MIST_Untitled
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Contents

1 C++

1.1 template

```
/*
c++:
ios_base::sync_with_stdio(false);
cin.tie(nullptr), cout.tie(nullptr);

python:
import sys
input = sys.stdin.readline
sys.stdout.write("------")
*/
```

1.2 random

```
#define accuracy chrono::steady_clock::
    now().time_since_epoch().count()
mt19937 rng(accuracy);

ll rand(ll l, ll r) {
```

```
  uniform_int_distribution<ll> ludo(l,
      r);
  return ludo(rng);
}
```

1.3 gp_hash

```
#include<ext/pb_ds/assoc_container.hpp>
#include<ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <typename p, typename q> using
    ht = gp_hash_table<p, q>;
```

1.4 pbds

```
#include<ext/pb_ds/assoc_container.hpp>
#include<ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <typename T>
using o_set = tree<T, null_type, less<T
    >, rb_tree_tag,
    tree_order_statistics_node_update>;
// find_by_order(k) - returns an
    iterator to the k-th largest element
     (0 indexed);
// order_of_key(k)- the number of
    elements in the set that are
    strictly smaller than k;
```

1.5 debug

```
string to_string(const string &s) {
    return '"' + s + '"'; }
string to_string(const char *s) {
    return to_string(string(s)); }
string to_string(const char c) { return
     "'" + string(1, c) + "'"; }
string to_string(bool b) { return b ? "
    true" : "false"; }
template <typename A, typename B>
    string to_string(pair<A, B> p) {
  return "(" + to_string(p.first) + ",
      " + to_string(p.second) + ")";
}
template <typename A> string to_string(
    A v) {
  string res = "{";
  for (const auto &x : v) {
    res += to_string(x) + ", ";
  }
  res += "}";
  return res;
}
void debug_out() { cerr << endl; }
template <typename Head, typename...
    Tail> void debug_out(Head H, Tail...
     T) {
  cerr << " " << to_string(H);
  debug_out(T...);
}
#define dbg(...)
    \
  cerr << __LINE__ << ": [" << #
      __VA_ARGS__ << "] = ", debug_out(
      __VA_ARGS__)
```

1.6 stress

```
#!/usr/bin/env bash
wrong="solution"
correct="brute"
gen="gen"
g++ -g solution.cpp -DONPC -o "$wrong"
g++ -g brute.cpp -DONPC -o "$correct"
g++ -g gen.cpp -DONPC -o "$gen"

for ((testNum=0;testNum<$1;testNum++))
do
        ./$gen 2>/dev/null > stdinput
        ./$correct < stdinput 2>/dev/
            null > outSlow
        ./$wrong < stdinput 2>/dev/null
            > outWrong
        H1=`md5sum outWrong`
        H2=`md5sum outSlow`
        if !(cmp -s "outWrong" "outSlow
            ")
        then
        echo "Error found!"
        echo "Input:"
        cat stdinput
        echo "Wrong Output:"
        cat outWrong
        echo "Slow Output:"
        cat outSlow
        exit
        fi
done
echo Passed $1 tests
# Usage: ./contest.sh times
```

1.7 vscode

```
{
  "key" : "f5",
  "command" : "workbench.action.
      terminal.sendSequence",
  "args" : {
    "text" : "g++ ${
        fileBasenameNoExtension}.cpp -o
         ${fileBasenameNoExtension} &&
        ./ ${fileBasenameNoExtension} <
        in.txt> out.txt\n "
  }
}
```

2 Dsa

2.1 KMP

```
vector<ll> createLPS(string pattern) {
  ll n = pattern.length(), idx = 0;
  vector<ll> lps(n);
  for (ll i = 1; i < n;) {
    if (pattern[idx] == pattern[i]) {
      lps[i] = idx + 1;
      idx++, i++;
    } else {
      if (idx != 0)
        idx = lps[idx - 1];
      else
        lps[i] = idx, i++;
    }
  }
  return lps;
}
ll kmp(string text, string pattern) {
  ll cnt_of_match = 0, i = 0, j = 0;
  vector<ll> lps = createLPS(pattern);
  while (i < text.length()) {
    if (text[i] == pattern[j])
      i++, j++; // i = text, j =
          pattern
    else {
      if (j != 0)
        j = lps[j - 1];
      else
        i++;
    }
    if (j == pattern.length()) {
      cnt_of_match++;
      // the index where match found ->
          (i - pattern.length());
      j = lps[j - 1];
    }
  }
  return cnt_of_match;
}
```

2.2 Hashing

```
const ll N = 2e5 + 5;
const ll MOD1 = 127657753, MOD2 =
    987654319;
const ll p1 = 137, p2 = 277;
ll ip1, ip2;
pair<ll, ll> pw[N], ipw[N];
void prec() {
  pw[0] = {1, 1};
  for (ll i = 1; i < N; i++) {
    pw[i].first = 1LL * pw[i - 1].first
        * p1 % MOD1;
    pw[i].second = 1LL * pw[i - 1].
        second * p2 % MOD2;
  }
  ip1 = binaryExp(p1, MOD1 - 2, MOD1);
  ip2 = binaryExp(p2, MOD2 - 2, MOD2);
  ipw[0] = {1, 1};
  for (ll i = 1; i < N; i++) {
    ipw[i].first = 1LL * ipw[i - 1].
        first * ip1 % MOD1;
    ipw[i].second = 1LL * ipw[i - 1].
        second * ip2 % MOD2;
  }
}
struct Hashing {
  ll n;
  string s;                    // 0 -
      indexed
  vector<pair<ll, ll>> hs; // 1 -
      indexed
  Hashing() {}
  Hashing(string _s) {
    n = _s.size();
    s = _s;
    hs.emplace_back(0, 0);
    for (ll i = 0; i < n; i++) {
      pair<ll, ll> p;
      p.first = (hs[i].first + 1LL * pw
          [i].first * s[i] % MOD1) %
          MOD1;
      p.second = (hs[i].second + 1LL *
          pw[i].second * s[i] % MOD2) %
          MOD2;
      hs.push_back(p);
    }
  }
  pair<ll, ll> get_hash(ll l, ll r) {
    // 1 - indexed
    assert(1 <= l && l <= r && r <= n);
    pair<ll, ll> ans;
    ans.first =
        (hs[r].first - hs[l - 1].first
            + MOD1) * 1LL * ipw[l - 1].
            first % MOD1;
    ans.second = (hs[r].second - hs[l -
        1].second + MOD2) * 1LL *
                 ipw[l - 1].second %
                     MOD2;
    return ans;
  }
  pair<ll, ll> get_hash() { return
      get_hash(1, n); }
};
```

2.3 BigInteger

```
struct BigInteger {
  string str;
  // Constructor to initialize
  // BigInteger with a string
  BigInteger(string s) { str = s; }
  // Overload + operator to add
  // two BigInteger objects
  BigInteger operator+(const BigInteger
      &b) {
    string a = str, c = b.str;
    ll alen = a.length(), clen = c.
        length();
    ll n = max(alen, clen);
    if (alen > clen)
      c.insert(0, alen - clen, '0');
    else if (alen < clen)
      a.insert(0, clen - alen, '0');
    string res(n + 1, '0');
    ll carry = 0;
    for (ll i = n - 1; i >= 0; i--) {
      ll digit=(a[i -'0')+(c[i]-'
          0')
      +carry;
      carry = digit / 10;
      res[i + 1] = digit % 10 + '
          0';
    }
    if (carry == 1) {
      res[0] = '1';
      return BigInteger(res);
    } else
      return BigInteger(res.
          substr(1));
  }
  // Overload - operator to subtract
  // first check which number is
     greater and then subtract
  BigInteger operator-(const BigInteger
      &b) {
    string a = str;
    string c = b.str;
    ll alen = a.length(), clen = c.
        length();
    ll n = max(alen, clen);
    if (alen > clen)
      c.insert(0, alen - clen, '0
          ');
    else if (alen < clen)
      a.insert(0, clen - alen, '0
          ');
    if (a < c) {
      swap(a, c);
      swap(alen, clen);
    }
    string res(n, '0');
    ll carry = 0;
    for (ll i = n - 1; i >= 0; i--) {
      ll digit = (a[i] - '0') - (
          c[i] - '0') - carry;
      if (digit < 0)
        digit += 10, carry = 1;
      else
        carry = 0;
      res[i] = digit + '0';
    }
    // remove leading zeros
    ll i = 0;
    while (i < n && res[i] == '0')
      i++;
    if (i == n)
      return BigInteger("0");
    return BigInteger(res.substr(i));
  }
  // Overload * operator to multiply
  // two BigInteger objects
  BigInteger operator*(const BigInteger
      &b) {
    string a = str, c = b.str;
    ll alen = a.length(), clen = c.
        length();
    ll n = alen + clen;
    string res(n, '0');
    for (ll i = alen - 1; i >= 0; i--)
        {
      ll carry = 0;
      for (ll j = clen - 1; j >=
          0; j--) {
        ll digit =
            (a[i] - '0') * (c[j -
                '0']) + (res[i +
                j + 1] - '0') +
                carry;
        carry = digit / 10;
        res[i + j + 1] = digit %
            10 + '0';
      }
      res[i] += carry;
    }
    ll i = 0;
    while (i < n && res[i] == '0')
```

```
            i++;
    if (i == n)
            return BigInteger("0");
    return BigInteger(res.substr(i));
  }
  // Overload << operator to output
  // BigInteger object
  friend ostream &operator<<(ostream &
      out, const BigInteger &b) {
    out << b.str;
    return out;
  }
};
```

2.4 Kadane

```
// return maximum subarray sum.
ll kadense(ll arr[], ll n) {
  ll mxsm = arr[0], curr_s = arr[0];
  for (ll i = 1; i < n; i++) {
    curr_s = max(arr[i], curr_s + arr[i
        ]);
    mxsm = max(mxsm, curr_s);
  }
  return mxsm;
}
```

2.5 Segement_tree

```
class SEGMENT_TREE {
public:
  vector<ll> v;
  vector<ll> seg;
  SEGMENT_TREE(ll n) {
    v.resize(n + 5);
    seg.resize(4 * n + 5);
  }
  //! initially: ti = 1, low = 1, high
      = n
  //(number of elements in the array);
  void build(ll ti, ll low, ll high) {
    if (low == high) {
      seg[ti] = v[low];
      return;
    }
    ll mid = (low + high) / 2;
    build(2 * ti, low, mid);
    build(2 * ti + 1, mid + 1, high);
    seg[ti] = (seg[2 * ti] + seg[2 * ti
        + 1]);
  }
  //! initially: ti = 1, low = 1, high
      = n
  //(number of elements in the array),
  //(ql & qr)=user input in 1 based
      index;
  ll find(ll ti, ll tl, ll tr, ll ql,
      ll qr) {
    if (tl > qr || tr < ql) {
      return 0;
    }
    if (tl >= ql and tr <= qr)
      return seg[ti];
    ll mid = (tl + tr) / 2;
    ll l = find(2 * ti, tl, mid, ql, qr
        );
    ll r = find(2 * ti + 1, mid + 1, tr
        , ql, qr);
    return (l + r);
  }
  //! initially: ti = 1, tl = 1, tr = n
  //(number of elements in the array),
  // id = user input in 1 based
      indexing,
  // val = updated value;
  void update(ll ti, ll tl, ll tr, ll
      id, ll val) {
    if (id > tr or id < tl)
      return;
    if (id == tr and id == tl) {
      seg[til] = val;
      return;
    }
    ll mid = (tl + tr) / 2;
    update(2 * ti, tl, mid, id, val);
    update(2 * ti + 1, mid + 1, tr, id,
        val);
    seg[ti] = (seg[2 * ti] + seg[2 * ti
        + 1]);
  }
};
// use 1 based indexing;
```

2.6 Fenwick_tree

```
struct FenwickTree {
  vector<ll> bit; // binary indexed
      tree
  ll n;
  FenwickTree(ll n) {
    this->n = n;
    bit.assign(n, 0);
  }
  FenwickTree(vector<ll> a) :
      FenwickTree(a.size()) {
    for (size_t i = 0; i < a.size(); i
        ++)
      add(i, a[i]);
  }
  ll sum(ll r) {
    ll ret = 0;
    for (; r >= 0; r = (r & (r + 1)) -
        1)
      ret += bit[r];
    return ret;
  }
  ll sum(ll l, ll r) { return sum(r) -
      sum(l - 1); }
  void add(ll idx, ll delta) {
    for (; idx < n; idx = idx | (idx +
        1))
      bit[idx] += delta;
  }
};

// minimum
struct FenwickTreeMin {
  vector<ll> bit;
  ll n;
  const ll INF = (ll)1e9;
  FenwickTreeMin(ll n) {
    this->n = n;
    bit.assign(n, INF);
  }
  FenwickTreeMin(vector<ll> a) :
      FenwickTreeMin(a.size()) {
    for (size_t i = 0; i < a.size(); i
        ++)
      update(i, a[i]);
  }
  ll getmin(ll r) {
    ll ret = INF;
    for (; r >= 0; r = (r & (r + 1)) -
        1)
      ret = min(ret, bit[r]);
    return ret;
  }
  void update(ll idx, ll val) {
    for (; idx < n; idx = idx | (idx +
        1))
      bit[idx] = min(bit[idx], val);
  }
};
```

2.7 Segment_tree_lazy

```
class SEGMENT_TREE {
public:
  vector<ll> v;
  vector<ll> seg;
  vector<ll> lazy;
  SEGMENT_TREE(ll n) {
    v.resize(n + 5, 0);
    seg.resize(4 * n + 5, 0);
    lazy.resize(4 * n + 5, 0);
  }
```

```
  void pull(ll ti) { seg[ti] = (seg[2 *
      ti] & seg[2 * ti + 1]); }
  void push(ll ti, ll tl, ll tr) {
    if (lazy[ti] == 0)
      return;
    seg[ti] |= lazy[ti];
    if (tl != tr) {
      lazy[2 * ti] |= lazy[ti];
      lazy[2 * ti + 1] |= lazy[ti];
    }
    lazy[ti] = 0;
  }
  //! llially: ti = 1, low = 1, high =
     n(number of elements in the array)
     ;
  void build(ll ti, ll low, ll high) {
    lazy[ti] = 0;
    if (low == high) {
      seg[ti] = v[low];
      return;
    }
    ll mid = (low + high) / 2;
    build(2 * ti, low, mid);
    build(2 * ti + 1, mid + 1, high);
    pull(ti);
  }
  //! llially: ti = 1, low = 1, high =
     n(number of elements in the array)
     , (ql
  //! & qr) = user input in 1 based
     indexing;
  ll query(ll ti, ll tl, ll tr, ll ql,
      ll qr) {
    push(ti, tl, tr);
    if (tl > qr || tr < ql) {
      return (1LL << 32) - 1;
    }
    if (tl >= ql and tr <= qr)
      return seg[ti];
    ll mid = (tl + tr) / 2;
    ll l = query(2 * ti, tl, mid, ql,
        qr);
    ll r = query(2 * ti + 1, mid + 1,
        tr, ql, qr);
    return (l & r);
  }
  //! llially: ti = 1, tl = 1, tr = n(
     number of elements in the array),
     id =
  //! user input in 1 based indexing,
     val = updated value;
  void update(ll ti, ll tl, ll tr, ll
      idL, ll idR, ll val) {
    push(ti, tl, tr);
    if (idR < tl or tr < idL)
      return;
    if (idL <= tl and tr <= idR) {
      lazy[ti] |= val;
      push(ti, tl, tr);
      return;
    }
    ll mid = (tl + tr) / 2;
    update(2 * ti, tl, mid, idL, idR,
        val);
    update(2 * ti + 1, mid + 1, tr, idL
        , idR, val);
    pull(ti);
  }
  // use 1 based indexing for input and
      queries and update;
};
```

2.8 Trie

```
const ll N = 26;
class Node {
public:
  ll EoW;
  Node *child[N];
  Node() {
    EoW = 0;
    for (ll i = 0; i < N; i++)
      child[i] = NULL;
  }
};
void insert(Node *node, string s) {
  for (size_t i = 0; i < s.size(); i++)
      {
    ll r = s[i] - 'A';
    if (node->child[r] == NULL)
      node->child[r] = new Node();
    node = node->child[r];
  }
  node->EoW += 1;
}
ll search(Node *node, string s) {
  for (size_t i = 0; i < s.size(); i++)
      {
    ll r = s[i] - 'A';
    if (node->child[r] == NULL)
      return 0;
  }
  return node->EoW;
}
void prll(Node *node, string s = "") {
  if (node->EoW)
    cout << s << "\n";
  for (ll i = 0; i < N; i++) {
    if (node->child[i] != NULL) {
      char c = i + 'A';
      prll(node->child[i], s + c);
    }
  }
}
bool isChild(Node *node) {
  for (ll i = 0; i < N; i++)
    if (node->child[i] != NULL)
      return true;
  return false;
}
bool isJunc(Node *node) {
  ll cnt = 0;
  for (ll i = 0; i < N; i++) {
    if (node->child[i] != NULL)
      cnt++;
  }
  if (cnt > 1)
    return true;
  return false;
}
ll trie_delete(Node *node, string s, ll
    k = 0) {
  if (node == NULL)
    return 0;
  if (k == (ll)s.size()) {
    if (node->EoW == 0)
      return 0;
    if (isChild(node)) {
      node->EoW = 0;
      return 0;
    }
    return 1;
  }
  ll r = s[k] - 'A';
  ll d = trie_delete(node->child[r], s,
      k + 1);
  ll j = isJunc(node);
  if (d)
    delete node->child[r];
  if (j)
    return 0;
  return d;
}
void delete_trie(Node *node) {
  for (ll i = 0; i < 15; i++) {
    if (node->child[i] != NULL)
      delete_trie(node->child[i]);
  }
  delete node;
}
```

2.9 DSU

```
class DisjollSet {
```

```
  vector<ll> par, sz, minElmt, maxElmt,
      cntElmt;
public:
  DisjollSet(ll n) {
    par.resize(n + 1);
    sz.resize(n + 1, 1);
    minElmt.resize(n + 1);
    maxElmt.resize(n + 1);
    cntElmt.resize(n + 1, 1);
    for (ll i = 1; i <= n; i++)
      par[i] = minElmt[i] = maxElmt[i]
          = i;
  }
  ll findUPar(ll u) {
    if (u == par[u])
      return u;
    return par[u] = findUPar(par[u]);
  }
  void unionBySize(ll u, ll v) {
    ll pU = findUPar(u);
    ll pV = findUPar(v);
    if (pU == pV)
      return;
    if (sz[pU] < sz[pV])
      swap(pU, pV);
    par[pV] = pU;
    sz[pU] += sz[pV];
    cntElmt[pU] += cntElmt[pV];
    minElmt[pU] = min(minElmt[pU],
        minElmt[pV]);
    maxElmt[pU] = max(maxElmt[pU],
        maxElmt[pV]);
  }
  ll getMinElementIntheSet(ll u) {
      return minElmt[findUPar(u)]; }
  ll getMaxElementIntheSet(ll u) {
      return maxElmt[findUPar(u)]; }
  ll getNumofElementIntheSet(ll u) {
      return cntElmt[findUPar(u)]; }
};
```

2.10 HLD

```
ll par[N], sub_tree_sz[N], heavy[N],
    wt_from_parent[N], depth[N], head[N
    ],
    position[N];
vector<pair<ll, ll>> gd[N];

// HLD part start
ll dfs(ll node, ll p) {
  par[node] = p;
  sub_tree_sz[node] = 1;
  heavy[node] = -1;

  for (auto [v, w] : gd[node]) {
    if (v == p)
      continue;
    depth[v] = depth[node] + 1;
    wt_from_parent[v] = w;
    sub_tree_sz[node] += dfs(v, node);
    if (heavy[node] == -1 ||
        sub_tree_sz[v] > sub_tree_sz[
        heavy[node]]) {
      heavy[node] = v;
    }
  }
  return sub_tree_sz[node];
}
ll pos;
void decompose(ll node, ll hd) {
  head[node] = hd;
  position[node] = ++pos;
  if (heavy[node] != -1) {
    decompose(heavy[node], hd);
  }
  for (auto [v, w] : gd[node]) {
    if (v != par[node] && v != heavy[
        node]) {
      decompose(v, v);
    }
  }
}
// HLD part end

// in main function
ll n, m;
cin >> n;
SEGMENT_TREE seg(n); // Lazy if needed
vector<ll> edge_u(n), edge_v(n),
    edge_node(n);

for (int i = 1; i < n; i++) {
  ll u, v, wt = 1;
  cin >> u >> v >> wt;
  gd[u].push_back({v, wt});
  gd[v].push_back({u, wt});
  edge_u[i] = u;
  edge_v[i] = v;
}

dfs(1, -1);
pos = 0;
decompose(1, 1);

for (int i = 1; i <= n; i++) {
  // seg.v[position[i]] = val[i];  //
      for node value
  seg.v[position[i]] = wt_from_parent[i
      ]; // for edge value
}

// work on a specific edge
for (int i = 1; i < n; i++) {
  ll u = edge_u[i], v = edge_v[i];
  edge_node[i] = (depth[u] > depth[v])
      ? u : v;
}

seg.build(1, 1, n);

auto updatePath = [&](ll u, ll v, ll x)
    {
  while (head[u] != head[v]) {
    if (depth[head[u]] < depth[head[v
        ]])
      swap(u, v);
    seg.update(1, 1, n, position[head[u
        ]], position[u], x);
    u = par[head[u]];
  }
  if (depth[u] > depth[v])
    swap(u, v);
  // edge value
  if (u != v) {
    seg.update(1, 1, n, position[u] +
        1, position[v], x);
  }
  // node value
  // seg.update(1, 1, n, position[u],
      position[v], x);
};

auto querypath = [&](ll u, ll v) {
  ll ans = -inf;
  while (head[u] != head[v]) {
    if (depth[head[u]] < depth[head[v
        ]])
      swap(u, v);
    ans = max(ans, seg.query(1, 1, n,
        position[head[u]], position[u]))
        ;
    u = par[head[u]];
  }
  if (depth[u] > depth[v])
    swap(u, v);
  // upward + downward
  if (u != v) {
    ans = max(ans, seg.query(1, 1, n,
        position[u] + 1, position[v]));
  }
  // only upward
//   ans = max(ans, seg.query(1, 1, n,
    position[u], position[v])); // for
    node value
  return ans;
};

seg.update(1, 1, n, position[edge_node[
    s]], position[edge_node[s]], x); //
    single point update. if path update
```

```
    need call update path
cout << querypath(x, s) << '\n';
```

2.11 Manacher

```
struct Manacher {
  vector<ll> p[2];
  string s;
  // p[1][i] = (max odd length
     palindrome centered at i) / 2 [
     floor division]
  // p[0][i] = same for even, it
     considers the right center
  // e.g. for s = "abbabba", p[1][3] =
     3, p[0][2] = 2
  Manacher(string s) {
    this->s = s;
    ll n = s.size();
    p[0].resize(n + 1);
    p[1].resize(n);
    for (ll z = 0; z < 2; z++) {
      for (ll i = 0, l = 0, r = 0; i <
         n; i++) {
        ll t = r - i + !z;
        if (i < r)
          p[z][i] = min(t, p[z][l + t])
             ;
        ll L = i - p[z][i], R = i + p[z
           ][i] - !z;
        while (L >= 1 && R + 1 < n && s
           [L - 1] == s[R + 1])
          p[z][i]++, L--, R++;
        if (R > r)
          l = L, r = R;
      }
    }
  }
  bool is_palindrome(ll l, ll r) {
    ll mid = (l + r + 1) / 2, len = r -
       l + 1;
    return 2 * p[len % 2][mid] + len %
       2 >= len;
  }
  string get_palin(ll i, bool odd =
     true) {
    ll len = p[odd][i];
    return s.substr(i - len, 2 * len +
       1 - !odd);
  }
};
```

2.12 2D prefix Sum

```
pref[i][j] = a[i][j] + pref[i - 1][j] +
    pref[i][j - 1] - pref[i - 1][j -
   1];
Sum of region = pref[row2 + 1][col2 +
   1] - pref[row2 + 1][col1] - pref[
   row1][col2 + 1] + pref[row1][col1];
```

2.13 CRT

```
class CRT {
  typedef long long vlong;
  typedef pair<vlong, vlong> pll;
  vector<pll> equations;

public:
  void clear() { equations.clear(); }
  vlong extended_euclid(vlong a, vlong
     b, vlong &x, vlong &y) {
    if (b == 0) {
      x = 1;
      y = 0;
      return a;
    }
    vlong x1, y1;
    vlong d = extended_euclid(b, a % b,
       x1, y1);
    x = y1;
    y = x1 - y1 * (a / b);
    return d;
  }
  vlong inverse(vlong a, vlong m) {
    vlong x, y;
    vlong g = extended_euclid(a, m, x,
       y);
    if (g != 1)
      return -1;
    return (x % m + m) % m;
  }
  /** Add equation of the form x = r (
     mod m) */
  void addEquation(vlong r, vlong m) {
     equations.push_back({r, m}); }
  pll solve() {
    if (equations.size() == 0)
      return {-1, -1};

    vlong a1 = equations[0].first;
    vlong m1 = equations[0].second;
    a1 %= m1;
    for (int i = 1; i < equations.size
       (); i++) {
      vlong a2 = equations[i].first;
      vlong m2 = equations[i].second;

      vlong g = __gcd(m1, m2);
      if (a1 % g != a2 % g)
        return {-1, -1};
      vlong p, q;
      extended_euclid(m1 / g, m2 / g, p
         , q);

      vlong mod = m1 / g * m2;
      vlong x = ((__int128)a1 * (m2 / g
         ) % mod * q % mod +
                 (__int128)a2 * (m1 / g
                    ) % mod * p % mod)
                    %
                mod;
      a1 = x;
      if (a1 < 0)
        a1 += mod;
      m1 = mod;
    }
    return {a1, m1};
  }
};
```

2.14 Intersect two arithmetic progression

```
using T = __int128;
// ax + by = __gcd(a, b)
// returns __gcd(a, b)
T extended_euclid(T a, T b, T &x, T &y)
    {
  T xx = y = 0;
  T yy = x = 1;
  while (b) {
    T q = a / b;
    T t = b;
    b = a % b;
    a = t;
    t = xx;
    xx = x - q * xx;
    x = t;
    t = yy;
    yy = y - q * yy;
    y = t;
  }
  return a;
}
pair<T, T> CRT(T a1, T m1, T a2, T m2)
   {
  T p, q;
  T g = extended_euclid(m1, m2, p, q);
  if (a1 % g != a2 % g)
    return make_pair(0, -1);
  T m = m1 / g * m2;
  p = (p % m + m) % m;
  q = (q % m + m) % m;
  return make_pair((p * a2 % m * (m1 /
     g) % m + q * a1 % m * (m2 / g) % m
     ) % m,
                   m);
}
```

```
// intersecting AP of two APs: (a1 +
    d1x) and (a2 + d2x)
pair<ll, ll> intersect(ll a1, ll d1, ll
    a2, ll d2) {
  auto x = CRT(a1 % d1, d1, a2 % d2, d2
      );
  ll a = x.first, d = x.second;
  if (d == -1)
    return {0, 0}; // empty
  ll st = max(a1, a2);
  a = a < st ? a + ((st - a + d - 1) /
      d) : a; // while (a < st) a += d;
  return {a, d};
}
```

2.15 Find nth value in a recurrence relation in O(logn

```
[ 1, 1; 1, 0 ] ^ (n - 1) =
    [F(n), F(n - 1); F(n - 1), F(n - 2)
        ]
    // Function to multiply two 2x2
        matrices
    void multiply(vector<vector<int>> &
        mat1, vector<vector<int>> &mat2
        ) {
      // Perform matrix multiplication
      int x = mat1[0][0] * mat2[0][0] +
          mat1[0][1] * mat2[1][0];
      int y = mat1[0][0] * mat2[0][1] +
          mat1[0][1] * mat2[1][1];
      int z = mat1[1][0] * mat2[0][0] +
          mat1[1][1] * mat2[1][0];
      int w = mat1[1][0] * mat2[0][1] +
          mat1[1][1] * mat2[1][1];

      // Update matrix mat1 with the
          result
      mat1[0][0] = x;
      mat1[0][1] = y;
      mat1[1][0] = z;
      mat1[1][1] = w;
    }

// Function to perform matrix
    exponentiation
void matrixPower(vector<vector<int>> &
    mat1, int n) {
  // Base case for recursion
  if (n == 0 || n == 1)
    return;

  // Initialize a helper matrix
  vector<vector<int>> mat2 = {{1, 1},
      {1, 0}};

  // Recursively calculate mat1^(n/2)
  matrixPower(mat1, n / 2);

  // Square the matrix mat1
  multiply(mat1, mat1);

  // If n is odd, multiply by the
      helper matrix mat2
  if (n % 2 != 0) {
    multiply(mat1, mat2);
  }
}

// Function to calculate the nth
    Fibonacci number
// using matrix exponentiation
int nthFibonacci(int n) {
  if (n <= 1)
    return n;

  // Initialize the transformation
      matrix
  vector<vector<int>> mat1 = {{1, 1},
      {1, 0}};

  // Raise the matrix mat1 to the power
       of (n - 1)
  matrixPower(mat1, n - 1);

  // The result is in the top-left cell
       of the matrix
  return mat1[0][0];
}
```

2.16 All_solution_of_ax+by_equal_c

```
// a*x+b*y=c. returns valid x and y if
    possible.
// all solutions are of the form (x0 +
    k * b / g, y0 - k * b / g)
bool find_any_solution(ll a, ll b, ll c
    , ll &x0, ll &y0, ll &g) {
  if (a == 0 and b == 0) {
    if (c)
      return false;
    x0 = y0 = g = 0;
    return true;
  }
  g = extended_euclid(abs(a), abs(b),
      x0, y0);
  if (c % g != 0)
    return false;
  x0 *= c / g;
  y0 *= c / g;
  if (a < 0)
    x0 *= -1;
  if (b < 0)
    y0 *= -1;
  return true;
}
void shift_solution(ll &x, ll &y, ll a,
    ll b, ll cnt) {
  x += cnt * b;
  y -= cnt * a;
}
// returns the number of solutions
    where x is in the range[minx, maxx]
    and y is
// in the range[miny, maxy]
ll find_all_solutions(ll a, ll b, ll c,
    ll minx, ll maxx, ll miny, ll maxy)
    {
  ll x, y, g;
  if (find_any_solution(a, b, c, x, y,
      g) == 0)
    return 0;
  if (a == 0 and b == 0) {
    assert(c == 0);
    return 1LL * (maxx - minx + 1) * (
        maxy - miny + 1);
  }
  if (a == 0) {
    return (maxx - minx + 1) * (miny <=
        c / b and c / b <= maxy);
  }
  if (b == 0) {
    return (maxy - miny + 1) * (minx <=
        c / a and c / a <= maxx);
  }
  a /= g, b /= g;
  ll sign_a = a > 0 ? +1 : -1;
  ll sign_b = b > 0 ? +1 : -1;
  shift_solution(x, y, a, b, (minx - x)
      / b);
  if (x < minx)
    shift_solution(x, y, a, b, sign_b);
  if (x > maxx)
    return 0;
  ll lx1 = x;
  shift_solution(x, y, a, b, (maxx - x)
      / b);
  if (x > maxx)
    shift_solution(x, y, a, b, -sign_b)
        ;
  ll rx1 = x;
  shift_solution(x, y, a, b, -(miny - y
      ) / a);
  if (y < miny)
    shift_solution(x, y, a, b, -sign_a)
        ;
  if (y > maxy)
    return 0;
  ll lx2 = x;
  shift_solution(x, y, a, b, -(maxy - y
      ) / a);
```

```
  if (y > maxy)
    shift_solution(x, y, a, b, sign_a);
  ll rx2 = x;
  if (lx2 > rx2)
    swap(lx2, rx2);
  ll lx = max(lx1, lx2);
  ll rx = min(rx1, rx2);
  if (lx > rx)
    return 0;
  return (rx - lx) / abs(b) + 1;
}
int32_t main() {
  ios_base::sync_with_stdio(0);
  cin.tie(0);
  int t, cs = 0;
  cin >> t;
  while (t--) {
    ll a, b, c, x1, x2, y1, y2;
    cin >> a >> b >> c >> x1 >> x2 >>
        y1 >> y2;
    cout << "Case " << ++cs << ": "
         << find_all_solutions(a, b, -c
             , x1, x2, y1, y2) << '\n';
  }
  return 0;
}
```

2.17 all soln of linear eq

```
struct Combi {
  int n;
  vector<ll> facts, finvs, invs;
  Combi(int _n) : n(_n), facts(_n),
      finvs(_n), invs(_n) {
    facts[0] = finvs[0] = 1;
    invs[1] = 1;
    for (int i = 2; i < n; i++)
      invs[i] = invs[mod % i] * (-mod /
          i);
    for (int i = 1; i < n; i++) {
      facts[i] = facts[i - 1] * i;
      finvs[i] = finvs[i - 1] * invs[i
          ];
    }
  }
  inline ll fact(int n) { return facts[
      n]; }
  inline ll finv(int n) { return finvs[
      n]; }
  inline ll inv(int n) { return invs[n
      ]; }
  inline ll ncr(int n, int k) {
    return n < k ? 0 : facts[n] * finvs
        [k] * finvs[n - k];
  }
};
Combi C(N);
// returns the number of solutions to
    the equation
// x_1 + x_2 + ... + x_n = s and 0 <= l
    <= x_i <= r
ll yo(int n, int s, int l, int r) {
  if (s < l * n)
    return 0;
  s -= l * n;
  r -= l;
  ll ans = 0;
  for (int k = 0; k <= n; k++) {
    ll cur = C.ncr(s - k - k * r + n -
        1 + 1, n - 1 + 1) * C.ncr(n, k);
    if (k & 1)
      ans -= cur;
    else
      ans += cur;
  }
  return ans;
}
int32_t main() {
  ios_base::sync_with_stdio(0);
  cin.tie(0);
  cout << yo(3, 3, 0, 1) << '\n';
  return 0;
}
```

2.18 Subset sum sqrt(n)

```
// Sum of elements <= N implies that
    every element is <= N
vector<int> freq(N + 1, 0);
for (int i = 0; i < N; i++) {
  int x;
  cin >> x;
  freq[x]++;
}
vector<pair<int, int>> compressed;
for (int i = 1; i <= N; i++) {
  if (freq[i] > 0)
    compressed.emplace_back(i, freq[i])
        ;
}
vector<int> dp(N + 1, 0);
dp[0] = 1;
for (const auto &[w, k] : compressed) {
  vector<int> ndp = dp;
  for (int p = 0; p < w; p++) {
    int sum = 0;
    for (int multiple = p, count = 0;
        multiple <= N; multiple += w,
        count++) {
      if (count > k) {
        sum -= dp[multiple - w * count
            ];
        count--;
      }
      if (sum > 0)
        ndp[multiple] = 1;
      sum += dp[multiple];
    }
  }
  swap(dp, ndp);
}
cout << "Possible subset sums are:\n";
for (int i = 0; i <= N; i++) {
  if (dp[i] > 0)
    cout << i << " ";
}
}
```

2.19 small giant (a^x = b mod m, find x, given other)

```
// Returns minimum x for which a ^ x %
    m = b % m, a and m are coprime.
int solve(int a, int b, int m) {
  a %= m, b %= m;
  int n = sqrt(m) + 1;
  int an = 1;
  for (int i = 0; i < n; ++i)
    an = (an * 1ll * a) % m;
  unordered_map<int, int> vals;
  for (int q = 0, cur = b; q <= n; ++q)
      {
    vals[cur] = q;
    cur = (cur * 1ll * a) % m;
  }
  for (int p = 1, cur = 1; p <= n; ++p)
      {
    cur = (cur * 1ll * an) % m;
    if (vals.count(cur)) {
      int ans = n * p - vals[cur];
      return ans;
    }
  }
  return -1;
}
// Returns minimum x for which a ^ x %
    m = b % m.
int solve(int a, int b, int m) {
  a %= m, b %= m;
```

```
  int k = 1, add = 0, g;
  while ((g = gcd(a, m)) > 1) {
    if (b == k)
      return add;
    if (b % g)
      return -1;
    b /= g, m /= g, ++add;
    k = (k * 1ll * a / g) % m;
  }
  int n = sqrt(m) + 1;
  int an = 1;
  for (int i = 0; i < n; ++i)
    an = (an * 1ll * a) % m;
  unordered_map<int, int> vals;
  for (int q = 0, cur = b; q <= n; ++q)
       {
    vals[cur] = q;
    cur = (cur * 1ll * a) % m;
  }
  for (int p = 1, cur = k; p <= n; ++p)
       {
    cur = (cur * 1ll * an) % m;
    if (vals.count(cur)) {
      int ans = n * p - vals[cur] + add
          ;
      return ans;
    }
  }
  return -1;
}
```

2.20 Gaussian Elimination

```
class GaussianElimination {
public:
  GaussianElimination(vector<vector<
      double>> matrix, vector<double>
      results)
      : matrix(matrix), results(results
          ), n(matrix.size()) {}
  void solve() {
    fElim();
    bSub();
  }
  vector<vector<double>> matrix;
  vector<double> results, solution;
  ll n;
  void fElim() {
    for (ll i = 0; i < n; ++i) {
      ll maxRow = i;
      for (ll k = i + 1; k < n; ++k)
        if (abs(matrix[k][i]) > abs(
            matrix[maxRow][i]))
          maxRow = k;
      swap(matrix[i], matrix[maxRow]);
      swap(results[i], results[maxRow])
          ;
      for (ll k = i + 1; k < n; ++k) {
        double factor = matrix[k][i] /
            matrix[i][i];
        for (ll j = i; j < n; ++j)
          matrix[k][j] -= factor *
              matrix[i][j];
        results[k] -= factor * results[
            i];
      }
    }
  }
  void bSub() {
    solution.resize(n);
    for (ll i = n - 1; i >= 0; --i) {
      solution[i] = results[i];
      for (ll j = i + 1; j < n; ++j)
        solution[i] -= matrix[i][j] *
            solution[j];
      solution[i] /= matrix[i][i];
    }
  }
};
```

2.21 Grundy

```
ll calculateGrundy(ll n, vector<ll> &
    grundy, const vector<ll> &moves) {
  if (grundy[n] != -1)
    return grundy[n];
  unordered_set<ll> s;
  for (ll move : moves) {
    if (n >= move) {
      s.insert(calculateGrundy(n - move
          , grundy, moves));
    }
  }
  ll g = 0;
  while (s.count(g))
    g++;
  return grundy[n] = g;
}
vector<ll> computeGrundy(ll maxN, const
    vector<ll> &moves) {
  vector<ll> grundy(maxN + 1, -1);
  grundy[0] = 0;
  for (ll i = 1; i <= maxN; ++i) {
    calculateGrundy(i, grundy, moves);
  }
  return grundy;
}
```

3 Dynamic Programming

3.1 LCS

```
/*
Fact about LCS:
1. Longest Increasing Substring
To solve this, we just care about when
    two char equals. Rest of the things
    should be neglected.
2. Longest Palindromic Subsequence(LPS)
To solve this, we just take a new
    string which is the reverse of the
    original string. Then just call the
    LCS function to find LPS.
3. Minimum insertions to make a string
    palindrome To solve this, we just
    basically do string length - LPS.
    Why this?
     Let's take an example: string s =
         aabca; Let's say aca is our LPS.
          Now we find how many char we
         need to insert to make the
         string palindrome while our LPS
         is fixed.
     a ab c a now to make the string
         palindrome we just need to
         insert the reverse of ab after c
         . So the new string looks like a
          ab c ba a
4. Minimum Number of Deletions and
    Insertions to make the string equals
     To solve this we just find the LCS
    of those string then just do: n + m
    - 2 * LCS.length() where n, m =
    strings length
*/
```

3.2 MCM

```
// TC: O(n ^ 3)
const ll N = 1005;
vector<ll> v;
ll dp[N][N], mark[N][N];
ll MCM(ll i, ll j) {
  if (i == j)
    return dp[i][j] = 0;
  if (dp[i][j] != -1)
    return dp[i][j];
  ll mn = INT_MAX;
  for (ll k = i; k < j; k++) {
    ll x = mn;
    mn = min(mn, MCM(i, k) + MCM(k + 1,
        j) + v[i - 1] * v[k] * v[j]);
    if (x != mn)
      mark[i][j] = k;
```

```
  }
  return dp[i][j] = mn;
}

void print_order(ll i, ll j) {
  if (i == j)
    cout << "X" << i;
  else {
    cout << "(";
    print_order(i, mark[i][j]);
    print_order(mark[i][j] + 1, j);
    cout << ")";
  }
}
// memset(dp, -1, sizeof dp);
// print_order(1, n);
```

3.3 LIS_length

```
vector<ll> v = {7, 3, 5, 3, 6, 2, 9,
    8};
vector<ll> seq;
/*
here we basically check is the current
    element from v is greater than the
    last element of the sequence. if it
    is then push it to the seq array and
     if not then replace that index
    value. let's take an example:
v = 7 3 5 3 6 2 9 8
1st iteration seq = 7;
2nd iteration seq = 3;
3rd iteration seq = 3 5;
4th iteration seq = 3 3;
5th iteration seq = 3 3 6;
6th iteration seq = 2 3 6;
7th iteration seq = 2 3 6 9;
8th iteration seq = 2 3 6 8;
*/
for (auto i : v) {
  auto id = lower_bound(seq.begin(),
      seq.end(), i);
  if (id == seq.end())
    seq.push_back(i);
  else
    seq[id - seq.begin()] = i;
}
cout << seq.size() << endl;
```

3.4 LCIS

```
ll a[100] = {0}, b[100] = {0}, f[100] =
    {0};
ll n = 0, m = 0;
ll main(void) {
  cin >> n;
  for (ll i = 1; i <= n; i++)
    cin >> a[i];
  cin >> m;
  for (ll i = 1; i <= m; i++)
    cin >> b[i];
  for (ll i = 1; i <= n; i++) {
    ll k = 0;
    for (ll j = 1; j <= m; j++) {
      if (a[i] > b[j] && f[j] > k)
        k = f[j];
      else if (a[i] == b[j] && k + 1 >
          f[j])
        f[j] = k + 1;
    }
  }
  ll and = 0;
  for (ll i = 1; i <= m; i++)
    if (f[i] > ans)
      ans = f[i];
  cout << and << endl;
  return 0;
}
```

3.5 SOS DP

```
// sum over subsets
for (int i = 0; i < B; i++) {
  for (int mask = 0; mask < (1 << B);
      mask++) {
    if ((mask & (1 << i)) != 0) {
      f[mask] += f[mask ^ (1 << i)];
    }
  }
}
// sum over supersets
for (int i = 0; i < B; i++) {
  for (int mask = (1 << B) - 1; mask >=
      0; mask--) {
    if ((mask & (1 << i)) == 0)
      g[mask] += g[mask ^ (1 << i)];
  }
}
// submask
for (int mask = 1; mask < (1 << 5);
    mask++) {
  for (int submask = mask; submask > 0;
      submask = ((submask - 1) & mask))
      {
    int subset = mask ^ submask;
  }
}
/**
    * we have to use SOS dp te idea is
      that we can iterate over all
      mask and using there subset get
      the actual value
    * the first is we have to find for
      each value in the array how many
       y we have so that x|y = x
    * if a binary is 1010 all possible
      y is 1000 0010 0000
    * so basically the subset of x
    * using sos dp we can find some
      properties from somethings
      subset
    * the idea is that we will consider
       first what if we can't change
      any bit to manipulate x what
      would be the ans
    * the 1 bit , 2 bit...log2(x)bit
      change from the right(LSB)
    * we will check what are the
      changes between binary
    * for example in the property of x|
      y = x
    * dp[i][x] = from right i bit can
      change and for x value
    * so initially this is the
      frequency of the x
    * then dp[1][x] will be dp[0][x] if
       1 bit is 0, because we
      obviously can't change it
      otherwise it will not be same
      after OR'ed
    * but if 1th bit is 1
    * dp[1][x] = dp[0][x] + dp[0][
      different with the mask or (x
      -(1<<i-1))] -> here is the right
       side part we basically needs to
       find in order to complete the
      code thats it
    * for example another case is x&y =
       x;
    * now lets see when is this
      property hold
    * if x is 1010
    * for 0 we can use other value
    * for 1 we can use only 1
    * now lets see 1010 -> for 1 bit
      change the subset(superset) is
      1011
    * for 2 bit change 1011 1010
    * for 3 bit change 1110 1011 1111
      1010
    * so if you see what are the diff
      berween bit 2 change and bit 3
      change it is just the 1110 and
      1111 now this is 1011+0100 ->
      1111 and 1010 + 0100 -> 1110 so
      each element in 2 bit change +
```

```
        (1<<(bit-1)); <--- this is the
        observation part
    * and now final one is x&y !=0
    *  if 1010 ->
    * for 1 bit change 1010 1011
    * for 2 bit change 1010 1000 1001
       1011
    * for 3 bit change 1010 1000 1001
       1011 1100 ..
    * for 4 bit change 1010 1000 0010
       ...
    * for 1010 just before mx bit all
       subset of 101 will be discared,
       if use any subset of 0101 like
       0100 it will not change anything
        cause 1010 & 0100 is 1000 the
       change is not on 0 bits
    * now how the code goes?
    * we iterate over all x
    * the we iterate from 0 bit change
       to max bit change  and store in
       dp
    * TC of sos dp is (b*2^b);
**/
```

3.6 BS optimization

```
bitset<100005> bs = 1;
for (auto i : a) {
  bs |= (bs << i);
  // if previous 1 value pos is
     possible now ith bit or ith sm is
     also possible
}
cout << bs.count() - 1 << endl;
for (ll i = 1; i <= 100003; i++)
  if (bs[i])
    cout << i << " ";
cout << endl;
```

4 Graph

4.1 Dijkstra

```
// TC: O(V + ElogV)
typedef pair<ll, ll> pairi;
ll N = 20000 + 5;
vector<vector<pairi>> adj(N);
vector<ll> dis(N, inf), parent(N);

void dijkstra(ll src) {
  priority_queue<pairi, vector<pairi>,
     greater<pairi>> pq;
  dis[src] = 0;
  pq.push({0, src});
  while (pq.size()) {
    auto top = pq.top();
    pq.pop();
    for (auto i : adj[top.second]) {
      ll v = i.first;
      ll wt = i.second;
      if (dis[v] > dis[top.second] + wt
          ) {
        dis[v] = dis[top.second] + wt;
        pq.push({dis[v], v});
        parent[v] = top.second
      }
    }
  }
}
ll node = n;
while (parent[node] != node) {
  path.push_back(node);
  node = parent[node];
}
path.push_back(1);
```

4.2 BellmanFord

```
// TC : O(V.E)
vector<ll> dist;
vector<ll> parent;
vector<vector<pair<ll, ll>>> adj;
// resize the vectors from main
   function

void bellmanFord(ll n, ll src) {
  dist[src] = 0;
  for (ll step = 0; step < n; step) {
    for (ll i = 1; i <= n; i++) {
      for (auto it : adj[i]) {
        ll u = i;
        ll v = it.first;
        ll wt = it.second;
        if (dist[u] != inf && ((dist[u]
            + wt) < dist[v])) {
          if (step == n - 1) {
            cout << "Negative cycle
                found\n ";
            return;
          }
          dist[v] = dist[u] + wt;
          parent[v] = u;
        }
      }
    }
  }
  for (ll i = 1; i <= n; i++)
    cout << dist[i] << " ";
  cout << endl;
}
```

4.3 FloydWarshall

```
// TC : O(n ^ 3)
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<ll> VI;
typedef vector<VI> VVI;

bool FloydWarshall(VVT &w, VVI &prev) {
  ll n = w.size();
  prev = VVI(n, VI(n, -1));

  for (ll k = 0; k < n; k++) {
    for (ll i = 0; i < n; i++) {
      for (ll j = 0; j < n; j++) {
        if (w[i][j] > w[i][k] + w[k][j
            ]) {
          w[i][j] = w[i][k] + w[k][j];
          prev[i][j] = k;
        }
      }
    }
  }
  // check for negative weight cycles
  for (ll i = 0; i < n; i++)
    if (w[i][i] < 0)
      return false;
  return true;
}
```

4.4 Toposort

```
// TC : O(V + E)
map<ll, vector<ll>> adj;
map<ll, ll> degree;
set<ll> nodes;
vector<ll> ans;
// adj: graph input, degree: cnt
   indegree,
// node: unique nodes, ans: path
ll c = 0;
void topo_sort() {
  queue<ll> qu;
  // traverse all the nodes and check
     if its degree is 0 or not..
  for (ll i : nodes) {
    if (degree[i] == 0)
      qu.push(i);
  }
  while (!qu.empty()) {
    ll top = qu.front();
    qu.pop();
```

```
    ans.push_back(top);
    for (ll i : adj[top]) {
      degree[i]--;
      if (degree[i] == 0) {
        qu.push(i);
      }
    }
  }
}
```

4.5 Kruskal

```
// TC : O(ElogE)
typedef pair<ll, ll> edge;
class Graph {
  vector<pair<ll, edge>> G, T;
  vector<ll> parent;
  ll cost = 0;

public:
  Graph(ll n) {
    for (ll i = 0; i < n; i++)
      parent.push_back(i);
  }
  void add_edges(ll u, ll v, ll wt) { G
      .push_back({wt, {u, v}}); }
  ll find_set(ll n) {
    if (n == parent[n])
      return n;
    else
      return find_set(parent[n]);
  }

  void union_set(ll u, ll v) { parent[u
      ] = parent[v]; }

  void kruskal() {
    sort(G.begin(), G.end());
    for (auto it : G) {
      ll uRep = find_set(it.second.
          first);
      ll vRep = find_set(it.second.
          second);
      if (uRep != vRep) {
        cost += it.first;
        T.push_back(it);
        union_set(uRep, vRep);
      }
    }
  }
  ll get_cost() { return cost; }
  void print() {
    for (auto it : T)
      cout << it.second.first << " " <<
          it.second.second << ": " <<
          it.first << endl;
  }
};

// g.add_edges(u, v, wt);
// g.kruskal();
```

4.6 Prims

```
// TC: O(ElogV)
typedef pair<ll, ll> pll;
class Prims {
  map<ll, vector<pll>> graph;
  map<ll, ll> visited;

public:
  void addEdge(ll u, ll v, ll w) {
    graph[u].push_back({v, w});
    graph[v].push_back({u, w});
  }

  vector<ll> path(pll start) {
    vector<ll> ans;
    priority_queue<pll, vector<pll>,
        greater<pll>> pq;
    // cost vs node
    pq.push({start.second, start.first
        });
    while (!pq.empty()) {
      pair<ll, ll> curr = pq.top();
      pq.pop();
      if (visited[curr.second])
        continue;
      visited[curr.second] = 1;
      ans.push_back(curr.second);
      for (auto i : graph[curr.second])
          {
        if (visited[i.first])
          continue;
        pq.push({i.second, i.first});
      }
    }
    return ans;
  }
};
```

4.7 LCA

```
// TC: preprocessing O(nlogn), each
    query O(logn)
ll n, l;
vector<vector<ll>> adj;
ll timer;
vector<ll> tin, tout;
vector<vector<ll>> up;

void dfs(ll v, ll p) {
  tin[v] = ++timer;
  up[v][0] = p;
  for (ll i = 1; i <= l; ++i)
    up[v][i] = up[up[v][i - 1]][i - 1];

  for (ll u : adj[v]) {
    if (u != p)
      dfs(u, v);
  }

  tout[v] = ++timer;
}

bool is_ancestor(ll u, ll v) { return
    tin[u] <= tin[v] && tout[u] >= tout[
    v]; }

ll lca(ll u, ll v) {
  if (is_ancestor(u, v))
    return u;
  if (is_ancestor(v, u))
    return v;
  for (ll i = l; i >= 0; --i) {
    if (!is_ancestor(up[u][i], v))
      u = up[u][i];
  }
  return up[u][0];
}

void preprocess(ll root) {
  tin.resize(n);
  tout.resize(n);
  timer = 0;
  l = ceil(log2(n));
  up.assign(n, vector<ll>(l + 1));
  dfs(root, root);
}
```

4.8 Rerooting

```
namespace reroot {
const auto exclusive = [](const auto &a
    , const auto &base,
                              const auto &
                                  merge_into
                                  , int
                                  vertex) {
  int n = (int)a.size();
  using Aggregate = decay_t<decltype(
      base)>;
  vector<Aggregate> b(n, base);
  for (int bit = (int)__lg(n); bit >=
      0; --bit) {
    for (int i = n - 1; i >= 0; --i)
      b[i] = b[i >> 1];
    int sz = n - (n & !bit);
    for (int i = 0; i < sz; ++i) {
      int index = (i >> bit) ^ 1;
      b[index] = merge_into(b[index], a
          [i], vertex, i);
```

```
    }
  }
  return b;
};
// MergeInto : Aggregate * Value *
    Vertex(int) * EdgeIndex(int) ->
    Aggregate
// Base : Vertex(int) -> Aggregate
// FinalizeMerge : Aggregate * Vertex(
    int) * EdgeIndex(int) -> Value
const auto rerooter = [](const auto &g,
     const auto &base,
                         const auto &
                             merge_into,
                              const auto
                              &
                             finalize_merge
                             ) {
  int n = (int)g.size();
  using Aggregate = decay_t<decltype(
      base(0))>;
  using Value = decay_t<decltype(
      finalize_merge(base(0), 0, 0))>;
  vector<Value> root_dp(n), dp(n);
  vector<vector<Value>> edge_dp(n),
      redge_dp(n);

  vector<int> bfs, parent(n);
  bfs.reserve(n);
  bfs.push_back(0);
  for (int i = 0; i < n; ++i) {
    int u = bfs[i];
    for (auto v : g[u]) {
      if (parent[u] == v)
        continue;
      parent[v] = u;
      bfs.push_back(v);
    }
  }

  for (int i = n - 1; i >= 0; --i) {
    int u = bfs[i];
    int p_edge_index = -1;
    Aggregate aggregate = base(u);
    for (int edge_index = 0; edge_index
        < (int)g[u].size(); ++
       edge_index) {
      int v = g[u][edge_index];
      if (parent[u] == v) {
        p_edge_index = edge_index;
        continue;
      }
      aggregate = merge_into(aggregate,
          dp[v], u, edge_index);
    }
    dp[u] = finalize_merge(aggregate, u
        , p_edge_index);
  }

  for (auto u : bfs) {
    dp[parent[u]] = dp[u];
    edge_dp[u].reserve(g[u].size());
    for (auto v : g[u])
      edge_dp[u].push_back(dp[v]);
    auto dp_exclusive = exclusive(
        edge_dp[u], base(u), merge_into,
         u);
    redge_dp[u].reserve(g[u].size());
    for (int i = 0; i < (int)
        dp_exclusive.size(); ++i)
      redge_dp[u].push_back(
          finalize_merge(dp_exclusive[i
          ], u, i));
    root_dp[u] = finalize_merge(
        n > 1 ? merge_into(dp_exclusive
            [0], edge_dp[u][0], u, 0) :
           base(u), u,
        -1);
    for (int i = 0; i < (int)g[u].size
        (); ++i) {
      dp[g[u][i]] = redge_dp[u][i];
    }
  }

  return make_tuple(move(root_dp), move
      (edge_dp), move(redge_dp));
};
} // namespace reroot

int main() {
  ll n;
  cin >> n;
  vector<vector<ll>> g(n);
  // everything should be 0 based.

  using Aggregate = int;
  using Value = int;

  auto base = [](int vertex) ->
      Aggregate {
    // task here
  };
  auto merge_into = [](Aggregate
      vertex_dp, Value neighbor_dp, int
      vertex, int edge_index) ->
      Aggregate {
    // task here
  };
  auto finalize_merge = [](Aggregate
      vertex_dp, int vertex, int
      edge_index) -> Value {
    // task here
  };

  auto [reroot_result, edge_dp,
      redge_dp] = reroot::rerooter(g,
      base, merge_into, finalize_merge);
}
```

4.9 Centroid_Tree

```
  const ll n = 1e5;
  vector<ll> sz(n + 5), dead(n + 5);
  function<void(ll, ll)> calculate_sz =
       [&](ll u, ll p) {
    sz[u] = 1;
    for (auto v : adj[u]) {
      if (v != p and !dead[v]) {
        calculate_sz(v, u);
        sz[u] += sz[v];
      }
    }
    return;
  };

  function<ll(ll, ll, ll)>
      find_centroid = [&](ll u, ll p, ll
       total) -> ll {
    for (auto v : adj[u]) {
      if (v != p and !dead[v] and 2 *
          sz[v] > total)
        return find_centroid(v, u,
            total);
    }
    return u;
  };

  function<void(ll)> decompose = [&] (
      ll u) -> void {
    // if needed change the parameter
    calculate_sz(u, -1);
    ll center = find_centroid(u, -1, sz
        [u]);
    // calculate the ans here
    dead[center] = 1;
    for(auto v : adj[center]) {
        if(!dead[v])
            decompose(v);
    }
  };
  // call decompose only
  decompose(1);
```

4.10 Euler_ckt

```
unordered_map<ll, ll> Start, End, Val;
unordered_map<ll, pair<ll, ll>> Range;
ll start = 0;
void dfs(ll node) {
  visited[node] = true;
  Start[node] = start++;
```

```
  for (auto child : adj[node]) {
    if (!visited[child])
      dfs(child);
  }
  End[node] = start - 1;
}
dfs(1);
vector<ll> FlatArray(start + 5);
for (auto i : Start) {
  FlatArray[i.second] = Val[i.first];
  Range[i.first] = {i.second, End[i.
      first]};
}
```

4.11 Min Cost Max Flow

```
#include <bits/stdc++.h>
using namespace std;
const int N = 3e5 + 9;
// Works for both directed, undirected
    and with negative cost too
// doesn't work for negative cycles
// for undirected edges just make the
    directed flag false
// Complexity: O(min(E^2 *V log V, E
    logV * flow))
using T = long long;
const T inf = 1LL << 61;
struct MCMF {
  struct edge {
    int u, v;
    T cap, cost;
    int id;
    edge(int _u, int _v, T _cap, T
        _cost, int _id) {
      u = _u;
      v = _v;
      cap = _cap;
      cost = _cost;
      id = _id;
    }
  };
  int n, s, t, mxid;
  T flow, cost;
  vector<vector<int>> g;
  vector<edge> e;
  vector<T> d, potential, flow_through;
  vector<int> par;
  bool neg;
  MCMF() {}
  MCMF(int _n) { // 0-based indexing
    n = _n + 10;
    g.assign(n, vector<int>());
    neg = false;
    mxid = 0;
  }
  void add_edge(int u, int v, T cap, T
      cost, int id = -1,
                bool directed = true) {
    if (cost < 0)
      neg = true;
    g[u].push_back(e.size());
    e.push_back(edge(u, v, cap, cost,
        id));
    g[v].push_back(e.size());
    e.push_back(edge(v, u, 0, -cost,
        -1));
    mxid = max(mxid, id);
    if (!directed)
      add_edge(v, u, cap, cost, -1,
          true);
  }
  bool dijkstra() {
    par.assign(n, -1);
    d.assign(n, inf);
    priority_queue<pair<T, T>, vector<
        pair<T, T>>, greater<pair<T, T
        >>> q;
    d[s] = 0;
    q.push(pair<T, T>(0, s));
    while (!q.empty()) {
      int u = q.top().second;
      T nw = q.top().first;
      q.pop();
      if (nw != d[u])
        continue;
      for (int i = 0; i < (int)g[u].
          size(); i++) {
        int id = g[u][i];
        int v = e[id].v;
        T cap = e[id].cap;
        T w = e[id].cost + potential[u]
            - potential[v];
        if (d[u] + w < d[v] && cap > 0)
            {
          d[v] = d[u] + w;
          par[v] = id;
          q.push(pair<T, T>(d[v], v));
        }
      }
    }
    for (int i = 0; i < n; i++) {
      if (d[i] < inf)
        d[i] += (potential[i] -
            potential[s]);
    }
    for (int i = 0; i < n; i++) {
      if (d[i] < inf)
        potential[i] = d[i];
    }
    return d[t] != inf; // for max flow
        min cost
    // return d[t] <= 0; // for min
        cost flow
  }
  T send_flow(int v, T cur) {
    if (par[v] == -1)
      return cur;
    int id = par[v];
    int u = e[id].u;
    T w = e[id].cost;
    T f = send_flow(u, min(cur, e[id].
        cap));
    cost += f * w;
    e[id].cap -= f;
    e[id ^ 1].cap += f;
    return f;
  }
  // returns {maxflow, mincost}
  pair<T, T> solve(int _s, int _t, T
      goal = inf) {
    s = _s;
    t = _t;
    flow = 0, cost = 0;
    potential.assign(n, 0);
    if (neg) {
      // Run Bellman-Ford to find
          starting potential on the
          starting graph
      // If the starting graph (before
          pushing flow in the residual
          graph) is a
      // DAG, then this can be
          calculated in O(V + E) using
          DP: potential(v) =
      // min({potential[u] + cost[u][v
          ]}) for each u -> v and
          potential[s] = 0
      d.assign(n, inf);
      d[s] = 0;
      bool relax = true;
      for (int i = 0; i < n && relax; i
          ++) {
        relax = false;
        for (int u = 0; u < n; u++) {
          for (int k = 0; k < (int)g[u
              ].size(); k++) {
            int id = g[u][k];
            int v = e[id].v;
            T cap = e[id].cap, w = e[id
                ].cost;
            if (d[v] > d[u] + w && cap
                > 0) {
              d[v] = d[u] + w;
              relax = true;
            }
          }
        }
```

```
      }
      for (int i = 0; i < n; i++)
        if (d[i] < inf)
          potential[i] = d[i];
    }
    while (flow < goal && dijkstra())
      flow += send_flow(t, goal - flow)
          ;
    flow_through.assign(mxid + 10, 0);
    for (int u = 0; u < n; u++) {
      for (auto v : g[u]) {
        if (e[v].id >= 0)
          flow_through[e[v].id] = e[v ^
              1].cap;
      }
    }
    return make_pair(flow, cost);
  }
};
int main() {
  ios_base::sync_with_stdio(0);
  cin.tie(0);
  int n;
  cin >> n;
  assert(n <= 10);
  MCMF F(2 * n);
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
      int k;
      cin >> k;
      F.add_edge(i, j + n, 1, k, i * 20
          + j);
    }
  }
  int s = 2 * n + 1, t = s + 1;
  for (int i = 0; i < n; i++) {
    F.add_edge(s, i, 1, 0);
    F.add_edge(i + n, t, 1, 0);
  }
  auto ans = F.solve(s, t).second;
  long long w = 0;
  set<int> se;
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
      int p = i * 20 + j;
      if (F.flow_through[p] > 0) {
        se.insert(j);
        w += F.flow_through[p];
      }
    }
  }
  assert(se.size() == n && w == n);
  cout << ans << '\n';
  return 0;
}
```

4.12 SCC

```
unordered_map<ll, vector<ll>> adj,
    InvAdj;
stack<ll> order;
unordered_map<ll, bool> visited;
unordered_map<ll, vector<ll>> all_scc;
unordered_map<ll, ll> compId;
void dfs_for_start(ll curr) {
  visited[curr] = 1;
  for (auto i : adj[curr])
    if (!visited[i])
      dfs_for_start(i);
  order.push(curr);
}
vector<ll> curr_comp;
void dfs_for_scc(ll curr) {
  visited[curr] = 1;
  for (auto i : InvAdj[curr])
    if (!visited[i])
      dfs_for_scc(i);
  curr_comp.push_back(curr);
}
inline void scc() {
  ll n, e, u, v;
  cin >> n >> e;
  for (ll i = 0; i < e; i++) {
    cin >> u >> v;
    adj[u].push_back(v);
    InvAdj[v].push_back(u);
  }
  for (ll i = 1; i <= n; i++)
    if (!visited[i])
      dfs_for_start(i);
  visited.clear();
  while (!order.empty()) {
    if (!visited[order.top()]) {
      curr_comp.clear();
      dfs_for_scc(order.top());
      ll sz = all_scc.size() + 1;
      all_scc[sz] = curr_comp;
      for (auto i : curr_comp)
        compId[i] = sz;
    }
    order.pop();
  }
}
// no.of ways and min cost of
    connecting the sccs
const ll MOD = 1e9 + 7, N = 1e5 + 2,
    INF = 1e18 + 2;
ll n, m, comp[N];
vector<ll> adj[N], rev[N];
bitset<N> vis;
void DFS1(ll u, stack<ll> &TS) {
  vis[u] = true;
  for (ll v : adj[u])
    if (!vis[v])
      DFS1(v, TS);
  TS.push(u);
}
void DFS2(ll u, const ll scc_no, ll &
    min_cost, ll &ways, vector<ll> &cost
    ) {
  vis[u] = true;
  comp[u] = scc_no;
  for (ll v : rev[u])
    if (!vis[v]) {
      if (min_cost == cost[v])
        ++ways;
      else if (min_cost > cost[v]) {
        ways = 1;
        min_cost = cost[v];
      }
      DFS2(v, scc_no, min_cost, ways,
          cost);
    }
}
signed main() {
  FIO cin >> n;
  vector<ll> cost(n + 1);
  for (ll i = 1; i <= n; ++i)
    cin >> cost[i];
  cin >> m;
  while (m--) {
    ll u, v;
    cin >> u >> v;
    adj[u].push_back(v);
    rev[v].push_back(u);
  }
  ll tot = 0, ways = 1;
  stack<ll> TS;
  for (ll i = 1; i <= n; ++i)
    if (!vis[i])
      DFS1(i, TS);
  vis.reset();
  ll scc_no = 0;
  while (!TS.empty()) {
    ll u = TS.top();
    TS.pop();
    if (!vis[u]) {
      ll tmp_cst = cost[u], tmp_ways =
          1;
      DFS2(u, ++scc_no, tmp_cst,
          tmp_ways, cost);
      tot += tmp_cst;
      ways = (ways * tmp_ways) % MOD;
    }
  }
  cout << tot << ' ' << ways;
} // TC: O(V+E)
```

4.13 0-1 BFS

```
vector<ll> d(n, INF);
d[s] = 0;
deque<ll> q;
q.push_front(s);
while (!q.empty()) {
  ll v = q.front();
  q.pop_front();
  for (auto edge : adj[v]) {
    ll u = edge.first;
    ll w = edge.second;
    if (d[v] + w < d[u]) {
      d[u] = d[v] + w;
      if (w == 1)
        q.push_back(u);
      else
        q.push_front(u);
    }
  }
}
```

4.14 Hull

```
// Convex Hull
#pragma GCC target("avx2")
#pragma GCC optimize("O3")
#pragma GCC optimize("unroll-loops")
#include <bits/stdc++.h>
using namespace std;

typedef long long int ll;
typedef long double ld;
typedef pair<ll, ll> pl;
typedef vector<ll> vl;
typedef complex<ll> pt;

#define G(x) ll x; cin >> x;
#define F(i, l, r) for (ll i = l; i < (
    r); ++i)
#define A(a) (a).begin(), (a).end()
#define CRS(a, b) (conj(a) * (b)).Y
#define K first
#define V second
#define X real()
#define Y imag()
#define N 100010

namespace std {
bool operator<(pt a, pt b) { return a.X
    == b.X ? a.Y < b.Y : a.X < b.X; }
} // namespace std

bool in_hull(pt p, vector<pt> &hu,
    vector<pt> &hd) {
  if (p == *hu.begin() || p == *hd.
      begin())
    return false; // change to true if
        border counts as inside
  if (p < *hu.begin() || *hd.begin() <
      p)
    return false;
  auto u = upper_bound(A(hu), p);
  auto d = lower_bound(hd.rbegin(), hd.
      rend(), p);
  return CRS(*u - p, *(u - 1) - p) > 0
      &&
          CRS(*(d - 1) - p, *d - p) >
              0; // change to >= if
                 border counts as "
                 inside"
}

void do_hull(vector<pt> &pts, vector<pt
    > &h) {
  for (pt p : pts) {
    while (h.size() > 1 && CRS(h.back()
        - p, h[h.size() - 2] - p) <=
                                  0) //
                                     change
                                      to <
                                      0 if
                                     border
                                     points
                                     included
      h.pop_back();
    h.push_back(p);
  }
}

pair<vector<pt>, vector<pt>> get_hull(
    vector<pt> &pts) {
  vector<pt> hu, hd;
  sort(A(pts)), do_hull(pts, hu);
  reverse(A(pts)), do_hull(pts, hd);
  return {hu, hd};
}

vector<pt> full_hull(vector<pt> &pts) {
  auto h = get_hull(pts);
  h.K.pop_back(), h.V.pop_back();
  for (pt p : h.V)
    h.K.push_back(p);
  return h.K;
}

int main() {
  G(n) vector<pt> v;
  F(i, 0, n) { G(x) G(y) v.push_back({x
      , y}); }
  vector<pt> h = full_hull(v);
}
```

4.15 Dynamic Hull

```
// Dynamic Convex Hull
#pragma GCC target("avx2")
#pragma GCC optimize("O3")
#pragma GCC optimize("unroll-loops")
#include <bits/stdc++.h>
using namespace std;

typedef long long int ll;
typedef long double ld;
typedef pair<ll, ll> pl;
typedef vector<ll> vl;
typedef complex<ll> pt;

#define G(x) ll x; cin >> x;
#define F(i, l, r) for (ll i = l; i < (
    r); ++i)
#define A(a) (a).begin(), (a).end()
#define CRS(a, b) (conj(a) * (b)).Y
#define X real()
#define Y imag()
#define N 100010

namespace std {
bool operator<(pt a, pt b) { return a.X
    == b.X ? a.Y < b.Y : a.X < b.X; }
} // namespace std

// helper function for dyn_in_hull
bool in(pt p, set<pt> &h) {
  if (h.empty() || p < *h.begin() || *h
      .rbegin() < p)
    return false;
  auto i = h.upper_bound(p), j = i--;
  return CRS(*j - p, *i - p) > 0; //
      change to >= if border counts as "
      inside"
}

// returns true if p contained in
    dynamic hull hu / hd
bool in_hull(pt p, set<pt> &hu, set<pt>
    &hd) { return in(p, hu) && in(-p,
    hd); }

// helper function for dyn_add
void fix_bad(set<pt>::iterator i, set<
    pt> &h, bool l) {
  if (i == --h.begin() || i == h.end())
    return;
  pt p = *i;
  h.erase(p);
  if (!in(p, h))
    h.insert(p);
  else
    fix_bad(l ? --h.lower_bound(p) : h.
```

```
        upper_bound(p), h, l);
}

// helper function for dyn_add_to_hull
void add(pt p, set<pt> &h) {
  if (in(p, h))
    return;
  h.insert(p);
  fix_bad(--h.lower_bound(p), h, true);
  fix_bad(h.upper_bound(p), h, false);
}

// adds p to dynamic hull hu / hd
void add_to_hull(pt p, set<pt> &hu, set
    <pt> &hd) { add(p, hu), add(-p, hd);
     }

int main() {
  G(n) set<pt> hu, hd;
  F(i, 0, n) { G(x) G(y) add_to_hull({x
      , y}, hu, hd); }
}
```

4.16 Count SImple Cycle

```
void findNumberOfSimpleCycles(int N,
    vector<vector<int>> adj) {
  int ans = 0;
  int dp[(1 << N)][N];
  memset(dp, 0, sizeof dp);
  for (int mask = 0; mask < (1 << N);
      mask++) {
    int nodeSet = __builtin_popcountll(
        mask);
    int firstSetBit = __builtin_ffsl(
        mask);
    if (nodeSet == 1)
      dp[mask][firstSetBit] = 1;
    else {
      for (int j = firstSetBit + 1; j <
          N; j++) {
        if ((mask & (1 << j))) {
          int newNodeSet = mask ^ (1 <<
              j);
          for (int k = 0; k < N; k++) {
            if ((newNodeSet & (1 << k))
                && adj[k][j]) {
              dp[mask][j] += dp[
                  newNodeSet][k];
              if (adj[j][firstSetBit]
                  && nodeSet > 2)
                ans += dp[mask][j];
            }
          }
        }
      }
    }
  }
  cout << ans << endl;
}
```

5 Misc

5.1 Max Pos and Next Greater

```
const ll MXX = 1e5 + 5;
ll mxtree[4 * MXX], arr[MXX];
void maxtree(ll idx, ll left, ll right)
    {
  if (left == right)
    mxtree[idx] = left;
  else {
    ll mid = (left + right) / 2;
    maxtree(idx * 2, left, mid);
    maxtree(idx * 2 + 1, mid + 1, right
        );
    ll left = mxtree[idx * 2];
    ll right = mxtree[idx * 2 + 1];
    if (arr[left] < arr[right])
      mxtree[idx] = right;
    else
      mxtree[idx] = left;
  }
}
ll mxPos(ll idx, ll tleft, ll tright,
    ll qleft, ll qright) {
  if (qleft > qright)
    return -1;
  if (qleft == tleft and qright ==
      tright)
    return mxtree[idx];
  ll tmid = (tleft + tright) / 2;
  ll left = mxPos(idx * 2, tleft, tmid,
      qleft, min(qright, tmid));
  ll right = mxPos(idx * 2 + 1, tmid +
      1, tright, max(qleft, tmid + 1),
      qright);
  ll ans;
  if (left == -1)
    ans = right;
  else if (right == -1)
    ans = left;
  else if (arr[left] < arr[right])
    ans = right;
  else
    ans = left;
  return ans;
}
ll main() {
  ll t = 1, n, q, a, b;
  cin >> t;
  while (t--) {
    cin >> n >> q;
    for (ll i = 0; i < n; i++)
      cin >> arr[i];
    stack<ll> stk;
    ll nge[n];
    stk.push(0);
    for (ll i = 1; i < n; i++) {
      while (stk.size() and arr[stk.top
          ()] < arr[i]) {
        nge[stk.top()] = i;
        stk.pop();
      }
      stk.push(i);
    }
    while (stk.size()) {
      nge[stk.top()] = n;
      stk.pop();
    }
    ll ans[n];
    ans[n - 1] = 0;
    for (ll i = n - 2; i >= 0; i--) {
      ll tmp = nge[i];
      if (tmp == n)
        ans[i] = 0;
      else
        ans[i] = ans[tmp] + 1;
    }
    maxtree(1, 0, n - 1);
    for (ll i = 0; i < q; i++) {
      cin >> a >> b;
      if (a > b)
        swap(a, b);
      cout << ans[mxPos(1, 0, n - 1, a
          - 1, b - 1)] << "\n";
    }
  }
}
```

5.2 Knight Move

```
ll X[8]={2,1,-1,-2,-2,-1,1,2};
ll Y[8]={1,2,2,1,-1,-2,-2,-1};
```

5.3 MatrixExpo

```
typedef long long LL;
LL arr[60][60], res[60][60], tmp
    [60][60], m;

void matMul(LL a[][60], LL b[][60], LL
    mod) {
  for (ll i = 0; i < m; i++)
    for (ll j = 0; j < m; j++) {
      tmp[i][j] = 0;
```

```
      for (ll k = 0; k < m; k++) {
        tmp[i][j] += (a[i][k] * b[k][j
            ]) % mod;
        tmp[i][j] %= mod;
      }
    }
}
void power(LL n, LL mod) {
  for (ll i = 0; i < m; i++)
    for (ll j = 0; j < m; j++)
      if (i == j)
        res[i][j] = 1;
      else
        res[i][j] = 0;

  while (n) {
    if (n & 1) {
      matMul(res, arr, mod);
      for (ll i = 0; i < m; i++)
        for (ll j = 0; j < m; j++)
          res[i][j] = tmp[i][j];
      n--;
    } else {
      matMul(arr, arr, mod);
      for (ll i = 0; i < m; i++)
        for (ll j = 0; j < m; j++)
          arr[i][j] = tmp[i][j];
      n /= 2;
    }
  }
}
```

5.4 Ternary Search

```
double ternary_search(double l, double
    r) {
  double eps = 1e-9; // error limit
  while (r - l > eps) {
    double m1 = l + (r - l) / 3, m2 = r
        - (r - l) / 3;
    double f1 = f(m1), f2 = f(m2); //
        evaluates the function at m1, m2
    if (f1 < f2)
      l = m1;
    else
      r = m2;
  }
  return f(l); // return the maximum of
      f(x) in [l, r]
}
```

6 Number Theory

6.1 Leap_year

```
bool isLeap(ll n) {
  if (n % 100 == 0)
    return (n % 400 == 0);
  else
    return (n % 4 == 0);
}
// leap year between l and r
ll calNum(ll y) { return (y / 4) - (y /
    100) + (y / 400); }
ll leapNum(ll l, ll r) { return calNum(
    r) - calNum(--l); }
```

6.2 Two Line Intersection

```
ll cross(ll x1, ll y1, ll x2, ll y2, ll
    x3, ll y3) {
  return (x2 - x1) * (y3 - y1) - (y2 -
      y1) * (x3 - x1);
}
bool intersect(ll x1, ll y1, ll x2, ll
    y2, ll x3, ll y3, ll x4, ll y4) {
  ll c1 = cross(x1, y1, x2, y2, x3, y3)
      , c2 = cross(x1, y1, x2, y2, x4,
      y4),
      c3 = cross(x3, y3, x4, y4, x1, y1)
          , c4 = cross(x3, y3, x4, y4, x2
          , y2);
  if ((!c1 && min(x1, x2) <= x3 && x3
      <= max(x1, x2) && min(y1, y2) <=
      y3 &&
        y3 <= max(y1, y2)) |
       (!c2 && min(x1, x2) <= x4 && x4
          <= max(x1, x2) && min(y1, y2)
          <= y4 &&
        y4 <= max(y1, y2)) |
       (!c3 && min(x3, x4) <= x1 && x1
          <= max(x3, x4) && min(y3, y4)
          <= y1 &&
        y1 <= max(y3, y4)) |
       (!c4 && min(x3, x4) <= x2 && x2
          <= max(x3, x4) && min(y3, y4)
          <= y2 &&
        y2 <= max(y3, y4)))
    return true;
  return (c1 > 0) != (c2 > 0) && (c3 >
      0) != (c4 > 0);
}
```

6.3 Binary_exponentiation

```
ll binaryExp(ll base, ll power, ll MOD
    = mod) {
  ll res = 1;
  while (power) {
    if (power & 1)
      res = (res * base) % MOD;
    base = ((base % MOD) * (base % MOD)
        ) % MOD;
    power /= 2;
  }
  return res;
}

/*
task: a ^ b ^ c
binaryExp(a, binaryExp(b, c, mod - 1),
    mod)
*/
```

6.4 Count_divisor

```
ll maxVal = 1e6 + 1;
vector<ll> countDivisor(maxVal, 0);
void countingDivisor() {
  for (ll i = 1; i < maxVal; i++)
    for (ll j = i; j < maxVal; j += i)
      countDivisor[j]++;
}
// TC: nlog(n)
// count the number of divisors of all
    numbers in a range.
```

6.5 Check_prime

```
bool prime(ll n) {
  if (n < 2)
    return false;
  if (n <= 3)
    return true;
  if (!(n % 2) || !(n % 3))
    return false;
  for (ll i = 5; i * i <= n; i += 6) {
    if (!(n % i) || !(n % (i + 2)))
      return false;
  }
  return true;
}
// TC: sqrt(n) / 6;
```

6.6 SPF

```
// smallest prime factor using seive
const ll N = 1e7 + 5;
ll spf[N];
void smallestPrimeFactorUsingSeive() {
  for (ll i = 2; i < N; i++) {
    if (spf[i] == 0) {
      for (ll j = i; j < N; j += i) {
        if (spf[j] == 0)
```

```
          spf[j] = i;
      }
    }
  }
}
// smallest factor of a number
ll factor(ll n) {
  ll a;
  if (n % 2 == 0)
    return 2;
  for (a = 3; a * a <= n; a += 2) {
    if (n % a == 0)
      return a;
  }
  return n;
}
// complete factorization
ll r;
while (n > 1) {
  r = factor(n);
  cout << r << '\n';
  n /= r;
}
```

6.7 Seive

```
const ll N = 1e7 + 5;
ll prime[N];
void sieveOfEratosthenes() {
  for (ll i = 2; i < N; i++)
    prime[i] = 1;
  for (ll i = 4; i < N; i += 2)
    prime[i] = 0;
  for (ll i = 3; i * i < N; i++) {
    if (prime[i]) {
      for (ll j = i * i; j < N; j += i
          * 2)
        prime[j] = 0;
    }
  }
}
```

6.8 Optimize_seive

```
vector<ll> sieve(const ll N, const ll Q
    = 17, const ll L = 1 << 15) {
  static const ll rs[] = {1, 7, 11, 13,
      17, 19, 23, 29};
  struct P {
    P(ll p) : p(p) {}
    ll p;
    ll pos[8];
  };
  auto approx_prime_count = [](const ll
      N) -> ll {
    return N > 60184 ? N / (log(N) -
        1.1) : max(1., N / (log(N) -
        1.11)) + 1;
  };
  const ll v = sqrt(N), vv = sqrt(v);
  vector<bool> isp(v + 1, true);
  for (ll i = 2; i <= vv; ++i)
    if (isp[i]) {
      for (ll j = i * i; j <= v; j += i
          )
        isp[j] = false;
    }
  const ll rsize = approx_prime_count(N
      + 30);
  vector<ll> primes = {2, 3, 5};
  ll psize = 3;
  primes.resize(rsize);

  vector<P> sprimes;
  size_t pbeg = 0;
  ll prod = 1;
  for (ll p = 7; p <= v; ++p) {
    if (!isp[p])
      continue;
    if (p <= Q)
      prod *= p, ++pbeg, primes[psize
          ++] = p;
    auto pp = P(p);
    for (ll t = 0; t < 8; ++t) {
      ll j = (p <= Q) ? p : p * p;
      while (j % 30 != rs[t])
        j += p << 1;
      pp.pos[t] = j / 30;
    }
    sprimes.push_back(pp);
  }
  vector<unsigned char> pre(prod, 0xFF)
      ;
  for (size_t pi = 0; pi < pbeg; ++pi)
      {
    auto pp = sprimes[pi];
    const ll p = pp.p;
    for (ll t = 0; t < 8; ++t) {
      const unsigned char m = ~(1 << t)
          ;
      for (ll i = pp.pos[t]; i < prod;
          i += p)
        pre[i] &= m;
    }
  }
  const ll block_size = (L + prod - 1)
      / prod * prod;
  vector<unsigned char> block(
      block_size);
  unsigned char *pblock = block.data();
  const ll M = (N + 29) / 30;
  for (ll beg = 0; beg < M; beg +=
      block_size, pblock -= block_size)
      {
    ll end = min(M, beg + block_size);
    for (ll i = beg; i < end; i += prod
        ) {
      copy(pre.begin(), pre.end(),
          pblock + i);
    }
    if (beg == 0)
      pblock[0] &= 0xFE;
    for (size_t pi = pbeg; pi < sprimes
        .size(); ++pi) {
      auto &pp = sprimes[pi];
      const ll p = pp.p;
      for (ll t = 0; t < 8; ++t) {
        ll i = pp.pos[t];
        const unsigned char m = ~(1 <<
            t);
        for (; i < end; i += p)
          pblock[i] &= m;
        pp.pos[t] = i;
      }
    }
    for (ll i = beg; i < end; ++i) {
      for (ll m = pblock[i]; m > 0; m
          &= m - 1) {
        primes[psize++] = i * 30 + rs[
            __builtin_ctz(m)];
      }
    }
  }
  assert(psize <= rsize);
  while (psize > 0 && primes[psize - 1]
      > N)
    --psize;
  primes.resize(psize);
  return primes;
}
// it takes 500ms for generating prime
    upto 1e9
```

6.9 nth_prime_number

```
vector<ll> nth_prime;
const ll MX = 86200005;
bitset<MX> visited;
void optimized_prime() {
  nth_prime.push_back(2);
  for (ll i = 3; i < MX; i += 2) {
    if (visited[i])
      continue;
    nth_prime.push_back(i);
    if (lll * i * i > MX)
```

```
      continue;
    for (ll j = i * i; j < MX; j += i +
        i)
      visited[j] = true;
  }
}
```

6.10 nCr

```
// 1:
// more space, less time
const ll MAX = 1e7 + 5;
vector<ll> fact(MAX), ifact(MAX), inv(
    MAX);
void factorial() {
  inv[1] = fact[0] = ifact[0] = 1;
  for (ll i = 2; i < MAX; i++)
    inv[i] = inv[mod % i] * (mod - mod
        / i) % mod;
  for (ll i = 1; i < MAX; i++)
    fact[i] = (fact[i - 1] * i) % mod;
  for (ll i = 1; i < MAX; i++)
    ifact[i] = ifact[i - 1] * inv[i] %
        mod;
}
ll nCr(ll n, ll r) {
  if (r < 0 || r > n)
    return 0;
  return (ll)fact[n] * ifact[r] % mod *
      ifact[n - r] % mod;
}
// 2:
// less space, more time
const ll MAX = 1e7 + 10;
vector<ll> fact(MAX), inv(MAX);
void factorial() {
  fact[0] = 1;
  for (ll i = 1; i < MAX; i++)
    fact[i] = (i * fact[i - 1]) % mod;
}
ll binaryExp(ll a, ll n, ll M = mod){};
    // needs to implement
void inverse() {
  for (ll i = 0; i < MAX; ++i)
    inv[i] = binaryExp(fact[i], mod -
        2);
}
ll nCr(ll a, ll b) {
  if (a < b or a < 0 or b < 0)
    return 0;
  ll de = (inv[b] * inv[a - b]) % mod;
  return (fact[a] * de) % mod;
}
// 3:
// nCr mod m where m is not prime
ll C_mod_p(ll n, ll k, ll p) {
  if (k > n)
    return 0;
  vector<ll> fac(p);
  fac[0] = 1;
  for (int i = 1; i < p; i++)
    fac[i] = fac[i - 1] * i % p;

  ll res = 1;
  while (n || k) {
    ll ni = n % p, ki = k % p;
    if (ki > ni)
      return 0;
    res = res * fac[ni] % p * modInv(
        fac[ki], p) % p * modInv(fac[ni
        - ki], p) %
          p;
    n /= p;
    k /= p;
  }
  return res;
}
// compute nCr mod composite m (non-
    prime)
ll nCr_mod_m(ll n, ll k, ll m) {
  // Step 1: factorize m
  vector<int> primes;
  int tmp = m;
  for (int i = 2; i * i <= tmp; i++) {
    if (tmp % i == 0) {
      primes.push_back(i);
      while (tmp % i == 0)
        tmp /= i;
    }
  }
  if (tmp > 1)
    primes.push_back(tmp);

  // Step 2: compute result mod each
      prime
  vector<ll> rem, mod;
  for (int p : primes) {
    rem.push_back(C_mod_p(n, k, p));
    mod.push_back(p);
  }
  // Step 3: Chinese Remainder Theorem
      (combine)
  ll res = 0;
  for (int i = 0; i < (int)mod.size();
      i++) {
    ll Mi = m / mod[i];
    ll invMi = binaryExp(Mi, mod[i] -
        2, mod[i]); // modular inverse
    res = (res + rem[i] * Mi % m *
        invMi % m) % m;
  }
  return res;
}
```

6.11 Factorial_mod

```
// n! mod p : Here P is mod value
// For binaryExp we call 1.6 function
ll factmod(ll n, ll p) {
  ll res = 1;
  while (n > 1) {
    res = (res * binaryExp(p - 1, n / p
        , p)) % p;
    for (ll i = 2; i <= n % p; ++i)
      res = (res * i) % p;
    n /= p;
  }
  return (res % p);
}
```

6.12 PHI

```
// the positive integers less than or
    equal to n that are relatively prime
     to n.
ll phi(ll n) {
  ll result = n;
  for (ll i = 2; i * i <= n; i++) {
    if (n % i == 0) {
      while (n % i == 0)
        n /= i;
      result -= result / i;
    }
  }
  if (n > 1)
    result -= result / n;
  return result;
}
// PHI of 1 to N
const int N = 1e6 + 9;
int phi[N];
int phiS[N];
void totient() {
  for (int i = 1; i < N; i++)
    phi[i] = i;
  for (int i = 2; i < N; i++) {
    if (phi[i] == i) {
      for (int j = i; j < N; j += i)
        phi[j] -= phi[j] / i;
    }
  }
  phiS[0] = phi[0];
  for (int i = 1; i < N; i++)
    phiS[i] = phiS[i - 1] + phi[i];
}
```

6.13 Catalan

```
void catalan(ll n) {
  ll res = 1;
  cout << res << " ";
  for (ll i = 1; i < n; i++) {
    res = (res * (4 * i - 2)) / (i + 1)
        ;
    cout << res << " ";
  }
}
```

6.14 Extended_GCD

```
// return {x,y} such that ax + by = gcd
    (a,b)
ll extended_euclid(ll a, ll b, ll &x,
    ll &y) {
  if (b == 0) {
    x = 1;
    y = 0;
    return a;
  }
  ll x1, y1;
  ll d = extended_euclid(b, a % b, x1,
      y1);
  x = y1;
  y = x1 - y1 * (a / b);
  return d;
}
ll inverse(ll a, ll m) {
  ll x, y;
  ll g = extended_euclid(a, m, x, y);
  if (g != 1)
    return -1;
  return (x % m + m) % m;
}
```

6.15 Large Mod

```
ll mod(string &num, ll a) {
  ll res = 0;
  for (ll i = 0; i < num.length(); i++)
    res = (res * 10 + num[i] - '0') % a
        ;
  return res;
}
```

6.16 Factorial_Divisor

```
ll factorialDivisors(ll n) {
  ll result = 1;
  for (ll i = 0; i < allPrimes.size();
      i++) {
    ll p = allPrimes[i];
    ll exp = 0;
    while (p <= n) {
      exp = exp + (n / p);
      p = p * allPrimes[i];
    }
    result = result * (exp + 1);
  }
  return result;
}
```

6.17 Number_conversion

```
// 10 - ary to m - ary
char a[16] = {'0', '1', '2', '3', '4',
    '5', '6', '7', '8', '9', 'A', 'B', '
    C', 'D', 'E', 'F'};
string tenToM(ll n, ll m) {
  ll temp = n;
  string result = "";
  while (temp != 0) {
    result = a[temp % m] + result;
    temp /= m;
  }
  return result;
}
// m - ary to 10 - ary
string num = "0123456789ABCDE";
ll mToTen(string n, ll m) {
  ll multi = 1;
  ll result = 0;
  for (ll i = n.size() - 1; i >= 0; i
      --) {
    result += num.find(n[i]) * multi;
    multi *= m;
  }
  return result;
}
```

6.18 Number_of_1_in_bit_till_N

```
ll cntOnes(ll n) {
  ll cnt = 0;
  for (ll i = 1; i <= n; i <<= 1) {
    ll x = (n + 1) / (i << 1);
    cnt += x * i;
    if ((n + 1) % i && n & i)
      cnt += (n + 1) % i;
  }
  return cnt;
}
```

6.19 Disarrangement

```
ll disarrange(ll n) {
  if (n == 1)
    return 0;
  if (n == 2)
    return 1;
  return (n - 1) * (disarrange(n - 1) +
      disarrange(n - 2));
}
// D(n) = (n!)/e
```

6.20 Millar_Rabin

```
bool check_composite(ll n, ll a, ll d,
    ll s) {
  ll x = binaryExp(a, d, n);
  if (x == 1 || x == n - 1)
    return false;
  for (ll r = 1; r < s; r++) {
    x = (u128)x * x % n;
    if (x == n - 1)
      return false;
  }
  return true;
};
bool MillerRabin(ll n, ll iter = 5) {
  // returns true if n is probably
     prime, else returns false.
  if (n < 4)
    return n == 2 || n == 3;
  ll s = 0;
  ll d = n - 1;
  while ((d & 1) == 0) {
    d >>= 1;
    s++;
  }
  for (ll i = 0; i < iter; i++) {
    ll a = 2 + rand() % (n - 3);
    if (check_composite(n, a, d, s))
      return false;
  }
  return true;
}
```

6.21 Modular_operation

```
// Addition :
ll mod_add(ll a, ll b, ll MOD = mod) {
  a = a % MOD, b = b % MOD;
  return (((a + b) % MOD) + MOD) % MOD;
}
// Subtraction :
ll mod_sub(ll a, ll b, ll MOD = mod) {
  a = a % MOD, b = b % MOD;
  return (((a - b) % MOD) + MOD) % MOD;
}
// Multiplication :
ll mod_mul(ll a, ll b, ll MOD = mod) {
```

```
  a = a % MOD, b = b % MOD;
  return (((a * b) % MOD) + MOD) % MOD;
}
// Division :
ll mminvprime(ll a, ll b) { return
    binaryExp(a, b - 2, b); }
ll mod_div(ll a, ll b, ll MOD = mod) {
  a = a % MOD, b = b % MOD;
  return (mod_mul(a, mminvprime(b, MOD)
      , MOD) + MOD) % MOD;
}
```

6.22 MSLCM

```
// For a given number N, maximum sum
    LCM indicates the set of numbers
    whose LCM
// is N and summation is maximum. Let,
    MSLCM(N) denote this maximum sum of
// numbers. Given the value of N you
    will have to find the value:
    sumation of
// MSLCM(i) from i to 2n
ll MSLCM(ll n) {
  ll l = 1, r, val, ret = 0;
  while (l <= n) {
    val = n / l, r = n / val;
    ret += val * ((l + r) * (r - l + 1)
        / 2);
    l = r + 1;
  }
  return ret - 1;
}
```

6.23 Find numbers in between [L, R] which are divisible by all Array elements

```
void solve(ll *arr, ll N, ll L, ll R) {
  ll LCM = arr[0];
  for (ll i = 1; i < N; i++) {
    LCM = (LCM * arr[i]) / (__gcd(LCM,
        arr[i]));
  }
  if ((LCM < L && LCM * 2 > R) || LCM >
      R) {
    return;
  }
  ll k = (L / LCM) * LCM;
  if (k < L)
    k = k + LCM;
  for (ll i = k; i <= R; i = i + LCM)
    cout << i << ' ';
}
```

Maximum Value (N)	Number with Most Divisors (n)	Number of Divisors ($\tau(n)$)
10^3	83,160	128
10^6	720,720	240
10^7	$9,609,600$	640
10^8	$98,280,000$	672
10^9	$735,134,400$	1,344
10^{10}	$7,242,460,800$	2,688
10^{11}	$73,346,256,000$	5,376
10^{12}	$936,966,912,400$	10,752

7 Mathematics

7.1 Area Formulas

Rectangle	length × width
Square	side^2
Triangle	$\frac{1}{2}$ × base × height
Parallelogram	base × height
Pyramid (no base)	$\frac{1}{2}$×(perimeter of base)×(slant height)
Polygon	$\frac{1}{2}\left\lvert\sum_{i=1}^{n}(x_i y_{i+1} - x_{i+1}y_i)\right\rvert$ $a + \frac{b}{2} - 1$ (for lattice coordinates)

a = interior lattice pts, b = boundary pts.

7.2 Volume Formulas

Cube	side^3
Rectangular Prism	length × width × height
Cylinder	π × radius^2 × height
Sphere	$\frac{4}{3}\pi$ × radius^3
Pyramid	$\frac{1}{3}$×(base area)×(height)

7.3 Surface Area Formulas

Cube	$6 \times \text{side}^2$
Rectangular Prism	$2(lw + lh + wh)$ (l = length, w = width, h = height)
Cylinder	$2\pi r(r + h)$
Sphere	$4\pi r^2$
Pyramid	base area $+ \frac{1}{2} \times$ (perimeter) × (slant height)

7.4 Triangles

Semiperimeter	$s = \frac{a+b+c}{2}$
Area	$A = \sqrt{s(s-a)(s-b)(s-c)}$
Circumradius	$R = \frac{abc}{4A}$
Inradius	$r = \frac{A}{s}$
Median	$m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$
Angle bisector	$s_a = \sqrt{\frac{bc}{1-(a/(b+c))^2}}$

Side lengths: a, b, c.

7.5 Sum Equations

$$\sum_{i=k}^{n} c^i = \frac{c^{n+1} - c^k}{c-1} \qquad \sum_{i=1}^{n} i = \frac{n(n+1)}{2}$$

$$\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6} \qquad \sum_{i=1}^{n} i^3 = \left(\frac{n(n+1)}{2}\right)^2$$

$$\sum_{i=1}^{n} i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30} \qquad \sum_{i=1}^{n}(2i-1) = n^2$$

7.6 Trigonometry

Sine law	$\frac{\sin\alpha}{a} = \frac{\sin\beta}{b} = \frac{\sin\gamma}{c} = \frac{1}{2R}$
Cosine law	$a^2 = b^2 + c^2 - 2bc\cos\alpha$
Tangent law	$\frac{a+b}{a-b} = \frac{\tan(\frac{\alpha+\beta}{2})}{\tan(\frac{\alpha-\beta}{2})}$
$\sin(A \pm B)$	$\sin A\cos B \pm \cos A\sin B$
$\cos(A \pm B)$	$\cos A\cos B \mp \sin A\sin B$
$\tan(A \pm B)$	$\frac{\tan A \pm \tan B}{1 \mp \tan A\tan B}$
$\sin 2\theta$	$2\sin\theta\cos\theta$
$\cos 2\theta$	$\cos^2\theta - \sin^2\theta$
$\tan 2\theta$	$\frac{2\tan\theta}{1-\tan^2\theta}$
Half-angle	$\sin(\frac{\theta}{2}) = \pm\sqrt{\frac{1-\cos\theta}{2}}$, $\cos(\frac{\theta}{2}) = \pm\sqrt{\frac{1+\cos\theta}{2}}$, $\tan(\frac{\theta}{2}) = \frac{1-\cos\theta}{\sin\theta}$
Sum identities	$\sin r + \sin w = 2\sin(\frac{r+w}{2})\cos(\frac{r-w}{2})$ $\cos r + \cos w = 2\cos(\frac{r+w}{2})\cos(\frac{r-w}{2})$
General	$(V + W)\tan\left(\frac{r-w}{2}\right) = (V - W)\tan\left(\frac{r+w}{2}\right)$
	$a\cos x + b\sin x = r\cos(x-\varphi)$
	$a\sin x - b\cos x = r\sin(x-\varphi)$
	where $r = \sqrt{a^2+b^2}$, $\varphi = \operatorname{atan2}(b, a)$

7.7 Logarithmic Basic

$\log_b 1 = 0$	$\log_b b = 0$
$b^{\log_b a} = a$	$x^{\log_b y} = y^{\log_b x}$
$\log_a b = \frac{1}{\log_b a}$	$\log_a x = \frac{\log_b x}{\log_b a}$
$\log_b(AB) = \log_b A + \log_b B$	
$\log_b(\frac{A}{B}) = \log_b A - \log_b B$	
$\log_a c = \log_a b \times \log_b c$	
$\log_b(A^x) = x\log_b A$	

7.8 Series

Catalan: $C_n = \frac{1}{n+1}\binom{2n}{n}$, $C_n = \sum_{k=0}^{n} C_k C_{n-k}$
Arithmetic: $a_n = a + (n-1)d$, $s_n = \frac{n}{2}(2a + (n-1)d)$
Geometric: $a_n = ar^{n-1}$, $s_n = \frac{a(1-r^n)}{1-r}$
Derangements: $D_n = n!\sum_{k=0}^{n}\frac{(-1)^k}{k!}$, $D_n = \left\lfloor \frac{n!}{e} + \frac{1}{2} \right\rfloor$
Fibonacci: $f_n = \frac{\phi^n - (1-\phi)^n}{\sqrt{5}}$, $\phi = \frac{1+\sqrt{5}}{2}$

7.9 Pick's Theorem

$A = I - \frac{1}{2}B + 1$ (I = interior points, B = boundary points)

7.10 Stars and Bars

Number of solutions of $x_1 + \cdots + x_k = n$:
$\binom{n-1}{k-1}$ when $x_i > 0$; $\binom{n+k-1}{k-1}$ when $x_i \geq 0$.

7.11 Facts

$\lceil \frac{a}{b} \rceil = \lfloor \frac{a-1}{b} \rfloor + 1$
Sum l to r: $\frac{l+r}{2}(r - l + 1)$
$\left\lfloor \frac{\lfloor n/a \rfloor}{b} \right\rfloor = \lfloor \frac{n}{ab} \rfloor$

7.12 LCM

$\text{lcm}(a,n) + \text{lcm}(n-a,n) = \frac{n^2}{\gcd(a,n)}$

$\text{SUM} = \frac{n}{2}\left(\sum_{d|n} \varphi(d)d + 1\right)$