

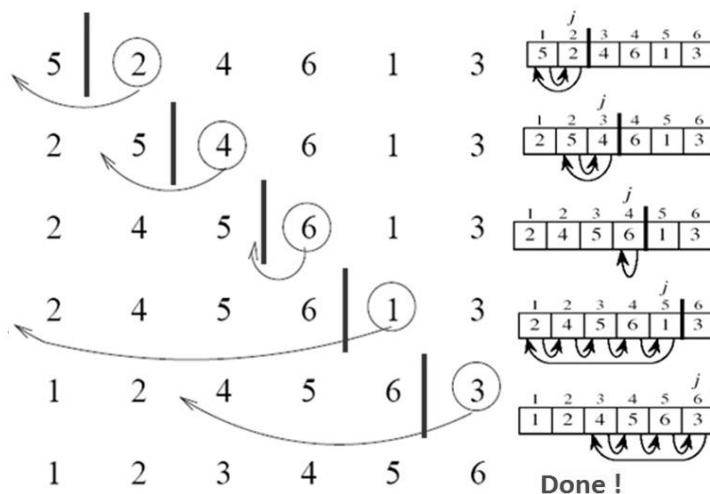
Heaven's Light is Our Guide
Computer Science & Engineering
Rajshahi University of Engineering & Technology

Lab Manual

Module- 09
Course Title: Sessional based on CSE 1201
Course No. : CSE 1202

Experiment No. 9**Name of the Experiment:** Sorting and Searching**Duration:** 1 cycle**Background Study:** Chapter 9 (Theory and Problems of Data Structures Written by Seymour Lipschutz)

[Contain:
 Insertion Sort & its Complexity
 Selection Sort & its Complexity
 Merge Sort & its Complexity
 Radix Sort & its Complexity
]

9.1 Insertion Sort

1. Select $TEMP \leftarrow A[j]$
2. Shift Data, and copy $TEMP$ to proper place

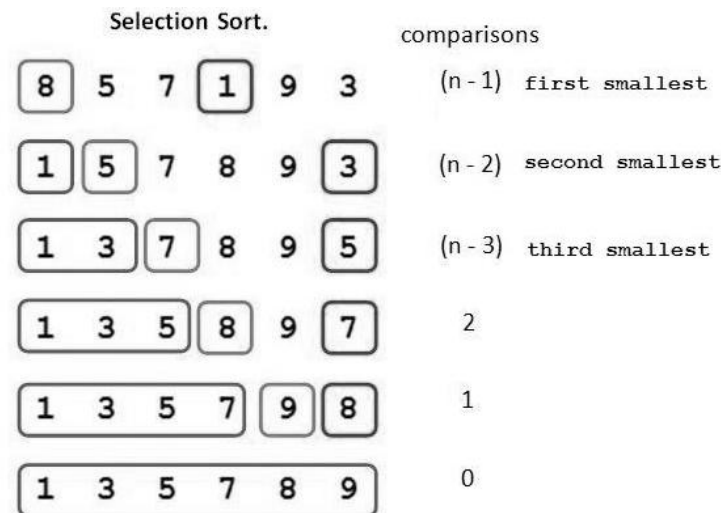
Problem I: Insertion Sort**Algorithm 9.1: INSERTION(A, N)**

This algorithm sorts the array A with N elements

1. Set $A[0] := -\infty$
2. Repeat Steps 3 to 5 for $K = 2, 3, \dots, N$:
3. Set $TEMP := A[K]$ and $PTR := K - 1$
4. Repeat while $TEMP < A[PTR]$:
5. (a) Set $A[PTR + 1] := A[PTR]$.
6. (b) Set $PTR := PTR - 1$.
7. [End of loop]
8. Set $A[PTR + 1] := TEMP$.
9. [End of Step 2]
10. Return.

Flow Chart: Draw a flow chart.**Complexity of Insertion Sort:**Worst Case: $O(n^2)$ Average Case: $O(n^2)$

9.2 Selection Sort



Total comparisons = $n(n-1)/2$

$\sim O(n^2)$

1. Find Min from A[2:N], and Interchange with First Element.
2. Find Min from A[3:N], and Interchange with second Element.
3. So on.

Problem II: Selection Sort

Algorithm 9.2(a): GETMIN(A,K,N,LOC)

An array A is in memory. The procedure finds the location LOC of the smallest element among A[K], A[K+1], ..., A[N].

1. Set MIN:=A[K] and LOC:=K.
2. Repeat for J=K+1 to N:
 - If MIN>A[J], then:
 - Set MIN:=A[J] and LOC:=J.
 - [End of If statement]
- [End of loop]
3. Return.

Algorithm 9.2(a): SELECTION(A,N)

This algorithm sorts the array A with N elements

1. Repeat Steps 2 and 3 for K=1,2, ..., N-1
2. **CALL GETMIN(A,K,N,LOC).**
3. Set TEMP:=A[K], A[K]:=A[LOC] and A[LOC]:=TEMP.
4. [End of Step 1 loop]
5. Exit

Flow Chart: Draw a flow chart.

Complexity of Selection Sort:

Worst Case: $O(n^2)$

Average Case: $O(n^2)$

9.3 Merge Sort

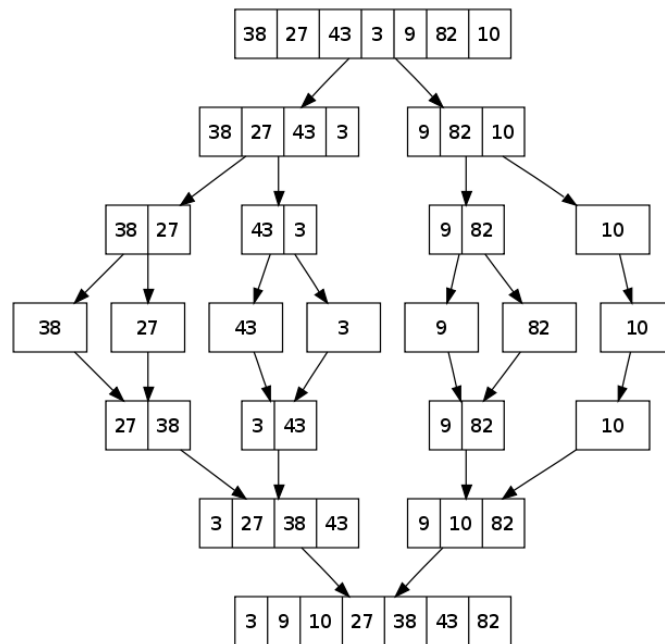


Fig. Merge Sort

Problem III: Merge Sort

Algorithm 9.3(a): MERGE(A,R,LBA,B,S,LBB,C,LBC)

Let A be a sorted array with R elements and lower bound LBA, B be a sorted array with S elements and lower bound LBB and C has lower bound LBC. The procedure merges A and B into an array C with R+S elements.

1. Set NA:=LBA, NB:=LBB, PTR:=LBC, UBA:=LBA+R-1, UBB:=LBB+S-1.
2. Repeat while NA≤UBA and NB≤UBB
 - If A[NA]<B[NB], then:
 - (a) Set C[PTR]:=A[NA].
 - (b) Set PTR:=PTR+1 and NA:=NA+1.
 - Else:
 - (a) Set C[PTR]:=B[NB].
 - (b) Set PTR:=PTR+1 and NB:=NB+1.
 - [End of If Statement]
- [End of loop]
3. If NA>UBA, then:
 - Repeat for K=0,1,2,..., UBB-NB
 - Set C[PTR+K]:=B[NB+K].
 - [End of loop]
- Else:
 - Repeat for K=0,1,2,..., UBA-NA
 - Set C[PTR+K]:=A[NA+K].
 - [End of loop]
- [End of If statement]
4. Exit.

Algorithm 9.3(b): MERGEPASS(A,N,L,B)

The N elements array A is composed of sorted subarrays where each subarray has L elements except possibly the last subarray, which may have fewer than L elements. The procedure merges the pairs of subarrays of A and assigns them to the array B.

1. Q:=INT(N/(2*L)), S:=2*L*Q and R:=N-S.
2. Repeat for J=1,2,..., Q:
 - (a) Set LB:=(1+2)*J-2)*L.
 - (b) Call MERGE(A,L,LB,A,L,LB+L,B,LB)

- [End of Loop]
3. If $R \leq L$, then:
 - Repeat for $J=1, 2, \dots, R$:
 - Set $B(S+J) := A(S+J)$.
 - [End of loop]
 - Else:
 - Call **MERGE**($A, L, S+1, A, R, L+S+1, B, S+1$).
 - [End of If statement]
 4. Return.

Algorithm 9.3(c): MERGESORT(A,N)

The algorithm sorts the array A with N elements using an auxiliary array B

1. Set $L := 1$
2. Repeat Steps 3 to 5 while $L < N$:
3. Call **MERGE**PASS(A, N, L, B).
4. Call **MERGE**PASS($B, N, 2*L, A$).
5. Set $L := 4*L$.
- [End of Step 2 loop]
6. Exit.

Flow Chart: Draw a flow chart.

Complexity:

Worst Case: $O(n \log n)$

Average Case: $O(n \log n)$

Extra Memory: $O(n)$.

9.4 Radix Sort

LSD Radix Sorting:
Sort by the last digit, then by the middle and the first one

362	291	207	207
436	362	436	253
291	253	253	291
487	436	362	362
207	487	487	397
253	207	291	436
397	397	397	487

MSD Radix Sorting:
Sort by the first digit, then sort each of the groups by the next digit

237	237	216	211
318	216	211	216
216	211	237	237
462	268	268	268
211	318	318	318
268	462	462	460
460	460	460	462

A:

62	36	91	87	53	97
0	1	2	3	4	5

d = 10,

PASS-I:

STEP-1: Initialize

COUNT:

0	0	0	0	0	0	0	0	0	0
0	1	2	3	4	5	6	7	8	9

STEP 2: Consider LSB of each Number. (2,6,1,7,3,7)

COUNT:

0	1	1	1	0	0	1	2	0	0
0	1	2	3	4	5	6	7	8	9

STEP 3: Now decide how to store the number**COUNT:**

0	1	2	3	3	3	4	6	6	6
0	1	2	3	4	5	6	7	8	9

STEP 4:**Now consider 97, LSB = 7 , Let Auxiliary array B****COUNT:**

0	1	2	3	3	3	4	5(6-1)	6	6
0	1	2	3	4	5	6	7	8	9

B:

									97
0	1	2	3	4	5				

Now consider 53, LSB = 3**COUNT:**

0	1	2	2(3-1)	3	3	4	5	6	6
0	1	2	3	4	5	6	7	8	9

B:

		53							97
0	1	2	3	4	5				

Now consider 87, LSB = 7**COUNT:**

0	1	2	2	3	3	4	4(5-1)	6	6
0	1	2	3	4	5	6	7	8	9

B:

		53		87					97
0	1	2	3	4	5				

Now consider 91, LSB = 1**COUNT:**

0	0	2	2	3	3	4	5	6	6
0	1	2	3	4	5	6	7	8	9

B:

91		53		87					97
0	1	2	3	4	5				

Now consider 36, LSB = 6**COUNT:**

0	0	2	2	3	3	3	5	6	6
0	1	2	3	4	5	6	7	8	9

B:

91		53	36	87					97
0	1	2	3	4	5				

Now consider 62, LSB = 2**COUNT:**

0	0	2	1	3	3	3	5	6	6
0	1	2	3	4	5	6	7	8	9

B:

91	62	53	36	87					97
0	1	2	3	4	5				

STEP 5: Copy from B to A:**A:**

91	62	53	36	87					97
0	1	2	3	4	5				

PASS-II:**STEP1: Again, Initialize****COUNT:**

0	0	0	0	0	0	0	0	0	0
0	1	2	3	4	5	6	7	8	9

STEP 2: Consider Next DIGIT of each Number. (6,3,9,8,5,9)**COUNT:**

0	0	0	1	0	1	1	0	1	2
0	1	2	3	4	5	6	7	8	9

STEP 3: Now decide how to store the number**COUNT:**

0	0	0	1	1	2	3	3	4	6
0	1	2	3	4	5	6	7	8	9

STEP 4:**Now consider 97, DIGIT = 9 , Let Auxiliary array B****COUNT:**

0	0	0	1	1	2	3	3	4	5
0	1	2	3	4	5	6	7	8	9

B:

									97
0	1	2	3	4	5	6	7	8	9

Now consider 87, LSB = 8**COUNT:**

0	0	0	1	1	2	3	3	3	5
0	1	2	3	4	5	6	7	8	9

B:

			87						97
0	1	2	3	4	5	6	7	8	9

Now consider 36, DIGIT = 3**COUNT:**

0	0	0	0	1	2	3	3	3	5
0	1	2	3	4	5	6	7	8	9

B:

36			87						97
0	1	2	3	4	5	6	7	8	9

Now consider 53, DIGIT = 5**COUNT:**

0	0	0	0	1	1	3	3	3	5
0	1	2	3	4	5	6	7	8	9

B:

36	53		87						97
0	1	2	3	4	5	6	7	8	9

Now consider 62, DIGIT = 6**COUNT:**

0	0	0	0	1	1	2	3	3	5
0	1	2	3	4	5	6	7	8	9

B:

36	53	62	87						97
0	1	2	3	4	5	6	7	8	9

Now consider 91, DIGIT = 9

COUNT:

0	0	0	0	1	1	2	3	3	4
0	1	2	3	4	5	6	7	8	9

B:

36	53	62	87	91	97
0	1	2	3	4	5

STEP 5: Copy from B to A:

A:

36	53	62	87	91	97
0	1	2	3	4	5

Problem IV: Radix Sort to sort a set of integer number

Algorithm 9.4(a): GETMAX(A,N,MAX)

An array A is in memory. The procedure returns largest element among A[1], a[2],..., A[N].

1. Set MAX:=A[1].
2. Repeat for J=2 to N:
 - If MAX<A[J], then:
 - Set MAX:=A[J].
 - [End of If statement]
- [End of loop]
3. Return.

Algorithm 9.4(b): COUNTSORT(A,N.EXP)

EXP is used to find the digit of each integer number. This procedure sorts the integer array A according to the digit represented by EXP using an auxiliary array B.

1. Repeat for J=0,1,, 9:
 - Set COUNT(J):=0.
- [End of loop]
2. Repeat for J=1,2,, N:
 - Set COUNT((A[j]/EXP)%10):= COUNT((A[j]/EXP)%10)+1.
- [End of loop]
3. Repeat for J=1,2,, 9:
 - Set COUNT(J):= COUNT(J) + COUNT(J-1).
- [End of loop]
4. Repeat for J=N,N-1,, 1:
 - Set COUNT((A[j]/EXP)%10):= COUNT((A[j]/EXP)%10)-1.
 - Set B(COUNT((A[j]/EXP)%10)):=A[J].
- [End of loop]
5. Repeat for J=1,2,, N:
 - Set A(J):=B(J).
- [End of loop]
6. Exit.

Algorithm 9.4(c): RADIXSORT(A,N)

The algorithm sorts the array A with N elements.

1. Call GETMAX(A,N,MAX).
2. Set EXP:=1.
3. Repeat Steps 4 and 5 while (MAX/EXP)>0
4. Call COUNTSORT(A,N.EXP).
5. EXP:=EXP*10.
- [End of Step 3 loop]
6. Exit.

Flow Chart: Draw a flow chart.

Complexity:

Let d be the radix, s be the total pass (i.e. max length of items) and n be the total number of items into a list.

Hence the number of comparisons for this algorithm is bounded as follows:

$$C(n) \leq d * s * n$$

d is constant for a specific case (e.g. $d = 10$ for decimal digits) so d is independent of n .

In worst case, $s = n$, then $O(n^2)$

In best case, $s = \log_d n$, then $O(n \log n)$.

Exercise:**[1] Supplementary Problem: Searching, Hashing****MORE PROBLEMS**

1. Programming (or supplementary) Problems of Chapter 9 of "Data Structures" by Seymour Lipschutz.

LAB REPORT: You have to submit all assigned problems in next lab.