

SYNTAX

Kind	κ	$::=$	Type KHole $S_\kappa(\tau)$ $\Pi_{t::\kappa_I, \kappa_B}$
User Types	$\hat{\tau}$	$::=$	t bse $\tau_I \oplus \tau_B$ \emptyset^u $\langle \hat{\tau} \rangle^u$ $\lambda t::\text{Type}.\hat{\tau}$ $\tau_I \ \tau_B$
Internal Types	τ	$::=$	t bse $\tau_I \oplus \tau_B$ \emptyset^u $\langle \hat{\tau} \rangle^u$ $\langle \emptyset \rangle^u$ $\lambda t::\kappa.\tau$ $\tau_I \ \tau_B$
Base Types	bse	$::=$	Int Float Bool
BinOp	\oplus	$::=$	\times $+$ \rightarrow
Type Pattern			
User Expression			
Internal Expression			

DECLARATIVES

$\Delta; \Phi \vdash \tau ::> \kappa$ τ has principal (well formed) kind κ

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{bse} ::> S_{\text{Type}}(\text{bse})} \text{PK-Base} \quad \frac{\Delta; \Phi_1, t::\kappa_1, \Phi_2 \vdash \text{OK}}{\Delta; \Phi \vdash t ::> S_\kappa(t)} \text{PK-Var} \quad \frac{\Delta; \Phi \vdash \tau_I :: \text{Type} \quad \Delta; \Phi \vdash \tau_B :: \text{Type}}{\Delta; \Phi \vdash \tau_I \oplus \tau_B ::> S_{\text{Type}}(\tau_I \oplus \tau_B)} \text{PK-}\oplus \quad \frac{\Delta_1, u::\kappa_1, \Delta_2; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \emptyset^u ::> S_\kappa(\emptyset^u)} \text{PK-EHole} \quad \frac{\Delta_1, u::\kappa_1, \Delta_2; \Phi \vdash \text{OK} \quad \Delta; \Phi \vdash \tau::\kappa_I}{\Delta; \Phi \vdash \langle \hat{\tau} \rangle^u ::> S_\kappa(\langle \hat{\tau} \rangle^u)} \text{PK-NEHole} \quad \frac{\Delta_1, u::\kappa_1, \Delta_2; \Phi \vdash \text{OK} \quad t \notin \Phi}{\Delta; \Phi \vdash \langle \emptyset \rangle^u ::> S_\kappa(\langle \emptyset \rangle^u)} \text{PK-Unbound} \quad \frac{\Delta; \Phi_s, t::\kappa_I \vdash \tau ::> \kappa_B}{\Delta; \Phi \vdash \lambda t::\kappa_I.\tau ::> S_{\Pi_{t::\kappa_I, \kappa_B}}(\lambda t::\kappa_I.\tau)} \text{PK-}\lambda$$

$$\frac{\Delta; \Phi \vdash \tau_I ::> \kappa \quad \Delta; \Phi \vdash \kappa \ \Pi \ \Pi_{t::\kappa_I, \kappa_B} \quad \Delta; \Phi \vdash \tau_B :: \kappa_I}{\Delta; \Phi \vdash \tau_I \ \tau_B ::> [\tau_B / t] \kappa_B} \text{PK-Ap}$$

$\Delta; \Phi \vdash \tau::\kappa$ τ is well formed at kind κ

$$\frac{\Delta; \Phi \vdash \tau ::> S_\kappa(\tau)}{\Delta; \Phi \vdash \tau::\kappa} \text{WFAK-1}$$

$$\frac{\Delta; \Phi \vdash \tau ::> \kappa_I \quad \Delta; \Phi \vdash \kappa_I \lesssim \kappa}{\Delta; \Phi \vdash \tau::\kappa} \text{WFAK-Subeump}$$

$$\frac{\Delta; \Phi \vdash \tau ::> \kappa}{\Delta; \Phi \vdash \tau::\kappa} \text{WFAK-Reit}$$

$$\frac{\Delta; \Phi \vdash \tau::\kappa}{\Delta; \Phi \vdash \tau::S_\kappa(\tau)} \text{WFAK-Self}$$

$$\frac{\Delta; \Phi \vdash \tau::\Pi_{t::\kappa_B, \kappa_I} \quad \Delta; \Phi \vdash \Pi_{t::\kappa_B, \kappa_I} \lesssim \Pi_{t::\kappa_B, \kappa_B}}{\Delta; \Phi \vdash \tau::\Pi_{t::\kappa_I, \kappa_B}} \text{WFAK-}\Pi\text{CSKTrans}$$

$$\frac{\Delta; \Phi \vdash \tau::S_\kappa(\tau_I) \quad \Delta; \Phi \vdash \tau_I::\kappa}{\Delta; \Phi \vdash \tau::\kappa} \text{WFAK-Flatten}$$

$\Delta; \Phi \vdash \kappa \ \Pi \ \Pi_{t::\kappa_I, \kappa_B}$ κ has matched Π -kind $\Pi_{t::\kappa_I, \kappa_B}$

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{KHole} \ \Pi \ \Pi_{t::\text{KHole}, \text{KHole}} \ \text{KHole}} \dashv\text{-KHole}$$

$$\frac{\Delta; \Phi \vdash \kappa \equiv S_{\text{KHole}}(\tau)}{\Delta; \Phi \vdash \kappa \ \Pi \ \Pi_{t::S_{\text{KHole}}(\tau), S_{\text{KHole}}(\tau \ t)} \dashv\text{-SKHole}}$$

$$\frac{\Delta; \Phi \vdash \kappa \equiv \Pi_{t::\kappa_I, \kappa_B}}{\Delta; \Phi \vdash \kappa \ \Pi \ \Pi_{t::\kappa_I, \kappa_B}} \dashv\text{-}\Pi$$

$\Delta; \Phi \vdash \kappa_I \equiv \kappa_B$ κ_I is equivalent to κ_B

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \equiv \kappa} \text{KEquiv-Ref1}$$

$$\frac{\Delta; \Phi \vdash \kappa_B \equiv \kappa_I}{\Delta; \Phi \vdash \kappa_I \equiv \kappa_B} \text{KEquiv-Symm}$$

$$\frac{\Delta; \Phi \vdash \kappa_I \equiv \kappa_B \quad \Delta; \Phi \vdash \kappa_B \equiv \kappa_C}{\Delta; \Phi \vdash \kappa_I \equiv \kappa_C} \text{KEquiv-Trans}$$

$$\frac{\Delta; \Phi \vdash \tau::S_\kappa(\tau_I)}{\Delta; \Phi \vdash S_{S_\kappa(\tau_I)}(\tau) \equiv S_\kappa(\tau_I)} \text{KEquiv-SKindSkind}$$

$$\frac{\Delta; \Phi \vdash \tau::\Pi_{t::\kappa_I, \kappa_B}}{\Delta; \Phi \vdash S_{\Pi_{t::\kappa_I, \kappa_B}}(\tau) \equiv \Pi_{t::\kappa_I, S_{[t_I / t] \kappa_B}}(\tau \ t_I)} \text{KEquiv-SKind}\Pi$$

$$\frac{\Delta; \Phi \vdash \kappa_I \equiv \kappa_B \quad \Delta; \Phi, t::\kappa_I \vdash \kappa_B \equiv \kappa_I}{\Delta; \Phi \vdash \Pi_{t::\kappa_I, \kappa_B} \equiv \Pi_{t::\kappa_B, \kappa_I}} \text{KEquiv-}\Pi$$

$$\frac{\Delta; \Phi \vdash \tau_I \stackrel{\kappa_I}{\equiv} \tau_B \quad \Delta; \Phi \vdash \kappa_I \equiv \kappa_B}{\Delta; \Phi \vdash S_{\kappa_I}(\tau_I) \equiv S_{\kappa_B}(\tau_B)} \text{KEquiv-SKind}$$

$\Delta; \Phi \vdash \kappa_I \lesssim \kappa_B$ κ_I is a consistent subkind of κ_B

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \text{KHole} \lesssim \kappa} \text{CSK-KHoleL}$$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim \text{KHole}} \text{CSK-KHoleR}$$

$$\frac{\Delta; \Phi \vdash S_{\text{KHole}}(\tau) \text{ OK} \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi \vdash S_{\text{KHole}}(\tau) \lesssim \kappa} \text{CSK-SKindSkoleL}$$

$$\frac{\Delta; \Phi \vdash \kappa \text{ OK} \quad \Delta; \Phi \vdash S_{\text{KHole}}(\tau) \text{ OK}}{\Delta; \Phi \vdash \kappa \lesssim S_{\text{KHole}}(\tau)} \text{CSK-SKindSkoleR}$$

$$\frac{\Delta; \Phi \vdash \kappa_I \equiv \kappa_B}{\Delta; \Phi \vdash \kappa_I \lesssim \kappa_B} \text{CSK-KEquiv}$$

$$\frac{\Delta; \Phi \vdash \kappa_I \equiv \kappa_B \quad \Delta; \Phi \vdash \kappa_B \lesssim \kappa_I \quad \Delta; \Phi \vdash \kappa_I \equiv \kappa_B}{\Delta; \Phi \vdash \kappa_I \lesssim \kappa_B} \text{CSK-Normal}$$

$$\frac{\Delta; \Phi \vdash S_\kappa(\tau) \text{ OK}}{\Delta; \Phi \vdash S_\kappa(\tau) \lesssim \kappa} \text{CSK-SKind}$$

$$\frac{\Delta; \Phi \vdash \kappa_B \lesssim \kappa_I \quad \Delta; \Phi, t::\kappa_B \vdash \kappa_B \lesssim \kappa_I}{\Delta; \Phi \vdash \Pi_{t::\kappa_I, \kappa_B} \lesssim \Pi_{t::\kappa_B, \kappa_I}} \text{CSK-}\Pi$$

$$\frac{\Delta; \Phi \vdash \kappa_I \lesssim \kappa_B \quad \Delta; \Phi \vdash \tau_I \stackrel{\kappa_I}{\equiv} \tau_B}{\Delta; \Phi \vdash S_{\kappa_I}(\tau_I) \lesssim S_{\kappa_B}(\tau_B)} \text{CSK-?}$$

$\Delta; \Phi \vdash \tau_I \stackrel{\kappa_I}{\equiv} \tau_B$ τ_I is provably equivalent to τ_B at kind κ

$$\frac{\Delta; \Phi \vdash \tau::\kappa}{\Delta; \Phi \vdash \tau \equiv \tau} \text{EquivAK-Ref1}$$

$$\frac{\Delta; \Phi \vdash \tau_B \stackrel{\kappa}{\equiv} \tau_I}{\Delta; \Phi \vdash \tau_I \stackrel{\kappa}{\equiv} \tau_B} \text{EquivAK-Symm}$$

$$\frac{\Delta; \Phi \vdash \tau_I \stackrel{\kappa}{\equiv} \tau_B \quad \Delta; \Phi \vdash \tau_B \stackrel{\kappa}{\equiv} \tau_I}{\Delta; \Phi \vdash \tau_I \stackrel{\kappa}{\equiv} \tau_I} \text{EquivAK-Trans}$$

$$\frac{\Delta; \Phi \vdash \tau_I ::> \kappa_I \quad \Delta; \Phi \vdash \kappa_I \equiv S_\kappa(\tau_B)}{\Delta; \Phi \vdash \tau_I \stackrel{\kappa}{\equiv} \tau_B} \text{EquivAK-SKind}$$

$$\frac{\Delta; \Phi \vdash \tau_I::\Pi_{t::\kappa_I, \kappa_B} \quad \Delta; \Phi \vdash \tau_B::\Pi_{t::\kappa_I, \kappa_I} \quad \Delta; \Phi, t::\kappa_I \vdash \tau_I \ t \stackrel{\kappa_B}{\equiv} \tau_B \ t}{\Delta; \Phi \vdash \tau_I \stackrel{\kappa}{\equiv} \tau_B} \text{EquivAK-}\Pi$$

$$\frac{\Delta; \Phi \vdash \tau_I \stackrel{\Pi_{t::\kappa_I, \kappa_B}}{\equiv} \tau_B \quad \Delta; \Phi \vdash \tau_B \stackrel{\kappa_I}{\equiv} \tau_I}{\Delta; \Phi \vdash \tau_I \ \tau_B \stackrel{[\tau_B / t] \kappa_B}{\equiv} \tau_B \ \tau_I} \text{EquivAK-Ap}$$

$$\frac{\Delta; \Phi \vdash \tau_I \stackrel{\text{Type}}{\equiv} \tau_B \quad \Delta; \Phi \vdash \tau_B \stackrel{\text{Type}}{\equiv} \tau_I}{\Delta; \Phi \vdash \tau_I \oplus \tau_B \stackrel{\text{Type}}{\equiv} \tau_B \oplus \tau_I} (2)$$

$$\frac{\Delta; \Phi \vdash \tau_I \stackrel{S_\kappa(\tau)}{\equiv} \tau_B \quad (1) \quad \Delta; \Phi \vdash \kappa_I \equiv \kappa_B \quad \Delta; \Phi, t::\kappa_I \vdash \tau_I \ t \stackrel{\kappa}{\equiv} \tau_B \quad (3)}{\Delta; \Phi \vdash \lambda t::\kappa_I.\tau_I \equiv \lambda t::\kappa_B.\tau_B} (3)$$

$$\frac{\Delta; \Phi \vdash \tau_I \stackrel{\kappa_I}{\equiv} \tau_B \quad \Delta; \Phi \vdash \kappa_I \equiv \kappa}{\Delta; \Phi \vdash \tau_I \equiv \tau_B} (4)$$

$\Delta; \Phi \vdash \kappa \text{ OK}$ κ is well formed

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{Type OK}} \text{KWF-Type}$$

$$\frac{\Delta; \Phi \vdash \text{OK}}{\Delta; \Phi \vdash \text{KHole OK}} \text{KWF-KHole}$$

$$\frac{\Delta; \Phi \vdash \tau::\kappa}{\Delta; \Phi \vdash S_\kappa(\tau) \text{ OK}} \text{KWF-SKind}$$

$$\frac{\Delta; \Phi, t::\kappa_I \vdash \kappa_B \text{ OK}}{\Delta; \Phi \vdash \Pi_{t::\kappa_I, \kappa_B} \text{ OK}} \text{KWF-}\Pi$$

$\Delta; \Phi \vdash \text{OK}$ Context is well formed

$$\frac{}{\cdot, \cdot \vdash \text{OK}} \text{CWF-Nil}$$

$$\frac{t \notin \Delta \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta; \Phi, t::\kappa \vdash \text{OK}} \text{CWF-TypVar}$$

$$\frac{u \notin \Delta \quad \Delta; \Phi \vdash \kappa \text{ OK}}{\Delta, u::\kappa; \Phi \vdash \text{OK}} \text{CWF-Hole}$$

METATHEORY

Lemma 1 (COK). If $\Delta; \Phi \vdash \mathcal{J}$, then $\Delta; \Phi \vdash \text{OK}$ in a subderivation (where $\Delta; \Phi \vdash \mathcal{J} \neq \Delta; \Phi \vdash \text{OK}$)

Proof. By simultaneous induction on derivations.
No interesting cases. □

Lemma 2 (Exchange).

If $\Delta; \Phi_1, t_{L1}::\kappa_{L1}, t_{L2}::\kappa_{L2}, \Phi_2 \vdash \mathcal{J}$ and $\Delta; \Phi_1, t_{L2}::\kappa_{L2}, t_{L1}::\kappa_{L1}, \Phi_2 \vdash \text{OK}$, then $\Delta; \Phi_1, t_{L2}::\kappa_{L2}, t_{L1}::\kappa_{L1}, \Phi_2 \vdash \mathcal{J}$

Proof. By induction on derivations.
No interesting cases.

(Only rules with Φ extended in the consequent are interesting, which is only CWF-TypVar, but when \mathcal{J} is CWF, Exchange is identity) □

Corollary 3 (Marked-Exchange).

If $\Delta; \Phi, t_{L1}::\kappa_{L1}, t_{L2}::\kappa_{L2} \vdash \mathcal{J}$ and $\Delta; \Phi, t_{L2}::\kappa_{L2}, t_{L1}::\kappa_{L1} \vdash \text{OK}$, then $\Delta; \Phi, t_{L2}::\kappa_{L2}, t_{L1}::\kappa_{L1} \vdash \mathcal{J}$

Lemma 4 (Weakening).

If $\Delta; \Phi \vdash \mathcal{J}$ and $\Delta; \Phi, t_L::\kappa_L \vdash \text{OK}$ and $t_L \notin \mathcal{J}$ and $\forall t \in \kappa_L, t \notin \mathcal{J}$, then $\Delta; \Phi, t_L::\kappa_L \vdash \mathcal{J}$

Proof. By induction on derivations.

When applying Weakening in the induction, check that the left premiss is always a subderivation, and check variable exclusion conditions are satisfied.

(PoS = premiss of subderivation)

Weakening

[illegible]

□

Lemma 5 (OK-PK). *If $\Delta; \Phi \vdash \tau ::> \kappa$, then $\Delta; \Phi \vdash \kappa$ OK*

Lemma 6 (OK-WFaK). *If $\Delta; \Phi \vdash \tau :: \kappa$, then $\Delta; \Phi \vdash \kappa$ OK*

Lemma 7 (OK-MatchPi). *If $\Delta; \Phi \vdash \kappa \overset{\blacktriangleright}{\Pi} \Pi_{t::\kappa_1, \kappa_2}$, then $\Delta; \Phi \vdash \kappa$ OK and $\Delta; \Phi \vdash \Pi_{t::\kappa_1, \kappa_2}$ OK*

Lemma 8 (OK-KEquiv). *If $\Delta; \Phi \vdash \kappa_I \equiv \kappa_2$, then $\Delta; \Phi \vdash \kappa_I$ OK and $\Delta; \Phi \vdash \kappa_2$ OK*

Lemma 9 (OK-CSK). *If $\Delta; \Phi \vdash \kappa_I \lesssim \kappa_2$, then $\Delta; \Phi \vdash \kappa_I$ OK and $\Delta; \Phi \vdash \kappa_2$ OK*

Lemma 10 (OK-EquivAK). *If $\Delta; \Phi \vdash \tau_I \overset{\triangle}{=} \tau_2$, then $\Delta; \Phi \vdash \tau_I :: \kappa$ and $\Delta; \Phi \vdash \tau_2 :: \kappa$ and $\Delta; \Phi \vdash \kappa$ OK*

Lemma 11 (OK-Substitution).

If $\Delta; \Phi \vdash \tau_L :: \kappa_{L1}$ and $\Delta; \Phi, t_L :: \kappa_{L1} \vdash \kappa_{L2}$ OK, then $\Delta; \Phi \vdash [\tau_L/t_L] \kappa_{L2}$ OK

(induction on $\Delta; \Phi, t_L :: \kappa_{L1} \vdash \kappa_{L2}$ OK)

Lemma 12 (K-Substitution).

If $\Delta; \Phi \vdash \tau_{L1} :: \kappa_{L1}$ and $\Delta; \Phi, t_L :: \kappa_{L1} \vdash \tau_{L2} :: \kappa_{L2}$, then $\Delta; \Phi \vdash [\tau_{L1}/t_L] \tau_{L2} :: [\tau_{L1}/t_L] \kappa_{L2}$

(induction on $\Delta; \Phi, t_L :: \kappa_{L1} \vdash \tau_{L2} :: \kappa_{L2}$)

Proof. By simultaneous induction on derivations.

The interesting cases per lemma:

(PoS = premiss of subderivation)	OK-PK.	PK-Base	$\Delta; \Phi \vdash \mathbf{bse} :: \mathbf{S_{Type}}(\mathbf{bse})$ $\Delta; \Phi \vdash \mathbf{bse} :: \mathbf{Type}$ $\Delta; \Phi \vdash \mathbf{S_{Type}}(\mathbf{bse})$ OK $\Delta; \Phi \vdash \text{OK}$	by (9) by (10) by (43) by premiss bad by (10) by (43)
		*		
		*		
		PK-Ap		
	OK-WFaK.	(12)	$\Delta; \Phi \vdash \tau_2 :: \kappa$ $\Delta; \Phi \vdash \mathbf{S}_{\kappa}(\tau_2)$ OK	
		*	$\Delta; \Phi \vdash \tau \ t ::> \kappa$ $\Delta; \Phi, t_L :: \kappa_{L1} \vdash \text{OK}$ $\Delta; \Phi \vdash \kappa_{L1}$ OK $\Delta; \Phi \vdash \text{OK}$	
	OK-KEquiv.	(22)	$\Delta; \Phi, t_L :: \kappa_{L1} \vdash \text{OK}$ $\Delta; \Phi \vdash \kappa_{L1}$ OK	premiss (41) by subderivation premiss (46)
	OK-Substitution.	(41)	$\Delta; \Phi, t_L :: \kappa_{L1} \vdash \tau :: \kappa$ $\Delta; \Phi, t_L :: \kappa_{L1} \vdash \text{OK}$ $\Delta; \Phi \vdash \kappa_{L1}$ OK $\Delta; \Phi \vdash \text{OK}$	by OK-KWF by (41) and degenerate subst premiss (43) by OK-WFaK
		*	$\Delta; \Phi \vdash \tau_{L1}/t_L \text{Type}$ OK	by subderivation premiss (46)
		*	$\Delta; \Phi, t_L :: \kappa_{L1} \vdash \tau :: \kappa$ $\Delta; \Phi, t_L :: \kappa_{L1} \vdash \text{OK}$ $\Delta; \Phi \vdash \kappa_{L1}$ OK $\Delta; \Phi \vdash \text{OK}$	by OK-KWF
		(43)	$\Delta; \Phi \vdash [\tau_{L1}/t_L] \tau :: [\tau_{L1}/t_L] \kappa$ $\Delta; \Phi \vdash [\tau_{L1}/t_L] \mathbf{S}_{\kappa}(\tau)$ OK	by K-Substitution on premiss by (43)
		*		
		*		

□

Lemma 13 (PK-Unicity). *If $\Delta; \Phi \vdash \tau_L ::> \kappa_{L1}$ and $\Delta; \Phi \vdash \tau_L ::> \kappa_{L2}$ then κ_{L1} is κ_{L2}*

Lemma 14. *If $\Delta; \Phi \vdash \tau ::> \kappa_I$ and $\Delta; \Phi \vdash \tau :: \kappa_2$, then $\Delta; \Phi \vdash \kappa_I \lesssim \kappa_2$*

Lemma 15. *If $\Delta; \Phi \vdash \kappa_I \lesssim \mathbf{S}_{\kappa_2}(\tau)$, then $\Delta; \Phi \vdash \kappa_I \lesssim \kappa_2$*