

Data Structures in Java - Homework 7

Problem 1

- 5.1 Given input {4371, 1323, 6173, 4199, 4344, 9679, 1989} and a hash function $h(x) = x \bmod 10$, show the resulting:
- Separate chaining hash table.
 - Hash table using linear probing.
 - Hash table using quadratic probing.

Answer:

(a) Index	(b) Index	(c) Index
0	0	0
1	1	1
2	2	2
3	3	3
4	4	4
5	5	5
6	6	6
7	7	7
8	8	8
9	9	9

Problem 2

- 5.2 Show the result of rehashing the hash tables in Exercise 5.1.
Use Table size of 23

Answer:

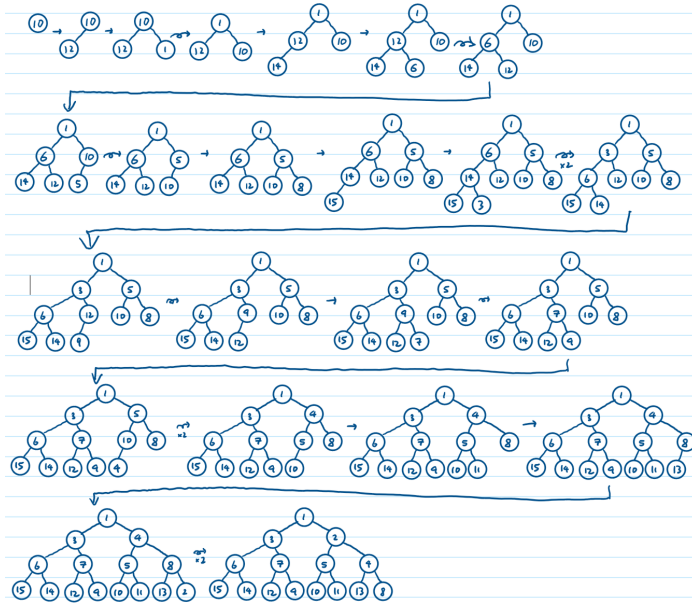
Same for all (a), (b) and (c):

Index				
0	8	16		
1	9	17		
2	10	18		
3	11	19	9679	
4	12	20	4344	
5	13	21		
6	14	22		
7	15			

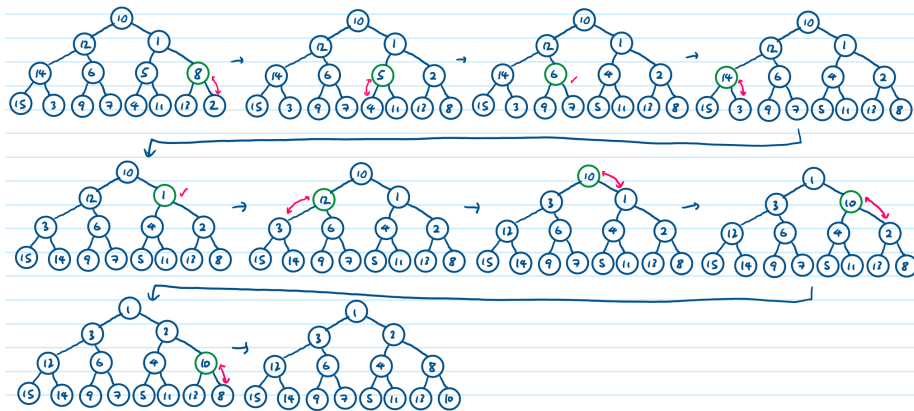
Problem 3

- 6.2 a. Show the result of inserting 10, 12, 1, 14, 6, 5, 8, 15, 3, 9, 7, 4, 11, 13, and 2, one at a time, into an initially empty binary heap.
 b. Show the result of using the linear-time algorithm to build a binary heap using the same input.

a) Answer:



b) Answer:



Problem 4

6.16 Suppose that binary heaps are represented using explicit links. Give a simple algorithm to find the tree node that is at implicit position i .

Answer:

Breadth-first traversal method and incrementing counter.

```
public BinaryNode get(int i)
{
    if (root == null || i < 1) return null;

    ArrayList<BinaryNode> queue = new ArrayList<BinaryNode>();
    queue.add(root); // access to root through class

    int currPos = 0;

    while (!queue.isEmpty()) {
        currPos++;
        BinaryNode currNode = queue.remove(0);

        if (currPos == i) {
            return currNode;
        }

        if (currNode.left != null) {
            queue.add(currNode.left);
        }
        if (currNode.right != null) {
            queue.add(currNode.right);
        }
    }

    return null;
}
```

Problem 5

6.18 A **min-max heap** is a data structure that supports both `deleteMin` and `deleteMax` in $O(\log N)$ per operation. The structure is identical to a binary heap, but the heap-order property is that for any node, X , at even depth, the element stored at X is smaller than the parent but larger than the grandparent (where this makes sense), and for any node X at odd depth, the element stored at X is larger than the parent but smaller than the grandparent. See Figure 6.57.

- a. How do we find the minimum and maximum elements?
- *b. Give an algorithm to insert a new node into the min-max heap.

a) **Answer:**

Minimum is the root. Maximum is the larger child of the root.

b) **Answer:**

Assuming heap is implemented in an array with extra space and the variable `size` is accessible through the class:

```
public void insert(Comparable item) {
    heap[size+1] = item; // assuming there is extra space
    size++;

    int idx = size;
    if (idx == 1) return;
    int parentIdx = idx/2;
    int depth = (int) (Math.log(idx) / Math.log(2));

    if (isEvenLevel(idx)) {
        if (item.compareTo(heap[parentIdx]) > 0)
        { // if item (even) is larger than parent (odd)
            swap(idx, parentIdx); // swap them
            percolateUpEven(parentIdx); // percolate the item up through odd levels
        } else { // otherwise
            percolateUpOdd(idx); // percolate up through even levels
        }
    } else { // if item is at odd level
        if (item.compareTo(heap[parentIdx]) < 0)
        { // if item (odd) is smaller than parent
            swap(idx, parentIdx); // swap them
            percolateUpEven(parentIdx); // percolate item up through even levels
        } else { // otherwise
            percolateUpOdd(idx); // percolate up through odd levels
        }
    }
}
```

Problem 6

7.1 Sort the sequence 3, 1, 4, 1, 5, 9, 2, 6, 5 using insertion sort.

Original: [3, 1, 4, 1, 5, 9, 2, 6, 5]

Swap 1: [1, 3 | 4, 1, 5, 9, 2, 6, 5]

Swap 2: [1, 3, 4 | 1, 5, 9, 2, 6, 5]

Swap 3: [1, 1, 3, 4 | 5, 9, 2, 6, 5]

Swap 4: [1, 1, 3, 4, 5 | 9, 2, 6, 5]

Swap 5: [1, 1, 3, 4, 5, 9 | 2, 6, 5]

Swap 6: [1, 1, 2, 3, 4, 5, 9 | 6, 5]

Swap 7: [1, 1, 2, 3, 4, 5, 6, 9 | 5]

Swap 8: [1, 1, 2, 3, 4, 5, 5, 6, 9]