

• Multiplication of two polynomials requires n^2 multiplications and $(n-1)$ additions.

• karatsuba

$$f = f_0 + x f_1$$

$$g = g_0 + x g_1$$

$$f \cdot g = f_0 g_0 + (f_0 g_1 + g_1 f_0) x + f_1 g_1 x^2 \quad n=2$$

$$f_0 g_1 + g_1 f_0 = (f_0 + f_1)(g_0 + g_1) - f_0 g_0 - f_1 g_1$$

3 multiplication , 4 additions

It is under the assumption that the cost of multiplication is more expensive than addition. Karatsuba reduce the number of multiplication in cost of increasing number of additions.

Karatsuba(F, G) # size = n .

If $n=1$:

return $F \cdot G$.

Else:

$$F = F_0 + F_1 x^{\lceil n/2 \rceil}$$

$$G = G_0 + G_1 x^{\lceil n/2 \rceil}$$

$$M_0 = \text{Karatsuba}(F_0, G_0)$$

$$M_1 = \text{Karatsuba}(F_0 + F_1, G_0 + G_1)$$

$$M_2 = \text{Karatsuba}(F_1, G_1)$$

$$M = M_0 + (M_1 - M_0 - M_2) x^{n/2} + M_2 x^n$$

return M .

Ex

$$F = 2 + x + 2x^2 + 3x^3$$

$$G = 1 + 4x + 3x^2 + x^3$$

$$\boxed{2/72}$$

$$F = \underbrace{(2+x)}_{F_0} + \underbrace{(2+3x)}_{F_1} x^2$$

$\rightarrow \deg \leq n = 2$

$$n = 4$$

$$m = 2$$

$$G = \underbrace{(1+4x)}_{G_0} + \underbrace{(3+x)}_{G_1} x^2$$

$$(F_0 + F_1)(G_0 + G_1)$$

$= H_0 + H_1x + H_2x^2$

$$H = F \cdot G = H_0 + H_1x^2 + H_2x^4$$

$$F \cdot G = \underbrace{(2+x)(1+4x)}_{H_0} + \underbrace{[(2+x)(3+x) + (2+3x)(1+4x)]}_{H_1} x^2 + \underbrace{[(2+3x)(3+x)]}_{H_2} x^4$$

$$H_0 = 4x^2 + 2x + 2$$

$$H_2 = 3x^2 + 4x + 6$$

$$H_1 = 2 + x + 6x^2 - 4x^2 - 2x - 2 - 3x^2 - 4x - 6$$

$$= 6x^2 + 2x + 1$$

$$\begin{array}{lcl} x^2 & -1 \bmod 7 & = 6 \\ x & -5 \bmod 7 & = 2 \\ & -6 \bmod 7 & = 1 \end{array}$$

$$H_0 = 4x^2 + 2x + 2$$

$$H_0 + H_1x^2 + H_2x^4$$

$$H_1x^2 = x^2 + 2x^3 + 6x^4$$

$$= 2 + 2x + 5x^2 + 2x^3 + 5x^4 + 4x^5 + 3x^6$$

$$H_2x^4 = 6x^4 + 4x^5 + 3x^6$$

$$x^4 \rightarrow 12 \bmod 7 = 5$$

Master theorem

$m = 3$ (number of recursive call)

$$C(n) \leq m C(n/p) + T(n)$$

$$p = q = 2$$

$$T(n) = 4n$$

$$k(n) \leq 3k(n/2) + 4n$$

$$\underline{q < m}$$

$$C(n) = k(n) = O\left(n^{\log_p(m/q)} T(n)\right) = O\left(n^{\log_2(3/2)} 4n\right)$$

- Polynomial division.

Input: f , degree n

g , degree m where $m < n$

Output: q, r where $f = q \cdot g + r$
 $\deg(r) < \deg(g)$

Example

$$f = 4x^3 + 2x^2 + 5x + 1$$

$$g = x^2 - 1$$

$$f/g = ?$$

$$\begin{array}{r} 4x+2 \\ x^2-1 \overline{) 4x^3+2x^2+5x+1} \\ \underline{-(4x^3-4x)} \longrightarrow q \\ 2x^2+9x+1 \\ \underline{-(2x^2+0x-2)} \\ 9x+3 \longrightarrow r \end{array}$$

$$\begin{aligned} (q \cdot g) + r &= (x^2-1)(4x+2) + (9x+3) \\ &= 4x^3+2x^2-4x-2+9x+3 \\ &= 4x^3+2x^2+5x+1 \quad \checkmark \end{aligned}$$

Polynomial division

Input: F, G , Output: Q, R

$$R := F ; Q := 0$$

While ($\deg(R) \geq \deg(G)$) $O(n)$

$$t = \frac{\text{lead}(F)}{\text{lead}(g)} \times \frac{(\deg(R)-\deg(G))}{m-n} \quad O(m-n)$$

$$Q := Q + t$$

$$R := R - t \cdot G$$

$$O(n(m-n))$$