

# Hazel Phi: 11-type-constructors

July 29, 2021

## SYNTAX

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Kind	$\kappa$	$::=$	<b>Type</b>   <b>KHole</b>   $S_{\kappa}(\tau)$   $\Pi_{t::\kappa_1}.\kappa_2$
User Types	$\hat{\tau}$	$::=$	$t$   <b>bse</b>   $\tau_1 \oplus \tau_2$   $\langle \rangle^u$   $\langle \hat{\tau} \rangle^u$   $\lambda t::\text{Type}.\hat{\tau}$   $\tau_1 \tau_2$
Internal Types	$\tau$	$::=$	$t$   <b>bse</b>   $\tau_1 \oplus \tau_2$   $\langle \rangle^u$   $\langle \tau \rangle^u$   $\langle t \rangle^u$   $\lambda t::\kappa.\tau$   $\tau_1 \tau_2$
Base Types	<b>bse</b>	$::=$	<b>Int</b>   <b>Float</b>   <b>Bool</b>
BinOp	$\oplus$	$::=$	$\times$   $+$   $\rightarrow$
Type Pattern			
User Expression			
Internal Expression			

## DECLARATIVES

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$\Delta; \Phi \vdash \tau ::> \kappa$   $\tau$  has principal (well formed) kind  $\kappa$

$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{bse} ::> S_{\text{Type}}(\text{bse})} \text{PK-Base}$	$\frac{\Delta; \Phi_1, t::\kappa, \Phi_2 \vdash \text{OK}}{\Delta; \Phi \vdash t ::> S_{\kappa}(t)} \text{PK-Var}$
$\frac{\Delta; \Phi \vdash \tau_1 :: \text{Type} \quad \Delta; \Phi \vdash \tau_2 :: \text{Type}}{\Delta; \Phi \vdash \tau_1 \oplus \tau_2 ::> S_{\text{Type}}(\tau_1 \oplus \tau_2)} \text{PK-}\oplus$	$\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \langle \rangle^u ::> S_{\kappa}(\langle \rangle^u)} \text{PK-EHole}$
$\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \text{OK} \quad \Delta; \Phi \vdash \tau :: \kappa_1}{\Delta; \Phi \vdash \langle \tau \rangle^u ::> S_{\kappa}(\langle \tau \rangle^u)} \text{PK-NEHole}$	$\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \text{OK} \quad t \notin \Phi}{\Delta; \Phi \vdash \langle t \rangle^u ::> S_{\kappa}(\langle t \rangle^u)} \text{PK-Unbound}$
$\frac{\Delta; \Phi, t::\kappa_1 \vdash \tau ::> \kappa_2}{\Delta; \Phi \vdash \lambda t::\kappa_1.\tau ::> S_{\Pi_{t::\kappa_1}.\kappa_2}}(\lambda t::\kappa_1.\tau)} \text{PK-}\lambda$	
$\frac{\Delta; \Phi \vdash \tau_1 ::> \kappa \quad \Delta; \Phi \vdash \kappa \Pi \Pi_{t::\kappa_1}.\kappa_2 \quad \Delta; \Phi \vdash \tau_2 :: \kappa_1}{\Delta; \Phi \vdash \tau_1 \tau_2 ::> [\tau_2/t]\kappa_2} \text{PK-Ap}$	

$\Delta; \Phi \vdash \tau :: \kappa$   $\tau$  is well formed at kind  $\kappa$

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \tau :: > \mathbf{S}_{\kappa}(\tau)}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-1} \qquad \frac{\Delta; \Phi \vdash \tau :: > \kappa_1 \quad \Delta; \Phi \vdash \kappa_1 \lesssim \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Subsump} \\
\\
\frac{\Delta; \Phi \vdash \tau :: > \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Reit} \qquad \frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \tau :: \mathbf{S}_{\kappa}(\tau)} \text{WFaK-Self} \\
\frac{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_3} \cdot \kappa_4 \quad \Delta; \Phi \vdash \Pi_{t::\kappa_3} \cdot \kappa_4 \lesssim \Pi_{t::\kappa_1} \cdot \kappa_2}{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_1} \cdot \kappa_2} \text{WFaK-PCSKTrans} \\
\frac{\Delta; \Phi \vdash \tau :: \mathbf{S}_{\kappa}(\tau_1) \quad \Delta; \Phi \vdash \tau_1 :: \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Flatten}
\end{array}$$

$\Delta; \Phi \vdash \kappa \dashv \vdash \Pi_{t::\kappa_1} \cdot \kappa_2$   $\kappa$  has matched  $\Pi$ -kind  $\Pi_{t::\kappa_1} \cdot \kappa_2$

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \mathbf{KHole} \dashv \vdash \Pi_{t::\mathbf{KHole}} \cdot \mathbf{KHole}} \dashv \vdash \text{-KHole} \qquad \frac{\Delta; \Phi \vdash \kappa \equiv \mathbf{SKHole}(\tau)}{\Delta; \Phi \vdash \kappa \dashv \vdash \Pi_{t::\mathbf{SKHole}(\tau)} \cdot \mathbf{SKHole}(\tau \ t)} \dashv \vdash \text{-SKHole} \\
\\
\frac{\Delta; \Phi \vdash \kappa \equiv \Pi_{t::\kappa_1} \cdot \kappa_2}{\Delta; \Phi \vdash \kappa \dashv \vdash \Pi_{t::\kappa_1} \cdot \kappa_2} \dashv \vdash \text{-}\Pi
\end{array}$$

$\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2$   $\kappa_1$  is equivalent to  $\kappa_2$

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \equiv \kappa} \text{KEquiv-Ref1} \qquad \frac{\Delta; \Phi \vdash \kappa_2 \equiv \kappa_1}{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2} \text{KEquiv-Symm} \\
\\
\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_3 \quad \Delta; \Phi \vdash \kappa_3 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2} \text{KEquiv-Trans} \\
\\
\frac{\Delta; \Phi \vdash \tau :: \mathbf{S}_{\kappa}(\tau_1)}{\Delta; \Phi \vdash \mathbf{S}_{\kappa}(\tau_1)(\tau) \equiv \mathbf{S}_{\kappa}(\tau_1)} \text{KEquiv-SKind}_{\mathbf{SKind}} \qquad \frac{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_1} \cdot \kappa_2}{\Delta; \Phi \vdash \mathbf{S}_{\Pi_{t::\kappa_1} \cdot \kappa_2}(\tau) \equiv \Pi_{t::\kappa_1} \cdot \mathbf{S}_{[t_1/t]\kappa_2}(\tau \ t_1)} \text{KEquiv-SKind}_{\Pi} \\
\\
\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \quad \Delta; \Phi, t::\kappa_1 \vdash \kappa_3 \equiv \kappa_4}{\Delta; \Phi \vdash \Pi_{t::\kappa_1} \cdot \kappa_2 \equiv \Pi_{t::\kappa_3} \cdot \kappa_4} \text{KEquiv-}\Pi \qquad \frac{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa_1} \tau_2 \quad \Delta; \Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta; \Phi \vdash \mathbf{S}_{\kappa_1}(\tau_1) \equiv \mathbf{S}_{\kappa_2}(\tau_2)} \text{KEquiv-SKind}
\end{array}$$

$\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$   $\kappa_1$  is a consistent subkind of  $\kappa_2$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \mathbf{KHole} \lesssim \kappa} \text{CSK-KHoleL} \quad \frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \mathbf{KHole}} \text{CSK-KHoleR}$$

$$\frac{\Delta; \Phi \vdash \mathbf{SKHole}(\tau) \text{ OK} \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \mathbf{SKHole}(\tau) \lesssim \kappa} \text{CSK-SKind}_{\mathbf{KHoleL}}$$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK} \quad \Delta; \Phi \vdash \mathbf{SKHole}(\tau) \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \mathbf{SKHole}(\tau)} \text{CSK-SKind}_{\mathbf{KHoleR}}$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \text{CSK-KEquiv} \quad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_3 \quad \Delta; \Phi \vdash \kappa_3 \lesssim \kappa_4 \quad \Delta; \Phi \vdash \kappa_4 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \text{CSK-Normal}$$

$$\frac{\Delta; \Phi \vdash \mathbf{S}_\kappa(\tau) \text{ OK}}{\Delta; \Phi \vdash \mathbf{S}_\kappa(\tau) \lesssim \kappa} \text{CSK-SKind} \quad \frac{\Delta; \Phi \vdash \kappa_3 \lesssim \kappa_1 \quad \Delta; \Phi, t::\kappa_3 \vdash \kappa_2 \lesssim \kappa_4}{\Delta; \Phi \vdash \Pi_{t::\kappa_1}.\kappa_2 \lesssim \Pi_{t::\kappa_3}.\kappa_4} \text{CSK-}\Pi$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2 \quad \Delta; \Phi \vdash \tau_1 \equiv^{\kappa_1} \tau_2}{\Delta; \Phi \vdash \mathbf{S}_{\kappa_1}(\tau_1) \lesssim \mathbf{S}_{\kappa_2}(\tau_2)} \text{CSK-?}$$

$\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2$   $\tau_1$  is provably equivalent to  $\tau_2$  at kind  $\kappa$

$$\frac{\Delta; \Phi \vdash \tau::\kappa}{\Delta; \Phi \vdash \tau \equiv \tau} \text{EquivAK-Ref1} \quad \frac{\Delta; \Phi \vdash \tau_2 \equiv^{\kappa} \tau_1}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} \text{EquivAK-Symm}$$

$$\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_3 \quad \Delta; \Phi \vdash \tau_3 \equiv^{\kappa} \tau_1}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} \text{EquivAK-Trans}$$

$$\frac{\Delta; \Phi \vdash \tau_1 ::> \kappa_1 \quad \Delta; \Phi \vdash \kappa_1 \equiv \mathbf{S}_\kappa(\tau_2)}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} \text{EquivAK-SKind}$$

$$\frac{\Delta; \Phi \vdash \tau_1 :: \Pi_{t::\kappa_1}.\kappa_3 \quad \Delta; \Phi \vdash \tau_2 :: \Pi_{t::\kappa_1}.\kappa_4 \quad \Delta; \Phi, t::\kappa_1 \vdash \tau_1 \equiv^{\kappa_2} \tau_2 \quad t}{\Delta; \Phi \vdash \tau_1 \equiv^{\Pi_{t::\kappa_1}.\kappa_2} \tau_2} \text{EquivAK-}\Pi$$

$$\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\Pi_{t::\kappa_1}.\kappa_2} \tau_3 \quad \Delta; \Phi \vdash \tau_2 \equiv^{\kappa_1} \tau_4}{\Delta; \Phi \vdash \tau_1 \tau_2 \equiv^{[\tau_2/t]\kappa_2} \tau_3 \tau_4} \text{EquivAK-Ap}$$

$$\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\mathbf{S}_\kappa(\tau)} \tau_2}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} (1)$$

$$\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\text{Type}} \tau_3 \quad \Delta; \Phi \vdash \tau_2 \equiv^{\text{Type}} \tau_4}{\Delta; \Phi \vdash \tau_1 \oplus \tau_2 \equiv^{\text{Type}} \tau_3 \oplus \tau_4} (2)$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \quad \Delta; \Phi, t::\kappa_1 \vdash \tau_1 \equiv^{\kappa} \tau_2}{\Delta; \Phi \vdash \lambda t::\kappa_1.\tau_1 \equiv^{\Pi_{t::\kappa_1}.\kappa} \lambda t::\kappa_2.\tau_2} (3)$$

$$\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa_1} \tau_2 \quad \Delta; \Phi \vdash \kappa_1 \equiv \kappa}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} (4)$$

$\Delta; \Phi \vdash \kappa \text{ OK}$   $\kappa$  is well formed

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{Type OK}} \text{KWF-Type}$$

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{KHole OK}} \text{KWF-KHole}$$

$$\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \text{S}_\kappa(\tau) \text{ OK}} \text{KWF-SKind}$$

$$\frac{\Delta; \Phi, \underline{t :: \kappa_1} \vdash \kappa_2 \text{ OK}}{\Delta; \Phi \vdash \Pi_{t :: \kappa_1} \cdot \kappa_2 \text{ OK}} \text{KWF-}\Pi$$

$\Delta; \Phi \vdash \text{OK}$  Context is well formed

$$\frac{}{\cdot; \cdot \vdash \text{OK}} \text{CWF-Nil}$$

$$\frac{t \notin \Phi \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi, \underline{t :: \kappa} \vdash \text{OK}} \text{CWF-TypVar}$$

$$\frac{u \notin \Delta \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta, u :: \kappa; \Phi \vdash \text{OK}} \text{CWF-Hole}$$

## METATHEORY

**Lemma 1** (COK). *If  $\Delta; \Phi \vdash \mathcal{J}$ , then  $\Delta; \Phi \vdash \text{OK}$*

*Proof.* By simultaneous induction on derivations.

No interesting cases. □

**Lemma 2** (Exchange).

*If  $\Delta; \Phi_1, \underline{t_{L1} :: \kappa_{L1}}, \underline{t_{L2} :: \kappa_{L2}}, \Phi_2 \vdash \mathcal{J}$  and  $\Delta; \Phi_1, \underline{t_{L2} :: \kappa_{L2}}, \underline{t_{L1} :: \kappa_{L1}}, \Phi_2 \vdash \text{OK}$ , then  $\Delta; \Phi_1, \underline{t_{L2} :: \kappa_{L2}}, \underline{t_{L1} :: \kappa_{L1}}, \Phi_2 \vdash \mathcal{J}$*

*Proof.* By induction on derivations.

No interesting cases.

(Only rules with  $\Phi$  extended in the consequent are interesting, which is only CWF-TypVar, but when  $\mathcal{J}$  is CWF, Exchange is identity) □

**Corollary 3** (Marked-Exchange).

*If  $\Delta; \Phi, \underline{t_{L1} :: \kappa_{L1}}, \underline{t_{L2} :: \kappa_{L2}} \vdash \mathcal{J}$  and  $\Delta; \Phi, \underline{t_{L2} :: \kappa_{L2}}, \underline{t_{L1} :: \kappa_{L1}} \vdash \text{OK}$ , then  $\Delta; \Phi, \underline{t_{L2} :: \kappa_{L2}}, \underline{t_{L1} :: \kappa_{L1}} \vdash \mathcal{J}$*

**Lemma 4** (Weakening).

*If  $\Delta; \Phi \vdash \mathcal{J}$  and  $\Delta; \Phi, \underline{t_L :: \kappa_L} \vdash \text{OK}$ , then  $\Delta; \Phi, \underline{t_L :: \kappa_L} \vdash \mathcal{J}$*

**Lemma 5** (OK-PK). *If  $\Delta; \Phi \vdash \tau :: > \kappa$ , then  $\Delta; \Phi \vdash \kappa \text{ OK}$*

**Lemma 6** (OK-WFaK). *If  $\Delta; \Phi \vdash \tau :: \kappa$ , then  $\Delta; \Phi \vdash \kappa \text{ OK}$*

**Lemma 7** (OK-MatchPi). *If  $\Delta; \Phi \vdash \kappa \blacktriangleright \Pi_{t :: \kappa_1} \cdot \kappa_2$ , then  $\Delta; \Phi \vdash \kappa \text{ OK}$  and  $\Delta; \Phi \vdash \Pi_{t :: \kappa_1} \cdot \kappa_2 \text{ OK}$*

**Lemma 8** (OK-KEquiv). *If  $\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2$ , then  $\Delta; \Phi \vdash \kappa_1 \text{ OK}$  and  $\Delta; \Phi \vdash \kappa_2 \text{ OK}$*

**Lemma 9** (OK-CSK). *If  $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$ , then  $\Delta; \Phi \vdash \kappa_1 \text{ OK}$  and  $\Delta; \Phi \vdash \kappa_2 \text{ OK}$*

**Lemma 10** (OK-EquivAK). *If  $\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2$ , then  $\Delta; \Phi \vdash \tau_1 :: \kappa$  and  $\Delta; \Phi \vdash \tau_2 :: \kappa$  and  $\Delta; \Phi \vdash \kappa \text{ OK}$*

**Lemma 11** (OK-Substitution).

*If  $\Delta; \Phi \vdash \tau_L :: \kappa_{L1}$  and  $\Delta; \Phi, \underline{t_L :: \kappa_{L1}} \vdash \kappa_{L2} \text{ OK}$ , then  $\Delta; \Phi \vdash [\tau_L / t_L] \kappa_{L2} \text{ OK}$*

*(induction on  $\Delta; \Phi, \underline{t_L :: \kappa_{L1}} \vdash \kappa_{L2} \text{ OK}$ )*

**Lemma 12** (K-Substitution).

*If  $\Delta; \Phi \vdash \tau_{L1} :: \kappa_{L1}$  and  $\Delta; \Phi, \underline{t_L :: \kappa_{L1}} \vdash \tau_{L2} :: \kappa_{L2}$ , then  $\Delta; \Phi \vdash [\tau_{L1} / t_L] \tau_{L2} :: [\tau_{L1} / t_L] \kappa_{L2}$*

*(induction on  $\Delta; \Phi, \underline{t_L :: \kappa_{L1}} \vdash \tau_{L2} :: \kappa_{L2}$ )*

*Proof.* By simultaneous induction on derivations.

The interesting cases per lemma:

## Weakening

	$\frac{\Delta; \Phi, \underline{t::\kappa_1}, \underline{t_L::\kappa_L} \vdash \tau ::> \kappa_2}{\Delta; \Phi, \underline{t_L::\kappa_L}, \underline{t::\kappa_1} \vdash \tau ::> \kappa_2} \text{Weakening}$	$\frac{\Delta; \Phi, \underline{t_L::\kappa_L}, \underline{t::\kappa_1} \vdash \text{OK}}{\Delta; \Phi, \underline{t_L::\kappa_L}, \underline{t::\kappa_1} \vdash \tau ::> \kappa_2} \text{CWF-TypVar}$	
	$\frac{\Delta; \Phi, \underline{t_L::\kappa_L}, \underline{t::\kappa_1} \vdash \tau ::> \kappa_2}{\Delta; \Phi, \underline{t_L::\kappa_L} \vdash \lambda t::\kappa_1. \tau ::> \mathbf{S}_{\Pi_{t::\kappa_1}. \kappa_2}(\lambda t::\kappa_1. \tau)} \text{Marked-Exchange}$		
	$\frac{\Delta; \Phi, \underline{t_L::\kappa_L} \vdash \lambda t::\kappa_1. \tau ::> \mathbf{S}_{\Pi_{t::\kappa_1}. \kappa_2}(\lambda t::\kappa_1. \tau)}{\Delta; \Phi, \underline{t_L::\kappa_L} \vdash \lambda t::\kappa_1. \tau ::> \mathbf{S}_{\Pi_{t::\kappa_1}. \kappa_2}(\lambda t::\kappa_1. \tau)} \text{PK-}\lambda$		
Weakening.	$\Delta; \Phi \vdash \kappa_1 \text{ OK}$ $\Delta; \Phi, \underline{t_L::\kappa_L} \vdash \text{OK}$ $\Delta; \Phi, \underline{t_L::\kappa_L} \vdash \kappa_1 \text{ OK}$ $\Delta; \Phi, \underline{t_L::\kappa_L}, \underline{t::\kappa_1} \vdash \text{OK}$ $\Delta; \Phi, \underline{t::\kappa_1}, \underline{t_L::\kappa_L} \vdash \text{OK}$ $\Delta; \Phi, \underline{t::\kappa_1}, \underline{t_L::\kappa_L} \vdash \tau ::> \kappa_2$ $\Delta; \Phi, \underline{t_L::\kappa_L}, \underline{t::\kappa_1} \vdash \tau ::> \kappa_2$ $\Delta; \Phi, \underline{t_L::\kappa_L} \vdash \lambda t::\kappa_1. \tau ::> \mathbf{S}_{\Pi_{t::\kappa_1}. \kappa_2}(\lambda t::\kappa_1. \tau)$	by subderivation premiss by IH by Weakening on subderivation premiss by CWF-TypVar by ? by Weakening on premiss by Marked-Exchange by PK- $\lambda$	
OK-PK.	PK-Base $\Delta; \Phi \vdash \mathbf{bse}::\mathbf{S}_{\text{Type}}(\mathbf{bse})$ $\Delta; \Phi \vdash \mathbf{bse}::\mathbf{Type}$ $\Delta; \Phi \vdash \mathbf{S}_{\text{Type}}(\mathbf{bse}) \text{ OK}$ $\Delta; \Phi \vdash \text{OK}$	by (9) by (10) by (43) by premiss	
	PK-Ap	bad	
OK-WFaK.	(12) $\Delta; \Phi \vdash \tau_2::\kappa$ $\Delta; \Phi \vdash \mathbf{S}_{\kappa}(\tau_2) \text{ OK}$	by (10) by (43)	
OK-KEquiv.	(22) $\Delta; \Phi \vdash \tau t ::> \kappa$		
OK-Substitution.	(41) $\Delta; \Phi, \underline{t_L::\kappa_{L1}} \vdash \text{OK}$ $\Delta; \Phi \vdash \kappa_{L1} \text{ OK}$ $\Delta; \Phi \vdash \text{OK}$ $\Delta; \Phi \vdash [\tau_L/t_L] \mathbf{Type} \text{ OK}$ (43) $\Delta; \Phi, \underline{t_L::\kappa_{L1}} \vdash \tau::\kappa$ $\Delta; \Phi, \underline{t_L::\kappa_{L1}} \vdash \text{OK}$ $\Delta; \Phi \vdash \kappa_{L1} \text{ OK}$ $\Delta; \Phi \vdash \text{OK}$ $\Delta; \Phi \vdash [\tau_{L1}/t_L] \tau::[\tau_{L1}/t_L] \kappa$ $\Delta; \Phi \vdash [\tau_L/t_L] \mathbf{S}_{\kappa}(\tau) \text{ OK}$	premiss (41) by subderivation premiss (46) by OK-KWF by (41) and degenerate subst premiss (43) by OK-WFaK by subderivation premiss (46) by OK-KWF by K-Substitution on premiss by (43)	

□

**Lemma 13** (PK-Unicity). *If  $\Delta; \Phi \vdash \tau_L ::> \kappa_{L1}$  and  $\Delta; \Phi \vdash \tau_L ::> \kappa_{L2}$  then  $\kappa_{L1}$  is  $\kappa_{L2}$*

**Lemma 14.** *If  $\Delta; \Phi \vdash \tau ::> \kappa_1$  and  $\Delta; \Phi \vdash \tau::\kappa_2$ , then  $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$*

**Lemma 15.** *If  $\Delta; \Phi \vdash \kappa_1 \lesssim \mathbf{S}_{\kappa_2}(\tau)$ , then  $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$*