

July 30, 2021

SYNTAX

Kind	$\kappa ::= \text{Type} \mid \text{KHole} \mid \text{S}_\kappa(\tau) \mid \Pi_{t::\kappa_I, \kappa_B}$
User Types	$\hat{\tau} ::= t \mid \text{bse} \mid \tau_I \oplus \tau_B \mid \emptyset^u \mid \langle \hat{\tau} \rangle^u \mid \lambda t::\text{Type}.\hat{\tau} \mid \tau_I' \tau_B'$
Internal Types	$\tau ::= t \mid \text{bse} \mid \tau_I \oplus \tau_B \mid \emptyset^u \mid \langle \hat{\tau} \rangle^u \mid \langle \hat{\tau} \rangle^u \mid \lambda t::\kappa.\tau \mid \tau_I \tau_B$
Base Types	$\text{bse} ::= \text{Int} \mid \text{Float} \mid \text{Bool}$
BinOp	$\oplus ::= \times \mid + \mid \rightarrow$
Type Pattern	
User Expression	
Internal Expression	

DECLARATIVES

 $\Delta; \Phi \vdash \tau ::> \kappa$ τ has principal (well formed) kind κ

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{bse} ::> \text{S}_{\text{Type}}(\text{bse})} \text{PK-Base} \quad \frac{\Delta; \Phi_1, t::\kappa_1, \Phi_2 \vdash \text{OK}}{\Delta; \Phi \vdash t ::> \text{S}_\kappa(t)} \text{PK-Var} \quad \frac{\Delta; \Phi \vdash \tau_I :: \text{Type} \quad \Delta; \Phi \vdash \tau_B :: \text{Type}}{\Delta; \Phi \vdash \tau_I \oplus \tau_B ::> \text{S}_{\text{Type}(\tau_I \oplus \tau_B)}} \text{PK-}\oplus$$

$$\frac{\Delta_1, u::\kappa_1, \Delta_2; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \emptyset^u ::> \text{S}_\kappa(\emptyset^u)} \text{PK-EHole} \quad \frac{\Delta_1, u::\kappa_1, \Delta_2; \Phi \vdash \text{OK} \quad \Delta; \Phi \vdash \tau::\kappa_I}{\Delta; \Phi \vdash \langle \hat{\tau} \rangle^u ::> \text{S}_\kappa(\langle \hat{\tau} \rangle^u)} \text{PK-NEHole}$$

$$\frac{\Delta_1, u::\kappa_1, \Delta_2; \Phi \vdash \text{OK} \quad t \notin \Phi}{\Delta; \Phi \vdash \langle \hat{t} \rangle^u ::> \text{S}_\kappa(\langle \hat{t} \rangle^u)} \text{PK-Unbound}$$

$$\frac{\Delta; \Phi, t::\kappa_I \vdash \tau ::> \kappa_B}{\Delta; \Phi \vdash \lambda t::\kappa_I.\tau ::> \text{S}_{\Pi_{t::\kappa_I, \kappa_B}}(\lambda t::\kappa_I.\tau)} \text{PK-}\lambda$$

$$\frac{\Delta; \Phi \vdash \tau_I ::> \kappa \quad \Delta; \Phi \vdash \kappa \overset{\blacktriangleright}{\Pi} \Pi_{t::\kappa_I, \kappa_B} \quad \Delta; \Phi \vdash \tau_B :: \kappa_I}{\Delta; \Phi \vdash \tau_I \tau_B ::> [\tau_B / t] \kappa_B} \text{PK-Ap}$$

 $\Delta; \Phi \vdash \tau :: \kappa$ τ is well formed at kind κ

$$\frac{\Delta; \Phi \vdash \tau ::> \text{S}_\kappa(\tau)}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-1}$$

$$\frac{\Delta; \Phi \vdash \tau ::> \kappa_I \quad \Delta; \Phi \vdash \kappa_I \lesssim \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Subeump}$$

$$\frac{\Delta; \Phi \vdash \tau ::> \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Reit}$$

$$\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \tau :: \text{S}_\kappa(\tau)} \text{WFaK-Self}$$

$$\frac{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_B, \kappa_I} \quad \Delta; \Phi \vdash \Pi_{t::\kappa_B, \kappa_I} \lesssim \Pi_{t::\kappa_B, \kappa_B}}{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_I, \kappa_B}} \text{WFaK-}\Pi\text{CSKTrans}$$

$$\frac{\Delta; \Phi \vdash \tau :: \text{S}_\kappa(\tau_I) \quad \Delta; \Phi \vdash \tau_I :: \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Flatten}$$

 $\Delta; \Phi \vdash \kappa \overset{\blacktriangleright}{\Pi} \Pi_{t::\kappa_I, \kappa_B}$ κ has matched Π -kind $\Pi_{t::\kappa_I, \kappa_B}$

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{KHole} \overset{\blacktriangleright}{\Pi} \Pi_{t::\text{KHole}, \text{KHole}}} \text{-KHole}$$

$$\frac{\Delta; \Phi \vdash \kappa \equiv \text{S}_{\text{KHole}}(\tau)}{\Delta; \Phi \vdash \kappa \overset{\blacktriangleright}{\Pi} \Pi_{t::\text{S}_{\text{KHole}}(\tau), \text{S}_{\text{KHole}}(\tau \ t)}} \text{-SKHole}$$

$$\frac{\Delta; \Phi \vdash \kappa \equiv \Pi_{t::\kappa_I, \kappa_B}}{\Delta; \Phi \vdash \kappa \overset{\blacktriangleright}{\Pi} \Pi_{t::\kappa_I, \kappa_B}} \text{-}\Pi$$

 $\Delta; \Phi \vdash \kappa_I \equiv \kappa_B$ κ_I is equivalent to κ_B

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \equiv \kappa} \text{KEquiv-Ref1}$$

$$\frac{\Delta; \Phi \vdash \kappa_B \equiv \kappa_I}{\Delta; \Phi \vdash \kappa_I \equiv \kappa_B} \text{KEquiv-Symm}$$

$$\frac{\Delta; \Phi \vdash \kappa_I \equiv \kappa_B \quad \Delta; \Phi \vdash \kappa_B \equiv \kappa_C}{\Delta; \Phi \vdash \kappa_I \equiv \kappa_C} \text{KEquiv-Trans}$$

$$\frac{\Delta; \Phi \vdash \tau :: \text{S}_\kappa(\tau_I)}{\Delta; \Phi \vdash \text{S}_{\kappa(\tau_I)}(\tau) \equiv \text{S}_\kappa(\tau_I)} \text{KEquiv-SKind}_{\text{SKind}}$$

$$\frac{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_I, \kappa_B}}{\Delta; \Phi \vdash \text{S}_{\Pi_{t::\kappa_I, \kappa_B}}(\tau) \equiv \Pi_{t::\kappa_I, \text{S}_{\Pi_{t::\kappa_I, \kappa_B}}(\tau \ t_I)}} \text{KEquiv-SKind}_{\Pi}$$

$$\frac{\Delta; \Phi \vdash \kappa_I \equiv \kappa_B \quad \Delta; \Phi, t::\kappa_I \vdash \kappa_B \equiv \kappa_I}{\Delta; \Phi \vdash \Pi_{t::\kappa_I, \kappa_B} \equiv \Pi_{t::\kappa_B, \kappa_I}} \text{KEquiv-}\Pi$$

$$\frac{\Delta; \Phi \vdash \tau_I \overset{\kappa_I}{\equiv} \tau_B \quad \Delta; \Phi \vdash \kappa_I \equiv \kappa_B}{\Delta; \Phi \vdash \text{S}_{\kappa_I}(\tau_I) \equiv \text{S}_{\kappa_B}(\tau_B)} \text{KEquiv-SKind}$$

 $\Delta; \Phi \vdash \kappa_I \lesssim \kappa_B$ κ_I is a consistent subkind of κ_B

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \text{KHole} \lesssim \kappa} \text{CSK-KHoleL}$$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \text{KHole}} \text{CSK-KHoleR}$$

$$\frac{\Delta; \Phi \vdash \text{S}_{\text{KHole}}(\tau) \text{ OK} \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \text{S}_{\text{KHole}}(\tau) \lesssim \kappa} \text{CSK-SKind}_{\text{SKindL}}$$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK} \quad \Delta; \Phi \vdash \text{S}_{\text{KHole}}(\tau) \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \text{S}_{\text{KHole}}(\tau)} \text{CSK-SKind}_{\text{SKindR}}$$

$$\frac{\Delta; \Phi \vdash \kappa_I \equiv \kappa_B}{\Delta; \Phi \vdash \kappa_I \lesssim \kappa_B} \text{CSK-KEquiv}$$

$$\frac{\Delta; \Phi \vdash \kappa_I \equiv \kappa_B \quad \Delta; \Phi \vdash \kappa_B \lesssim \kappa_I \quad \Delta; \Phi \vdash \kappa_I \equiv \kappa_B}{\Delta; \Phi \vdash \kappa_I \lesssim \kappa_B} \text{CSK-Normal}$$

$$\frac{\Delta; \Phi \vdash \text{S}_\kappa(\tau) \text{ OK}}{\Delta; \Phi \vdash \text{S}_\kappa(\tau) \lesssim \kappa} \text{CSK-SKind}$$

$$\frac{\Delta; \Phi \vdash \kappa_B \lesssim \kappa_I \quad \Delta; \Phi, t::\kappa_B \vdash \kappa_B \lesssim \kappa_I}{\Delta; \Phi \vdash \Pi_{t::\kappa_I, \kappa_B} \lesssim \Pi_{t::\kappa_B, \kappa_I}} \text{CSK-}\Pi$$

$$\frac{\Delta; \Phi \vdash \kappa_I \lesssim \kappa_B \quad \Delta; \Phi \vdash \tau_I \overset{\kappa_I}{\equiv} \tau_B}{\Delta; \Phi \vdash \text{S}_{\kappa_I}(\tau_I) \lesssim \text{S}_{\kappa_B}(\tau_B)} \text{CSK-?}$$

 $\Delta; \Phi \vdash \tau_I \overset{\kappa}{\equiv} \tau_B$ τ_I is provably equivalent to τ_B at kind κ

$$\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \tau \equiv \tau} \text{EquivAK-Ref1}$$

$$\frac{\Delta; \Phi \vdash \tau_B \overset{\kappa}{\equiv} \tau_I}{\Delta; \Phi \vdash \tau_I \overset{\kappa}{\equiv} \tau_B} \text{EquivAK-Symm}$$

$$\frac{\Delta; \Phi \vdash \tau_I \overset{\kappa}{\equiv} \tau_B \quad \Delta; \Phi \vdash \tau_B \overset{\kappa}{\equiv} \tau_I}{\Delta; \Phi \vdash \tau_I \overset{\kappa}{\equiv} \tau_B} \text{EquivAK-Trans}$$

$$\frac{\Delta; \Phi \vdash \tau_I ::> \kappa_I \quad \Delta; \Phi \vdash \kappa_I \equiv \text{S}_\kappa(\tau_B)}{\Delta; \Phi \vdash \tau_I \overset{\kappa}{\equiv} \tau_B} \text{EquivAK-SKind}$$

$$\frac{\Delta; \Phi \vdash \tau_I :: \Pi_{t::\kappa_I, \kappa_B} \quad \Delta; \Phi \vdash \tau_B :: \Pi_{t::\kappa_I, \kappa_B} \quad \Delta; \Phi, t::\kappa_I \vdash \tau_I \ t \overset{\kappa_B}{\equiv} \tau_B \ t}{\Delta; \Phi \vdash \tau_I \overset{\kappa}{\equiv} \tau_B} \text{EquivAK-}\Pi$$

$$\frac{\Delta; \Phi \vdash \tau_I \overset{\Pi_{t::\kappa_I, \kappa_B}}{\equiv} \tau_B \quad \Delta; \Phi \vdash \tau_B \overset{\kappa_I}{\equiv} \tau_I}{\Delta; \Phi \vdash \tau_I \tau_B \overset{[\tau_B / t] \kappa_B}{\equiv} \tau_B \tau_I} \text{EquivAK-Ap}$$

$$\frac{\Delta; \Phi \vdash \tau_I \overset{\text{Type}}{\equiv} \tau_B \quad \Delta; \Phi \vdash \tau_B \overset{\text{Type}}{\equiv} \tau_I}{\Delta; \Phi \vdash \tau_I \oplus \tau_B \overset{\text{Type}}{\equiv} \tau_B \oplus \tau_I} (2)$$

$$\frac{\Delta; \Phi \vdash \tau_I \overset{\text{S}_\kappa(\tau)}{\equiv} \tau_B \quad (1) \quad \Delta; \Phi \vdash \kappa_I \equiv \kappa_B \quad \Delta; \Phi, t::\kappa_I \vdash \tau_I \ t \overset{\kappa}{\equiv} \tau_B \quad (3)}{\Delta; \Phi \vdash \lambda t::\kappa_I.\tau_I \overset{\Pi_{t::\kappa_I, \kappa}}{\equiv} \lambda t::\kappa_B.\tau_B} (3)$$

$$\frac{\Delta; \Phi \vdash \tau_I \overset{\kappa_I}{\equiv} \tau_B \quad \Delta; \Phi \vdash \kappa_I \equiv \kappa}{\Delta; \Phi \vdash \tau_I \equiv \tau_B} (4)$$

 $\Delta; \Phi \vdash \kappa \text{ OK}$ κ is well formed

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{Type OK}} \text{KWF-Type}$$

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{KHole OK}} \text{KWF-KHole}$$

$$\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \text{S}_\kappa(\tau) \text{ OK}} \text{KWF-SKind}$$

$$\frac{\Delta; \Phi, t::\kappa_I \vdash \kappa_B \text{ OK}}{\Delta; \Phi \vdash \Pi_{t::\kappa_I, \kappa_B} \text{ OK}} \text{KWF-}\Pi$$

 $\Delta; \Phi \vdash \text{OK}$ Context is well formed

$$\frac{}{\cdot, \cdot \vdash \text{OK}} \text{CWF-Nil}$$

$$\frac{t \notin \Delta \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi, t::\kappa \vdash \text{OK}} \text{CWF-TypVar}$$

$$\frac{u \notin \Delta \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta, u::\kappa; \Phi \vdash \text{OK}} \text{CWF-Hole}$$

METATHEORY

Lemma 1 (COK). If $\Delta; \Phi \vdash \mathcal{J}$, then $\Delta; \Phi \vdash \text{OK}$
Proof. By simultaneous induction on derivations.
No interesting cases.

Lemma 2 (Exchange).

If $\Delta; \Phi_1, t_{L1}::\kappa_{L1}, t_{L2}::\kappa_{L2}, \Phi_2 \vdash \mathcal{J}$ and $\Delta; \Phi_1, t_{L2}::\kappa_{L2}, t_{L1}::\kappa_{L1}, \Phi_2 \vdash \text{OK}$, then $\Delta; \Phi_1, t_{L2}::\kappa_{L2}, t_{L1}::\kappa_{L1}, \Phi_2 \vdash \mathcal{J}$
Proof. By induction on derivations.
No interesting cases.

(Only rules with Φ extended in the consequent are interesting, which is only CWF-TypVar, but when \mathcal{J} is CWF, Exchange is identity)

Corollary 3 (Marked-Exchange).

If $\Delta; \Phi, t_{L1}::\kappa_{L1}, t_{L2}::\kappa_{L2} \vdash \mathcal{J}$ and $\Delta; \Phi, t_{L2}::\kappa_{L2}, t_{L1}::\kappa_{L1} \vdash \text{OK}$, then $\Delta; \Phi, t_{L2}::\kappa_{L2}, t_{L1}::\kappa_{L1} \vdash \mathcal{J}$
Lemma 4 (Weakening).

If $\Delta; \Phi \vdash \mathcal{J}$ and $\Delta; \Phi, t_L::\kappa_L \vdash \text{OK}$ and $t_L \notin \mathcal{J}$ and $\forall t \in \kappa_L, t \notin \mathcal{J}$, then $\Delta; \Phi, t_L::\kappa_L \vdash \mathcal{J}$
Lemma 5 (OK-PK). If $\Delta; \Phi \vdash \tau ::> \kappa$, then $\Delta; \Phi \vdash \kappa \text{ OK}$
Lemma 6 (OK-WFaK). If $\Delta; \Phi \vdash \tau :: \kappa$, then $\Delta; \Phi \vdash \kappa \text{ OK}$
Lemma 7 (OK-MatchPi). If $\Delta; \Phi \vdash \kappa \overset{\blacktriangleright}{\Pi} \Pi_{t::\kappa_I, \kappa_B}$, then $\Delta; \Phi \vdash \kappa \text{ OK}$ and $\Delta; \Phi \vdash \Pi_{t::\kappa_I, \kappa_B} \text{ OK}$
Lemma 8 (OK-KEquiv). If $\Delta; \Phi \vdash \kappa_I \equiv \kappa_B$, then $\Delta; \Phi \vdash \kappa_I \text{ OK}$ and $\Delta; \Phi \vdash \kappa_B \text{ OK}$
Lemma 9 (OK-CSK). If $\Delta; \Phi \vdash \kappa_I \lesssim \kappa_B$, then $\Delta; \Phi \vdash \kappa_I \text{ OK}$ and $\Delta; \Phi \vdash \kappa_B \text{ OK}$
Lemma 10 (OK-EquivAK). If $\Delta; \Phi \vdash \tau_I \overset{\kappa}{\equiv} \tau_B$, then $\Delta; \Phi \vdash \tau_I :: \kappa$ and $\Delta; \Phi \vdash \tau_B :: \kappa$ and $\Delta; \Phi \vdash \kappa \text{ OK}$
Lemma 11 (OK-Substitution).

If $\Delta; \Phi \vdash \tau_{L1}::\kappa_{L1}$ and $\Delta; \Phi, t_L::\kappa_{L1} \vdash \kappa_{L2} \text{ OK}$, then $\Delta; \Phi \vdash [\tau_{L1} / t_L] \kappa_{L2} \text{ OK}$
(induction on $\Delta; \Phi, t_L::\kappa_{L1} \vdash \kappa_{L2} \text{ OK}$)

Lemma 12 (K-Substitution).

If $\Delta; \Phi \vdash \tau_{L1}::\kappa_{L1}$ and $\Delta; \Phi, t_L::\kappa_{L1} \vdash \tau_{L2}::\kappa_{L2}$, then $\Delta; \Phi \vdash [\tau_{L1} / t_L] \tau_{L2}::[\tau_{L1} / t_L] \kappa_{L2}$
(induction on $\Delta; \Phi, t_L::\kappa_{L1} \vdash \tau_{L2}::\kappa_{L2}$)

Proof. By simultaneous induction on derivations.

The interesting cases per lemma:

(PoS = premiss of subderivation)

Weakening

$\frac{\frac{\overline{\Delta; \Phi, \underline{t_L} :: \kappa_L} \vdash \text{OK}}{\Delta; \Phi \vdash \kappa_L \text{ OK}} \text{PoS} \quad \frac{\overline{\Delta; \Phi, l :: \kappa_I} \vdash \tau ::> \kappa_g}{\Delta; \Phi, l :: \kappa_I \vdash \text{OK}} \text{premiss} \text{COK}}{\Delta; \Phi, l :: \kappa_I \vdash \kappa_L \text{ OK}} \text{Weakening}$	$\frac{\frac{\overline{\Delta; \Phi, \underline{t_L} :: \kappa_L} \vdash \text{OK}}{\underline{t_L} \notin \Phi} \text{PoS} \quad \frac{\overline{\Delta; \Phi, \underline{t_L} :: \kappa_L} \vdash \text{OK}}{\underline{t_L} \notin \Phi} \text{PoS} \quad \frac{\overline{t_L \notin \mathcal{J}} \text{IH} \quad \overline{t \in \mathcal{J}}}{t_L \neq t} \quad \frac{\overline{t_L \notin \mathcal{J}} \text{IH}}{\underline{t_L} \notin \kappa_I}}{\underline{t_L} \notin \Phi, l :: \kappa_I} \text{CWF-TypVar}$	$\frac{\frac{\overline{\Delta; \underline{t_L} :: \kappa_I} \vdash \text{OK}}{\Delta; \Phi \vdash \kappa_I \text{ OK}} \text{premiss} \text{PoS} \quad \frac{\overline{\Delta; \Phi, \underline{t_L} :: \kappa_L} \vdash \text{OK}}{\Delta; \Phi, \underline{t_L} :: \kappa_L \vdash \kappa_I \text{ OK}} \text{premiss} \text{Weakening}}{\Delta; \Phi, \underline{t_L} :: \kappa_L \vdash \kappa_I \text{ OK}} \text{Weakening}$	$\frac{\frac{\overline{\Delta; \Phi, l :: \kappa_I} \vdash \tau ::> \kappa_g}{\Delta; \Phi, l :: \kappa_I \vdash \text{OK}} \text{premiss} \text{COK} \quad \frac{\overline{t_L \notin \mathcal{J}} \text{IH} \quad \overline{t \in \mathcal{J}}}{t \neq t_L} \quad \frac{\overline{\forall i \in \kappa_L, i \notin \mathcal{J}} \text{IH} \quad \overline{t \in \mathcal{J}}}{t \notin \kappa_L}}{\Delta; \Phi, \underline{t_L} :: \kappa_L, l :: \kappa_I \vdash \text{OK}} \text{CWF-TypVar}$	$\frac{\Delta; \Phi, \underline{t_L} :: \kappa_L, l :: \kappa_I \vdash \tau ::> \kappa_g}{\Delta; \Phi, \underline{t_L} :: \kappa_L, l :: \kappa_I \vdash \tau ::> \kappa_g} \text{Marked-Exchange}$	$\frac{\Delta; \Phi, \underline{t_L} :: \kappa_L, l :: \kappa_I \vdash \tau ::> \kappa_g}{\Delta; \Phi, \underline{t_L} :: \kappa_L \vdash \lambda l :: \kappa_I. \tau ::> \mathbf{S}_{\text{new } \kappa_I, \kappa_g}(\lambda l :: \kappa_I. \tau)} \text{PK-}\lambda$
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Weakening.			$\Delta; \Phi \vdash \kappa_I$ OK $\Delta; \Phi, t_L :: \kappa_L \vdash \text{OK}$ $\Delta; \Phi, t_L :: \kappa_L \vdash \kappa_I$ OK $\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_I \vdash \text{OK}$ $\Delta; \Phi, t :: \kappa_I, t_L :: \kappa_L \vdash \text{OK}$ $\Delta; \Phi, t :: \kappa_I, t_L :: \kappa_L \vdash \tau ::> \kappa_B$ $\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_I \vdash \tau ::> \kappa_B$ $\Delta; \Phi, t_L :: \kappa_L \vdash \lambda t :: \kappa_I. \tau ::> \mathbf{S}_{\text{H}_{\text{IO}, \kappa_I, \kappa_B}}(\lambda t :: \kappa_I. \tau)$ $\Delta; \Phi \vdash \mathbf{bse} :: \mathbf{S_{Type}}(\mathbf{bse})$ $\Delta; \Phi \vdash \mathbf{bse} :: \mathbf{Type}$ $\Delta; \Phi \vdash \mathbf{S_{Type}}(\mathbf{bse})$ OK $\Delta; \Phi \vdash \text{OK}$	by subderivation premiss by IH by Weakening on subderivation premiss by CWF-TypVar by ? by Weakening on premiss by Marked-Exchange by PK- λ by (9) by (10) by (43) by premiss bad by (10) by (43)
OK-PK.	PK-Base			
	*			
	*			
OK-WFaK.	PK-Ap (12)			
	*			
OK-KEquiv.	(22)			
OK-Substitution.	(41)			
	*			
	*			
	(43)			
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	*			
	*			
	*		$\Delta; \Phi \vdash \tau_B :: \kappa$ $\Delta; \Phi \vdash \mathbf{S}_{\kappa}(\tau_B)$ OK $\Delta; \Phi \vdash \tau \ell ::> \kappa$ $\Delta; \Phi, t_L :: \kappa_{L,I} \vdash \text{OK}$ $\Delta; \Phi \vdash \kappa_{L,I}$ OK $\Delta; \Phi \vdash \text{OK}$ $\Delta; \Phi \vdash [\tau_L / t_L] \mathbf{Type}$ OK $\Delta; \Phi, t_L :: \kappa_{L,I} \vdash \tau :: \kappa$ $\Delta; \Phi, t_L :: \kappa_{L,I} \vdash \text{OK}$ $\Delta; \Phi \vdash \kappa_{L,I}$ OK $\Delta; \Phi \vdash \text{OK}$ $\Delta; \Phi \vdash [\tau_{L,I} / t_L] \tau :: [\tau_{L,I} / t_L] \kappa$ $\Delta; \Phi \vdash [\tau_L / t_L] \mathbf{S}_{\kappa}(\tau)$ OK	premiss (41) by subderivation premiss (46) by OK-KWF by (41) and degenerate subst premiss (43) by OK-WFaK by subderivation premiss (46) by OK-KWF by K-Substitution on premiss by (43)

Lemma 13 (PK-Unicity). *If $\Delta; \Phi \vdash \tau_L ::> \kappa_{L,1}$ and $\Delta; \Phi \vdash \tau_L ::> \kappa_{L,2}$ then $\kappa_{L,1}$ is $\kappa_{L,2}$*

Lemma 14. *If $\Delta; \Phi \vdash \tau ::> \kappa_I$ and $\Delta; \Phi \vdash \tau :: \kappa_B$, then $\Delta; \Phi \vdash \kappa_I \lesssim \kappa_B$*

Lemma 15. *If $\Delta; \Phi \vdash \kappa_I \lesssim \mathbf{S}_{\kappa_B}(\tau)$, then $\Delta; \Phi \vdash \kappa_I \lesssim \kappa_B$*

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