

Hazel Phi: 9-type-aliases

July 13, 2021

SYNTAX

BinOp	\oplus	$::=$	$\times \mid + \mid \rightarrow$
Kind	κ	$::=$	Type \mid KHole \mid $\mathbf{S}_{\kappa}(\tau) \mid \Pi_{t::\kappa_1}.\kappa_2$
Base Types	bse	$::=$	Int \mid Float \mid Bool
User Types	$\hat{\tau}$	$::=$	$t \mid \mathbf{bse} \mid \hat{\tau}_1 \oplus \hat{\tau}_2 \mid \langle \rangle^u \mid \langle \hat{\tau} \rangle^u \mid \lambda t::\mathbf{Type}.\hat{\tau} \mid \hat{\tau}_1 \hat{\tau}_2$
Internal Types	τ	$::=$	$t \mid \mathbf{bse} \mid \tau_1 \oplus \tau_2 \mid \langle \rangle^u \mid \langle \tau \rangle^u \mid \lambda t::\kappa.\tau \mid \tau_1 \tau_2$
Type Pattern			
User Expression			
Internal Expression			

DECLARATIVES

$\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$ κ_1 is a consistent subkind of κ_2

$$\begin{array}{c}
 \overline{\Delta; \Phi \vdash \mathbf{KHole} \lesssim \kappa} \qquad \overline{\Delta; \Phi \vdash \kappa \lesssim \mathbf{KHole}} \\
 \\
 \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \quad \frac{\Delta; \Phi \vdash \tau::\kappa}{\Delta; \Phi \vdash \mathbf{S}_{\kappa}(\tau) \lesssim \kappa} \quad \frac{\Delta; \Phi \vdash \kappa_3 \lesssim \kappa_1 \quad \Delta; \Phi, t::\kappa_3 \vdash \kappa_2 \lesssim \kappa_4}{\Delta; \Phi \vdash \Pi_{t::\kappa_1}.\kappa_2 \lesssim \Pi_{t::\kappa_3}.\kappa_4}
 \end{array}$$

$\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2$ κ_1 is equivalent to κ_2

$$\begin{array}{c}
 \overline{\Delta; \Phi \vdash \kappa \equiv \kappa} \quad \frac{\Delta; \Phi \vdash \kappa_2 \equiv \kappa_1}{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2} \quad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_3 \quad \Delta; \Phi \vdash \kappa_3 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2} \\
 \\
 \frac{\Delta; \Phi \vdash \tau_1 \overset{\kappa}{\equiv} \tau_2}{\Delta; \Phi \vdash \mathbf{S}_{\kappa}(\tau_1) \equiv \mathbf{S}_{\kappa}(\tau_2)} \quad \frac{\Delta; \Phi \vdash \tau::\mathbf{S}_{\kappa}(\tau_1)}{\Delta; \Phi \vdash \mathbf{S}_{\mathbf{S}_{\kappa}(\tau_1)}(\tau) \equiv \mathbf{S}_{\kappa}(\tau_1)} \quad \frac{\Delta; \Phi \vdash \tau::\Pi_{t::\kappa_1}.\kappa_2}{\Delta; \Phi \vdash \mathbf{S}_{\Pi_{t::\kappa_1}.\kappa_2}(\tau) \equiv \Pi_{t::\kappa_1}.\mathbf{S}_{\kappa_2}(\tau \ t)} \\
 \\
 \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \quad \Delta; \Phi, t::\kappa_1 \vdash \kappa_3 \equiv \kappa_4}{\Delta; \Phi \vdash \Pi_{t::\kappa_1}.\kappa_2 \equiv \Pi_{t::\kappa_3}.\kappa_4}
 \end{array}$$

$\boxed{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2}$ τ_1 is equivalent to τ_2 at kind κ

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \tau \equiv^{\kappa} \tau} \quad \frac{\Delta; \Phi \vdash \tau_2 \equiv^{\kappa} \tau_1}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} \quad \frac{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_3 \quad \Delta; \Phi \vdash \tau_3 \equiv^{\kappa} \tau_1}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} \\
\\
\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\text{Type}} \tau_3 \quad \Delta; \Phi \vdash \tau_2 \equiv^{\text{Type}} \tau_4}{\Delta; \Phi \vdash \tau_1 \oplus \tau_2 \equiv^{\text{Type}} \tau_3 \oplus \tau_4} \quad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \quad \Delta; \Phi, t :: \kappa_1 \vdash \tau_1 \equiv^{\kappa} \tau_2}{\Delta; \Phi \vdash \lambda t :: \kappa_1. \tau_1 \equiv^{\Pi t :: \kappa_1. \kappa} \lambda t :: \kappa_2. \tau_2} \\
\\
\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\Pi t :: \kappa_1. \kappa_2} \tau_3 \quad \Delta; \Phi \vdash \tau_2 \equiv^{\kappa_1} \tau_4}{\Delta; \Phi \vdash \tau_1 \tau_2 \equiv^{[\tau_1/t] \kappa_2} \tau_3 \tau_4} \\
\\
\frac{\Delta; \Phi \vdash \tau_1 :: \Pi t :: \kappa_1. \kappa_3 \quad \Delta; \Phi \vdash \tau_2 :: \Pi t :: \kappa_1. \kappa_4 \quad \Delta; \Phi, t :: \kappa_1 \vdash \tau_1 t \equiv^{\kappa_2} \tau_2 t}{\Delta; \Phi \vdash \tau_1 \equiv^{\Pi t :: \kappa_1. \kappa_2} \tau_2} \\
\\
\frac{\Delta; \Phi \vdash \tau_1 :: \mathbf{S}_{\kappa}(\tau_2) \quad \Delta; \Phi \vdash \tau_2 :: \kappa}{\Delta; \Phi \vdash \tau_1 \equiv^{\mathbf{S}_{\kappa}(\tau_2)} \tau_2} \\
\\
\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa_1} \tau_2 \quad \Delta; \Phi \vdash \kappa_1 \equiv \kappa}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2}
\end{array}$$

$\boxed{\Delta; \Phi \vdash \tau_1 \equiv \tau_2}$ τ_1 is equivalent to τ_2 at “top” kind

$$\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2}{\Delta; \Phi \vdash \tau_1 \equiv \tau_2} \quad \frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \tau \equiv \tau} \quad \frac{\Delta; \Phi \vdash \tau_2 \equiv \tau_1}{\Delta; \Phi \vdash \tau_1 \equiv \tau_2} \quad \frac{\Delta; \Phi \vdash \tau_1 \equiv \tau_3 \quad \Delta; \Phi \vdash \tau_3 \equiv \tau_2}{\Delta; \Phi \vdash \tau_1 \equiv \tau_2}$$

$\boxed{\Delta; \Phi \vdash \tau :: > \kappa}$ τ has principal kind κ

$$\frac{}{\Delta; \Phi \vdash \mathbf{bse} :: > \mathbf{S}_{\text{Type}}(\mathbf{bse})} \quad \frac{t :: \kappa \in \Phi}{\Delta; \Phi \vdash t :: > \mathbf{S}_{\kappa}(t)}$$

$\boxed{\Delta; \Phi \vdash \kappa \blacktriangleright \Pi t :: \kappa_1. \kappa_2}$ κ has matched Π -kind $\Pi t :: \kappa_1. \kappa_2$

$$\frac{}{\Delta; \Phi \vdash \mathbf{KHole} \blacktriangleright \Pi t :: \mathbf{KHole}. \mathbf{KHole}} \quad \frac{\Delta; \Phi \vdash \kappa \equiv \Pi t :: \kappa_1. \kappa_2}{\Delta; \Phi \vdash \kappa \blacktriangleright \Pi t :: \kappa_1. \kappa_2}$$

$\boxed{\Delta; \Phi \vdash \tau :: \kappa}$ τ is well formed at kind κ