

Hazel Phi: 11-type-constructors

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SYNTAX

Kind	$\kappa ::= \text{Type} \mid \text{KHole} \mid \text{S}_\kappa(\tau) \mid \Pi_{t::\kappa_1.\kappa_2}$
User Types	$\hat{\tau} ::= t \mid \mathbf{bse} \mid \tau_1 \oplus \tau_2 \mid \emptyset^u \mid \langle \hat{\tau} \rangle^u \mid \lambda t::\text{Type}.\hat{\tau} \mid \tau_1 \hat{\tau}_2$
Internal Types	$\tau ::= t \mid \mathbf{bse} \mid \tau_1 \oplus \tau_2 \mid \emptyset^u \mid \langle \tau \rangle^u \mid \langle t \rangle^u \mid \lambda t::\kappa.\tau \mid \tau_1 \tau_2$
Base Types	$\mathbf{bse} ::= \text{Int} \mid \text{Float} \mid \text{Bool}$
BinOp	$\oplus ::= \times \mid + \mid \rightarrow$
Type Pattern	
User Expression	
Internal Expression	

DECLARATIVES

$\Delta; \Phi \vdash \tau ::> \kappa$ τ has principal (well formed) kind κ

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \mathbf{bse} ::> \text{S}_{\text{Type}}(\mathbf{bse})} \text{PK-Base} \quad \frac{\Delta; \Phi_1, t::\kappa, \Phi_2 \vdash \text{OK}}{\Delta; \Phi \vdash t ::> \text{S}_\kappa(t)} \text{PK-Var} \quad \frac{\Delta; \Phi \vdash \tau_1 :: \text{Type} \quad \Delta; \Phi \vdash \tau_2 :: \text{Type}}{\Delta; \Phi \vdash \tau_1 \oplus \tau_2 ::> \text{S}_{\text{Type}}(\tau_1 \oplus \tau_2)} \text{PK-}\oplus \quad \frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \emptyset^u ::> \text{S}_\kappa(\emptyset^u)} \text{PK-EHole} \quad \frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \text{OK} \quad \Delta; \Phi \vdash \tau::\kappa_1}{\Delta; \Phi \vdash \langle \tau \rangle^u ::> \text{S}_\kappa(\langle \tau \rangle^u)} \text{PK-NEHole} \\
\\
\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \text{OK} \quad t \notin \Phi}{\Delta; \Phi \vdash \langle t \rangle^u ::> \text{S}_\kappa(\langle t \rangle^u)} \text{PK-Unbound} \quad \frac{\Delta; \Phi, t::\kappa_1 \vdash \tau ::> \kappa_2}{\Delta; \Phi \vdash \lambda t::\kappa_1.\tau ::> \text{S}_{\Pi_{t::\kappa_1}.\kappa_2}(\lambda t::\kappa_1.\tau)} \text{PK-}\lambda \quad \frac{\Delta; \Phi \vdash \tau_1 ::> \kappa \quad \Delta; \Phi \vdash \kappa \multimap \Pi_{t::\kappa_1}.\kappa_2 \quad \Delta; \Phi \vdash \tau_2 :: \kappa_1}{\Delta; \Phi \vdash \tau_1 \tau_2 ::> [\tau_2/t]\kappa_2} \text{PK-Ap}
\end{array}$$

$\Delta; \Phi \vdash \tau :: \kappa$ τ is well formed at kind κ

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \tau ::> \text{S}_\kappa(\tau)}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-1} \quad \frac{\Delta; \Phi \vdash \tau ::> \kappa_1 \quad \Delta; \Phi \vdash \kappa_1 \lesssim \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Subsump} \\
\\
\frac{\Delta; \Phi \vdash \tau ::> \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Reit} \quad \frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \tau :: \text{S}_\kappa(\tau)} \text{WFaK-Self} \quad \frac{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_3}.\kappa_4 \quad \Delta; \Phi \vdash \Pi_{t::\kappa_3}.\kappa_4 \lesssim \Pi_{t::\kappa_1}.\kappa_2}{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_1}.\kappa_2} \text{WFaK-}\Pi\text{CSKTrans} \quad \frac{\Delta; \Phi \vdash \tau :: \text{S}_\kappa(\tau_1) \quad \Delta; \Phi \vdash \tau_1 :: \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Flatten}
\end{array}$$

$\Delta; \Phi \vdash \kappa \multimap \Pi_{t::\kappa_1}.\kappa_2$ κ has matched Π -kind $\Pi_{t::\kappa_1}.\kappa_2$

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{KHole} \multimap \Pi_{t::\text{KHole}}.\text{KHole}} \multimap\text{-KHole} \quad \frac{\Delta; \Phi \vdash \kappa \equiv \text{SKHole}(\tau)}{\Delta; \Phi \vdash \kappa \multimap \Pi_{t::\text{SKHole}(\tau)}.\text{SKHole}(\tau \ t)} \multimap\text{-SKHole} \quad \frac{\Delta; \Phi \vdash \kappa \equiv \Pi_{t::\kappa_1}.\kappa_2}{\Delta; \Phi \vdash \kappa \multimap \Pi_{t::\kappa_1}.\kappa_2} \multimap\text{-}\Pi
\end{array}$$

$\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2$ κ_1 is equivalent to κ_2

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \equiv \kappa} \text{KEquiv-Ref1} \quad \frac{\Delta; \Phi \vdash \kappa_2 \equiv \kappa_1}{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2} \text{KEquiv-Symm} \quad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_3 \quad \Delta; \Phi \vdash \kappa_3 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2} \text{KEquiv-Trans} \\
\\
\frac{\Delta; \Phi \vdash \tau :: \text{S}_\kappa(\tau_1)}{\Delta; \Phi \vdash \text{S}_{\text{S}_\kappa(\tau_1)}(\tau) \equiv \text{S}_\kappa(\tau_1)} \text{KEquiv-SKind}_{\text{SKind}} \quad \frac{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_1}.\kappa_2}{\Delta; \Phi \vdash \text{S}_{\Pi_{t::\kappa_1}.\kappa_2}(\tau) \equiv \Pi_{t::\kappa_1}.\text{S}_{[t_1/t]\kappa_2}(\tau \ t_1)} \text{KEquiv-SKind}_{\Pi} \quad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \quad \Delta; \Phi, t::\kappa_1 \vdash \kappa_3 \equiv \kappa_4}{\Delta; \Phi \vdash \Pi_{t::\kappa_1}.\kappa_2 \equiv \Pi_{t::\kappa_3}.\kappa_4} \text{KEquiv-}\Pi \quad \frac{\Delta; \Phi \vdash \tau_1 \equiv \tau_2 \quad \Delta; \Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta; \Phi \vdash \text{S}_{\kappa_1}(\tau_1) \equiv \text{S}_{\kappa_2}(\tau_2)} \text{KEquiv-SKind}
\end{array}$$

$\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$ κ_1 is a consistent subkind of κ_2

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \text{KHole} \lesssim \kappa} \text{CSK-KHoleL} \quad \frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \text{KHole}} \text{CSK-KHoleR} \quad \frac{\Delta; \Phi \vdash \text{SKHole}(\tau) \text{ OK} \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \text{SKHole}(\tau) \lesssim \kappa} \text{CSK-SKind}_{\text{KHoleL}} \quad \frac{\Delta; \Phi \vdash \kappa \text{ OK} \quad \Delta; \Phi \vdash \text{SKHole}(\tau) \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \text{SKHole}(\tau)} \text{CSK-SKind}_{\text{KHoleR}} \\
\\
\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \text{CSK-KEquiv} \quad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_3 \quad \Delta; \Phi \vdash \kappa_3 \lesssim \kappa_4}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \text{CSK-Normal} \quad \frac{\Delta; \Phi \vdash \text{S}_\kappa(\tau) \text{ OK}}{\Delta; \Phi \vdash \text{S}_\kappa(\tau) \lesssim \kappa} \text{CSK-SKind} \quad \frac{\Delta; \Phi \vdash \kappa_3 \lesssim \kappa_1 \quad \Delta; \Phi, t::\kappa_3 \vdash \kappa_2 \lesssim \kappa_4}{\Delta; \Phi \vdash \Pi_{t::\kappa_1}.\kappa_2 \lesssim \Pi_{t::\kappa_3}.\kappa_4} \text{CSK-}\Pi
\end{array}$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2 \quad \Delta; \Phi \vdash \tau_1 \equiv \tau_2}{\Delta; \Phi \vdash \text{S}_{\kappa_1}(\tau_1) \lesssim \text{S}_{\kappa_2}(\tau_2)} \text{CSK-?}$$

$\Delta; \Phi \vdash \tau_1 \equiv \tau_2$ τ_1 is provably equivalent to τ_2 at kind κ

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \tau \equiv \tau} \text{EquivAK-Ref1} \quad \frac{\Delta; \Phi \vdash \tau_2 \equiv \tau_1}{\Delta; \Phi \vdash \tau_1 \equiv \tau_2} \text{EquivAK-Symm} \quad \frac{\Delta; \Phi \vdash \tau_1 \equiv \tau_3 \quad \Delta; \Phi \vdash \tau_3 \equiv \tau_1}{\Delta; \Phi \vdash \tau_1 \equiv \tau_2} \text{EquivAK-Trans} \\
\\
\frac{\Delta; \Phi \vdash \tau_1 ::> \kappa_1 \quad \Delta; \Phi \vdash \kappa_1 \equiv \text{S}_\kappa(\tau_2)}{\Delta; \Phi \vdash \tau_1 \equiv \tau_2} \text{EquivAK-SKind} \quad \frac{\Delta; \Phi \vdash \tau_1 :: \Pi_{t::\kappa_1}.\kappa_3 \quad \Delta; \Phi \vdash \tau_2 :: \Pi_{t::\kappa_1}.\kappa_4 \quad \Delta; \Phi, t::\kappa_1 \vdash \tau_1 \ t \equiv \tau_2 \ t}{\Delta; \Phi \vdash \tau_1 \equiv \tau_2} \text{EquivAK-}\Pi \quad \frac{\Delta; \Phi \vdash \tau_1 \equiv \tau_3 \quad \Delta; \Phi \vdash \tau_2 \equiv \tau_4}{\Delta; \Phi \vdash \tau_1 \tau_2 \equiv \tau_3 \tau_4} \text{EquivAK-Ap} \\
\\
\frac{\Delta; \Phi \vdash \tau_1 \equiv \tau_2}{\Delta; \Phi \vdash \tau_1 \equiv \tau_2} \text{(1)} \quad \frac{\Delta; \Phi \vdash \tau_1 \equiv \tau_2 \quad \Delta; \Phi, t::\kappa_1 \vdash \tau_1 \ t \equiv \tau_2 \ t}{\Delta; \Phi \vdash \lambda t::\kappa_1.\tau_1 \equiv \lambda t::\kappa_2.\tau_2} \text{(3)} \quad \frac{\Delta; \Phi \vdash \tau_1 \equiv \tau_2 \quad \Delta; \Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta; \Phi \vdash \tau_1 \equiv \tau_2} \text{(4)} \\
\\
\frac{\Delta; \Phi \vdash \tau_1 \equiv \tau_3 \quad \Delta; \Phi \vdash \tau_2 \equiv \tau_4}{\Delta; \Phi \vdash \tau_1 \oplus \tau_2 \equiv \tau_3 \oplus \tau_4} \text{(2)}
\end{array}$$

$\Delta; \Phi \vdash \kappa \text{ OK}$ κ is well formed

$$\begin{array}{c}
\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{Type} \text{ OK}} \text{KWF-Type} \quad \frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{KHole} \text{ OK}} \text{KWF-KHole} \quad \frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \text{S}_\kappa(\tau) \text{ OK}} \text{KWF-SKind} \quad \frac{\Delta; \Phi, t::\kappa_1 \vdash \kappa_2 \text{ OK}}{\Delta; \Phi \vdash \Pi_{t::\kappa_1}.\kappa_2 \text{ OK}} \text{KWF-}\Pi
\end{array}$$

$\Delta; \Phi \vdash \text{OK}$ Context is well formed

$$\begin{array}{c}
\frac{}{.; \vdash \text{OK}} \text{CWF-Nil} \quad \frac{t \notin \Delta \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi, t::\kappa \vdash \text{OK}} \text{CWF-TypVar} \quad \frac{u \notin \Delta \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta, u::\kappa; \Phi \vdash \text{OK}} \text{CWF-Hole}
\end{array}$$

METATHEORY

Lemma 1 (COK). *If $\Delta; \Phi \vdash \mathcal{J}$, then $\Delta; \Phi \vdash \text{OK}$*

Proof. By simultaneous induction on derivations.
No interesting cases. □

Lemma 2 (Exchange).

If $\Delta; \Phi_1, t_{L1}::\kappa_{L1}, t_{L2}::\kappa_{L2}, \Phi_2 \vdash \mathcal{J}$ and $\Delta; \Phi_1, t_{L2}::\kappa_{L2}, t_{L1}::\kappa_{L1}, \Phi_2 \vdash \text{OK}$, then $\Delta; \Phi_1, t_{L2}::\kappa_{L2}, t_{L1}::\kappa_{L1}, \Phi_2 \vdash \mathcal{J}$

Proof. By induction on derivations.
No interesting cases.

(Only rules with Φ extended in the consequent are interesting, which is only CWF-TypVar, but when \mathcal{J} is CWF, Exchange is identity) □

Corollary 3 (Marked-Exchange).

If $\Delta; \Phi, t_{L1}::\kappa_{L1}, t_{L2}::\kappa_{L2} \vdash \mathcal{J}$ and $\Delta; \Phi, t_{L2}::\kappa_{L2}, t_{L1}::\kappa_{L1} \vdash \text{OK}$, then $\Delta; \Phi, t_{L2}::\kappa_{L2}, t_{L1}::\kappa_{L1} \vdash \mathcal{J}$

Lemma 4 (Weakening).

If $\Delta; \Phi \vdash \mathcal{J}$ and $\Delta; \Phi, t_L::\kappa_L \vdash \text{OK}$, then $\Delta; \Phi, t_L::\kappa_L \vdash \mathcal{J}$

Lemma 5 (OK-PK). *If $\Delta; \Phi \vdash \tau ::> \kappa$, then $\Delta; \Phi \vdash \kappa \text{ OK}$*

Lemma 6 (OK-WFaK). *If $\Delta; \Phi \vdash \tau :: \kappa$, then $\Delta; \Phi \vdash \kappa \text{ OK}$*

Lemma 7 (OK-MatchPi). *If $\Delta; \Phi \vdash \kappa \Pi_{\Pi}^{\blacktriangleright} \Pi_{t::\kappa_1}.\kappa_2$, then $\Delta; \Phi \vdash \kappa$ OK and $\Delta; \Phi \vdash \Pi_{t::\kappa_1}.\kappa_2$ OK*

Lemma 8 (OK-KEquiv). *If $\Delta; \Phi \vdash \kappa_I \equiv \kappa_2$, then $\Delta; \Phi \vdash \kappa_I$ OK and $\Delta; \Phi \vdash \kappa_2$ OK*

Lemma 9 (OK-CSK). *If $\Delta; \Phi \vdash \kappa_I \lesssim \kappa_2$, then $\Delta; \Phi \vdash \kappa_I$ OK and $\Delta; \Phi \vdash \kappa_2$ OK*

Lemma 10 (OK-EquivAK). *If $\Delta; \Phi \vdash \tau_I \stackrel{\kappa}{\equiv} \tau_2$, then $\Delta; \Phi \vdash \tau_I::\kappa$ and $\Delta; \Phi \vdash \tau_2::\kappa$ and $\Delta; \Phi \vdash \kappa$ OK*

Lemma 11 (OK-Substitution).

*If $\Delta; \Phi \vdash \tau_L::\kappa_{L1}$ and $\Delta; \Phi, t_L::\kappa_{L1} \vdash \kappa_{L2}$ OK, then $\Delta; \Phi \vdash [\tau_L/t_L]\kappa_{L2}$ OK
(induction on $\Delta; \Phi, t_L::\kappa_{L1} \vdash \kappa_{L2}$ OK)*

Lemma 12 (K-Substitution).

*If $\Delta; \Phi \vdash \tau_{L1}::\kappa_{L1}$ and $\Delta; \Phi, t_L::\kappa_{L1} \vdash \tau_{L2}::\kappa_{L2}$, then $\Delta; \Phi \vdash [\tau_{L1}/t_L]\tau_{L2}::[\tau_{L1}/t_L]\kappa_{L2}$
(induction on $\Delta; \Phi, t_L::\kappa_{L1} \vdash \tau_{L2}::\kappa_{L2}$)*

Proof. By simultaneous induction on derivations.

The interesting cases per lemma:

Weakening

$$\frac{\frac{\frac{\frac{\frac{\frac{\frac{\Delta; \Phi, t_L :: \kappa_L \vdash \text{OK}}{\Delta; \Phi \vdash \kappa_L \text{ OK}} \text{ IH}}{\Delta; \Phi, t :: \kappa_I \vdash \tau ::> \kappa_2} \text{ premiss}}{\Delta; \Phi, t :: \kappa_I \vdash \tau ::> \kappa_2} \text{ premiss of subderivation}}{\Delta; \Phi, t :: \kappa_I \vdash \tau ::> \kappa_2} \text{ premiss}}{\frac{\frac{\frac{\frac{\frac{\frac{\Delta; \Phi, t :: \kappa_I \vdash \tau ::> \kappa_2}{\Delta; \Phi, t :: \kappa_I \vdash \text{OK}} \text{ COK}}{\Delta; \Phi, t :: \kappa_I, t_L :: \kappa_L \vdash \tau ::> \kappa_2} \text{ Weakening}}{\Delta; \Phi, t :: \kappa_I, t_L :: \kappa_L \vdash \tau ::> \kappa_2} \text{ Weakening}}{\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_I \vdash \tau ::> \kappa_2} \text{ Marked-Exchange}}{\Delta; \Phi, t_L :: \kappa_L \vdash \lambda t :: \kappa_I. \tau ::> \mathbf{S}_{\Pi_0 :: \kappa_I, \kappa_2}(\lambda t :: \kappa_I. \tau)} \text{ PK-}\lambda}$$

Weakening.		$\begin{array}{l} \Delta; \Phi \vdash \kappa_I \text{ OK} \\ \Delta; \Phi, t_L :: \kappa_L \vdash \text{OK} \\ \Delta; \Phi, t_L :: \kappa_L \vdash \kappa_I \text{ OK} \\ \Delta; \Phi, t_L :: \kappa_L, t :: \kappa_I \vdash \text{OK} \\ \Delta; \Phi, t :: \kappa_I, t_L :: \kappa_L \vdash \text{OK} \\ \Delta; \Phi, t :: \kappa_I, t_L :: \kappa_L \vdash \tau ::> \kappa_2 \\ \Delta; \Phi, t_L :: \kappa_L, t :: \kappa_I \vdash \tau ::> \kappa_2 \\ \Delta; \Phi, t_L :: \kappa_L \vdash \lambda t :: \kappa_I. \tau ::> \mathbf{S}_{\Pi_E :: \kappa_I, \kappa_2}(\lambda t :: \kappa_I. \tau) \\ \Delta; \Phi \vdash \mathbf{bse} :: \mathbf{S}_{\text{Type}}(\mathbf{bse}) \\ \Delta; \Phi \vdash \mathbf{bse} :: \mathbf{Type} \\ \Delta; \Phi \vdash \mathbf{S}_{\text{Type}}(\mathbf{bse}) \text{ OK} \\ \Delta; \Phi \vdash \text{OK} \end{array}$	by subderivation premiss by IH by Weakening on subderivation premiss by CWF-TypVar by ? by Weakening on premiss by Marked-Exchange by PK- λ by (9) by (10) by (43) by premiss bad by (10) by (43)
OK-PK.	PK-Base		
	*		
	*		
	PK-Ap		
OK-WFaK.	(12)	$\begin{array}{l} \Delta; \Phi \vdash \tau_2 :: \kappa \\ \Delta; \Phi \vdash \mathbf{S}_{\kappa}(\tau_2) \text{ OK} \\ \Delta; \Phi \vdash \tau \ t ::> \kappa \\ \Delta; \Phi, t_L :: \kappa_{LI} \vdash \text{OK} \\ \Delta; \Phi \vdash \kappa_{LI} \text{ OK} \\ \Delta; \Phi \vdash \text{OK} \\ \Delta; \Phi \vdash [\tau_L / t_L] \mathbf{Type} \text{ OK} \\ \Delta; \Phi, t_L :: \kappa_{LI} \vdash \tau :: \kappa \\ \Delta; \Phi, t_L :: \kappa_{LI} \vdash \text{OK} \\ \Delta; \Phi \vdash \kappa_{LI} \text{ OK} \\ \Delta; \Phi \vdash \text{OK} \\ \Delta; \Phi \vdash [\tau_{LI} / t_L] \tau :: [\tau_{LI} / t_L] \kappa \\ \Delta; \Phi \vdash [\tau_L / t_L] \mathbf{S}_{\kappa}(\tau) \text{ OK} \end{array}$	by subderivation premiss (41) by subderivation premiss (46) by OK-KWF by (41) and degenerate subst premiss (43) by OK-WFaK by subderivation premiss (46) by OK-KWF by K-Substitution on premiss by (43)
OK-KEquiv.	(22)		
OK-Substitution.	(41)		
	*		
	*		
	(43)		
	*		
	*		

□

Lemma 13 (PK-Unicity). *If $\Delta; \Phi \vdash \tau_L ::> \kappa_{LI}$ and $\Delta; \Phi \vdash \tau_L ::> \kappa_{L2}$ then κ_{LI} is κ_{L2}*

Lemma 14. *If $\Delta; \Phi \vdash \tau ::> \kappa_I$ and $\Delta; \Phi \vdash \tau :: \kappa_2$, then $\Delta; \Phi \vdash \kappa_I \lesssim \kappa_2$*

Lemma 15. *If $\Delta; \Phi \vdash \kappa_I \lesssim \mathbf{S}_{\kappa_2}(\tau)$, then $\Delta; \Phi \vdash \kappa_I \lesssim \kappa_2$*