

Hazel Phi: 11-type-constructors

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SYNTAX

Kind	κ	$::=$	Type KHole $\mathbf{S}_{\kappa}(\tau)$ $\Pi_{t::\kappa_1}.\kappa_2$
User Types	$\hat{\tau}$	$::=$	t bse $\tau_1 \oplus \tau_2$ $\langle \rangle^u$ $\langle \hat{\tau} \rangle^u$ $\lambda t::\mathbf{Type}.\hat{\tau}$ $\tau_1 \tau_2$
Internal Types	τ	$::=$	t bse $\tau_1 \oplus \tau_2$ $\langle \rangle^u$ $\langle \tau \rangle^u$ $\langle t \rangle^u$ $\lambda t::\kappa.\tau$ $\tau_1 \tau_2$
Base Types	bse	$::=$	Int Float Bool
BinOp	\oplus	$::=$	\times $+$ \rightarrow
Type Pattern			
User Expression			
Internal Expression			

DECLARATIVES

$\boxed{\Delta; \Phi \vdash \tau ::> \kappa}$ τ has principal (well formed) kind κ

$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \mathbf{bse} ::> \mathbf{S}_{\mathbf{Type}}(\mathbf{bse})} \text{PK-Base}$	$\frac{\Delta; \Phi_1, t::\kappa, \Phi_2 \vdash \text{OK}}{\Delta; \Phi \vdash t ::> \mathbf{S}_{\kappa}(t)} \text{PK-Var}$
$\frac{\Delta; \Phi \vdash \tau_1 :: \mathbf{Type} \quad \Delta; \Phi \vdash \tau_2 :: \mathbf{Type}}{\Delta; \Phi \vdash \tau_1 \oplus \tau_2 ::> \mathbf{S}_{\mathbf{Type}}(\tau_1 \oplus \tau_2)} \text{PK-}\oplus$	$\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \langle \rangle^u ::> \mathbf{S}_{\kappa}(\langle \rangle^u)} \text{PK-EHole}$
$\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \text{OK} \quad \Delta; \Phi \vdash \tau :: \kappa_1}{\Delta; \Phi \vdash \langle \tau \rangle^u ::> \mathbf{S}_{\kappa}(\langle \tau \rangle^u)} \text{PK-NEHole}$	$\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \text{OK} \quad t \notin \Phi}{\Delta; \Phi \vdash \langle t \rangle^u ::> \mathbf{S}_{\kappa}(\langle t \rangle^u)} \text{PK-Unbound}$
$\frac{\Delta; \Phi, t::\kappa_1 \vdash \tau ::> \kappa_2}{\Delta; \Phi \vdash \lambda t::\kappa_1.\tau ::> \mathbf{S}_{\Pi_{t::\kappa_1}.\kappa_2}(\lambda t::\kappa_1.\tau)} \text{PK-}\lambda$	
$\frac{\Delta; \Phi \vdash \tau_1 ::> \kappa \quad \Delta; \Phi \vdash \kappa \Pi_{t::\kappa_1}.\kappa_2 \quad \Delta; \Phi \vdash \tau_2 :: \kappa_1}{\Delta; \Phi \vdash \tau_1 \tau_2 ::> [\tau_2/t]\kappa_2} \text{PK-Ap}$	

$\Delta; \Phi \vdash \tau :: \kappa$ τ is well formed at kind κ

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \tau :: > \mathbf{S}_{\kappa}(\tau)}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-1} \qquad \frac{\Delta; \Phi \vdash \tau :: > \kappa_1 \quad \Delta; \Phi \vdash \kappa_1 \lesssim \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Subsump} \\
\\
\frac{\Delta; \Phi \vdash \tau :: > \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Reit} \qquad \frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \tau :: \mathbf{S}_{\kappa}(\tau)} \text{WFaK-Self} \\
\frac{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_3} \cdot \kappa_4 \quad \Delta; \Phi \vdash \Pi_{t::\kappa_3} \cdot \kappa_4 \lesssim \Pi_{t::\kappa_1} \cdot \kappa_2}{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_1} \cdot \kappa_2} \text{WFaK-PCSKTrans} \\
\frac{\Delta; \Phi \vdash \tau :: \mathbf{S}_{\kappa}(\tau_1) \quad \Delta; \Phi \vdash \tau_1 :: \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Flatten}
\end{array}$$

$\Delta; \Phi \vdash \kappa \dashv \vdash \Pi_{t::\kappa_1} \cdot \kappa_2$ κ has matched Π -kind $\Pi_{t::\kappa_1} \cdot \kappa_2$

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \mathbf{KHole} \dashv \vdash \Pi_{t::\mathbf{KHole}} \cdot \mathbf{KHole}} \dashv \vdash \text{-KHole} \qquad \frac{\Delta; \Phi \vdash \kappa \equiv \mathbf{SKHole}(\tau)}{\Delta; \Phi \vdash \kappa \dashv \vdash \Pi_{t::\mathbf{SKHole}(\tau)} \cdot \mathbf{SKHole}(\tau \ t)} \dashv \vdash \text{-SKHole} \\
\\
\frac{\Delta; \Phi \vdash \kappa \equiv \Pi_{t::\kappa_1} \cdot \kappa_2}{\Delta; \Phi \vdash \kappa \dashv \vdash \Pi_{t::\kappa_1} \cdot \kappa_2} \dashv \vdash \text{-}\Pi
\end{array}$$

$\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2$ κ_1 is equivalent to κ_2

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \equiv \kappa} \text{KEquiv-Ref1} \qquad \frac{\Delta; \Phi \vdash \kappa_2 \equiv \kappa_1}{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2} \text{KEquiv-Symm} \\
\\
\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_3 \quad \Delta; \Phi \vdash \kappa_3 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2} \text{KEquiv-Trans} \\
\\
\frac{\Delta; \Phi \vdash \tau :: \mathbf{S}_{\kappa}(\tau_1)}{\Delta; \Phi \vdash \mathbf{S}_{\kappa}(\tau_1)(\tau) \equiv \mathbf{S}_{\kappa}(\tau_1)} \text{KEquiv-SKind}_{\mathbf{SKind}} \qquad \frac{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_1} \cdot \kappa_2}{\Delta; \Phi \vdash \mathbf{S}_{\Pi_{t::\kappa_1} \cdot \kappa_2}(\tau) \equiv \Pi_{t::\kappa_1} \cdot \mathbf{S}_{[t_1/t]\kappa_2}(\tau \ t_1)} \text{KEquiv-SKind}_{\Pi} \\
\\
\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \quad \Delta; \Phi, t::\kappa_1 \vdash \kappa_3 \equiv \kappa_4}{\Delta; \Phi \vdash \Pi_{t::\kappa_1} \cdot \kappa_2 \equiv \Pi_{t::\kappa_3} \cdot \kappa_4} \text{KEquiv-}\Pi \qquad \frac{\Delta; \Phi \vdash \tau_1 \stackrel{\kappa_1}{\equiv} \tau_2 \quad \Delta; \Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta; \Phi \vdash \mathbf{S}_{\kappa_1}(\tau_1) \equiv \mathbf{S}_{\kappa_2}(\tau_2)} \text{KEquiv-SKind}
\end{array}$$

$\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$ κ_1 is a consistent subkind of κ_2

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \mathbf{KHole} \lesssim \kappa} \text{CSK-KHoleL} \quad \frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \mathbf{KHole}} \text{CSK-KHoleR}$$

$$\frac{\Delta; \Phi \vdash \mathbf{SKHole}(\tau) \text{ OK} \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \mathbf{SKHole}(\tau) \lesssim \kappa} \text{CSK-SKind}_{\mathbf{KHoleL}}$$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK} \quad \Delta; \Phi \vdash \mathbf{SKHole}(\tau) \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \mathbf{SKHole}(\tau)} \text{CSK-SKind}_{\mathbf{KHoleR}}$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \text{CSK-KEquiv} \quad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_3 \quad \Delta; \Phi \vdash \kappa_3 \lesssim \kappa_4 \quad \Delta; \Phi \vdash \kappa_4 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \text{CSK-Normal}$$

$$\frac{\Delta; \Phi \vdash \mathbf{S}_\kappa(\tau) \text{ OK}}{\Delta; \Phi \vdash \mathbf{S}_\kappa(\tau) \lesssim \kappa} \text{CSK-SKind} \quad \frac{\Delta; \Phi \vdash \kappa_3 \lesssim \kappa_1 \quad \Delta; \Phi, t::\kappa_3 \vdash \kappa_2 \lesssim \kappa_4}{\Delta; \Phi \vdash \Pi_{t::\kappa_1}.\kappa_2 \lesssim \Pi_{t::\kappa_3}.\kappa_4} \text{CSK-}\Pi$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2 \quad \Delta; \Phi \vdash \tau_1 \equiv^{\kappa_1} \tau_2}{\Delta; \Phi \vdash \mathbf{S}_{\kappa_1}(\tau_1) \lesssim \mathbf{S}_{\kappa_2}(\tau_2)} \text{CSK-?}$$

$\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2$ τ_1 is provably equivalent to τ_2 at kind κ

$$\frac{\Delta; \Phi \vdash \tau::\kappa}{\Delta; \Phi \vdash \tau \equiv \tau} \text{EquivAK-Ref1} \quad \frac{\Delta; \Phi \vdash \tau_2 \equiv^{\kappa} \tau_1}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} \text{EquivAK-Symm}$$

$$\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_3 \quad \Delta; \Phi \vdash \tau_3 \equiv^{\kappa} \tau_1}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} \text{EquivAK-Trans}$$

$$\frac{\Delta; \Phi \vdash \tau_1 ::> \kappa_1 \quad \Delta; \Phi \vdash \kappa_1 \equiv \mathbf{S}_\kappa(\tau_2)}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} \text{EquivAK-SKind}$$

$$\frac{\Delta; \Phi \vdash \tau_1 :: \Pi_{t::\kappa_1}.\kappa_3 \quad \Delta; \Phi \vdash \tau_2 :: \Pi_{t::\kappa_1}.\kappa_4 \quad \Delta; \Phi, t::\kappa_1 \vdash \tau_1 \equiv^{\kappa_2} \tau_2 \quad t}{\Delta; \Phi \vdash \tau_1 \equiv^{\Pi_{t::\kappa_1}.\kappa_2} \tau_2} \text{EquivAK-}\Pi$$

$$\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\Pi_{t::\kappa_1}.\kappa_2} \tau_3 \quad \Delta; \Phi \vdash \tau_2 \equiv^{\kappa_1} \tau_4}{\Delta; \Phi \vdash \tau_1 \tau_2 \equiv^{[\tau_2/t]\kappa_2} \tau_3 \tau_4} \text{EquivAK-Ap}$$

$$\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\mathbf{S}_\kappa(\tau)} \tau_2}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} (1)$$

$$\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\text{Type}} \tau_3 \quad \Delta; \Phi \vdash \tau_2 \equiv^{\text{Type}} \tau_4}{\Delta; \Phi \vdash \tau_1 \oplus \tau_2 \equiv^{\text{Type}} \tau_3 \oplus \tau_4} (2)$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \quad \Delta; \Phi, t::\kappa_1 \vdash \tau_1 \equiv^{\kappa} \tau_2}{\Delta; \Phi \vdash \lambda t::\kappa_1.\tau_1 \equiv^{\Pi_{t::\kappa_1}.\kappa} \lambda t::\kappa_2.\tau_2} (3)$$

$$\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa_1} \tau_2 \quad \Delta; \Phi \vdash \kappa_1 \equiv \kappa}{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa} \tau_2} (4)$$

$\Delta; \Phi \vdash \kappa \text{ OK}$ κ is well formed

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{Type OK}} \text{KWF-Type}$$

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{KHole OK}} \text{KWF-KHole}$$

$$\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \text{S}_\kappa(\tau) \text{ OK}} \text{KWF-SKind}$$

$$\frac{\Delta; \Phi, \underline{t :: \kappa_1} \vdash \kappa_2 \text{ OK}}{\Delta; \Phi \vdash \Pi_{t :: \kappa_1} \cdot \kappa_2 \text{ OK}} \text{KWF-}\Pi$$

$\Delta; \Phi \vdash \text{OK}$ Context is well formed

$$\frac{}{\cdot; \cdot \vdash \text{OK}} \text{CWF-Nil}$$

$$\frac{t \notin \Phi \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi, \underline{t :: \kappa} \vdash \text{OK}} \text{CWF-TypVar}$$

$$\frac{u \notin \Delta \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta, u :: \kappa; \Phi \vdash \text{OK}} \text{CWF-Hole}$$

METATHEORY

Lemma 1 (COK). *If $\Delta; \Phi \vdash \mathcal{J}$, then $\Delta; \Phi \vdash \text{OK}$*

Proof. By simultaneous induction on derivations.

No interesting cases. □

Lemma 2 (Exchange).

If $\Delta; \Phi_1, \underline{t_{L1} :: \kappa_{L1}}, \underline{t_{L2} :: \kappa_{L2}}, \Phi_2 \vdash \mathcal{J}$ and $\Delta; \Phi_1, \underline{t_{L2} :: \kappa_{L2}}, \underline{t_{L1} :: \kappa_{L1}}, \Phi_2 \vdash \text{OK}$, then $\Delta; \Phi_1, \underline{t_{L2} :: \kappa_{L2}}, \underline{t_{L1} :: \kappa_{L1}}, \Phi_2 \vdash \mathcal{J}$

Proof. By simultaneous induction on derivations.

No interesting cases.

(Only rules with Φ extended in the consequent are interesting, which is only CWF-TypVar, but when \mathcal{J} is CWF, Exchange is identity) □

Lemma 3 (Weakening).

If $\Delta; \Phi \vdash \mathcal{J}$ and $\Delta; \Phi, \underline{t_L :: \kappa_L} \vdash \text{OK}$, then $\Delta; \Phi, \underline{t_L :: \kappa_L} \vdash \mathcal{J}$

Lemma 4 (OK-PK). *If $\Delta; \Phi \vdash \tau :: > \kappa$, then $\Delta; \Phi \vdash \kappa \text{ OK}$*

Lemma 5 (OK-WFaK). *If $\Delta; \Phi \vdash \tau :: \kappa$, then $\Delta; \Phi \vdash \kappa \text{ OK}$*

Lemma 6 (OK-MatchPi). *If $\Delta; \Phi \vdash \kappa \Pi_{t :: \kappa_1} \cdot \kappa_2$, then $\Delta; \Phi \vdash \kappa \text{ OK}$ and $\Delta; \Phi \vdash \Pi_{t :: \kappa_1} \cdot \kappa_2 \text{ OK}$*

Lemma 7 (OK-KEquiv). *If $\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2$, then $\Delta; \Phi \vdash \kappa_1 \text{ OK}$ and $\Delta; \Phi \vdash \kappa_2 \text{ OK}$*

Lemma 8 (OK-CSK). *If $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$, then $\Delta; \Phi \vdash \kappa_1 \text{ OK}$ and $\Delta; \Phi \vdash \kappa_2 \text{ OK}$*

Lemma 9 (OK-EquivAK). *If $\Delta; \Phi \vdash \tau_1 \equiv^\kappa \tau_2$, then $\Delta; \Phi \vdash \tau_1 :: \kappa$ and $\Delta; \Phi \vdash \tau_2 :: \kappa$ and $\Delta; \Phi \vdash \kappa \text{ OK}$*

Lemma 10 (OK-Substitution).

*If $\Delta; \Phi \vdash \tau_L :: \kappa_{L1}$ and $\Delta; \Phi, \underline{t_L :: \kappa_{L1}} \vdash \kappa_{L2} \text{ OK}$, then $\Delta; \Phi \vdash [\tau_L / t_L] \kappa_{L2} \text{ OK}$
(induction on $\Delta; \Phi, \underline{t_L :: \kappa_{L1}} \vdash \kappa_{L2} \text{ OK}$)*

Lemma 11 (K-Substitution).

*If $\Delta; \Phi \vdash \tau_{L1} :: \kappa_{L1}$ and $\Delta; \Phi, \underline{t_L :: \kappa_{L1}} \vdash \tau_{L2} :: \kappa_{L2}$, then $\Delta; \Phi \vdash [\tau_{L1} / t_L] \tau_{L2} :: [\tau_{L1} / t_L] \kappa_{L2}$
(induction on $\Delta; \Phi, \underline{t_L :: \kappa_{L1}} \vdash \tau_{L2} :: \kappa_{L2}$)*

Proof. By simultaneous induction on derivations.

The interesting cases per lemma:

Weakening.	PK-λ	$\Delta; \Phi, t_L :: \kappa_2 \vdash \text{OK}$ $\Delta; \Phi, t :: \kappa_1 \vdash \text{OK}$ $\Delta; \Phi \vdash \kappa_1 \text{ OK}$ $\Delta; \Phi, t_L :: \kappa_L \vdash \kappa_1 \text{ OK}$ $\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_1 \vdash \text{OK}$ $\Delta; \Phi, t :: \kappa_1, t_L :: \kappa_L \vdash \tau :: > \kappa_2$ $\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_1 \vdash \tau :: > \kappa_2$	by IH by COK on premiss by premiss by Weakening on premiss
OK-PK.	PK-Base	$\Delta; \Phi \vdash \text{bse} :: \text{S}_{\text{Type}}(\text{bse})$ $\Delta; \Phi \vdash \text{bse} :: \text{Type}$	by (9) by (10)
	*	$\Delta; \Phi \vdash \text{S}_{\text{Type}}(\text{bse}) \text{ OK}$	by (43)
	*	$\Delta; \Phi \vdash \text{OK}$	by premiss
	PK-App		bad
OK-WFaK.	(12)	$\Delta; \Phi \vdash \tau_2 :: \kappa$	by (10)
	*	$\Delta; \Phi \vdash \text{S}_{\kappa}(\tau_2) \text{ OK}$	by (43)
OK-KEquiv.	(22)	$\Delta; \Phi \vdash \tau t :: > \kappa$	
OK-Substitution.	(41)	$\Delta; \Phi, t_L :: \kappa_{L1} \vdash \text{OK}$ $\Delta; \Phi \vdash \kappa_{L1} \text{ OK}$	premiss (41) by subderivation premiss (46)
	*	$\Delta; \Phi \vdash \text{OK}$	by OK-KWF
	*	$\Delta; \Phi \vdash [\tau_L / t_L] \text{Type} \text{ OK}$	by (41) and degenerate subst
	(43)	$\Delta; \Phi, t_L :: \kappa_{L1} \vdash \tau :: \kappa$ $\Delta; \Phi, t_L :: \kappa_{L1} \vdash \text{OK}$ $\Delta; \Phi \vdash \kappa_{L1} \text{ OK}$	premiss (43) by OK-WFaK by subderivation premiss (46)
	*	$\Delta; \Phi \vdash \text{OK}$	by OK-KWF
	*	$\Delta; \Phi \vdash [\tau_{L1} / t_L] \tau :: [\tau_{L1} / t_L] \kappa$ $\Delta; \Phi \vdash [\tau_L / t_L] \text{S}_{\kappa}(\tau) \text{ OK}$	by K-Substitution on premiss
			by (43)

□

Lemma 12 (PK-Unicity). *If $\Delta; \Phi \vdash \tau_L :: > \kappa_{L1}$ and $\Delta; \Phi \vdash \tau_L :: > \kappa_{L2}$ then κ_{L1} is κ_{L2}*

Lemma 13. *If $\Delta; \Phi \vdash \tau :: > \kappa_1$ and $\Delta; \Phi \vdash \tau :: \kappa_2$, then $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$*

Lemma 14. *If $\Delta; \Phi \vdash \kappa_1 \lesssim \text{S}_{\kappa_2}(\tau)$, then $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$*