

CSE2312-001, 002 (Spring 2022)

Homework #5

Notes:

With this homework, we continue writing assembly functions for the RPi 3b/3b+.

All numbers are in base-10 unless otherwise noted.

If part of a problem is not solvable, explain why in the answer area.

The target date to complete this homework set is April 5, 2022.

This homework set will not be graded, but please solve all of the problems to prepare for the quizzes and exams.

1. Suppose that BUSINESS2 structure is defined as:

```
typedef struct _BUSINESS2
{
    uint32_t taxId;
    char name[27];
    char direction;
    char street[35];
    uint32_t addNo;
    char city[30];
    char state[3];
    uint32_t zip;
} BUSINESS2;
```

Show the relative offset of each field in the structure from the beginning of the structure for the unpacked (default alignment) case:

For Packed, take proposed starting location, divide by size of variable type.
If no remainder, no padding is needed, otherwise add padding bytes
equal to the remainder.

uint32_t taxId; is 0 mod 4 = 0, yes => valid start	0 -> 3
char name[27]; is 4 mod 1 = 0, yes => valid start	4 -> 30
char direction; is 31 mod 1 = 0, yes => valid start	31 -> 31
char street[35] is 32 mod 1 = 0, yes => valid start	32 -> 66
uint32_t addNo; is 67 mod 4 = 0, no => need padding +1	68 -> 71
char city[30]; is 72 mod 1 = 0, yes => valid start	72 -> 102
char state[3]; is 102 mod 1 = 0, yes => valid start	102 -> 104
uint32_t zip; is 105 mod 4 = 0, no => need padding +3	108 -> 111

since we started at zero, we add 1 so total = 112
Un-packed (STANDARD!)
total = 112

Show the relative offset of each field in the structure from the beginning of the structure for the packed case:

uint32_t taxId; 4	0 -> 3
char name[27]; 27	4 -> 30
char direction; 1	31 -> 31
char street[35]; 35	32 -> 66
uint32_t addNo; 4	67 -> 70
char city[30]; 30	71 -> 100
char state[3]; 3	101 -> 103
uint32_t zip; 4	104 -> 107
total = 108 byte	

2. Write assembly functions that implement the following C functions:

- a. `int32_t sumS32(const int32_t x[], uint32_t count)`
 // returns sum of the values in the array (x) containing count entries.
- b. `int32_t dotpS32(const int32_t x[], const int32_t y[], uint32_t count)`
 // returns the dot product of the values in the arrays (x and y)
 containing count entries.
- c. `uint32_t countAboveLimit(const int32_t x[], int32_t limit, uint32_t count)`
 // returns number of values in the array (x) containing count entries
 that are > limit
- d. `int32_t findCityAligned (const char city[], const BUSINESS2 business[], uint32_t
count)`
 // returns the index of the first entry in the array (business) containing
count entries which matches the requested city. If the city is not
found, return a value of -1. You can assume that C default alignment is
used for this problem.
- e. `int32_t findCityPacked (const char city[], const BUSINESS2 business[], uint32_t
count)`
 // returns the index of the first entry in the array (business) containing
count entries which matches the requested city. If the city is not
found, return a value of -1. You can assume that C packing is used for
this problem.

a. 0.125

32b hex value = 0x3E000000

32b hex value = 0xC1790000

32b hex value = 0x3EAAAAAA

32b hex value = 0x437B4200

32b hex value = 0x45000000

4. Assume float $x = 1048576$.

a. Calculate the smallest positive number that can be added to x that will not be lost in the mantissa.

b. In general, what is the ratio of the large to the smallest single-precision floating point number that can be added together without a loss of accuracy? Consider the cases you add a number to $x = 1048576$ or $x = 2097151$.

a) Since 104857's (220) mantissa is all zeros, the smallest we could add would set the least significant bit to 1.

S	Exp	Mantissa	v
0	10010011	000000000000000000000000	

changing to:

0	10010011	000000000000000000000001	
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With a mantissa of 23 digits we have $2^{(20-23)}$ so 2^{-3}

$2^{-3} = 1/8$ or 0.125

b)

For any given number expressed as 2^n , the smallest value we could add would be $2^{(n-23)}$ without losing accuracy.

So, as a ratio, of $2^n : 2^{(n-23)}$ then substituting n with 23, we get $(2^{23}) : (2^0)$ or $(2^{23}) : 1$