

Hazel PHI: 10-modules

June 29, 2021

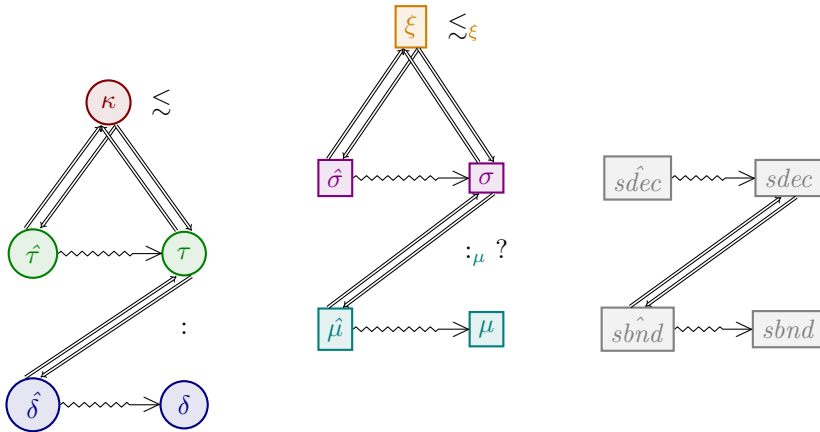
prerequisites

- Hazel PHI: 9-type-aliases-redux
 - github
 - current commit: 4410cd565ce717707e580e44f64868d3175fe2a6
- (optional) Hazel PHI: 1-labeled-tuples
 - github
 - current commit: 0a7d0b53ee7286d03ea3be13a7ac91a86f1c90b1

how to read

800000	kinds	D08000	signature kind
008000	types (constructors)	800080	signatures
000080	terms	008080	modules

notes



external typ/sig/mod syntax not written out yet (waiting for construction dust to settle); patterns not handled yet– will be left till end.

syntax

kind	κ	::=	Type	kind of types
			$S(\tau)$	singleton kind
			$KHole$	kind hole
			$\Pi_{t::\kappa_1}.\kappa_2$	dependent function kind

HTyp	τ	::=	t bse $\tau_1 \oplus \tau_2$ $[\tau]$ $\lambda t :: \kappa. \tau$ $\tau_1 \tau_2$ $\{lab_1 \hookrightarrow \tau_1, \dots lab_n \hookrightarrow \tau_n\}$ $\mu.lab$ $\langle \rangle$ $\langle \tau \rangle$	type variable base type type binop list type type function type application labelled product type (record) module type projection empty type hole nonempty type hole
base type	bse	::=	\mathbf{Int} \mathbf{Float} \mathbf{Bool}	
HTyp BinOp	\oplus	::=	\times $+$ \rightarrow	
external expression	$\hat{\delta}$::=	\dots x $\mathbf{signature} \ s = \hat{\sigma} \ \mathbf{in} \ \hat{\delta}$ $\mathbf{module} \ m = \hat{\mu} \ \mathbf{in} \ \hat{\delta}$ $\mathbf{module} \ m :_{\mu} s = \hat{\mu} \ \mathbf{in} \ \hat{\delta}$ $\mathbf{functor} \ \mathbf{something} = \mathbf{something} \ \mathbf{in} \ \hat{\delta}$ $\hat{\mu}.lab$	module term projection
internal expression	δ	::=	\dots x $\mathbf{signature} \ s = \sigma \ \mathbf{in} \ \delta$ $\mathbf{module} \ m :_{\mu} s = \mu \ \mathbf{in} \ \delta$ $\mathbf{functor} \ \mathbf{something} = \mathbf{something} \ \mathbf{in} \ \delta$ $\mu.lab$	module term projection
signature kind	ξ	::=	$\mathbf{SSigKind}(\sigma)$ $\mathbf{SigKHole}$	
signature	σ	::=	s $\{sdecs\}$ $\Pi_{m :_{\mu} \sigma_1}. \sigma_2$ $\langle \rangle$ $\langle s \rangle$	signature variable structure signature functor signature empty signature hole nonempty signature hole
module	μ	::=	m $\{sbnds\}$ $\lambda m :_{\mu} \sigma. \mu$ $\mu_1 \ \mu_2$ $\mu.lab$ $\langle \rangle$ $\langle \mu \rangle$	module variable structure functor functor application submodule projection empty module hole nonempty module hole
signature declarations	$sdecs$::=	\cdot $sdec, sdecs$	
signature declaration	$sdec$::=	$\mathbf{type} \ lab$ $\mathbf{type} \ lab = \tau$ $\mathbf{val} \ lab : \tau$ $\mathbf{module} \ lab :_{\mu} \sigma$ $\mathbf{functor} \ lab :_{\mu} \sigma$	
structure bindings	$sbnds$::=	\cdot $sbnd, sbnds$	
structure binding	$sbnd$::=	$\mathbf{type} \ t = \tau$ $\mathbf{let} \ x : \tau = \delta$ $\mathbf{module} \ m = \mu$ $\mathbf{module} \ m :_{\mu} s = \mu$	

| functor $m;_{\mu}s = \mu$

contexts

$\Delta, ?; \Gamma, x::\tau; \Phi, t::\kappa; \Xi, m;_{\mu}\sigma; \Psi, s::_{\sigma}\xi$

statics

scratch

$\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$ κ_1 is a consistent subkind of κ_2

KCSubsumption

$\frac{test}{test}$

$test$

$\Delta; \Phi; \Xi; \Psi \vdash \xi_1 \lesssim_{\xi} \xi_2$ ξ_1 is a consistent sub signature kind of ξ_2

nameMe

$\exists sdec_x \in sdec_{s_1} \text{ st } \Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\{sdec_x\}) \lesssim_{\xi} \text{SSigKind}(\{sdec_2\})$
 $\Delta; \Phi, \text{type}(sdec_2); \Xi, \text{submodule}(sdec_2); \Psi \vdash \{sdec_{s_1}\} \lesssim_{\xi} \{sdec_{s_2}\}$

$\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\{sdec_{11}, sdec_{12}, sdec_{13} \text{ as } sdec_{s_1}\}) \lesssim_{\xi} \text{SSigKind}(\{sdec_2, sdec_{s_2}\})$

singleType

$\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\{\text{type } lab = \tau\}) \lesssim_{\xi} \text{SSigKind}(\{\text{type } lab\})$

singleType2

$\Delta; \Phi; \Xi; \Psi \vdash \tau_1 \equiv \tau_2$

$\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\{\text{type } lab = \tau_1\}) \lesssim_{\xi} \text{SSigKind}(\{\text{type } lab = \tau_2\})$

singleType3

$\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\{\text{type } lab\}) \lesssim_{\xi} \text{SSigKind}(\{\text{type } lab\})$

singleVa

$\Delta; \Phi; \Xi; \Psi \vdash \tau_1 \equiv \tau_2$

$\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\{\text{val } lab:\tau_1\}) \lesssim_{\xi} \text{SSigKind}(\{\text{val } lab:\tau_2\})$

singleMod

$\Delta; \Phi; \Xi; \Psi \vdash \sigma_1 \Leftarrow \text{SSigKind}(\sigma_2)$

$\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\{\text{module } lab:_{\mu}\sigma_1\}) \lesssim_{\xi} \text{SSigKind}(\{\text{module } lab:_{\mu}\sigma_2\})$

nil

$\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\{sdec_{s_1}\}) \lesssim_{\xi} \text{SSigKind}(\{\cdot\})$

varprop

$s::_{\sigma}\xi \in \Psi$

$\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(s) \lesssim_{\xi} \xi$

nameMe

$\Delta; \Phi; \Xi; \Psi \vdash \sigma_1 \Leftarrow \text{SSigKind}(\sigma_2)$

$\Delta; \Phi; \Xi; \Psi \vdash \text{SSigKind}(\sigma_1) \lesssim_{\xi} \text{SSigKind}(\sigma_2)$

CSubSigKindHoleL

$\Delta; \Phi; \Xi; \Psi \vdash \text{SigKHole} \lesssim_{\xi} \xi$

CSubSigKindHoleR

$\Delta; \Phi; \Xi; \Psi \vdash \xi \lesssim_{\xi} \text{SigKHole}$

$\Delta; \Phi; \Xi; \Psi \vdash \sigma \Rightarrow \xi$ σ synthesizes signature kind ξ

SynSigKndVar

$s::_{\sigma}\xi \in \Psi$

$\Delta; \Phi; \Xi; \Psi \vdash s \Rightarrow \text{SSigKind}(s)$

SynSigKndVarFail

$s \notin \text{dom}(\Psi)$

$\Delta; \Phi; \Xi; \Psi \vdash s \Rightarrow \text{SigKHole}$

$\{sdec_{s_1}\} \text{ well formed?}$

$\vdash \{sdec_{s_1}\} \Rightarrow \text{SSigKind}(\{sdec_{s_1}\})$

SynSigKndSigHole

$u::_{\sigma}\xi \in \Delta$

$\Delta; \Phi; \Xi; \Psi \vdash \langle \langle s \rangle \rangle^u \Rightarrow \xi$

SynSigKndSigHole

$u::_{\sigma}\xi \in \Delta$ $\Delta; \Phi; \Xi; \Psi \vdash s \Rightarrow \xi_1$

$\Delta; \Phi; \Xi; \Psi \vdash \langle \langle s \rangle \rangle^u \Rightarrow \xi$

$\Delta; \Phi; \Xi; \Psi \vdash \sigma \leftarrow \xi$ σ analyzes against signature kind ξ

$$\frac{\text{Sub} \quad \Delta; \Phi; \Xi; \Psi \vdash \sigma \Rightarrow \xi_1 \quad \Delta; \Phi; \Xi; \Psi \vdash \xi_1 \lesssim_{\xi} \xi}{\Delta; \Phi; \Xi; \Psi \vdash \sigma \leftarrow \xi}$$

elab

$\Gamma; \Phi; \Xi \vdash \hat{\delta} \Rightarrow \tau \rightsquigarrow \delta \dashv \Delta$ $\hat{\delta}$ synthesizes type τ and elaborates to δ with hole context Δ

$$\begin{array}{c} \dots \\ \text{SynElabLetMod} \\ \frac{\Gamma; \Phi; \Xi \vdash \hat{\mu} \Rightarrow \sigma \rightsquigarrow \mu \dashv \Delta_1 \quad \Gamma; \Phi; \Xi, m;_{\mu} \sigma \vdash \hat{\delta} \Rightarrow \tau \rightsquigarrow \delta \dashv \Delta_2}{\Gamma; \Phi; \Xi \vdash \text{module } m = \hat{\mu} \text{ in } \hat{\delta} \Rightarrow \tau \rightsquigarrow \text{module } m = \mu \text{ in } \delta \dashv \Delta_1 \cup \Delta_2} \end{array}$$

$$\begin{array}{c} \text{SynElabLetModAnn} \\ \frac{\Phi; \Xi \vdash \hat{\sigma} \Rightarrow \xi \rightsquigarrow \sigma \dashv \Delta_1 \quad \Gamma; \Phi; \Xi \vdash \hat{\mu} \leftarrow \sigma \rightsquigarrow \mu \dashv \Delta_2 \quad \Gamma; \Phi; \Xi, m;_{\mu} \sigma \vdash \hat{\delta} \Rightarrow \tau \rightsquigarrow \delta \dashv \Delta_3}{\Gamma; \Phi; \Xi \vdash \text{module } m;_{\mu} \hat{\sigma} = \hat{\mu} \text{ in } \hat{\delta} \Rightarrow \tau \rightsquigarrow \text{module } m;_{\mu} \sigma = \mu \text{ in } \delta \dashv \Delta_1 \cup \Delta_2 \cup \Delta_3} \end{array}$$

$$\begin{array}{c} \text{SynElabModTermPrj} \\ \frac{\Gamma; \Phi; \Xi \vdash \hat{\mu} \Rightarrow \sigma \rightsquigarrow \mu \dashv \Delta \quad \Phi; \Xi \vdash \sigma \Rightarrow \xi}{\Gamma; \Phi; \Xi \vdash \hat{\mu}.lab \Rightarrow \tau \rightsquigarrow \mu.lab \dashv \Delta} \end{array}$$

$\Phi; \Xi \vdash \hat{\tau} \Rightarrow \kappa \rightsquigarrow \tau \dashv \Delta$ $\hat{\tau}$ synthesizes kind κ and elaborates to τ with hole context Δ

$$\begin{array}{c} \dots \\ \text{SynElabModTypPrj} \\ \frac{\Phi; \Xi \vdash m \Rightarrow \sigma \rightsquigarrow m \dashv \Delta \quad \text{something } \sigma \kappa}{\Phi; \Xi \vdash m.lab \Rightarrow \kappa \rightsquigarrow m.lab \dashv \Delta} \end{array}$$

$\Phi; \Xi \vdash \hat{\tau} \leftarrow \kappa \rightsquigarrow \tau \dashv \Delta$ $\hat{\tau}$ analyzes against kind κ and elaborates to τ with hole context Δ

$\Gamma; \Phi; \Xi \vdash \hat{\mu} \Rightarrow \sigma \rightsquigarrow \mu \dashv \Delta$ $\hat{\mu}$ synthesizes signature σ and elaborates to μ with hole context Δ

$$\begin{array}{c} \text{SynElabModVar} \\ \frac{m;_{\mu} \sigma \in \Xi}{\Gamma; \Phi; \Xi \vdash m \Rightarrow \sigma \rightsquigarrow m \dashv \cdot} \end{array}$$

$$\begin{array}{c} \text{SynElabModVarFail} \\ \frac{m \notin \text{dom}(\Xi)}{\Gamma; \Phi; \Xi \vdash m \Rightarrow \emptyset \rightsquigarrow (m)^u \dashv u;_{\mu} \emptyset} \end{array}$$

SynElabConsStruct

$$\begin{array}{c} \Gamma; \Phi; \Xi \vdash sbnd \Rightarrow sdec \rightsquigarrow sbnd \dashv \Delta_1 \\ \frac{\Gamma, \text{val}(sdec); \Phi, \text{type}(sdec); \Xi, \text{submodule}(sdec) \vdash \{sbnds\} \Rightarrow \{sdec\} \rightsquigarrow \{sbnds\} \dashv \Delta_2}{\Gamma; \Phi; \Xi \vdash \{sbnd, sbnds\} \Rightarrow \{sdec, sdec\} \rightsquigarrow \{sbnd, sbnds\} \dashv \Delta_1 \cup \Delta_2} \end{array}$$

SynElabNilStruct

$$\frac{}{\Gamma; \Phi; \Xi \vdash \{\cdot\} \Rightarrow \{\cdot\} \rightsquigarrow \{\cdot\} \dashv \cdot}$$

SynElabEmptyModHole

$$\frac{}{\Gamma; \Phi; \Xi \vdash \emptyset^u \Rightarrow \emptyset \rightsquigarrow \emptyset^u \dashv u;_{\mu} \emptyset}$$

SynElabNonemptyModHole

$$\frac{}{\Gamma; \Phi; \Xi \vdash (m)^u \Rightarrow \emptyset \rightsquigarrow (m)^u \dashv u;_{\mu} \emptyset}$$

functor stuff

$\Gamma; \Phi; \Xi \vdash \hat{\mu} \leftarrow \sigma \rightsquigarrow \mu \dashv \Delta$ $\hat{\mu}$ analyzes against signature σ and elaborates to μ with hole context Δ

AnaElabModSubsumption

$$\frac{\Gamma; \Phi; \Xi \vdash \hat{\mu} \Rightarrow \sigma \rightsquigarrow \mu \dashv \Delta}{\Gamma; \Phi; \Xi \vdash \hat{\mu} \leftarrow \sigma \rightsquigarrow \mu \dashv \Delta}$$

$\boxed{\Gamma; \Phi; \Xi \vdash \hat{sbnd} \Rightarrow sdec \rightsquigarrow sbnd \dashv \Delta}$ \hat{sbnd} synthesizes declaration $sdec$ and elaborates to $sbnd$ with hole context Δ

SynElabTypeSbnd

$$\frac{\Phi; \Xi \vdash \hat{\tau} \Rightarrow \kappa \rightsquigarrow \tau \dashv \Delta}{\Gamma; \Phi; \Xi \vdash \text{type } t = \hat{\tau} \Rightarrow \text{type } t = \tau \rightsquigarrow \text{type } t = \tau \dashv \Delta}$$

SynElabValSbnd

$$\frac{\Phi; \Xi \vdash \hat{\tau} \Rightarrow \kappa \rightsquigarrow \tau \dashv \Delta_1 \quad \Gamma; \Phi; \Xi \vdash \hat{\delta} \Leftarrow \tau \rightsquigarrow \delta \dashv \Delta_2}{\Gamma; \Phi; \Xi \vdash \text{let } x:\hat{\tau} = \hat{\delta} \Rightarrow \text{val } x:\tau \rightsquigarrow \text{let } x:\tau = \delta \dashv \Delta_1 \cup \Delta_2}$$

SynElabModSbnd

$$\frac{\Gamma; \Phi; \Xi \vdash \hat{\mu} \Rightarrow \sigma \rightsquigarrow \mu \dashv \Delta}{\Gamma; \Phi; \Xi \vdash \text{module } m = \hat{\mu} \Rightarrow \text{module } m:_{\mu}\sigma \rightsquigarrow \text{module } m:_{\mu}\sigma = \mu \dashv \Delta}$$

SynElabModAnnSbnd

$$\frac{\Phi; \Xi \vdash \hat{\sigma} \Rightarrow \xi \rightsquigarrow \sigma_1 \dashv \Delta_1 \quad \Gamma; \Phi; \Xi \vdash \hat{\mu} \Rightarrow \sigma_2 \rightsquigarrow \mu \dashv \Delta_2 \quad \Phi; \Xi; \Psi \vdash \sigma_2 \Leftarrow \xi}{\Gamma; \Phi; \Xi \vdash \text{module } m:_{\mu}\hat{\sigma} = \hat{\mu} \Rightarrow \text{module } m:_{\mu}\sigma_1 \rightsquigarrow \text{module } m:_{\mu}\sigma_1 = \mu \dashv \Delta_1 \cup \Delta_2}$$

$\boxed{\Gamma; \Phi; \Xi \vdash \hat{sbnd} \Leftarrow sdec \rightsquigarrow sbnd \dashv \Delta}$ \hat{sbnd} analyzes against declaration $sdec$ and elaborates to $sbnd$ with hole context Δ

$\boxed{\Gamma; \Phi; \Xi; \Psi \vdash \hat{sdec} \rightsquigarrow sdec \dashv \Delta}$ \hat{sdec} elaborates to $sdec$ with hole context Δ

opq

$$\frac{}{\Gamma; \Phi; \Xi; \Psi \vdash \text{type } lab \rightsquigarrow \text{type } lab \dashv \cdot}$$

trn

$$\frac{\Gamma; \Phi; \Xi; \Psi \vdash \hat{\tau} \Rightarrow \kappa \rightsquigarrow \tau \dashv \Delta}{\Gamma; \Phi; \Xi; \Psi \vdash \text{type } lab = \hat{\tau} \rightsquigarrow \text{type } lab = \tau \dashv \Delta}$$

val

$$\frac{\Gamma; \Phi; \Xi; \Psi \vdash \hat{\tau} \Rightarrow \kappa \rightsquigarrow \tau \dashv \Delta}{\Gamma; \Phi; \Xi; \Psi \vdash \text{val } lab:\hat{\tau} \rightsquigarrow \text{val } lab:\tau \dashv \Delta}$$

mod

$$\frac{\Gamma; \Phi; \Xi; \Psi \vdash \hat{\sigma} \Rightarrow \xi \rightsquigarrow \sigma \dashv \Delta}{\Gamma; \Phi; \Xi; \Psi \vdash \text{module } lab:_{\mu}\hat{\sigma} \rightsquigarrow \text{module } lab:_{\mu}\sigma \dashv \Delta}$$

we're going to need HOFunctors so we don't need to preclude users from typing a functor into a module and vice versa

$\boxed{\Phi; \Xi; \Psi \vdash \hat{\sigma} \Rightarrow \xi \rightsquigarrow \sigma \dashv \Delta}$ $\hat{\sigma}$ synthesizes signature kind ξ and elaborates to σ with hole context Δ

SynSigEmptyHole

$$\frac{}{\Phi; \Xi; \Psi \vdash \mathbb{0}^u \Rightarrow \text{SigKHole} \rightsquigarrow \mathbb{0}^u \dashv u::_{\sigma}\text{SigKHole}}$$

SynSigNonEmptyHole

$\boxed{\Phi; \Xi \vdash \hat{\sigma} \Leftarrow \xi \rightsquigarrow \sigma \dashv \Delta}$ $\hat{\sigma}$ analyzes against signature kind ξ and elaborates to σ with hole context Δ

misc

$$\text{val}(sdec) = \begin{cases} lab:\tau & sdec \equiv \text{val } lab:\tau \\ \cdot & \text{otherwise} \end{cases}$$

$$\text{type}(sdec) = \begin{cases} lab::\text{Type} & sdec \equiv \text{type } lab \\ lab::\kappa & sdec \equiv \text{type } lab = \tau \\ & \text{where } \vdash \tau \Rightarrow \kappa \\ \cdot & \text{otherwise} \end{cases}$$

$$\text{submodule}(sdec) = \begin{cases} lab:_{\mu}\sigma & sdec \equiv \text{module } lab:_{\mu}\sigma \\ \cdot & \text{otherwise} \end{cases}$$