Hazel Phi: 9-type-aliases

July 20, 2021

SYNTAX

```
Kind \kappa ::= Type | KHole | S_{\kappa}(\tau) | \Pi_{t::\kappa_{1}}.\kappa_{2}
User Types \hat{\tau} ::= t | bse | \hat{\tau_{1}} \oplus \hat{\tau_{2}} | (||)^{u} | (\hat{\tau})^{u} | \lambda t::Type.\hat{\tau} | \hat{\tau_{1}} \hat{\tau_{2}}
Internal Types \tau ::= t | bse | \tau_{1} \oplus \tau_{2} | (||)^{u} | (\tau)^{u} | (t)^{u} | \lambda t::\kappa.\tau | \tau_{1} \tau_{2}
Base Types bse ::= Int | Float | Bool
BinOp \oplus ::= \times | + | \to
Type Pattern
User Expression
Internal Expression
```

DECLARATIVES

 $\Delta; \Phi \vdash \tau ::> \kappa \mid \tau \text{ has principal (well formed) kind } \kappa$

$$\frac{\Delta; \Phi \vdash \text{bse} ::> S_{\text{Type}}(\text{bse})}{\Delta; \Phi \vdash \text{bse} ::> S_{\text{Type}}(\text{bse})} (1) \qquad \frac{t :: \kappa \in \Phi}{\Delta; \Phi \vdash t ::> S_{\kappa}(t)} (2)$$

$$\frac{\Delta; \Phi \vdash \tau_{1} :: \text{Type} \qquad \Delta; \Phi \vdash \tau_{2} :: \text{Type}}{\Delta; \Phi \vdash \tau_{1} \oplus \tau_{2} ::> S_{\text{Type}}(\tau_{1} \oplus \tau_{2})} (3) \qquad \frac{u :: \kappa \in \Delta}{\Delta; \Phi \vdash (\|^{u} ::> S_{\kappa}((\|^{u})^{u})} (4)$$

$$\frac{u :: \kappa \in \Delta \qquad \Delta; \Phi \vdash \tau :: \kappa_{1}}{\Delta; \Phi \vdash (\|^{t})^{u} ::> S_{\kappa}((\|^{t})^{u})} (5) \qquad \frac{u :: \kappa \in \Delta \qquad t \notin \text{dom}(\Phi)}{\Delta; \Phi \vdash (\|^{t})^{u} ::> S_{\kappa}((\|^{t})^{u})} (6)$$

$$\frac{\Delta; \Phi \vdash \kappa_{1} \text{ OK} \qquad \Delta; \Phi, t :: \kappa_{1} \vdash \tau ::> \kappa_{2}}{\Delta; \Phi \vdash \lambda t :: \kappa_{1} \cdot \tau ::> S_{\Pi_{t :: \kappa_{1}} \cdot \kappa_{2}} (\lambda t :: \kappa_{1} \cdot \tau)} (7)$$

$$\frac{\Delta; \Phi \vdash \tau_{1} ::> \kappa \qquad \Delta; \Phi \vdash \kappa \prod_{\Pi} \Pi_{t :: \kappa_{1}} \cdot \kappa_{2} \qquad \Delta; \Phi \vdash \tau_{2} :: \kappa_{1}}{\Delta; \Phi \vdash \tau_{1} \tau_{2} ::> [\tau_{2}/t] \kappa_{2}} (8)$$

 $\Delta; \Phi \vdash \tau :: \kappa$ τ is well formed at kind κ

$$\frac{\Delta; \Phi \vdash \tau ::> \kappa}{\Delta; \Phi \vdash \tau ::\kappa} (9) \qquad \frac{\Delta; \Phi \vdash \tau ::S_{\kappa}(\tau_{1})}{\Delta; \Phi \vdash \tau ::\kappa} (10)$$

$$\frac{\Delta; \Phi \vdash \tau ::> \kappa_{1} \qquad \Delta; \Phi \vdash \kappa_{1} \lesssim \kappa}{\Delta; \Phi \vdash \tau ::\kappa} (11)$$

$$\frac{\Delta; \Phi \vdash \tau_{2} ::S_{\kappa}(\tau_{1})}{\Delta; \Phi \vdash \tau_{1} ::S_{\kappa}(\tau_{3})} (12)$$

$$\frac{\Delta; \Phi \vdash \tau_{1} ::S_{\kappa}(\tau_{3}) \qquad \Delta; \Phi \vdash \tau_{3} ::S_{\kappa}(\tau_{2})}{\Delta; \Phi \vdash \tau_{1} ::S_{\kappa}(\tau_{2})} (13)$$

$$\frac{\Delta; \Phi \vdash \tau_{1} ::S_{\kappa}(\tau_{2})}{\Delta; \Phi \vdash \tau_{1} ::S_{\kappa}(\tau_{2})} (14)$$

 $\Delta; \Phi \vdash \kappa \prod_{\Pi} \Pi_{t::\kappa_1}.\kappa_2$ κ has matched Π -kind $\Pi_{t::\kappa_1}.\kappa_2$

$$\frac{\Delta; \Phi \vdash \kappa \equiv \text{KHole}}{\Delta; \Phi \vdash \kappa \stackrel{\blacktriangleright}{\Pi} \Pi_{t::\text{KHole}}.\text{KHole}}$$
(15)
$$\frac{\Delta; \Phi \vdash \kappa \equiv \Pi_{t::\kappa_1}.\kappa_2}{\Delta; \Phi \vdash \kappa \stackrel{\blacktriangleright}{\Pi} \Pi_{t::\kappa_1}.\kappa_2}$$
(16)

 $\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2$ κ_1 is equivalent to κ_2

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \equiv \kappa} (17) \qquad \frac{\Delta; \Phi \vdash \kappa_2 \equiv \kappa_1}{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2} (18) \qquad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_3 \qquad \Delta; \Phi \vdash \kappa_3 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2} (19)$$

$$\frac{\Delta; \Phi \vdash \tau :: S_{\kappa}(\tau_1)}{\Delta; \Phi \vdash S_{S_{\kappa}(\tau_1)}(\tau) \equiv S_{\kappa}(\tau_1)} (20) \qquad \frac{\Delta; \Phi \vdash \tau :: KHole}{\Delta; \Phi \vdash S_{KHole}(\tau) \equiv KHole} (21)$$

$$\frac{\Delta; \Phi \vdash \tau :: \Pi_{t :: \kappa_1} \cdot \kappa_2}{\Delta; \Phi \vdash S_{\Pi_{t :: \kappa_1} \cdot \kappa_2}(\tau) \equiv \Pi_{t :: \kappa_1} \cdot S_{\kappa_2}(\tau t)} (22) \qquad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \qquad \Delta; \Phi, t :: \kappa_1 \vdash \kappa_3 \equiv \kappa_4}{\Delta; \Phi \vdash \Pi_{t :: \kappa_1} \cdot \kappa_2 \equiv \Pi_{t :: \kappa_3} \cdot \kappa_4} (23)$$

$$\frac{\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2}{\Delta; \Phi \vdash S_{\kappa}(\tau_1)} \equiv S_{\kappa}(\tau_2)} (24)$$

 $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$ κ_1 is a consistent subkind of κ_2

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \text{ KHole}} \lesssim \kappa \tag{25} \qquad \frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \text{ KHole}} \tag{26}$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \tag{27} \qquad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_3 \qquad \Delta; \Phi \vdash \kappa_3 \lesssim \kappa_4 \qquad \Delta; \Phi \vdash \kappa_4 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \tag{28}$$

$$\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash S_{\kappa}(\tau) \lesssim \kappa} \tag{29} \qquad \frac{\Delta; \Phi \vdash \kappa_3 \lesssim \kappa_1 \qquad \Delta; \Phi, t :: \kappa_3 \vdash \kappa_2 \lesssim \kappa_4}{\Delta; \Phi \vdash \Pi_{t :: \kappa_1} . \kappa_2 \lesssim \Pi_{t :: \kappa_3} . \kappa_4} \tag{30}$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2 \qquad \Delta; \Phi \vdash \pi_1 \stackrel{\kappa_1}{\equiv} \tau_2}{\Delta; \Phi \vdash S_{\kappa_1}(\tau_1) \lesssim S_{\kappa_2}(\tau_2)} \tag{31}$$

 $\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{=} \tau_2 \mid \tau_1 \text{ is equivalent to } \tau_2 \text{ at kind } \kappa$

$$\frac{\Delta; \Phi \vdash \tau_1 :: \mathbf{S}_{\kappa}(\tau_2)}{\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{=} \tau_2}$$
 (32)

$$\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \tau := \pi} (33) \qquad \frac{\Delta; \Phi \vdash \tau_2 \stackrel{\kappa}{\equiv} \tau_1}{\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2} (34) \qquad \frac{\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_3}{\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2} (35) \qquad \Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2 \qquad (35)$$

$$\frac{\Delta; \Phi \vdash \tau_1 \stackrel{\pi}{\equiv} \tau_3}{\equiv} \tau_3 \qquad \Delta; \Phi \vdash \tau_2 \stackrel{\pi}{\equiv} \tau_4 \qquad (36) \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \qquad \Delta; \Phi, t :: \kappa_1 \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \qquad \Delta; \Phi \vdash \tau_2 \stackrel{\kappa_1}{\equiv} \tau_4 \qquad (38)$$

$$\frac{\Delta; \Phi \vdash \tau_1 \qquad \Xi}{\equiv} \tau_3 \qquad \Delta; \Phi \vdash \tau_2 \stackrel{\kappa_1}{\equiv} \tau_4 \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \tau_1 \equiv \kappa_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \kappa_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \kappa_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \kappa_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash \tau_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \kappa \qquad \Delta; \Phi \vdash \tau_1 \equiv \tau_2 \qquad \Delta; \Phi \vdash$$

 $\Delta; \Phi \vdash \kappa \text{ OK}$ κ is well formed

$$\frac{\Delta; \Phi \vdash \mathsf{Type} \; \mathsf{OK}}{\Delta; \Phi \vdash \mathsf{KHole} \; \mathsf{OK}} \; \frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \mathsf{S}_{\kappa}(\tau) \; \mathsf{OK}} \; (43)}{\Delta; \Phi \vdash \kappa_{1} \; \mathsf{OK} \quad \Delta; \Phi \vdash \kappa_{2} \; \mathsf{OK}}$$

$$\frac{\Delta; \Phi \vdash \kappa_{1} \; \mathsf{OK} \quad \Delta; \Phi \vdash \kappa_{2} \; \mathsf{OK}}{\Delta; \Phi \vdash \Pi_{t :: \kappa_{1}} . \kappa_{2} \; \mathsf{OK}} \; (44)}$$

METATHEORY

Contexts implicitly assumed to be wellformed in the usual way.

Definition 1 (size of kinds a la Stone).

$$\begin{array}{lll} size(\texttt{Type}) & = & 1 \\ size(\texttt{KHole}) & = & 1 \\ size(\texttt{S}_{\kappa}(\tau)) & = & size(\kappa) + 2 \\ size(\texttt{\Pi}_{t::\kappa_1}.\kappa_2) & = & size(\kappa_1) + size(\kappa_2) + 2 \end{array}$$

Lemma 1. If Δ ; $\Phi \vdash \tau ::> \kappa$, then Δ ; $\Phi \vdash \kappa$ OK

Lemma 2. If Δ ; $\Phi \vdash \tau :: \kappa$, then Δ ; $\Phi \vdash \kappa$ OK

Lemma 3. If Δ ; $\Phi \vdash \kappa_1 \equiv \kappa_2$, then Δ ; $\Phi \vdash \kappa_1$ OK and Δ ; $\Phi \vdash \kappa_2$ OK

Lemma 4. If $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$, then $\Delta; \Phi \vdash \kappa_1$ OK and $\Delta; \Phi \vdash \kappa_2$ OK

Lemma 5. If Δ ; $\Phi \vdash \tau ::> \kappa_1$ and Δ ; $\Phi \vdash \tau :: \kappa_2$, then Δ ; $\Phi \vdash \kappa_1 \lesssim \kappa_2$