

Longest Common Subsequence (LCS)

String 1: a b c f g h i j

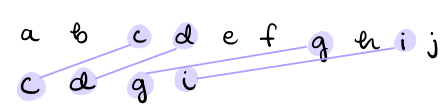
String 2: c d g i

Some characters common between string 1 and string 2.

which set of chars already in string 2 IN SEQUENCE are also in string 1?

↳ select longest common subsequence

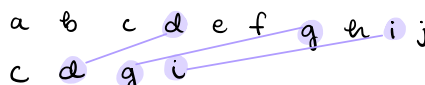
String 1: a b c d e f g h i j
String 2: c d g i



Subsequence 1

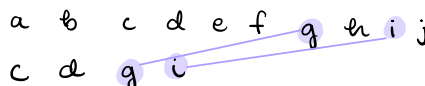
longest
cdgi
length 4

String 1: a b c d e f g h i j
String 2: c d g i



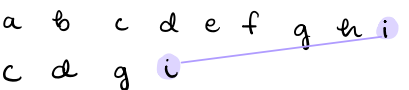
Subsequence 2 dgi

String 1: a b c d e f g h i j
String 2: c d g i



Subsequence 3 gi

String 1: a b c d e f g h i j
String 2: c d g i



Subsequence 4 i

String 1: a b c d e f g h i j
String 2: c d g i

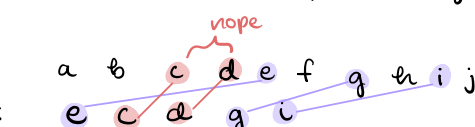


Subsequence 4 cdg

What if we added e?

limitation: NO INTERSECTION, cannot go across

String 1: a b c d e f g h i j
String 2: e c d g i



so egi LCS

PRACTICE:

String 1: a b d a c e *nope*
 String 2: b a b c e

String 1: a b d a c e ✓
 String 2: b a b c e ✓

OR

String 1: a b d a c e ✓
 String 2: b a b c e ✓

LCS = 4 but there are different subsequences

Tabulation Method

String 1: l o n g e s t
 String 2: s t o n e

	0	1	2	3	4	5	6
A	l	o	n	g	e	s	t
B	s	t	o	n	e		
	0	1	2	3	4		

		l	o	n	g	e	s	t
	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0
s 1	0	0	0	0	0	0	1	1
t 2	0	0	0	0	0	0	1	2
o 3	0	0	1	1	1	1	1	2
n 4	0	0	1	2	2	2	2	2
e 5	0	0	1	2	2	3	3	3

match = max of diagonal + 1

otherwise = max of diagonal
 ↳ previous column and row

		l	o	n	g	e	s	t
	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0
s 1	0	0	0	0	0	0	1	1
t 2	0	0	0	0	0	0	1	2
o 3	0	0	1	1	1	1	1	2
n 4	0	0	1	2	2	2	2	2
e 5	0	0	1	2	2	3	3	3

← if same # to left

↖ otherwise

● = included in LCS

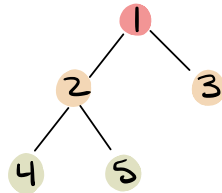
LCS = one

Breadth First Search

BFS & DFS: graph traversal search

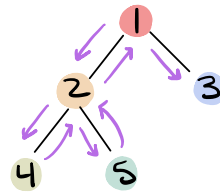
- discover the vertex (visit)
- explore the vertex
- select any vertex as source

BFS:



Queue
1, 2, 3, 4, 5

DFS:



Stack
1, 2, 4, 5, 3

$\xrightarrow{\text{vertex}}$
 $\xrightarrow{\text{edge}}$
 $\xrightarrow{\text{source}}$
 BFS(V, E, S)
 for each vertex v in $V - \{S\}$ \leftarrow exclude source
 do
 $\text{color}[v] \leftarrow \text{white}$ \leftarrow initially all vertices are white
 $d[v] \leftarrow \infty$ $\leftarrow d[v]$: distance of vertex v from source
 $\pi[v] \leftarrow \text{NULL}$ $\leftarrow \pi[v]$: parent of vertex v
 $\text{color}[S] = \text{Grey}$
 $d[S] \leftarrow 0$ \leftarrow distance of source from source = 0
 $\pi[S] \leftarrow \text{NULL}$
 $Q \leftarrow \{ \}$
 ENQUEUE(Q, S) \leftarrow Queue source now
 while Q is not empty
 do $v \leftarrow \text{DEQUEUE}(Q)$
 for each u adjacent to v
 do if $\text{color}[u] \leftarrow \text{White}$ \leftarrow all white adjacent nodes to S
 then $\text{color}[u] \leftarrow \text{Grey}$ \leftarrow change color to Grey
 $d[u] \leftarrow d[v] + 1$ \leftarrow increase $d[v]$ by one
 $\pi[u] \leftarrow v$
 ENQUEUE(Q, u) \leftarrow add the now Grey node to queue
 $\text{color}[v] = \text{Black}$ \leftarrow after dequeued = black

$\text{color}[v]$

- white (vertex has not been discovered)
- Grey (vertex has been discovered & is being processed)
- Black (has been completely processed)

1. all white adjacent nodes to S change color to grey
2. increase $d[v]$ by one
3. add now grey nodes to the queue
4. change color of dequeued nodes to black

1, 5, 4, 2, 7, 6, 3

