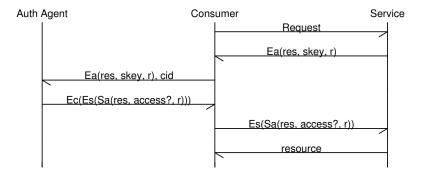
# Protocols Steven Allen April 18, 2013

# 1 Analysis Variables

- R Number of resources
- ${f G}$  Number of groups per person
- ${f F}$  Number of friends per person
- S Number of services
- M Average members per group
- r Request rate
- g Group change rate
- p Resource post rate

# 2 Obvious PK

## 2.1 Protocol



## 2.2 Costs

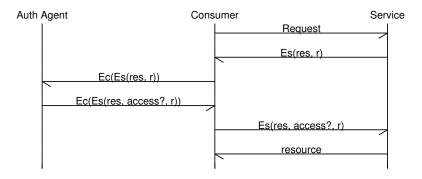
Op	Auth Agent	Consumer	Service
Encryption	$\Theta(r)$	$\Theta(r)$	$\Theta(r)$
Storage	$\Theta(RG + RS)$	$\Theta(1)$	$\Theta(R)$
Transfer	$\Theta(r)$	$\Theta(r)$	$\Theta(r)$

# 2.3 Analysis

Perfect privacy.

# 3 Obvious Shared Secret

## 3.1 Protocol



## 3.2 Costs

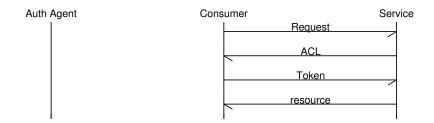
Op	Auth Agent	Consumer	Service
Encryption	$\Theta(r)$	$\Theta(r)$	$\Theta(r)$
Storage	$\Theta(RG + RS)$	$\Theta(F)$	$\Theta(R)$
Transfer	$\Theta(r)$	$\Theta(r)$	$\Theta(r)$

# 3.3 Analysis

Perfect privacy.

# 4 Public Key ACL (Basic)

## 4.1 Protocol



### 4.2 Costs

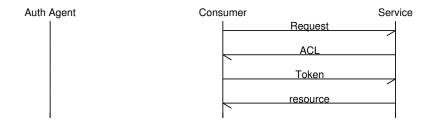
Op	Auth Agent	Consumer	Service
Encryption	$\Theta(g)$	$\Theta(rM)$	0
Storage	$\Theta(GM)$	$\Theta(1)$	$\Theta(RM)$
Transfer	$\Theta(g)$	$\Theta(rM)$	$\Theta(rM)$

## 4.3 Analysis

The ACL is of the form:  $\{E_{c_1}(t), E_{c_2}(t), ...\}$ . Consumers are able to determine the size of the ACL group.

# 5 Public Key ACL (Per Group)

## 5.1 Protocol



### 5.2 Costs

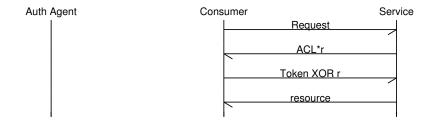
Op	Auth Agent	Consumer	Service
Encryption	$\Theta(g)$	$\Theta(rM)$	0
Storage	$\Theta(GM)$	$\Theta(1)$	$\Theta(GM)$
Transfer	$\Theta(g)$	$\Theta(rM)$	$\Theta(rM)$

## 5.3 Analysis

Both the consumers and the services are able to group content.

# 6 Public Key ACL (per group, enhanced)

## 6.1 Protocol



#### 6.2 Costs

Op	Auth Agent	Consumer	Service
Encryption	$\Theta(g)$	$\Theta(rM)$	0
Storage	$\Theta(GM)$	$\Theta(1)$	$\Theta(GM)$
Transfer	$\Theta(g)$	$\Theta(rM)$	$\Theta(rM)$

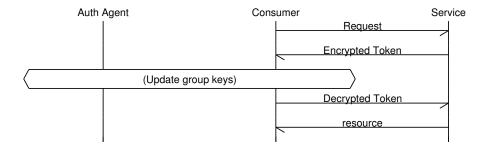
## 6.3 Analysis

Using Goldwasser-Micali for pk encryption, compute  $r * ACL = \{E_{c_1}(t \oplus r), E_{c_2}(t \oplus r), ...\}$ . This allows us to hide the group from the consumer.

With this encryption scheme, the consumers can't group content but the server still can.

## 7 El-Gamal Variant

#### 7.1 Protocol



#### 7.2 Costs

Op	Auth Agent	Consumer	Service
Encryption	$\Theta(p+g)$	$\Theta(r)$	$\Theta(r+p)$
Storage	$\Theta(GM)$	$\Theta(GF)$	$\Theta(R)$
Transfer	$\Theta(g+p)$	$\Theta(r)$	$\Theta(r+p)$

## 7.3 Description

This scheme uses a variant of ElGamal that allows for multiple uncorrelatable (to be proven) public keys per private key.

A loose description of the encryption scheme follows (there are a few other requirements but are not necessary to get the gist of it).

The consumer key consists of two parameters: x, w. The service keys are defined as  $g^{ax}, g^{a-w}$  where a is chosen randomly and uniquely per resource posted. Under this scheme, encryption is defined as  $E_a(m) = (g^r, g^{ar-wr}, mg^{axr}) = (j, k, l)$  and decryption is defined as  $D_a(j, k, l) = l/(kj^w)^x$ .

The token that the service sends the consumer is  $E_a(m)$ , and the decrypted token is m.

To prevent man-in-the-middle attacks, m should encode information about the server requesting authentication (IP address, a hash of their SSL key,

etc).

Also, as ElGamal (and therefore this scheme) is malleable, the message will have to be properly padded to ensure that it isn't modified.

# 8 Privacy Analysis

In this system, the server is unable to identify users, or even which groups guard a piece of content. However, users are able to tell if they are in the same group by comparing private group keys. Unfortunately, I was unable to come up with a way around this.