

# Chapter 3

## Data Link Layer

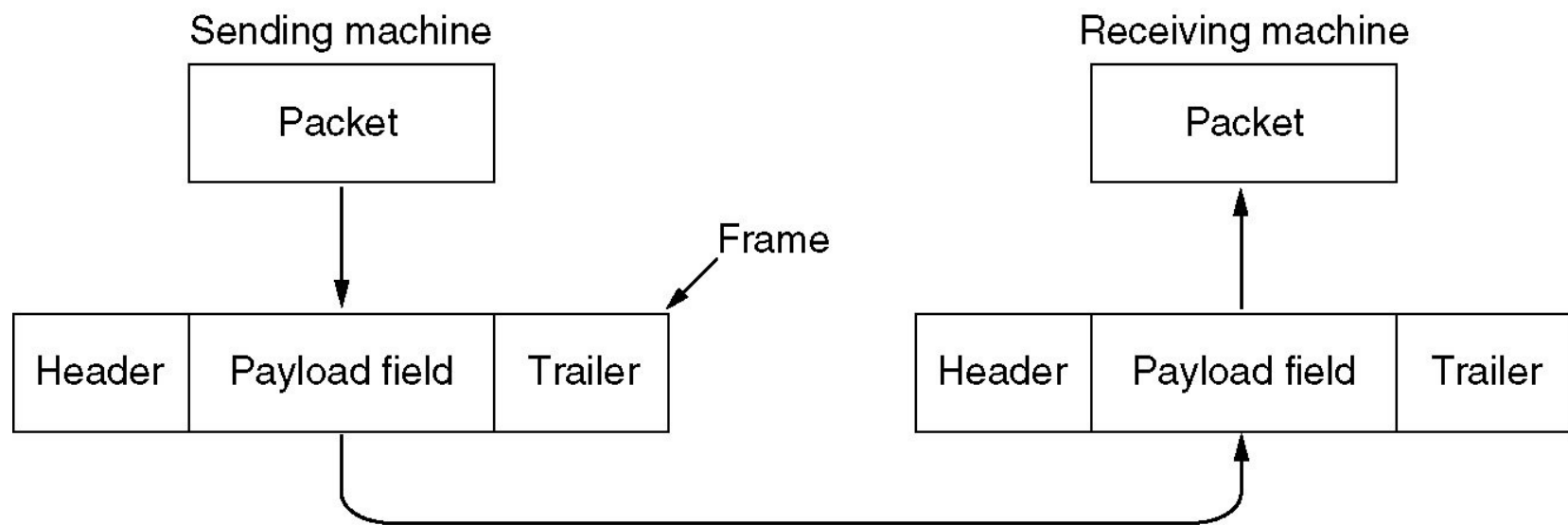
# Data Link Layer Design Issues

- Services Provided to the Network Layer
- Framing
- Error Control
- Flow Control

# Functions of the Data Link Layer

- Provide service interface to the network layer
- Dealing with transmission errors
- Regulating data flow
  - Slow receivers not swamped by fast senders

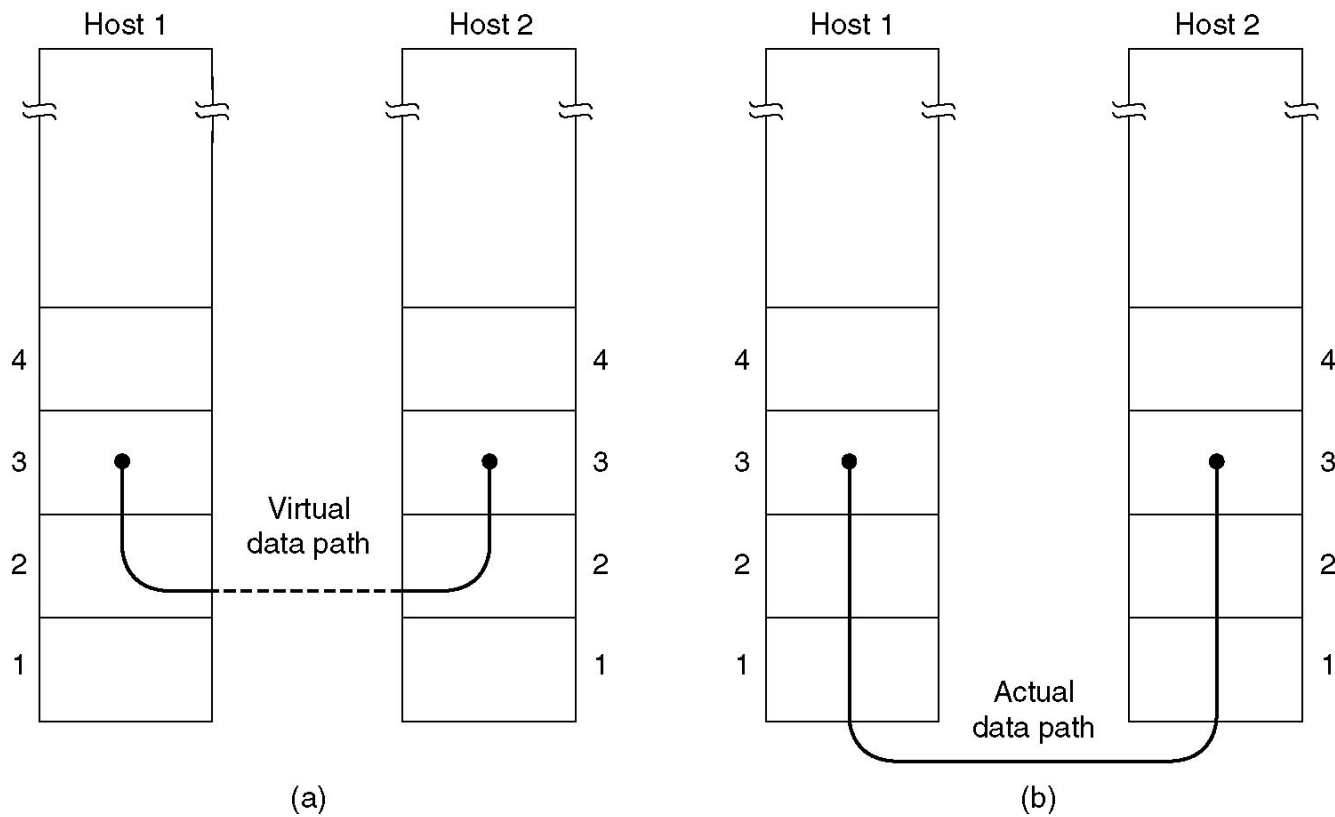
# Relationship between packets and frames.



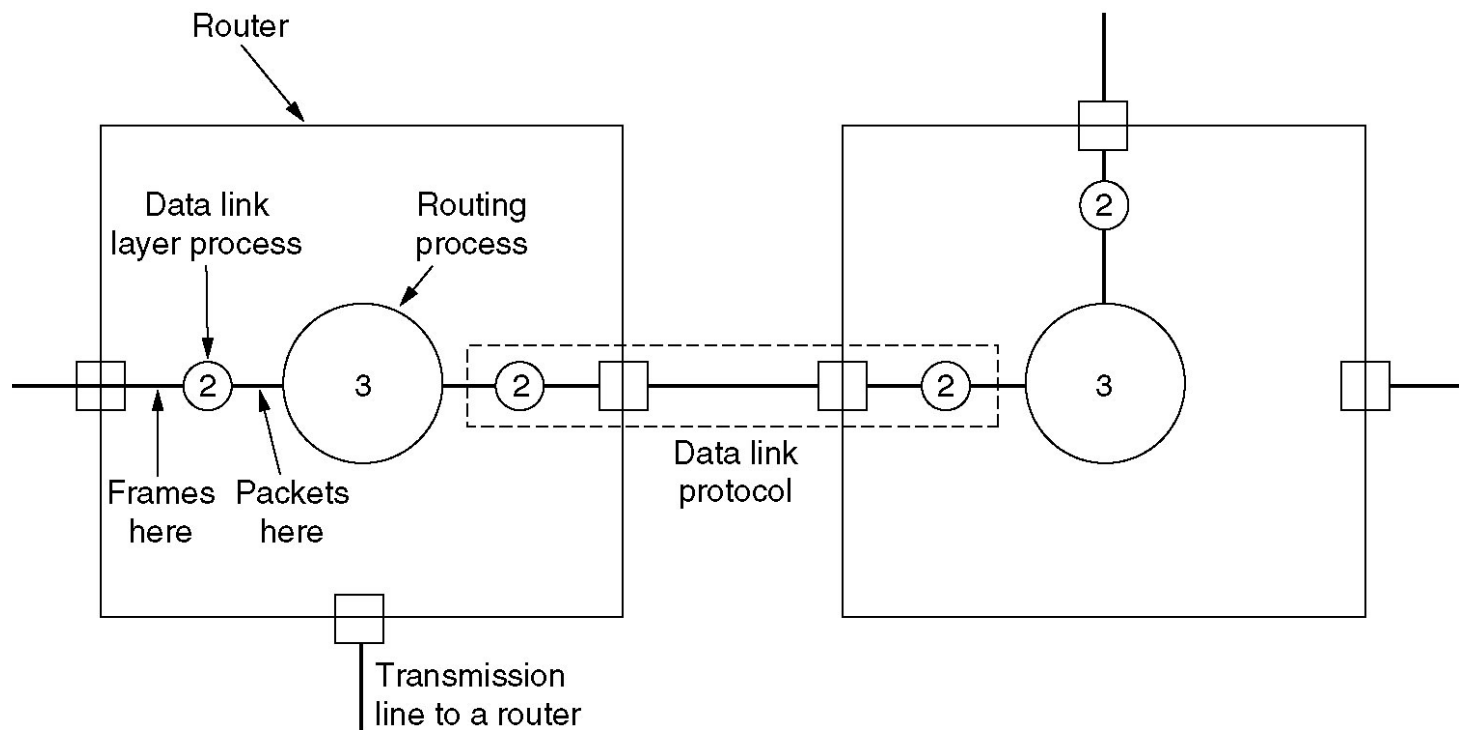
# Service Provided to the Network Layer

- Transmission of network layer packets to destination
- Type of data transmission
  - No Ack No Connection
    - No reliability
  - Ack but No Connection
    - Reliable but duplication of packets
  - Ack and Connection
    - Reliable data transmission and the no duplication
    - Three phases of data transmission

- (a) Virtual communication.
- (b) Actual communication.



# Placement of the data link protocol.



# Framing

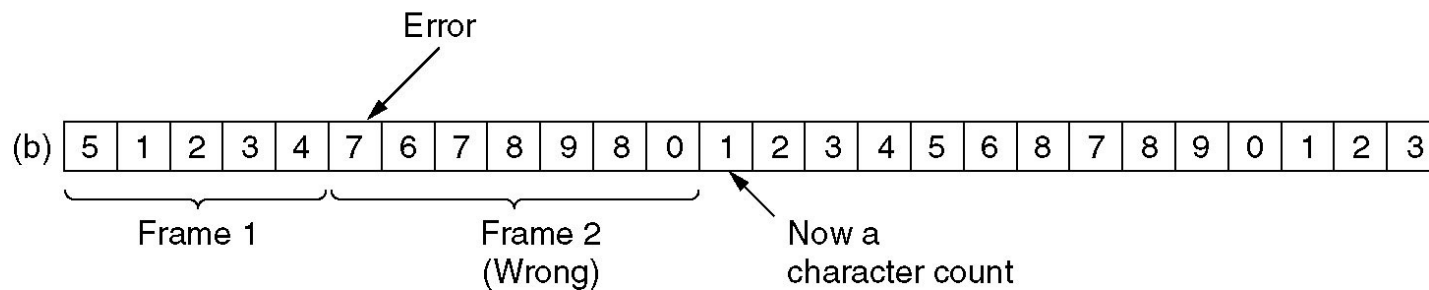
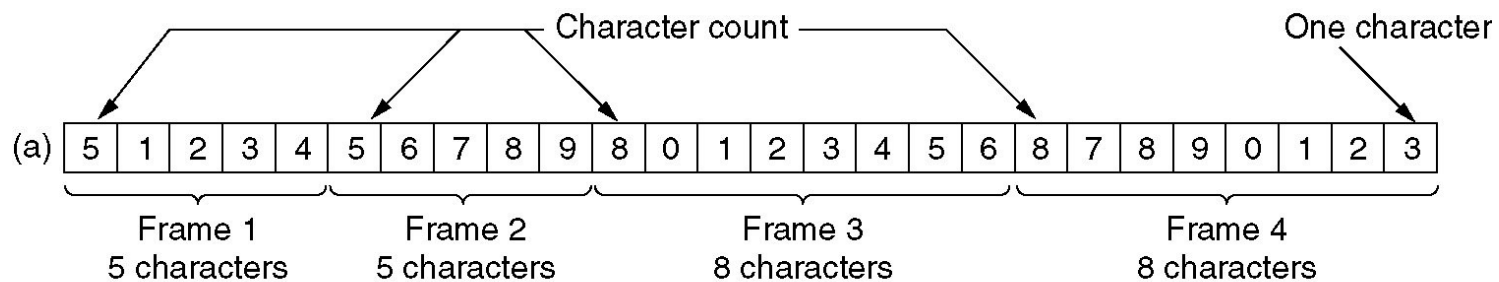
- Frame is the Basic unit of data streams
- Packets are broken into frames
- Frame size depends on the layer below DLL
- Checksum is computed and appended in the frame
- Generally used
  - Hybrid of Character count plus any one



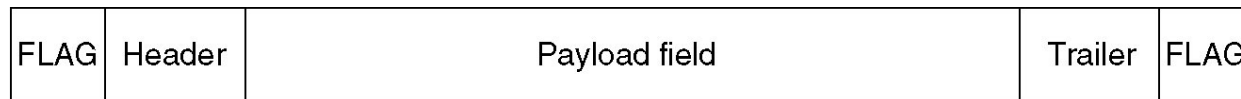
# End of Frame Detection

- Character count
  - Count can be garbled by a transmission error
- Starting and ending character with character stuffing
  - DLE STX for starting DLE ETX for ending
  - This sequences might be in floating point data
  - Character stuffing by two DLE for an actual DLE
- Starting and Ending flags, with bit stuffing
  - 01111110
  - When five 1s a 0 is stuffed
- Physical layer coding violations
  - Applicable when physical medium contains some redundancy

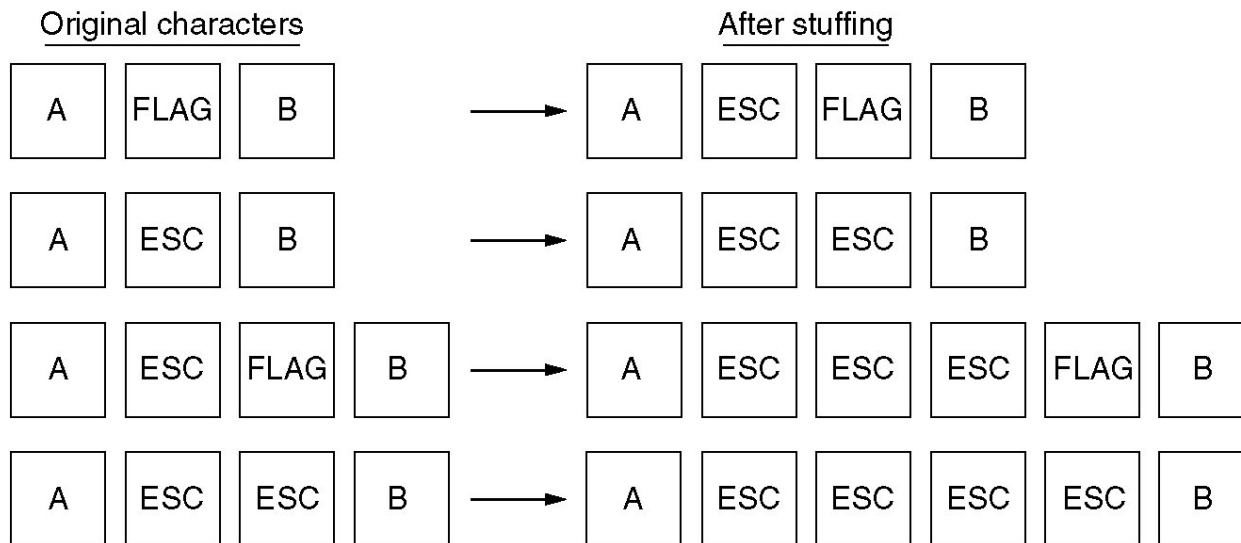
A character stream. (a) Without errors. (b)  
With one error.



- (a) A frame delimited by flag bytes.
- (b) Four examples of byte sequences before and after stuffing.



(a)



(b)

(a) 0 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 0 0 1 0

(b) 0 1 1 0 1 1 1 1 1 0 1 1 1 1 1 0 1 1 1 1 1 0 1 0 0 1 0

Stuffed bits

(c) 0 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 0 0 1 0

## Bit stuffing

- (a) The original data.
- (b) The data as they appear on the line.
- (c) The data as they are stored in receiver's memory after destuffing.

# Error Control

- Ack is the solution
- A frame is sent and sender waits for Ack
- If Ack is not back within the deadline
  - Retransmission occurs
  - Not an infinite number of transmission
  - Detection of Receiver problem
- Sequence numbers for preventing duplications

# Flow Control

- Sender is faster than Receiver
- Flow control to throttle the sender
- Sender must know whether the receiver is ready
- May be some receiver window of particular size

# Error Detection and Correction

Hopefully Already  
Studied in other  
Courses

# Elementary Data Link Protocols

- DLL accepts a packet from the network layer
- Frames are created
  - By adding header and trailer
  - Header : *frame\_kind, seq\_nr, ack*
- Frame is sent to the destination by calling *to\_physical\_layer( )*



# Elementary Data Link Protocols

- *wait\_for\_event( )* is waiting in the receiver of DLL for the frames
- *from\_physical\_layer( )* extracts the frame
- Checksum is computed
- The frame is sent to the Network layer of receiver
- An Ack is sent to the sender of DLL

# Elementary Data Link Protocols

- Sender waits for Ack from receiver
- Retransmits the frame
  - If time out occurs
- The Seq number
  - 0 to MAX\_SEQ (inclusive)
  - Incremented after each frame transmission

# Sliding Window Protocols

- Features
  - Piggybacking
  - Sequence Number for each frame
  - Windows for sender and receiver
- Examples
  - A One-Bit Sliding Window Protocol
  - A Protocol Using Go Back N
  - A Protocol Using Selective Repeat

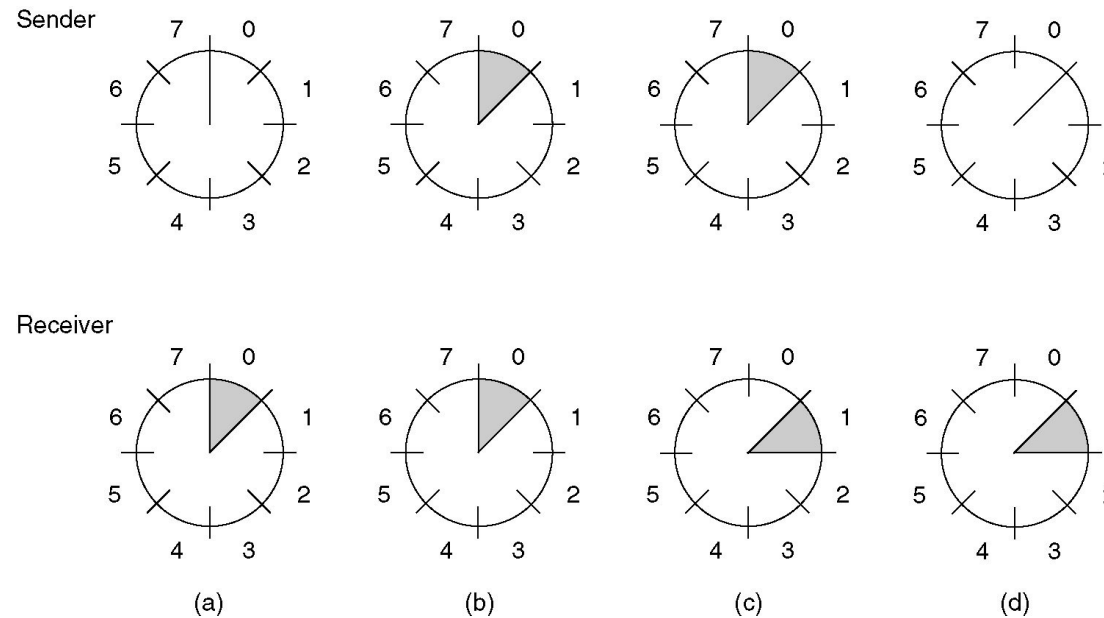
# Piggybacking

- ACK frames consume lots of BW
- ACK frame gets a free ride on the data frames
- Takes time if receiver has no data frames
  - Causes lots of retransmission
  - Should not wait for free ride

# Windows

- Seq number 0 to  $2^n-1$
- If  $n=1$  then two Seq Number 0 and 1
  - Called Stop and Wait Protocol
- Sending Window
  - The frames are being sent
  - Needs buffer to hold frames
- Receiving Window
  - permitted frames to accept

# Example of Sliding Window Protocols



A sliding window of size 1, with a 3-bit sequence number.

(a) Initially.

(b) After the first frame has been sent.

(c) After the first frame has been received.

(d) After the first acknowledgement has been received.

# Why Window?

- Prevents Duplication of packets in the network Layer
- ACK can be destroyed
- Frames will be retransmitted
- The Receiver must discard those duplication

# One Bit Sliding Window Protocol

- Stop and Wait Protocol
- Sender Window
  - *Next frame to send* (A variable)
  - Initially 0
- Receiver Window
  - *Frame expected* (A variable)
  - Initially 0



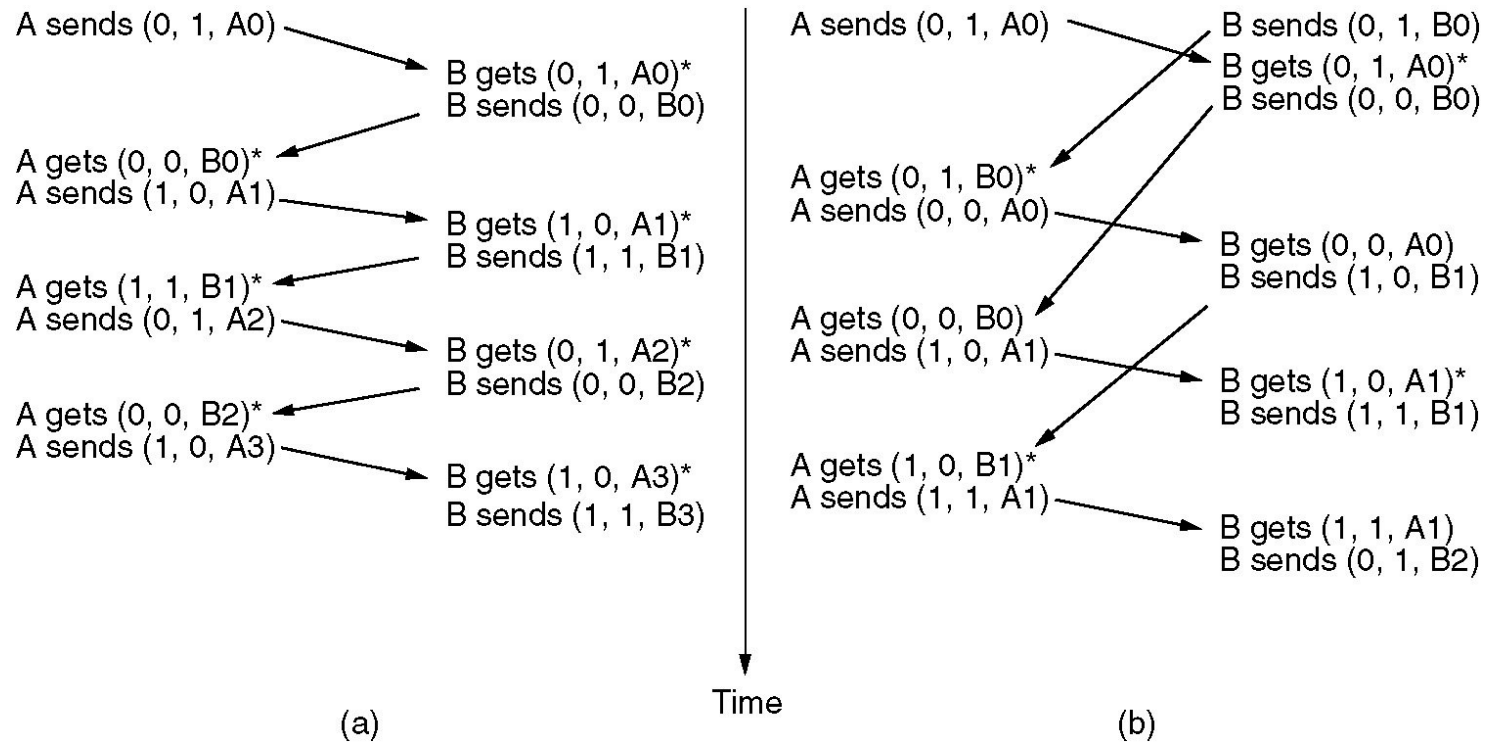
# One Bit Sliding Window Protocol

- $Seq = Next\_frame\_to\_send$
- $Ack = Inverse(frame\_expected)$ , I.e., previous received frame without error
- incoming frame is accepted if
  - $Seq == frame\_expected$
  - $frame\_expected$  is incremented
- New Frame is Sent if
  - $Ack = Next\_frame\_to\_send$
  - $Next\_frame\_to\_send$  is incremented
  - New frame is brought from Network Layer

# Scenario of One Bit Sliding Window Protocol

- When Sender and Receiver starts simultaneously
  - Initially  $seq=0$  and  $ack=1$
- Peculiar situation if both side sends initial packet
  - Duplication of each frame
  - Does not make any disaster

# Scenario One-Bit Sliding Window Protocol (2)



Two scenarios for protocol (a) Normal case. (b) Abnormal case. The notation is (seq, ack, packet number). An asterisk indicates where a network layer accepts a packet.

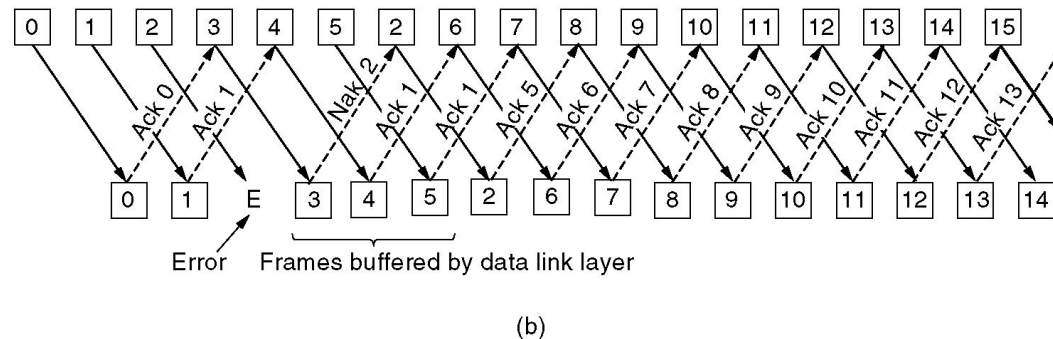
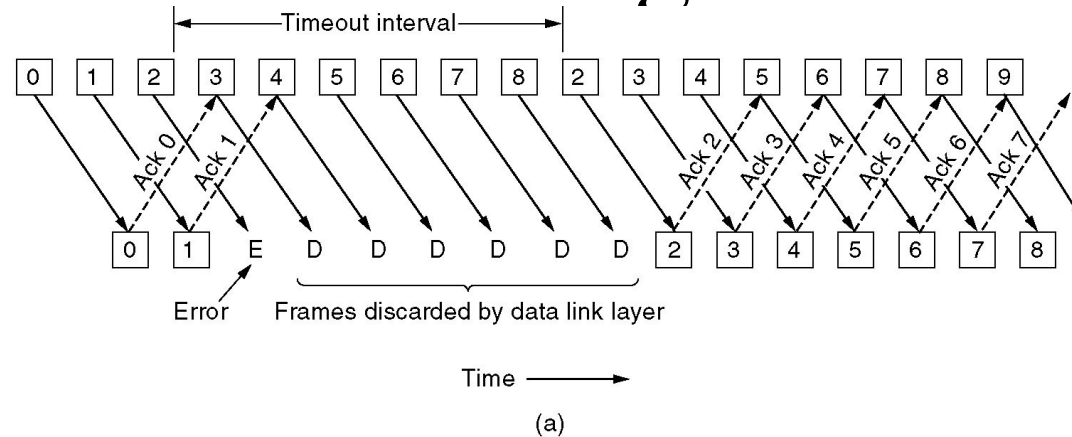
# Larger Window Size

- One bit Sliding Window is costly
  - When large propagation delay
  - Sender window should be bigger
- Pipelining technology should be used
  - 50% efficiency can be achieved
  - 50% loss for Acks
  - Two Approaches when error occurs

# Go Back n

- Consecutive frames are discarded
- Timeout occurs
- Retransmission of all the frames from the lost frame
- Channel bandwidth is not utilized

# A Protocol Using Go Back N



Pipelining and error recovery. Effect on an error when

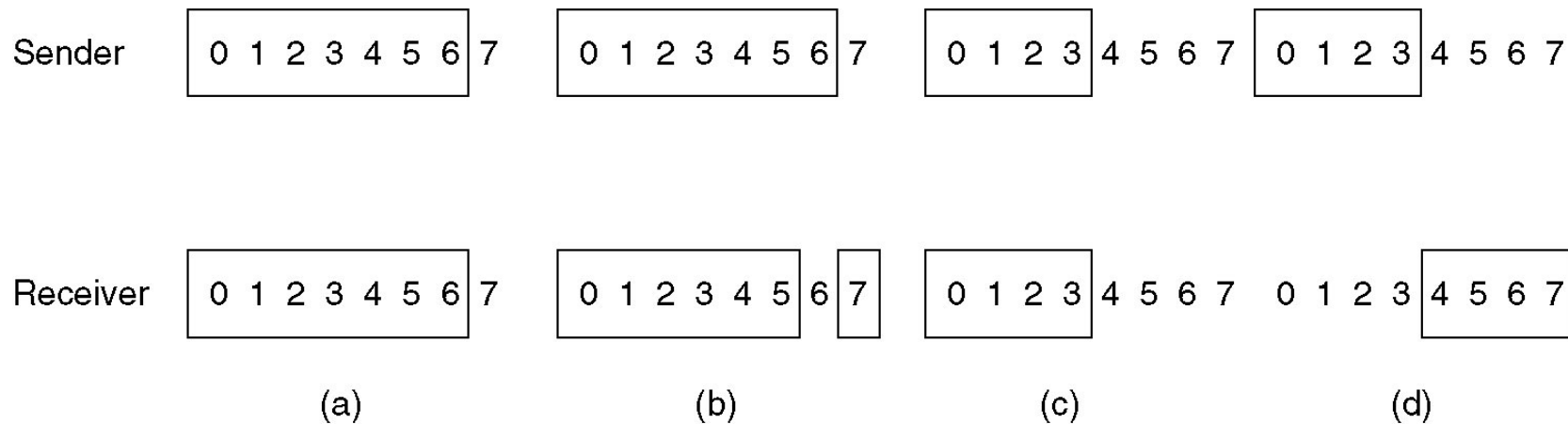
(a) Receiver's window size is 1.

(b) Receiver's window size is large.

# Selective Repeat

- Only the corrupted frame is discarded
- Good frames are accepted and buffered but not sent to network layer
- Only the last frame accepted by Net Layer is Acknowledged
- Timeout and retransmission of the faulty packet only

## Scenario of A Sliding Window Protocol Using Selective Repeat



- (a) Initial situation with a window size seven.
- (b) After seven frames sent and received, but not acknowledged.
- (c) Initial situation with a window size of four.
- (d) After four frames sent and received, but not acknowledged.



# Problem in Selective Repeat

- Maximum window size: MAX\_SEQ
- Seq Numbers : 0 to MAX\_SEQ
- A Faulty Scenario
  - Sender send 0 to 7
  - Gets all Acks, Again sends 0 7
  - Gets an Ack of 7
  - Which 7 is it? Might be for the first batch if all frames are lost
- Should be from 0 to 6

# More Problems

- Problems for Non Sequential Receipt of frames
  - Sender Sends 0 to 6
  - Receiver gets everything
  - Receiver is ready to get 7,0....5
  - All acks lost
  - Retransmission of 0 frame of first batch
  - Will be accepted as 0 frame of 2<sup>nd</sup> batch

# Solutions

- Should not be any overlapping
- Window Size  $(MAX\_SEQ+1)/2$
- Windows are
  - 0 to  $(MAX\_SEQ+1)/2-1$
  - $(MAX\_SEQ+1)/2$  to  $MAX\_SEQ$

# Negative Ack

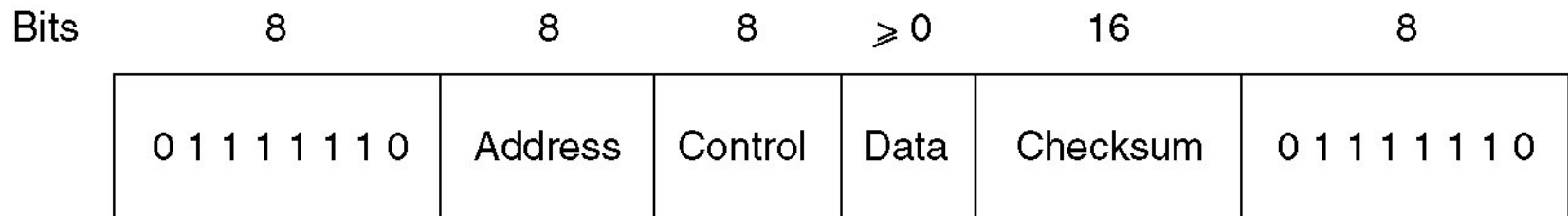
- Notifies the incorrect frames
- Retransmission if NO\_ACK is received
- Saves Channel Bandwidth
  - As the sender does not wait for retransmission

# Example Data Link Protocols

- HDLC – High-Level Data Link Control
- The Data Link Layer in the Internet  
–PPP

# High-Level Data Link Control

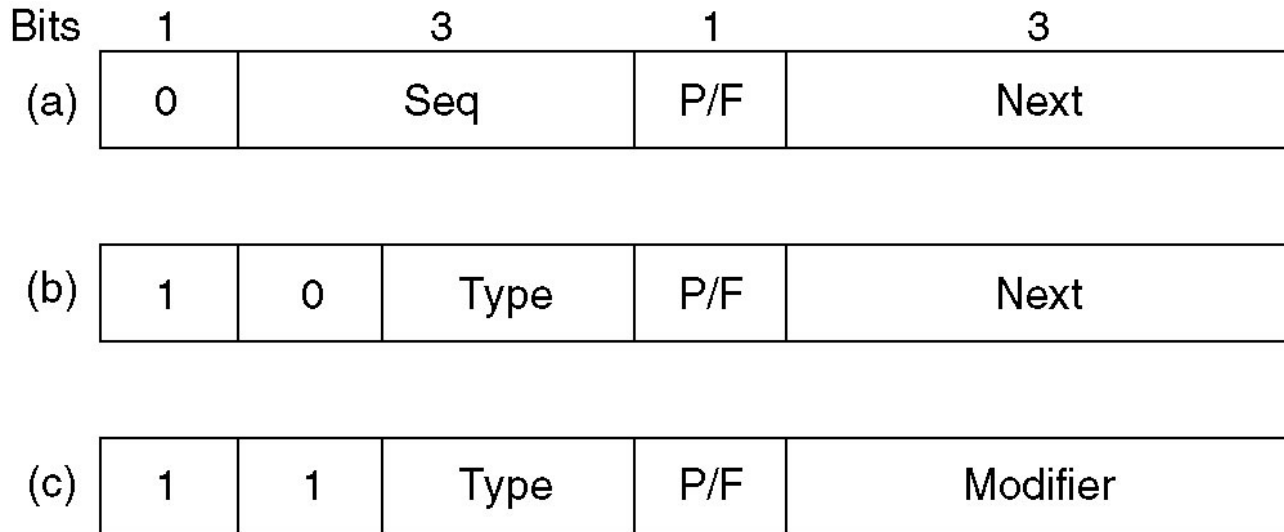
Frame format for bit-oriented protocols.



# Fields of HDLC protocol

- The frame is started and delimited by 01111110.
  - Bit stuffing is used if the sequence is present in the payload
- Address field is not desired as it is a protocol between two single connected entities
  - Required when several terminals are connected with the server
- Control field
  - Seq number, type and ack
- Data field may contain any number of bytes
  - Large data means degraded performance of checksum (CRC coding)

# Control Frame of HDLC



Control field of

(a) An information frame.

(b) A supervisory frame.

(c) An unnumbered frame.



# Details of Control Fields

- Sequence number is 3 bit
- Next field is the acknowledgement
  - May be piggybacked or not
  - May be the last frame number received
  - Or, the frame expected
- P/F bit is for polling the other party for sending information
  - The computer is polling the terminal by setting P/F as P
  - The terminal will send back the required data
    - The terminal data frame will be with F set.
    - Other frames will contain P/F set as P.

# Supervisory frame

- Use of P/F bit when the sender requires Ack
  - Receiver will not wait for reverse traffic
  - The receiver will send a supervisory frame
- 4 Types of Supervisory frame
  - Type 0, Ack frame without any data, Called Receive Ready
  - Type 1, negative Ack for transmission error
  - Type 2, Receive not ready, Acks all but not including Next.
    - Asks not to send any more frame
  - Type 3, Selective reject
    - Some frames have been discarded from buffer for sudden shortage of buffer

# Unnumbered Frames

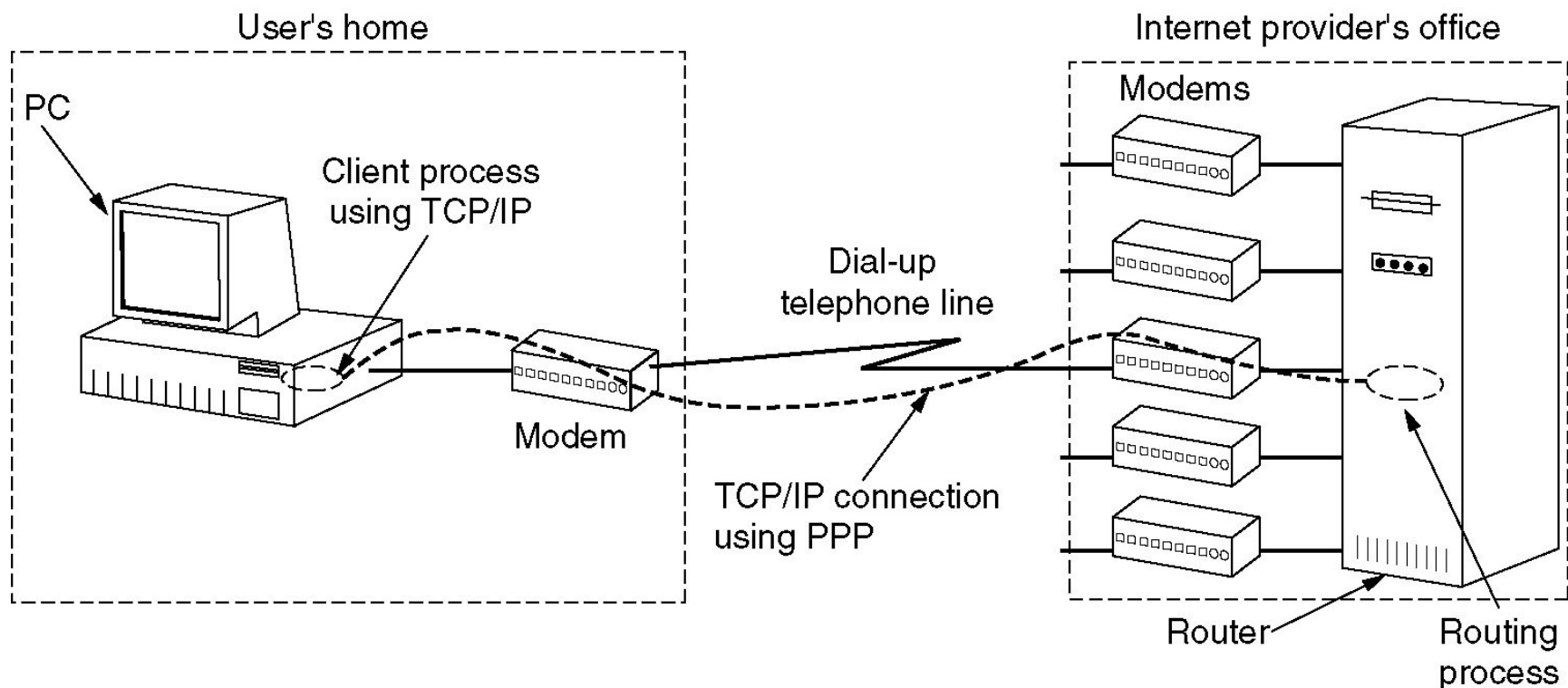
- Used for Control purpose
- Can carry data for unreliable connections
- 5 bits available, but not all are used
- Some of the examples are as follows
  - For allowing Disconnection
  - The notification for being live
  - Notification of frame rejection for unknown format
- Control frames can be lost
  - Ack is essential
  - No buffering

# Motivation to Point to Point lines in the Internet

- Sometime the routers are interconnected by LAN, not PP
- Connection to the outside world
  - Through point to point leased lines
  - Communication subnet is built over the PP leased lines
- Dialing up the ISP for the Internet
  - Through modem

# The Data Link Layer in the Internet

A home personal computer acting as an internet host.



# Point to Point Protocol (PPP)

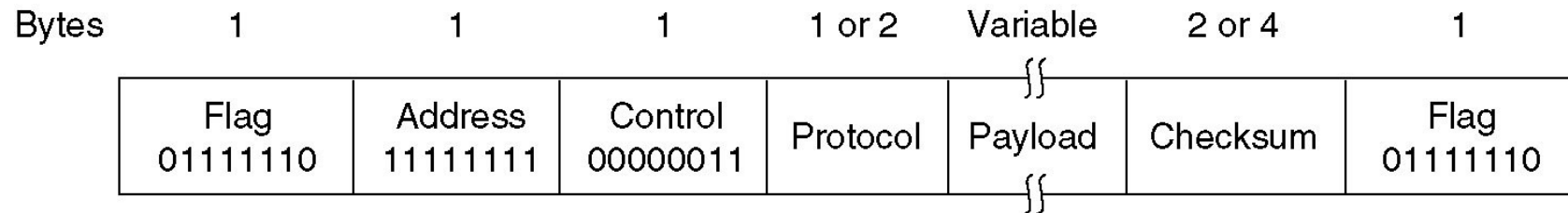
- Three features
  - Framing method
  - Link control protocol (LCP) for bringing lines
    - up
    - testing
    - Negotiating
    - Down
  - Way to negotiate network layer option
    - NCP (Network Control Protocol)
- The last two are the higher layer functionality

# PPP scenario

- PC calls the providers router via modem
- Provider's modem answered and physical connection established
  - Exchanges of LCP packets
  - Selecting PPP parameters
- Now NCP packets are exchanged for setting IP addresses
  - The provider might have pool of IP addresses
- NCP tears down the connection and frees up IP address
- LCP shuts down the DLL connection
- Computer tells the modem to hang up.

# PPP frame format

The PPP full frame format for unnumbered mode operation.





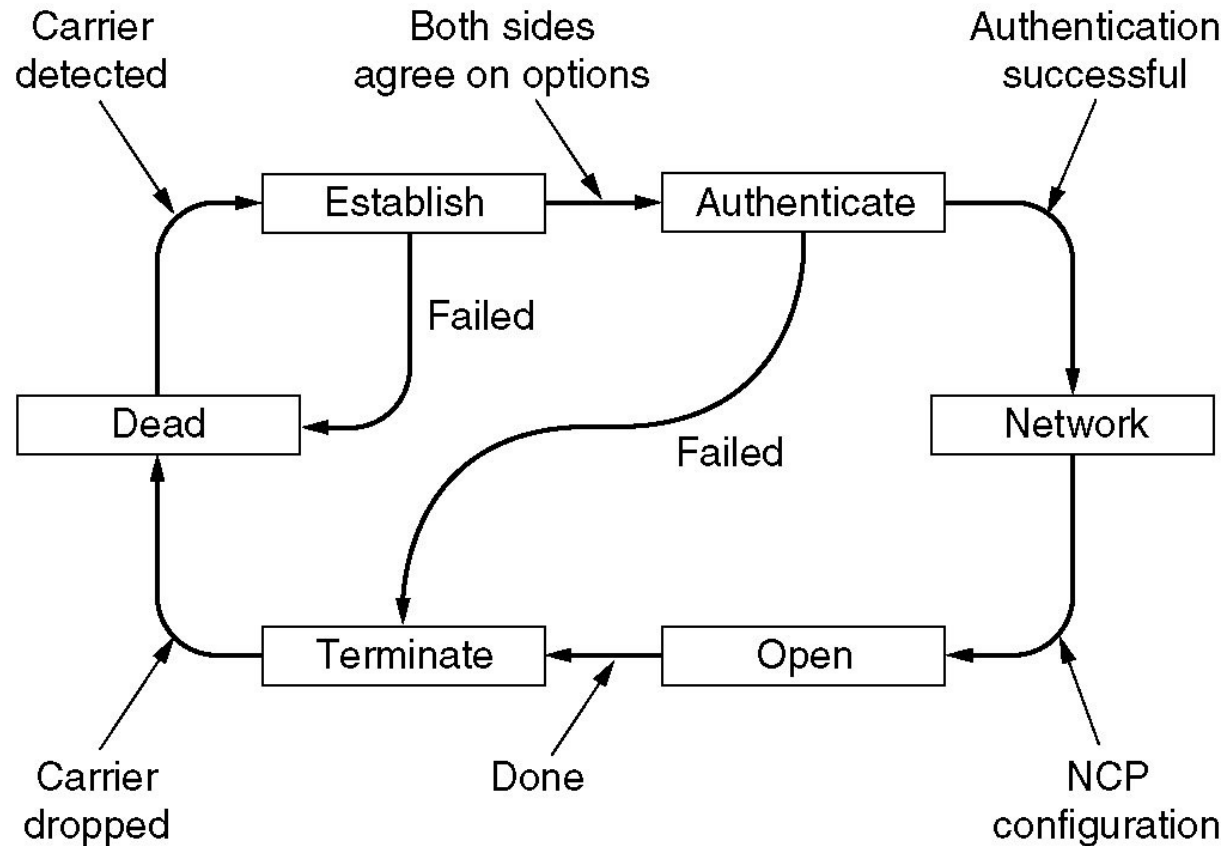
# PPP frame details

- Little different from HDLC
  - Not bit oriented, character oriented
  - Byte stuffing if the starting and ending is in the data
  - Modem can send integral number of bytes
    - Fractional numbers are not allowed
- Address is always 1111111, this is for compatibility with HDLC
  - Indicates all stations will accept the data

## PPP frame details (Cont.)

- Control field is the same as HDLC
  - But by default it is unnumbered, no reliable communication
  - For reliable transmission the same seq number and ack is used
- Protocol field specifies the type of payload
  - LCP, NCP, IP, IPX etc
- The length of Protocol and checksum can be negotiated down

# PPP – Point to Point Protocol (2)



A simplified phase diagram for bring a line up and down.