### Query Optimization: From 0 to 10 (and up to 5.7)

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dbahire.com/pleu15

### **Agenda - First 3 hours**

1. Introduction	5. FULLTEXT Search
2. Access Types and Basic Indexing Techniques	6. Joins
3. Break	7. Subqueries
4. Multi-Column Indexing	8. Query Profiling

#### **Agenda - Last 3 hours**

1. General Optimizer Improvements	6. SQL Mode Changes
2. Computed/Virtual Columns	7. GIS Improvements and JSON Type
3. Query Rewrite Plugins	8. Results and Conclusions
4. Break	9. Q&A
5. Optimizer Hints	

### Query Optimization: From 0 to 10 (and up to 5.7) INTRODUCTION

### This is me fighting bad query performance



- Sr. Database Administrator at Wikimedia Foundation
- Used to work as a trainer for Oracle (MySQL), as a Consultant (Percona) and as a Freelance administrator (DBAHire.com)

#### **MySQL Versions**

- 5.1 no longer has official support
- I will be showing you the results on mysql versions 5.5-5.7/10.1
- MySQL 5.7 and MariaDB 10.1 in RC with great new features

### Recently Added Features Related to Query Optimization

- Index Condition Pushdown
- Subquery Optimizations materialization and semijoin)
- IN-to-EXISTS/EXISTS-to-IN
- JOIN-to-WHERE
- Multi-Range Read
- Batched Key Access
- Persistent InnoDB Statistics
- UNION ALL optimization
- Improved GIS support

- EXPLAIN FORMAT=JSON
- EXPLAIN INSERT/UPDATE/DELETE
- Hash Join
- New optimizer hints
- New cost-based optimizer
- Optimizer Trace
- Filesort optimizations
- Virtual/computed columns and "functional indexes"
- New JSON type

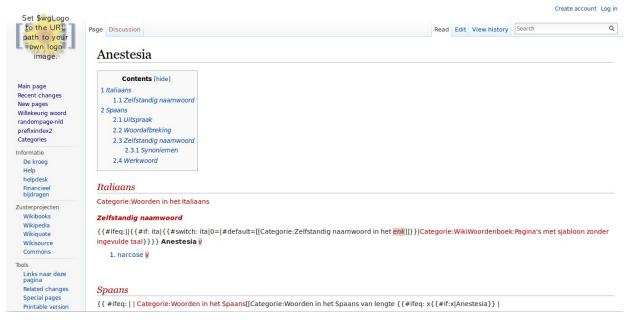
#### I Already Mentioned Some of Those Topics during the Last Years





 Check my presentations here: <a href="http://www.slideshare.net/jynus/">http://www.slideshare.net/jynus/</a>

### Example Application (I)



 Wiktionary (and all Wikimedia project's data) is licensed under the Creative Commons BY-SA-2.5 License and is Copyright its Contributors

#### Example Application (II)



OSM Database is licensed under the Open
 DataBase License and is Copyright OpenStreetMap
 Contributors

#### Install the example databases

- Downloads and instructions at: http://dbahire.com/pleu15
  - Requirements: a MySQL or MariaDB installation (MySQL Sandbox is suggested)
  - The wiktionary and OSM extracts
- Import them by doing:
   \$ bzcat <file> | mysql <database>

Query Optimization: From 0 to 10 (and up to 5.7)

## ACCESS TYPES AND BASIC INDEXING TECHNIQUES

#### **EXPLAIN**

- Essential to understand the execution plan of our queries
  - Works on SELECTs, INSERTs, UPDATEs,
     REPLACEs, DELETEs and connections
  - Fully documented on: http://dev.mysql.com/doc/refman/5.6/en/expla in-output.html

#### **EXPLAIN Example**

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE page_title
= 'Dutch';
+-----+
key_len | ref | rows | Extra
+----+----+-----
+-----+
| 1 | SIMPLE | page | ALL | NULL
                       I NULL I NULL
Difficult to see
+----+
                      something
1 row in set (0.00 sec)
```

# EXPLAIN Example (vertical format)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page title = 'Dutch'\G
********
                               ********
         id: 1
 select type: SIMPLE
       table: page
        type: ALL
possible keys: NULL
        kev: NULL
     key len: NULL
         ref: NULL
                                          Use \G for
        rows: 90956
                                           vertical
       Extra: Using where
1 row in set (0.00 sec)
                                          formatting
```

#### **EXPLAIN Example (id)**

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page title = 'Dutch'\G
********
                               ********
                      Indicates hierarchy
 select type: SIMPLE
                      level, not execution
       table: page
                             order
        type: ALL
possible keys: NULL
        key: NULL
     key len: NULL
         ref: NULL
        rows: 90956
       Extra: Using where
1 row in set (0.00 sec)
```

#### EXPLAIN Example (select\_type)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page title = 'Dutch'\G
********
                               ********
                         1. row
         id: 1
 select type: SIMPLE
                        Not a subquery or a
       table: page
                               UNION
        type: ALL
possible keys: NULL
        key: NULL
     key len: NULL
         ref: NULL
        rows: 90956
       Extra: Using where
1 row in set (0.00 sec)
```

#### **EXPLAIN Example (table)**

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page title = 'Dutch'\G
********
                               *********
         id: 1
 select type: SIMPLE
       table: page
                       Table scanned for
        type: ALL
                            this step
possible keys: NULL
        key: NULL
     key len: NULL
         ref: NULL
        rows: 90956
       Extra: Using where
1 row in set (0.00 sec)
```

#### **EXPLAIN Example (type)**

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page title = 'Dutch'\G
********
                               *********
          id: 1
 select type: SIMPLE
       table: page
                     'All rows are read for
        type: ALL -
                        this table (FULL
possible keys: NULL
                         TABLE SCAN)
         key: NULL
     key len: NULL
         ref: NULL
        rows: 90956
       Extra: Using where
1 row in set (0.00 sec)
```

#### EXPLAIN Example (rows)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page title = 'Dutch'\G
********
                               *********
         id: 1
 select type: SIMPLE
       table: page
        type: ALL
possible keys: NULL
        kev: NULL
     key len: NULL
         ref: NULL
                         Estimated number
        rows: 90956
                         of rows to be read
       Extra: Using where
                          (all table rows)
1 row in set (0.00 sec)
```

#### How to improve performance?

```
MariaDB [nlwiktionary]> SELECT * FROM page
WHERE page_title = 'Dutch';
1 row in set (0.11 sec)
```

Let's add an index on page.page\_title:

```
MariaDB [nlwiktionary]> ALTER TABLE page ADD INDEX page_title (page_title);
Query OK, 0 rows affected (0.19 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

### Index creation results (type)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page title = 'Dutch'\G
********
                               ********
         id: 1
 select type: SIMPLE
       table: page
                       type: ref means that an
        type: ref
                      equality comparison will
possible keys: page ...
                       be checked against an
        key: page titl
                      index and several results
     key len: 257
        ref: const
                          could be returned
        rows: 1
       Extra: Using index condition
1 row in set (0.00 sec)
```

# Index creation results (possible\_keys and key)

```
MariaDB [nlwiktionary] > EXPLAIN SELECT * FROM page WHERE
page title = 'Dutch'\G
********
                               *********
          id: 1
 select type: SIMPLE
       table: page
        type: ref
                            index(es) that the
possible keys: page title
                           optimizer considered
        key: page title
     key len: 257
                          potentially useful, and
         ref: const
                            final index chosen
        rows: 1
       Extra: Using index condition
1 row in set (0.00 sec)
```

#### Index creation results (ref)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page title = 'Dutch'\G
********
                               *********
         id: 1
 select type: SIMPLE
       table: page
        type: ref
possible keys: page title
        key: page title
     key len: 257
        ref: const
                   Index is compared
        rows: 1
       Extra: Using i with a constant, not
1 row in set (0.00 sec with another table
```

#### Index creation results (rows)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page title = 'Dutch'\G
********
                               *********
          id: 1
 select type: SIMPLE
       table: page
        type: ref
possible keys: page title
        key: page title
     key len: 257
         ref: const
        rows: 1
                    Only 1 row read. In this
       Extra: Usus
                    case, estimation is exact
1 row in set (0.00 se
                     (thanks to index dive)
```

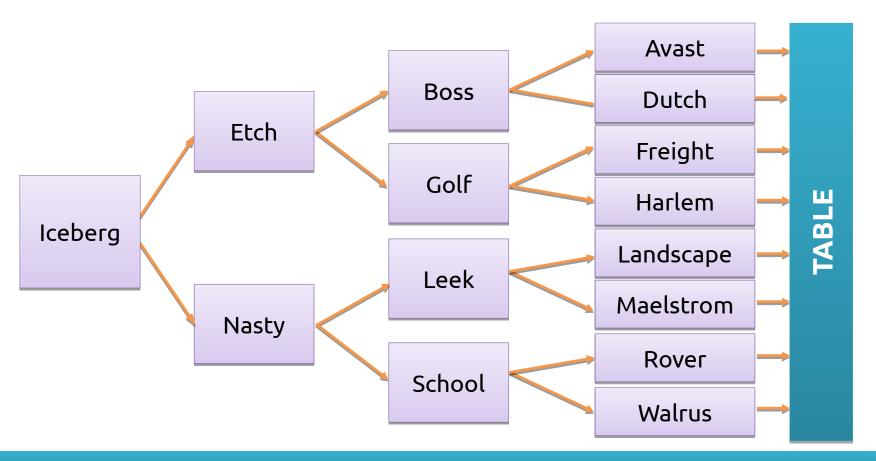
# Index creation results (query time)

```
MariaDB [nlwiktionary] > SELECT * FROM page
WHERE page_title = 'Dutch';
1 row in set (0.00 sec) Query time has been reduced substantially
```

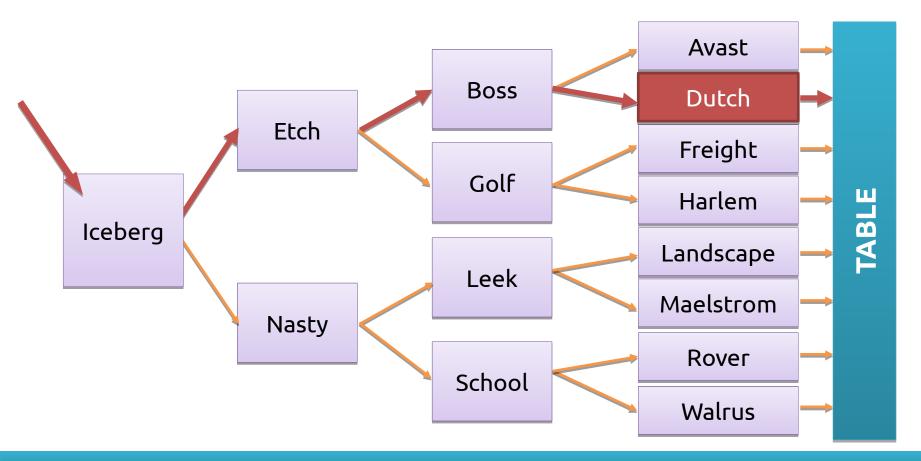
#### Types of indexes

- BTREE
  - B-TREE in MyISAM, B+TREE in InnoDB
- HASH
  - Only available for MEMORY and NDB
- FULLTEXT
  - Inverted indexes in MyISAM and InnoDB
- SPATIAL
  - RTREEs in MyISAM and InnoDB

#### Finding "Dutch" with a BTREE



#### Finding "Dutch" with a BTREE



#### Do indexes always work? (1/2)

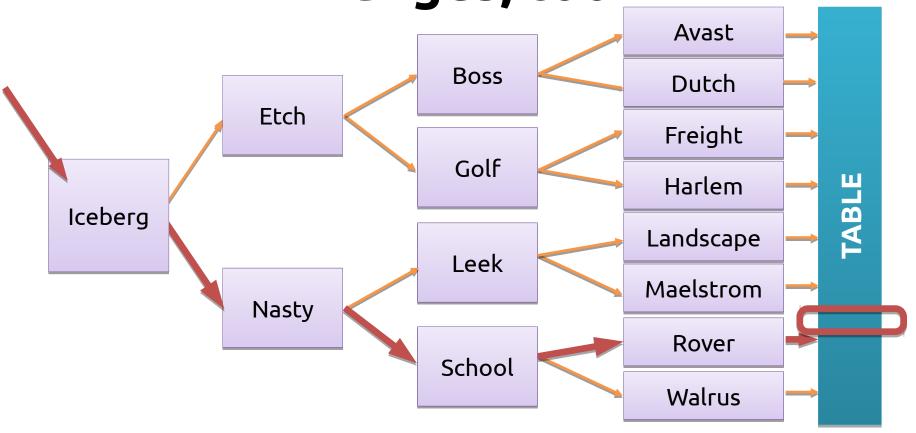
 Can we use an index to make this query faster?

```
SELECT * FROM page WHERE page_title like 'Spa%';
```

#### It is a range

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page title like 'Spa%'\G
********
                                ********
          id: 1
 select type: SIMPLE
       table: page
        type: range
                              Despite not being an
possible keys: page_title
                             equality, we can use the
         key: page title
                             index to find the values
     key len: 257
         ref: NULL
                                    quickly
        rows: 94
       Extra: Using index condition
1 row in set (0.00 sec)
```

BTREE Indexes can be used for ranges, too



### Do indexes always work? (2/2)

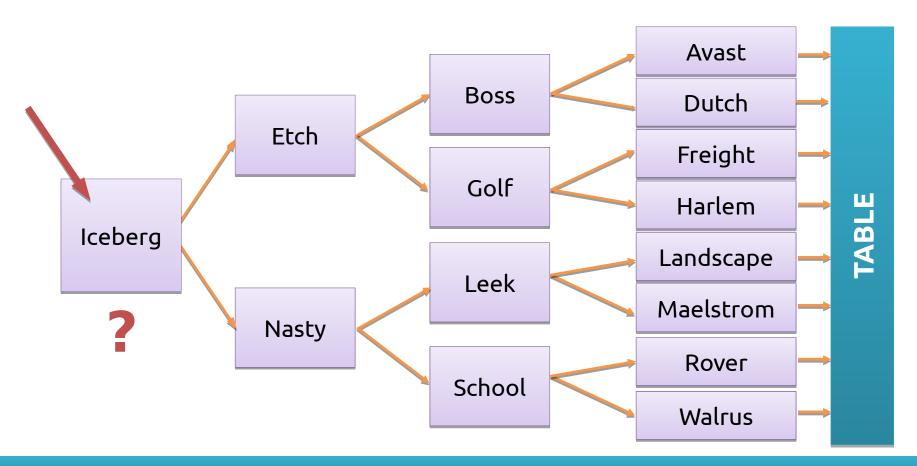
What about this other query?

```
SELECT * FROM page WHERE page_title like '%utch';
```

#### Let's check with EXPLAIN

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page title like '%utch'\G
********
                                ********
          id: 1
 select type: SIMPLE
                          No index can be used
       table: page
        type: ALL
                            for filtering. A full
possible keys: NULL
                              table scan is
         key: NULL
                               performed.
     key len: NULL
         ref: NULL
        rows: 93189
       Extra: Using where
1 row in set (0.00 sec)
```

#### **BTREE Index**



#### Btree indexes usage

- Filtering
  - Equality (operator '=')
  - Ranges (BETWEEN ... AND, >, <, >=, <=, like 'prefix%')
  - "EXISTS" operators: IN, OR on the same column
- Ordering
  - ORDER BY (indexed columns)
  - GROUP BY (indexed columns)
- Returning values directly from the index
  - Covering index
  - Functions like max(), min(), etc.

#### type: const

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision WHERE
rev id = 2\G
********
                                  *********
          id: 1
 select type: SIMPLE
                              'const' is a special case of 'ref',
       table: revision
        type: const <==
                               when the index can assure
possible keys: PRIMARY
                               that only 1 results can be
         key: PRIMARY
                               returned (equality + primary
     key len: 4
                               key or unique key). It is faster.
         ref: const
        rows: 1
       Extra:
1 row in set (0.00 sec)
```

#### type: NULL

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision WHERE
rev id = -1\G
********
                                 *********
                          1. row
          id: 1
 select_type: SIMPLE
       table: NULL
                           'NULL' is not really a plan,
        type: NULL
                           just an optimization that
possible keys: NULT
                           allow discarding
         key: NULL
     key_len: NULL
                           immediately impossible
         ref: NULL
                            conditions
        rows: NULL
       Extra: Impossible WHERE noticed after reading const
tables
1 row in set (0.00 sec)
```

#### type: ref\_or\_null

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM user WHERE
user email token = '0' OR user email token IS NULL\G
id: 1
 select type: SIMPLE
                              Equivalent to 'ref', but
       table: user
                              also has into account
       type: ref or null
                                 NULL values
possible keys: user email token
        key: user email token
     key len: 33
        ref: const
       rows: 2
       Extra: Using index condition; Using where
1 row in set (0.00 sec)
```

#### type: range (using IN / OR)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page title IN ('Dutch', 'English', 'Spanish')\G
id: 1
 select type: SIMPLE
                             Despite being a range, its
       table: page
                             execution is very
        type: range
possible_keys: page_title
                             different from ranges
        key: page title
                             using like, between or
     key len: 257
                            inequality operators
        ref: NULL
        rows: 4
       Extra: Using index condition
1 row in set (0.00 sec)
```

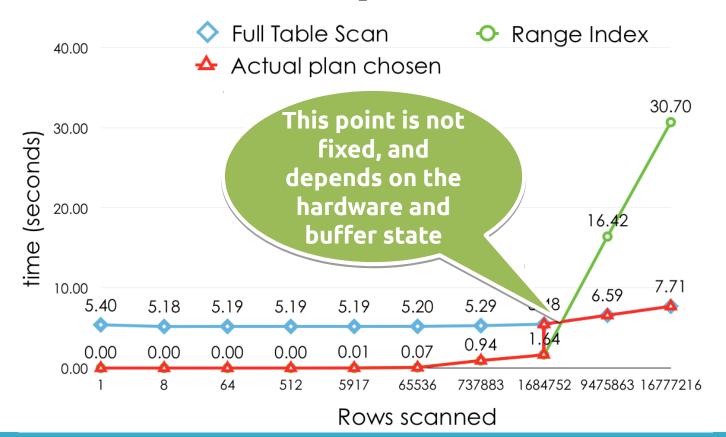
### Is this a bug? (1/2)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page namespace = 2\G
********
                                *********
          id: 1
 select type: SIMPLE
       table: page
        type: ref
possible keys: name title
                                 An index is used to
         key: name title
                               return pages with ns=2
     key len: 4
         ref: const
        rows: 45
       Extra:
1 row in set (0.00 sec)
```

#### Is this a bug? (2/2)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page WHERE
page namespace = 0\G
********
                                *********
          id: 1
 select type: SIMPLE
       table: page
        type: ALL
possible keys: name_title
                          The index is not used
         key: NULL <
                               with ns=0
     key len: NULL
         ref: NULL
        rows: 7493
       Extra: Using where
1 row in set (0.00 sec)
```

### Using an index is sometimes suboptimal



# What index should we add to make this query faster?

```
MariaDB [nlwiktionary] > EXPLAIN SELECT * FROM revision WHERE
left(rev timestamp, 6) = '201509'\G
id: 1
 select type: SIMPLE
      table: revision
       type: ALL
possible keys: NULL
        key: NULL
    key len: NULL
        ref: NULL
       rows: 163253
      Extra: Using where
1 row in set (0.00 sec)
```

### The table has already an index on rev\_timestamp

```
MariaDB [nlwiktionary]> SHOW CREATE TABLE revision\G

************************

Table: revision

Create Table: CREATE TABLE `revision` (
...

KEY `rev_timestamp` (`rev_timestamp`),
```

#### We need to rewrite the query

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision WHERE
rev timestamp >= '201509' and rev timestamp < '201510'\G
id: 1
 select type: SIMPLE
      table: revision
       type: range
possible keys: rev timestamp
        key: rev timestamp
     key len: 14
        ref: NULL
       rows: 7
      Extra: Using index condition
1 row in set (0.00 sec)
```

### This transformation is not trivial or even possible in all cases

```
MariaDB [nlwiktionary] > EXPLAIN SELECT * FROM revision WHERE
substr(rev timestamp, 5, 2) = '09'\G
*******
                          1. row ***************
          id: 1
 select type: SIMPLE
       table: revision
        type: ALL
possible keys: NULL
                          Can you think a way to
         key: NULL
                           improve this query?
     key len: NULL
         ref: NULL
        Rows: 173154
       Extra: Using where
1 row in set (0.00 sec)
```

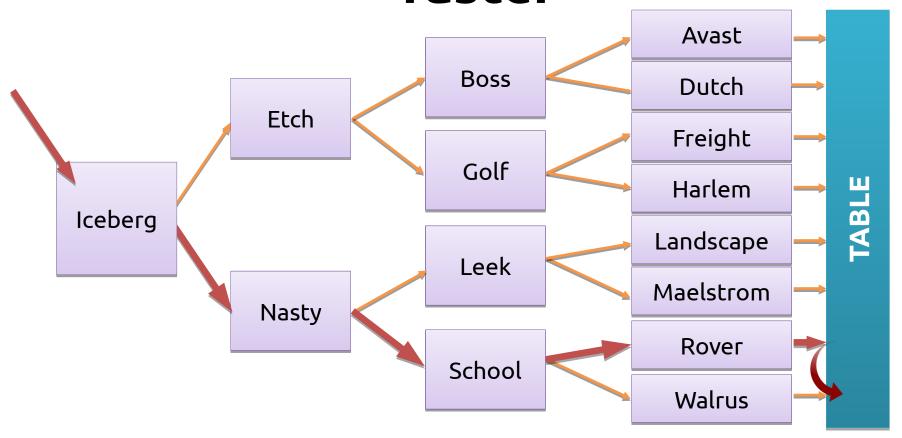
#### Indexes for Ordering

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page ORDER BY
page touched DESC LIMIT 10\G
*******
                                ********
          id: 1
 select type: SIMPLE
       table: page
        type: ALL
possible keys: NULL
         key: NULL
     key len: NULL
         ref: NULL
                            "Using filesort" indicates that an
        rows: 8720
                               ordering is needed before
       Extra: Using filesort
                                 returning the results
1 row in set (0.00 sec)
```

### If that is frequent, we can create an index on page\_touched...

```
[nlwiktionary] > ALTER TABLE page ADD INDEX page page touched(page touched);
Query OK, 0 rows affected (0.30 sec)
Records: 0 Duplicates: 0 Warnings: 0
[nlwiktionary]> EXPLAIN SELECT * FROM page ORDER BY page_touched DESC LIMIT
10\G
                          1. FOW ****************
********
          id: 1
 select type: SIMPLE
                                The index does not produce any
       table: page
        type: index
                                      advantage for filtering
possible_keys: NULL
         key: page_page_touched
     key len: 14
                                 However, it is very effective by
         ref: NULL
        rows: 10
                                 helping avoiding the sort phase
       Fxtra:
1 row in set (0.00 sec)
```

### It can return data in index order faster



### Indexes and GROUP BY (no indexes)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT rev_page, count(*) FROM
revision IGNORE INDEX(rev_page_id, page_timestamp,
page user timestamp) GROUP BY rev page\G
id: 1
 select type: SIMPLE
       table: revision
        type: ALL
possible keys: NULL
        key: NULL
                           Without indexes, a temporary
     key len: NULL
        ref: NULL
                          table is created to order results
        rows: 201094
       Extra: Using temporary; Using filesort
1 row in set (0.00 sec)
```

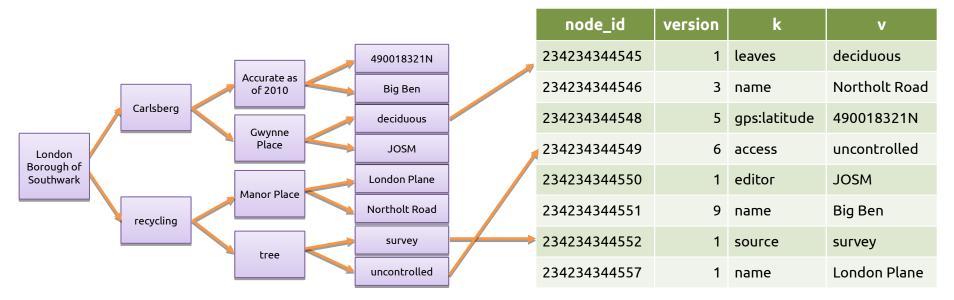
### Trick: ORDER BY NULL avoids filesort

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision GROUP BY substr(rev_timestamp, 5, 2)
= '09' G
id: 1
                                     There is no good index in this
 select type: SIMPLE
      table: revision
       type: ALL
                                                   case
possible keys: NULL
       key: NULL
    key len: NULL
       ref: NULL
       rows: 196824
      Extra: Using temporary; Using filesort
1 row in set (0.00 sec)
MariaDB [nlwiktionary] > EXPLAIN SELECT * FROM revision GROUP BY substr(rev_timestamp, 5, 2)
= '09' ORDER BY NULL\G
The advantage is not too big, but
       rows: 196871
      Extra: Using temporary
                                          it avoids the filesort
1 row in set (0.00 sec)
```

# Indexes and GROUP BY (rev\_page\_id)

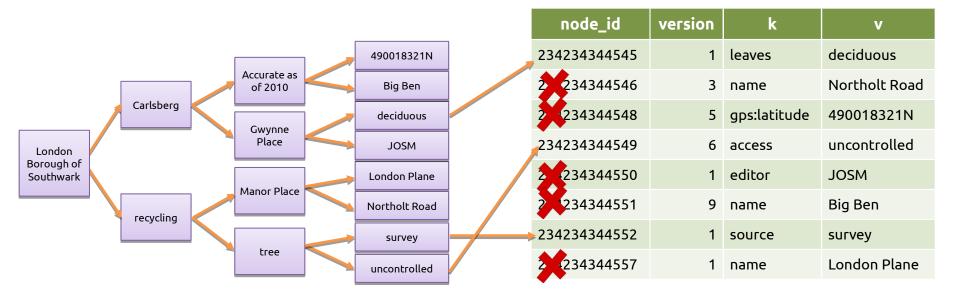
```
MariaDB [nlwiktionary] > EXPLAIN SELECT rev page, count(*)
revision GROUP BY rev page\G
********
                                 ********
          id: 1
 select type: SIMPLE
                            The index does not produce any
       table: revision
                            advantage for filtering (there is
        type: index
                                   no WHERE clause)
possible keys: NULL
         key: rev page id
     key len: 8
         ref: NULL
                             However, thanks to it we avoid a
        rows: 192388
                               sort and a temporary table
       Extra: Using index
1 row in set (0.00 sec)
```

### type: index, loose index scan and covering index (1/3)



With 'type:index', all rows are read in index order (full index scan)

### type: index, loose index scan and covering index (2/3)

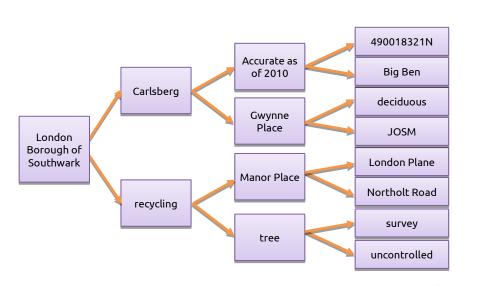


If we have in addition 'Using index for group-by' we have the loose index scan optimization

#### Loose Index Scan Example

```
MariaDB [nlwiktionary] > EXPLAIN SELECT rev page,
max(rev timestamp) FROM revision GROUP BY rev page\G
id: 1
 select type: SIMPLE
      table: revision
       type: range
possible keys: NULL
        key: page timestamp
     key len: 4
        ref: NULL
       rows: 9769
      Extra: Using index for group-by
1 row in set (0.00 sec)
```

type: index, loose index scan and covering index (3/3)





If we have in addition 'Using index' we have the covering index optimization

#### Covering Index Example (1/3)

```
MariaDB [nlwiktionary]> ALTER TABLE revision DROP INDEX rev page id, drop index page timestamp,
drop index page user timestamp;
                                                Let's start with no indexes
Query OK, 0 rows affected (0.05 sec)
Records: 0 Duplicates: 0 Warnings: 0
MariaDB [nlwiktionary]> EXPLAIN SELECT count(DISTINCT rev user) FROM revision WHERE rev page =
790\G
id: 1
 select type: SIMPLE
      table: revision
       type: ALL
possible keys: NULL
        key: NULL
     key len: NULL
        ref: NULL
       rows: 218384
      Extra: Using where
1 row in set (0.00 sec)
MariaDB [nlwiktionary] > SELECT count(DISTINCT rev user) FROM revision WHERE rev page = 790\G
count(DISTINCT rev user): 1
1 row in set (0.06 sec)
```

### Covering Index Example (2/3)

```
MariaDB [nlwiktionary] > ALTER TABLE revision ADD INDEX revision rev page(rev page);
Query OK, 0 rows affected (0.44 sec)
Records: 0 Duplicates: 0 Warnings: 0
MariaDB [nlwiktionary] > EXPLAIN SELECT count(DISTINCT rev user) FROM revision WHERE rev page =
790\G
id: 1
                                  Adding an index on rev_page
 select type: SIMPLE
      table: revision
                                    increases the speed due to
       type: ref
possible keys: revision rev page
                                         improved filtering
       key: revision rev page
    key len: 4
        ref: const
       rows: 4863
      Extra:
1 row in set (0.00 sec)
MariaDB [nlwiktionary]> SELECT count(DISTINCT rev_user) FROM revision WHERE rev_page = 790\G
count(DISTINCT rev user): 1
1 row in set (0.01 sec)
```

#### Covering Index Example (3/3)

```
MariaDB [nlwiktionary]> ALTER TABLE revision ADD INDEX revision_rev_page_rev_user(rev_page,
rev user):
Query OK, 0 rows affected (1.48 sec)
Records: 0 Duplicates: 0 Warnings: 0
MariaDB [nlwiktionary]> EXPLAIN SELECT count(DISTINCT rev user) FROM revision WHERE rev page =
790\G
id: 1
                                             rev_page, rev_user does not
 select type: SIMPLE
      table: revision
                                                   increase the index
       type: ref
possible keys: revision rev page, revision rev page rev i
                                              selectiveness, but allow to
       key: revision rev page rev user
    key len: 4
                                              return results directly from
        ref: const
                                                         the index
       rows: 4863
      Extra: Using index
1 row in set (0.00 sec)
MariaDB [nlwiktionary] > SELECT count(DISTINCT rev user) FROM revision WHERE rev page = 790\G
count(DISTINCT rev user): 1
                                      The speed difference can be huge
1 row in set (0.00 sec)
```

Query Optimization: From 0 to 10 (and up to 5.7)

#### **MULTI-COLUMN INDEXES**

# In many cases, conditions are applied on more than one column MariaDB [nlwiktionary]> EXPLAIN SELECT \* FROM revision WHERE

```
rev_page = 790 and rev_timestamp < '2008'\G
id: 1
 select type: SIMPLE
      table: revision
                    Assuming there were no previously
       type: ALL
                     created indexes, which would the
possible keys: NULL
                            optimal one be?
        kev: NULL
     key len: NULL
        ref: NULL
       rows: 686822
      Extra: Using where
1 row in set (0.00 sec)
```

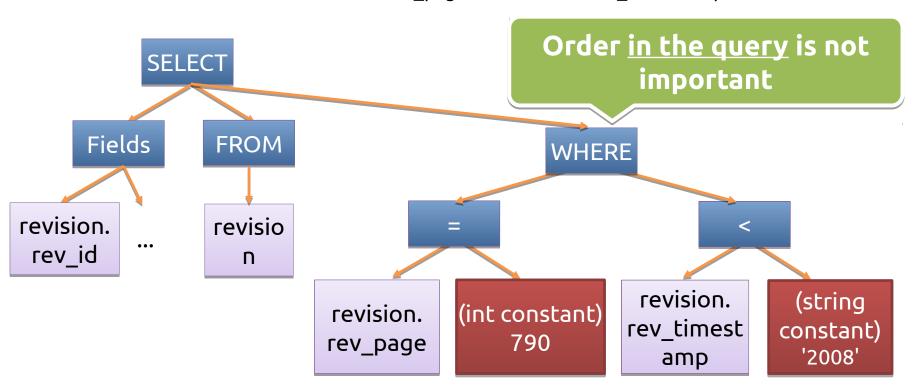
#### Options for indexes

- 1 index on column (rev\_page)
- 1 index on column (rev\_timestamp)
- 2 indexes, 1 on (rev\_page) and another on (rev\_timestamp)
- 1 multi-column index on (rev\_page, rev\_timestamp)
- 1 multi-column index on (rev\_timestamp,rev\_page)

Are these last 2 different from each other? Would it depend on the query order?

#### A brief reminder about query parsing

SELECT \* FROM revision WHERE rev\_page = 790 and rev\_timestamp < '2008'



#### Index on (rev\_page)

```
MariaDB [nlwiktionary]> ALTER TABLE revision ADD INDEX rev_page (rev_page);
Query OK, 0 rows affected (2.31 sec)
Records: 0 Duplicates: 0 Warnings: 0
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision WHERE rev_page = 790 and
rev timestamp < '2008'\G
id: 1
 select type: SIMPLE
                                          Query time improves
       table: revision
                                          significantly with this
       type: ref
possible keys: rev page
                                          index
        key: rev page
     key len: 4
        ref: const
                                          Less rows are
       rows: 4863
                                          scanned
      Extra: Using where
1 row in set (0.00 sec)
```

#### Adding (rev\_timestamp)

```
MariaDB [nlwiktionary] > ALTER TABLE revision ADD INDEX rev timestamp
(rev timestamp);
Query OK, 0 rows affected (1.77 sec)
Records: 0 Duplicates: 0 Warnings: 0
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision WHERE rev_page = 790
and rev timestamp < '2008'\G
In general, only one index can
         id: 1
 select type: SIMPLE
                                        be used per table access
       table: revision
       type: ref
                                        rev_page is preferred
possible keys: rev_page,rev_timestamp
        key: rev_page
                                         over rev_timestamp
     key len: 4
        ref: const
        rows: 4863
       Extra: Using where
1 row in set (0.01 sec)
```

# Forcing the use of (rev\_timestamp)

```
MariaDB [nlwiktionary] > EXPLAIN SELECT * FROM revision FORCE
INDEX(rev timestamp) WHERE rev page = 790 and rev timestamp <</pre>
'2008'\G
id: 1
 select type: SIMPLE
       table: revision
                                 It is a range
       type: range
                                   access
possible_keys: rev timestamp
        key: rev timestamp
     key len: 14
                             A lot more accessed rows
        ref: NULL
        rows: 343411
       Extra: Using index condition; Using where
1 row in set (0.00 sec)
```

# Adding (rev\_page, rev\_timestamp)

```
MariaDB [nlwiktionary] > ALTER TABLE revision ADD INDEX
rev page rev timestamp(rev page, rev timestamp);
Query OK, 0 rows affected (1.59 sec)
Records: 0 Duplicates: 0 Warnings: 0
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision WHERE rev_page = 790
and rev timestamp < '2008'\G
id: 1
 select type: SIMPLE
       table: revision
        type: range
possible_keys: rev_page,rev_timestamp,rev_page_rev_timestamp
        key: rev page rev timestamp
     key len: 18
                                        Reduced number of
        ref: NULL
        rows: 1048
                                            rows scanned
       Extra: Using index condition
1 row in set (0.00 sec)
```

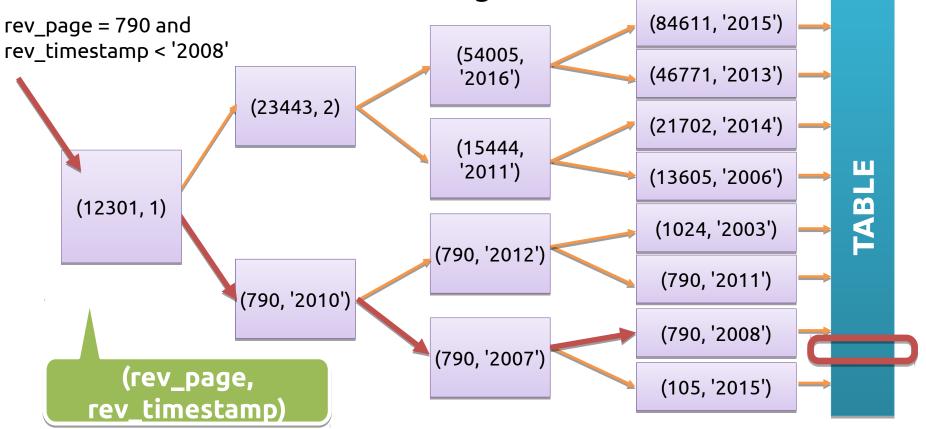
### Is (rev\_timestamp, rev\_page) a better option?

```
MariaDB [nlwiktionary] > ALTER TABLE revision ADD INDEX rev_timestamp_rev_page
(rev timestamp, rev page);
Query OK, 0 rows affected (1.76 sec)
Records: 0 Duplicates: 0 Warnings: 0
MariaDB [nlwiktionary] > EXPLAIN SELECT * FROM revision WHERE rev_page = 790 and
rev_timestamp < '2008'\G</pre>
id: 1
 select type: SIMPLE
       table: revision
        type: range
possible_keys: rev_page,rev_timestamp,rev_page_rev_timestamp,rev_timestamp_rev_page
         key: rev_page_rev_timestamp
     key len: 18
                                              Previous index is still
         ref: NULL
        rows: 1048
                                                 preferred, why?
       Extra: Using index condition
1 row in set (0.00 sec)
```

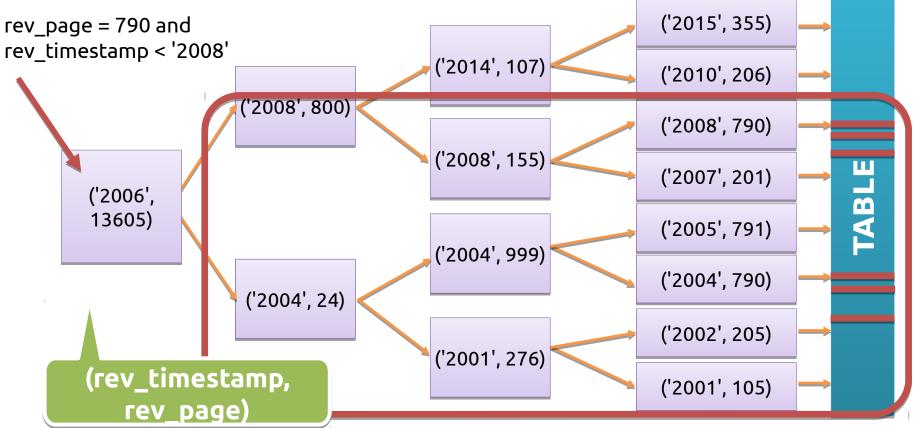
### Forcing (rev\_timestamp, rev\_page)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision FORCE
INDEX(rev timestamp rev page) WHERE rev page = 790 and
rev timestamp < '2008'\G
********
                                 ********
          id: 1
 select type: SIMPLE
       table: revision
        type: range
possible keys: rev_timestamp_rev_page
         key: rev_timestamp rev page
     kev len: 18
                                Only the first column is being
         ref: NULL
                                used effectively for filtering
        rows: 343411
       Extra: Using index condition
1 row in set (0.00 sec)
```

A compound index produces a single tree ordered by the 2 values



The alternative index cannot be used in such an effective way



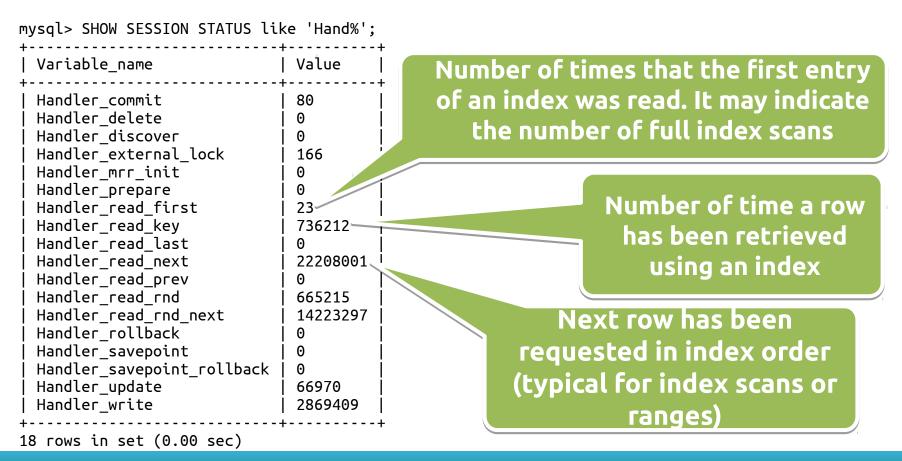
#### Order and column selection

- Range access using >, <, >=, <=, BETWEEN can only be filtered once effectively, at the end of an index
- When selecting indexes, prefer columns with high cardinality (very selective)
  - The optimal index can depend on the constants used

#### Use of "Handler\_\*" statistics

- They are post-execution statistics at row level
  - Unlike EXPLAIN's "rows" column, they are exact, not a guess
  - They allow to compare query execution performance in a deterministic way, independently of the execution time

### "Handler" Stats (indexed)



### "Handler" Stats (unindexed)

mysql> SHOW SESSION STATUS like 'Hand%'; Variable name Value Handler commit 80 Handler delete Handler\_discover Handler external lock 166 Handler mrr init Handler\_prepare Handler\_read\_first 23 Handler read key 736212 Handler read last Handler read next 22208001 Handler read prev Handler read rnd 665215 Handler read rnd next 14223297 Handler rollback Handler savepoint Handler\_savepoint\_rollback Handler update 66970 Handler write 2869409 18 rows in set (0.00 sec)

A row has been requested in a specific position (typical for joins or order by without indexes)

Request tp read the next row in "table order" (typical for full table scans)

**Insertions in SELECTS may indicate temporary tables** 

## Comparing statistics of the previous indexes (no indexes)

```
MariaDB [nlwiktionary]> FLUSH STATUS;
Query OK, 0 rows affected (0.00 sec)
```

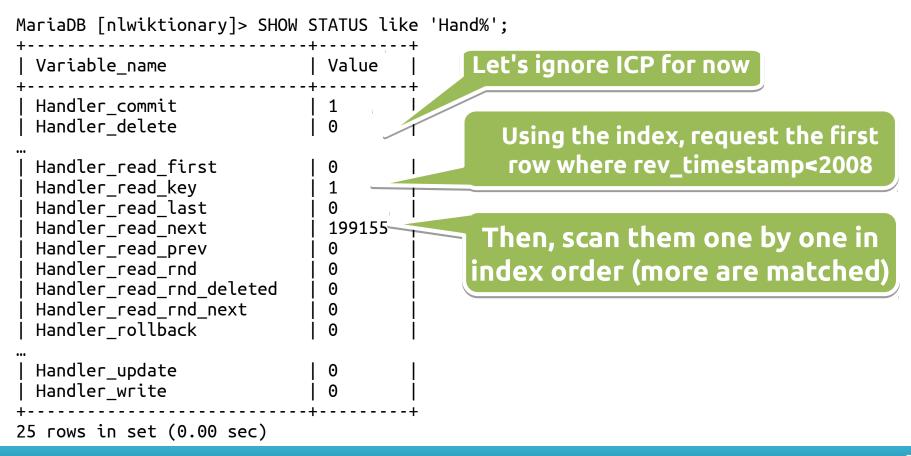
```
MariaDB [nlwiktionary]> SELECT * FROM revision IGNORE INDEX(rev_page, rev_timestamp, rev_page_rev_timestamp, rev_timestamp_rev_page) WHERE rev_page = 790 and rev_timestamp < '2008'; 1049 rows in set (0.58 sec)
```

```
MariaDB [nlwiktionary]> SHOW STATUS like 'Hand%';
 Variable_name
                          | Value
 Handler_commit
 Handler delete
 Handler read first
 Handler read key
 Handler read last
 Handler_read_next
 Handler_read prev
 Handler read rnd
 Handler_read_rnd_deleted
                                                           Typical result for a
 Handler read rnd next
                            820305
 Handler_rollback
                            0
                                                               full table scan
 Handler update
 Handler_write
25 rows in set (0.00 sec)
```

### Index on (rev\_page)

```
MariaDB [nlwiktionary]> SHOW STATUS like 'Hand%';
 Variable name
                            | Value
 Handler_commit
 Handler delete
                                           Using the index, request the
                                            first row with rev_page=790
 Handler_read_first
                             0
 Handler read key
 Handler read last
                             0
 Handler read next
                             4864
                                          Then, scan them one by one in
 Handler read prev
                             0
 Handler_read_rnd
                                                     index order
 Handler_read_rnd_deleted
 Handler read rnd next
                             0
 Handler rollback
 Handler_update
 Handler_write
25 rows in set (0.01 sec)
```

### Index on (rev\_timestamp)



# Index on (rev\_page, rev\_timestamp)

```
MariaDB [nlwiktionary]> SHOW STATUS like 'Hand%';
  Variable name
                            | Value
                                             With both conditions covered, we
  Handler_commit
                                              can find the actual first row that
 Handler delete
                                              matches the condition using the
 Handler_read_first
                             0
                                                            index
 Handler read key
                             1
 Handler read last
                             0
 Handler_read_next
                             1049
                                          Rows scanned == Rows returned
 Handler read prev
 Handler_read_rnd
 Handler_read_rnd_deleted
                             0
 Handler read rnd next
                             0
 Handler rollback
 Handler_update
  Handler_write
25 rows in set (0.00 sec)
```

## Index on (rev\_timestamp, rev\_page), no ICP

```
MariaDB [nlwiktionary] > SHOW STATUS like 'Hand%';
  Variable name
                              | Value
  Handler_commit
  Handler delete
  Handler_read_first
                                0
  Handler read key
  Handler read last
  Handler_read_next
                                199155
  Handler read prev
  Handler_read_rnd
  Handler_read_rnd_deleted
  Handler read rnd next
  Handler rollback
  Handler update
                                0
  Handler_write
25 rows in set (0.00 sec)
```

Assuming no ICP, exact same results than with (rev\_timestamp). The extra column does not help. Also, EXPLAIN's row count was very off.

#### Redundant Indexes

- Creating all 4 previous indexes in production is not a great idea
  - "Left-most index prefix" allows, for example (rev\_page, rev\_timestamp) doing everything you can do with (rev\_page)
  - If two indexes have equal selectivity,
     MySQL chooses the shortest one

#### "Left-most index" Example

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision FORCE
INDEX(rev_page_rev_timestamp) WHERE rev page = 790\G
id: 1
 select type: SIMPLE
       table: revision
       type: ref
possible keys: rev page rev times tames
        key: rev page rev time
                             Only the first column is used
     key len: 4 —
        ref: const
       rows: 4863
      Extra:
1 row in set (0.00 sec)
```

#### **Duplicate Indexes**

- It is very easy to create indexes with the same exact definition (same columns and ordering)
  - Set a convention for index naming (e.g tablename\_column1\_column2\_idx) – MySQL does not allow 2 indexes with the same identifier
  - Since MySQL 5.6, an warning is thrown if a duplicate index is created

#### pt-duplicate-index-checker

```
$ pt-duplicate-key-checker h=127.0.0.1,P=5621,u=msandbox,p=msandbox
[...]
# rev_timestamp is a left-prefix of rev_timestamp_rev_page
                                                              Simple tool to check
# Key definitions:
   KEY `rev_timestamp` (`rev_timestamp`),
                                                                   redundant and
   KEY `rev_timestamp_rev_page` (`rev_timestamp`,`rev_page`)
# Column types:
                                                                 duplicate indexes
    `rev timestamp` binary(14) not null default '\0\0\0\0\0\0\0\0
    `rev_page` int(10) unsigned not null
# To remove this duplicate index, execute:
ALTER TABLE `nlwiktionary`.`revision` DROP INDEX `rev timestamp`:
# rev_page is a left-prefix of rev_page_rev_timestamp
# Key definitions:
   KEY `rev_page` (`rev_page`),
   KEY `rev_page_rev_timestamp` (`rev_page`,`rev_timestamp`),
# Column types:
    `rev_page` int(10) unsigned not null
    `rev timestamp` binary(14) not null default '\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0\0
# To remove this duplicate index, execute:
ALTER TABLE 'nlwiktionary'.'revision' DROP INDEX 'rev page';
  # Size Duplicate Indexes
                       15478317
# Total Duplicate Indexes
# Total Indexes
                       285
```

### "OR"-style conditions over the same column

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision WHERE
rev_page = 790 OR rev_page = 795 OR rev page = 1024\G
id: 1
                               Equivalent to:
 select type: SIMPLE
                        SELECT * FROM revision WHERE
       table: revision
                          rev_page IN (790, 795, 1024)
        type: range
possible keys: rev_page, rev_page_rev_earner
        kev: rev page
     key len: 4
        ref: NULL
        rows: 4890
       Extra: Using index condition; Using where
1 row in set (0.01 sec)
```

### Handlers on "IN" / "OR" conditions over the same column

4	L
Variable_name	Value
Handler_commit   Handler_delete	1
" Handler_prepare Handler_read_first Handler_read_key Handler_read_last Handler_read_next Handler_read_prev Handler_read_rnd Handler_read_rnd_deleted Handler_read_rnd_next Handler_rollback	0   0   3   0   4891   0   0   0   0   0   0   0   0   0
Handler_update   Handler_write  +	0
25 rows in set (0.00 sec)	

Despite identifying themselves as "range"s, the execution is slightly different, one index dive (similar to a ref) is done per value. This can be an issue in conditions with thousands of items.

### "OR"-style conditions over different columns

- We cannot use a single index efficiently for both conditions
  - We can scan both conditions separatelly and mix the results, discarding duplicates
  - Or use an index for one condition and not for the other
  - Index merge allows to use two indexes for a single table access simultaneously

### Index Merge Example

```
MariaDB [nlwiktionary] > EXPLAIN SELECT * FROM revision WHERE
rev_page = 790 or rev_timestamp < '2004'\G</pre>
id: 1
  select type: SIMPLE
       table: revision
        type: index merge
possible_keys: rev_page,rev_timestamp,rev_page_rev_timestamp,
rev timestamp rev page
         key: rev page, rev timestam
                                    Both indexes are used, then
     kev len: 4.14
                                    combined using the "union"
         ref: NULL
                                           operation
        rows: 4871
       Extra: Using sort_union(rev_page,rev_timestamp); Using
where
1 row in set (0.00 sec)
```

#### **INDEX\_MERGE** Issues

- Sometimes it is faster to to execute the sentence using UNION:
  - This is specially true with (UNION ALL) since MySQL 5.7, if you do not care about duplicates
- There are also interseccion merges, but multi-column indexes are preferred

### Disabling optimizer features (I)

 The optimizer\_switch variable allows enabling and disablig globally or per session many query optimizer features:

# Deshabilitar características del optimizador (II)

```
MariaDB [nlwiktionary] > SET optimizer switch='index merge sort union=off';
Query OK, 0 rows affected (0.00 sec)
MariaDB [nlwiktionary] > EXPLAIN SELECT * FR
                                          This will only have effect for
or rev timestamp < '2004'\G
                                               the current session.
****** 1. TOW ******
          id: 1
 select type: SIMPLE
       table: revision
        type: ALL
possible keys:
rev page, rev timestamp, rev page rev timestamp, rev timestamp rev page
         key: NULL
     key len: NULL
         ref: NULL
        rows: 686822
       Extra: Using where
1 row in set (0.00 sec)
```

### What happens if we have two ranges?

- As seen previously, we cannot use efficiently two range types on the same table access. Alternatives:
  - Use only one index for the most selective column
  - Use index condition pushdown to get an advantage
  - Change one of the two ranges into a discrete "IN" comparison/bucketize with a new column
  - Use quadtrees or R-TREEs (spatial indexing)

### Example of Bucketizing (I)

```
MariaDB [nlwiktionary]> EXPLAIN SELECT count(*) FROM revision
WHERE rev timestamp < '2008' AND rev len > 5500\G
id: 1
 select type: SIMPLE
       table: revision
        type: ALL
possible keys: rev timestamp,rev_timestamp_rev_page
        key: NULL
     key len: NULL
                              Looks like only an index on
        ref: NULL
        rows: 686822
                             (rev_timestamp) or (rev_len)
       Extra: Using where
                             would be useful as we have 2
1 row in set (0.00 sec)
                                       ranges.
```

### Example of Bucketizing (II)

```
MariaDB [nlwiktionary] > ALTER TABLE revision ADD COLUMN
rev len cat int;
Query OK, 0 rows affected (38.28 sec)
Records: 0 Duplicates: 0 Warnings: 0
MariaDB [nlwiktionary]> UPDATE revision set rev_len_cat =
IF(rev len < 10000, rev_len div 1000, 10);</pre>
Query OK, 820308 rows affected (15.19 sec)
Rows matched: 820308 Changed: 820308 Warnings: 0
MariaDB [nlwiktionary]> ALTER TABLE revision ADD INDEX
rev_len_cat_rev_timestamp (rev_len_cat, rev_timestamp);
Query OK, 0 rows affected (2.11 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

### Example of Bucketizing (III)

```
MariaDB [nlwiktionary] > EXPLAIN SELECT count(*) FROM revision WHERE
rev_timestamp < '2008' AND rev_len > 5500 AND rev_len_cat IN (5, 6,
7, 8, 9, 10)\G
******* 1. row
                                *********
          id: 1
 select type: SIMPLE
       table: revision
        type: range
possible kevs:
rev_timestamp,rev_timestamp_rev_page,rev_len_cat_rev_timestamp
         key: rev len cat rev timestamp
     key len: 19
         ref: NULL
                                 We did some transformations
        rows: 4442
                                 to both the structure and the
       Extra: Using where
                                             query.
1 row in set (0.00 sec)
```

### Example of Index Condition Pushdown

```
MariaDB [nlwiktionary] > ALTER TABLE revision ADD INDEX
rev_len_rev_timestamp(rev_len, rev_timestamp);
Query OK, 0 rows affected (1.77 sec)
Records: 0 Duplicates: 0 Warnings: 0
MariaDB [nlwiktionary]> SET optimizer_switch='index_condition_pushdown=on'; EXPLAIN
SELECT * FROM revision WHERE rev_timestamp < '2008' AND rev_len > 5500\G
Query OK, 0 rows affected (0.00 sec)
id: 1
                                       Index condition pushdown
 select type: SIMPLE
                                      (ICP) eanbles the engines to
       table: revision
       type: range
                                      use extra parts of the index
possible keys: rev timestamp, rev timestamp
        key: rev len rev timestamp
                                       while avoiding costly row
     key len: 5
                                      movements to and from the
        ref: NULL
       rows: 38744
                                                SQL layer
       Extra: Using index condition
```

1 row in set (0.00 sec)

#### ICP Issues

- Differences in execution time is more significative when the extra column condition is very selective (getting 5x the original performance)
- ICP is ignored when using covering Index, making the performance worse

## Does LIMIT improve the performance? (I)

```
MariaDB [nlwiktionary] > EXPLAIN SELECT * FROM page ORDER BY page touched\G
type: ALL
possible keys: NULL
       key: NULL
    key len: NULL
       ref: NULL
      rows: 90956
      Extra: Using filesort
MariaDB [nlwiktionary] > EXPLAIN SELECT * FROM page ORDER BY page_touched LIMIT 10\G
type: index
                                      In some cases it can be
possible keys: NULL
       key: page_page_touched
                                        essential to allow
    key_len: 14
       ref: NULL
                                      effective usage of the
      rows: 10
                                             indexes
      Extra:
1 row in set (0.00 sec)
```

## Does LIMIT improve the performance? (II)

```
MariaDB [nlwiktionary] > EXPLAIN SELECT * FROM revision ORDER BY rev_comment\G
[...] type: ALL
possible_keys: NULL
       key: NULL
    key len: NULL
       ref: NULL
       rows: 817636
      Extra: Using filesort
1 row in set (0.00 sec)
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision ORDER BY rev_comment LIMIT 10\G
table: revision
                               In other cases, it has no
      type: ALL
possible keys: NULL
                                effect on the scanned
       key: NULL
    key len: NULL
                                  rows (just on the
       ref: NULL
                                   returned ones)
       rows: 817636
      Extra: Using filesort
1 row in set (0.00 sec)
```

# Does LIMIT improve the performance? (I)

```
MariaDB [nlwiktionary] > EXPLAIN SELECT * FROM page ORDER BY page_title LIMIT 100\G
type: index
possible_keys: NULL
       key: page_title
    key len: 257
       ref: NULL
      rows: 100
      Extra:
1 row in set (0.00 sec)
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM page ORDER BY page_title LIMIT 10000, 100\G
type: ALL
possible_keys: NULL
                                         In this case,
       kev: NULL
    key len: NULL
                                    performance will vary
       ref: NULL
                                      depending on the
       rows: 90956
      Extra: Using filesort
                                      offset (not ideal)
1 row in set (0.00 sec)
```

## Can we filter and sort at the same time using indexes?

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision WHERE
rev comment='' ORDER BY rev timestamp ASC\G
id: 1
 select type: SIMPLE
       table: revision
                      This query is slow
        type: ALI
possible keys: NULL
                     because a) the full
        key: NULL
                         table scan
     key len: NULL
        ref: NULL
        rows: 817636
                                       b) Required sort
       Extra: Using where; Using filesort
                                        after filtering
1 row in set (0.00 sec)
```

# Adding and index on (rev\_comment, rev\_timestamp)

```
MariaDB [nlwiktionary] > ALTER TABLE revision ADD INDEX
rev comment rev timestamp (rev comment, rev timestamp);
Query OK, 0 rows affected (3.19 sec)
Records: 0 Duplicates: 0 Warnings: 0
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision WHERE rev_comment=''
ORDER BY rev timestamp ASC\G
id: 1
 select type: SIMPLE
       table: revision
                                   Both type: ALL and
       type: ref
possible keys: rev comment rev times.
                                       filesort have
         key: rev comment rev timestar
                                       disappeared
     key len: 769
        ref: const
        rows: 266462
       Extra: Using index condition; Using where
1 row in set (0.00 sec)
```

### This is not always possible

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision WHERE
rev len > 5500 ORDER BY rev timestamp ASC\G
id: 1
 select type: SIMPLE
                              The range makes impossible
       table: revision
                               to use the index optimally
        type: range
possible keys: rev len rev timest
                              for the the ORDER BY: either
         key: rev_len_rev_timest
                               we filter (rev_len) or sort
     key len: 5
                                   (rev_timestamp)
         ref: NULL
        rows: 38744
       Extra: Using index condition; Using fivesort
1 row in set (0.00 sec)
```

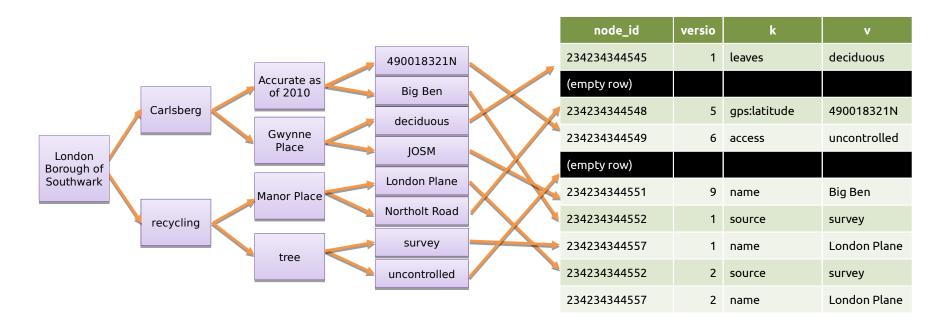
## A Strange Game. The Only Winning Move is Not to Play

```
mysql-5.6.21 (osm) > SELECT * FROM nodes FORCE
INDEX(version idx) WHERE version < 15 ORDER BY changeset_id;</pre>
/* type: range, Using filesort */
2859673 rows in set (30.58 sec)
mysql-5.6.21 (osm) > SELECT * FROM nodes FORCE
INDEX(changeset id idx) WHERE version < 15 ORDER BY</pre>
changeset id:
/* type: index */
2859673 rows in set (30.92 sec)
mysql-5.6.21 (osm) > SELECT * FROM nodes WHERE version < 15
ORDER BY changeset id;
/* type: ALL, Using filesort */
2859673 rows in set (16.54 sec)
```

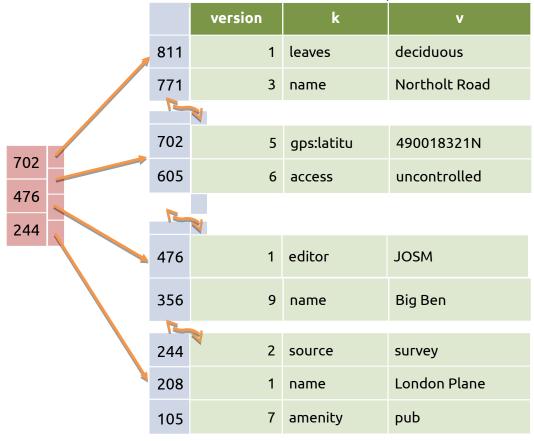
#### **MyISAM Internals**

Index (part of revision.MYI)

Data (revision.MYD)

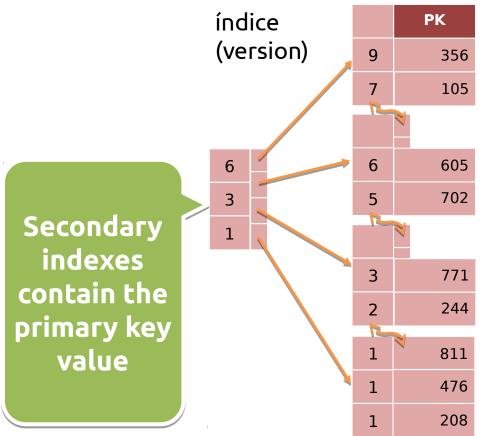


#### InnoDB Internals (PRIMARY)

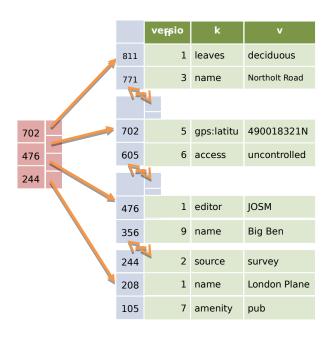


Data clustered always using the primary key

### InnoDB Internals (Secondary)



PK / Datos



## Consequences of using InnoDB (I)

- Every table should have a primary key
  - If one is not defined, MySQL will choose an available NOT NULL unique key
  - If that is not possible, an internal 6-byte row identifier will be generated (not user-accesible)

# Consequences of using InnoDB (II)

- Inserting in primary key order is much faster
  - Less fragmentation/page-split
  - Usage of "batch" mode, improving insert speed
- Using auto-increment keys as primary keys can be a good idea for InnoDB

# Consequences of using InnoDB (III)

- A very long primary key may increment substantially the size of secondary keys
  - Int or bigint types are recommended instead of UUIDs or other long strings

#### Differences in size

```
mysql-5.6.21 (osm) > CREATE
                              ■ .ibd size (no secondary indexes) ■ .ibd size (with secondary indexes)
TABLE pk int (id int
PRIMARY KEY auto increment,
                                  42500000
a int.
b int.
c int,
                                  34000000
d int);
                                                                                            32,505,856
Query OK, 0 rows affected
(0.16 \text{ sec})
                                  25500000
mysql-5.6.21 (osm) > CREATE
TABLE pk_uuid (id char(36)
                                                                               20,971,520
                                   17000000
PRIMARY KEY,
a int,
                                                              13,631,488
b int,
                                                12,582,912
c int,
                                    8500000
d int);
Query OK, 0 rows affected
(0.04 \text{ sec})
                                                         pk int
                                                                                       pk uuid
```

# Extended primary key optimization

- As the primary key is part of all seconday keys, this can be used "for free":
  - For row filtering (since MySQL 5.6)
  - To return results in primary key order
  - To avoid reading data from the table (covering index)

## Extended Primary Key Example

```
mysql-5.6.21 (osm) > EXPLAIN SELECT node id FROM nodes WHERE
changeset id = 24284 and node id <> 146472942\G
id: 1
 select type: SIMPLE
      table: nodes
       type: range
possible keys: PRIMARY,changeset_id_idx
        key: changeset id idx
     key len: 16
        ref: NULL
       rows: 50
      Extra: Using where; Using index
1 row in set (0.07 sec)
```

Query Optimization: From 0 to 10 (and up to 5.7)

#### **FULLTEXT SEARCH**

# Fuzzy Search of "gloucester/Gloucester's/etc"

"Typical" way to solve this:

```
mysql-5.7.5 (osm) >
SELECT way id as id, v
FROM way tags
WHERE v like '%gloucester%';
425 rows in set (0.46 sec)
                                   Too slow
```

#### Let's Add an Index

```
mysql-5.7.5 (osm) > ALTER TABLE way_tags ADD INDEX(v);
Query OK, 0 rows affected (6.44 sec)
Records: 0 Duplicates: 0 Warnings: 0
                                        Still slow, why?
mysql-5.7.5 (osm) > SELECT ...;
425 rows in set (0.38 sec)
mysql-5.7.5 (osm) > EXPLAIN SELECT way as type, way_id as
id, v FROM way_tags WHERE v like '%gloucester%';
 id |.| type | possible_keys | key | key_len | ref | rows | filtered | Extra
 1 row in set, 1 warning (0.01 sec)
```

#### Fulltext Index

```
mysql-5.7.5 (osm) > ALTER TABLE way_tags add FULLTEXT index(v);
Query OK, 0 rows affected (3.20 sec)
Records: 0 Duplicates: 0 Warnings: 0
mysql-5.7.5 (osm) > SELECT ...;
425 rows in set (0.00 sec)
mysql-5.7.5 (osm) > EXPLAIN SELECT 'way' as type, way_id as id, v
FROM way tags WHERE MATCH(v) AGAINST ('+gloucester*' IN BOOLEAN
MODE):
   ------
1 row in set, 1 warning (0.00 sec)
```

#### Newer Fulltext Optimizations

```
mysql-5.5.40 (osm) > EXPLAIN SELECT count(*) FROM way_tags_myisam WHERE MATCH(v) AGAINST('gloucester');
 1 row in set (0.00 sec)
mysql-5.5.40 (osm) > SHOW STATUS like 'Hand%';
Handler read first
 Handler_read_key
 Handler_read_last
 Handler_read_next
Handler_read_prev
Handler_read_rnd
Handler_read_rnd_next
Handler_update
Handler_write
16 rows in set (0.00 sec)
```

# Newer Fulltext Optimizations (cont.)

```
mysql-5.7.5 (osm) > EXPLAIN SELECT count(*) FROM way_tags WHERE MATCH(v) AGAINST('gloucester');
   id | select_type | table | partitions | type | possible_keys | key | key_len | ref | rows | filtered | Extra
   1 | SIMPLE | NULL | NULL | NULL | NULL | NULL | NULL |
                                         | NULL | NULL | NULL | Select tables optimized away |
  1 row in set, 1 warning (0.00 sec)
mysql-5.7.5 (osm) > SHOW STATUS like 'Hand%';
Variable_name
                                                            counting
Handler commit
directly
Handler read first
                                                            from the
Handler_read_key
Handler read last
Handler read next
                                                            FULLTEXT
Handler read prev
| Handler_read_rnd
                                                              index
Handler_read_rnd_next
Handler update
.
| Handler_write
18 rows in set (0.00 sec)
```

#### **Open Issues and Limitations**

- No postfix support (wildcards)
- Simple Ranking (and different from MyISAM)
- No stemming support
- Some multi-language limitations

More on FULLTEXT InnoDB support:

http://www.drdobbs.com/database/full-text-search-with-innodb/231902587

#### **Alternatives**

- Apache Lucene
  - Solr
  - Elasticsearch
- Sphinx
  - SphinxSE

Query Optimization: From 0 to 10 (and up to 5.7)

JOINS

## (Block) Nested Loop Join

• Until MySQL 5.5 there was only one algorithm to execute a JOIN:

				node_id	version	k	v
node_id	version	lat	lon	1	1	name	Big Benn
1	1	52	0.5	1	1	tourism	attraction
1	2	52	0.5	1	2	name	Big Ben
2	1	51	1	1	2	tourism	attraction
3	1	53	1.5	3	1	name	London Eye

## Extra Access type: eq\_ref

```
mysql-5.6.21 (osm) > EXPLAIN SELECT * FROM nodes JOIN node tags
                USING(node id, version) WHERE node tags.v= 'Big Ben'\G
id: 1
 select type: SIMPLE
      table: node tags
       type: ref
possible keys: PRIMARY, v idx
       key: v idx
    key len: 767
        ref: const
       rows: 1
      Extra: Using where; Using index
id: 1
 select type: SIMPLE
      table: nodes
       type: eq ref
possible keys: PRIMARY, version idx
        key: PRIMARY
    key len: 16
        ref: osm.node tags.node id,osm.node tags.version
       rows: 1
      Extra: NULL
2 rows in set (0.00 sec)
```

eq\_ref is similar to ref, but allows faster JOINS because, by using a unique key, it only has to search one row for each previous result

#### JOIN Optimization

- Two main goals:
  - Perform an effective filtering on each table access, if possible using indexes
  - Perform the access in the most efficient table order
- When joining 3 or more tables in a star schema, the "covering index" strategy can have a huge impact

#### **Example: optimize this JOIN (I)**

```
SELECT n.node id, n.latitude, n.longitude
  FROM way nodes w n
  JOIN way tags w t
   ON w_n.way_id = w_t.way_id and
      w_n.version = w_t.version
  JOIN nodes n
    ON w n.node id = n.node id
  JOIN node tags n t
    ON n.node_id = n_t.node_id and
       n.version = n_t.version
WHERE w_t.k = 'building' and
       n_t.k = 'entrance' and
       n t.v = 'main';
```

We start without secondary indexes

## Example: optimize this JOIN (II)

```
id: 1
       id: 1
 select type: SIMPLE
                                         select type: SIMPLE
                                              table: n t
     table: w t
                                               type: ref
      type: index
                                        possible keys: PRIMARY
possible_keys: PRIMARY
                                                key: PRIMARY
      key: PRIMARY
                                             key len: 8
    key len: 783
                                                ref: osm.w n.node id
       ref: NULL
                                               rows: 1
      rows: 1335702
                                              Extra: Using where
     Extra: Using where; Using index
id: 1
       id: 1
                                         select type: SIMPLE
 select type: SIMPLE
                                              table: n
     table: w_n
                                               type: eq ref
      type: ref
                                        possible keys: PRIMARY
possible keys: PRIMARY
                                                key: PRIMARY
      key: PRIMARY
                                             key len: 16
    key len: 16
                                                ref: osm.w n.node id,osm.n t.version
      ref: osm.w_t.way_id,osm.w_t.version
                                               rows: 1
      rows: 3
                                              Extra: Using index
     Extra: NULL
                                        4 rows in set (0.01 sec)
mysql-5.6.21 (osm) > SELECT ...
```

858 rows in set (9.00 sec)

## Example: optimize this JOIN (III)

```
mysql-5.6.21 (osm) > ALTER TABLE way_tags ADD
INDEX k_idx(k);
Query OK, 0 rows affected (4.80 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

Creating an index on way\_tags.k

## Example: optimize this JOIN (IV)

```
id: 1
 select type: SIMPLE
      table: w t
       type: ref
possible_keys: PRIMARY,k_idx
        key: k idx
     key len: 767
        ref: const
       rows: 452274
      Extra: Using where; Using index
******************* 2. FOW **************
         id: 1
 select_type: SIMPLE
      table: w_n
       type: ref
possible keys: PRIMARY
        key: PRIMARY
     key len: 16
        ref: osm.w_t.way_id,osm.w_t.version
       rows: 3
       Fxtra: NULL
mysql-5.6.21 (osm) > SELECT ...
858 rows in set (8.58 sec)
```

```
id: 1
 select type: SIMPLE
      table: n t
      type: ref
possible keys: PRIMARY
       key: PRIMARY
    key len: 8
       ref: osm.w n.node id
      rows: 1
      Extra: Using where
id: 1
 select type: SIMPLE
      table: n
      type: eq ref
possible keys: PRIMARY
       key: PRIMARY
    key len: 16
       ref: osm.w n.node id,osm.n t.version
      rows: 1
      Extra: NULL
4 rows in set (0.00 sec)
                   It seems like the index
                      is not very useful
```

## Example: optimize this JOIN (V)

```
mysql-5.6.21 (osm) > ALTER TABLE node_tags ADD
INDEX k_idx(k);
Query OK, 0 rows affected (2.82 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

The order does not seem to be adequate, let's try adding an index to start by accessing node\_tags

#### **Example: optimize this JOIN (VI)**

```
id: 1
        id: 1
                                            select type: SIMPLE
 select type: SIMPLE
                                                                            It keeps using
                                                  table: n t
      table: w t
                                                  type: ref
      type: ref
                                                                              the wrong
                                          possible keys: PRIMARY,k_idx
possible_keys: PRIMARY,k_idx
                                                   key: PRIMARY
                                                                             order, even if
       key: k idx
                                                kev len: 8
    key len: 767
                                                   ref: osm.w n.node id
                                                                            we delete the
       ref: const
                                                  rows: 1
       rows: 452274
                                                  Extra: Using where
                                                                               w_t.k_idx
                                           ****** 4. FOW ******
      Extra: Using where; Using index
******************* 2. FOW **************
                                                    id: 1
                                                                                  index
                                            select type: SIMPLE
        id: 1
                                                  table: n
 select_type: SIMPLE
                                                  type: eq ref
      table: w n
                                          possible keys: PRIMARY
      type: ref
                                                   key: PRIMARY
possible keys: PRIMARY
                                                kev len: 16
       key: PRIMARY
                                                   ref: osm.w n.node id,osm.n t.version
    key len: 16
                                                  rows: 1
       ref: osm.w_t.way_id,osm.w_t.version
                                                  Extra: NULL
                                           4 rows in set (0.00 sec)
      rows: 3
      Fxtra: NULL
```

mysql-5.6.21 (osm) > SELECT ... 858 rows in set (7.33 sec)

## Example: optimize this JOIN (VII)

```
SELECT STRAIGHT_JOIN n.node_id, n.latitude, n.longitude
  FROM node tags n t
  JOIN nodes n
    ON n.node id = n_t.node_id and
       n.version = n t.version
  JOIN way_nodes w_n
    ON w n.node id = n.node id
  JOIN way tags w t
    ON w n.way id = w t.way id and
       w n.version = w t.version
 WHERE w_t.k = 'building' and
       n t.k = 'entrance' and
       n t.v = 'main';
```

Let's see why rewriting it into this query

## Example: optimize this JOIN (VIII)

```
id: 1
                                                    id: 1
                                              select type: SIMPLE
                                                                There is no index
 select type: SIMPLE
                                                  table: w_n
      table: n t
                                                               on w_n that would
                                                   type: ALL
       type: ref
possible keys: PRIMARY,k_idx
                                            possible_keys: PRIMARY
                                                                 allow efficient
                                                   key: NULL
        key: k idx
                                                key_len: NULL
    key len: 767
                                                                       access
        ref: const
                                                   ref: NULL
       rows: 2390
                                                   rows: 3597858
      Extra: Using index condition; Using where
                                                  Extra: Using where; Using join buffer (Block
Nested Loop)
                                            id: 1
                                                    id: 1
 select type: SIMPLE
      table: n
                                              select type: SIMPLE
                                                  table: w_t
       type: eq ref
possible keys: PRIMARY
                                                   type: eq ref
        key: PRIMARY
                                            possible keys: PRIMARY
    key len: 16
                                                   key: PRIMARY
        ref: osm.n t.node id,osm.n t.version
                                                key len: 783
       rows: 1
                                                   ref: osm.w_n.way_id,osm.w_n.version,const
      Extra: NULL
                                                   rows: 1
                                                  Extra: Using where; Using index
                                            4 rows in set (0.00 sec)
```

## Example: optimize this JOIN (IX)

```
mysql-5.6.21 (osm) > ALTER TABLE way_nodes ADD
INDEX node_id_idx(node_id);
Query OK, 0 rows affected (17.77 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

#### **Example: optimize this JOIN (X)**

```
id: 1
 select type: SIMPLE
                        Now it starts by
     table: n_t
                         the right table
      type: ref
possible keys: PRIMARY,k idx
                             (without
       key: k_idx
    key len: 767
                       STRAIGHT_JOIN)
       ref: const
      rows: 2390
      Extra: Using index condition; Using where
id: 1
 select type: SIMPLE
     table: n
      type: eq_ref
possible keys: PRIMARY
       key: PRIMARY
    key len: 16
       ref: osm.n t.node id,osm.n t.version
      rows: 1
      Extra: NULL
mysql-5.6.21 (osm) > SELECT ...
858 rows in set (0.73 sec)
```

```
id: 1
 select type: SIMPLE
      table: w n
       tvpe: ref
possible keys: PRIMARY, node id idx
        key: node id idx
     kev len: 8
        ref: osm.n t.node id
       rows: 1
       Extra: Using index
************** 4. row ************
         id: 1
 select type: SIMPLE
      table: w t
       type: eq_ref
possible keys: PRIMARY
        key: PRIMARY
     key len: 783
        ref: osm.w n.way id,osm.w n.version,const
       rows: 1
       Extra: Using where; Using index
4 rows in set (0.04 sec)
```

#### **Example: optimize this JOIN (XI)**

```
id: 1
        id: 1
                                              select type: SIMPLE
 select type: SIMPLE
                                                   table: w n
     table: n t
                                                    type: ref
      type: ref
                                            possible_keys: PRIMARY,node_id_idx
possible_keys: PRIMARY,k_idx,k_v_i
                                                    key: node id idx
                        An index on (k,v) is
       key: k v idx
                                                 key len: 8
    key len: 1534
                                                    ref: osm.n t.node id
                             even better
       ref: const,const
                                                    rows: 1
      rows: 900
                                                   Extra: Using index
                                            Extra: Using where; Using index
                                                     id: 1
select type: SIMPLE
       id: 1
                                                   table: w t
 select_type: SIMPLE
                                                    type: eq ref
     table: n
                                            possible keys: PRIMARY
      type: eq_ref
                                                    key: PRIMARY
possible keys: PRIMARY
                                                 key len: 783
       key: PRIMARY
                                                    ref: osm.w n.way id,osm.w n.version,const
    key len: 16
                                                    rows: 1
       ref: osm.n_t.node_id,osm.n_t.version
                                                   Extra: Using where; Using index
                                             4 rows in set (0.00 sec)
      rows: 1
      Fxtra: NULL
mysql-5.6.21 (osm) > SELECT ...
858 rows in set (0.02 sec)
```

#### "New" JOIN methods

- MySQL 5.6 added:
  - Batch Key Access
- MariaDB has since 5.3:
  - Batch Key Access
  - Hash Joins
  - Slightly modified versions of the above ones (with "incremental" buffers to join 3 or more tables)

#### Multi-range read

- This optimization orders results obtained from a secondary key in primary key/physical order before accessing the rows
  - It may help exection time of queries when disk-bound
  - It requires tunning of the read\_rnd\_buffer\_size (size of the buffer used for ordering the results)
- BKA JOINs are based on the mrr optimization

#### MRR Example (I)

```
mysql-5.6.21 (osm) > EXPLAIN SELECT * FROM nodes WHERE timestamp >=
'2013-07-01 00:00:00' AND timestamp < '2014-01-01 00:00:00'\G
id: 1
 select type: SIMPLE
      table: nodes
       type: range
possible_keys: nodes_timestamp_idx
        key: nodes timestamp idx
     key len: 5
        ref: NULL
       rows: 429684
      Extra: Using index condition; Using MRR
1 row in set (0.02 sec)
```

## MRR example (II)

```
[restart]
mysql> SET optimizer switch='mrr=off';
mysql> SELECT * FROM nodes WHERE timestamp >= '2013-07-01 00:00:00' AND timestamp <
'2014-01-01 00:00:00';
205617 rows in set (5.16 sec)
mysql> SELECT * FROM nodes WHERE timestamp >= '2013-07-01 00:00:00' AND timestamp <
'2014-01-01 00:00:00';
205617 rows in set (0.60 sec)
[restart]
mysql> SET read rnd buffer size=50 * 1024 * 1024;
mysql> SELECT * FROM nodes WHERE timestamp >= '2013-07-01 00:00:00' AND timestamp <
'2014-01-01 00:00:00';
                                                            "Cold" results are
205617 rows in set (2.39 sec)
mysql> SELECT * FROM nodes WHERE timestamp >= '2013-07
                                                        significantly better with
'2014-01-01 00:00:00';
                                                         mrr (but it can impact
205617 rows in set (0.73 sec)
                                                             negatively, too)
```

## **Batch Key Access**

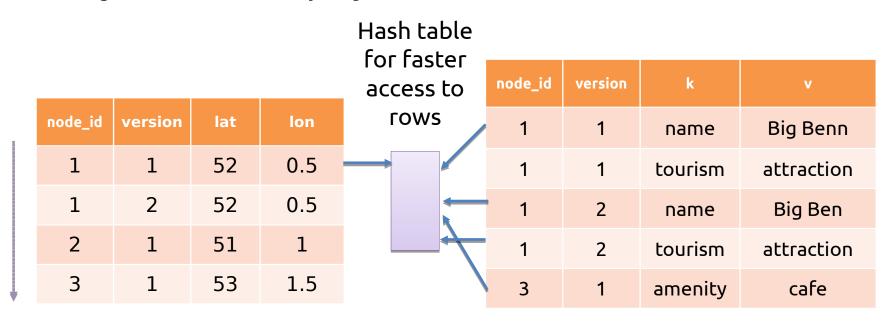


node_id	version	lat	lon	order
1	1	52	0.5	
1	2	52	0.5	
2	1	51	1	
3	1	53	1.5	

node_id	version	k	v	
1	1	name	Big Benn	
1	1	tourism	attraction	
1	2	name	Big Ben	
1	2	tourism	attraction	
3 1		amenity	cafe	

#### **Hash Joins**

Only work for equi-joins



#### MySQL Configuration

• BKA requires changes of default optimizer configuration:

```
mysql-5.7.5 (osm) > SET optimizer_switch= 'mrr=on';
mysql-5.7.5 (osm) > SET optimizer_switch= 'mrr_cost_based=off';
mysql-5.7.5 (osm) > SET optimizer_switch= 'batch_key_access=on';
```

- Additionally, configuring the join\_buffer\_size adequately

#### **MariaDB** configuration

```
mariadb-10.0.14 (osm) > SET optimizer_switch = 'join_cache_incremental=on';
mariadb-10.0.14 (osm) > SET optimizer_switch = 'join_cache_hashed=on';
mariadb-10.0.14 (osm) > SET optimizer_switch = 'join_cache_bka=on';
```

- Enabled by default

```
mariadb-10.0.14 (osm) > SET join_cache_level = 3 (for hash joins)
mariadb-10.0.14 (osm) > SET join_cache_level = 5 (for BKA)
```

- Also, configure join\_buffer\_size appropriately.
- Hash joins, like BKA, are highly dependent on disk-bound DBs to be effective due to the extra overhead

## Nested Loop Join (cold buffers buffer\_pool=100MB, join\_buffer=4M)

```
mariadb-10.0.14 (osm) > EXPLAIN SELECT
changeset id, count(*) FROM changesets JOIN nodes
on changesets.id = nodes.changeset id GROUP BY
visible\G
************** 1. row ************
          id: 1
 select type: SIMPLE
       table: changesets
        type: index
possible keys: PRIMARY
         key: PRIMARY
     kev len: 8
         ref: NULL
        rows: 69115
       Extra: Using index; Using temporary; Using
filesort
```

```
id: 1
 select type: SIMPLE
      table: nodes
       type: ref
possible keys: changeset_id
        key: changeset id
    kev len: 8
        ref: osm.changesets.id
       rows: 19
      Extra:
2 rows in set (0.00 sec)
mariadb-10.0.14 (osm) > SELECT visible, count(*)
FROM changesets JOIN nodes on changesets.id =
nodes.changeset id GROUP BY visible;
1 row in set (32.86 sec)
```

# Hash Join (cold buffers, buffer\_pool=100M, join\_buffer=4M)

```
mariadb-10.0.14 (osm) > EXPLAIN SELECT
                                            id: 1
changeset id, count(*) FROM changesets JOIN
                                              select type: SIMPLE
nodes on changesets.id = nodes.changeset_id
                                                   table: nodes
GROUP BY visible\G
                                                    type: hash ALL
******* possible keys: changeset id
                                                    key: #hash#changeset id
         id: 1
                                                 key len: 8
 select_type: SIMPLE
                                                     ref: osm.changesets.id
       table: changesets
                                                    rows: 2781732
                                                   Extra: Using join buffer (flat, BNLH join)
        type: index
                                            2 rows in set (0.00 sec)
possible keys: PRIMARY
         key: PRIMARY
                                            mariadb-10.0.14 (osm) > SELECT visible, count(*)
     key len: 8
                                            FROM changesets JOIN nodes on changesets.id =
                                            nodes.changeset_id GROUP BY visible;
         ref: NULL
        rows: 69115
                                             Extra: Using index; Using temporary;
Using filesort
                                            1 row in set (6.66 sec)
```

Query Optimization: From 0 to 10 (and up to 5.7)

#### **SUBQUERIES**

# Access types: unique\_subquery/index\_subquery

```
mysql-5.6.21 (osm) > EXPLAIN SELECT ********* 2. row **********
* FROM node tags WHERE v = 'Big Ben'
                                              id: 2
and node id NOT IN (SELECT node id
                                      select_type: DEPENDENT SUBQUERY
FROM nodes WHERE tile < 100000000)\G
                                            table: nodes
******* 1. row *******
                                             type: index subquery
           id: 1
                                    possible keys: PRIMARY, nodes tile idx
 select type: PRIMARY
                                              key: PRIMARY
        table: node tags
                                                                 Unique
                                          key len: 8
         type: ref
                                                                subquery
                                             ref: func
possible keys: v idx
                                                               is similar,
         key: v idx
                                             rows: 1
      key len: 767
                                                               but using a
                                            Extra: Using where
          ref: const
                                    2 rows in set (0.00 sec)
                                                                unique or
         rows: 1
                                                                primary
        Extra: Using where; Using
index
                                                                   kev
```

#### Subqueries in MySQL

- MySQL versions traditionally had very bad press regarding subquries
  - It was common to recommend rewriting them (when possible) into JOINS
- Since MySQL 5.6, its query execution plans have improved significantly

## Lazy Materialization of derived tables

- Option available since MySQL 5.6
  - Improves the execution time of EXPLAIN (it no longer needs to execute subqueries)
  - Derived tables can be indexed automatically at execution time to improve its performance

#### Derived Table Example

mysql-5.5.40 (osm) > EXPLAIN SELECT count(\*) FROM (SELECT \* FROM nodes WHERE VISIBLE = 1) n JOIN changesets ON n.changeset\_id = changesets.id;

id	select_type	table	type	possible_keys		key_len	ref	rows	Extra
1	PRIMARY   PRIMARY   DERIVED	<derived2>   changesets   nodes  </derived2>	ALL   eq_ref   ALL	NULL PRIMARY NULL	NULL PRIMARY NULL	NULL   8   NULL	NULL n.changeset_id   NULL	2865312 1 2865521	Using index     Using where

3 rows in set (1.42 sec)

mysql-5.6.21 (osm) > EXPLAIN SELECT count(\*) FROM (SELECT \* FROM nodes WHERE VISIBLE = 1) n JOIN changesets ON n.changeset\_id = changesets.id;

id	select_type	table	type	possible_keys	key	key_len	ref	rows	Extra
1     1     2	PRIMARY PRIMARY DERIVED	changesets <derived2> nodes</derived2>	index   ref   ALL	PRIMARY   <auto_key0>   NULL</auto_key0>	PRIMARY   <auto_key0>   NULL</auto_key0>	8 8 NULL	   NULL   osm.changesets.id   NULL	70917 40 2853846	Using index   NULL   Using where

3 rows in set (0.00 sec)

Subquery is not executed

Autogenerated index

## A Common 5.5 Performance Problem

```
mysql-5.5.40 (osm) > EXPLAIN SELECT * FROM nodes
                       WHERE nodes.changeset id IN (
                         SELECT changesets.id
                          FROM changesets
                          JOIN users
                            ON changesets.user id = users.id and users.display name = 'Steve');
  1 | PRIMARY
                    nodes
                                                                               NULL
                             | cons | PRIMARY, users display name idx | users display name idx | 767 | const |
  2 | DEPENDENT SUBQUERY | users
                                                                                                   1 | Using index
  2 | DEPENDENT SUBQUERY | changesets | eq ref |
                                                                                                   1 | Usina where
3 rows in set (0.00 sec)
mysql-5.5.40 (osm) > SELECT ...;
                                                             This means that
                                                             the subquery is
 node id | latitude | longitude | changeset id | visit
                                                             executed almost
     99890 | 515276425 | -1497621 |
                                           552
    109174 | 515364532 | -1457329 |
                                          1875 |
                                                              3 million times
    276538 | 515324296 |
                        -2094688
                                          810
    442987 | 515449207 |
                        -1275650
                                          1941 |
                                                          2006-01
    442988 | 515449741 |
                        -1272860
                                          1941 |
    498803 | 515438432 |
                                                          2006-02-03 21:55:1/ | 2062268628
                        -1269436 l
                                          2171 l
  138212838 | 513010180 |
                         -1699929 l
                                        7757299
                                                          2011-04-03 18:14:14 | 2062220563
7 rows in set (2.60 sec)
```

#### Semijoin Optimization

- The only way to execute certain IN subqueries was to execute them with poor strategy
  - This forced to rewrite certain queries into JOINS or scalar subqueries, when possible
- There are now several additional automatic options:
  - Convert to a JOIN
  - Materialization (including index creation)
  - FirstMatch
  - LooseScan
  - Duplicate Weedout

### The Previous Query is Not a Problem in 5.6/5.7/MariaDB 5.3+

```
mysql-5.6.21 (osm) > EXPLAIN SELECT * FROM nodes
                        WHERE nodes.changeset id IN (
                         SELECT changesets.id
                           FROM change sets
                           JOIN users
                            ON changesets.user_id = users.id and users.display_name = 'Steve');
 id | select type | table | type | possible_keys | key | key_len | ref
                     | const | PRIMARY,users display name idx | users display name idx | 767 | const
                                                                                    | 1 | Using index |
                                                | NULL | NULL | NULL | 70917 | Using where | | changeset_id | 8 | osm.changesets id | 21 | NULL |
 1 | SIMPLE | changesets | ALL | PRIMARY
                     | ref | changeset id
 1 | SIMPLE
             nodes
3 rows in set (0.00 sec)
mysql-5.6.21 (osm) > SELECT ...;
                                                                                      Executed as
+------
 node id | latitude | longitude | changeset id | visible | timestamp | tile | version |
                                                                                        a regular
    99890 | 515276425 | -1497621 | 552 | 1 | 2005-10-25 00:35:24 | 2062268512 |
                                                                                           JOIN
7 rows in set (0.02 \text{ sec})
```

#### First Match Strategy

```
mysql-5.7.5 (osm) > EXPLAIN SELECT * FROM changesets WHERE id IN (SELECT changeset id FROM nodes)\G
id: 1
 select type: SIMPLE
      table: changesets
  partitions: NULL
      type: ALL
possible_keys: PRIMARY
       key: NULL
    key len: NULL
       ref: NULL
       rows: 70917
   filtered: 100.00
      Extra: NULL
                                                            It is converting
id: 1
                                                             the ref into an
 select type: SIMPLE
      table: nodes
                                                              eq_ref, shot-
  partitions: NULL
      type: ref
                                                              circuiting the
possible_keys: changeset_id
       key: changeset_id
                                                                execution
    key len: 8
       ref: osm.changesets.id
       rows: 33
   filtered: 100.00
      Extra: Using index; FirstMatch(changesets)
2 rows in set, 1 warning (0.00 sec)
```

# Enabling and disabling materialization, semijoin, etc

Query Optimization: From 0 to 10 (and up to 5.7)

#### **QUERY PROFILING**

### Which Queries Should I Optimize First?

- My two favorite methods:
  - pt-query-digest
  - PERFORMANCE SCHEMA
- I prefer pt-query-digest for long-term reports, PERFORMANCE\_SCHEMA for more real-time evaluation and fine-tuning
  - Also, PERFORMANCE\_SCHEMA was not "ready" until MySQL 5.6

#### pt-query-digest

- It is a 3rd party tool written in Perl, originally created by Baron Schwartz
- It requires activation of the slow log:

```
- SET GLOBAL slow_query_log = 1;
```

- SET long\_query\_time = 0;

Be careful with extra IO and latency!

- In Percona Server and MariaDB it can provide extra information:
  - SHOW GLOBAL VARIABLES like 'log\_slow\_verbosity';

#### pt-query-digest Execution (I)

```
# 1094.7s user time, 9.4s system time, 141.22M rss, 205.21M vsz
# Current date: Wed Jul 1 07:32:28 2015
# Hostname: db1018
# Files: STDIN
# Overall: 4.66M total, 640 unique, 53.47 QPS, 0.02x concurrency ____
# Time range: 2015-06-30 07:00:10 to 2015-07-01 07:11:37
# Attribute
                  total
                           min
                                        avg
                                               95% stddev median
                                 max
# =========
# Exec time
                1320s 1us 3s 283us 332us 3ms 152us
                            0 13ms 51us 93us 39us 52us
# Lock time
                 238s
                5.02M 0 4.16k 1.13 1.96 8.69 0.99
# Rows sent
             9.50M
                            0 232.93k 2.14 3.89 261.15 0.99
# Rows examine
# Merge passes
                            0
# Query size
                 1.06G
                           17 67.89k
                                      243.89 511.45
                                                    368.99 192.76
# Boolean:
# Filesort 8% yes, 91% no
                                                    Actual execution
# Full scan
             94% yes, 5% no
# Priority que 3% yes, 96% no
```

on Wikipedia production servers

# Tmp table on

# Tmp table 29% yes, 70% no

1% yes, 98% no

#### pt-query-digest Execution (II)

# Profile									
# Ra	nk	Query ID	Response	time	Calls	R/Call	V/M	Item	
# ==	==	=======================================	=======	=====	======	=====	=====	======	=====
#	1	0xSANITIZED	242.2765	18.4%	691005	0.0004	0.00	SELECT	revision page user
#	2	0×SANITIZED	204.7052	15.5%	80863	0.0025	0.01	SELECT	revision page user
#	3	0xSANITIZED	162.8476	12.3%	1025179	0.0002	0.00	<b>SELECT</b>	page
#	4	0xSANITIZED	68.1164	5.2%	93928	0.0007	0.01	<b>SELECT</b>	revision page user
#	5	0xSANITIZED	66.8302	5.1%	354562	0.0002	0.00	<b>SELECT</b>	page revision
#	6	0xSANITIZED	57.0374	4.3%	211631	0.0003	0.00	<b>SELECT</b>	page revision
#	7	0xSANITIZED	44.0751	3.3%	6925	0.0064	0.07	SELECT	page categorylinks
category									
#	8	0xSANITIZED	35.0655	2.7%	9689	0.0036	0.00	<b>SELECT</b>	text
#	9	0xSANITIZED	29.4363	2.2%	152259	0.0002	0.00	<b>SELECT</b>	page
#	10	0xSANITIZED	24.1864	1.8%	176927	0.0001	0.00	<b>SELECT</b>	msg_resource
#	11	0xSANITIZED	23.7016	1.8%	144807	0.0002	0.00	<b>SELECT</b>	page_restrictions
#	12	0xSANITIZED	16.6547	1.3%	10135	0.0016	0.03	<b>SELECT</b>	revision
#	13	0xSANITIZED	15.0564	1.1%	263809	0.0001	0.00	SET	

#### pt-query-digest Execution (III)

```
# Ouerv 1: 7.93 OPS. 0.00x concurrency. ID 0xSANITIZED at byte 1553864032
# This item is included in the report because it matches -limit.
# Scores: V/M = 0.00
# Time range: 2015-06-30 07:00:10 to 2015-07-01 07:11:37
                                                 95% stddev median
# Attribute
             pct total
# Count
              14 691005
# Exec time
                   242s
                                        350us
                         163us
                                  91ms
                                               348us
                                                       563us
                                                              301us
# Lock time
                    63s
                          47us
                                         91us
                                               103us
                                                       14us
                                                              84us
# Rows sent
              12 657.18k
                                         0.97
                                                0.99
                                                       0.16
                                                              0.99
# Rows examine 6 657.18k
                                         0.97
                                                0.99
                                                       0.16
                                                              0.99
                                                       9.22 511.45
# Query size
            31 345.42M
                                   749 524.16 537.02
# String:
# Databases
             itwiki (225976/32%), enwiktiona... (219461/31%)... 15 more
# Hosts
# Users
             wikiuser
# Query time distribution
  1us
  10us
        1ms
  10ms #
# 100ms
    1s
# 10s+
# Tables
    SHOW TABLE STATUS FROM `enwiktionary` LIKE 'revision'\G
    SHOW CREATE TABLE `enwiktionary`.`revision`\G
    SHOW TABLE STATUS FROM `enwiktionary` LIKE 'page'\G
    SHOW CREATE TABLE 'enwiktionary'.'page'\G
    SHOW TABLE STATUS FROM `enwiktionary` LIKE 'user'\G
    SHOW CREATE TABLE 'enwiktionary'.'user'\G
# EXPLAIN /*!50100 PARTITIONS*/
SELECT /* Revision::fetchFromConds SANITIZED */ * FROM `revision`
INNER JOIN 'page' ON ((page id = rev page)) LEFT JOIN 'user' ON
((rev_user != 0) AND (user_id = rev_user)) WHERE page_namespace = '0' AND
page title = 'SANITIZED' AND (rev id=page latest) LIMIT 1\G
```

#### PERFORMANCE\_SCHEMA

- Monitoring schema (engine) enabled by default since MySQL 5.6
  - performance\_schema = 1 (it is not dynamic)
- Deprecates the old query profiling
- It is way more user-friendly when combined with the <u>SYS schema/ps\_helper</u> (a set of views and stored procedures created by Mark Leith)
  - Included by default since 5.7.7

# Installation of the SYS Schema for 5.6/MariaDB

```
$ git clone https://github.com/MarkLeith/mysql-sys.git
Cloning into 'mysql-sys'...
remote: Counting objects: 926, done.
remote: Compressing objects: 100% (73/73), done.
remote: Total 926 (delta 35), reused 6 (delta 2)
Receiving objects: 100% (926/926), 452.19 KiB | 225.00
KiB/s, done.
Resolving deltas: 100% (584/584), done.
$ cd mysal-sys/
$ ~/sandboxes/msb 5 6 24/use < sys 56.sql</pre>
```

### Example Usage: Discovering Unused Indexes

```
mysql-5.7.8 (osm) > SELECT * FROM sys.schema unused indexes LIMIT 5;
 object schema | object name
                                   index name
                                   acls k idx
                acls
 osm
                                   changeset_tags_id_idx
                changeset_tags
 osm
                                 | current_nodes_timestamp_idx
                current_nodes
 osm
                current_nodes
                                 | current_nodes_tile_idx
 osm
                current relations | current relations timestamp idx
 osm
5 rows in set (0.04 sec)
mysql-5.7.8 (osm) > SELECT * FROM current nodes WHERE tile = 100;
mysql-5.7.8 (osm) > SELECT * FROM sys.schema unused indexes LIMIT 5;
 object schema | object name
                                   index name
                acls
                                   acls_k_idx
 osm
                changeset_tags
                                   changeset tags id idx
 osm
                current nodes
                                   current nodes timestamp idx
 osm
```

current relations timestamp idx

With enough activity, it can help us clean up our schema

5 rows in set (0.03 sec)

osm

osm

current relations |

| current relations | changeset id

# Example Usage: Slow Queries (ordered by server time)

```
mysql-5.7.8 (osm) > SELECT * FROM sys.statement_analysis LIMIT 10\G
********* 1. TOW ********
           query: SELECT `way id` AS
                                                rows examined: 20152155
`id` , `v` FROM `way tags` WHERE `v` LIKE ?
                                            rows_examined_avg: 1343477
              db: osm
                                                rows_affected: 0
       full scan: *
                                            rows affected avg: 0
      exec_count: 15
                                                   tmp tables: 0
       err_count: 0
                                              tmp_disk_tables: 0
      warn count: 0
                                                  rows_sorted: 0
   total_latency: 7.83 s
                                            sort_merge_passes: 0
     max_latency: 1.33 s
                                                       digest:
     avg_latency: 521.84 ms
                                            21f90695b1ebf20a5f4d4c1e5e860f58
    lock_latency: 17.94 ms
                                                   first_seen: 2014-11-01 17:04:51
       rows sent: 6779
                                                    last seen: 2014-11-01 17:05:22
    rows sent avg: 452
```

### **Example Usage: Top Queries Creating Temporary Tables**

```
mysql-5.7.8 (osm) > SELECT * FROM
sys.statements_with_temp_tables WHERE db = 'osm' LIMIT 10\G
query: SELECT ? AS TYPE , `node_id` A ... gs` WHERE `k` = ? AND `v` = ?
                 db: osm
           exec count: 11
        total latency: 7.57 s
     memory tmp tables: 11
       disk tmp tables: 0
avg tmp tables per query: 1
 tmp_tables_to_disk_pct: 0
           first seen: 2014-11-01 17:33:55
            last seen: 2014-11-01 17:34:45
              digest: 5e6e82799b7c7c0e5c57cfe63eb98d5d
```

### Example Usage: Top Queries Creating Temporary Tables (cont.)

mysql-5.7.8 (osm) > SELECT DIGEST\_TEXT FROM performance\_schema.events\_statements\_summary\_by\_digest
WHERE digest = '5e6e82799b7c7c0e5c57cfe63eb98d5d'\G

DIGEST\_TEXT: SELECT ? AS TYPE , `node\_id` AS `id` FROM `node\_tags` WHERE `k` = ? AND `v` = ? UNION SELECT ? AS TYPE , `way\_id` AS `id` FROM `way\_tags` WHERE `k` = ? AND `v` = ? UNION SELECT ? AS TYPE , `relation\_id` AS `id` FROM `relation\_tags` WHERE `k` = ? AND `v` = ? 1 row in set (0.00 sec)

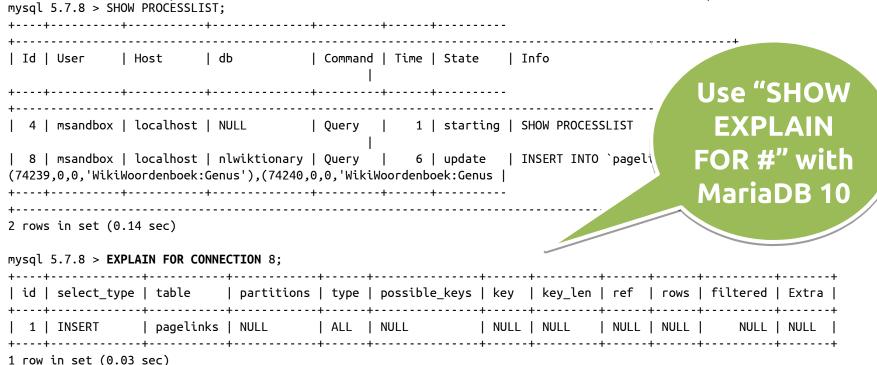
mysql-5.7.8 (osm) > EXPLAIN SELECT 'node' as type, node\_id as id FROM node\_tags WHERE k='amenity'
and v='cafe' UNION SELECT 'way' as type, way\_id as id FROM way\_tags WHERE k='amenity' and v='cafe'
UNION SELECT 'relation' as type, relation\_id as id FROM relation\_tags WHERE k='amenity' and
v='cafe';

<sup>4</sup> rows in set, 1 warning (0.01 sec)

Query Optimization: From 0 to 10 (and up to 5.7)

## GENERAL OPTIMIZER IMPROVEMENTS

# EXPLAIN FOR CONNECTION (SHOW EXPLAIN FOR)



#### MySQL 5.7 Optimizer Tweaks

```
mysql-5.5.40 (osm) > SELECT *
                       FROM nodes
                       JOIN node_tags
                         ON node tags.node id = nodes.node id
                      WHERE nodes.latitude
                            BETWEEN 517000000 and 520000000\G
59 rows in set (1.37 sec)
mysql-5.5.40 (osm) > SELECT STRAIGHT_JOIN *
                                                     This condition is
                       FROM nodes
                                                      very selective,
                       JOIN node_tags
                         ON node tags.node id =
                                                      but there is no
                      WHERE nodes.latitude
                                                      index available
                            BETWEEN 517000000 and
59 rows in set (0.86 sec)
```

## Why does it take the wrong table order?

```
mysql-5.5.40 (osm) > EXPLAIN EXTENDED SELECT * FROM nodes JOIN node_tags ON node_tags.node_id = nodes.node id WHERE
nodes.latitude BETWEEN 517000000 and 520000000\G
id: 1
 select_type: SIMPLE
      table: node_tags
       type: ALL
possible keys: PRIMARY
        key: NULL
    key len: NULL
        ref: NULL
       rows: 839031
    filtered: 100.00
      Extra:
id: 1
 select type: SIMPLE
      table: nodes
       type: ref
possible keys: PRIMARY
        key: PRIMARY
    key len: 8
        ref: osm.node tags.node id
       rows: 1
    filtered: 100.00
      Extra: Using where
2 rows in set, 1 warning (0.00 sec)
```

# Unindexed Columns Are not Accounted

mysql-5.5.40 (osm) > EXPLAIN EXTENDED SELECT STRAIGHT\_JOIN \* FROM nodes JOIN node\_tags ON node tags.node id = nodes.node id WHERE nodes.latitude BETWEEN 517000000 and 520000000\G \*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\* 1. row \*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\* id: 1 select type: SIMPLE The optimizer table: nodes type: ALL assumes that all possible keys: **PRIMARY** key: NULL rows will be key len: NULL ref: NULL rows: 2865521 returned filtered: 100.00 Extra: Using where id: 1 select type: SIMPLE table: node tags type: ref possible keys: PRIMARY key: PRIMARY key len: 8 ref: osm.nodes.node id rows: 1 filtered: 100.00 Extra:

2 rows in set, 1 warning (0.00 sec)

#### What's new in 5.7?

mysql-5.7.5 (osm) > EXPLAIN SELECT \* FROM nodes JOIN node tags ON node tags.node id = nodes.node id WHERE nodes.latitude BETWEEN 517000000 and 520000000\G id: 1 select\_type: SIMPLE It is actually table: nodes partitions: NULL 0.002%, but type: ALL possible keys: PRIMARY it is good key: NULL key len: NULL ref: NULL enough rows: 2773853 filtered: 11.11 Extra: Using where \* 2. FOW \* id: 1 select type: SIMPLE table: node\_tags partitions: NULL type: ref possible\_keys: PRIMARY key: PRIMARY key len: 8 ref: osm.nodes.node\_id **rows:** 3 filtered: 100.00 Extra: NULL

2 rows in set, 1 warning (0.00 sec)

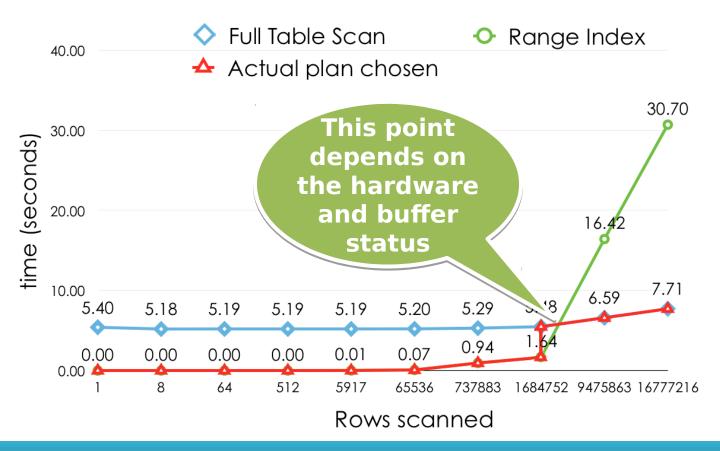
# If things go wrong -our Friend optimizer\_switch

```
mysql-5.7.8 (osm) > SET optimizer_switch='condition_fanout_filter=off';
Query OK, 0 rows affected (0.00 sec)
```

## Improved EXPLAIN FORMAT=JSON

```
"nested_loop": [
                                                  Cost
       "table": {
                                             information
         "table name": "nodes",
                                                is now
         "access_type": "ALL",
         "possible_keys": [
                                               included
           "PRIMARY"
         "rows_examined_per_scan": 2773853,
         "rows produced_per_join": 308175,
                                                        Engine
         "filtered": 11.11.
                                                    statistics are
         "cost info": {
                                                   now floats for
           "read cost": "512783.58",
           "eval cost": "61635.01",
                                                      improved
           "prefix_cost": "574418.60",
                                                      precision
           "data_read_per_join": "21M"
         },
```

#### **Optimizer Cost Tuning**



#### **Configurable Costs**

```
mysql 5.7.8> SELECT * FROM mysql.server cost;
 cost_name | cost_value | last update
                                                  comment
 disk_temptable_create_cost |
                             NULL | 2015-09-20 20:48:10 | NULL
                                                                The unit is
 disk_temptable_row_cost |
                             NULL | 2015-09-20 20:48:10 |
                                                   NULL
                                                                "read of a
 key compare cost
                             NULL | 2015-09-20 20:48:10 | NULL
 memory temptable create cost |
                             NULL | 2015-09-20 20:48:10 |
                                                   NULL
                                                                  random
 memory_temptable_row_cost |
                             NULL | 2015-09-20 20:48:10 | NULL
                                                                data page"
 row evaluate cost
                             NULL | 2015-09-20 20:48:10 | NULL
6 rows in set (0.00 sec)
mysql 5.7.8> SELECT * FROM mysql.engine cost;
 l comment
 default | 0 | io_block_read_cost | NULL | 2015-09-20 20:48:10 | NULL
 default |
                   0 | memory_block_read_cost | NULL | 2015-09-20 20:48:10 | NULL
2 rows in set (0.00 sec)
```

#### Changing Costs Example

```
mysql-5.7.5 (osm) > EXPLAIN FORMAT=JSON SELECT 'node' as type, node_id as id FROM
node tags WHERE k='amenity' and v='cafe' UNION SELECT 'way' as type, way id as id FROM
way tags WHERE k='amenity' and v='cafe' UNION SELECT 'relation' as type, relation_id as
id FROM relation tags WHERE k='amenity' and v='cafe'\setminus G
            "cost info": {
              "query cost": "22567.80"
1 row in set, 1 warning (0.00 sec)
mysql-5.7.5 (osm) > UPDATE mysql.server_cost SET cost value = 10;
mysql-5.7.5 (osm) > FLUSH OPTIMIZER_COSTS;
<session restart>
mysql-5.7.5 (osm) > EXPLAIN ...\G
            "cost info": {
              "query cost": "661371.00"
1 row in set, 1 warning (0.00 sec)
```

More info on usage in the manual: <a href="http://dev.mysql.com/doc/refman/5.7/en/cost-model.html">http://dev.mysql.com/doc/refman/5.7/en/cost-model.html</a>

## Still waiting for...

- Utility to analyze the underlying technology (HD, SSD, memory) and filling up the tables automatically
- Buffer "hotness"-aware statistics
- Statistics for non-indexed columns/histograms
  - MariaDB has histograms since 10

### Engine-Independent Statistics/Histograms in MariaDB 10

```
mariadb-10.0.14 (osm) > SET histogram size = 255;
Query OK, 0 rows affected (0.00 sec)
mariadb-10.0.14 (osm) > SET use_stat_tables = 2;
Query OK, 0 rows affected (0.01 sec)
mariadb-10.0.14 (osm) > ANALYZE TABLE node_tags;
| Table | Op | Msg type | Msg text
+-----
| osm.node_tags | analyze | status | Engine-independent statistics collected |
osm.node tags | analyze | status | OK
2 rows in set (3.01 sec)
mariadb-10.0.14 (osm) > ANALYZE TABLE nodes:
| Table | Op | Msg_type | Msg_text
  -----
osm.nodes | analyze | status | Engine-independent statistics collected |
osm.nodes | analyze | status | OK
+-----
2 rows in set (32.52 sec)
mariadb-10.0.14 (osm) > SET optimizer use condition selectivity = 3; (or 4 for histograms)
Query OK, 0 rows affected (0.00 sec)
```

#### **Better Stats!**

```
mariadb-10.0.14 (osm) > EXPLAIN EXTENDED SELECT * FROM nodes JOIN node tags ON node tags.node id = nodes.node id WHERE
nodes.latitude BETWEEN 517000000 and 520000000\G
id: 1
 select type: SIMPLE
      table: nodes
       type: ALL
possible_keys: PRIMARY
                                       Much better
       key: NULL
    key_len: NULL
       ref: NULL
                                        estimation
       rows: 2865312
   filtered: 0.39
      Extra: Using where
id: 1
 select_type: SIMPLE
      table: node tags
      type: ref
possible_keys: PRIMARY
       key: PRIMARY
    key len: 8
       ref: osm.nodes.node_id
       rows: 3
   filtered: 100.00
      Extra:
2 rows in set, 1 warning (0.01 sec)
```

## Other Changes/Bugs Fixed

- UNION ALL does not create temporary tables, returns tables faster
- (a, b) IN ((1, 2), (2, 3)) can use index ranges
- EXPLAIN EXTENDED is now the default behavior
- I.C. Pushdown support for partitioned tables
- IGNORE clause meaning has been standardized between sentence types
- Increased default for optimizer\_search\_depth

Query Optimization: From 0 to 10 (and up to 5.7)

## **COMPUTED/VIRTUAL COLUMNS**

#### **Syntax**

```
ALTER TABLE nodes

ADD COLUMN lon DECIMAL (10, 7)

as (longitude/10000000) VIRTUAL,

ADD COLUMN lat DECIMAL (9, 7)

as (latitude/10000000) VIRTUAL;
```

- They can be used to simplify SELECTS, calculating values on the fly
- Non accessed rows are not calculated

#### **Functional Indexes**

• Before 5.7.8:

```
mysql-5.7.5 (osm) > ALTER TABLE nodes add index(lon);
ERROR 1951 (HY000): Key/Index cannot be defined on a non-
stored virtual column.
```

Now:

```
mysql-5.7.8 (osm) > ALTER TABLE nodes add index(lon);
Query OK, 0 rows affected (16.54 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

 This effectively is an implementation of functional indexes, allowing to solve previous query optimization issues without the overhead of an extra column

### Do you remember this query?

```
MariaDB [nlwiktionary]> EXPLAIN SELECT * FROM revision WHERE
substr(rev timestamp, 5, 2) = '09'\G
id: 1
 select type: SIMPLE
      table: revision
       type: ALL
possible keys: NULL
                        Can you think a way to
        key: NULL
                         improve this query?
     key len: NULL
        ref: NULL
       Rows: 173154
      Extra: Using where
1 row in set (0.00 sec)
```

#### Now we can solve it like this

```
mysql-5.7.8> ALTER TABLE revision ADD COLUMN rev_month tinyint AS
(substr(rev timestamp, 5, 2)) VIRTUAL, ADD INDEX rev month (rev month);
Query OK, 820308 rows affected (3 min 29.48 sec)
Records: 820308 Duplicates: 0 Warnings: 0
mysql-5.7.8> EXPLAIN SELECT * FROM revision WHERE rev_month = 9\G
id: 1
 select type: SIMPLE
       table: revision
                                 The column does not
  partitions: NULL
       type: ref
                                 take space, only the
possible keys: rev month
        key: rev_month
                                         index
     key len: 2
        ref: const
       rows: 104112
    filtered: 100.00
       Extra: NULL
1 row in set, 1 warning (0.01 sec)
```

#### **Stored Columns**

```
mysql-5.7.8 (osm) > ALTER TABLE nodes CHANGE lat lat DECIMAL (9, 7) as
(latitude/10000000) STORED;
ERROR 1954 (HY000): 'Changing the STORED status' is not supported for virtual columns.
They dropped the 'yet':-)
```

mysql-5.7.8 (osm) > ALTER TABLE nodes DROP COLUMN lat, ADD COLUMN lat DECIMAL
(9, 7) as (latitude/10000000) STORED;
Query OK, 2865312 rows affected (4 min 51.05 sec)
Records: 2865312 Duplicates: 0 Warnings: 0

MariaDB uses
the
'PERMANENT'
keyword

#### **Features and Limitations**

- Virtual "non-stored" columns cannot be indexed
- "Stored columns" can be PRIMARY, UNIQUE, FOREIGN and MULTI keys
- It cannot contain subqueries or other tables or rows references
- It cannot contain used-defined stored functions
- It can contain server and user-defined variables
- It can be computed based on other virtual columns

Query Optimization: From 0 to 10 (and up to 5.7)

#### **QUERY REWRITE PLUGINS**

# New APIs for Query Rewriting

- One for pre-parsing rewriting
- Another for post-parsing rewriting

### **Example Plugin Installation**

```
$ ~/sandboxes/msb_5_7_8/my sql <
./share/install_rewriter.sql</pre>
```

More on this: <a href="http://mysqlserverteam.com/the-query-rewrite-plugins/">http://mysqlserverteam.com/the-query-rewrite-plugins/</a>

## **Query Rewriting Setup**

```
mysql-5.7.8 > INSERT INTO query rewrite.rewrite rules (pattern, replacement) VALUES ('SELECT ?', 'SELECT ?
+ 1');
Query OK, 1 row affected (0.01 sec)
mysql-5.7.8 > CALL query rewrite.flush rewrite rules();
Query OK, 0 rows affected (0.01 sec)
mysql-5.7.8 > SELECT 1;
+----+
| 1 + 1 |
1 2 1
+----+
1 row in set, 1 warning (0.00 sec)
mysql-5.7.8 > SHOW WARNINGS;
+-----+
| Level | Code | Message
| Note | 1105 | Query 'SELECT 1' rewritten to 'SELECT 1 + 1' by a query rewrite plugin |
+-----+
1 row in set (0.00 sec)
```

#### Considerations

- It cannot correct malformed queries- pattern and replacement must be syntactically correct
- Useful for query optimizations for 3rd party applications
- User is notified that a rewrite happened with a note-level message
- Low overhead (5%); specially for untouched queries
- You can do stupid things like:
   mysql-5.7.5 ((none)) > INSERT INTO
   query\_rewrite.rewrite\_rules( pattern, replacement ) VALUES
   ( 'SELECT 1', 'DROP TABLE test.test' );
   Query OK, 1 row affected (0.01 sec)

Query Optimization: From 0 to 10 (and up to 5.7)

#### **OPTIMIZER HINTS**

## New functionality since 5.7.7

MySQL accepts hints for query execution with the syntax /\*+... \*/:

```
mysql-5.7.8> EXPLAIN SELECT /*+ NO_ICP(revision)*/ * FROM revision WHERE rev_comment like
'jaime%' AND rev_timestamp > '2008'\G
id: 1
 select type: SIMPLE
       table: revision
  partitions: NULL
       type: range
possible keys:
rev_timestamp,rev_timestamp_rev_page,rev_timestamp_2,rev_comment_rev_timestamp
         key: rev comment rev timestamp
     key len: 783
        ref: NULL
        rows: 1
    filtered: 50.00
       Extra: Using where
1 row in set, 1 warning (0.00 sec)
```

More info on this feature:

https://www.percona.com/blog/2015/04/30/optimizer-hints-mysql-5-7-7-missed-manual/

#### **Syntax**

- The syntax is identical to the one Oracle
   Database uses, but does not deprecate yet the
   old hint syntax (USE INDEX, STRAIGHT\_JOIN, ...)
  - "planned"
- It has some overlap with optimizer\_switch
  - Although it has lower granularity (table instead of statement)

#### Max Query Execution

Indicate it in miliseconds:

```
mysql-5.7.8 > SELECT /*+ MAX_EXECUTION_TIME(1784)*/ SLEEP(10)\G
**************************
SLEEP(2000): 1
1 row in set (1.78 sec)
```

 max\_statement\_time syntax and variable seems to have been removed since 5.7.8

#### Other hints

- BKA / NO\_BKA
- BNL/NO\_BNL
- MRR / NO\_MRR
- NO\_ICP
- NO\_RANGE\_OPTIMIZATION
- QB\_NAME (controlls the query block where to apply the hint to)

Query Optimization: From 0 to 10 (and up to 5.7)

#### **SQL MODE CHANGES**

#### Default SQL Mode Changes

- MySQL 5.5 and earlier
  - \_\_ , ,
- MySQL 5.6 (from 5.6.6):
  - 'NO\_ENGINE\_SUBSTITUTION' is the default
  - NO\_ENGINE\_SUBSTITUTION and STRICT\_TRANS\_TABLES were suggested in upstream default my.cnf
- MySQL 5.7 (from 5.7.5):
  - 'ONLY\_FULL\_GROUP\_BY,STRICT\_TRANS\_TABLES,NO\_ENGINE\_SUBST ITUTION' is the new default

More on this: <a href="http://www.tocker.ca/2014/01/14/making-strict-sql">http://www.tocker.ca/2014/01/14/making-strict-sql</a> mode-the-default.html

#### Stricter Defaults

```
mysql> CREATE TABLE `test` (
  `id` int(11) NOT NULL,
  `type` enum('movie','album','videogame') NOT NULL,
  'date' datetime NOT NULL,
  PRIMARY KEY ('id')
 ENGINE=InnoDB:
mysql> INSERT INTO test (type, date) VALUES ('tv show', -1);
```



## MySQL 5.5/5.6 with default settings

```
mysql> INSERT INTO test (type, date) VALUES ('tv show', -1);
Query OK, 1 row affected, 3 warnings (0.00 sec)
mysql> SHOW WARNINGS;
  | Level | Code | Message
| Warning | 1265 | Data truncated for column 'type' at row 1
| Warning | 1264 | Out of range value for column 'date' at row 1
+----+
3 rows in set (0.00 sec)
mysql> SELECT * FROM test;
+---+
| id | type | date
+---+
 0 | 0000-00-00 00:00:00 |
1 row in set (0.00 sec)
```

#### MySQL 5.7

```
mysql> INSERT INTO test (type, date) VALUES ('tv show', -1);
ERROR 1265 (01000): Data truncated for column 'type' at row 1
mysql> INSERT INTO test (type, date) VALUES ('videogame', -1);
ERROR 1292 (22007): Incorrect datetime value: '-1' for column 'date' at row
mysql> INSERT INTO test (type, date) VALUES ('movie', now());
ERROR 1364 (HY000): Field 'id' doesn't have a default value
mysql> INSERT INTO test (id, type, date) VALUES (1, 'videogame', now());
Query OK, 1 row affected (0.01 sec)
```

### GROUP BY Behavior Changes

+	+	+
way_id	count(*)	node_id
+	++	+
155339744	1187	1558095871
243986064	1156	2604713337
87136668	1128	1013304944
148812873	852	1618837453
149200774	835	34921158
183618216	826	1940223096
273858696	824	1267549776
261584374	770	2669122104
227880171	704	2240011804
193564006	684	1808872763
+	++	+

Nondeterministic behavior

10 rows in set (1.24 sec)

## With ONLY\_FULL\_GROUP\_BY (default in 5.7)

```
mysql-5.7.5 (osm) > SELECT way_id, count(*), node_id
                    FROM way nodes
                    GROUP BY way id
                    ORDER BY count(*) DESC
                    LIMIT 10;
ERROR 1055 (42000): Expression #3 of SELECT list is not in GROUP
BY clause and contains nonaggregated column
'osm.way_nodes.node_id' which is not functionally dependent on
columns in GROUP BY clause; this is incompatible with
sql_mode=only_full_group_by
```

# Problem with ONLY\_FULL\_GROUP\_BY in MySQL <= 5.6

```
mysql-5.6.21 (osm) > SET SQL_mode='ONLY_FULL_GROUP_BY';
Query OK, 0 rows affected (0.00 sec)
mysql-5.6.21 (osm) > SELECT
                        u.id as `user id`,
                        u.display name as `user name`,
                        wunt(*) as `# changesets`
     Functional KOM users u
    dependenc
                      IN changesets c
                        ON u.id = c.user id
                     GROUP BY u.id
                    ORDER BY count(*) DESC
                     LIMIT 10;
ERROR 1055 (42000): 'osm.u.display name' isn't in GROUP BY
```

More on this: <a href="http://rpbouman.blogspot.com/2014/09/mysql-575-group-by-respects-functional.html">http://rpbouman.blogspot.com/2014/09/mysql-575-group-by-respects-functional.html</a>

### 5.7 Aims for SQL99 Compliance

```
mysql-5.7.5 (osm) > SELECT u.id as `user id`, u.display name as `user name`,
count(*) as `# changesets` FROM users u JOIN changesets c ON u.id =
c.user id GROUP BY u.id ORDER BY count(*) DESC LIMIT 10;
 user id | user name
                          | # changesets
    31257
            Ed Avis
                                    4333
            Welshie
                                    2696
      508
           Amaroussi
  1016290
                                    2351
   352985 | ecatmur
                                    1450
   142807 |
           SDavies
                                    1342
            Steve Chilton
                                    1182
      736 l
      346 | Tom Chance
                                    1175
    38784 | Tom Morris
                                    1165
            UrbanRambler
    88718 l
                                    1151
            Harry Wood
     1611
                                    1091
10 rows in set (0.09 sec)
```

## **Backward Compatibility**

- Some ORMs and frameworks change the default SQL Mode:
  - Ruby on Rails 4+ sets STRICT\_ALL\_TABLES
  - Drupal 7+ sets TRADITIONAL
  - Mediawiki will set TRADITIONAL, ONLY\_FULL\_GROUP\_BY
- Other applications do not work in standardcompliance modes:
  - Wordpress used to not work in strict mode (fixed):
     <a href="http://www.xaprb.com/blog/2013/03/15/wordpress-and-mysqls-strict-mode/">http://www.xaprb.com/blog/2013/03/15/wordpress-and-mysqls-strict-mode/</a>
  - Cacti does not work with ONLY\_FULL\_GROUP\_BY,NO\_ZERO\_DATE

#### **Deprecated Modes**

- ERROR\_FOR\_DIVISION\_BY\_ZERO, NO\_ZERO\_DATE, and NO\_ZERO\_IN\_DATE are deprecated and do nothing
  - Use STRICT\_TRANS\_TABLES or STRICT\_ALL\_TABLES, which include those modes and produce an error

Query Optimization: From 0 to 10 (and up to 5.7)

## GIS IMPROVEMENTS & JSON TYPES

#### Find the Closest Starbucks

```
mysql-5.7.5 (osm) > SET @lat:=51.49353; SET @lon:=-0.18340;
mysql-5.7.5 (osm) > SELECT n.node id,
                          n.longitude/10000000 as longitude,
                          n.latitude/10000000 as latitude.
                                                                                  You are here
                          sqrt(pow((latitude/10000000 - @lat) * 111257.67, 2)
                                pow((longitude/10000000 - @lon) * 69450.32, 2))
                          as `distance in metres`
                     FROM nodes n
                      JOIN node tags n t1
                       ON n.node id = n t1.node id
                     JOIN node tags n t2
                       ON n.node id = n t2.node id
                    WHERE
                          n t1.k = 'amenity' and
                          n t1.v = 'cafe' and
                          n t2.k = 'name' and
                          n t2.v = 'Starbucks'
                 ORDER BY 'distance in metres' ASC
                    LIMIT 1;
 node id | longitude | latitude | distance in metres
 699693936 l
               -0.1823 | 51.4945 |
                                     130.9792513096838
1 row in set (0.20 sec)
```

### This Query is Slow

```
mysql-5.7.5 (osm) > EXPLAIN SELECT n.node id,
n.longitude/10000000 as longitude, n.latitude/10000000
as latitude, sqrt(pow((latitude/10000000-
@lat)*111257.67, 2) + pow((longitude/10000000-
@lon)*69450.32, 2)) as `distance in metres` FROM nodes
n JOIN node tags n t1 ON n.node id = n t1.node id JOIN
node tags n t2 ON n.node id = n t2.node id WHERE
n t1.k = 'amenity' and <math>n t1.v = 'cafe' and n t2.k = 'amenity'
'name' and n t2.v = 'Starbucks' ORDER BY `distance in
metres` ASC LIMIT 1\G
id: 1
  select type: SIMPLE
       table: n t1
  partitions: NULL
        type: ALL
possible keys: PRIMARY
         key: NULL
     key len: NULL
         ref: NULL
        rows: 832040
    filtered: 0.00
       Extra: Using where; Using temporary; Using
file sort
```

```
****************** 2. FOW **************
         id: 1
 select type: SIMPLE
       table: n_t2
  partitions: NULL
        type: ref
possible_keys: PRIMARY
        key: PRIMARY
     key len: 8
        ref: osm.n t1.node id
        rows: 3
    filtered: 1.41
       extra: Using where
id: 1
 select type: SIMPLE
       table: n
  partitions: NULL
        type: ref
possible keys: PRIMARY
        key: PRIMARY
     key len: 8
        ref: osm.n_t1.node_id
        rows: 1
    filtered: 100.00
       Extra: NULL
3 rows in set, 1 warning (0.00 sec)
```

#### Can We Optimize it?

We could add a bounding box:

```
mysql-5.7.5 (osm) > ALTER TABLE nodes add index(latitude, longitude);
mysql-5.7.5 (osm) >
      SELECT n.node_id,
       n.longitude/10000000 as longitude,
       n.latitude/10000000 as latitude,
       sqrt(pow((latitude/10000000 - @lat) * 111257.67, 2) +
           pow((longitude/10000000 - @lon) * 69450.32, 2))
           as `distance in metres`
       FROM nodes n
       JOIN node_tags n_t1
          ON n.node id = n t1.node id
       JOIN node tags n t2
          ON n.node_id = n_t2.node_id
                                                                                     This is not a
       WHFRF
                                                                                   square, only an
          n_t1.k = 'amenity' and
          n_t1.v = 'cafe' and
                                                                                   approximation
          n t2.k = 'name' and
          n t2.v = 'Starbucks' and
          n.latitude BETWEEN ((@lat - 1000/111257.67) * 10000000)
                         AND ((@lat + 1000/111257.67) * 10000000) and
          n.longitude BETWEEN ((@lon - 1000/69450.32) * 10000000)
                         AND ((@lon + 1000/69450.32) * 10000000)
       ORDER BY 'distance in metres' ASC
```

LIMIT 1:

mycal-5 7 5 (ocm) > EXPLATA SELECT

#### We Create an Index... and Force It

id   	select_type	table 	partitions 	type   	possible   _keys			SQL ignore ly created			Extra	
1	SIMPLE	n_t1	NULL 	ALL   	PRIMARY	NO		why?		00	Using where;     Using temporary;	
1     	SIMPLE	   n 	   NULL 	   ref   	PRIMARY,     latitude	PRIMARY	8	osm.n_t1.node_id	1	5.00	Using filesort     Using where   	
1	SIMPLE	n_t2	NULL	l ref l	PRIMARY	PRIMARY	Q İ	osm.n t1.node id	j 3	1 1 11	Using where	
	s in set, 1 was		·	++	+	<del>-</del>	+		+			
	-		·	++	+	<del>-</del>	+		-+	+	Ostrig where	
	-5.7.5 (osm) >	EXPLAIN	SELECT F	ROM node	s n FORCE I	NDEX(latit	+	·;	+	filtered	······	
/sql-	·5.7.5 (osm) >	EXPLAIN	SELECT F	ROM node +   type 	es n FORCE I	NDEX(latite +   key 	ude) -+   key_   len	·;	+		Extra   Using where;   Using index;	
/sql-	select_type	EXPLAIN +   table 	SELECT F +	ROM node +   type 	s n FORCE I +   possible   _keys +	NDEX(latite +   key 	ude) -+   key_   len	.; -+	+	filtered	Extra	

examined

not great

the filtering

1 row in set (0.09 se-

# Multiple Range Scans Cannot Be Optimized with BTREE Indexes

- We need quadtrees or R-TREE Indexes for indexing in multiple dimensions
  - The later are implemented in MySQL with the name "SPATIAL indexes", as they only apply to GIS types
- Spatial indexing is available for the first time for InnoDB tables on MySQL 5.7.5

# Creating a Spatial Index

```
mysql-5.7.5 (osm) > ALTER TABLE nodes
                    ADD COLUMN coord GEOMETRY NOT NULL:
Query OK, 0 rows affected (21.80 sec)
Records: 0 Duplicates: 0 Warnings: 0
mysql-5.7.5 (osm) > UPDATE nodes
                    SET coord = point(longitude/10000000,
                                      latitude/10000000);
Query OK, 2865312 rows affected (34.66 sec)
Rows matched: 2865312 Changed: 2865312 Warnings: 0
mysql-5.7.5 (osm) > ALTER TABLE nodes add SPATIAL
index(coord);
                                                      This is new
Query OK, 0 rows affected (1 min 50.00 sec)
                                                        in 5.7
Records: 0 Duplicates: 0 Warnings: 0
```

### **New Query**

```
mysql> SET @area := envelope(linestring(POINT(@lon - 500/69450.32, @lat - 500/111257.67), POINT(@lon +
500/69450.32, @lat + 500/111257.67)));
mysql> SELECT n.node id,
              x(n.coord) as longitude,
              y(n.coord) as latitude,
              st distance(POINT(@lon, @lat), coord) as distance
       FROM nodes n
       JOIN node tags n t1
           ON n.node_id = n_t1.node_id
       JOIN node tags n t2
           ON n.node id = n t2.node id
      WHFRF
           n t1.k = 'amenity' and
           n t1.v = 'cafe' and
           n t2.k = 'name' and
           n t2.v = 'Starbucks' and
           st_within(coord, @area)
      ORDER BY st distance(POINT(@lon, @lat), coord) ASC
      LIMIT 1;
```

We can use any shape we want thanks to 5.6 **improvements** 

#### **Better Performance**

```
mysql-5.7.5 (osm) > SELECT ...;
  node id | longitude | latitude | distance
  699693936 | -0.1822879 | 51.4944808 | 0.0014631428672541478
                                                                             This field is
                                                                                almost
1 row in set (0.02 sec)
                                                                                useless
mysql-5.7.5 (osm) > EXPLAIN SELECT ...;
 id | select | table | parti | type | possible | key | key | ref
                                                           | rows | filtered | Extra
                                             | _len |
    | _type |
                 | tions |
                              | keys
  1 | SIMPLE | n
                 | NULL | range | PRIMARY | coord | 34 | NULL
                                                                  100.00 | Using where;
                                                                          | Using filesort
                              coord,
  1 | SIMPLE | n t1 | NULL | ref
                             | PRIMARY | PRIMARY | 8
                                                 | osm.n. | 3 |
                                                                  1.41 | Using where
                                                  | node_id |
  1 | SIMPLE | n_t2 | NULL | ref | PRIMARY | PRIMARY | 8
                                                 osm.n. | 3 | 1.41 | Using where
                                                   | node_id |
3 rows in set, 1 warning (0.00 sec)
```

# **Better Filtering**

mysql-5.7.5 (osm) > SHOW STATUS LIKE 'Hand%';

#### Not using the index:

#### Using the BTREE index: Using the SPATIAL index:

,	
Variable_name	Value
+	
Handler_commit	1
Handler_delete	0
Handler_discover	0
Handler_external_lock	6
Handler_mrr_init	0
Handler_prepare	0
Handler_read_first	1
Handler_read_key	1914
Handler_read_last	0
Handler_read_next	1954
Handler_read_prev	0
Handler_read_rnd	1
Handler_read_rnd_next	833426
Handler_rollback	0
Handler_savepoint	0
Handler_savepoint_rollback	0
Handler_update	0
Handler_write	1
+	+

Using the SPATIAL I				
Variable_name	+   Value			
Handler_commit Handler_delete Handler_discover Handler_external_lock Handler_mrr_init Handler_prepare Handler_read_first Handler_read_last Handler_read_last Handler_read_prev Handler_read_rnd Handler_read_rnd Handler_rollback Handler_savepoint Handler_update Handler_write	1			
8 rows in set (0 00 sec)				

18 rows in set (0.00 sec)

#### **Geohash Functions**

- Useful to index coordinates with a BTREE
  - It could be specially useful combined with indexed STORED columns (emulating quadtrees)

More on Geohashing: <a href="http://mysqlserverteam.com/geohash-functions/">http://mysqlserverteam.com/geohash-functions/</a>

# **GeoJSON Functions**

```
mysql-5.7.5 (osm) > SELECT nm.v, ST_AsGeoJson(n.coord)
                         FROM node tags n t
                         JOIN nodes n USING (node id, version)
                         JOIN node tags nm USING (node id, version)
                        WHERE n t.k='tourism' AND
                              n t.v='attraction' AND
                              nm.k='name';
                        | ST AsGeoJson(n.coord)
                   | {"type":"Point","coordinates":[-0.1560632,51.3821745]}
 BedZED
 Blewcoat School | {"type":"Point","coordinates":[-0.1360484,51.4983179]}
 Camden / Buck Street Market | {"type":"Point","coordinates":[-0.143193,51.5400398]}
 Camden Lock Village | {"type":"Point","coordinates":[-0.1447181,51.5416552]}
 Wimbledon Windmill
                         [ "type":"Point","coordinates":[-0.2315468,51.4376583]]
```

# **GeoJSON Functions (cont.)**

```
$ ./sandboxes/msb_5_7_5/use osm -B -e "SET SESSION group_concat_max_len = 10000; SELECT
CONCAT('{\"type\":\"FeatureCollection\", \"features\":[ ',
GROUP_CONCAT(CONCAT('{\"type\":\"Feature\", \"geometry\":', ST_AsGeoJson(n.coord),
',\"properties\":{\"name\":\"',nm.v,'\"}}')), ' ]}') FROM node_tags n_t JOIN nodes n USING
(node_id, version) JOIN node_tags nm USING (node_id, version) WHERE n_t.k='tourism' and
n t.v='attraction' AND nm.k='name'"
```

#### http://geojsonlint.com/



# **Open Issues**

 The SRID can be set and retrieved, but all operations are done in squared euclidean coordinates:

```
mysql-5.7.5 (osm) > SET @p1 := GeomFromText('POINT(-1 51)',
4326);
Query OK, 0 rows affected (0.00 sec)
mysql-5.7.5 (osm) > SET @p2 := GeomFromText('POINT(0 51)', 4326);
Query OK, 0 rows affected (0.00 sec)
mysql-5.7.5 (osm) > SET @p3 := GeomFromText('POINT(-1 52)',
4326);
Query OK, 0 rows affected (0.00 sec)
mysql-5.7.5 (osm) > SELECT srid(@p1);
 srid(@p1)
       4326
1 row in set (0.00 sec)
```

## **New JSON Native Data Type**

 Since 5.7.8, MySQL allows columns defined with the JSON data type:

```
mysql-5.7.8 > CREATE TABLE json test(id int PRIMARY KEY
auto increment, content JSON);
Query OK, 0 rows affected (0.03 sec)
mysql-5.7.8 > INSERT INTO json test (content) VALUES ('{"type":
"correct ison"}');
                                                             They get
Query OK, 1 row affected (0.00 sec)
                                                            validated on
                                                              insert
mysql-5.7.8 > INSERT INTO json_test (content) VALUES ('{"type :
"incorrect json}');
ERROR 3140 (22032): Invalid JSON text: "Missing a closing quotation
mark in string." at position 24 in value (or column) '{"type":
"incorrect ison}'.
```

#### **JSON functions**

- MySQL includes almost all functions to manipulate JSON that you may think of:
  - Validation test: JSON\_TYPE
  - Object creation: JSON\_ARRAY, JSON\_MERGE, ...
  - Searching: JSON\_EXTRACT
  - Modifying: JSON\_SET, JSON\_INSERT, ...

# Indexing JSON

- JSON Columns cannot be indexed:
   mysql [localhost] {msandbox} (test) > ALTER TABLE
   json\_test ADD INDEX(content);
   ERROR 3152 (42000): JSON column 'content' cannot be
   used in key specification.
- However, they can be compared with regular fields and use indexes thanks to virtual columns

Query Optimization: From 0 to 10 (and up to 5.7)

#### **CONCLUSIONS**

# 5.7/10.1 About to be released

- Both are currently in Release Candidate
- Unless you are desperate for a feature, skip the first releases (or backport it to your current version)

#### MySQL 5.7 New Features

- MySQL 5.6 seemed Percona Server-inspired
- MySQL 5.7 seems MariaDB/Galera-inspired
  - Competition is <u>always good for consumer</u>



#### Many Optimizer Advantages Have to Be Manually Enabled

Modifying on a per-query basis:

```
SET optimizer_switch='batched_key_access=on'; SET join_cache_level=8; # for MariaDB
```

in order to take advantage of them make the features useless unless you are fine-tuning

- I expect that to change in the future

# I Herby Declare MyISAM as Dead

- All major MyISAM-only features are now on MySQL 5.7
  - FULLTEXT
  - GIS
  - Transportable tables
- There are still reasons to use MyISAM
  - MyISAM is still required for the mysql schema and non-durable temporary tables (WIP)



#### **Benchmarks**

- Do not trust first party benchmarks
  - In fact, do not trust 3rd party benchmarks either
- Only care about the performance of your application running on your hardware





#### **Not to Miss**

- Official track:
   MySQL at Wikipedia: How we
   do relational data at the
   Wikimedia Foundation
- Do you want to do query optimization for a website with 20 Billion views per month?

http://grnh.se/0y4pxm



# Thank You for Attending!

- Do not forget, after the session finishes, to please login with your Percona Live account and "Rate This Session"
- Special thanks to in order by rand() to: Morgan Tocker, Sean Pringle, David Hildebrandt, Bill Karwin, Domas Mituzas, Mark Callaghan, Shlomi Noach, Mark Bergsma, Valerii Kravchuk, Miguel Ángel Nieto, Dimitri Kravtchuk, Olav Sandstå and the whole Wikimedia Team, and all people at the MariaDB, Percona and MySQL/Oracle teams, and the Percona Live Organization and Sponsors