## 

Structure 表示硬件上的物理thread实体

Cnt	
Fe_active	
Al_active	
Available	
Thread_using	
Cid	
Context[MAX_CONTEXT_NUM]	表示当发生mispredict的情况下,function model可以执行的最大的错误分支路径
	包括 rat_table, partial_flag Context_s Tid Phytid Prev_cid Next_cid Depth Prev_rob Prev_store_tail Al_active Rat Mapping Partial_flags Simrunner Stew Bigstew Ras_stack Ras_tos Call_depth Global_history Target VA Last_br CpuframeworkUop
Clonenum	Reached_rat
Deletenum	
Context_active_list	
Context_free_list	
Fe_mode	
Fe_prev_mode	
Fe_current_instr_disasm	
Fe_current_instr_bytes	
Fe_fetch_plr	
Plr_id	
Fe_fetch_stall	
Fe_icache_itlb_stall	
Fe_fetch_stall_reason	
Fe_stall_btb_notstalled	
Fe_uq_stall	
Is_64bit_mode	
Fe_transition_point	
Fe_restart	
Fe_sild_mode	
Fe_seen_prefix_bubble	
Uq_is_empty	
Fe_current_instr	CpuframeworkInstr_s
Fe_prev_first_uop	CpuframeworkUop_s
Cycle_last_clear	
Fe_sb_occupied	Sb=stream-buffer
Fe_sb_excellent_line	
Fe_stream_sent	
Fe_mtf_lip	

Fe_il1_line		
Fe_seen_ms		
Fetch_ul2_checked		
Fe_sild_addr		
Fe_in_irq		
Fe_excellent_address	Va	
Fe_excellent_address_paddr	Pa	
Fe_last_btflush	Va	
Fe_btflush_queue		
Inst_had_il1_miss		
Fe_previous_pos_in_line		
Fe_previous_line_number		
Fe_trace_plr		
Ms_ignore_fusing		
Itlb_miss		
Itlb_write		
Miss_in_itlb		
Tid		
Uops_fetched		
Uops_scheduled		
Uops_retired		
Al_stall		
Al_stall_bubble		
Mwaiting		
Block_younger_loads		
Look_for_store_unlock		
Lock_uop_num		
Lock_paddr	VA?	
Lock_vaddr	Va	
Num_locks		
Last_uop_retire_time		
Last_inst_retire_time		
Last_retired_uop_alloc_time		
Last_retired_uop_fetch_time		
Earliest_next_retire_time		
Last_retired_uop_action	Beuflush_action_t	
Stalled_fro_rob_entries		
Physical thread selection policy		

Physical thread selection policy

Physical thread selection policy
POLICY\_INDEPENDENT
POLICY\_PHYTHREAD\_INTERLEAVE
POLICY\_PHYTHREAD\_ROUNDROBIN
POLICY\_ICOUNT
POLICY\_ICOUNT
POLICY\_OLDEST\_FIRST
POLICY\_LEAST\_RECENTLY\_SERVICED

POLICY\_REPLICATE

POLICY\_TIME\_INTERLEAVE POLICY\_THREAD\_UNAWARE

## Logical thread => 对应core.h thread[threadnum] var 表示os看到的logic CPU的实体

State	THREAD_STATE THREAD_ACTIVE THREAD_READY THREAD_STALLED THREAD_ACTIVE_IDLE
Active	
Active_stall	
Phytid	
Last_phytid	
Num_stall_events	
Imiss_stall	
Cycle_last_end_drain	
Cycle_last_activate	
Cycle_last_inactivate	
Uop_fetch_last_inactivate	
Uop_retire_last_activate	
Cycles_last_act	
Retry_cycle	
Al_prevent_switch	

```
VA
Miss_vaddr
                  PA
Miss_paddr
Core的physical thread的轮询模板
ThreadState
struct ThreadState {
                          // 当前T正在处理的physical thread
 int phytid;
                         // 当前T开始轮询前的初始physical thread, 这个变量由global_thread_priority或是g_phytid[all_core][phytid_idx]决定
 int save_g_phytid;
                         // 当前T正在处理的physical thread处理了多少请求
 int count;
                       // 当前T已经有多少个physical thread已经处理过了
 int phytids_used_count;
 int phytids_used[MAX_THREADS]; // 当前T哪些physical thread处理了,哪些没有处理
 float min_used;
                          // 资源使用情况,只记录占用量最小的,(float)used_resource/limit_resource
g_phytid[all_core][phytid_idx]是个函数内的static变量,用于保存当前T执行完后,下一个T用于获取初始physical thread的信息,更新规则为:如果当前T有处理请求,则把
最后服务的physical thread记录到phytid_idx(仅仅是轮询时用的索引变量,与physical thread无关)的索引中
目前physical thread轮询支持的Policy
enum {
 POLICY_INDEPENDENT;
 // 遍历所有的physical thread,只要满足处理条件全部处理
 POLICY_PHYTHREAD_INTERLEAVE;
 // 找出满足条件的physical thread,处理一次,如果没有处理,则查找下个满足条件的physical thread,直到有一个可以处理或是遍历所有physical thread
 POLICY_PHYTHREAD_ROUNDROBIN;
 // 当前T遍历所有的满足条件的physical thread
 POLICY_ICOUNT;
 //首先找出使用资源最少的phyiscal thread进行处理;如果ThreadState.count已经处理过了,则当前T结束,否则轮询下个资源使用最少的physical thread
 POLICY_OLDEST_FIRST;
 // 同PHYTHREAD_ROUNDROBIN
 POLICY_LEAST_RECENT_SERVICED;
 //同PHYTHREAD_ROUNDROBIN
 POLICY_REPLICATED;
 // 同PHYTHREAD_ROUNDROBIN
 POLICY_TIME_INTERLEAVE;
  // 只处理当前初始选择的physical thread,无论处理条件是否满足
 POLICY_THREAD_UNWARE;
 // 同PHYTHREAD ROUNDROBIN
};
```