

Hazel Phi: 9-type-aliases

July 21, 2021

SYNTAX

| | | | |
|---------------------|----------------|-------|--|
| Kind | κ | $::=$ | $\mathbf{Type} \mid \mathbf{KHole} \mid \mathbf{S}_\kappa(\tau) \mid \Pi_{t::\kappa_1}.\kappa_2$ |
| User Types | $\hat{\tau}$ | $::=$ | $t \mid \mathbf{bse} \mid \tau_1 \oplus \tau_2 \mid \langle \rangle^u \mid \langle \hat{\tau} \rangle^u \mid \lambda t::\mathbf{Type}.\hat{\tau} \mid \tau_1 \tau_2$ |
| Internal Types | τ | $::=$ | $t \mid \mathbf{bse} \mid \tau_1 \oplus \tau_2 \mid \langle \rangle^u \mid \langle \tau \rangle^u \mid \langle t \rangle^u \mid \lambda t::\kappa.\tau \mid \tau_1 \tau_2$ |
| Base Types | \mathbf{bse} | $::=$ | $\mathbf{Int} \mid \mathbf{Float} \mid \mathbf{Bool}$ |
| BinOp | \oplus | $::=$ | $\times \mid + \mid \rightarrow$ |
| Type Pattern | | | |
| User Expression | | | |
| Internal Expression | | | |

DECLARATIVES

$\Delta; \Phi \vdash \tau ::> \kappa$ τ has principal (well formed) kind κ

$$\frac{\Delta; \Phi \vdash \mathbf{OK}}{\Delta; \Phi \vdash \mathbf{bse} ::> \mathbf{S}_{\mathbf{Type}}(\mathbf{bse})} \quad (1)$$

$$\frac{\Delta; \Phi_1, t::\kappa, \Phi_2 \vdash \mathbf{OK}}{\Delta; \Phi \vdash t ::> \mathbf{S}_\kappa(t)} \quad (2)$$

$$\frac{\Delta; \Phi \vdash \mathbf{OK} \quad \Delta; \Phi \vdash \tau_1::\mathbf{Type} \quad \Delta; \Phi \vdash \tau_2::\mathbf{Type}}{\Delta; \Phi \vdash \tau_1 \oplus \tau_2 ::> \mathbf{S}_{\mathbf{Type}}(\tau_1 \oplus \tau_2)} \quad (3)$$

$$\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \mathbf{OK}}{\Delta; \Phi \vdash \langle \rangle^u ::> \mathbf{S}_\kappa(\langle \rangle^u)} \quad (4)$$

$$\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \mathbf{OK} \quad \Delta; \Phi \vdash \tau::\kappa_1}{\Delta; \Phi \vdash \langle \tau \rangle^u ::> \mathbf{S}_\kappa(\langle \tau \rangle^u)} \quad (5)$$

$$\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \mathbf{OK} \quad t \notin \text{dom}(\Phi)}{\Delta; \Phi \vdash \langle t \rangle^u ::> \mathbf{S}_\kappa(\langle t \rangle^u)} \quad (6)$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \mathbf{OK} \quad \Delta; \Phi, t::\kappa_1 \vdash \tau ::> \kappa_2}{\Delta; \Phi \vdash \lambda t::\kappa_1.\tau ::> \mathbf{S}_{\Pi_{t::\kappa_1}.\kappa_2}(\lambda t::\kappa_1.\tau)} \quad (7)$$

$$\frac{\Delta; \Phi \vdash \tau_1 ::> \kappa \quad \Delta; \Phi \vdash \kappa \blacktriangleright \Pi_{t::\kappa_1}.\kappa_2 \quad \Delta; \Phi \vdash \tau_2::\kappa_1}{\Delta; \Phi \vdash \tau_1 \tau_2 ::> [\tau_2/t]\kappa_2} \quad (8)$$

$\Delta; \Phi \vdash \tau :: \kappa$ τ is well formed at kind κ

$$\frac{\Delta; \Phi \vdash \tau :: > \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \quad (9)$$

$$\frac{\Delta; \Phi \vdash \tau :: \mathbf{S}_{\kappa}(\tau_I) \quad \Delta; \Phi \vdash \tau_I :: \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \quad (10)$$

$$\frac{\Delta; \Phi \vdash \tau :: > \kappa_I \quad \Delta; \Phi \vdash \kappa_I \lesssim \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \quad (11)$$

$$\frac{\Delta; \Phi \vdash \tau_2 :: \mathbf{S}_{\kappa}(\tau_1)}{\Delta; \Phi \vdash \tau_1 :: \mathbf{S}_{\kappa}(\tau_2)} \quad (12)$$

$$\frac{\Delta; \Phi \vdash \tau_I :: \mathbf{S}_{\kappa}(\tau_3) \quad \Delta; \Phi \vdash \tau_3 :: \mathbf{S}_{\kappa}(\tau_2)}{\Delta; \Phi \vdash \tau_I :: \mathbf{S}_{\kappa}(\tau_2)} \quad (13)$$

$$\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \tau :: \mathbf{S}_{\kappa}(\tau)} \quad (14)$$

$\Delta; \Phi \vdash \kappa \blacktriangleright_{\Pi} \Pi_{t::\kappa_I} \cdot \kappa_2$ κ has matched Π -kind $\Pi_{t::\kappa_I} \cdot \kappa_2$

$$\frac{\Delta; \Phi \vdash \kappa \equiv \mathbf{KHole}}{\Delta; \Phi \vdash \kappa \blacktriangleright_{\Pi} \Pi_{t::\mathbf{KHole}} \cdot \mathbf{KHole}} \quad (15)$$

$$\frac{\Delta; \Phi \vdash \kappa \equiv \Pi_{t::\kappa_I} \cdot \kappa_2}{\Delta; \Phi \vdash \kappa \blacktriangleright_{\Pi} \Pi_{t::\kappa_I} \cdot \kappa_2} \quad (16)$$

$\Delta; \Phi \vdash \kappa_I \equiv \kappa_2$ κ_I is equivalent to κ_2

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \equiv \kappa} \quad (17)$$

$$\frac{\Delta; \Phi \vdash \kappa_2 \equiv \kappa_I}{\Delta; \Phi \vdash \kappa_I \equiv \kappa_2} \quad (18)$$

$$\frac{\Delta; \Phi \vdash \kappa_I \equiv \kappa_3 \quad \Delta; \Phi \vdash \kappa_3 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_I \equiv \kappa_2} \quad (19)$$

$$\frac{\Delta; \Phi \vdash \tau :: \mathbf{S}_{\kappa}(\tau_I)}{\Delta; \Phi \vdash \mathbf{S}_{\mathbf{S}_{\kappa}(\tau_I)}(\tau) \equiv \mathbf{S}_{\kappa}(\tau_I)} \quad (20)$$

$$\frac{\Delta; \Phi \vdash \tau :: \mathbf{KHole}}{\Delta; \Phi \vdash \mathbf{S}_{\mathbf{KHole}}(\tau) \equiv \mathbf{KHole}} \quad (21)$$

$$\frac{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_I} \cdot \kappa_2}{\Delta; \Phi \vdash \mathbf{S}_{\Pi_{t::\kappa_I} \cdot \kappa_2}(\tau) \equiv \Pi_{t::\kappa_I} \cdot \mathbf{S}_{\kappa_2}(\tau \ t)} \quad (22)$$

$$\frac{\Delta; \Phi \vdash \kappa_I \equiv \kappa_2 \quad \Delta; \Phi, t::\kappa_I \vdash \kappa_3 \equiv \kappa_4}{\Delta; \Phi \vdash \Pi_{t::\kappa_I} \cdot \kappa_2 \equiv \Pi_{t::\kappa_3} \cdot \kappa_4} \quad (23)$$

$$\frac{\Delta; \Phi \vdash \tau_I \overset{\kappa}{\equiv} \tau_2}{\Delta; \Phi \vdash \mathbf{S}_{\kappa}(\tau_I) \equiv \mathbf{S}_{\kappa}(\tau_2)} \quad (24)$$

$\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$ κ_1 is a consistent subkind of κ_2

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \text{KHole} \lesssim \kappa} \quad (25)$$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \text{KHole}} \quad (26)$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \quad (27) \quad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_3 \quad \Delta; \Phi \vdash \kappa_3 \lesssim \kappa_4 \quad \Delta; \Phi \vdash \kappa_4 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \quad (28)$$

$$\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \mathbf{S}_\kappa(\tau) \lesssim \kappa} \quad (29)$$

$$\frac{\Delta; \Phi \vdash \kappa_3 \lesssim \kappa_1 \quad \Delta; \Phi, t :: \kappa_3 \vdash \kappa_2 \lesssim \kappa_4}{\Delta; \Phi \vdash \Pi_{t :: \kappa_1} \cdot \kappa_2 \lesssim \Pi_{t :: \kappa_3} \cdot \kappa_4} \quad (30)$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2 \quad \Delta; \Phi \vdash \tau_1 \equiv^{\kappa_1} \tau_2}{\Delta; \Phi \vdash \mathbf{S}_{\kappa_1}(\tau_1) \lesssim \mathbf{S}_{\kappa_2}(\tau_2)} \quad (31)$$

$\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2$ τ_1 is equivalent to τ_2 at kind κ

$$\frac{\Delta; \Phi \vdash \tau_1 :: \mathbf{S}_\kappa(\tau_2)}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} \quad (32)$$

$$\begin{aligned} & \frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \tau \equiv^{\kappa} \tau} \quad (33) \quad \frac{\Delta; \Phi \vdash \tau_2 \equiv^{\kappa} \tau_1}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} \quad (34) \quad \frac{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_3 \quad \Delta; \Phi \vdash \tau_3 \equiv^{\kappa} \tau_1}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} \quad (35) \\ & \frac{\Delta; \Phi \vdash \tau_1 \equiv^{\text{Type}} \tau_3 \quad \Delta; \Phi \vdash \tau_2 \equiv^{\text{Type}} \tau_4}{\Delta; \Phi \vdash \tau_1 \oplus \tau_2 \equiv^{\text{Type}} \tau_3 \oplus \tau_4} \quad (36) \quad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \quad \Delta; \Phi, t :: \kappa_1 \vdash \tau_1 \equiv^{\kappa} \tau_2}{\Delta; \Phi \vdash \lambda t :: \kappa_1 \cdot \tau_1 \equiv \lambda t :: \kappa_2 \cdot \tau_2} \quad (37) \\ & \frac{\Delta; \Phi \vdash \tau_1 \equiv^{\Pi_{t :: \kappa_1} \cdot \kappa_2} \tau_3 \quad \Delta; \Phi \vdash \tau_2 \equiv^{\kappa_1} \tau_4}{\Delta; \Phi \vdash \tau_1 \tau_2 \equiv^{\tau_1 / t \cdot \kappa_2} \tau_3 \tau_4} \quad (38) \\ & \frac{\Delta; \Phi \vdash \tau_1 :: \Pi_{t :: \kappa_1} \cdot \kappa_3 \quad \Delta; \Phi \vdash \tau_2 :: \Pi_{t :: \kappa_1} \cdot \kappa_4 \quad \Delta; \Phi, t :: \kappa_1 \vdash \tau_1 \equiv^{\kappa_2} \tau_2 \quad t}{\Delta; \Phi \vdash \tau_1 \equiv^{\Pi_{t :: \kappa_1} \cdot \kappa_2} \tau_2} \quad (39) \\ & \frac{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa_1} \tau_2 \quad \Delta; \Phi \vdash \kappa_1 \equiv \kappa}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} \quad (40) \end{aligned}$$

$\Delta; \Phi \vdash \kappa \text{ OK}$ κ is well formed

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{Type} \text{ OK}} \quad (41)$$

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{KHole} \text{ OK}} \quad (42)$$

$$\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \mathbf{S}_\kappa(\tau) \text{ OK}} \quad (43)$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \text{ OK} \quad \Delta; \Phi, t :: \kappa_1 \vdash \kappa_2 \text{ OK}}{\Delta; \Phi \vdash \Pi_{t :: \kappa_1} \cdot \kappa_2 \text{ OK}} \quad (44)$$

$\boxed{\Delta; \Phi \vdash \text{OK}}$ Context is well formed

$$\frac{}{\vdash \text{OK}} \quad (45)$$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi, t::\kappa \vdash \text{OK}} \quad (46)$$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta, u::\kappa; \Phi \vdash \text{OK}} \quad (47)$$

Variables implicitly assumed to be fresh as necessary

METATHEORY

Lemma 1 (substitution). *If $\Delta; \Phi \vdash \tau_1::\kappa_1$ and $\Delta; \Phi, t::\kappa_1 \vdash \tau_2::\kappa_2$, then $\Delta; \Phi \vdash [\tau_1/t]\tau_2::\kappa_2$*

Lemma 2. *If $\Delta; \Phi \vdash \tau::> \kappa$, then $\Delta; \Phi \vdash \text{OK}$ and $\Delta; \Phi \vdash \kappa \text{ OK}$*

Lemma 3. *If $\Delta; \Phi \vdash \tau::\kappa$, then $\Delta; \Phi \vdash \text{OK}$ and $\Delta; \Phi \vdash \kappa \text{ OK}$*

Lemma 4. *If $\Delta; \Phi \vdash \kappa \xrightarrow{\Pi} \Pi_{t::\kappa_1}.\kappa_2$, then $\Delta; \Phi \vdash \text{OK}$ and $\Delta; \Phi \vdash \kappa \text{ OK}$ and $\Delta; \Phi \vdash \Pi_{t::\kappa_1}.\kappa_2 \text{ OK}$*

Lemma 5. *If $\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2$, then $\Delta; \Phi \vdash \text{OK}$ and $\Delta; \Phi \vdash \kappa_1 \text{ OK}$ and $\Delta; \Phi \vdash \kappa_2 \text{ OK}$*

Lemma 6. *If $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$, then $\Delta; \Phi \vdash \text{OK}$ and $\Delta; \Phi \vdash \kappa_1 \text{ OK}$ and $\Delta; \Phi \vdash \kappa_2 \text{ OK}$*

Lemma 7. *If $\Delta; \Phi \vdash \kappa \text{ OK}$, then $\Delta; \Phi \vdash \text{OK}$*

Proof. By simultaneous rule induction/length of proof.

The interesting cases per lemma:

| | | |
|----------|--|------------|
| L1. (1) | $\Delta; \Phi \vdash \text{bse}::\text{S}_{\text{Type}}(\text{bse})$ | by (9) |
| | $\Delta; \Phi \vdash \text{bse}::\text{Type}$ | by (10) |
| * | $\Delta; \Phi \vdash \text{S}_{\text{Type}}(\text{bse}) \text{ OK}$ | by (43) |
| * | $\Delta; \Phi \vdash \text{OK}$ | by premiss |
| (8) | | bad |
| L2. (12) | $\Delta; \Phi \vdash \tau_2::\kappa$ | by (10) |
| * | $\Delta; \Phi \vdash \text{S}_{\kappa}(\tau_2) \text{ OK}$ | by (43) |
| L4. (22) | $\Delta; \Phi \vdash \tau t::> \kappa$ | |

□

Lemma 8. *If $\Delta; \Phi \vdash \tau::> \kappa_1$ and $\Delta; \Phi \vdash \tau::\kappa_2$, then $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$*