Hazel Phi: 9-type-aliases

July 13, 2021

Syntax

$$
\begin{array}{rlll}
\text{BinOp} & \oplus & ::= & \times \mid + \mid \to \\
\text{Kind} & \kappa & ::= & \mathtt{Type} \mid \mathtt{KHole} \mid \mathtt{S}_{\kappa}(\tau) \mid \Pi_{t::\kappa_1}.\kappa_2 \\
\text{Base Types} & \mathtt{bse} & ::= & \mathtt{Int} \mid \mathtt{Float} \mid \mathtt{Bool} \\
\text{User Types} & \hat{\tau} & ::= & t \mid \mathtt{bse} \mid \hat{\tau_1} \oplus \hat{\tau_2} \mid (\!|\,|\!)^{\mathtt{u}} \mid (\!|\hat{\tau}|\!)^{\mathtt{u}} \mid \lambda t{::}\mathtt{Type}.\hat{\tau} \mid \hat{\tau_1}\ \hat{\tau_2} \\
\text{Internal Types} & \hat{\tau} & ::= & t \mid \mathtt{bse} \mid \tau_1 \oplus \tau_2 \mid (\!|\,|\!)^{\mathtt{u}} \mid (\!|\tau|\!)^{\mathtt{u}} \mid \lambda t{::}\kappa.\tau \mid \tau_1\ \tau_2 \\
\text{Type Pattern} & & & \\
\text{User Expression} & & & \\
\text{Internal Expression} & & &
\end{array}
$$

Declaratives

$\boxed{\Delta;\Phi \vdash \kappa_1 \lesssim \kappa_2}$ κ_1 is a consistent subkind of κ_2

$$
\frac{}{\Delta;\Phi \vdash \mathtt{KHole} \lesssim \kappa}
\qquad
\frac{}{\Delta;\Phi \vdash \kappa \lesssim \mathtt{KHole}}
\qquad
\frac{\Delta;\Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta;\Phi \vdash \kappa_1 \lesssim \kappa_2}
\qquad
\frac{\Delta;\Phi \vdash \tau{::}\kappa}{\Delta;\Phi \vdash \mathtt{S}_{\kappa}(\tau) \lesssim \kappa}
$$

$\boxed{\Delta;\Phi \vdash \kappa_1 \equiv \kappa_2}$ κ_1 is equivalent to κ_2

$$
\frac{}{\Delta;\Phi \vdash \kappa \equiv \kappa}
\qquad
\frac{\Delta;\Phi \vdash \kappa_2 \equiv \kappa_1}{\Delta;\Phi \vdash \kappa_1 \equiv \kappa_2}
\qquad
\frac{\Delta;\Phi \vdash \kappa_1 \equiv \kappa_3 \qquad \Delta;\Phi \vdash \kappa_3 \equiv \kappa_2}{\Delta;\Phi \vdash \kappa_1 \equiv \kappa_2}
$$

$$
\frac{\Delta;\Phi \vdash \tau_1 \overset{\kappa}{\equiv} \tau_2}{\Delta;\Phi \vdash \mathtt{S}_{\kappa}(\tau_1) \equiv \mathtt{S}_{\kappa}(\tau_2)}
\qquad
\frac{\Delta;\Phi \vdash \tau{::}\mathtt{S}_{\kappa}(\tau_1)}{\Delta;\Phi \vdash \mathtt{S}_{\mathtt{S}_{\kappa}(\tau_1)}(\tau) \equiv \mathtt{S}_{\kappa}(\tau_1)}
\qquad
\frac{\Delta;\Phi \vdash \tau{::}\Pi_{t::\kappa_1}.\kappa_2}{\Delta;\Phi \vdash \mathtt{S}_{\Pi_{t::\kappa_1}.\kappa_2}(\tau) \equiv \Pi_{t::\kappa_1}.\mathtt{S}_{\kappa_2}(\tau\ t)}
$$

$\boxed{\Delta;\Phi \vdash \tau_1 \overset{\kappa}{\equiv} \tau_2}$ τ_1 is equivalent to τ_2 at kind κ

$$\frac{\Delta;\Phi \vdash \tau :: \kappa}{\Delta;\Phi \vdash \tau \overset{\kappa}{\equiv} \tau} \qquad \frac{\Delta;\Phi \vdash \tau_2 \overset{\kappa}{\equiv} \tau_1}{\Delta;\Phi \vdash \tau_1 \overset{\kappa}{\equiv} \tau_2} \qquad \frac{\Delta;\Phi \vdash \tau_1 \overset{\kappa}{\equiv} \tau_3 \qquad \Delta;\Phi \vdash \tau_3 \overset{\kappa}{\equiv} \tau_1}{\Delta;\Phi \vdash \tau_1 \overset{\kappa}{\equiv} \tau_2}$$

$$\frac{\Delta;\Phi \vdash \tau_1 \overset{\mathtt{Type}}{\equiv} \tau_3 \qquad \Delta;\Phi \vdash \tau_2 \overset{\mathtt{Type}}{\equiv} \tau_4}{\Delta;\Phi \vdash \tau_1 \oplus \tau_2 \overset{\mathtt{Type}}{\equiv} \tau_3 \oplus \tau_4} \qquad \frac{\Delta;\Phi \vdash \tau_1 :: \mathtt{S}_{\kappa}(\tau_2) \qquad \Delta;\Phi \vdash \tau_2 :: \kappa}{\Delta;\Phi \vdash \tau_1 \overset{\mathtt{S}_{\kappa}(\tau_2)}{\equiv} \tau_2}$$

$$\frac{\Delta;\Phi \vdash \tau_1 \overset{\kappa_1}{\equiv} \tau_2 \qquad \Delta;\Phi \vdash \kappa_1 \equiv \kappa}{\Delta;\Phi \vdash \tau_1 \overset{\kappa}{\equiv} \tau_2}$$

$\boxed{\Delta;\Phi \vdash \tau_1 \equiv \tau_2}$ τ_1 is equivalent to τ_2 at “top” kind

$$\frac{\Delta;\Phi \vdash \tau_1 \overset{\kappa}{\equiv} \tau_2}{\Delta;\Phi \vdash \tau_1 \equiv \tau_2} \qquad \frac{\Delta;\Phi \vdash \tau :: \kappa}{\Delta;\Phi \vdash \tau \equiv \tau} \qquad \frac{\Delta;\Phi \vdash \tau_2 \equiv \tau_1}{\Delta;\Phi \vdash \tau_1 \equiv \tau_2} \qquad \frac{\Delta;\Phi \vdash \tau_1 \equiv \tau_3 \qquad \Delta;\Phi \vdash \tau_3 \equiv \tau_2}{\Delta;\Phi \vdash \tau_1 \equiv \tau_2}$$

$\boxed{\Delta;\Phi \vdash \tau ::> \kappa}$ τ has principal kind κ

$$\frac{}{\Delta;\Phi \vdash \mathtt{bse} ::> \mathtt{S}_{\mathtt{Type}}(\mathtt{bse})}$$