

Harvard University  
Computer Science 121

Problem Set 8

Due Tuesday, November 23, 2010 at 1:20 PM.

Submit a single PDF (lastname+ps8.pdf) of your solutions to cs121+ps8@seas.harvard.edu

**LATE PROBLEM SETS WILL NOT BE ACCEPTED.**

See syllabus for collaboration policy.

**Name**

Problem set by !!! Your Name Here !!!

with collaborator !!! Collaborators' names here !!!!

PROBLEM 1 (3+3+3+3 points)

For each of the following statements, state whether it is true or false, and prove your assertion. For parts A-C, if the statement is true, you should state what the actual parameters  $n_0$  and  $c$  are, and why.

- (A)  $4n^3 + 2n^2 = O(n^3 + n)$ .
- (B)  $\log_2 n = \Theta(\log_{10} n^2)$ .
- (C)  $2^{n/2} = \Omega(2^n)$ .
- (D)  $n^{100001} = o(1.00001^n)$ . (*Hint:* you may want to recall l'Hôpital's rule.)

PROBLEM 2 (10 points)

Throughout the class we have focused on computability properties of *languages*. But often we want to understand computational problems that are not YES/NO problems, for example computing a function. In this problem we will see how such problems can be translated into “equivalent” languages (so focusing on languages is not a big restriction).

Given any function  $f : \Sigma^* \rightarrow \Delta^*$ , show how to define a language  $L_f$  such that (a) any algorithm to compute  $f$  can be computably transformed into an algorithm that decides  $L_f$ , and (b) conversely, any algorithm that decides  $L_f$  can be transformed into an algorithm that computes  $f$ .

PROBLEM 3 (8+12 points)

Prove that the class  $P$  is closed under

- (A) Concatenation.
- (B) Kleene star. (*Hint:* Use dynamic programming. Look at the algorithm we gave in class for recognizing context-free languages via Chomsky Normal Form.)

PROBLEM 4 (10 points)

Show that there is a polynomial time algorithm which, given an NFA  $N$  and string  $w$ , determines whether  $N$  accepts  $w$ . Assume a multitape TM model of computation and analyze the degree of the polynomial as a function of both  $|\langle N \rangle|$  and  $|w|$ .

Note that converting  $N$  to a DFA won't do the trick, because that step alone would be exponential in  $|\langle N \rangle|$ .

PROBLEM 5 (10 points)

Prove that  $\text{ALL}_{\text{DFA}} = \{\langle D \rangle : D \text{ is a DFA and } L(D) = \Sigma^*\}$  is in P.

PROBLEM 6 (Challenge!! 3 points)

It is known (though not trivial) that testing whether a binary number represents a prime number is in P. However, it is currently unknown whether or not a number (given in binary) can be *factored* in polynomial time. Explicitly construct a TM  $M$  that factors numbers such that  $M$  runs in polynomial time iff there exists a TM that factors numbers in polynomial time. We want you to give us a correspondence between existence and construction.  $M$  should factor a number no matter what, but do it in polynomial time if in fact there exists *some* TM that factors numbers in polynomial time.  $M$  should take a binary number  $n$  as input and then halt with the (unique) prime factorization of  $n$  written on its tape.