

Hazel Assistant Calculus WIP

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Abstract

The hazelnut assistant calculus provides an extensible framework for type- and value-directed completion and refactoring support in a structured editing context.

CCS Concepts: • Software and its engineering → General programming languages.

Keywords: live programming, code completion, refactoring, GUIs

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wazzzzzaaaaaaaaaaaaaaaaaaaaaa

1 Assistant Calculus

blah blah blah types

TODOs:

- get cursor icons from hazelnut paper
- get right arrow for bidi
- basic zipper cases
- remaining zipper cases? do i need to actually include mirror cases?
- var + varapp
- NOTE: we basically need a construct expression action for varapp
- proj
- base case for hole
- base cases for non-empty holes, incld:
- delete + act for general hexps
- simple wrap for exprs incld. non-empty holes
- complex (n-ary) wraps
- iterated wraps? with cutoffs? (:jean-shorts-emoji)
- are there non-empty hole suggests distinct from arbitrary expr suggests? don't think so

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$\boxed{\Gamma \vdash \hat{e} \Rightarrow \tau \curvearrowright A}$ e synthesizes τ , suggesting actions A

$\boxed{\Gamma \vdash \hat{e} \Leftarrow \tau \curvearrowright A}$ e analyzes against τ , suggesting actions A

$$\frac{\Gamma \vdash e_1 \Leftarrow \tau_2 \quad \Gamma \vdash \hat{e}_2 \Leftarrow \tau_1 \curvearrowright A}{\Gamma \vdash (\hat{e}_1, e_2) \Rightarrow \tau_1 \times \tau_2 \curvearrowright A}$$

$$\frac{\Gamma \vdash \hat{e} \Rightarrow \tau_1 \curvearrowright A \quad \tau_1 \blacktriangleright \tau_2 \rightarrow \tau \quad \Gamma \vdash e_2 \Leftarrow \tau_2}{\Gamma \vdash \hat{e}_1(e_2) \Rightarrow \tau \curvearrowright A}$$

$$\frac{\tau \blacktriangleright \tau_1 \rightarrow \tau_2 \quad \Gamma \vdash \hat{e} \Leftarrow \tau \curvearrowright A}{\Gamma \vdash \lambda x. \hat{e} \Leftarrow \tau \curvearrowright A}$$

Figure 1. Suggestion Zipper Cases

- for all: change type to type consistency
- for all: add numerical subscripts to types where missing
- fig 4: make it consistency not equal
- fig 5: change proj judgement to analysis
- fig 4: change varapp proj to have x in gamma, not gamma comma x
- consider matched product, matched arrow to suggest for unknown types
- above: not sure i want matched arrow type in fig 5? feels weird to suggest an unknown typed var for fn than value...
- rankings:
- privilege more specific types
- read contextual modal types
- replace constructs with construct expressions
- AppProj needs a better treatment for selection (should be first non-empty hole)... chain a separate action?
- implementation: implement ENTER vs TAB
- add keyboard shortcuts for swap etc
- proj1 and proj2: make tau in product another underscore
- intro rule for type=hole
- fig 4 var and appproj case: type consistency not premise: put in
- read: polarity: noam zalburger. bob harper blog post
- what does it mean to synthesize 'action a': sensibility theorem:
- well-typed in tsynthetid case, checked against tau in analytic case

$$\begin{array}{c}
\text{Suggest Hole Analytic} \\
\frac{\text{Intros}(\tau) \rightsquigarrow A_{\text{intros}} \quad \text{Elims}(\Gamma, \tau) \rightsquigarrow A_{\text{elims}}}{\Gamma \vdash \triangleright \langle \rangle \triangleleft \Leftarrow \tau \rightsquigarrow A_{\text{intros}} \cup A_{\text{elims}}} \\
\\
\text{Suggest Expr Analytic} \\
\frac{\text{Wraps}(e, \tau) \rightsquigarrow A_{\text{wraps}} \quad \text{Replaces}(\Gamma, \tau) \rightsquigarrow A_{\text{replaces}}}{\Gamma \vdash \triangleright e \triangleleft \Leftarrow \tau \rightsquigarrow A_{\text{wraps}} \cup A_{\text{replaces}}} \\
\\
\text{Replacement} \\
\frac{\Gamma \vdash \triangleright \langle \rangle \triangleleft \Leftarrow \tau \rightsquigarrow A}{\text{Replaces}(\Gamma, \tau) \rightsquigarrow \{\text{del} ; \alpha \mid \alpha \in A\}} \\
\\
\text{Wrapping (simple)} \\
\frac{\Gamma \vdash e \Rightarrow \tau_e \quad \tau_e \sim \tau'}{\text{Wraps}(e, \tau) \rightsquigarrow \{\text{construct } f(\triangleright e \triangleleft) \mid f : \tau' \rightarrow \tau \in \Gamma\}}
\end{array}$$

Figure 2. Suggestion base cases

$$\begin{array}{c}
\text{IntrosTriv} \\
\frac{\tau \sim 1}{\text{Intros}(\tau) \rightsquigarrow \{\text{construct } \triangleright \langle \rangle \triangleleft\}} \\
\\
\text{IntrosProd} \\
\frac{\tau \sim \tau_1 \times \tau_2}{\text{Intros}(\tau) \rightsquigarrow \{\text{construct } (\triangleright \langle \rangle \triangleleft, \langle \rangle)\}} \\
\\
\text{IntrosArrow} \\
\frac{\tau \sim \tau_1 \rightarrow \tau_2}{\text{Intros}(\tau) \rightsquigarrow \{\text{construct } \lambda x. \triangleright \langle \rangle \triangleleft\}} \\
\\
\text{IntrosSum} \\
\frac{\tau \sim \tau_1 + \tau_2}{\text{Intros}(\tau) \rightsquigarrow \{\text{construct } \text{inl}(\triangleright \langle \rangle \triangleleft), \text{construct } \text{inr}(\triangleright \langle \rangle \triangleleft)\}}
\end{array}$$

Figure 3. Introduction suggestions

$$\begin{array}{c}
\text{Suggest Elims} \\
\frac{\text{ElimCase} \rightsquigarrow A_{\text{case}} \quad \text{Var}(\Gamma, \tau) \rightsquigarrow A_{\text{var}} \quad \text{AppProj}(\Gamma, \tau) \rightsquigarrow A_{\text{appproj}}}{\text{Elims}(\Gamma, \tau) \rightsquigarrow A_{\text{case}} \cup A_{\text{var}} \cup A_{\text{appproj}}} \\
\\
\text{ElimCase} \\
\frac{}{\text{ElimCase} \rightsquigarrow \{\text{construct } \text{case}(\triangleright \langle \rangle \triangleleft, x. \langle \rangle, y. \langle \rangle)\}} \\
\\
\text{Var} \\
\frac{\tau \sim \tau'}{\text{Var}(\Gamma, \tau) \rightsquigarrow \{\text{construct } \triangleright x \triangleleft \mid x : \tau' \in \Gamma\}} \\
\\
\text{AppProj} \\
\frac{\tau \sim \tau'}{\text{AppProj}(\Gamma, \tau) \rightsquigarrow \{\text{construct } \triangleright e \triangleleft \mid x : \tau' \in \Gamma \wedge \Gamma \vdash x \rightsquigarrow e \Leftarrow \tau\}}
\end{array}$$

Figure 4. Elimination suggestions

$$\begin{array}{c}
\boxed{\Gamma \vdash e \rightsquigarrow e' \Leftarrow \tau} \quad e, \text{ applied to 0 or more holes,} \\
\text{and projected 0 or more times,} \\
\text{yields an expr analyzing to } \tau \\
\frac{\text{AppProjBase} \quad \Gamma \vdash e \Leftarrow \tau}{\Gamma \vdash e \rightsquigarrow e \Leftarrow \tau} \\
\\
\text{App} \\
\frac{\Gamma \vdash e \Rightarrow \tau' \quad \tau' \blacktriangleright _ \rightarrow _ \quad \Gamma \vdash e(\langle \rangle) \rightsquigarrow e' \Leftarrow \tau}{\Gamma \vdash e \rightsquigarrow e' \Leftarrow \tau} \\
\\
\text{Proj1} \\
\frac{\Gamma \vdash e \Rightarrow \tau' \quad \tau' \blacktriangleright_{\times} \tau \times _ \quad \Gamma \vdash \pi_1 e \rightsquigarrow e' \Leftarrow \tau}{\Gamma \vdash e \rightsquigarrow e' \Leftarrow \tau} \\
\\
\text{Proj2} \\
\frac{\Gamma \vdash e \Rightarrow \tau' \quad \tau' \blacktriangleright_{\times} _ \times \tau \quad \Gamma \vdash \pi_2 e \rightsquigarrow e' \Leftarrow \tau}{\Gamma \vdash e \rightsquigarrow e' \Leftarrow \tau}
\end{array}$$

Figure 5. Supporting elimination judgments

References