Blockchain Technology and Applications

CS 731

Consensus in Bitcoin

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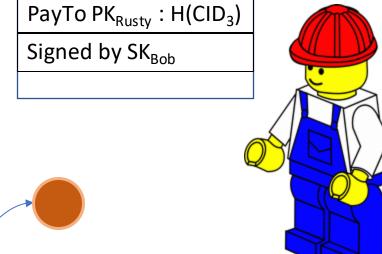
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Consensus

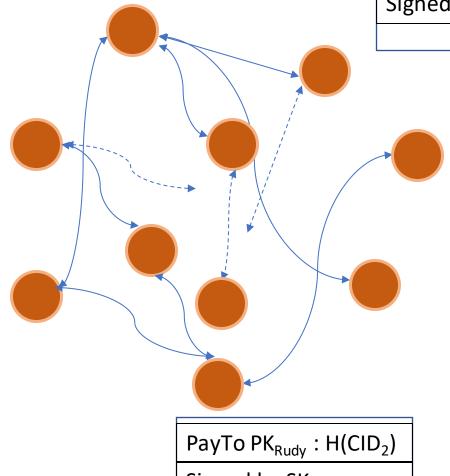
Broadcast transactions to the p2p network





PayTo PK_{Bob} : H(CID₁)

Signed by SK_{Alice}

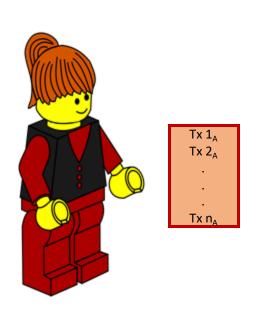


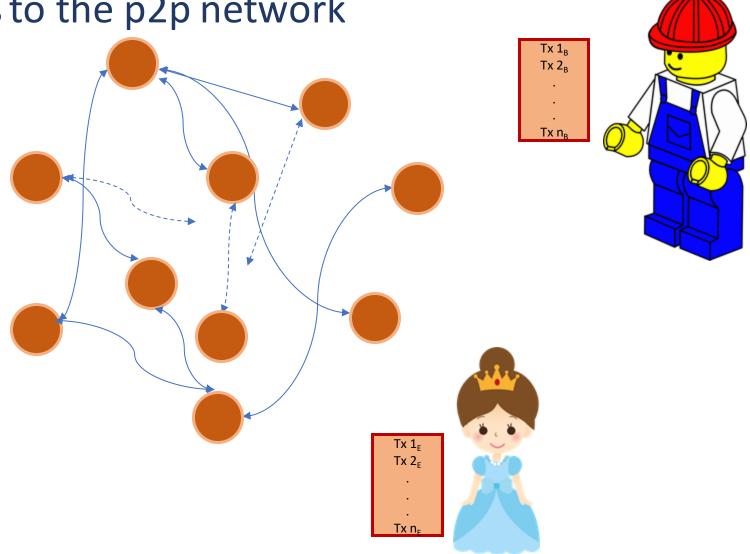
Signed by SK_{Eva}



Consensus

Broadcast transactions to the p2p network





- A simplified Bitcoin consensus algorithm
- 1. New transactions are broadcast to all nodes
- 2. Each node composes blocks with these transactions
- 3. In each round, a *random* node is selected
 - a. This node's block is selected as the next block
- 4. Other nodes accept this block
 - a. Only if all the transactions are valid
 - b. Funds are unspent
 - c. Signatures are valid, etc.
- 5. Nodes express their acceptance of the block by including its hash into the next block they create

Consensus in Bitcoin Lack of Identity

- There is no persistent identities in Bitcoin
 - Bitcoin uses cryptographic methods as identities
 - Generating public-key is trivial
 - Any node have multiple identities or public-keys
 - Anonymity was a design objective of Bitcoin

Lack of Identity

- Building a consensus using identities is simpler
 - Like a lottery system
 - Is not tied to real world identities
 - But there is an identity associated to each individual
 - At each round the node with certain bit arrangement will choose the next block
 - Can be held accountable if something goes wrong
- Owners can prove their ownership of their identities later

Consensus in Bitcoin Sybil attack

- It is trivial and cheap to create multiple identities
- An attacker can create an enormous number of identities
- If we use a simple random node choosing method
 - The attacker will gain disproportionate advantage
 - Can block the other nodes' blocks forever
- This is called *Sybil attack*, and the nodes *Sybil nodes*
- So, the random node choosing algorithm has to be intelligent
- Robust against Sybil attack

- For now let's assume
 - The random node choosing algorithm is intelligent
 - Works without persistent identities
- Lot to assume
 - We will show how these assumptions are satisfied

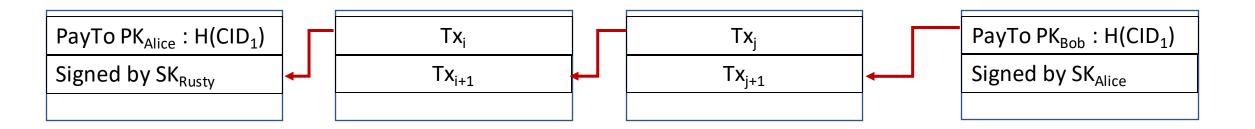
- Let's say the a node or an user Alice is chosen to propose the next block
- Assume Alice is dishonest
- What Alice can do?

Dishonest nodes

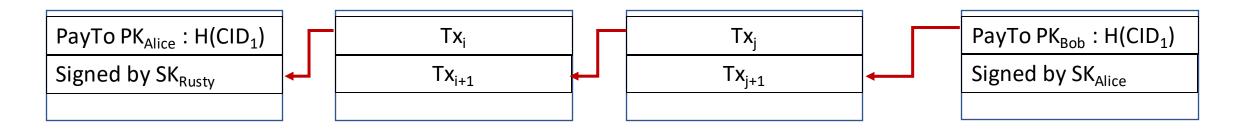
- Stealing Bitcoins?
- No Secured by digital signatures!
- Denial-of-Service?
- Alice excludes Bob transactions from her block
- Alice may prevent for one or more rounds
- Eventually, an honest node will be picked
- Bob's transaction will be included in the proposed block
- Minor inconvenience

Dishonest nodes

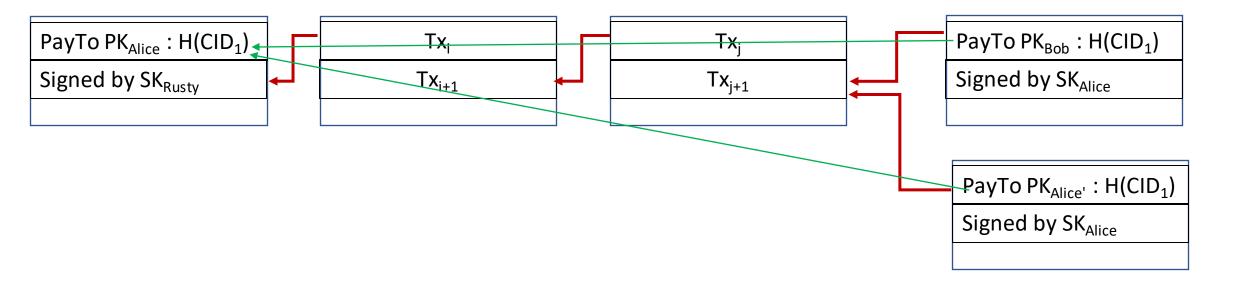
- Double-spend attack?
- Alice purchases service from Bob and pays in coins
- Broadcasts this transaction to the network
- Later, Alice pays the same coin to one of her accounts



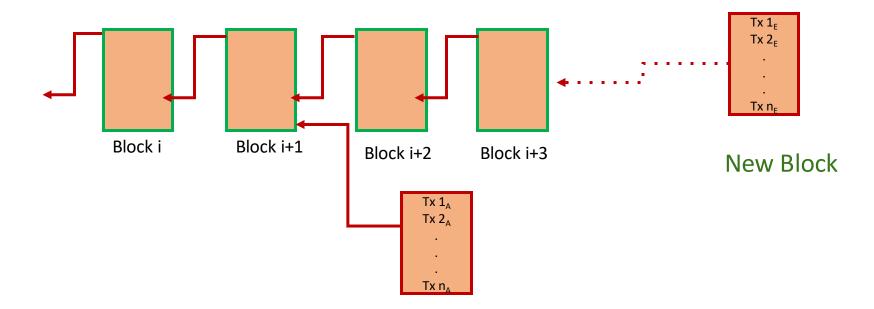
- Zero confirmation transaction
 - Bob ships the item as soon as he sees this transaction
- This can be included by an honest node



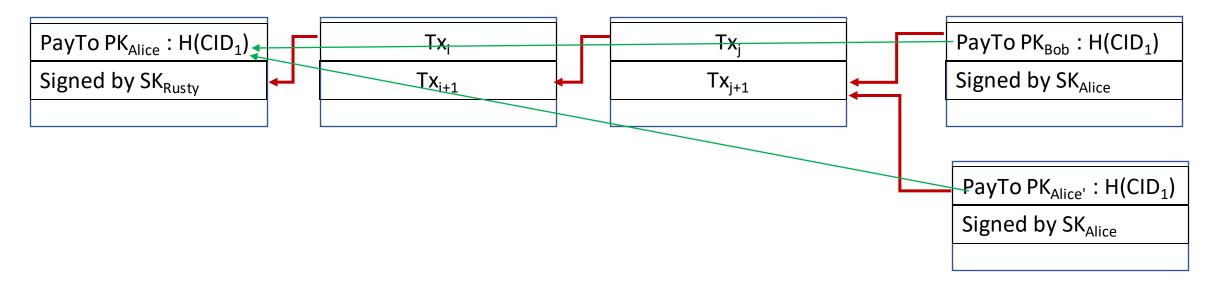
- Suppose Alice is chosen as the next random node
- She creates a payment to herself and ignore the previous node



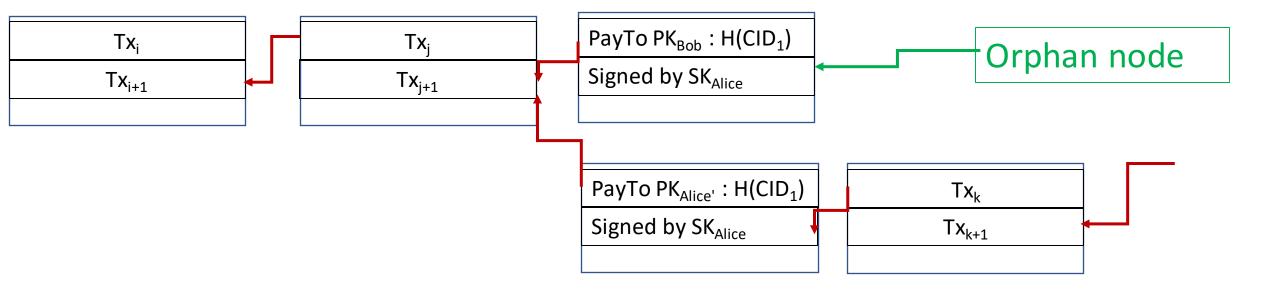
- Principle of "Extending longest valid chain"
- Honest nodes always adds their blocks to the longest valid chain



- Technically both chains are valid
- Network latency
 - Some nodes may hear Alice--> Alice transaction before Alice--> Bob transactions

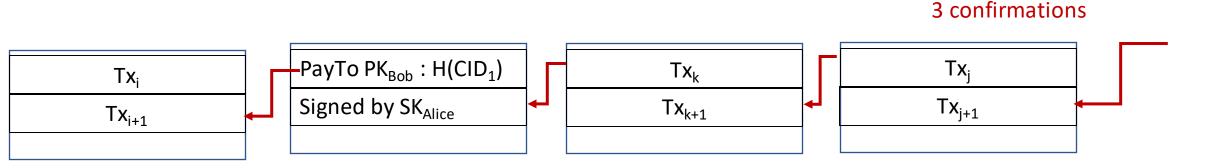


- Honest nodes build on the double-spend branch
- Alice can even bribe the next node
- The double-spend branch gets included in the chain
 - Successful attack



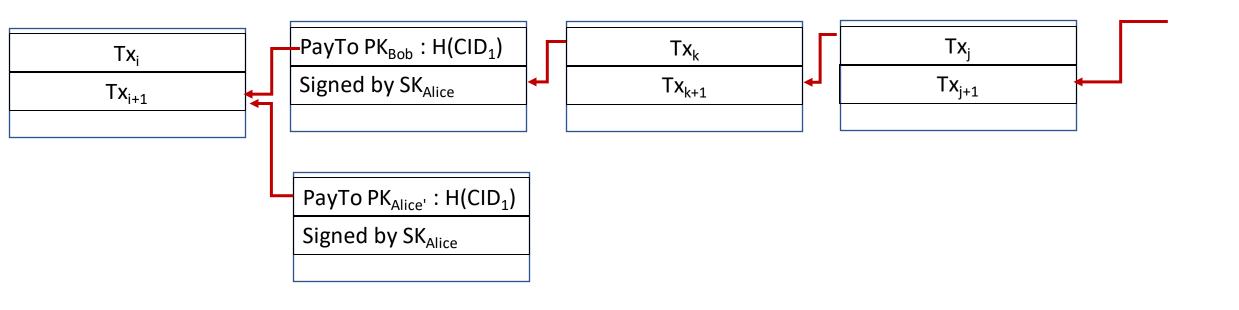
Bob's view

- Bob does not ship the item at the first confirmation
- Waits for more confirmations
- Only ships after enough confirmation
 - De facto standard is 6 in Bitcoin



Bob's view

- Alice can try to double spend
- Honest nodes will reject the block
- Longest valid chain



Bob's view

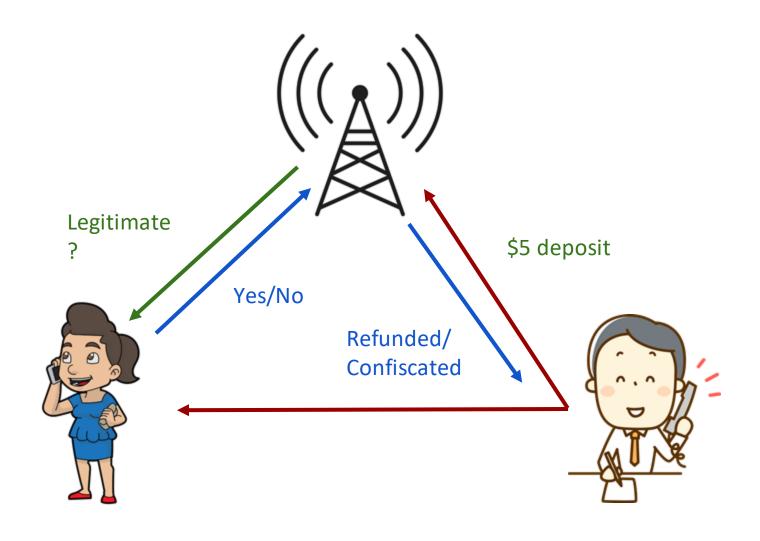
- Double-spend transaction probability decreases exponentially
 - With each confirmation
- Never 100% chance that the transaction is a success
- Virtually impossible after 6 transactions
 - Takes 1 hour!!!!!

- Recap
 - Protection against invalid transactions are cryptographic
 - Protection against double-spend is purely by consensus

Breaking away from traditional assumptions

- How Bitcoin overcomes impossibility results?
- Introduces randomness
 - Choose a node randomly somehow
 - No start/end time
 - Consensus happens over a long time
 - Divergence at the end of possible
 - Divergence probability decreases exponentially with time
- Incentives
 - Award/penalty for behaving honestly/dishonestly

Blocking Spam Callers



Incentives

- Our assumptions
- The process can pick a random node
 - At least 50% of the time honest node
- Can we punish the nodes who double-spends?
 - Difficult since there is no identity
- Alternatively, can we reward the honest nodes?
- Currency incentives
 - Possible since Bitcoin is a currency
 - Can award some Bitcoins for honest behaviour

Incentives

- Block reward
 - A node which creates/proposes a block can add a transaction
 - Awards certain number of bitcoins to itself
 - Coinbase transaction

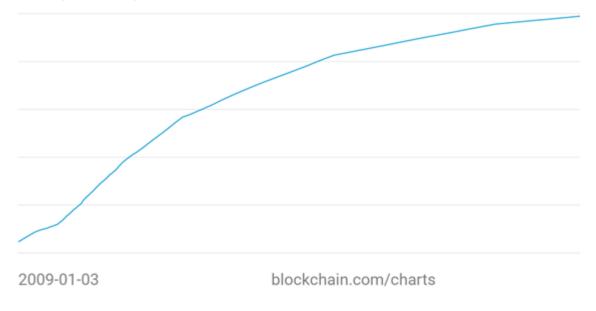
CID_x<--CreateCoin()

PayTo PK_{Bob} : H(CID_X)

Incentives

- 50 BTC in 2009
- Block reward halves every 210k blocks
 - Roughly each 4 years
- Currently 6.25 bitcoins
 - Last halved on May 11, 2020 Bitcoins in circulation

19,266,662.50 BTC



Incentives

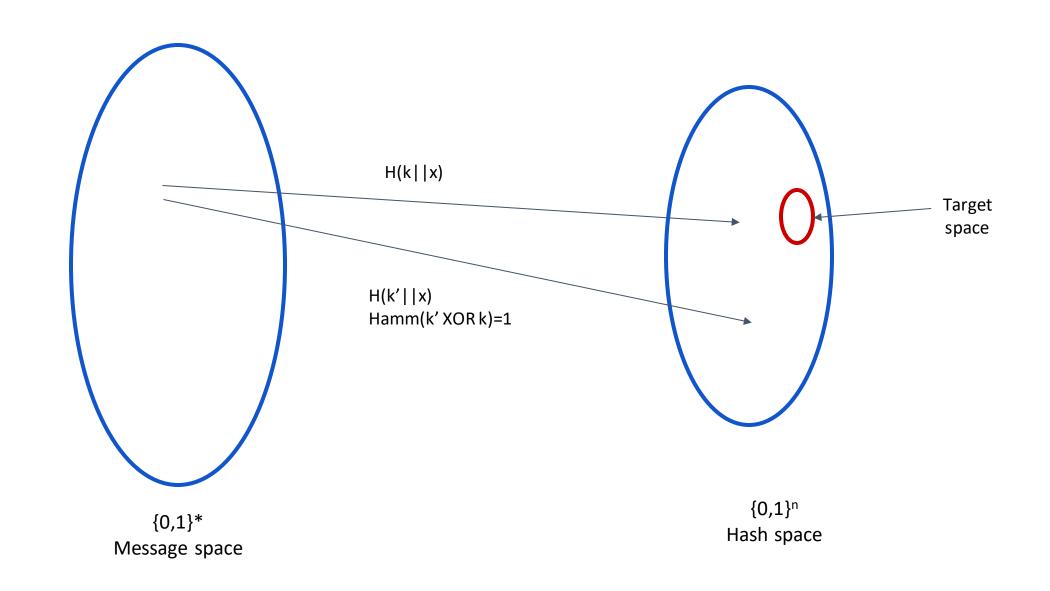
- Total number of bitcoins limited to 21 million
 - To limit the number of coins in circulation
 - Stop inflation
- Finite sum
 - Will end in 2040 approximately
- Then what?
 - Transaction fees
- Nodes may ask for some small fee for a transaction to be included in the block
- Total input Bitcoins > Total output Bitcoins
 - The difference is transaction fee

Proof-of-work

- Incentives solves the problem of reward/penalty
- How to pick a random node?
- Incentives introduces more problem
- Everybody wants to run Bitcoin nodes and get reward
- Sybil attacks!!
- Select nodes in proportion of ownership of resources
 - Hard to monopolize
- Computing resource --> Proof-of work
- Ownership of coins --> Proof-of-stake

Later

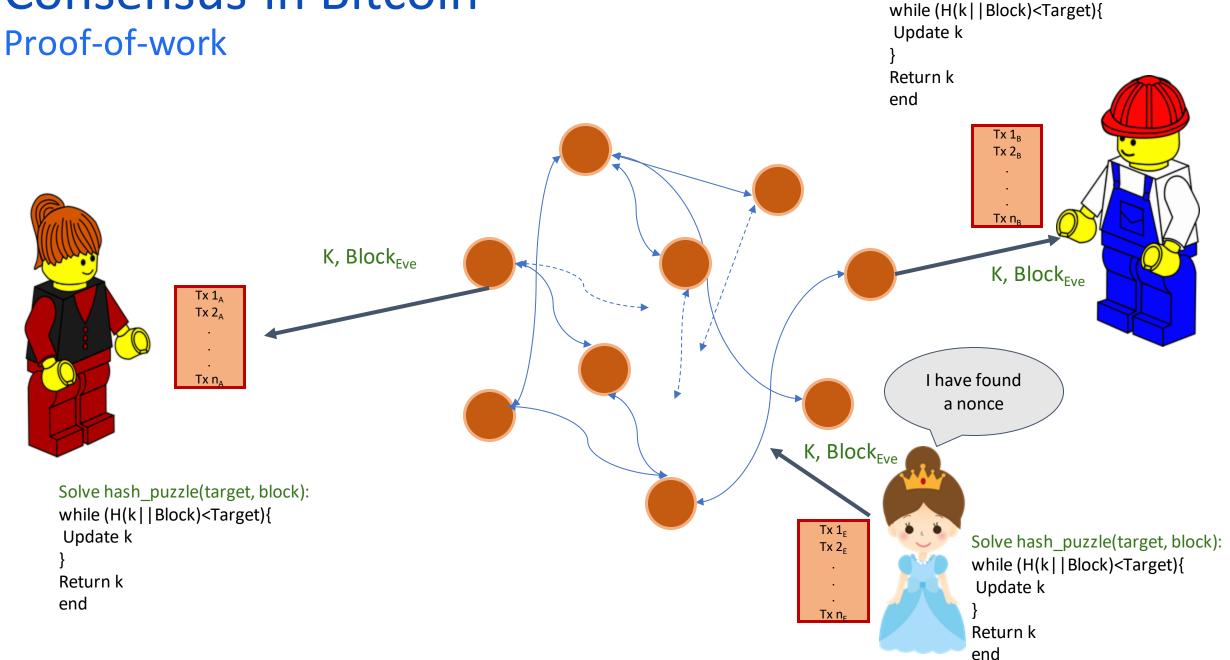
- Hash puzzles
- For fixed, x and find k (also called nonce), such that,
 - The hash i.e. H(k | |x) has some specific values
 - E.g. t number of bits out of n bits are zero



- As an example
- Let's assume the hash space has only 16 bits
 - We set our target circle with those hashes where the first 4 bits are set to 0 or 1
 - The probability of finding such values is 2¹²/2¹⁶
 - If the hash function is puzzle friendly we have to try 2¹⁶/2¹² different values of k in the worst case to find a hash in the target space
 - Or in other words, we have to perform this much hashes or work to find a hash in the target area
- The amount of hashes to perform to find the desired nonce is also called *difficulty* of the puzzle

- In Bitcoin, proof-of-work is achieved by solving hash puzzles
- Hashing requires computation power
 - Often requires powerful GPUs, ASICs, etc.
 - Requires significant time and monetary investment
- Difficult to monopolize
 - For a node to be selected with more than .5 probability the node has to control 50% of the global hashing power
- Apart from the incentive to behave honestly,
 - If a node behaves dishonestly all the work done to solve the puzzle and therefore the money required to do that is wasted
 - So it also serves as a penalty for dishonest behaviour

- A block has many data
 - Prev hash
 - Root hash
 - Data related to transactions
- Node must find a nonce such that
 - H(nonce | | tx | | prev_hash | | root_hash | | ...) < target
- Recall, for puzzle-friendly hash functions only way to find the nonce is by brute-force
- Trivial to verify



Solve hash_puzzle(target, block):

- We are allowing nodes to compete
- The nodes will be selected in the proportion of their %age of total global computing power
 - It is difficult to monopolize the total hash rate
- Sybil attack resistant
 - Nodes can still create many new identities
 - No extra advantage unless the hash rate is increased
 - Financial restriction
- Upholds our assumption on choice of random nodes

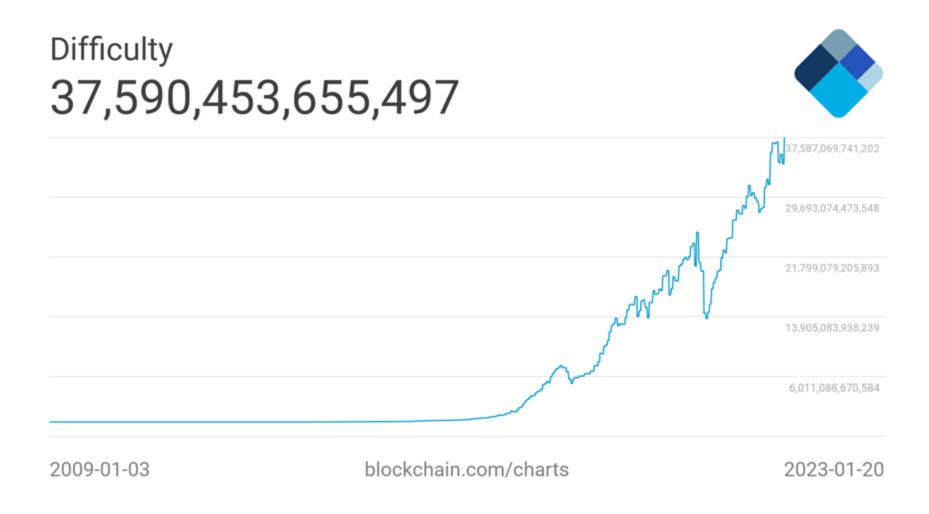
- Cost: the probability a miner is going to get the next block is proportional to the fraction of global hash power it controls
 - Alice controls .1% of hash power, so Alice is going to win roughly one block out of 1000 blocks
 - Parameterizable cost
- To sum up everything,
 - Roughly the profit from mining
 - Mining reward mining cost
 - Mining reward = block reward + transaction fees
 - Mining cost = hardware cost + electricity + cooling + real estate,
 etc.

- Solving this hash puzzle is also called mining
- They adjust the difficulty such that considering the total hash power of the network
 - The block finding time follows a Poisson distribution
 - It takes 10 minutes to find a new block on an average
 - Some blocks can be found earlier and some can be found later

- Nodes can leave or join network
- Hash rate of the network increases or decreases
- The p2p network automatically chooses new target
- So the difficulty should be adjusted to keep
 - To maintain the 10 minutes average block time
 - Every 2016 blocks
 - Every two weeks
- Why?
 - Caching
 - Can store the previous nonces
 - Global hash rate increases or decreases

Proof-of-work

Hash puzzles are difficult to compute



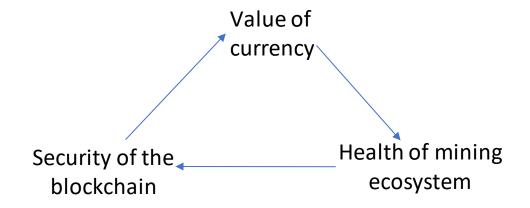
- Hashing requires computation power
 - Often requires powerful GPUs, ASICs, etc.
 - Requires significant time and monetary investment
- The cost of being dishonest is too high
- Proof-of-work + cost of mining ensures
 - Majority of miners, weighted by hash powers is honest
 - Random selection in the proportion of hash power will pick honest nodes
- It creates a stable equilibrium that nobody can get paid higher by being dishonest
- Still an active area of research

Breaking away from traditional assumptions

- Works in practice
- No theoretical explanation!!
 - Catching up fast
- Necessary for security and reliability

Bootstrapping

- The value of a cryptocurrency depends on three ideas
 - Value of the currency
 - Security of the blockchain
 - Health of mining eco-system
- Each one of these is dependent on the others



- Crucial for a new cryptocurrency to succeed
- Bootstrapping

51% attack

- Steal coins from existing address? X
- Suppress some transactions?
 - From the block chain
 - From the P2P network
- Change the block reward?
- Destroy confidence in Bitcoin?

Bitcoin: Fully decentralized?

- Designed to be fully decentralized.
- Anyone can join and leave
 - No central authority
- Mining
 - Technically open for everyone
 - High capital cost
 - Concentrated in few regions
- Rules and updates
 - Few trusted groups/users
- Hybrid
- Another example : SMTP

Bitcoin: Fully decentralized?

- Based on data provided by World Population Review, the current hash rates of the leading countries in Bitcoin mining, as of 2023, are as follows:
- United States: 35.4%
- Kazakhstan: 18.1%
- Russia: 11.23%
- Canada: 9.55%
- Ireland: 4.68%
- Malaysia: 4.58%
- Germany: 4.48%
- Iran: 3.1%

The End!!