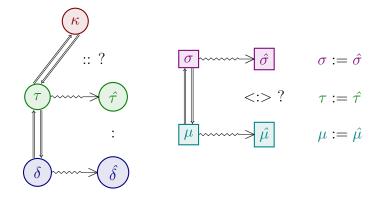
Hazel PHI: 10-modules

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how to read

```
\begin{array}{cccc} 800000 & kinds & 800080 & signatures \\ 008000 & types (constructors) & 008080 & modules \\ 000080 & terms & \end{array}
```

notes



syntax

```
\mathrm{HTyp} \ \hat{\tau} \ ::= \ t
                                                                                                  type variable
                                                                                                        base type
                        bse
                        \hat{	au_1} \oplus \hat{	au_2}
                                                                                                      type binop
                                                                                                          list type
                         \lambda t :: \kappa . \hat{\tau}
                                                                                                  type function
                        \hat{	au_1} \hat{	au_2}
                                                                                             type application
                        \{lab_1 \hookrightarrow \hat{\tau_1}, \dots \, lab_n \hookrightarrow \hat{\tau_n}\} labelled product type (record)
                         \hat{\mu}.lab
                                                                                  module type projection
                                                                                             empty type hole
                         (|\hat{\tau}|)
                                                                                       nonempty type hole
```

```
base type bse ::= Int
        HTyp BinOp \oplus ::= \times
external expression \delta ::= ...
                                                        \begin{array}{l} | \text{ signature } s = \sigma \text{ in } \delta \\ | \text{ module } m = \mu \text{ in } \delta \\ | \text{ module } m < :> s = \mu \text{ in } \delta \\ | \text{ functor something = something in } \delta \end{array}
                                                           \mu.lab
                                                                                                                                                           module term projection
internal expression \hat{\delta} ::= ...
                                                        \mid x \\ \mid \text{ signature } s = \hat{\sigma} \text{ in } \hat{\delta} \\ \mid \text{ module } m <:> s = \hat{\mu} \text{ in } \hat{\delta} 
                                                          | functor something = something in \hat{\delta}
                                                                                                                                                           module term projection
        signature \hat{\sigma} ::= s
                                                                                           signature variable
                                          \begin{array}{lll} := & s & \text{signature variable} \\ & \left\{ sdecs \right\} & \text{structure signature} \\ & \left\| \Pi_{m<:>\hat{\sigma_1}}.\hat{\sigma_2} \right\| & \text{functor signature} \\ & \left\| \left\| \right\| & \text{empty signature hole} \\ & \left\| \left( \hat{\sigma} \right) \right\| & \text{nonempty signature hole} \end{array}
      signature declarations sdecs ::= \cdot
                                                                             sdec, sdecs
        signature declaration sdec ::= type lab
                                                                        \begin{array}{c|c} & \text{type } lab = \hat{\tau} \\ & \text{val } lab : \hat{\tau} \\ & \text{module } lab <:> \hat{\sigma} \\ & \text{functor } lab <:> \hat{\sigma} \end{array} 
        structure bindings sbnds ::= \cdot
                                                                         sbnd, sbnds
```

```
structure binding \mathit{sbnd} ::= type t = \hat{\tau}
                                                                                                          \begin{array}{l} | \text{ let } x : \hat{\tau} = \hat{\delta} \\ | \text{ module } m = \hat{\mu} \\ | \text{ module } m <:> s = \hat{\mu} \\ | \text{ functor } m <:> s = \hat{\mu} \end{array}
```

contexts

$$\Gamma, x : \hat{\tau}; \Phi, t :: \kappa; \Xi, m <:> \hat{\sigma}; \Delta, ?$$

statics

 $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$ κ_1 is a consistent subkind of κ_2

KCSubsumption

testtest

elab

 δ synthesizes type τ and elaborates to $\hat{\delta}$ with hole context Δ $\Gamma; \Phi; \Xi \vdash \delta \Rightarrow \tau \leadsto \hat{\delta} \dashv \Delta$

SynElabLetMod

$$\frac{\Gamma;\Phi;\Xi\vdash\mu\ \Rightarrow\ \sigma\leadsto\hat{\mu}\dashv\Delta_1\qquad \Gamma;\Phi;\Xi,m<:>\sigma\vdash\delta\ \Rightarrow\ \tau\leadsto\hat{\delta}\dashv\Delta_2}{\Gamma;\Phi;\Xi\vdash \mathrm{module}\ m=\mu\ \mathrm{in}\ \delta\ \Rightarrow\ \tau\leadsto\mathrm{module}\ m=\hat{\mu}\ \mathrm{in}\ \hat{\delta}\dashv\Delta_1\cup\Delta_2}$$

 $\Gamma; \Phi; \Xi \vdash \mu \Rightarrow \sigma \leadsto \hat{\mu} \dashv \Delta$ μ synthesizes signature σ and elaborates to $\hat{\mu}$ with hole context Δ

$$\frac{m <:> \hat{\sigma} \in \Xi}{\Gamma; \Phi; \Xi \vdash m \Rightarrow \sigma \leadsto m \dashv \cdot}$$

SynElabModVarFail

$$\frac{m <:> \hat{\sigma} \in \Xi}{\Gamma; \Phi; \Xi \vdash m \ \Rightarrow \ \sigma \leadsto m \ \dashv \cdot} \qquad \qquad \frac{m \not \in \mathsf{dom}(\Xi)}{\Gamma; \Phi; \Xi \vdash m \ \Rightarrow \ (\!\!\!\! \|) \leadsto (\!\!\! \| m)\!\!\!\! \|^u \ \dashv u <:> (\!\!\!\! \|)\!\!\!\! \|}$$

SynElabStruct

$$\Gamma; \Phi; \Xi \vdash sbnd \Rightarrow sdec \leadsto sbnd \dashv \Delta_1$$

$$\frac{\Gamma, \mathsf{val}(sdec); \Phi, \mathsf{type}(sdec); \Xi, \mathsf{submodule}(sdec) \vdash sbnds \ \Rightarrow \ sdecs \leadsto sbnds \dashv \Delta_2}{\Gamma; \Phi; \Xi \vdash \{sbnd, sbnds\} \ \Rightarrow \ \{sdec, sdecs\} \leadsto \{sbnd, sbnds\} \dashv \Delta_1 \cup \Delta_2}$$