## Hazel Phi: 9-type-aliases

July 15, 2021

## **SYNTAX**

## **DECLARATIVES**

 $\overline{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2}$   $\kappa_1$  is a consistent subkind of  $\kappa_2$ 

$$\frac{\Delta;\Phi \vdash \kappa \text{ OK}}{\Delta;\Phi \vdash \text{ KHole}} \lesssim \kappa \qquad \qquad \frac{\Delta;\Phi \vdash \kappa \text{ OK}}{\Delta;\Phi \vdash \kappa \lesssim \text{ KHole}}$$
 
$$\frac{\Delta;\Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta;\Phi \vdash \kappa_1 \lesssim \kappa_2} \qquad \frac{\Delta;\Phi \vdash \kappa_1 \equiv \kappa_3 \qquad \Delta;\Phi \vdash \kappa_3 \lesssim \kappa_4 \qquad \Delta;\Phi \vdash \kappa_4 \equiv \kappa_2}{\Delta;\Phi \vdash \kappa_1 \lesssim \kappa_2}$$
 
$$\frac{\Delta;\Phi \vdash \kappa_1 \lesssim \kappa_2}{\Delta;\Phi \vdash \mathbf{S}_\kappa(\tau) \lesssim \kappa} \qquad \frac{\Delta;\Phi \vdash \kappa_3 \lesssim \kappa_1 \qquad \Delta;\Phi,t::\kappa_3 \vdash \kappa_2 \lesssim \kappa_4}{\Delta;\Phi \vdash \mathbf{\Pi}_{t::\kappa_1}.\kappa_2 \lesssim \mathbf{\Pi}_{t::\kappa_3}.\kappa_4}$$
 
$$\frac{\Delta;\Phi \vdash \kappa_1 \lesssim \kappa_2 \qquad \Delta;\Phi \vdash \mathbf{\Pi}_{t::\kappa_1}.\kappa_2 \lesssim \mathbf{\Pi}_{t::\kappa_3}.\kappa_4}{\Delta;\Phi \vdash \mathbf{S}_{\kappa_1}(\tau_1) \lesssim \mathbf{S}_{\kappa_2}(\tau_2)}$$

 $\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2$   $\kappa_1$  is equivalent to  $\kappa_2$ 

$$\begin{array}{lll} \underline{\Delta}; \Phi \vdash \kappa \ \mathsf{OK} \\ \overline{\Delta}; \Phi \vdash \kappa \equiv \kappa \end{array} & \underline{\Delta}; \Phi \vdash \kappa_2 \equiv \kappa_1 \\ \underline{\Delta}; \Phi \vdash \kappa_1 \equiv \kappa_2 \end{array} & \underline{\Delta}; \Phi \vdash \kappa_1 \equiv \kappa_3 \qquad \underline{\Delta}; \Phi \vdash \kappa_3 \equiv \kappa_2 \\ \underline{\Delta}; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2 \\ \underline{\Delta}; \Phi \vdash \mathbf{S}_{\kappa}(\tau_1) \equiv \mathbf{S}_{\kappa}(\tau_2) \end{array} & \underline{\Delta}; \Phi \vdash \tau_{::} \mathbf{S}_{\kappa}(\tau_1) \\ \underline{\Delta}; \Phi \vdash \mathbf{S}_{\mathbf{S}_{\kappa}(\tau_1)}(\tau) \equiv \mathbf{S}_{\kappa}(\tau_1) \end{array} & \underline{\Delta}; \Phi \vdash \tau_{::} \mathbf{\Pi}_{t::\kappa_1}.\kappa_2 \\ \underline{\Delta}; \Phi \vdash \mathbf{S}_{\mathbf{\Pi}_{t::\kappa_1}.\kappa_2}(\tau) \equiv \mathbf{\Pi}_{t::\kappa_1}.\mathbf{S}_{\kappa_2}(\tau \ t) \end{array}$$
$$\underline{\Delta}; \Phi \vdash \kappa_1 \equiv \kappa_2 \qquad \underline{\Delta}; \Phi \vdash \mathbf{S}_{\mathbf{\Pi}_{t::\kappa_1}.\kappa_2}(\tau) \equiv \mathbf{\Pi}_{t::\kappa_1}.\mathbf{S}_{\kappa_2}(\tau \ t)$$

 $\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{=} \tau_2 \mid \tau_1 \text{ is equivalent to } \tau_2 \text{ at kind } \kappa$ 

$$\frac{\Delta; \Phi \vdash \tau_{I} :: S_{\kappa}(\tau_{2})}{\Delta; \Phi \vdash \tau_{I} \stackrel{\kappa}{\equiv} \tau_{2}}$$

$$\frac{\Delta; \Phi \vdash \tau_{I} :: K}{\Delta; \Phi \vdash \tau_{I} \stackrel{\kappa}{\equiv} \tau_{I}} \qquad \Delta; \Phi \vdash \tau_{I} \stackrel{\kappa}{\equiv} \tau_{I} \qquad \Delta; \Phi \vdash \tau_{I} \stackrel{\kappa}$$

 $\Delta; \Phi \vdash \tau ::> \kappa$   $\tau$  has principal kind  $\kappa$ 

$$\frac{t :: \kappa \in \Phi}{\Delta; \Phi \vdash \mathsf{bse} ::> \mathsf{S}_{\mathsf{Type}}(\mathsf{bse})} \qquad \frac{t :: \kappa \in \Phi}{\Delta; \Phi \vdash t ::> \mathsf{S}_{\kappa}(t)} \qquad \frac{\Delta; \Phi \vdash \tau_{1} :: \mathsf{Type} \qquad \Delta; \Phi \vdash \tau_{2} :: \mathsf{Type}}{\Delta; \Phi \vdash \tau_{1} \oplus \tau_{2} ::> \mathsf{S}_{\mathsf{Type}}(\tau_{1} \oplus \tau_{2})} \\ \frac{u :: \kappa \in \Delta}{\Delta; \Phi \vdash (||)^{\mathsf{u}} ::> \kappa} \qquad \frac{u :: \kappa \in \Delta \qquad \Delta; \Phi \vdash \tau :: \kappa_{1}}{\Delta; \Phi \vdash (|\tau|)^{\mathsf{u}} ::> \kappa} \qquad \frac{u :: \kappa \in \Delta \qquad t \notin \mathsf{dom}(\Phi)}{\Delta; \Phi \vdash (|t|)^{\mathsf{u}} ::> \kappa} \\ \frac{\Delta; \Phi \vdash \lambda t :: \kappa_{1} \vdash \tau ::> \kappa_{2}}{\Delta; \Phi \vdash \lambda t :: \kappa_{1} . \tau ::> \mathsf{S}_{\mathsf{\Pi}_{t :: \kappa_{1}} . \kappa_{2}}(\lambda t :: \kappa_{1} . \tau)} \\ \frac{\Delta; \Phi \vdash \tau_{1} ::> \kappa \qquad \Delta; \Phi \vdash \kappa \prod_{\mathsf{II}} \mathsf{\Pi}_{t :: \kappa_{1}} . \kappa_{2}}{\Delta; \Phi \vdash \tau_{1} \ \tau_{2} ::> [\tau_{2}/t] \kappa_{2}} \qquad \Delta; \Phi \vdash \tau_{2} :: \kappa_{1}}$$

 $\Delta; \Phi \vdash \kappa \prod_{\Pi} \Pi_{t::\kappa_1} \cdot \kappa_2$   $\kappa$  has matched  $\Pi$ -kind  $\Pi_{t::\kappa_1} \cdot \kappa_2$ 

$$\frac{\Delta; \Phi \vdash \kappa \equiv \Pi_{t::\kappa_1}.\kappa_2}{\Delta; \Phi \vdash \mathsf{KHole} \ \Pi_{t::\mathsf{KHole}}.\mathsf{KHole}} \qquad \frac{\Delta; \Phi \vdash \kappa \equiv \Pi_{t::\kappa_1}.\kappa_2}{\Delta; \Phi \vdash \kappa \ \Pi_{\mathsf{T}} \ \Pi_{t::\kappa_1}.\kappa_2}$$

 $\Delta; \Phi \vdash \tau :: \kappa$   $\tau$  is well formed at kind  $\kappa$ 

$$\frac{\Delta; \Phi \vdash \tau ::> \kappa}{\Delta; \Phi \vdash \tau ::\kappa} \frac{\Delta; \Phi \vdash \tau :: S_{\kappa}(\tau_{1})}{\Delta; \Phi \vdash \tau ::\kappa} \frac{\Delta; \Phi \vdash \tau ::> \kappa_{1}}{\Delta; \Phi \vdash \kappa_{1} \lesssim \kappa} \frac{\Delta; \Phi \vdash \kappa_{1} \lesssim \kappa}{\Delta; \Phi \vdash \tau ::\kappa}$$

$$\frac{\Delta; \Phi \vdash \tau_{2} :: S_{\kappa}(\tau_{1})}{\Delta; \Phi \vdash \tau_{1} :: S_{\kappa}(\tau_{2})} \frac{\Delta; \Phi \vdash \tau_{1} :: S_{\kappa}(\tau_{3})}{\Delta; \Phi \vdash \tau_{1} :: S_{\kappa}(\tau_{2})}$$