

Hazel PHI: 10-modules

June 29, 2021

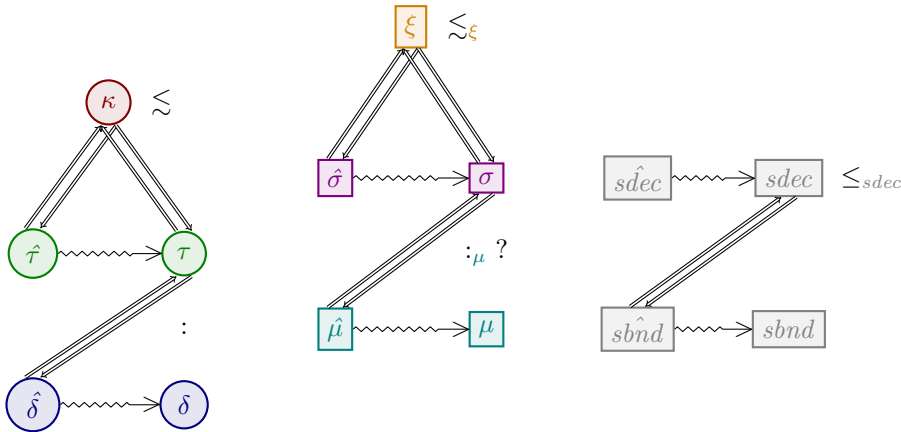
prerequisites

- Hazel PHI: 9-type-aliases-redux
 - github
 - current commit: 4410cd565ce717707e580e44f64868d3175fe2a6
- (optional) Hazel PHI: 1-labeled-tuples
 - github
 - current commit: 0a7d0b53ee7286d03ea3be13a7ac91a86f1c90b1

how to read

| | | | |
|--------|----------------------|--------|----------------|
| 800000 | kinds | D08000 | signature kind |
| 008000 | types (constructors) | 800080 | signatures |
| 000080 | terms | 008080 | modules |

notes



external typ/sig/mod syntax not written out yet (waiting for construction dust to settle); patterns not handled yet– will be left till end.

syntax

| | | | | |
|------|----------|-----|------------------------------|-------------------------|
| kind | κ | ::= | Type | kind of types |
| | | | $S(\tau)$ | singleton kind |
| | | | $KHole$ | kind hole |
| | | | $\Pi_{t::\kappa_1}.\kappa_2$ | dependent function kind |

| | | | | |
|------------------------|----------------|-----|--|---|
| HTyp | τ | ::= | t bse $\tau_1 \oplus \tau_2$ $[\tau]$ $\lambda t :: \kappa. \tau$ $\tau_1 \tau_2$ $\{lab_1 \hookrightarrow \tau_1, \dots lab_n \hookrightarrow \tau_n\}$ $\mu.lab$ $\langle \rangle$ $\langle \tau \rangle$ | type variable base type type binop list type type function type application labelled product type (record) module type projection empty type hole nonempty type hole |
| base type | bse | ::= | \mathbf{Int} \mathbf{Float} \mathbf{Bool} | |
| HTyp BinOp | \oplus | ::= | \times $+$ \rightarrow | |
| external expression | $\hat{\delta}$ | ::= | \dots x $\mathbf{signature} \ s = \hat{\sigma} \ \mathbf{in} \ \hat{\delta}$ $\mathbf{module} \ m = \hat{\mu} \ \mathbf{in} \ \hat{\delta}$ $\mathbf{module} \ m :_{\mu} s = \hat{\mu} \ \mathbf{in} \ \hat{\delta}$ $\mathbf{functor} \ \mathbf{something} = \mathbf{something} \ \mathbf{in} \ \hat{\delta}$ $\hat{\mu}.lab$ | module term projection |
| internal expression | δ | ::= | \dots x $\mathbf{signature} \ s = \sigma \ \mathbf{in} \ \delta$ $\mathbf{module} \ m :_{\mu} s = \mu \ \mathbf{in} \ \delta$ $\mathbf{functor} \ \mathbf{something} = \mathbf{something} \ \mathbf{in} \ \delta$ $\mu.lab$ | module term projection |
| signature kind | ξ | ::= | $\mathbf{SSigKind}(\sigma)$ $\mathbf{SigKHole}$ | |
| signature | σ | ::= | s $\{sdecs\}$ $\Pi_{m :_{\mu} \sigma_1}. \sigma_2$ $\langle \rangle$ $\langle s \rangle$ | signature variable structure signature functor signature empty signature hole nonempty signature hole |
| module | μ | ::= | m $\{sbnds\}$ $\lambda m :_{\mu} \sigma. \mu$ $\mu_1 \ \mu_2$ $\mu.lab$ $\langle \rangle$ $\langle \mu \rangle$ | module variable structure functor functor application submodule projection empty module hole nonempty module hole |
| signature declarations | $sdecs$ | ::= | \cdot $sdec, sdecs$ | |
| signature declaration | $sdec$ | ::= | $\mathbf{type} \ lab$ $\mathbf{type} \ lab = \tau$ $\mathbf{val} \ lab : \tau$ $\mathbf{module} \ lab :_{\mu} \sigma$ $\mathbf{functor} \ lab :_{\mu} \sigma$ | |
| structure bindings | $sbnds$ | ::= | \cdot $sbnd, sbnds$ | |
| structure binding | $sbnd$ | ::= | $\mathbf{type} \ t = \tau$ $\mathbf{let} \ x : \tau = \delta$ $\mathbf{module} \ m = \mu$ $\mathbf{module} \ m :_{\mu} s = \mu$ | |

| functor $m;_{\mu}s = \mu$

contexts

$\Delta, ?; \Gamma, x::\tau; \Phi, t::\kappa; \Xi, m;_{\mu}\sigma; \Psi, s::_{\sigma}\xi$

statics

scratch

$\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$ κ_1 is a consistent subkind of κ_2

KCSubsumption

$\frac{test}{test}$

$\frac{test}{test}$

$\Delta; \Phi; \Xi; \Psi \vdash \xi_1 \lesssim_{\xi} \xi_2$ ξ_1 is a consistent sub signature kind of ξ_2

nameMe

$$\frac{\exists sdec_x \in sdec_{s_1} \text{ st } \Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\{sdec_x\}) \lesssim_{\xi} \text{SSigKind}(\{sdec_2\}) \quad \Delta; \Phi, \text{type}(\Delta; \Phi; \Xi; \Psi, sdec_2); \Xi, \text{submodule}(sdec_2); \Psi \vdash \{sdec_{s_1}\} \lesssim_{\xi} \{sdec_{s_2}\}}{\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\{sdec_{11}, sdec_{12}, sdec_{13} \text{ as } sdec_{s_1}\}) \lesssim_{\xi} \text{SSigKind}(\{sdec_2, sdec_{s_2}\})}$$

single

$$\frac{\Delta; \Phi; \Xi; \Psi \vdash sdec_1 \leq_{sdec} sdec_2}{\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\{sdec_1\}) \lesssim_{\xi} \text{SSigKind}(\{sdec_2\})}$$

nil

$$\frac{}{\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\{sdec_{s_1}\}) \lesssim_{\xi} \text{SSigKind}(\{sdec_{s_2}\})}$$

varprop

$$\frac{s::_{\sigma}\xi \in \Psi}{\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(s) \lesssim_{\xi} \xi}$$

nameMe?delete?

$$\frac{\sigma_1 \neq s \quad \Delta; \Phi; \Xi; \Psi \vdash \sigma_1 \Leftarrow \text{SSigKind}(\sigma_2)}{\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\sigma_1) \lesssim_{\xi} \text{SSigKind}(\sigma_2)}$$

funct

$$\frac{\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\sigma_{21}) \lesssim_{\xi} \text{SSigKind}(\sigma_{11}) \quad \Delta; \Phi; \Xi; m;_{\mu}\sigma_{11}; \Psi \vdash \text{SSigKind}(\sigma_{12}) \lesssim_{\xi} \text{SSigKind}(\sigma_{22})}{\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\Pi_{m;_{\mu}\sigma_{11}}.\sigma_{12}) \lesssim_{\xi} \text{SSigKind}(\Pi_{m;_{\mu}\sigma_{21}}.\sigma_{22})}$$

holes

$$\frac{\Delta; \Phi; \Xi; \Psi \vdash \emptyset^u \Leftarrow \xi}{\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\emptyset^u) \lesssim_{\xi} \xi}$$

neholes

$$\frac{\Delta; \Phi; \Xi; \Psi \vdash (\emptyset^u) \Leftarrow \xi}{\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}((\emptyset^u)) \lesssim_{\xi} \xi}$$

CSubSigKindHoleL

$$\frac{}{\Delta; \Phi; \Xi; \Psi \vdash \text{SigKHole} \lesssim_{\xi} \xi}$$

CSubSigKindHoleR

$$\frac{}{\Delta; \Phi; \Xi; \Psi \vdash \xi \lesssim_{\xi} \text{SigKHole}}$$

$\Delta; \Phi; \Xi; \Psi \vdash \sigma \Rightarrow \xi$ σ synthesizes signature kind ξ

SynSigKindVar

$$\frac{s::_{\sigma}\xi \in \Psi}{\Delta; \Phi; \Xi; \Psi \vdash s \Rightarrow \text{SSigKind}(s)}$$

SynSigKindVarFail

$$\frac{s \notin \text{dom}(\Psi)}{\Delta; \Phi; \Xi; \Psi \vdash s \Rightarrow \text{SigKHole}}$$

$\{sdec_{s_1}\} \text{ wellformed?}$

$$\vdash \{sdec_{s_1}\} \Rightarrow \text{SSigKind}(\{sdec_{s_1}\})$$

SynSigKindSigHole

$$\frac{u::_{\sigma}\xi \in \Delta}{\Delta; \Phi; \Xi; \Psi \vdash \emptyset^u \Rightarrow \xi}$$

SynSigKindSigHole

$$\frac{u::_{\sigma}\xi \in \Delta \quad \Delta; \Phi; \Xi; \Psi \vdash s \Rightarrow \xi_1}{\Delta; \Phi; \Xi; \Psi \vdash (\emptyset^u) \Rightarrow \xi}$$

$\Delta; \Phi; \Xi; \Psi \vdash \sigma \Leftarrow \xi$ σ analyzes against signature kind ξ

Sub

$$\frac{\Delta; \Phi; \Xi; \Psi \vdash \sigma \Rightarrow \xi_1 \quad \Delta; \Phi; \Xi; \Psi \vdash \xi_1 \lesssim_{\xi} \xi}{\Delta; \Phi; \Xi; \Psi \vdash \sigma \Leftarrow \xi}$$

$\Delta; \Phi; \Xi; \Psi \vdash sdec_1 \leq_{sdec} sdec_2$ $sdec_1$ is a subsdec of $sdec_2$

singleType

$\Delta; \Phi; \Xi; \Psi \vdash \text{type } lab = \tau \leq_{sdec} \text{type } lab$

singleType2

$\Delta; \Phi; \Xi; \Psi \vdash \tau_1 \equiv \tau_2$
 $\Delta; \Phi; \Xi; \Psi \vdash \text{type } lab = \tau_1 \leq_{sdec} \text{type } lab = \tau_2$

singleType3

$\Delta; \Phi; \Xi; \Psi \vdash \text{type } lab \leq_{sdec} \text{type } lab$

singleVa

$\Delta; \Phi; \Xi; \Psi \vdash \tau_1 \equiv \tau_2$
 $\Delta; \Phi; \Xi; \Psi \vdash \text{val } lab:\tau_1 \leq_{sdec} \text{val } lab:\tau_2$

singleMod

$\Delta; \Phi; \Xi; \Psi \vdash \sigma_1 \Leftarrow \text{SSigKind}(\sigma_2)$
 $\Delta; \Phi; \Xi; \Psi \vdash \text{module } lab:\mu\sigma_1 \leq_{sdec} \text{module } lab:\mu\sigma_2$

elab

$\Gamma; \Phi; \Xi \vdash \hat{\delta} \Rightarrow \tau \rightsquigarrow \delta \dashv \Delta$ $\hat{\delta}$ synthesizes type τ and elaborates to δ with hole context Δ

...

SynElabLetMod
 $\Gamma; \Phi; \Xi \vdash \hat{\mu} \Rightarrow \sigma \rightsquigarrow \mu \dashv \Delta_1 \quad \Gamma; \Phi; \Xi, m:\mu\sigma \vdash \hat{\delta} \Rightarrow \tau \rightsquigarrow \delta \dashv \Delta_2$
 $\Gamma; \Phi; \Xi \vdash \text{module } m = \hat{\mu} \text{ in } \hat{\delta} \Rightarrow \tau \rightsquigarrow \text{module } m = \mu \text{ in } \delta \dashv \Delta_1 \cup \Delta_2$

SynElabLetModAnn

$\Phi; \Xi \vdash \hat{\sigma} \Rightarrow \xi \rightsquigarrow \sigma \dashv \Delta_1 \quad \Gamma; \Phi; \Xi \vdash \hat{\mu} \Leftarrow \sigma \rightsquigarrow \mu \dashv \Delta_2 \quad \Gamma; \Phi; \Xi, m:\mu\sigma \vdash \hat{\delta} \Rightarrow \tau \rightsquigarrow \delta \dashv \Delta_3$
 $\Gamma; \Phi; \Xi \vdash \text{module } m:\mu\hat{\sigma} = \hat{\mu} \text{ in } \hat{\delta} \Rightarrow \tau \rightsquigarrow \text{module } m:\mu\sigma = \mu \text{ in } \delta \dashv \Delta_1 \cup \Delta_2 \cup \Delta_3$

SynElabModTermPrj

$\Gamma; \Phi; \Xi \vdash \hat{\mu} \Rightarrow \sigma \rightsquigarrow \mu \dashv \Delta \quad \Phi; \Xi \vdash \sigma \Rightarrow \xi$
 $\Gamma; \Phi; \Xi \vdash \hat{\mu}.lab \Rightarrow \tau \rightsquigarrow \mu.lab \dashv \Delta$

$\Phi; \Xi \vdash \hat{\tau} \Rightarrow \kappa \rightsquigarrow \tau \dashv \Delta$ $\hat{\tau}$ synthesizes kind κ and elaborates to τ with hole context Δ

...

SynElabModTypPrj
 $\Phi; \Xi \vdash m \Rightarrow \sigma \rightsquigarrow m \dashv \Delta \quad \text{something}\sigma\kappa$
 $\Phi; \Xi \vdash m.lab \Rightarrow \kappa \rightsquigarrow m.lab \dashv \Delta$

$\Phi; \Xi \vdash \hat{\tau} \Leftarrow \kappa \rightsquigarrow \tau \dashv \Delta$ $\hat{\tau}$ analyzes against kind κ and elaborates to τ with hole context Δ

$\Gamma; \Phi; \Xi \vdash \hat{\mu} \Rightarrow \sigma \rightsquigarrow \mu \dashv \Delta$ $\hat{\mu}$ synthesizes signature σ and elaborates to μ with hole context Δ

SynElabModVar

$m:\mu\sigma \in \Xi$
 $\Gamma; \Phi; \Xi \vdash m \Rightarrow \sigma \rightsquigarrow m \dashv \cdot$

SynElabModVarFail

$m \notin \text{dom}(\Xi)$
 $\Gamma; \Phi; \Xi \vdash m \Rightarrow () \rightsquigarrow (m)^u \dashv u:\mu()$

SynElabConsStruct

$\Gamma; \Phi; \Xi \vdash sbnd \Rightarrow sdec \rightsquigarrow sbnd \dashv \Delta_1$
 $\Gamma, \text{val}(sdec); \Phi, \text{type}(\Delta_1; \Phi; \Xi; \Psi, sdec); \Xi, \text{submodule}(sdec) \vdash \{sbnds\} \Rightarrow \{sdec\} \rightsquigarrow \{sbnds\} \dashv \Delta_2$
 $\Gamma; \Phi; \Xi \vdash \{sbnd, sbnds\} \Rightarrow \{sdec, sdec\} \rightsquigarrow \{sbnd, sbnds\} \dashv \Delta_1 \cup \Delta_2$

SynElabNilStruct

$\Gamma; \Phi; \Xi \vdash \{\cdot\} \Rightarrow \{\cdot\} \rightsquigarrow \{\cdot\} \dashv \cdot$
 $\Gamma; \Phi; \Xi \vdash \hat{\mu} \Leftarrow \sigma \rightsquigarrow \mu \dashv \Delta$ $\hat{\mu}$ analyzes against signature σ and elaborates to μ with hole context Δ

SynElabEmptyModHole

$\Gamma; \Phi; \Xi \vdash ()^u \Rightarrow () \rightsquigarrow ()^u \dashv u:\mu()$

SynElabNonemptyModHole

$\Gamma; \Phi; \Xi \vdash (m)^u \Rightarrow () \rightsquigarrow (m)^u \dashv u:\mu()$

functor stuff

AnaElabModSubsumption

$\Gamma; \Phi; \Xi \vdash \hat{\mu} \Rightarrow \sigma \rightsquigarrow \mu \dashv \Delta$
 $\Gamma; \Phi; \Xi \vdash \hat{\mu} \Leftarrow \sigma \rightsquigarrow \mu \dashv \Delta$

$\boxed{\Gamma; \Phi; \Xi \vdash \hat{sbnd} \Rightarrow sdec \rightsquigarrow sbnd \dashv \Delta}$ \hat{sbnd} synthesizes declaration $sdec$ and elaborates to $sbnd$ with hole context Δ

SynElabTypeSbnd

$$\frac{\Phi; \Xi \vdash \hat{\tau} \Rightarrow \kappa \rightsquigarrow \tau \dashv \Delta}{\Gamma; \Phi; \Xi \vdash \text{type } t = \hat{\tau} \Rightarrow \text{type } t = \tau \rightsquigarrow \text{type } t = \tau \dashv \Delta}$$

SynElabValSbnd

$$\frac{\Phi; \Xi \vdash \hat{\tau} \Rightarrow \kappa \rightsquigarrow \tau \dashv \Delta_1 \quad \Gamma; \Phi; \Xi \vdash \hat{\delta} \Leftarrow \tau \rightsquigarrow \delta \dashv \Delta_2}{\Gamma; \Phi; \Xi \vdash \text{let } x:\hat{\tau} = \hat{\delta} \Rightarrow \text{val } x:\tau \rightsquigarrow \text{let } x:\tau = \delta \dashv \Delta_1 \cup \Delta_2}$$

SynElabModSbnd

$$\frac{\Gamma; \Phi; \Xi \vdash \hat{\mu} \Rightarrow \sigma \rightsquigarrow \mu \dashv \Delta}{\Gamma; \Phi; \Xi \vdash \text{module } m = \hat{\mu} \Rightarrow \text{module } m:_{\mu}\sigma \rightsquigarrow \text{module } m:_{\mu}\sigma = \mu \dashv \Delta}$$

SynElabModAnnSbnd

$$\frac{\Phi; \Xi \vdash \hat{\sigma} \Rightarrow \xi \rightsquigarrow \sigma_1 \dashv \Delta_1 \quad \Gamma; \Phi; \Xi \vdash \hat{\mu} \Rightarrow \sigma_2 \rightsquigarrow \mu \dashv \Delta_2 \quad \Phi; \Xi; \Psi \vdash \sigma_2 \Leftarrow \xi}{\Gamma; \Phi; \Xi \vdash \text{module } m:_{\mu}\hat{\sigma} = \hat{\mu} \Rightarrow \text{module } m:_{\mu}\sigma_1 \rightsquigarrow \text{module } m:_{\mu}\sigma_1 = \mu \dashv \Delta_1 \cup \Delta_2}$$

$\boxed{\Gamma; \Phi; \Xi \vdash \hat{sbnd} \Leftarrow sdec \rightsquigarrow sbnd \dashv \Delta}$ \hat{sbnd} analyzes against declaration $sdec$ and elaborates to $sbnd$ with hole context Δ

subsump

$$\frac{\Gamma; \Phi; \Xi; l\Psi \vdash \hat{sbnd} \Rightarrow sdec_1 \rightsquigarrow sbnd \dashv \Delta \quad \Delta; \Phi; \Xi; \Psi \vdash sdec_1 \leq_{sdec} sdec}{\Gamma; \Phi; \Xi; \Psi \vdash \hat{sbnd} \Leftarrow sdec \rightsquigarrow sbnd \dashv \Delta}$$

$\boxed{\Gamma; \Phi; \Xi; \Psi \vdash \hat{sdec} \rightsquigarrow sdec \dashv \Delta}$ \hat{sdec} elaborates to $sdec$ with hole context Δ

opq

$$\frac{}{\Gamma; \Phi; \Xi; \Psi \vdash \text{type } lab \rightsquigarrow \text{type } lab \dashv \cdot}$$

trn

$$\frac{\Gamma; \Phi; \Xi; \Psi \vdash \hat{\tau} \Rightarrow \kappa \rightsquigarrow \tau \dashv \Delta}{\Gamma; \Phi; \Xi; \Psi \vdash \text{type } lab = \hat{\tau} \rightsquigarrow \text{type } lab = \tau \dashv \Delta}$$

val

$$\frac{\Gamma; \Phi; \Xi; \Psi \vdash \hat{\tau} \Rightarrow \kappa \rightsquigarrow \tau \dashv \Delta}{\Gamma; \Phi; \Xi; \Psi \vdash \text{val } lab:\hat{\tau} \rightsquigarrow \text{val } lab:\tau \dashv \Delta}$$

mod

$$\frac{\Gamma; \Phi; \Xi; \Psi \vdash \hat{\sigma} \Rightarrow \xi \rightsquigarrow \sigma \dashv \Delta}{\Gamma; \Phi; \Xi; \Psi \vdash \text{module } lab:_{\mu}\hat{\sigma} \rightsquigarrow \text{module } lab:_{\mu}\sigma \dashv \Delta}$$

we're going to need HOFunctors so we don't need to preclude users from typing a functor into a module and vice versa

$\boxed{\Phi; \Xi; \Psi \vdash \hat{\sigma} \Rightarrow \xi \rightsquigarrow \sigma \dashv \Delta}$ $\hat{\sigma}$ synthesizes signature kind ξ and elaborates to σ with hole context Δ

SynSigEmptyHole

$$\frac{}{\Phi; \Xi; \Psi \vdash \emptyset^u \Rightarrow \text{SigKHole} \rightsquigarrow \emptyset^u \dashv u::_{\sigma}\text{SigKHole}}$$

SynSigNonEmptyHole

$\boxed{\Phi; \Xi \vdash \hat{\sigma} \Leftarrow \xi \rightsquigarrow \sigma \dashv \Delta}$ $\hat{\sigma}$ analyzes against signature kind ξ and elaborates to σ with hole context Δ

misc

$$\begin{aligned} \text{val}(sdec) &= \begin{cases} lab:\tau & sdec = \text{val } lab:\tau \\ \cdot & \text{otherwise} \end{cases} \\ \text{type}(cntxts, sdec) &= \begin{cases} lab::\text{Type} & sdec = \text{type } lab \\ lab::\kappa & sdec = \text{type } lab = \tau \\ & \text{where } cntxts \vdash \tau \Rightarrow \kappa \\ \cdot & \text{otherwise} \end{cases} \\ \text{submodule}(sdec) &= \begin{cases} lab:_{\mu}\sigma & sdec = \text{module } lab:_{\mu}\sigma \\ \cdot & \text{otherwise} \end{cases} \end{aligned}$$