

10_Transactions

Due Mar 11 at 3pm **Points** 9 **Questions** 9 **Time Limit** None **Allowed Attempts** 7

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Attempt History

	Attempt	Time	Score
KEPT	Attempt 3	1 minute	9 out of 9
LATEST	Attempt 3	1 minute	9 out of 9
	Attempt 2	less than 1 minute	7 out of 9
	Attempt 1	2 minutes	6 out of 9

⚠️ Answers will be shown after your last attempt

Score for this attempt: **9** out of 9

Submitted Mar 18 at 8:19am

This attempt took 1 minute.

Question 1

1 / 1 pts

Two transactions running at the same time can see each other's updates. What ACID property is violated?

- ☐ atomicity
- ☐ consistency
- ☒ isolation
- ☐ durability
- ☐ none of them

Question 2**1 / 1 pts**

A company stores a customer's address in the database. The customer moves and does not tell the company to update its database. What ACID property is violated?

- ☐ atomicity
- ☐ consistency
- ☐ isolation
- ☐ durability
- ☒ none of them

Even though the database is not consistent with the real-world, it is internally consistent.

Question 3**1 / 1 pts**

Is it possible for a transaction to be in the aborted and committed states at different times during its lifetime?

☐ yes

☒ no

Question 4**1 / 1 pts**

Is the schedule valid for the two transactions below?

Schedule:

T_1	T_2
<code>read(A,t)</code> <code>read(B,o)</code> <code>write(A,t)</code> <code>write(B,o)</code>	 <code>read(A,t)</code> <code>write(A,t)</code> <code>read(B,o)</code> <code>write(B,o)</code>

Transaction T1:

`read(A,t)`
`write(A,t)`
`read(B,o)`
`write(B,o)`

Transaction T2:

`read(A,t)`
`write(A,t)`
`read(B,o)`
`write(B,o)`

☐ yes

☒ no

Cannot reorder operations within a transaction in a schedule.

Question 5

1 / 1 pts

What consistency issue does this schedule have?

T_1	T_2
read(<u>A</u> , <u>t</u>)	
write(<u>A</u> , <u>t</u>)	read (<u>A</u> , <u>t</u>)
write(<u>B</u> , 10)	
write(<u>C</u> , <u>t</u>)	read(<u>B</u> , <u>u</u>)
	write(<u>C</u> , <u>t+u</u>)

- ☐ lost update
- ☐ dirty read
- ☐ incorrect summary
- ☐ none
- ☒ more than one

Lost update as T1 write on C is lost.

Incorrect summary as T2 computes sum with old value of A but updated value of B from T1.

Dirty read (even though T1 did not abort) as T2 read a value that was not committed.

Question 6

1 / 1 pts

Is this schedule conflict serializable?

T_1	T_2
read(A)	write(A)
read(B)	write(B)
read(C)	read(C)
write(C)	

☐ yes

☒ no

Conflict on C. Otherwise could re-organize so T2 comes after T1.

Question 7

1 / 1 pts

How many of the following statements are true?

- i) Every serial schedule is a strict schedule.
- ii) A serializable schedule may not be recoverable.
- iii) Every cascade-free schedule is also a strict schedule.
- iv) There are more recoverable schedules than cascade-free schedules.

☐ 0

☐ 1

☐ 2

☒ 3

True statements: I, ii, and iv

☐ 4

Question 8

1 / 1 pts

Is this schedule recoverable?

T_8	T_9
read(A) write(A)	read(A) commit
read(B) commit	

☐ yes

☒ no

T9 reads from T8 and commits before it.

Question 9

1 / 1 pts

Is this schedule cascade-free?

T_8	T_9
read(A) write(A)	
read(B) commit	read(B)
	commit

☒ yes

T9 reads B which is a committed value. Note that it would not be cascade-free if T9 did read(A).

☐ no

Quiz Score: **9** out of 9