Select a canonical, logical query tree for the given SQL statement. Do NOT apply any optimization.

You Answered



Correct Answer

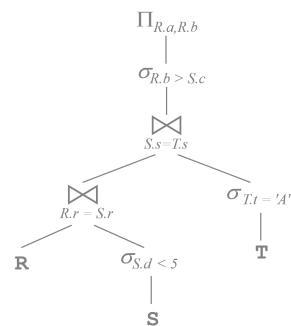
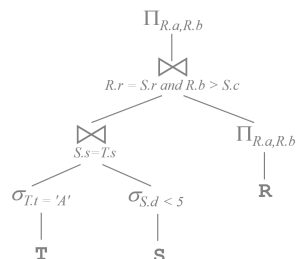


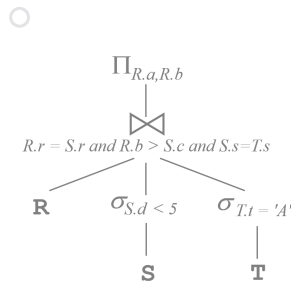
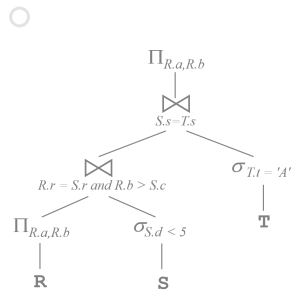
$$\sigma_{R.r=S.r \text{ and } S.s=T.s \text{ and } R.b>S.c \text{ and } T.t='A' \text{ and } S.d < 5}$$

☐

☐


Question 7

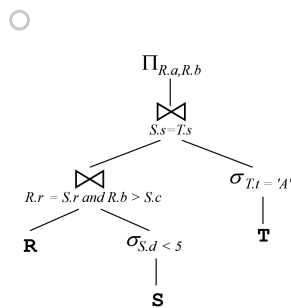
0 / 3 pts

Select the best, optimized logical query tree.

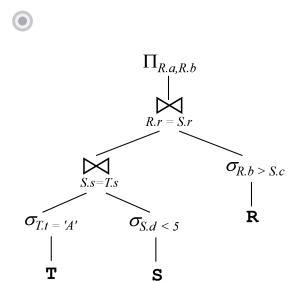
☐

☐




Correct Answer



You Answered



Given the following statistics and logical query tree, perform the requested calculations and physical optimization. Assume no spanning is used. **Answer the questions by giving the number of tuples and the number of blocks.**

Statistics:

A(a,d,e,f,g) - T(A)=30000 V(A,a)=10000 V(A,d)=1000

B(b,a,h) - T(B)=100000 V(B,a)=50000

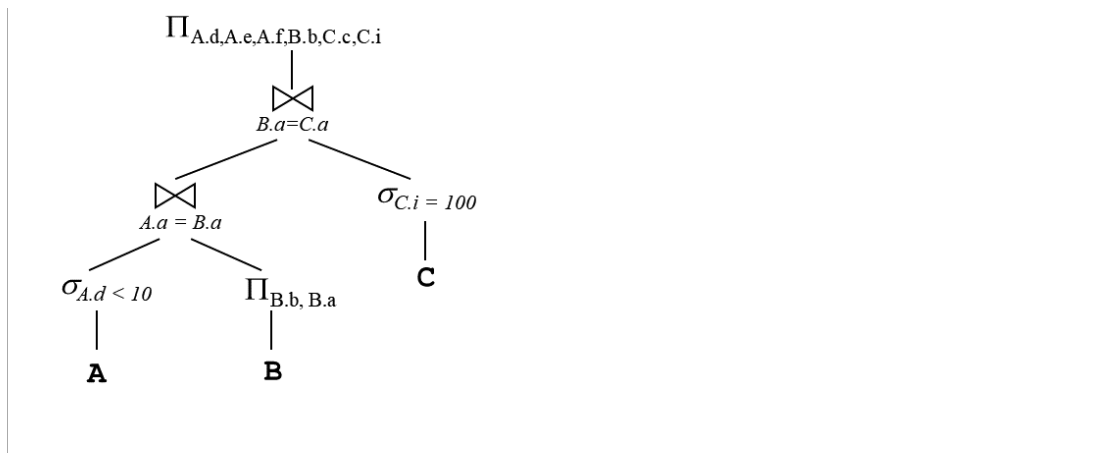
C(c,a,i) - T(C)=50000 V(C,a)=5000 V(C,i)=250

Block size is 1000 bytes with no spanning.

Attribute sizes: a=b=c=i=4 bytes ; h=1000 bytes ; d=6 bytes ; e=50 bytes ; f=g=20 bytes

There are indexes on A.a, A.d, B.a, C.a, and C.i.

Maximum memory size per operator is 100 blocks.



Unanswered

Question 8

0 / 1 pts

Calculate size of: $\sigma_{A.d < 10} (A)$ In tuples: In blocks:

Answer 1:

You Answered

(You left this blank)

Correct Answer

10000

Answer 2:

You Answered

(You left this blank)

Correct Answer

1000

Unanswered

Question 9

0 / 1 pts

Calculate size of: $\sigma_{C.i = 100} (C)$ In tuples: In blocks:

Answer 1:

You Answered

(You left this blank)

Correct Answer 200

Answer 2:

You Answered (You left this blank)

Correct Answer 3

Unanswered

Question 10

0 / 1 pts

Calculate size of: $\Pi_{B.b, B.a}(B)$ In tuples: In blocks:

Answer 1:

You Answered (You left this blank)

Correct Answer 100000

Answer 2:

You Answered (You left this blank)

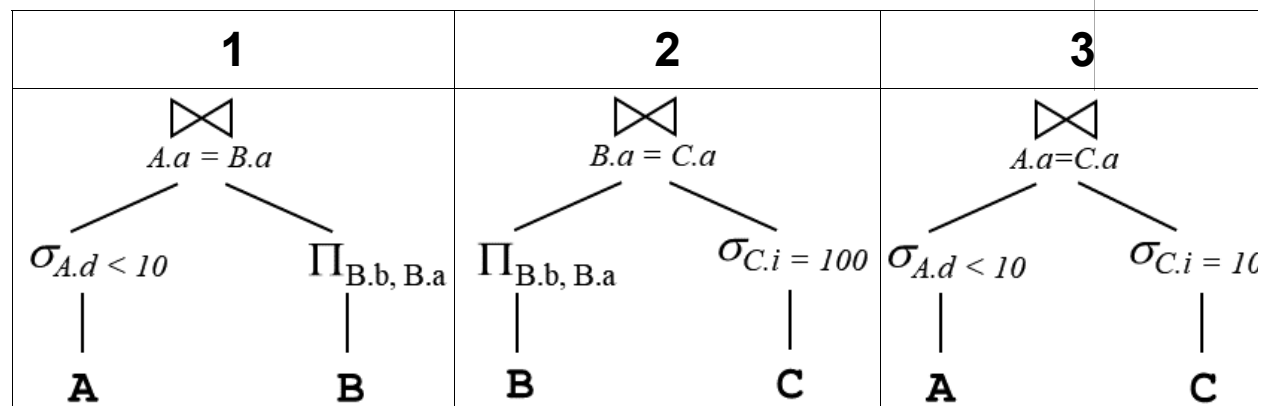
Correct Answer 800

Unanswered

Question 11

0 / 3 pts

Calculate the following three join sizes in tuples and blocks:



1) Join of A and B

In tuples: In blocks: **2) Join of B and C**In tuples: In blocks: **3) Join of A and C**In tuples: In blocks: **Answer 1:****You Answered**

(You left this blank)

Correct Answer

20000

Answer 2:**You Answered**

(You left this blank)

Correct Answer

2222

Answer 3:**You Answered**

(You left this blank)

Correct Answer

400

Answer 4:**You Answered**

(You left this blank)

Correct Answer

8

Correct Answer

7

Answer 5:**You Answered**

(You left this blank)

Correct Answer

200

Answer 6:**You Answered**

(You left this blank)

Correct Answer 25

Correct Answer 23

Question 12

1 / 1 pts

Apply cost-based optimization to answer the next few questions.

What is the best join order?

☐ A \bowtie B \bowtie C

Correct!

☒ A \bowtie (B \bowtie C)

☐ A \bowtie C \bowtie B

☐ B \bowtie (A \bowtie C)

☐ C \bowtie A \bowtie B

Question 13

0 / 1 pts

What join algorithms should be used?

☐ first join is hash join, second join is hash join

You Answered

☒ first join is hash join, second join is one-pass join

Correct Answer

☐ first join is one-pass join, second join is one-pass join

☐ first join is sort join, second join is sort join

☐ first join is one-pass join, second join is hash join

Question 14

0 / 1 pts

What is the best way to implement $\sigma_{C.i = 100}(C)$?

☐ table scan

Correct Answer

☐ index scan

You Answered

☒ one-pass join☐ operator is eliminated in physical tree

Question 15

0 / 1 pts

Consider the optimized, physical query tree. Select all of the **true** statements.

You Answered

☒ The operator $\Pi_{B.b, B.a}(B)$ was eliminated after cost-based optimization as it is not improve cost.

You Answered

☒ There is a selection operator that is above a join.☐ The selection $\sigma_{A.d < 10}(A)$ should use an index scan.

Correct Answer

☐ The join order is different than the original query plan (given at start of question just above Q8).

Correct Answer

☐ A projection operator will not use an index scan.

Unanswered

Question 16

0 / 5 pts

Determine if the following schedule is conflict serializable by drawing the precedence graph.

$r1(A); r2(A); w2(A); r3(B); w3(B); r4(C); w4(A); w2(B); r1(C); w1(D); w2(D); r4(A); w3(A)$

Answer the questions about your precedence graph.

nodes:

edges with label A:

edges with label B:

edges with label C:

edges with label D:

Is the schedule serializable (y/n)?

Enter # of nodes in a cycle (put 0 if no cycle):

If a cycle exists, is T2 in the cycle (y/n)?

Is there an edge T1-> T4 with label A (y/n)?

Is there an edge T3-> T1 with label A (y/n)?

Answer 1:

You Answered

(You left this blank)

Correct Answer

4

Answer 2:

You Answered

(You left this blank)

Correct Answer

6

Answer 3:

You Answered

(You left this blank)

Correct Answer

1

Answer 4:

You Answered

(You left this blank)

Correct Answer

0

Answer 5:

You Answered

(You left this blank)

Correct Answer

1

Answer 6:

You Answered

(You left this blank)

Correct Answer

n

Correct Answer

no

Answer 7:

You Answered

(You left this blank)

Correct Answer

2

Answer 8:

You Answered

(You left this blank)

Correct Answer

y

Correct Answer yes

Answer 9:

You Answered (You left this blank)

Correct Answer y

Correct Answer yes

Answer 10:

You Answered (You left this blank)

Correct Answer n

Correct Answer no

Answer the following questions on transaction schedules and schedule properties given these transactions:

```
T1: r1(A);r1(B);w1(A);c1;  
T2: r2(B);r2(A);w2(A);w2(B);c2;  
T3: r3(C);w3(B);c3;
```

Question 17

0 / 1 pts

TRUE or FALSE: This schedule is conflict-serializable and recoverable:

```
r3(C);w3(B);r2(B);r2(A);w2(A);c3;r1(A);w2(B);r1(B);w1(A);c1;c2;
```

You Answered ☒ True

Correct Answer ☐ False

Question 18

1 / 1 pts

TRUE or FALSE: This schedule is cascade-free but not strict:

```
r2(B);r3(C);r2(A);w2(A);w2(B);w3(B);c2;r1(A);c3;r1(B);w1(A);c1;
```

Correct! ☒ True

☐ False

(cascade free as always reads committed values, not strict as w3(B) overwrites w2(B))

Question 19

1 / 1 pts

TRUE or FALSE: It is possible to create a schedule that is view serializable but not conflict serializable from these transactions.

Correct!

☒ True

☐ False

r1(A); r1(B); r3(C); w3(B); r2(B); w1(B); r2(A); w2(A); w2(B);

(Not conflict serializable as cycle between T1 and T3 on B, view serializable to T1, T3, T2)

Question 20

0 / 1 pts

Select **ALL** of the properties of the following schedule:

r3(C); w3(B); r2(B); r2(A); w2(A); c3; r1(A); w2(B); r1(B); w1(A); c2; c1;

You Answered

☒ serial

Correct!

☒ conflict serializable

Correct!

☒ recoverable

You Answered

☒ cascade-free

☐ strict

Question 21

0.33 / 1 pts

Select **ALL** of the properties of the following schedule:

`r1(A); r1(B); r3(C); w3(B); r2(B); w1(B); c1; r2(A); c3; w2(A); w2(B); c2;`

Correct!

☒ non-serial

You Answered

☒ conflict serializable

Correct!

☒ view serializable

Correct Answer

☐ recoverable

☐ cascade-free

Unanswered

Question 22

Not yet graded / 0 pts

NOTE: This question shows as 0 marks but is worth 3 bonus marks.

Using the GitHub repository: <https://classroom.github.com/a/CbuniShh>
(<https://classroom.github.com/a/CbuniShh>)

Modify the file `src/textdb/question/PracticeQuestion.java`. Test with `src/textdb/junit/TestPracticeQueryTree.java`.

Write an optimized query plan for the SQL query:

```
SELECT n_name, c_custkey, c_name, o_orderkey, o_totalprice
FROM Nation AS N, Customer AS C, Orders AS O
WHERE N.n_nationkey = C.c_nationkey AND O.o_custkey = C.c_custkey
      AND n_name = 'CANADA' AND o_totalprice > 410000
```

Assumptions:

- T(Nation) = 25, T(Customer) = 15000, T(Orders) = 150000
- There is index on c_nationkey. No index on o_totalprice.
- Assume a selection above a TextFileScan will use an index if possible.
- Size of SELECT * FROM Orders WHERE o_totalprice > 410000 is 70.
- Only using BlockNestedLoopJoin. Smallest table is left input (input #0).

You will get 1 bonus mark for passing the first test that produces a CORRECT plan. You will get 2 more bonus marks for passing the second test that checks for the optimal query plan.

When done, commit/push your repository. Submit the URL of the repository on Canvas.

Your Answer: