

# Hazel Phi: 11-type-constructors

July 31, 2021

## SYNTAX

---

Kind	$\kappa$	$::=$	<b>Type</b>   <b>KHole</b>   $\mathbf{S}_{\kappa}(\tau)$   $\Pi_{t::\kappa_1}.\kappa_2$
User Types	$\hat{\tau}$	$::=$	$t$   <b>bse</b>   $\tau_1 \oplus \tau_2$   $\langle \rangle^u$   $\langle \hat{\tau} \rangle^u$   $\lambda t::\mathbf{Type}.\hat{\tau}$   $\tau_1 \tau_2$
Internal Types	$\tau$	$::=$	$t$   <b>bse</b>   $\tau_1 \oplus \tau_2$   $\langle \rangle^u$   $\langle \tau \rangle^u$   $\lambda t::\kappa.\tau$   $\tau_1 \tau_2$
Base Types	<b>bse</b>	$::=$	<b>Int</b>   <b>Float</b>   <b>Bool</b>
BinOp	$\oplus$	$::=$	$\times$   $+$   $\rightarrow$
Type Pattern			
User Expression			
Internal Expression			

## DECLARATIVES

---

$\boxed{\Delta; \Phi \vdash \tau ::> \kappa}$   $\tau$  has principal (well formed) kind  $\kappa$

$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \mathbf{bse} ::> \mathbf{S}_{\mathbf{Type}}(\mathbf{bse})} \text{PK-Base}$	$\frac{\Delta; \Phi_1, t::\kappa, \Phi_2 \vdash \text{OK}}{\Delta; \Phi \vdash t ::> \mathbf{S}_{\kappa}(t)} \text{PK-Var}$
$\frac{\Delta; \Phi \vdash \tau_1 :: \mathbf{Type} \quad \Delta; \Phi \vdash \tau_2 :: \mathbf{Type}}{\Delta; \Phi \vdash \tau_1 \oplus \tau_2 ::> \mathbf{S}_{\mathbf{Type}}(\tau_1 \oplus \tau_2)} \text{PK-}\oplus$	$\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \langle \rangle^u ::> \mathbf{S}_{\kappa}(\langle \rangle^u)} \text{PK-EHole}$
$\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \text{OK} \quad \Delta; \Phi \vdash \tau :: \kappa_1}{\Delta; \Phi \vdash \langle \tau \rangle^u ::> \mathbf{S}_{\kappa}(\langle \tau \rangle^u)} \text{PK-NEHole}$	$\frac{\Delta_1, u::\kappa, \Delta_2; \Phi \vdash \text{OK} \quad t \notin \Phi}{\Delta; \Phi \vdash \langle t \rangle^u ::> \mathbf{S}_{\kappa}(\langle t \rangle^u)} \text{PK-Unbound}$
$\frac{\Delta; \Phi, t::\kappa_1 \vdash \tau ::> \kappa_2}{\Delta; \Phi \vdash \lambda t::\kappa_1.\tau ::> \mathbf{S}_{\Pi_{t::\kappa_1}.\kappa_2}}(\lambda t::\kappa_1.\tau)} \text{PK-}\lambda$	
$\frac{\Delta; \Phi \vdash \tau_1 ::> \kappa \quad \Delta; \Phi \vdash \kappa \blacktriangleright \Pi_{t::\kappa_1}.\kappa_2 \quad \Delta; \Phi \vdash \tau_2 :: \kappa_1}{\Delta; \Phi \vdash \tau_1 \tau_2 ::> [\tau_2/t]\kappa_2} \text{PK-Ap}$	

$\Delta; \Phi \vdash \tau :: \kappa$   $\tau$  is well formed at kind  $\kappa$

$$\frac{\Delta; \Phi \vdash \tau :: > \mathbf{S}_{\kappa}(\tau)}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-1} \quad \frac{\Delta; \Phi \vdash \tau :: > \kappa_1 \quad \Delta; \Phi \vdash \kappa_1 \lesssim \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Subsump}$$

$$\frac{\Delta; \Phi \vdash \tau :: > \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Reit} \quad \frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \tau :: \mathbf{S}_{\kappa}(\tau)} \text{WFaK-Self}$$

$$\frac{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_3} \cdot \kappa_4 \quad \Delta; \Phi \vdash \Pi_{t::\kappa_3} \cdot \kappa_4 \lesssim \Pi_{t::\kappa_1} \cdot \kappa_2}{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_1} \cdot \kappa_2} \text{WFaK-PCSKTrans}$$

$$\frac{\Delta; \Phi \vdash \tau :: \mathbf{S}_{\kappa}(\tau_1) \quad \Delta; \Phi \vdash \tau_1 :: \kappa}{\Delta; \Phi \vdash \tau :: \kappa} \text{WFaK-Flatten}$$

$\Delta; \Phi \vdash \kappa \Pi \Pi_{t::\kappa_1} \cdot \kappa_2$   $\kappa$  has matched  $\Pi$ -kind  $\Pi_{t::\kappa_1} \cdot \kappa_2$

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{KHole} \Pi \Pi_{t::\text{KHole}} \cdot \text{KHole}} \Pi\text{-KHole} \quad \frac{\Delta; \Phi \vdash \kappa \equiv \mathbf{S}_{\text{KHole}}(\tau)}{\Delta; \Phi \vdash \kappa \Pi \Pi_{t::\mathbf{S}_{\text{KHole}}(\tau)} \cdot \mathbf{S}_{\text{KHole}}(\tau \ t)} \Pi\text{-SKHole}$$

$$\frac{\Delta; \Phi \vdash \kappa \equiv \Pi_{t::\kappa_1} \cdot \kappa_2}{\Delta; \Phi \vdash \kappa \Pi \Pi_{t::\kappa_1} \cdot \kappa_2} \Pi\text{-}\Pi$$

$\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2$   $\kappa_1$  is equivalent to  $\kappa_2$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \equiv \kappa} \text{KEquiv-Ref1} \quad \frac{\Delta; \Phi \vdash \kappa_2 \equiv \kappa_1}{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2} \text{KEquiv-Symm}$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_3 \quad \Delta; \Phi \vdash \kappa_3 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2} \text{KEquiv-Trans}$$

$$\frac{\Delta; \Phi \vdash \tau :: \mathbf{S}_{\kappa}(\tau_1)}{\Delta; \Phi \vdash \mathbf{S}_{\mathbf{S}_{\kappa}(\tau_1)}(\tau) \equiv \mathbf{S}_{\kappa}(\tau_1)} \text{KEquiv-SKind}_{\text{SKind}}$$

$$\frac{\Delta; \Phi \vdash \tau :: \Pi_{t::\kappa_1} \cdot \kappa_2}{\Delta; \Phi \vdash \mathbf{S}_{\Pi_{t::\kappa_1} \cdot \kappa_2}(\tau) \equiv \Pi_{t_1::\kappa_1} \cdot \mathbf{S}_{[t_1/t]\kappa_2}(\tau \ t_1)} \text{KEquiv-SKind}_{\Pi}$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \quad \Delta; \Phi, t::\kappa_1 \vdash \kappa_3 \equiv \kappa_4}{\Delta; \Phi \vdash \Pi_{t::\kappa_1} \cdot \kappa_2 \equiv \Pi_{t::\kappa_3} \cdot \kappa_4} \text{KEquiv-}\Pi$$

$$\frac{\Delta; \Phi \vdash \tau_1 \equiv^{\kappa_1} \tau_2 \quad \Delta; \Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta; \Phi \vdash \mathbf{S}_{\kappa_1}(\tau_1) \equiv \mathbf{S}_{\kappa_2}(\tau_2)} \text{KEquiv-SKind}$$

$\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$   $\kappa_1$  is a consistent subkind of  $\kappa_2$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \text{KHole} \lesssim \kappa} \text{CSK-KHoleL} \quad \frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \text{KHole}} \text{CSK-KHoleR}$$

$$\frac{\Delta; \Phi \vdash \text{SKHole}(\tau) \text{ OK} \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \text{SKHole}(\tau) \lesssim \kappa} \text{CSK-SKind}_{\text{KHoleL}}$$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK} \quad \Delta; \Phi \vdash \text{SKHole}(\tau) \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \text{SKHole}(\tau)} \text{CSK-SKind}_{\text{KHoleR}}$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \text{CSK-KEquiv} \quad \frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_3 \quad \Delta; \Phi \vdash \kappa_3 \lesssim \kappa_4 \quad \Delta; \Phi \vdash \kappa_4 \equiv \kappa_2}{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2} \text{CSK-Normal}$$

$$\frac{\Delta; \Phi \vdash \text{S}_\kappa(\tau) \text{ OK}}{\Delta; \Phi \vdash \text{S}_\kappa(\tau) \lesssim \kappa} \text{CSK-SKind} \quad \frac{\Delta; \Phi \vdash \kappa_3 \lesssim \kappa_1 \quad \Delta; \Phi, t::\kappa_3 \vdash \kappa_2 \lesssim \kappa_4}{\Delta; \Phi \vdash \Pi_{t::\kappa_1} \cdot \kappa_2 \lesssim \Pi_{t::\kappa_3} \cdot \kappa_4} \text{CSK-}\Pi$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2 \quad \Delta; \Phi \vdash \tau_1 \stackrel{\kappa_1}{\equiv} \tau_2}{\Delta; \Phi \vdash \text{S}_{\kappa_1}(\tau_1) \lesssim \text{S}_{\kappa_2}(\tau_2)} \text{CSK-?}$$

$\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2$   $\tau_1$  is provably equivalent to  $\tau_2$  at kind  $\kappa$

$$\frac{\Delta; \Phi \vdash \tau::\kappa}{\Delta; \Phi \vdash \tau \equiv \tau} \text{EquivAK-Ref1} \quad \frac{\Delta; \Phi \vdash \tau_2 \stackrel{\kappa}{\equiv} \tau_1}{\Delta; \Phi \vdash \tau_1 \equiv \tau_2} \text{EquivAK-Symm}$$

$$\frac{\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_3 \quad \Delta; \Phi \vdash \tau_3 \stackrel{\kappa}{\equiv} \tau_1}{\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2} \text{EquivAK-Trans}$$

$$\frac{\Delta; \Phi \vdash \tau_1 ::> \kappa_1 \quad \Delta; \Phi \vdash \kappa_1 \equiv \text{S}_\kappa(\tau_2)}{\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2} \text{EquivAK-SKind}$$

$$\frac{\Delta; \Phi \vdash \tau_1 :: \Pi_{t::\kappa_1} \cdot \kappa_3 \quad \Delta; \Phi \vdash \tau_2 :: \Pi_{t::\kappa_1} \cdot \kappa_4 \quad \Delta; \Phi, t::\kappa_1 \vdash \tau_1 \stackrel{\kappa_2}{t} \equiv \tau_2 \quad t}{\Delta; \Phi \vdash \tau_1 \stackrel{\Pi_{t::\kappa_1} \cdot \kappa_2}{\equiv} \tau_2} \text{EquivAK-}\Pi$$

$$\frac{\Delta; \Phi \vdash \tau_1 \stackrel{\Pi_{t::\kappa_1} \cdot \kappa_2}{\equiv} \tau_3 \quad \Delta; \Phi \vdash \tau_2 \stackrel{\kappa_1}{\equiv} \tau_4}{\Delta; \Phi \vdash \tau_1 \tau_2 \stackrel{[\tau_2/t]\kappa_2}{\equiv} \tau_3 \tau_4} \text{EquivAK-Ap}$$

$$\frac{\Delta; \Phi \vdash \tau_1 \stackrel{\text{S}_\kappa(\tau)}{\equiv} \tau_2}{\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2} (1)$$

$$\frac{\Delta; \Phi \vdash \tau_1 \stackrel{\text{Type}}{\equiv} \tau_3 \quad \Delta; \Phi \vdash \tau_2 \stackrel{\text{Type}}{\equiv} \tau_4}{\Delta; \Phi \vdash \tau_1 \oplus \tau_2 \stackrel{\text{Type}}{\equiv} \tau_3 \oplus \tau_4} (2)$$

$$\frac{\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2 \quad \Delta; \Phi, t::\kappa_1 \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2}{\Delta; \Phi \vdash \lambda t::\kappa_1. \tau_1 \stackrel{\Pi_{t::\kappa_1} \cdot \kappa}{\equiv} \lambda t::\kappa_2. \tau_2} (3)$$

$$\frac{\Delta; \Phi \vdash \tau_1 \stackrel{\kappa_1}{\equiv} \tau_2 \quad \Delta; \Phi \vdash \kappa_1 \equiv \kappa}{\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2} (4)$$

$\boxed{\Delta; \Phi \vdash \kappa \text{ OK}}$   $\kappa$  is well formed

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{Type OK}} \text{KWF-Type}$$

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{KHole OK}} \text{KWF-KHole}$$

$$\frac{\Delta; \Phi \vdash \tau :: \kappa}{\Delta; \Phi \vdash \text{S}_\kappa(\tau) \text{ OK}} \text{KWF-SKind}$$

$$\frac{\Delta; \Phi, t :: \kappa_1 \vdash \kappa_2 \text{ OK}}{\Delta; \Phi \vdash \Pi_{t :: \kappa_1} \cdot \kappa_2 \text{ OK}} \text{KWF-}\Pi$$

$\boxed{\Delta; \Phi \vdash \text{OK}}$  Context is well formed

$$\frac{}{\cdot; \cdot \vdash \text{OK}} \text{CWF-Nil}$$

$$\frac{t \notin \Phi \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi, t :: \kappa \vdash \text{OK}} \text{CWF-TypVar}$$

$$\frac{u \notin \Delta \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta, u :: \kappa; \Phi \vdash \text{OK}} \text{CWF-Hole}$$

## METATHEORY

---

subderivation preserving inferences:

- premiss
- COK (Context OK)
- PoS (premiss of subderivation)

**Lemma 1** (COK). *If  $\Delta; \Phi \vdash \mathcal{J}$ , then  $\Delta; \Phi \vdash \text{OK}$  in a subderivation (where  $\Delta; \Phi \vdash \mathcal{J} \neq \Delta; \Phi \vdash \text{OK}$ )*

*Proof.* By induction on derivations.

No interesting cases. □

**Lemma 2** (Exchange).

*If  $\Delta; \Phi_1, t_{L1} :: \kappa_{L1}, t_{L2} :: \kappa_{L2}, \Phi_2 \vdash \mathcal{J}$  and  $\Delta; \Phi_1, t_{L2} :: \kappa_{L2}, t_{L1} :: \kappa_{L1}, \Phi_2 \vdash \text{OK}$ , then  $\Delta; \Phi_1, t_{L2} :: \kappa_{L2}, t_{L1} :: \kappa_{L1}, \Phi_2 \vdash \mathcal{J}$*

*Proof.* By induction on derivations.

No interesting cases.

(Only rules with  $\Phi$  extended in the consequent are interesting, which is only CWF-TypVar, but when  $\mathcal{J}$  is CWF, Exchange is identity) □

**Corollary 3** (Marked-Exchange).

*If  $\Delta; \Phi, \underline{t_{L1} :: \kappa_{L1}}, \underline{t_{L2} :: \kappa_{L2}} \vdash \mathcal{J}$  and  $\Delta; \Phi, \underline{t_{L2} :: \kappa_{L2}}, \underline{t_{L1} :: \kappa_{L1}} \vdash \text{OK}$ , then  $\Delta; \Phi, \underline{t_{L2} :: \kappa_{L2}}, \underline{t_{L1} :: \kappa_{L1}} \vdash \mathcal{J}$*

*Proof.* Exchange when  $\Phi_2 = \cdot$  □

**Lemma 4** (Weakening).

*If  $\Delta; \Phi \vdash \mathcal{J}$  and  $\Delta; \Phi, \underline{t_L :: \kappa_L} \vdash \text{OK}$  and  $t_L \notin \mathcal{J}$  and  $\forall t \in \kappa_L, t \notin \mathcal{J}$ , then  $\Delta; \Phi, \underline{t_L :: \kappa_L} \vdash \mathcal{J}$*

*Proof.* see addendum □

**Lemma 5** (K-Substitution).

*If  $\Delta; \Phi \vdash \tau_{L1} :: \kappa_{L1}$  and  $\Delta; \Phi, \underline{t_L :: \kappa_{L1}} \vdash \tau_{L2} :: \kappa_{L2}$ , then  $\Delta; \Phi \vdash [\tau_{L1}/t_L] \tau_{L2} :: [\tau_{L1}/t_L] \kappa_{L2}$  (induction on  $\Delta; \Phi, \underline{t_L :: \kappa_{L1}} \vdash \tau_{L2} :: \kappa_{L2}$ )*

**Lemma 6** (OK-Substitution).

*If  $\Delta; \Phi \vdash \tau_L :: \kappa_{L1}$  and  $\Delta; \Phi, \underline{t_L :: \kappa_{L1}} \vdash \kappa_{L2} \text{ OK}$ , then  $\Delta; \Phi \vdash [\tau_L/t_L] \kappa_{L2} \text{ OK}$  (induction on  $\Delta; \Phi, \underline{t_L :: \kappa_{L1}} \vdash \kappa_{L2} \text{ OK}$ )*

**Lemma 7** (OK-PK). *If  $\Delta; \Phi \vdash \tau :: \kappa$ , then  $\Delta; \Phi \vdash \kappa \text{ OK}$*

**Lemma 8** (OK-WFaK). *If  $\Delta; \Phi \vdash \tau :: \kappa$ , then  $\Delta; \Phi \vdash \kappa$  OK*

**Lemma 9** (OK-MatchPi). *If  $\Delta; \Phi \vdash \kappa \blacktriangleright \Pi_{t::\kappa_1} \kappa_2$ , then  $\Delta; \Phi \vdash \kappa$  OK and  $\Delta; \Phi \vdash \Pi_{t::\kappa_1} \kappa_2$  OK*

**Lemma 10** (OK-KEquiv). *If  $\Delta; \Phi \vdash \kappa_1 \equiv \kappa_2$ , then  $\Delta; \Phi \vdash \kappa_1$  OK and  $\Delta; \Phi \vdash \kappa_2$  OK*

**Lemma 11** (OK-CSK). *If  $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$ , then  $\Delta; \Phi \vdash \kappa_1$  OK and  $\Delta; \Phi \vdash \kappa_2$  OK*

**Lemma 12** (OK-EquivAK). *If  $\Delta; \Phi \vdash \tau_1 \stackrel{\kappa}{\equiv} \tau_2$ , then  $\Delta; \Phi \vdash \tau_1 :: \kappa$  and  $\Delta; \Phi \vdash \tau_2 :: \kappa$  and  $\Delta; \Phi \vdash \kappa$  OK*

*Proof.* see addendum

□

*Proof.*

#### Weakening

By induction on derivations.

Note: When applying Weakening in the induction, check that the left premiss is always a subderivation, and check variable exclusion conditions are satisfied (usually checked elsewhere in the derivation).

$$\begin{array}{c}
\frac{\frac{\frac{\Delta; \Phi, t_L :: \kappa_L \vdash \text{OK}}{t_L \notin \Phi} \text{IH} \quad \frac{\frac{\frac{\overline{t_L \notin \mathcal{J}} \text{IH} \quad \overline{t \in \mathcal{J}}}{t_L \neq t} \quad \frac{\overline{t_L \notin \kappa_1} \text{IH}}{t_L \notin \kappa_1}}{t_L \notin \Phi, t :: \kappa_1} \text{PoS} \quad \frac{\frac{\Delta; \Phi, t_L :: \kappa_L \vdash \text{OK}}{\Delta; \Phi \vdash \kappa_L \text{ OK}} \text{IH} \quad \frac{\frac{\Delta; \Phi, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{premiss} \quad \frac{\Delta; \Phi, t :: \kappa_1 \vdash \text{OK}}{\Delta; \Phi, t :: \kappa_1 \vdash \kappa_L \text{ OK}} \text{COK}}{\Delta; \Phi, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{Weakening} \\
\frac{\Delta; \Phi, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{premiss} \quad \frac{\Delta; \Phi, t :: \kappa_1, t_L :: \kappa_L \vdash \text{OK}}{\Delta; \Phi, t :: \kappa_1, t_L :: \kappa_L \vdash \tau ::> \kappa_2} \text{CWF-TypVar} \\
\frac{\Delta; \Phi, t :: \kappa_1, t_L :: \kappa_L \vdash \tau ::> \kappa_2} \text{Weakening} \\
\frac{\frac{\frac{\Delta; \Phi, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{premiss} \quad \frac{\frac{\Delta; \Phi, t :: \kappa_1 \vdash \text{OK}}{t \notin \Phi} \text{COK} \quad \frac{\frac{\overline{t_L \notin \mathcal{J}} \text{IH} \quad \overline{t \in \mathcal{J}}}{t \neq t_L} \quad \frac{\frac{\overline{\forall t \in \kappa_L, t \notin \mathcal{J}} \text{IH} \quad \overline{t \in \mathcal{J}}}{t \notin \kappa_L} \text{PoS}}{t \notin \Phi, t_L :: \kappa_L} \text{PoS} \quad \frac{\frac{\frac{\Delta; \Phi, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{premiss} \quad \frac{\Delta; \Phi, t :: \kappa_1 \vdash \text{OK}}{\Delta; \Phi \vdash \kappa_1 \text{ OK}} \text{COK} \quad \frac{\Delta; \Phi, t_L :: \kappa_L \vdash \text{OK}}{\Delta; \Phi, t_L :: \kappa_L \vdash \kappa_1 \text{ OK}} \text{IH}}{\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_1 \vdash \text{OK}} \text{Weakening} \\
\frac{\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_1 \vdash \text{OK}}{\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{CWF-TypVar} \\
\frac{\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{Marked-Exchange} \\
\frac{\Delta; \Phi, t_L :: \kappa_L \vdash \lambda t :: \kappa_1. \tau ::> \mathbf{S}_{\Pi_{t :: \kappa_1}, \kappa_2}(\lambda t :: \kappa_1. \tau)}{\Delta; \Phi, t_L :: \kappa_L \vdash \tau ::> \kappa_2} \text{PK-}\lambda \\
\frac{\frac{\frac{\frac{\Delta; \Phi, t_L :: \kappa_L \vdash \text{OK}}{t_L \notin \Phi} \text{IH} \quad \frac{\frac{\overline{t_L \notin \mathcal{J}} \text{IH} \quad \overline{t \in \mathcal{J}}}{t_L \neq t} \quad \frac{\overline{t_L \notin \kappa_1} \text{IH}}{t_L \notin \kappa_1}}{t_L \notin \Phi, t :: \kappa_1} \text{PoS} \quad \frac{\frac{\Delta; \Phi, t_L :: \kappa_L \vdash \text{OK}}{\Delta; \Phi \vdash \kappa_L \text{ OK}} \text{IH} \quad \frac{\frac{\Delta; \Phi, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{premiss} \quad \frac{\Delta; \Phi, t :: \kappa_1 \vdash \text{OK}}{\Delta; \Phi, t :: \kappa_1 \vdash \kappa_L \text{ OK}} \text{COK}}{\Delta; \Phi, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{Weakening} \\
\frac{\Delta; \Phi, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{premiss} \quad \frac{\Delta; \Phi, t :: \kappa_1, t_L :: \kappa_L \vdash \text{OK}}{\Delta; \Phi, t :: \kappa_1, t_L :: \kappa_L \vdash \tau ::> \kappa_2} \text{CWF-TypVar} \\
\frac{\Delta; \Phi, t :: \kappa_1, t_L :: \kappa_L \vdash \tau ::> \kappa_2} \text{Weakening} \\
\frac{\frac{\frac{\Delta; \Phi, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{premiss} \quad \frac{\frac{\Delta; \Phi, t :: \kappa_1 \vdash \text{OK}}{t \notin \Phi} \text{COK} \quad \frac{\frac{\overline{t_L \notin \mathcal{J}} \text{IH} \quad \overline{t \in \mathcal{J}}}{t \neq t_L} \quad \frac{\frac{\overline{\forall t \in \kappa_L, t \notin \mathcal{J}} \text{IH} \quad \overline{t \in \mathcal{J}}}{t \notin \kappa_L} \text{PoS}}{t \notin \Phi, t_L :: \kappa_L} \text{PoS} \quad \frac{\frac{\frac{\Delta; \Phi, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{premiss} \quad \frac{\Delta; \Phi, t :: \kappa_1 \vdash \text{OK}}{\Delta; \Phi \vdash \kappa_1 \text{ OK}} \text{COK} \quad \frac{\Delta; \Phi, t_L :: \kappa_L \vdash \text{OK}}{\Delta; \Phi, t_L :: \kappa_L \vdash \kappa_1 \text{ OK}} \text{IH}}{\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_1 \vdash \text{OK}} \text{Weakening} \\
\frac{\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_1 \vdash \text{OK}}{\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{CWF-TypVar} \\
\frac{\Delta; \Phi, t_L :: \kappa_L, t :: \kappa_1 \vdash \tau ::> \kappa_2} \text{Marked-Exchange} \\
\frac{\Delta; \Phi, t_L :: \kappa_L \vdash \tau ::> \kappa_2} \text{Marked-Exchange} \\
\frac{\Delta; \Phi, t_L :: \kappa_L \vdash \tau ::> \kappa_2} \text{KEquiv-II} \quad \frac{}{\text{sim}} \text{CSK-II} \quad \frac{}{\text{sim}} \text{EquivAK-II} \quad \frac{}{\text{sim}} \text{KWF-II}
\end{array}$$

#### O?K-\*

By simultaneous induction on derivations.

The interesting cases per lemma:

**K-Substitution** by type size??

**OK-Substitution**

**OK-PK**

$$\begin{array}{c}
\frac{\frac{\Delta; \Phi \vdash \mathbf{bse} ::> \text{SType}(\mathbf{bse}) \text{ premiss} \quad \frac{\Delta; \Phi \vdash \mathbf{bse} :: \text{Type}}{\Delta; \Phi \vdash \text{SType}(\mathbf{bse}) \text{ OK}} \text{WFaK-1} \\
\frac{}{\Delta; \Phi \vdash \text{SType}(\mathbf{bse}) \text{ OK}} \text{KWF-SKind} \quad \frac{}{\Delta; \Phi \vdash [\tau_2 / t] \kappa_2 \text{ OK}} \text{OK-Substitution}
\end{array}$$

**OK-WFaK**

□

*Proof.* By simultaneous induction on derivations.

The interesting cases per lemma:

OK-PK.	PK-Base	$\Delta; \Phi \vdash \mathbf{bse} :: \mathbf{S_{Type}}(\mathbf{bse})$	by (9)
		$\Delta; \Phi \vdash \mathbf{bse} :: \mathbf{Type}$	by (10)
	*	$\Delta; \Phi \vdash \mathbf{S_{Type}}(\mathbf{bse}) \text{ OK}$	by (43)
	*	$\Delta; \Phi \vdash \text{OK}$	by premiss
	PK-Ap		bad
OK-WFaK.	(12)	$\Delta; \Phi \vdash \tau_2 :: \kappa$	by (10)
	*	$\Delta; \Phi \vdash \mathbf{S_{\kappa}}(\tau_2) \text{ OK}$	by (43)
OK-KEquiv.	(22)	$\Delta; \Phi \vdash \tau \ t :: > \kappa$	
OK-Substitution.	(41)	$\Delta; \Phi, t_L :: \kappa_{L1} \vdash \text{OK}$	premiss (41)
		$\Delta; \Phi \vdash \kappa_{L1} \text{ OK}$	by subderivation premiss (46)
	*	$\Delta; \Phi \vdash \text{OK}$	by OK-KWF
	*	$\Delta; \Phi \vdash [\tau_L / t_L] \mathbf{Type} \text{ OK}$	by (41) and degenerate subst
	(43)	$\Delta; \Phi, t_L :: \kappa_{L1} \vdash \tau :: \kappa$	premiss (43)
		$\Delta; \Phi, t_L :: \kappa_{L1} \vdash \text{OK}$	by OK-WFaK
		$\Delta; \Phi \vdash \kappa_{L1} \text{ OK}$	by subderivation premiss (46)
	*	$\Delta; \Phi \vdash \text{OK}$	by OK-KWF
		$\Delta; \Phi \vdash [\tau_{L1} / t_L] \tau :: [\tau_{L1} / t_L] \kappa$	by K-Substitution on premiss
	*	$\Delta; \Phi \vdash [\tau_L / t_L] \mathbf{S_{\kappa}}(\tau) \text{ OK}$	by (43)

□

**Lemma 13** (PK-Unicity). *If  $\Delta; \Phi \vdash \tau_L :: > \kappa_{L1}$  and  $\Delta; \Phi \vdash \tau_L :: > \kappa_{L2}$  then  $\kappa_{L1}$  is  $\kappa_{L2}$*

**Lemma 14.** *If  $\Delta; \Phi \vdash \tau :: > \kappa_1$  and  $\Delta; \Phi \vdash \tau :: \kappa_2$ , then  $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$*

**Lemma 15.** *If  $\Delta; \Phi \vdash \kappa_1 \lesssim \mathbf{S_{\kappa_2}}(\tau)$ , then  $\Delta; \Phi \vdash \kappa_1 \lesssim \kappa_2$*

ELABORATION

TODO