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Modified By YUSUF REZA HASNAT & ISTIAQUE AHMED ARIK

```
→Useful Things
→Fast I/O
C++:
ios base::sync with stdio(false),
cin.tie(nullptr), cout.tie(nullptr);
Python:
import sys
input = sys.stdin.readline
sys.stdout.write("----")
→ Random algorithm
#define accuracy
chrono::steady clock::now().time since epoch(
).count()
mt19937 rng(accuracy);
int rand(int 1, int r) {
   uniform int distribution<int> ludo(l, r);
   return ludo(rng);
→ Sublime Build Command
"shell cmd": "g++ -std=c++17 \"${file}\" -o
            \"${file base name}.exe\" &&
            \"${file path}/${file base name}
            .exe\" < in.txt > out.txt 2>
            error.txt",
"file regex": "^(...*?):([0-9]*):?([0-9]*)",
"working dir": "${file path}",
"selector": "source.c++, source.c",
"variants":
      "name": "Run with input/output",
      "shell cmd": "q++ \"${file}\" -o
      \"${file base name}.exe\" 2>
      error.txt && \"${file path}
      /${file base name}.exe\" < in.txt >
       out.txt 2>> error.txt"
  1
→ C++ code run without vscode extension
run a file which name is test.cpp
g++ .\test.cpp -o test && .\test.exe
```

1 Formula

1.1 Area Formulas

Rectangle: $Area = length \times width$ Square: $Area = side \times side$

Triangle: $Area = \frac{1}{2} \times base \times height$

Circle: $Area = \pi \times radius^2$

Parallelogram: $Area = base \times height$

Pyramid Base: $Area = \frac{1}{2} \times base \times slant \ height$

Polygon:

(a)
$$Area = \frac{1}{2} \left| \sum_{i=1}^{n-1} (x_i y_{i+1} - x_{i+1} y_i) \right|$$

(b) $Area = a + \frac{b}{2} - 1$ (for int coordinates) here, a=#int points inside polygon b=#int points outside polygon

1.2 Perimeter Formulas

Rectangle: $Perimeter = 2 \times (length + width)$

Square: $Perimeter = 4 \times side$

Triangle: $Perimeter = sum \ of \ all \ sides$ Circle: $Circumference = 2 \times \pi \times radius$

1.3 Volume Formulas

Cube: $Volume = side^3$

Rect Prism: $Volume = length \times width \times height$ Cylinder: $Volume = \pi \times radius^2 \times height$

Sphere: $Volume = \frac{4}{3} \times \pi \times radius^3$

Pyramid: $Volume = \frac{1}{3} \times base \ area \times height$

1.4 Surface Area Formulas

Cube: $Surface Area = 6 \times side^2$

Rectangular Prism:

Surface Area = $2 \times (length \times width + length \times height + width \times height)$

Cylinder:

Surface Area = $2 \times \pi \times \text{radius} \times (\text{radius} + \text{height})$

Sphere: $Surface Area = 4 \times \pi \times radius^2$

Pyramid:

Surface Area = base area +

 $\frac{1}{2}$ × perimeter of base × slant height

Triangles

Side lengths: a, b, c

Semiperimeter: $p = \frac{a+b+c}{2}$

Area: A = $\sqrt{(p(p-a)(p-b)(p-c))}$

Circumradius: $R = \frac{abc}{4A}$

Inradius: $r = \frac{A}{p}$

Length of median (divides triangle into two equal-area triangles):

$$ma = \frac{1}{2} * \sqrt{2b^2 + 2c^2 - a^2}$$

Length of bisector (divides angles in two):

$$sa = \sqrt{\frac{bc}{1 - \left(\frac{a}{b+c}\right)^2}}$$

1.5 Trigonometry

Law of sines:
$$sin \frac{\alpha}{a} = sin \frac{\beta}{b} = sin \frac{\gamma}{c} = \frac{1}{2R}$$

Law of cosines:
$$a^2 = b^2 + c^2 - 2bc \cos \alpha$$

Law of tangents:
$$\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$$

$$sin(A + B) = sin A cos B + cos A sin B$$

$$cos(A + B) = cos A cos B - sin A sin B$$

$$sin(A - B) = sin A cos B - cos A sin B$$

$$cos(A - B) = cos A cos B + sin A sin B$$

$$tan (A + B) = \frac{(tan A + tan B)}{(1 - tan A tan B)}$$

$$tan (A - B) = \frac{(tan A - tan B)}{(1 + tan A tan B)}$$

Double Angle and Half Angle Formulas: $\sin 2\theta = 2 \sin \theta \cos \theta$

$$\cos 2\theta = \cos^2 \theta - \sin^2 \theta$$

$$\tan 2\theta = \frac{(2 \tan \theta)}{(1 - \tan^2 \theta)}$$

$$sin\left(\frac{\theta}{2}\right) = \pm \sqrt{\frac{1-cos\theta}{2}}$$

$$cos\left(\frac{\theta}{2}\right) = \pm\sqrt{\frac{1+cos\theta}{2}}$$

 $tan\left(\frac{\theta}{2}\right) = \frac{(1-cos\theta)}{sin\theta}$

$$tan\left(\frac{\theta}{2}\right) = \frac{(1-\cos\theta)}{\sin\theta}$$

$$sin(r + w) = sin r cos w + cos r sin w$$

$$cos(r + w) = cos r cos w - sin r sin w$$

$$tan(r + w) = \frac{(tan r + tan w)}{(1 - tan r tan w)}$$

$$\sin r + \sin w = 2 \sin \left(\frac{r+w}{2}\right) \cos \left(\frac{r-w}{2}\right)$$

$$\cos r + \cos w = 2\cos\left(\frac{r+w}{2}\right)\cos\left(\frac{r-w}{2}\right)$$

$$(V + W) \tan(\frac{r-w}{2}) = (V - W) \tan(\frac{r+w}{2})$$

where V, W are lengths of sides opposite angles r, w.

$$a\cos x + b\sin x = r\cos(x - \varphi)$$

$$a \sin x - b \cos x = r \sin(x - \varphi)$$

where
$$r = \sqrt{a^2 + b^2}$$
, $\varphi = atan2(b, a)$

1.6 Sum

$$\frac{c^{k} + c^{k+1}}{c' + (c-1)} + \dots + c^{n} = c^{n+1} - c^{k}$$
// (c-1) for c \neq 1

$$1 + 2 + 3 + ... + n = \frac{n*(n+1)}{2}$$

$$1^{2} + 2^{2} + \ldots + n^{2} = \frac{n * (n+1) * (2n+1)}{6}$$

$$1^3 + 2^3 + \ldots + n^3 = \frac{n^2 * (n+1)^2}{4}$$

$$1^{4} + 2^{4} + \ldots + n^{4} = \frac{n * (n+1) * (2n+1) * (3n^{2} + 3n - 1)}{30}$$

sum of first n odd num = n^2

7 Logarithmic Basic

- $\bullet \quad \log_h 1 = 0$
- $\bullet \quad log_b b = 0$
- $log_h(AB) = log_h A + log_h B$
- $log_h(\frac{A}{B}) = log_h A log_h B$
- $\log_h A^x = x \log_h A$
- $log_a c = log_a b * log_b c$

- $\bullet \quad \log_a b = \frac{1}{\log_b a}$
- $\bullet \quad \log_a x = \frac{\log_b x}{\log_b a}$

1.8 Catalan Series

```
2 Number Theory
2.1 Prime number under 1000
        5 7 11 13 17 19 23 29 31 37
    3
41 43 47 53 59 61 67 71 73 79 83 89
97 101 103 107 109 113 127 131 137 139 149
151 157 163 167 173 179 181 191 193 197 199
211 223 227 229 233 239 241 251 257 263 269
271 277 281 283 293 307 311 313 317 331 337
347 349 353 359 367 373 379 383 389 397 401
409 419 421 431 433 439 443 449 457 461 463
467 479 487 491 499 503 509 521 523 541 547
557 563 569 571 577 587 593 599 601 607 613
617 619 631 641 643 647 653 659 661 673 677
683 691 701 709 719 727 733 739 743 751 757
761 769 773 787 797 809 811 821 823 827 829
839 853 857 859 863 877 881 883 887 907 911
919 929 937 941 947 953 967 971 977 983 991
997
2.2 Divisor Count
int maxVal = 1e6 + 1;
vector<int> countDivisor(maxVal, 0);
void countingDivisor(){
    for (int i = 1; i < maxVal; i++)
        for(int j= i; j<maxVal;j+= i)</pre>
            countDivisor[j]++;
// count the number of divisors of all
numbers in a range.
2.3 Leap year
bool isLeap(int n) {
    if (n%100==0)
        if (n%400==0) return true;
        else return false;
    if (n%4==0) return true;
    else return false;
2.4 Num of Leap year in between
int calNum(int year) {
return (year / 4) - (year / 100) +
 (year / 400);
int leapNum(int 1, int r) {
      return calNum(r) - calNum(l);
}
2.5 Print Calendar of any year
int dayNumber(int day, int month, int year) {
   static int t[]={0,3,2,5,0,3,5,1,4,6,2,4};
   year -= month < 3;
   return (year + year / 4 - year / 100 +
      year / 400 + t[month - 1] + day) % 7;
string getMonthName(int monthNumber) {
    string months[]={"January", "February",
        "March", "April", "May", "June", "July",
        "August", "September", "October",
        "November", "December"};
```

return (months[monthNumber]);

}

```
int numberOfDays(int monthNumber, int year){
    if (monthNumber==1 && isLeapYear(year))
        return 29;
    int monthDays[] = \{31, 28, 31, 30, 31, \}
               30, 31, 31, 30, 31, 30, 31};
    return (monthDays[monthNumber]);
void printCalendar(int year) {
    printf("Calendar - %d\n\n", year);
    int days;
    int current = dayNumber(1, 1, year);
    // i--> Iterate through all the months
    // j--> Iterate through all the days of
            the month - i
    for (int i = 0; i < 12; i++) {
        days = numberOfDays(i, year);
        cout << "
                          | " <<
            getMonthName(i).c str()
            << "|" << endl;
        printf(" Sun Mon Tue Wed Thu Fri
                 Sat\n");
        int k;
        for (k = 0; k < current; k++)
            printf(" ");
        for (int j = 1; j \le days; j++) {
            printf("%4d", j);
            if (++k > 6) {
                k = 0; cout << endl;
            }
        if (k)
            cout << endl;
        cout << "----\n";
        current = k;
} //Function call: printCalendar(year);
2.6 BINARY EXPONENTIATION: (a^b)
int binaryExp(int base,int power,int MOD =
mod) {
    int res = 1;
    while (power) {
        if (power & 1)
            res = (res * base) % MOD;
        base = ((base%MOD) * (base%MOD)) %MOD;
        power \neq 2;
    }
    return res;
2.7 BINARY EXPONENTIATION: (a^b^c)
int binaryExp(int base, int power, int
modulo) {
    int and = 1;
    while (power) {
        if (power % 2 == 1)
            ans = (ans * base) % modulo;
        base = (base * base) % modulo;
        power \neq 2;
    }
    return ans;
} //function call:
binaryExp(a, binaryExp(b, c, mod-1), mod)
```

```
2.8 Power
```

```
int x = (int) (pow(base, power) + 1e-18);
2.9 Check is prime number-O(sqrt(n))
bool prime(int n) {
    if (n<2) return false;
    if (n<=3) return true;
    if (!(n%2) || !(n%3)) return false;
    for (int i=5; i*i<=n; i+=6) {
        if (!(n%i) || !(n%(i+2)))
             return false;
    return true;
2.10 Prime factorization-O(sqrt(n))
// smallest prime factor of a number.
int factor(int n) {
    int a;
    if (n%2==0)
       return 2;
    for (a=3; a \le sqrt(n); a+=2) {
        if (n%a==0)
            return a;
    return n;
// complete factorization
int r;
while (n>1) {
    r = factor(n);
    printf("%d", r);
    n /= r;
// some facts about spf
suppose you have a number N = 120;
you represent it as N = 2^3 * 3^1 * 5^2
Now from this representation we can easily
calculate the number of divisors of number N.
Let's see how it works:
(i). we can take 2^3 in 4 different ways
      like 2^0, 2^1, 2^2, 2^3. In the same
      way we can take 3^1 in 2 ways(3^0,
      3^1) and 5^2 in 3 ways (5^0, 5^1, 5^2).
(ii). Total number of divisor is = 4 * 2 * 3
suppose, N = p_1^a \times p_2^b \times p_3^c
number of divisors = (a + 1) * (b + 1) * (c + 1)
```

```
sum of divisors
```

divisors.

$$\sigma(N) = \frac{p_1^{a+1} - 1}{p_1 - 1} * \frac{p_2^{b+1} - 1}{p_2 - 1} * \frac{p_3^{c+1} - 1}{p_3 - 1}$$

As like calculating the number of divisors,

we can also calculate the sum of all

```
2.11 Seive
```

2.12 Bitwise Seive (memory efficient)

```
const int N = 3125005;
int prime[3125005];
bool is set(int n, int pos) {
   if (n & (1 << pos))
       return true;
   return false;
int set bit(int n, int pos) {
   return (n | (1 << pos));
void sieve() {
   for (int i = 0; i < 3125005; i++)
       prime[i] = 0;
   prime[0] = set bit(0, 0);
   prime[0] = set bit(prime[0], 1);
    for (int i = 4; i \le N; i += 2) {
     prime[i/32]=set bit(prime[i/32],i%32);
    for (int i = 3; i * i <= N; i += 2) {
        if (!is set(prime[i/32],i%32)) {
            for (int j=i*i; j<=N; j+=2*i) {
               prime[j / 32] =
            set bit(prime[j / 32], j % 32);
       }
bool isPrime(int n) {
   if (!is set(prime[n / 32], n % 32))
       return true;
    return false;
void solve() {
   for (int i = 2; i < 100; i++) {
        if (isPrime(i)) {
           cout << i << " -> OK!" << endl;
   }
}
```

```
2.13 smallest prime factor(SPF) using Seive
                                                  2.16 PHI of N
const int N = 1e7 + 5;
                                                            if n = p_1^{a_1} * p_2^{a_2} * ... * p_k^{a_k} then
int spf[N];
                                                    \phi(n) = n * (1 - \frac{1}{p_1}) * (1 - \frac{1}{2}) * ... * (1 - \frac{1}{p_k})
void smallestPrimeFactorUsingSeive() {
    for (int i = 2; i < N; i++) {
                                                  int phi(int n) {
        if (spf[i] == 0) {
                                                      int result = n;
             for (int j = i; j < N; j += i) {
                                                      for (int i = 2; i * i <= n; i++) {
                 if (spf[j] == 0)
                                                          if (n % i == 0) {
                     spf[j] = i;
                                                              while (n % i == 0)
                                                                   n /= i;
        }
                                                              result -= result / i;
    }
                                                      if (n > 1)
2.14 nth prime number
                                                          result -= result / n;
// Time complexity O(log(logn))
                                                      return result;
vector<int> nth prime;
const int MX = 86200005;
bitset<MX> visited;
                                                  2.17 PHI of 1 to N
void optimized prime() {
                                                  const int N = 1e5 + 5;
    nth prime.push back(2);
                                                  vector<int> phi(N);
    for(int i=3; i<MX; i+=2){
                                                  void phi 1 to n() {
            if(visited[i])
                                                      for (int i = 0; i < N; i++)
                 continue;
                                                          phi[i] = i;
             nth prime.push back(i);
                                                      for (int i = 2; i < N; i++) {
             if(\overline{111*i*i} > MX)
                                                          if (phi[i] == i) {
                 continue;
                                                              for (int j = i; j < N; j += i)
             for (int j = i*i; j < MX; j+= i+i)
                                                                   phi[j] -= phi[j] / i;
                 visited[j] = true;
                                                      }
2.15 Modular Operation
                                                  Fact: Summation of phi of divisors of N is
Addition:
                                                  equal to N. For example N = 10.
int mod add(int a, int b, int MOD = mod) {
                                                  Divisors of 10 are 1, 2, 5, 10. Hence,
    a = a % MOD, b = b % MOD;
                                                  \phi(1) + \phi(2) + \phi(5) + \phi(10) = 1 + 1 + 4 + 4 = 10
    return (((a + b) % MOD) + MOD) % MOD;
                                                  2.18 nCr(more space, less time)
Subtraction:
                                                  int mod = 1e9 + 7;
int mod sub(int a, int b, int MOD = mod) {
                                                  const int MAX = 1e7 + 5;
    a = a % MOD, b = b % MOD;
                                                  vector<int> fact(MAX), ifact(MAX), inv(MAX);
    return (((a - b) % MOD) + MOD) % MOD;
                                                  void factorial() {
                                                      inv[1] = fact[0] = ifact[0] = 1;
Multiplication:
                                                      for (int i = 2; i < MAX; i++)
int mod mul(int a, int b, int MOD = mod) {
                                                          inv[i]=inv[mod%i]*(mod-mod/i)%mod;
    a = a % MOD, b = b % MOD;
                                                      for (int i = 1; i < MAX; i++)
    return (((a * b) % MOD) + MOD) % MOD;
                                                          fact[i] = (fact[i - 1] * i) % mod;
                                                      for (int i = 1; i < MAX; i++)
<u>Division:</u>
                                                          ifact[i]=ifact[i-1]*inv[i] % mod;
//call binary Exponential Function here.
int mminvprime(int a, int b) { return
binaryExp(a, b - 2, b);}
                                                  int nCr(int n, int r) {
//call modular multiplication here.
                                                      if (r < 0 | | r > n)
int mod div(int a, int b, int MOD = mod) {
                                                          return 0;
    a = a % MOD, b = b % MOD;
                                                      return (int)fact[n] * ifact[r] % mod *
    return (mod mul(a, mminvprime(b, MOD),
                                                  ifact[n - r] % mod;
MOD) + MOD) % MOD;
                                                  // first call factorial() function
//only for prime MOD
                                                  // then for nCr just call nCr(n,r)
```

```
2.19 nCr(less space, more time)
                                                 2.22 Binomial coefficient
const int MOD = 1e9 + 7;
                                                 #define MAXN 100 // largest n or m
const int MAX = 1e7+10;
                                                 long binomial coefficient(n,m){
vector<int> fact(MAX), inv(MAX);
                                                     int i, j;
void factorial(){
                                                     long bc[MAXN][MAXN];
    fact[0] = 1;
                                                     for (i=0; i \le n; i++) bc[i][0] = 1;
    for (int i = 1; i < MAX; i++)
                                                     for (j=0; j \le n; j++) bc[j][j] = 1;
        fact[i] = (i * fact[i - 1]) % MOD;
                                                     for (i=1; i<=n; i++)
                                                        for (j=1; j<i; j++)
//For binaryExp we call 1.6 function
                                                           bc[i][j]=bc[i-1][j-1]+bc[i-1][j];
void inverse(){
                                                     return bc[n][m];
    for (int i = 0; i < MAX; ++i)
        inv[i]=binaryExp(fact[i], MOD - 2);
                                                 2.23 10-ary to m-ary
                                                 char a[16]={'0','1','2','3','4','5','6','7',
                                                             '8','9','A','B','C','D','E','F'};
int nCr(int a, int b) {
    if (a < b \text{ or } a < 0 \text{ or } b < 0)
                                                 string tenToM(int n, int m){
        return 0;
                                                     int temp=n;
    int de = (inv[b] * inv[a - b]) % MOD;
                                                     string result="";
    return (fact[a] * de) % MOD;
                                                     while (temp!=0) {
                                                         result=a[temp%m]+result;
// nCr ends here
                                                         temp/=m;
int ModInv(int a, int M) {
                                                     }
    return binaryExp(a, M - 2, M);
                                                     return result;
                                                 2.24 m-ary to 10-ary
                                                 string num = "0123456789ABCDE";
2.20 Factorial mod
//n! mod p : Here P is mod value
                                                 int mToTen(string n, int m){
//For binaryExp we call 1.6 function
                                                     int multi=1;
int factmod (int n, int p) {
                                                     int result=0;
    int res = 1;
                                                     for (int i=n.size()-1; i>=0; i--)
    while (n > 1) {
                                                         result += num.find(n[i]) *multi;
        res=(res*binaryExp(p-1,n/p,p))%p;
                                                         multi*=m;
        for (int i=2; i<=n%p; ++i)
                                                     }
            res=(res*i) %p;
                                                     return result;
            n /= p;
                                                 2.26 Euler's totient function
    return int (res % p);
                                                 // the positive integers less than or equal
}
                                                 to n that are relatively prime to n.
                                                 int phi (int n) {
2.21 Generate combinations
                                                     int result = n;
// n>=m, choose M numbers from 1 to N.
                                                     for (int i=2; i*i<=n; ++i)
void combination(int n, int m) {
                                                         if(n %i==0){
    if (n<m) return;
                                                             while (n \%i==0)
    int a[50] = \{0\};
                                                                 n /= i;
    int k=0;
                                                             result -= result / i;
    for (int i=1; i<=m; i++) a[i]=i;
                                                         }
    while (true) {
                                                     if (n > 1)
        for (int i=1; i<=m; i++)
                                                         result -= result / n;
            cout << a[i] << " ";
                                                     return result;
        cout << endl;</pre>
                                                 2.27 EXT GCD
        k=m:
        while ((k>0) \&\& (n-a[k]==m-k)) k--;
                                                 // return \{x,y\} such that ax+by=qcd(a,b)
        if (k==0) break;
                                                 pair<int,int>ext gcd(int a, int b) {
                                                     if (b == 0)
        a[k]++;
        for (int i=k+1; i<=m; i++)
                                                         return {1, 0};
            a[i]=a[i-1]+1;
                                                         pair<int,int> tmp=ext gcd(b, a % b);
    }
                                                         return {tmp.second,
}
                                                           tmp.first - (a / b) * tmp.second};
                                                     }
                                                 }
```

2.28 Power Set

```
void printPowerSet(char* set, int setSz) {
     // Setsize of power set of a set with
     / setsize. n is (2^n-1)
    unsigned int powSetSz = pow(2, setSz);
    int i, j; // i as counter
    // Run from i 000..0 to 111..1
    for (i = 0; i < powSetSz; i++) {</pre>
        for (j = 0; j < setSz; j++) {
    //Check if jth bit in the counter is set
    //If set then print jth element from set
            if (i & (1 << j))
                cout << set[j];</pre>
        cout << endl;</pre>
    }
2.29 Number of Set Bit from 1 to N
int GLMSB(int n) {
    // GLMSB = Get Left Most Set Bit
    int pos = 0;
    while (n > 0) {
        pos++;
        n >>= 1;
    return pos;
}
int TotalSetBitsFrom1ToN(int n) {
    int id = GLMSB(n);
    int totalRep, mod, nearestPow;
    int totalSetBit = 0,addRem = 0,curr = 0;
    for (int i = 1; i <= id; ++i) {
        nearestPow = 1LL << i;</pre>
        if (nearestPow > n) {
            int lastPow = 1LL \ll (i - 1);
            mod = n % lastPow;
            totalSetBit += mod + 1;
        else {
            if (i == 1 && n % 2 == 1) {
                totalRep = (n+1) / nearestPow;
                mod = nearestPow % 2;
                addRem = 0;
            else {
                totalRep = n / nearestPow;
                mod = n % nearestPow;
                if (mod >= (nearestPow / 2))
                     addRem = mod -
                          (nearestPow/2) + 1;
                else
                     addRem = 0;
            curr = totalRep*(nearestPow / 2)
                    + addRem;
            totalSetBit += curr;
        }
    }
    return totalSetBit;
```

2.30 Legendre formula

```
// calculate the maximum power of a prime p
that divides n!
int legendre(int n, int p) {
   int ans = 0;
   while (n) {
        n /= p;
        ans += n;
   }
   return ans;
}
```

3 Algorithms

3.1 Biginteger Operation

```
struct BigInteger {
    string str;
    // Constructor to initialize
    // BigInteger with a string
    BigInteger(string s) { str = s; }
    // Overload + operator to add
    // two BigInteger objects
    BigInteger operator+(const BigInteger& b)
{
        string a = str, c = b.str;
        int alen=a.length(),clen=c.length();
        int n = max(alen, clen);
        if (alen > clen)
            c.insert(0, alen - clen, '0');
        else if (alen < clen)
            a.insert(0, clen - alen, '0');
        string res(n + 1, '0');
        int carry = 0;
        for (int^{-}i = n - 1; i >= 0; i--) {
            int digit=(a[i -'0')+(c[i]-'0')
                       +carry;
            carry = digit / 10;
            res[i + 1] = digit % 10 + '0';
        if (carry == 1) {
            res[0] = '1';
            return BigInteger(res);
        }
        else
           return BigInteger(res.substr(1));
    // Overload - operator to subtract
    // first check which number is greater
and then subtract
    BigInteger operator-(const BigInteger& b)
{
        string a = str;
        string c = b.str;
        int alen=a.length(),clen=c.length();
        int n = max(alen, clen);
        if (alen > clen)
            c.insert(0, alen - clen, '0');
        else if (alen < clen)
            a.insert(0, clen - alen, '0');
        if (a < c) {
            swap(a, c);
            swap(alen, clen);
        string res(n, '0');
        int carry = 0;
        for (int i = n - 1; i \ge 0; i--) {
            int digit = (a[i]-'0')-(c[i]-'0')
                         carry;
            if (digit < 0) {
                digit += 10;
                carry = 1;
            else {
                carry = 0;
            res[i] = digit + '0';
        // remove leading zeros
        int i = 0;
```

```
while (i < n \&\& res[i] == '0')
            i++;
        if (i == n)
            return BigInteger("0");
        return BigInteger(res.substr(i));
    }
    // Overload * operator to multiply
    // two BigInteger objects
    BigInteger operator*(const BigInteger& b)
        string a = str, c = b.str;
        int alen=a.length(),clen=c.length();
        int n = alen + clen;
        string res(n, '0');
        for (int i = alen - 1; i >= 0; i--) {
            int carry = 0;
            for(int j=clen-1; j>=0; j--) {
                 int digit = (a[i] - '0') *
            (c[j-'0')+(res[i+j+1]-'0')+carry;
                carry = digit / 10;
                res[i+j+1]=digit % 10 + '0';
            res[i] += carry;
        int i = 0;
        while (i < n \&\& res[i] == '0')
            i++;
        if (i == n)
            return BigInteger("0");
        return BigInteger(res.substr(i));
    }
    // Overload << operator to output
    // BigInteger object
    friend ostream& operator << (ostream& out,
const BigInteger& b) {
        out << b.str;
        return out;
};
3.2 Find rank k in array
int find(int 1, int r, int k) {
    int i=0, j=0, x=0, t=0;
    if (l==r) return a[l];
   x=a[(1+r)/2];
    t=a[x];
    a[x]=a[r];
    a[r]=t;
    i=1-1;
    for (int j=l; j<=r-1; j++)
        if (a[j] \le a[r]) {
            i++;
            t=a[i];
            a[i]=a[j];
            a[j]=t;
    t=a[i];
    a[i]=a[r];
    a[r]=t;
    if (i==k) return a[i];
    if (i<k) return find(i+1, r,k);</pre>
    return find(l, i-1, k);
}
```

```
3.3 InfixToPostFix
bool delim(char c) { return c == ' '; }
bool is op(char c) {
    return c == '+' || c == '-' || c == '*'
            || c == '/' || c == '^';
bool is unary(char c) {
      return c == '+' || c == '-';
int priority(char op) {
    if (op < 0) return 3;
    if (op == '+' || op == '-') return 1;
if (op == '*' || op == '/') return 2;
if (op == '^') return 4;
    return -1;
}
void process op(string& output, char op) {
    if (op < 0) {
        switch (-op) {
             case '+':
                 output += "+ ";
                 break;
             case '-':
                 output += "- ";
                 break;
         }
    else {
        switch (op) {
             case '+':
                 output += "+ ";
                 break:
             case '-':
                 output += "- ";
                 break;
             case '*':
                 output += "* ";
                 break;
             case '/':
                 output += "/ ";
                 break;
             case '^':
                 output += "^ ";
                 break;
         }
    }
}
string InfixToPostFix(string& s) {
    string output;
    stack<char> op;
    bool may_be_unary = true;
    for (int^{-}i = 0; i < (int)s.size(); i++){}
         if (delim(s[i]))
             continue;
         if (s[i] == '(') {
             op.push('(');
             may be unary = true;
        else if (s[i] == ')') {
             while (op.top() != '(') {
                process op(output, op.top());
                op.pop();
             op.pop();
             may_be_unary = false;
```

```
else if (is_op(s[i])) {
            char \overline{cur} op = s[i];
            if (may_be_unary &&
is_unary(cur_op))
                cur op = -cur op;
            while (!op.empty() &&
                     ((cur op >= 0 &&
priority(op.top()) >= priority(cur_op)) ||
                     (cur op < 0 \&\&
priority(op.top()) > priority(cur_op)))) {
                process_op(output, op.top());
                op.pop();
            op.push(cur op);
            may be unary = true;
        else {
            char number;
            while (i < (int)s.size() &&
isalnum(s[i]))
                number = s[i++];
            --i;
            output.push_back(number);
            output.push_back(' ');
            may be unary = false;
    while (!op.empty()) {
        process op(output, op.top());
        op.pop();
    return output;
}
```

```
3.4 Expression Parsing
bool delim(char c) { return c == ' '; }
bool is op(char c) { return c == '+' || c ==
'-' || c == '*' || c == '/'; }
bool is unary(char c) { return c == '+' || c
== '-';<sup>-</sup>}
int priority(char op) {
    if (op < 0) // unary operator
        return 3;
    if (op == '+' || op == '-')
        return 1;
    if (op == '*' || op == '/')
        return 2;
    return -1;
}
void process op(stack<int>& st, char op) {
    if (op < 0) {
        int l = st.top();
        st.pop();
        switch (-op) {
            case '+':
                 st.push(1);
                 break;
            case '-':
                 st.push(-1);
                 break:
        }
    }
    else {
        int r = st.top();
        st.pop();
        int l = st.top();
        st.pop();
        switch (op) {
            case '+':
                 st.push(l + r);
                 break;
            case '-':
                 st.push(l - r);
                 break;
            case '*':
                 st.push(l * r);
                 break;
            case '/':
                 st.push(1 / r);
                 break;
        }
    }
}
int evaluate(string& s) {
    stack<int> st;
    stack<char> op;
    bool may be unary = true;
    for (int i = 0; i < (int)s.size(); i++) {
        if (delim(s[i]))
            continue;
        if (s[i] == '(') {
            op.push('(');
            may be unary = true;
        else if (s[i] == ')') {
```

```
while (op.top() != '(') {
                 process op(st, op.top());
                 op.pop(\overline{)};
            op.pop();
            may be unary = false;
        else if (is op(s[i])) {
            char cur op = s[i];
            if (may_be_unary &&
is unary(cur op))
                 cur_op = -cur op;
            while (!op.empty() \&\&
                    ((cur op >= 0 &&
priority(op.top()) >= priority(cur op)) ||
                     (cur op < 0 &&
priority(op.top()) > priority(cur op)))) {
                 process op(st, op.top());
                 op.pop();
            op.push(cur op);
            may be unary = true;
        else {
            int number = 0;
            while (i < (int)s.size() &&
isalnum(s[i]))
                 number = number * 10 + s[i++]
- '0';
            --i;
            st.push(number);
            may be unary = false;
    }
    while (!op.empty()) {
        process op(st, op.top());
        op.pop();
    return st.top();
3.5 2D prefix sum
class NumMatrix {
    int row, col;
    vector<vector<int>> sums;
public:
    NumMatrix(vector<vector<int>> &matrix) {
        row = matrix.size();
        col = row>0 ? matrix[0].size() : 0;
        sums = vector<vector<int>>(row+1,
                vector<int>(col+1, 0));
        for(int i=1; i<=row; i++) {
    for(int j=1; j<=col; j++) {</pre>
                 sums[i][j] =
                 matrix[i-1][j-1] +
                 sums[i-1][j] +
                 sums[i][j-1] -
                 sums[i-1][j-1];
            }
        }
    }
```

int sumRegion(int row1, int col1,

```
int row2, int col2) {
              return sums[row2+1][col2+1] -
               sums[row2+1][col1] -
              sums[row1][col2+1] +
               sums[row1][col1];
    }
};
3.6 KMP Algorithm-O(n+m)
vector<int> createLPS(string pattern) {
    int n = pattern.length(), idx = 0;
    vector<int> lps(n);
    for (int i = 1; i < n;) {
        if (pattern[idx] == pattern[i]) {
            lps[i] = idx + 1;
            idx++, i++;
        else {
            if (idx != 0)
                idx = lps[idx - 1];
                lps[i] = idx, i++;
    }
    return lps;
int kmp(string text, string pattern) {
    int cnt of match = 0, i = 0, j = 0;
    vector<int> lps = createLPS(pattern);
    while (i < text.length()) {</pre>
        if (text[i] == pattern[j])
            i++, j++; // i->text, j->pattern
        else {
            if (j != 0)
                j = lps[j - 1];
            else
                i++;
        if (j == pattern.length()) {
            cnt of match++;
            // the index where match found ->
(i - pattern.length());
            j = lps[j - 1];
    return cnt of match;
3.7 Kadane's Algorithm O(n)
  return maximum subarray sum.
int maxSubArraySum(vector<int> &a) {
    int size = a.size();
    int maxTill = INT MIN, maxEnd = 0;
    for (int i = 0; i < size; i++) {
        maxEnd = maxEnd + a[i];
        if (maxTill < maxEnd)</pre>
            maxTill = maxEnd;
        if (maxEnd < 0)
            maxEnd = 0;
    return maxTill;
}
```

4 Data Structure

```
4.1 SEGMENT TREE
class SEGMENT TREE {
  public:
   vector<int> v;
   vector<int> seq;
   SEGMENT TREE(int n) {
        v.resize(n + 5);
        seq.resize(4 * n + 5);
   //! initially: ti = 1, low = 1, high = n
      //(number of elements in the array);
   void build(int ti, int low, int high) {
        if (low == high) {
            seg[ti] = v[low];
            return;
        int mid = (low + high) / 2;
        build(2 * ti, low, mid);
        build(2 * ti + 1, mid + 1, high);
        seg[ti] = (seg[2*ti]+seg[2*ti+1]);
    //! initially: ti = 1, low = 1, high = n
      //(number of elements in the array),
      //(ql & qr)=user input in 1 based
index;
   int find(int ti, int tl, int tr, int ql,
            int qr) {
        if (tl > qr || tr < ql) {
            return 0;
        if (tl >= ql and tr <= qr)
           return seg[ti];
        int mid = (tl + tr) / 2;
        int l = find(2*ti, tl, mid, ql, qr);
        int r = find(2*ti+1,mid+1,tr,ql,qr);
        return (l + r);
   //! initially: ti = 1, tl = 1, tr = n
      //(number of elements in the array),
      //id = user input in 1 based indexing,
      //val = updated value;
   void update(int ti, int tl, int tr, int
                id, int val) {
        if (id > tr or id < tl)
           return;
        if (id == tr and id == tl) {
            seg[til] = val;
            return;
        int mid = (tl + tr) / 2;
        update(2 * ti, tl, mid, id, val);
        update(2*ti+1,mid+1,tr,id,val);
        seg[til] = (seg[2*ti] + seg[2*ti + 1]);
   }
// use 1 based indexing for input and
//queries and update;
```

```
void update(int u, int st, int en, int l,
4.2 FENWICK TREE
                                                             int r, int x) {
// Sum
                                                     propagate(u, st, en);
struct FenwickTree {
    vector<int> bit; // binary indexed tree
                                                     if (r < st or en < 1)
                                                         return;
    FenwickTree(int n) {
                                                     else if (st \ge 1 \text{ and en } \le r) {
        this->n = n;
                                                         lz[u] += x;
        bit.assign(n, 0);
                                                         propagate(u, st, en);
    FenwickTree(vector<int>a):
                                                     else {
                   FenwickTree(a.size()) {
                                                         int mid = (st + en) >> 1;
                                                         update(2 * u, st, mid, l, r, x);
        for (size_t i=0; i < a.size(); i++)</pre>
            add(i, a[i]);
                                                         update(2*u + 1, mid+1, en, 1, r, x);
                                                         tree[u] = tree[2*u]+tree[2*u+1];
    int sum(int r) {
        int ret = 0;
        for (; r \ge 0; r = (r&(r + 1)) - 1)
                                                 int query(int u,int st,int en,int l,int r){
            ret += bit[r];
                                                     propagate(u, st, en);
                                                     if (r < st or en < 1)
        return ret;
                                                         return 0;
    int sum(int l, int r) {
                                                     else if (st >= 1 \text{ and en } <= r)
        return sum(r) - sum(1 - 1);
                                                         return tree[u];
    }
                                                     else {
                                                         int mid = (st + en) >> 1;
    void add(int idx, int delta) {
        for (; idx < n; idx = idx | (idx + 1))
                                                         int left=query(2*u, st, mid, l, r);
            bit[idx] += delta;
                                                         int right=query(2*u+1,mid+1,en,l,r);
    }
                                                         return left + right;
};
                                                     }
                                                 }
// minimum
struct FenwickTreeMin {
                                                 4.5 TRIE
    vector<int> bit;
                                                 class TrieNode {
    int n;
                                                    public:
    const int INF = (int)1e9;
                                                     int isEnd;
    FenwickTreeMin(int n) {
                                                     TrieNode *child[26];
        this->n = n;
                                                     TrieNode() {
        bit.assign(n, INF);
                                                         isEnd = 0;
                                                         for (int i = 0; i < 26; i++)
    FenwickTreeMin(vector<int> a) :
                                                             child[i] = NULL;
                   FenwickTreeMin(a.size()) {
        for (size t i=0; i < a.size(); i++)
                                                 };
            update(i, a[i]);
                                                 class Trie {
                                                     TrieNode *root;
    int getmin(int r) {
        int ret = INF;
                                                    public:
        for (; r>=0; r = (r & (r + 1)) - 1)
                                                     Trie() : root(new TrieNode()) {}
            ret = min(ret, bit[r]);
                                                     void insert(string word) {
        return ret;
                                                         TrieNode *curr = root;
                                                         for (char ch : word) {
                                                             if(curr->child[ch-'a'] == NULL)
    void update(int idx, int val) {
        for (; idx < n; idx = idx | (idx + 1))
                                                                 curr->child[ch - 'a'] =
            bit[idx] = min(bit[idx], val);
                                                                     new TrieNode();
    }
                                                             curr = curr->child[ch - 'a'];
};
                                                         }
                                                         curr->isEnd++;
4.3 SEGMENT TREE LAZY
const int N = 1e5 + 100;
                                                     bool search(string word) {
int tree[N << 2], 1z[N << 2];
                                                         TrieNode *curr = root;
void propagate(int u, int st, int en) {
                                                         for (char ch : word) {
    if (!lz[u])
                                                             if(curr->child[ch-'a'] == NULL)
        return;
    tree[u] += lz[u] * (en - st + 1);
                                                                 return false;
                                                             curr = curr->child[ch - 'a'];
    if (st != en) {
        lz[2 * u] += lz[u];
                                                         return curr->isEnd;
        lz[2 * u + 1] += lz[u];
                                                     bool startsWith(string prefix) {
    lz[u] = 0;
                                                         TrieNode *curr = root;
}
```

```
for (char ch : prefix) {
         if (curr->child[ch-'a']==NULL)
             return false;
         curr = curr->child[ch - 'a'];
     return true;
bool isJunc(TrieNode *curr) {
     for (int i = 0; i < 26; i++) {
         if (curr->child[i] != NULL)
             return true;
     return false;
 // 1 means junction delete kore asche
bool dlt(string s, int idx,
          TrieNode *curr) {
     if (idx >= s.size())
         return 0;
     if (idx == s.size() - 1) {
        if (isJunc(curr->child[s[idx] -
                                   'a'])) {
             curr->child[s[idx] -
                          'a']->isEnd = 0;
             return false;
        }
        else {
          delete curr->child[s[idx]-'a'];
          curr->child[s[idx]-'a']= NULL;
             return true;
     bool res = dlt(s, idx + 1,
              curr->child[s[idx] - 'a']);
     if (res) {
      if(isJunc(curr->child[s[idx]-'a']))
           return false;
      else if (!curr->child[s[idx] -
                'a']->isEnd) {
          delete curr->child[s[idx]-'a'];
          curr->child[s[idx]-'a']=NULL;
          return true;
     return false;
bool dlt(string s) {
     if (search(s)) {
         dlt(s, 0, root);
         return true;
     return false;
void print(string start, TrieNode *curr){
     if (curr->isEnd)
         cout << start << endl;</pre>
     for (int i = 0; i < 26; i++) {
         if (curr->child[i] != NULL) {
            start.push back(i + 'a');
            print(start, curr->child[i]);
            start.pop back();
     }
 void print() { print("", root); }
```

};

```
4.6 DSU
class DisjointSet{
   vector<int> par, sz, minElmt, maxElmt,
cntElmt;
  public:
   DisjointSet(int n) {
       par.resize(n + 1);
       sz.resize(n + 1, 1);
       minElmt.resize(n + 1);
       maxElmt.resize(n + 1);
        cntElmt.resize(n + 1, 1);
        for (int i = 1; i <= n; i++)
            par[i]=minElmt[i]=maxElmt[i]=i;
   int findUPar(int u) {
        if (u == par[u])
            return u;
        return par[u] = findUPar(par[u]);
   void unionBySize(int u, int v){
        int pU = findUPar(u);
        int pV = findUPar(v);
        if (pU == pV)
            return;
        if (sz[pU] < sz[pV])
            swap(pU, pV);
        par[pV] = pU;
        sz[pU] += sz[pV];
        cntElmt[pU] += cntElmt[pV];
       minElmt[pU] = min(minElmt[pU],
                           minElmt[pV]);
       maxElmt[pU] = max(maxElmt[pU],
                           maxElmt[pV]);
   int getMinElementIntheSet(int u) {
        return minElmt[findUPar(u)];
    int getMaxElementIntheSet(int u) {
        return maxElmt[findUPar(u)];
    int getNumofElementIntheSet(int u) {
       return cntElmt[findUPar(u)];
    }
};
4.7 Order Set
#include <ext/pb ds/assoc container.hpp>
#include <ext/pb ds/tree policy.hpp>
using namespace gnu pbds;
template <typename T> using o set = tree<T,</pre>
null type, less<T>, rb tree tag,
tree order statistics node update>;
// find by order(k) - returns an iterator to
// the k-th largest element (0 indexed);
// order of key(k)-the number of elements in
// the set that are strictly smaller than k;
```

5 Dynamic Programming

Many problems can be solved using LCS techniques.

- Longest Increasing Substring
 To solve this, we just care about when
 two char equals. Rest of the things
 should be neglected.
- Longest Palindromic Subsequence (LPS)
 To solve this, we just take a new
 string which is the reverse of the
 original string. Then just call the LCS
 function to find LPS.
- Minimum insertions to make a string palindrome
 To solve this, we just basically do string length LPS.
 Why this? Let's take an example: string s = aabca;
 Let's say aca is our LPS. Now we find how many char we need to insert to make the string palindrome while our LPS is fixed.
 - a ab c a now to make the string
 palindrome we just need to insert the
 reverse of ab after c. So the new
 string looks like a ab c ba a
- Minimum Number of Deletions and Insertions to make the string equals To solve this we just find the LCS of those string then just do: n + m - LCS.length() where n, m = strings length

// Added by HASNAT

```
5.2 MCM O(n^3)
const int N = 1005;
vector<int> v;
int dp[N][N], mark[N][N];
int MCM(int i, int j) {
    if (i == j)
        return dp[i][j] = 0;
    if (dp[i][j] != -1)
        return dp[i][j];
    int mn = INT MAX;
    for (int k = i; k < j; k++) {
        int x = mn;
        mn = min(mn, MCM(i, k) + MCM(k + 1,
j) + v[i - 1] * v[k] * v[j]);
        if (x != mn)
            mark[i][j] = k;
    }
    return dp[i][j] = mn;
void print order(int i, int j) {
    if (i == j)
        cout << "X" << i;
    else {
        cout << "(";
        print_order(i, mark[i][j]);
        print order(mark[i][j] + 1, j);
        cout << ")";
    }
// memset(dp, -1, sizeof dp);
// print order(1, n);
5.3 Length of LIS O(nlogn)
vector<int> v = \{7, 3, 5, 3, 6, 2, 9, 8\};
vector<int> seq;
here we basically check is the current
element from v is greater than the last
element of the sequence.
if it is then push it to the seq array and
if not then replace that index value.
let's take an example: v = 7 \ 3 \ 5 \ 3 \ 6 \ 2 \ 9 \ 8
1st iteration seq = 7;
2nd iteration seq = 3;
3rd iteration seq = 35;
4th iteration seq = 3 3;
5th iteration seq = 3 3 6;
6th iteration seq = 2 3 6;
7th iteration seq = 2 3 6 9;
8th iteration seq = 2 3 6 8;
*/
for (auto i : v) {
    auto id = lower bound(seq.begin(),
seq.end(), i);
    if (id == seq.end())
        seq.push back(i);
    else
        seq[id - seq.begin()] = i;
cout << seq.size() << endl;</pre>
// Edited by HASNAT
```

```
5.4 LCIS O(n * m)
                                                  5.6 SOS DP
int a[100] = \{0\}, b[100] = \{0\}, f[100] = \{0\};
                                                  // # of elements in the list for which you //
int n=0, m=0;
                                                  want to find the sum over all subsets
                                                  int n = 20;
int main(void){
    cin >> n;
                                                  // the list for which you want to find the //
    for (int i=1; i<=n; i++) cin >> a[i];
                                                  sum over all subsets
    cin >> m;
                                                  vector<int> a(1 << n);</pre>
    for (int i=1; i<=m; i++) cin >> b[i];
    for (int i=1; i<=n; i++) {
                                                  //answer for sum over subsets of each subset
        int k=0;
                                                  vector<int> sos(1 << n);</pre>
        for (int j=1; j <= m; j++) {
                                                  for (int i = 0; i < (1 << n); i++) {
            if (a[i]>b[j] \&\& f[j]>k)
                                                        // iterate over all other sets and
                   k=f[j];
            else if (a[i]==b[j] \&\& k+1>f[j])
                                                  checks whether they're a subset of i
                                                        for (int j = 0; j < (1 << n); j++) {
                   f[j]=k+1;
                                                               if ((i \& j) == j) {
        }
    }
                                                                    sos[i] += a[j];
    int and=0;
                                                               }
    for (int i=1; i<=m; i++)
                                                        }
        if (f[i]>ans) ans=f[i];
    cout << and << endl;
    return 0;
                                                  5.7 Depth and width of tree
                                                  int l[100] = \{0\}, int r[100] = \{0\};
5.5 Maximum submatrix
                                                  stack<int> mystack;
int a[150][150] = \{0\};
                                                  int n = 0, w = 0, d = 0;
int c[200] = \{0\};
                                                  int depth(int n) {
int maxarray(int n) {
                                                      if (l[n]==0 \&\& r[n]==0)
    int b=0, sum=-100000000;
                                                          return 1;
    for (int i=1; i<=n; i++) {
                                                      in depthl=depth(l[n]);
                                                      int depthr=depth(r[n]);
        if (b>0) b+=c[i];
        else b=c[i];
                                                      int dep=depthl>depthr ? depthl:depthr;
        if (b>sum) sum=b;
                                                      return dep+1;
    }
                                                  void width(int n) {
    return sum;
                                                      if (n \le d)
                                                          int t=0,x;
                                                          stack<int> tmpstack;
                                                          while (!mystack.empty()){
int maxmatrix(int n) {
                                                              x=mystack.top();
    int sum=-100000000, max=0;
    for (int i=1; i<=n; i++) {
                                                              mystack.pop();
        for (int j=1; j<=n; j++)
                                                               if (x!=0) {
            c[j]=0;
                                                                   t++;
        for (int j=i; j<=n; j++) {
                                                                   tmpstack.push(l[x]);
            for (int k=1; k \le n; k++)
                                                                   tmpstack.push(r[x]);
                                                               }
                 c[k] += a[j][k];
            max=maxarray(n);
                                                          }
            if (max>sum) sum=max;
                                                          w=w>t?w:t;
                                                          mystack=tmpstack;
                                                          width (n+1);
    return sum;
                                                      }
int main (void) {
                                                  int main(void) {
    int n=0;
                                                      cin >> n;
    cin >> n;
                                                      for (int i=1; i<=n; i++)
    for (int i=1; i<=n; i++)
                                                          cin >> l[i] >> r[i];
        for (int j=1; j<=n; j++)
                                                      d=depth(1);
            cin >> a[i][j];
                                                      mystack.push(1);
    cout << maxmatrix(n);</pre>
                                                      width (1);
                                                      cout << w << " " << d << endl;
    return 0;
                                                      return 0;
}
                                                  }
```

5.8 All possible SubArraySum in O(1)

```
bitset<100005> bs = 1;
    for (auto i : a)
    {
        bs |= (bs << i); // if previous 1
value pos is possible now ith bit or ith sm
is also possible
    }
    cout << bs.count() - 1 << endl;
    for (int i = 1; i <= 100003; i++)
        if (bs[i])
            cout << i << " ";
    cout << endl;</pre>
```

6 Graph Theory

```
6.1 SPFA - Optimal BF O(V * E)
int q[3001] = \{0\}; // queue for node
it d[1001] = {0}; // record shortest path
from start to ith node
bool f[1001] = \{0\};
int a[1001][1001]= {0}; // adjacency list
int w[1001][1001] = {0}; // adjacency matrix
int main(void) {
 int n=0, m=0;
 cin >> n >> m;
 for (int i=1; i<=m; i++) {
 int x=0, y=0, z=0;
cin >> x >> y >> z;
// node x to node y has weight z
a[x][0]++;
a[x][a[x][0]]=y;
 w[x][y]=z;
// for undirected graph
a[x][0]++;
a[y][a[y][0]]=x;
w[y][x]=z;
int s=0, e=0;
cin >> s >> e; // s: start, e: end
SPFA(s);
cout << d[e] << endl;</pre>
return 0;
}
void SPFA(int v0){
    int t,h,u,v;
    for (int i=0; i<1001; i++) d[i]=INT MAX;
    for (int i=0; i<1001; i++) f[i]=false;
    d[v0]=0;
   h=0;
   t=1;
    q[1]=v0;
   f[v0]=true;
    while (h!=t) {
        h++;
        if (h>3000) h=1;
        u=q[h];
        for (int j=1; j \le a[u][0]; j++) {
            v=a[u][j];
            if (d[u]+w[u][v]< d[v]) // change
to > if calculating longest path
            {
                d[v]=d[u]+w[u][v];
                if (!f[v]) {
                    t++;
                     if (t>3000) t=1;
                     q[t]=v;
                     f[v]=true;
                 }
            }
        f[u]=false;
    }
}
```

6.2 Dijkstra O(V + ElogV)

```
typedef pair<int, int> pairi;
int N = 20000 + 5;
vector<vector<pairi>> adj(N);
vector<int> dis(N, inf), parent(N);
void dijkstra(int src) {
    priority queue<pairi, vector<pairi>,
                  greater<pairi>> pq;
    dis[src] = 0;
    pq.push({0, src});
    while (pq.size()) {
        auto top = pq.top();
        pq.pop();
        for (auto i : adj[top.second]) {
            int v = i.first;
            int wt = i.second;
            if (dis[v]>dis[top.second]+wt) {
                dis[v]=dis[top.second]+wt;
                pq.push({dis[v], v});
        }
    }
}
```

6.3 BellmanFord O(V.E)

```
vector<int> dist;
vector<int> parent;
vector<vector<pair<int, int>>> adj;
// resize the vectors from main function
void bellmanFord(int num of nd, int src) {
    dist[src] = 0;
    for (int step=0;step<num of nd;step) {</pre>
        for (int i = 1; i<=num of nd; i++) {
            for (auto it : adj[i]) {
                int u = i;
                int v = it.first;
                int wt = it.second;
                if (dist[u] != inf &&
                ((dist[u] + wt) < dist[v])) {
                     if(step==num of nd - 1){
                         cout << "Negative
                              cycle found\n";
                         return;
                     dist[v] = dist[u] + wt;
                    parent[v] = u;
                }
            }
    for (int i = 1; i \le num of nd; i++)
        cout << dist[i] << "";
    cout << endl;
}
```

```
6.4 Floyd-Warshall algorithm O(n^3)
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool FloydWarshall (VVT &w, VVI &prev) {
 int n = w.size();
 prev = VVI (n, VI(n, -1));
  for (int k = 0; k < n; k++) {
   for (int i = 0; i < n; i++) {
      for (int j = 0; j < n; j++) {
        if (w[i][j] > w[i][k] + w[k][j]){
          w[i][j] = w[i][k] + w[k][j];
          prev[i][j] = k;
      }
    }
  }
// check for negative weight cycles
  for(int i=0;i<n;i++)</pre>
    if (w[i][i] < 0) return false;</pre>
  return true;
6.5 Topological sort
map<string, vector<string>> adj;
map<string, int> degree;
set<string> nodes;
vector<string> ans;
// adj: graph input, degree: cnt indegree,
// node: unique nodes, ans: path
int c = 0;
void topo sort() {
   queue<string> qu;
// traverse all the nodes and check if its
degree is 0 or not..
    for (string i : nodes) {
        if (degree[i] == 0) {
            qu.push(i);
    }
    while (!qu.empty()) {
        string top = qu.front();
        qu.pop();
        ans.push back(top);
        for (string i : adj[top]) {
            degree[i]--;
            if (degree[i] == 0) {
```

qu.push(i);

}

}

}

6.6 Kruskal O(ElogE)

```
typedef pair<int, int> edge;
class Graph {
    vector<pair<int, edge>> G, T;
    vector<int> parent;
    int cost = 0;
   public:
    Graph(int n) {
        for (int i = 0; i < n; i++)
            parent.push back(i);
    void add edges(int u,int v,int wt) {
        G.push back(\{wt, \{u, v\}\}\);
    int find set(int n) {
        if (n == parent[n])
            return n;
        else
            return find set(parent[n]);
    }
    void union set(int u, int v) {
        parent[u] = parent[v];
    void kruskal() {
        sort(G.begin(), G.end());
        for (auto it : G) {
         int uRep=find set(it.second.first);
        int vRep=find set(it.second.second);
            if (uRep != vRep) {
                cost += it.first;
                T.push back(it);
                union set(uRep, vRep);
        }
    }
    int get cost() { return cost; }
    void print() {
        for (auto it : T)
            cout << it.second.first << " "</pre>
                  << it.second.second << "->"
                  << it.first << endl;
    }
// g.add edges(u, v, wt);
// g.kruskal();
```

```
6.7 Prim - MST O(ElogV)
typedef pair<int, int> pii;
class Prims {
   map<int, vector<pii>>> graph;
   map<int, int> visited;
  public:
   void addEdge(int u, int v, int w) {
        graph[u].push back({v, w});
        graph[v].push back({u, w});
   vector<int> path(pii start) {
      vector<int> ans;
      priority queue<pii, vector<pii>,
                        greater<pii>>> pq;
                        // cost vs node
      pq.push((start.second, start.first));
      while (!pq.empty()) {
            pair<int, int> curr = pq.top();
            pq.pop();
            if (visited[curr.second])
                continue;
            visited[curr.second] = 1;
            ans.push back(curr.second);
            for (auto i:graph[curr.second]) {
               if (visited[i.first])
                   continue;
               pq.push({i.second, i.first});
            }
      return ans;
6.8 Eulerian circuit O(V+E)
unordered map<int, int> Start, End, Val;
unordered map<int, pair<int, int>> Range;
int start = 0;
void dfs(int node) {
   visited[node] = true;
   Start[node] = start++;
   for (auto child : adj[node]) {
        if (!visited[child])
            dfs(child);
   End[node] = start - 1;
dfs(1);
vector<int> FlatArray(start + 5);
for (auto i : Start) {
   FlatArray[i.second] = Val[i.first];
```

{i.second, End[i.first]};

Range[i.first]=

}

6.9 LCA struct LCA { vector<int> height, euler; vector<int> first, segtree; vector<bool> visited; int n; LCA(vector<vector<int>> &adj,int root=0) { n = adj.size();height.resize(n), first.resize(n); euler.reserve(n * 2); visited.assign(n, false); dfs(adj, root); int m = euler.size(); segtree.resize(m * 4); build(1, 0, m - 1); } void dfs(vector<vector<int>> &adj, int node, int h = 0) { visited[node] = true; height[node] = h;first[node] = euler.size(); euler.push back(node); for (auto to : adj[node]) { if (!visited[to]) { dfs(adj, to, h + 1);euler.push_back(node); void build(int node, int b, int e) { if (b == e)segtree[node] = euler[b]; else { int mid = (b + e) / 2;build(node << 1, b, mid);</pre> build(node << 1 | 1, mid+1, e);</pre> int l = segtree[node << 1];</pre> int r = segtree[node << 1 | 1];seqtree[node] = (height[l] < height[r]) ? l : r;int query(int node,int b,int e,int L,int R){ if $(b > R \mid\mid e < L)$ return -1; if (b >= L && e <= R)return segtree[node]; int mid = $(b + e) \gg 1$; int lf = query(node << 1, b, mid, L, R);</pre> int rg = query(node << 1 | 1, mid + 1, e, L, R);if (lf == -1) return rg; if (rg == -1) return lf; return height[lf]<height[rg]?lf: rg;</pre> int lca(int u, int v) { int left=first[u],right = first[v]; if (left > right) swap(left, right); return query(1, 0, euler.size() - 1, left, right); }

};

```
6.10 Min cost max flow
struct Edge{
    int from, to, capacity, cost;
vector<vector<int>> adj, cost, capacity;
const int INF = 1e9;
void shortest paths (int n, int v0,
vector<int>& d, vector<int>& p) {
    d.assign(n, INF);
    d[v0] = 0;
   vector<bool> inq(n, false);
    queue<int> q;
    q.push(v0);
    p.assign(n, -1);
   while (!q.empty()) {
        int u = q.front();
        q.pop();
        inq[u] = false;
        for (int v : adj[u]) {
            if (capacity[u][v] > 0 \&\& d[v] >
d[u] + cost[u][v]) {
                d[v] = d[u] + cost[u][v];
                p[v] = u;
                if (!inq[v]) {
                    inq[v] = true;
                    q.push(v);
            }
        }
int min cost flow(int N, vector<Edge> edges,
int K, int s, int t) {
    adj.assign(N, vector<int>());
    cost.assign(N, vector<int>(N, 0));
    capacity.assign(N, vector<int>(N, 0));
    for (Edge e : edges) {
        adj[e.from].push back(e.to);
        adj[e.to].push back(e.from);
        cost[e.from][e.to] = e.cost;
        cost[e.to][e.from] = -e.cost;
        capacity[e.from][e.to] = e.capacity;
    }
    int flow = 0;
    int cost = 0;
    vector<int> d, p;
    while (flow < K) {
        shortest paths (N, s, d, p);
        if (d[t] == INF)
            break;
        // find max flow on that path
        int f = K - flow;
        int cur = t;
        while (cur != s) {
            f = min(f,
capacity[p[cur]][cur]);
```

cur = p[cur];

// apply flow
flow += f;

cost += f * d[t];

```
cur = t;
        while (cur != s) {
            capacity[p[cur]][cur] -= f;
            capacity[cur][p[cur]] += f;
            cur = p[cur];
        }
    }
    if (flow < K)
        return -1;
    else
        return cost;
6.11 SCC
unordered map<int, vector<int>> adj, InvAdj;
stack<int> order;
unordered map<int, bool> visited;
unordered map<int, vector<int>> all scc;
unordered map<int, int> compId;
void dfs for start(int curr) {
    visited[curr] = 1;
    for (auto i : adj[curr])
        if (!visited[i])
            dfs for start(i);
    order.push(curr);
vector<int> curr comp;
void dfs for scc(int curr) {
    visited[curr] = 1;
    for (auto i : InvAdj[curr])
        if (!visited[i])
            dfs for scc(i);
    curr comp.push back(curr);
inline void scc() {
    int n, e, u, v;
    cin >> n >> e;
    for (int i = 0; i < e; i++) {
        cin >> u >> v;
        adj[u].push back(v);
        InvAdj[v].push back(u);
    for (int i = 1; i \le n; i++)
        if (!visited[i])
            dfs for start(i);
    visited.clear();
    while (!order.empty()){
        if (!visited[order.top()]){
            curr comp.clear();
            dfs for scc(order.top());
            int sz = all scc.size() + 1;
            all scc[sz] = curr comp;
            for (auto i : curr comp)
                compId[i] = sz;
        order.pop();
    }
no. of ways and min cost of connecting the
sccs
```

```
const int MOD = 1e9 + 7, N = 1e5 + 2, INF =
1e18 + 2;
int n, m, comp[N];
vector<int> adj[N], rev[N];
bitset<N> vis;
void DFS1(int u, stack<int> &TS) {
    vis[u] = true;
    for (int v : adj[u])
        if (!vis[v])
            DFS1(v, TS);
    TS.push(u);
void DFS2(int u, const int scc_no, int
&min cost, int &ways, vector<int> &cost) {
    vis[u] = true;
    comp[u] = scc no;
    for (int v : rev[u])
        if (!vis[v]) {
            if (\min cost == cost[v])
                ++ways;
            else if (min cost > cost[v]) {
                ways = 1;
                min cost = cost[v];
            DFS2(v, scc no, min cost, ways,
                 cost);
        }
signed main() {
    FIO cin >> n;
    vector<int> cost(n + 1);
    for (int i = 1; i \le n; ++i)
        cin >> cost[i];
    cin >> m;
    while (m--) {
        int u, v;
        cin >> u >> v;
        adj[u].push back(v);
        rev[v].push back(u);
    int tot = 0, ways = 1;
    stack<int> TS;
    for (int i = 1; i \le n; ++i)
        if (!vis[i])
            DFS1(i, TS);
    vis.reset();
    int scc no = 0;
    while (!TS.empty()) {
        int u = TS.top();
        TS.pop();
        if (!vis[u]) {
            int tmp cst = cost[u], tmp ways =
1;
            DFS2(u, ++scc no, tmp cst,
                 tmp ways, cost);
            tot += tmp cst;
            ways = (ways * tmp ways) % MOD;
    }
    cout << tot << ' ' << ways;
}//TC: O(V+E)
```

```
6.12 Biparitite
const int N=1000;
int adj[N][N];
int n,e;
bool isBicolored(int s) {
    int colorArray[n];
    for(int i=0;i<n;i++)</pre>
        colorArray[i]=-1; //init no color;
    queue<int>q;
    q.push(s);
    colorArray[s]=1; //assigning first color
    while(!q.empty()){
        int senior = q.front();
        q.pop();
        if (adj[senior] [senior] == 1)
            return false;
        for(int i=0;i<n;i++) {</pre>
            int junior=i;
            if(adj[senior][junior]==1){
if (colorArray[junior] == colorArray[senior])
//successor(child/junior) having same color
                     return false;
                 ///if(colorArray[junior]!=-1)
continue;
             ///not same color but have a
color
                else
if(colorArray[junior]==-1){
                                     ///No
color assigned
                     q.push(junior);
colorArray[junior]=!colorArray[senior];
///assigning diff color
}}}  return true;}
6.13 Two farthest node
vector<int>adj[30001];
map<pair<int,int>,int>weight;
map<int,int>vis,dis;
void dfs(int node)
    vis[node]=1;
    for(int i=0;i<adj[node].size();i++){</pre>
        int child=adj[node][i];
        if(vis[child]==1) continue;
dis[child] += dis[node] + weight[{node, child}];
        dfs(child);
    }
```

void reset()

}

for (int i=0; i<30001; i++) {

dis.clear(), weight.clear(), vis.clear();

adj[i].clear();

```
int main()
    int t; cin>>t;
    for (int p=1;p<=t;p++)
        int n,u,v,w; cin>>n;
        for (int i=0; i< n-1; i++) {
            cin>>u>>v>>w;
            adj[u].push back(v);
            adj[v].push back(u);
            weight[\{u,v\}]=w;
            weight [\{v,u\}] = w;
        dfs(0);
        int max dis=0, farthestVertex;
        map<int,int>::iterator i;
        for(i=did.begin();i!=did.end();i++){
            if(i->second>max dis){
                max dis=i->second;
                 farthestVertex=i->first;
            }
        }
        vis.clear();
        dis.clear();
        dfs(farthestVertex);
        max dis=0;
        for(i=did.begin();i!=did.end();i++){
            if(i->second>max dis){
                 max dis=i->second;
        court<<"Case "<<p<<":
"<<max dis<<"\n";
        reset();
7 Random Staff
7.4 Knight Moves
int X[8] = \{2, 1, -1, -2, -2, -1, 1, 2\};
int Y[8] = \{1, 2, 2, 1, -1, -2, -2, -1\};
7.5 bit count in O(1)
int BitCount(unsigned int u) {
     unsigned int uCount;
     uCount = u - ((u >> 1) & 033333333333 -
((u \gg 2) \& 011111111111);
     return ((uCount + (uCount >> 3)) &
030707070707) % 63;
}
7.6 Matrix Exponentiation
// A technique of computing a number raised
to a square matrix in a fast and efficient
manner.
// Uses properties of exponentiation and
binary numbers for fast computation.
```

```
// Running time:
                                                     }
// O(m<sup>3</sup>*log(n)) where m is the size of the
matrix and n is the power the matrix is being
raised to.
                                                 7.8 sqrt decomposition (MO's Algo)
//
                                                 // https://www.spoj.com/problems/DQUERY/
// INPUT:
                                                 #include <bits/stdc++.h>
// - size of matrix m
                                                 using namespace std;
// - the matrix A
                                                 const int SIZE 1 = 1e6 + 10, SIZE 2 = 3e4 +
// - the power n
// - modulo value mod
                                                 class query{
//
                                                 public:
// OUTPUT:
                                                     int l, r, indx;
// - the matrix A^n (all values mod m)
                                                 };
                                                 int block size, cnt = 0;
#include<bits/stdc++.h>
                                                 int frequency[SIZE 1], a[SIZE 2];
using namespace std;
                                                 void add(int indx) {
typedef long long LL;
                                                     ++frequency[a[indx]];
                                                     if (frequency[a[indx]] == 1)
LL arr[60][60], res[60][60], tmp[60][60], m;
                                                         ++cnt;
void matMul (LL a[][60], LL b[][60], LL mod)
                                                 void sub(int indx){
                                                     --frequency[a[indx]];
    for(int i=0; i<m; i++)
                                                     if (frequency[a[indx]] == 0)
        for (int j=0; j < m; j++)
                                                         --cnt;
            tmp[i][j] = 0;
                                                 bool comp(query a, query b) {
            for (int k=0; k < m; k++)
                                                     if (a.l / block size == b.l / block size)
                                                        return a.r < b.r;
                tmp[i][j] +=
                                                     return a.l / block size < b.l /</pre>
(a[i][k]*b[k][j])%mod;
                                                 block size;
                tmp[i][j] %= mod;
            }
                                                 signed main(){
        }
                                                     int n; cin >> n;
                                                     for (int i = 0; i < n; ++i) cin>>a[i];
void power(LL n, LL mod)
                                                    int q; cin >> q;
                                                    int ans[q] = {};
    for(int i=0; i<m; i++)
                                                     query Qur[q];
        for(int j=0; j<m; j++)
                                                     for (int i = 0; i < q; ++i) {
            if(i==j) res[i][j] = 1;
                                                         int l, r; cin>>l>>r;
            else res[i][j] = 0;
                                                         Qur[i].1 = 1 - 1;
                                                         Qur[i].r = r - 1;
    while(n)
                                                         Qur[i].indx = i;
        if(n&1)
                                                     block size = sqrt(n); // sqrt(q) dileo
                                                 hobe, but n is more accurate
            matMul(res,arr,mod);
            for(int i=0; i<m; i++)
                                                    sort(Qur, Qur + q, comp);
                for(int j=0; j<m; j++)</pre>
                                                     int ML = 0, MR = -1;
res[i][j] = tmp[i][j];
                                                     for (int i = 0; i < q; ++i) {
           n--;
                                                         int L = Qur[i].1;
        }
                                                         int R = Qur[i].r;
        else
            matMul(arr,arr,mod);
                                                         // fixing right pointer
            for(int i=0; i<m; i++)
                                                         while (MR < R) add(++MR);
                for(int j=0; j<m; j++)
                                                         while (MR > R) sub(MR--);
                                                         // fixing left pointer
arr[i][j] = tmp[i][j];
            n/=2;
                                                         while (ML < L) sub(ML++);
                                                         while (ML > L) add(--ML);
```

```
for (auto it : set1) {
        ans[Qur[i].indx] = cnt;
                                                         int left = t - it;
    for (int i = 0; i < q; ++i)
                                                         if (left < 0) continue;
        cout << and[i] << '\n';
}//sqrt(n)
                                                         int indx = les equal(set2, left);
                                                         int temp sum set2 = (indx != -1 ? (it
                                                 + set2[indx]) : 0);
7.9 Meet in the middle
                                                         if (temp sum set2 \le t)
#include <bits/stdc++.h>
                                                             ans = max(ans, temp sum set2);
using namespace std;
int les equal(vector<int> &s, int key) {
                                                     cout << ans;
                                                 }//TC: O(2^(LK+1))
    int size = s.size();
    int lo = 0, hi = size - 1, ans = 0;
                                                 7.10 PIE (inclusion - exclusion)
                                                 #include <bits/stdc++.h>
    while (hi \geq lo) {
                                                 using namespace std;
        int mid = lo + (hi - lo) / 2;
        if (s[mid] \le key) \{
                                                 inline int LCM(int a, int b) {
            ans = max(ans, mid);
                                                     return a * b / gcd(a, b);
            lo = mid + 1;
        else hi = mid - 1;
                                                 int PIE(int div[], int n, int num){
                                                     int sum = 0;
    return ans;
                                                     for (int msk=1; msk < (1 << n); ++msk) {
signed main(){
                                                         int bit cnt = 0;
                                                         int cur lcm = 1;
    FIO int n, n1, n2, t;
    cin >> n >> t;
                                                         for (int i = 0; i < n; ++i) {
    n1 = (n + 1) / 2;
                                                             if (msk & (1 << i)) {
    n2 = n / 2;
                                                                 ++bit cnt;
                                                                 cur lcm = LCM(cur lcm,
    int al[n1]; for(int &i: al) cin>>i;
                                                 div[i]);
    int a2[n2]; for(int &i: a2) cin>>i;
                                                             }
    vector<int> set1, set2;
    for(int mask=0; mask < (1<<n1); ++mask){</pre>
                                                         int cur = num / cur lcm;
        int temp sum = 0;
                                                         if (bit cnt & 1) sum += cur;
        for (int i = 0; i < n1; ++i) {
                                                         else sum -= cur;
            int f = 1 << i;
            if (f & mask)
                                                     return num - sum;
                temp_sum += a1[i];
        }
                                                 signed main() {
        set1.push_back(temp_sum);
                                                     int n, m;
                                                     while (cin >> n >> m) {
    for(int mask=0; mask < (1 << n2); ++mask){
        int temp sum = 0;
                                                         int a[m];
        for (int^{-}i = 0; i < n2; ++i){
                                                         for(int &i : a)cin >> i;
            int f = 1 << i;
                                                         cout << PIE(a, m, n) << '\n';
            if (f & mask)
                temp sum += a2[i];
                                                     }
        set2.push back(temp sum);
                                                 7.12 Binary Search
    sort(set2.begin(), set2.end());
                                                 11 lo=0, hi=mx; ///mx=max possible ans
                                                 while(lo<hi){
    // for(auto itr: set2) cout<<itr<<' ';</pre>
                                                     ll mid=(lo+hi+1)>>1;
    // cout<<'\n';
                                                     if(condition) ///valid condition->and
    // for(auto itr: set1) cout<<itr<<' ';</pre>
                                                 can be greater than or equal mid
    // cout<<'\n';
                                                         lo=mid;
```

else

int and = 0;

```
hi=mid-1; ///ans is less than mid
///or
while(lo<hi){</pre>
    11 mid=(lo+hi)>>1;
    if(condition) ///valid condition->and
can be less than or equal mid
        hi=mid;
    else
        lo=mid+1; ///ans is greater than mid
}
ll lo=0, hi=mx, esp=maxError;
while((hi-lo)>esp){
    11 \text{ mid}=(10+\text{hi}+\text{esp})/2.0;
    if (condition) lo=mid;
    else
              hi=mid-esp;
 while((hi-lo)>esp){
    11 \text{ mid}=(10+\text{hi})/2.0;
    if(condition) hi=mid;
    else
                lo=mid+esp;
7.13 Generating Permutations
```

```
int length, perm left to print;
bool placed[10000];
vector<char>perm;
void generate permutations(int curr length) {
    if (perm left to print==0) return;
    if(curr length==length) {
        for(int i=0;i<length;i++) {</pre>
            cout<<perm[i];
        cout<<"\n";
        perm left to print--;
    for(char ch='A';ch<('A'+length);ch++) {</pre>
        if(!placed[ch-'A']){
             perm.push_back(ch);
            placed[ch-'A']=true;
generate permutations(curr length+1);
            perm.pop back();
            placed[ch-'A']=false;
    }
}
int main() {
    ioi;
    int t; cin>>t;
    for(int tc=1;tc<=t;tc++) {</pre>
        cin>>length>>perm left to print;
        court<<"Case "<<tc<<":\n";
        generate permutations(0);
    }
}
```

```
7.14 N Queen optimal
```

```
// It just counts the number of ways to place
the order.
const int N = 32;
int mark[N][N];
char grid[N][N];
int n, cnt;
void fillup(int row, int col) {
   for (int i = 1; i < n - row + 1; i++) {
       mark[row + i][col]++;
        if (col - i >= 0)
            mark[row + i][col - i]++;
        if (col + i < n)
            mark[row + i][col + i]++;
void fillout(int row, int col) {
   for (int i = 1; i < n - row + 1; i++) {
       mark[row + i][col]--;
        if (col - i >= 0)
            mark[row + i][col - i]--;
        if (col + i < n)
            mark[row + i][col + i]--;
void find way(int row) {
   if (row == n) {
        cnt++;
        return;
   for (int j = 0; j < n; j++) {
        if (grid[row][j] == '*' or
mark[row][j])
            continue;
        fillup(row, j);
        find way(row + 1);
        fillout(row, j);
   }
// input in grid. call find way(0);
```