



Introduction to Arithmetic Operation

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Unsigned Addition/Subtraction

~~| | |
|-------|-------|
| 8 | 1000 |
+9	+1001
1	0001

8	01000
+9	+01001
-----	-----
17	10001

14	1110
- 5	- 0101
-----	-----
9	1001

~~| | | |
|-------|--------|-------|
| 5 | 0101 | 0101 |
- 14	- 1110	+0010
7		0111



Signed Addition/Subtraction

~~| | |
|-------|--------|
| 8 | 01000 |
+9	+01001
-15	10001

8	001000
+9	+001001
-----	-----
17	010001

~~| | | |
|-------|--------|--------|
| (-5) | 11011 | 11011 |
- 14	-01110	+10010
13		01101

(-5)	111011	111011
- 14	- 001110	+110010
-----	-----	-----
-19		101101



Signed/Unsigned Addition

- If N -bit addition/subtraction is required, usually we use
 - **Unsigned addition/subtraction**
 - **wire** [N-1:0] In1,In2;
 - **wire** [N:0] AddOut;
 - **assign** AddOut={1'b0,In1}+{1'b0,In2};
 - **Signed addition/subtraction**
 - **wire** [N-1:0] In1,In2;
 - **wire** [N:0] AddOut;
 - **assign** AddOut={In1[N-1],In1}+{In2[N-1],In2};

Unsigned Multiplication

❑ Multiplicand: $Y = (y_{M-1}, y_{M-2}, \dots, y_1, y_0)$

❑ Multiplier: $X = (x_{N-1}, x_{N-2}, \dots, x_1, x_0)$

❑ Product: $P = \left(\sum_{j=0}^{M-1} y_j 2^j \right) \left(\sum_{i=0}^{N-1} x_i 2^i \right) = \sum_{i=0}^{N-1} \sum_{j=0}^{M-1} x_i y_j 2^{i+j}$

						y_5	y_4	y_3	y_2	y_1	y_0	
						x_5	x_4	x_3	x_2	x_1	x_0	
						$x_0 y_5$	$x_0 y_4$	$x_0 y_3$	$x_0 y_2$	$x_0 y_1$	$x_0 y_0$	
					$x_1 y_5$	$x_1 y_4$	$x_1 y_3$	$x_1 y_2$	$x_1 y_1$	$x_1 y_0$		
				$x_2 y_5$	$x_2 y_4$	$x_2 y_3$	$x_2 y_2$	$x_2 y_1$	$x_2 y_0$			
			$x_3 y_5$	$x_3 y_4$	$x_3 y_3$	$x_3 y_2$	$x_3 y_1$	$x_3 y_0$				
		$x_4 y_5$	$x_4 y_4$	$x_4 y_3$	$x_4 y_2$	$x_4 y_1$	$x_4 y_0$					
	$x_5 y_5$	$x_5 y_4$	$x_5 y_3$	$x_5 y_2$	$x_5 y_1$	$x_5 y_0$						
	p_{11}	p_{10}	p_9	p_8	p_7	p_6	p_5	p_4	p_3	p_2	p_1	p_0

multiplicand
multiplier

partial
products

product



Two's Complement Multiplication

- A binary two's-complement number is formulated as

$$\begin{aligned} X &= -X_{N-1}2^{N-1} + X_{N-2}2^{N-2} + X_{N-3}2^{N-3} \cdots + X_02^0 \\ &= -X_{N-1}2^{N-1} + \sum_{i=0}^{N-2} X_i 2^i \end{aligned}$$

- Similarly for Y and the product of X and Y is

$$\begin{aligned} XY &= (-X_{N-1}2^{N-1} + \sum_{i=0}^{N-2} X_i 2^i)(-Y_{N-1}2^{N-1} + \sum_{j=0}^{N-2} Y_j 2^j) \\ &= X_{N-1}Y_{N-1}2^{2N-2} + \sum_{i=0}^{N-2} \sum_{j=0}^{N-2} X_i Y_j 2^{i+j} - Y_{N-1} \sum_{i=0}^{N-2} X_i 2^{N+i-1} - X_{N-1} \sum_{j=0}^{N-2} Y_j 2^{N+j-1} \end{aligned}$$

- The last two terms can both be expressed as

$$\begin{aligned} -\sum_{i=0}^{N-2} Y_{N-1} X_i 2^{N+i-1} &= -2^{2N-2} + (\sum_{i=0}^{N-2} (1 - Y_{N-1} X_i) 2^{N+i-1}) + 2^{N-1} \\ &= -2^{2N-2} + (\sum_{i=0}^{N-2} \overline{Y_{N-1} X_i} 2^{N+i-1}) + 2^{N-1} \end{aligned}$$

$$x_{N-1} y_{M-1} 2^{M+N-2}$$

$x_5 y_5$

$$p_{11} \quad p_{10} \quad p_9 \quad p_8 \quad p_7 \quad p_6 \quad p_5 \quad p_4 \quad p_3 \quad p_2 \quad p_1 \quad p_0$$

$$p_{11} \quad p_{10} \quad p_9 \quad p_8 \quad p_7 \quad p_6 \quad p_5 \quad p_4 \quad p_3 \quad p_2 \quad p_1 \quad p_0$$

Two's Complement Array Multiplication (2/2)

- Modified Baugh-Wooly multiplier

						y_5	y_4	y_3	y_2	y_1	y_0
						x_5	x_4	x_3	x_2	x_1	x_0
					1	$\overline{x_5 y_0}$	$x_0 y_4$	$x_0 y_3$	$x_0 y_2$	$x_0 y_1$	$x_0 y_0$
				$\overline{x_5 y_1}$		$x_1 y_4$	$x_1 y_3$	$x_1 y_2$	$x_1 y_1$	$x_1 y_0$	
			$\overline{x_5 y_2}$	$x_2 y_4$		$x_2 y_3$	$x_2 y_2$	$x_2 y_1$	$x_2 y_0$		
		$\overline{x_5 y_3}$	$x_3 y_4$	$x_3 y_3$		$x_3 y_2$	$x_3 y_1$	$x_3 y_0$			
	$\overline{x_5 y_4}$	$x_4 y_4$	$x_4 y_3$	$x_4 y_2$		$x_4 y_1$	$x_4 y_0$				
1	$x_5 y_5$	$\overline{x_4 y_5}$	$\overline{x_3 y_5}$	$\overline{x_2 y_5}$	$\overline{x_1 y_5}$	$\overline{x_0 y_5}$					
p_{11}	p_{10}	p_9	p_8	p_7	p_6	p_5	p_4	p_3	p_2	p_1	p_0



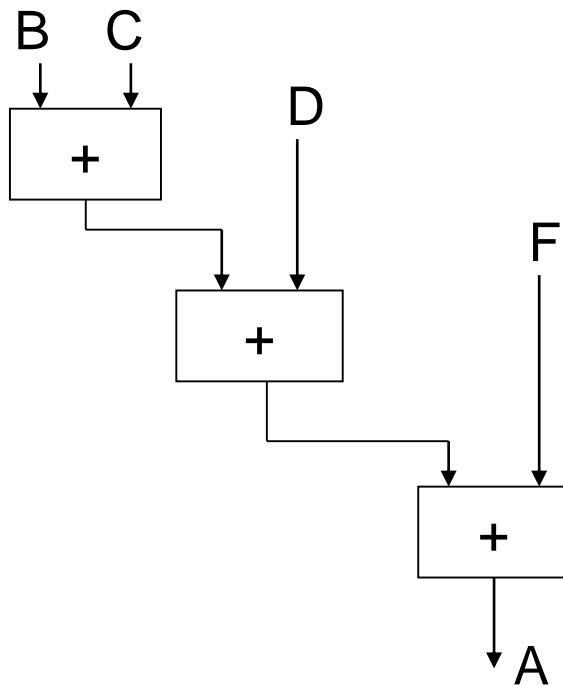
Signed/Unsigned Multiplication

- If N -bit multiplication is required, usually we use
 - **Unsigned multiplication**
 - **wire** [N-1:0] In1,In2;
 - **wire** [2N-1:0] MulOut;
 - **assign** MulOut=In1*In2;
 - **Signed multiplication**
 - **wire signed** [N-1:0] In1,In2;
 - **wire signed** [2N-1:0] MulOut;
 - **assign** MulOut=In1*In2;

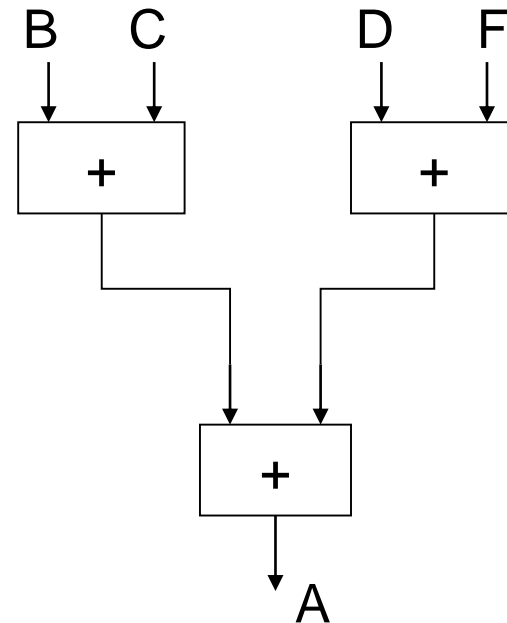


Using Parenthesis

$$A=B+C+D+F$$

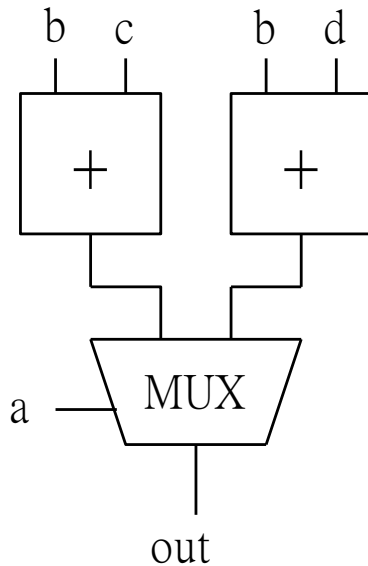


$$A=(B+C)+(D+F)$$



Resource Sharing (1/2)

```
always@(a or b or c or d)
  out =(a) ? (b+c):(b+d);
```

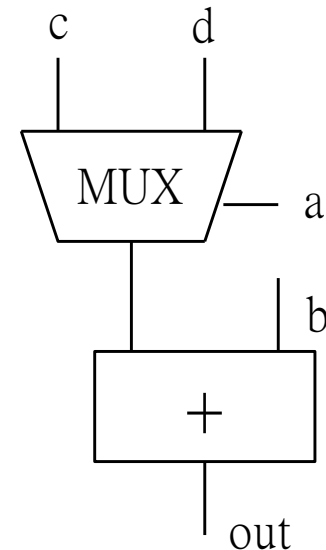


Keep sharable resource
in the

- ✓ same conditional statement
- ✓ same always block
- ✓ same module

without resource sharing

```
always@(a or b or c or d)
  if (a) out = b+c;
  else out = b+d ;
```



with resource sharing



Resource Sharing (2/2)

- The operators that can share resources must be in the mutual exclusive paths

```
if (sel1) out1=a1+b1;  
else begin  
  out1 = c1+d1 ;  
  if ( sel2) out2=a2+b2;  
  else out2=c2+d2;  
end
```

operator "c1+d1" and
"a2+b2" or "c2+d2" are **not**
in the mutual exclusive paths

```
if (sel1) out1=a1+b1;  
else begin  
out1 = c1+d1 ;  
  if ( sel2) out2=a2+b2;  
  else out2=c2+d2;  
end
```

All operators are in the
mutual exclusive paths

Resource sharing



Explicit Resource Sharing

■ Original

```
if (F)
    A=B+C;
else
    A=B+20;
```

■ Modified

```
if(F)
    T=C;
else
    T=20;
A=B+T;
```