

Faculty Development Program (FDP) On
Quantum Computing Leverages

June 16 to 20, 2025

Organized by LDRP Institute of Technology and Research, Gandhinagar

Foundation of Quantum Computing

June 16, 2025

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Acknowledgement

- **Udemy's** “*Quantum Computing and Introduction to Quantum Machine Learning*” October 2024.
- **NPTEL-AICTE FDP** on “*Introduction to Quantum Computing: Quantum Algorithms and Qiskit*” during July-August 2024.
- AICTE's **ATAL FDP** on “*Quantum Machine Learning*” at CHRIST UNIVERSITY Bangalore in Nov-2024.
- **YouTube:** “*Quantum Computing Course – Math and Theory for Beginners*” at freeCodeCamp.org (<https://youtu.be/tsbCSkvHhMo>)



Presentation Outline

- Introduction
- Essential Mathematics behind Quantum Computing
 - Complex Numbers, Linear Algebra, Matrices
- Quantum Bits (Qubits)
 - Single Qubit, Multiple Qubits, Operations on Qubits
- Quantum Phenomena
 - Entanglement, Phase kickback
- Quantum Algorithms



Introduction

- Applications of Quantum Computing
 - Artificial Intelligence, Machine Learning, Natural Language Processing (NLP)
 - Cybersecurity / Cryptography / Blockchain
 - Drug Discovery, Molecular Modelling and Chemical Research
 - Fundamental Science and Mathematics
 - Optimization (Logistic, Finance, Manufacturing, Supply Chain, Portfolio, Search, Flight path, Fuel Consumption, Radar signal processing, satellite positioning)
 - Modelling (Material, Finance, Climate, Weather)



Introduction

- Why Quantum Computing?
- Problem Hardness / Complexity [Time/Space]
 - N
 - $N \log (N)$
 - N^2
 - N^5
 - 2^N
 - 7^N



Introduction: Classical Vs Quantum Computing

Classical Computing	Quantum Computing
Binary Digits (Bits)	Quantum Bits (Qubits)
States: 0, 1	$ \psi\rangle = \alpha 0\rangle + \beta 1\rangle$, States: $ 0\rangle, 1\rangle$, α, β : Complex Numbers, +: Superposition
State: Deterministic	State: Probabilistic Probability of $ 0\rangle$: $ \alpha ^2$, Probability of $ 1\rangle$: $ \beta ^2$
Logical Gates: NOT, AND, OR, XOR	Gates: Pauli X Gate (Matrix/Vector representation)
Algebra: Boolean	Algebra: Linear
Operation: Irreversible	Operation: Reversible
State Space Size: Linear (N)	State Space Size: Exponential (2^N)



Numbers: Real, Imaginary, Complex

Real Numbers

$$x^2 = 4$$

$$\sqrt{x^2} = \sqrt{4}$$

$$x = \pm 2$$

Imaginary Numbers

$$x^2 = -4$$

$$\sqrt{x^2} = \sqrt{-4}$$

$$x = \pm \sqrt{-4}$$

$$x = \pm \sqrt{4} \sqrt{-1}$$

$$x = \pm 2\sqrt{-1}$$

Let $i = \sqrt{-1}$

$$x = \pm 2i$$



Numbers: Real, Imaginary, Complex

Complex Numbers = Real Numbers + Imaginary Numbers

$a + i b$ where $a, b \in \mathbb{R}$

E.g.

$$3 + 4i$$

$$-5 - 2i$$

$$\sqrt{3} + i\sqrt{5}$$



Numbers: Real, Imaginary, Complex

Operations on Complex Numbers

Addition:

$$(3+2i)+(2-4i)$$
$$=5-2i$$

Multiplication:

$$(3+2i)+(2-4i)$$
$$=5+4i-12i-12i^2$$
$$=5-8i-12(\sqrt{-1})^2$$
$$=5-8i+12$$
$$=17-8i$$



Numbers: Real, Imaginary, Complex

Complex Conjugate: Negate the imaginary part

$$(a+ib)$$

Example:

$$(-2+3i)$$

$$(a+ib)^*$$

$$(-2+3i)^*$$

$$=(a-ib)$$

$$=(-2-3i)$$



Numbers: Real, Imaginary, Complex

Complex Conjugate: Multiplication of any complex number with its complex conjugate will always be a real number

Example:

Number: $(-2+3i)$

$$(-2+3i) \times (-2+3i)^*$$

$$= (-2+3i) \times (-2-3i)$$

$$= 4 - 6i + 6i - 9i^2$$

$$= 4 - 9i^2$$

$$= 4 + 9$$

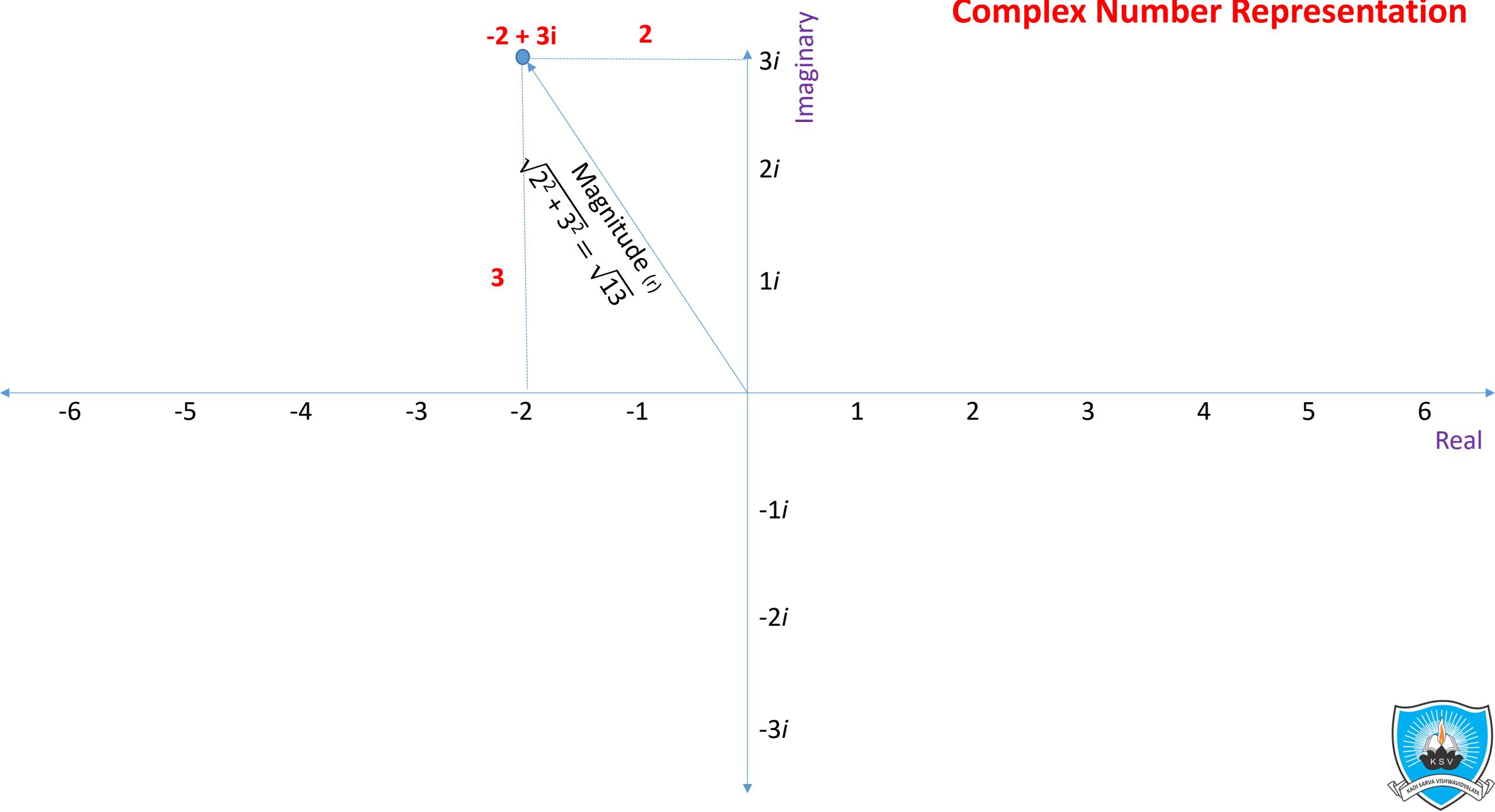
$$= 13$$

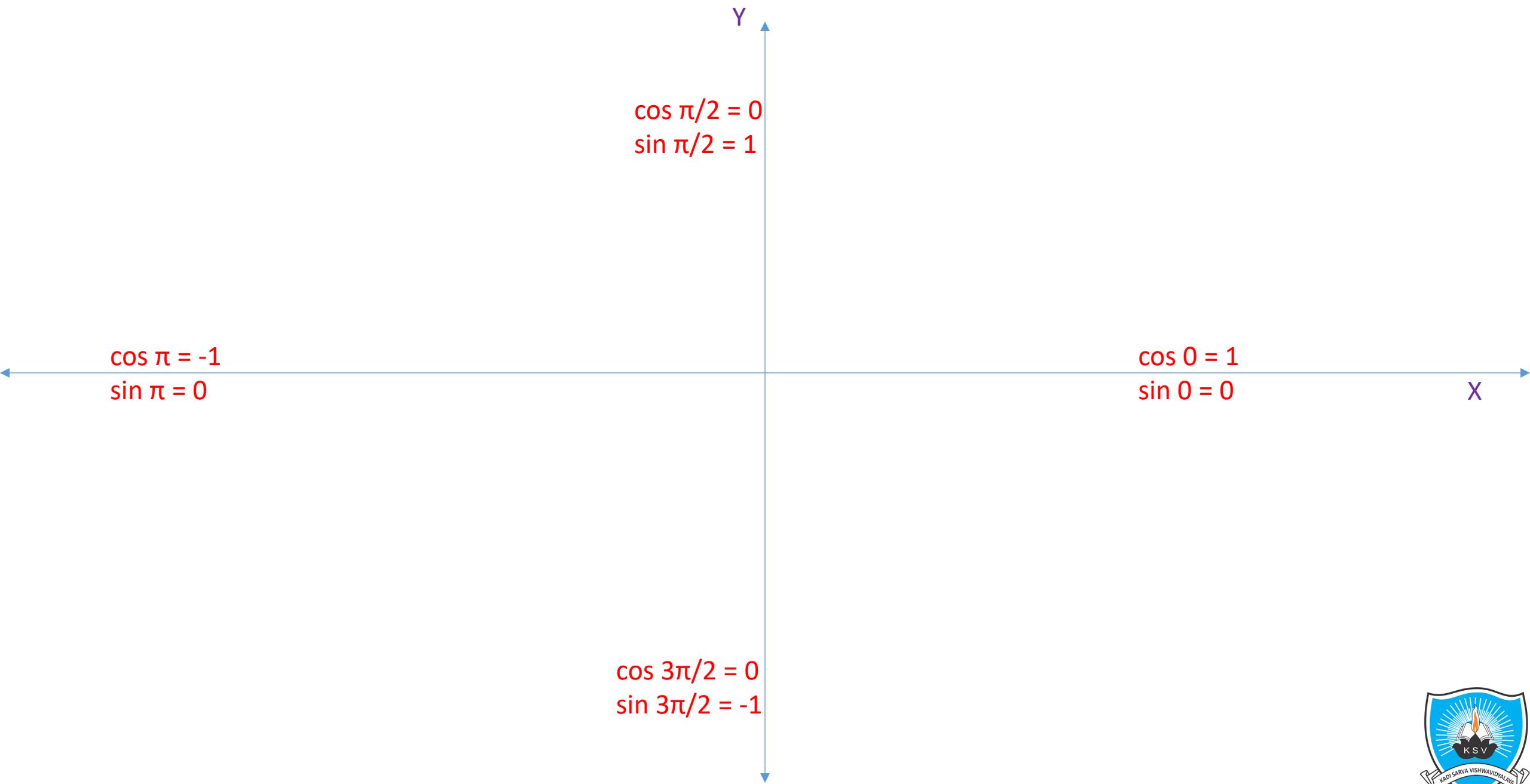


Complex Numbers on a Number Plane

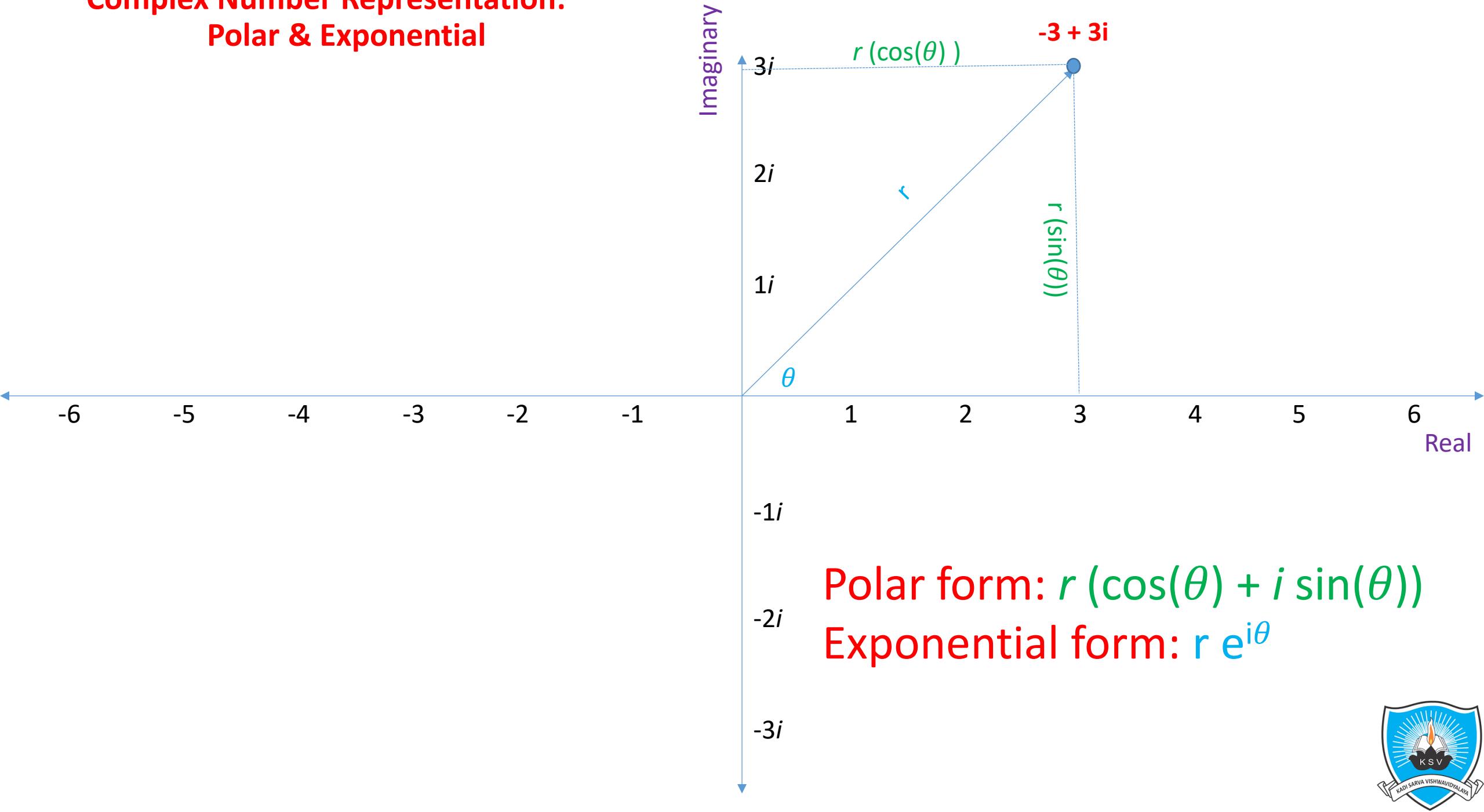


Complex Number Representation





Complex Number Representation: Polar & Exponential



Convert $1+i$ into polar form and exponential form

$r(\cos(\theta) + i \sin(\theta))$ Polar form

$= r e^{i\theta}$ Exponential form

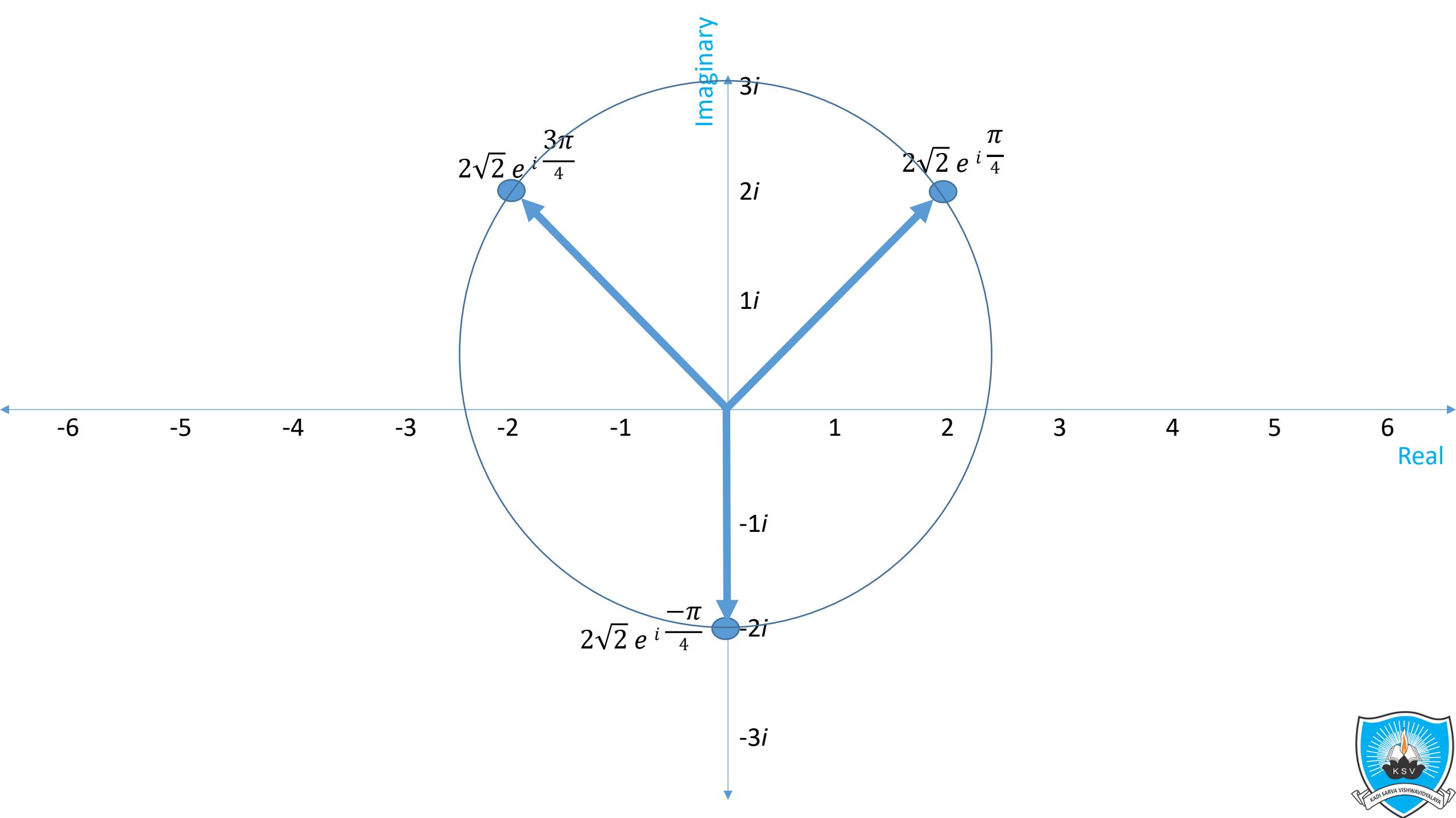


$$= \sqrt{1+1} \left(\cos\left(\frac{\pi}{4}\right) + i \sin\left(\frac{\pi}{4}\right) \right)$$

$$= \sqrt{2} \left(\cos\left(\frac{\pi}{4}\right) + i \sin\left(\frac{\pi}{4}\right) \right) \quad \text{Polar form}$$

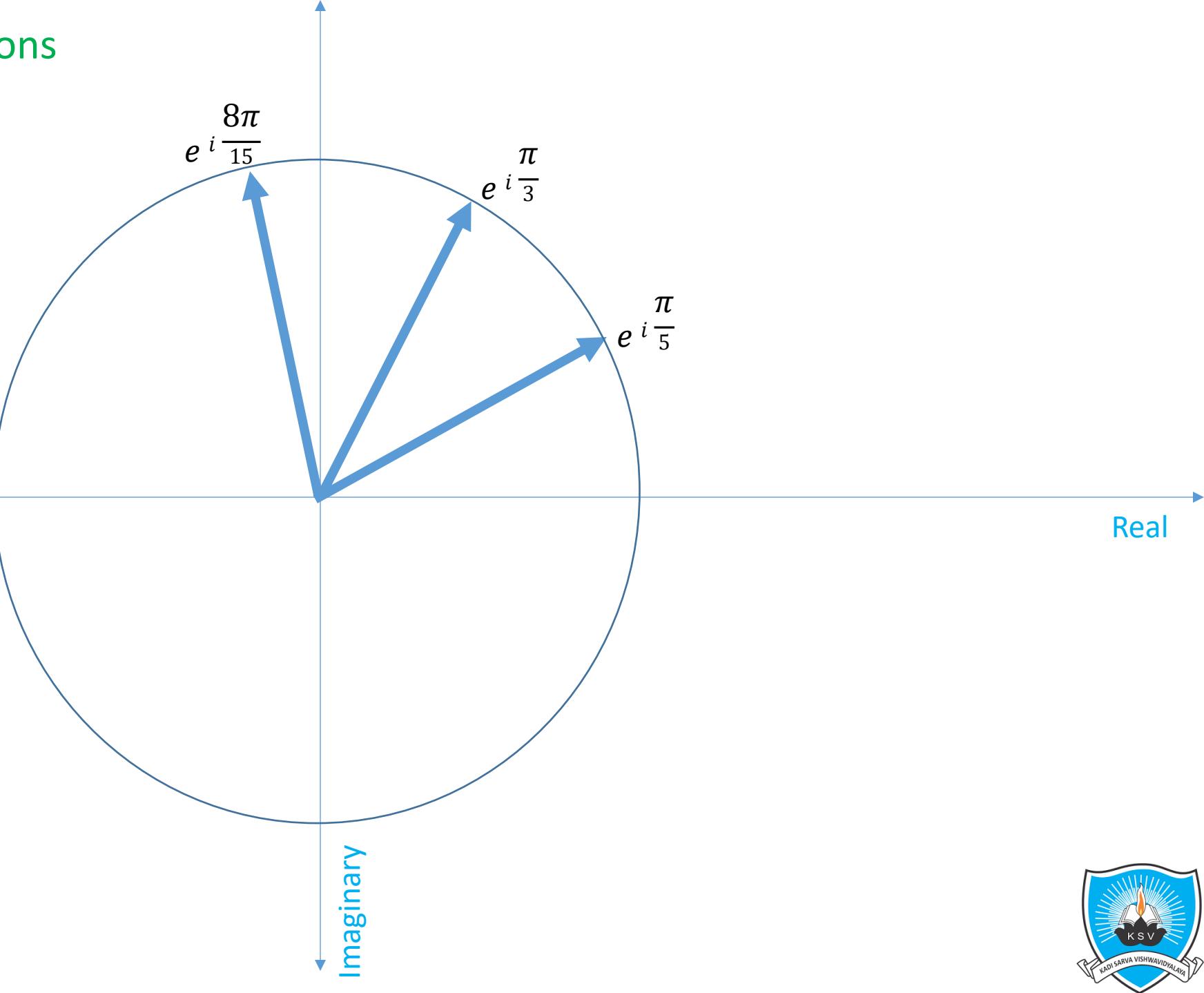
$$= \sqrt{2} e^{i \frac{\pi}{4}}$$

Exponential form



Complex Numbers Multiplications

$$\begin{aligned} e^{i\frac{\pi}{5}} \times e^{i\frac{\pi}{3}} &= e^{i(\frac{\pi}{5} + \frac{\pi}{3})} \\ &= e^{i\frac{8\pi}{15}} \end{aligned}$$



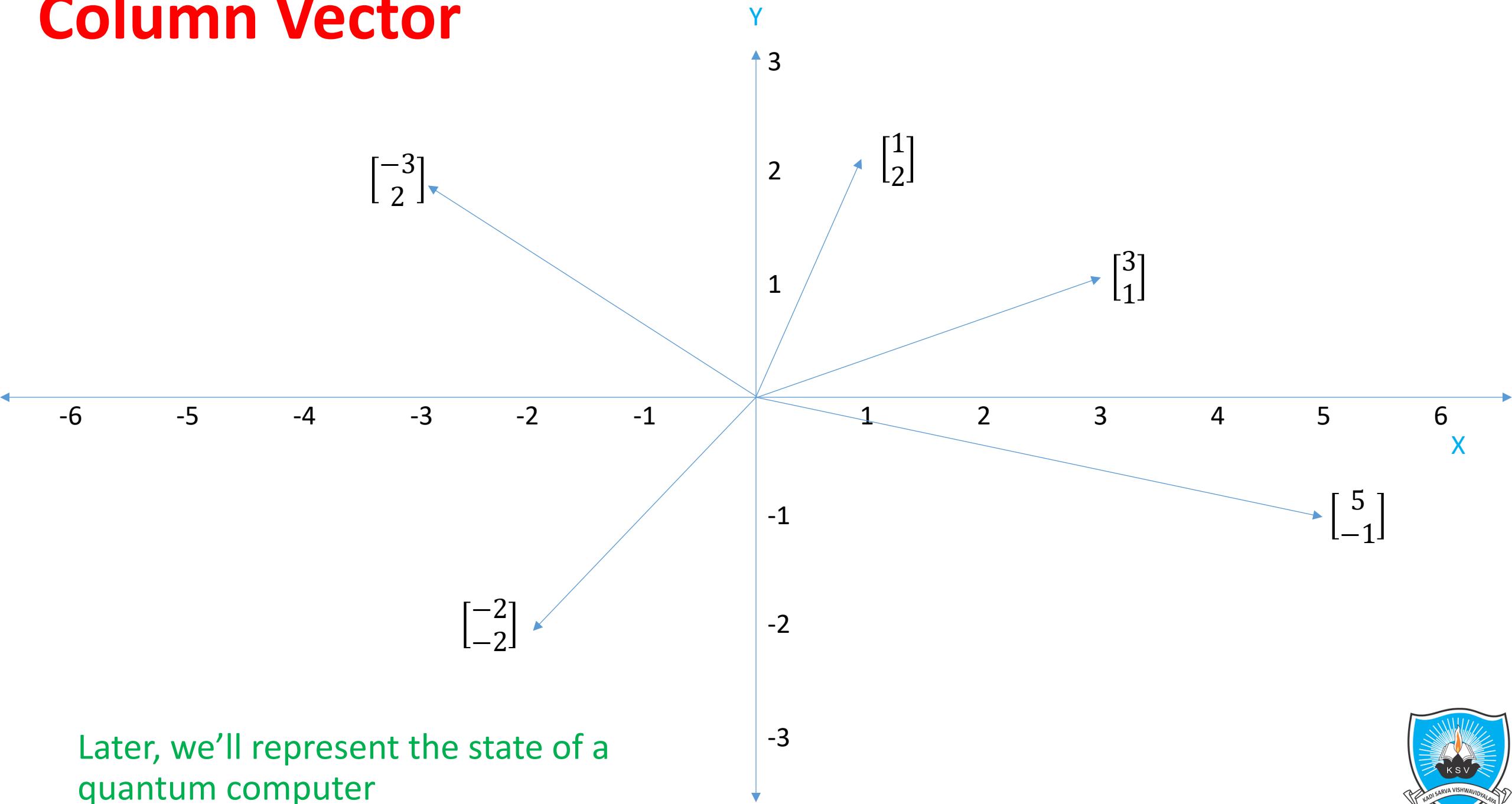
Matrices

Operation: Addition, Multiplication, Scalar

Single column matrix: Column Vector



Column Vector

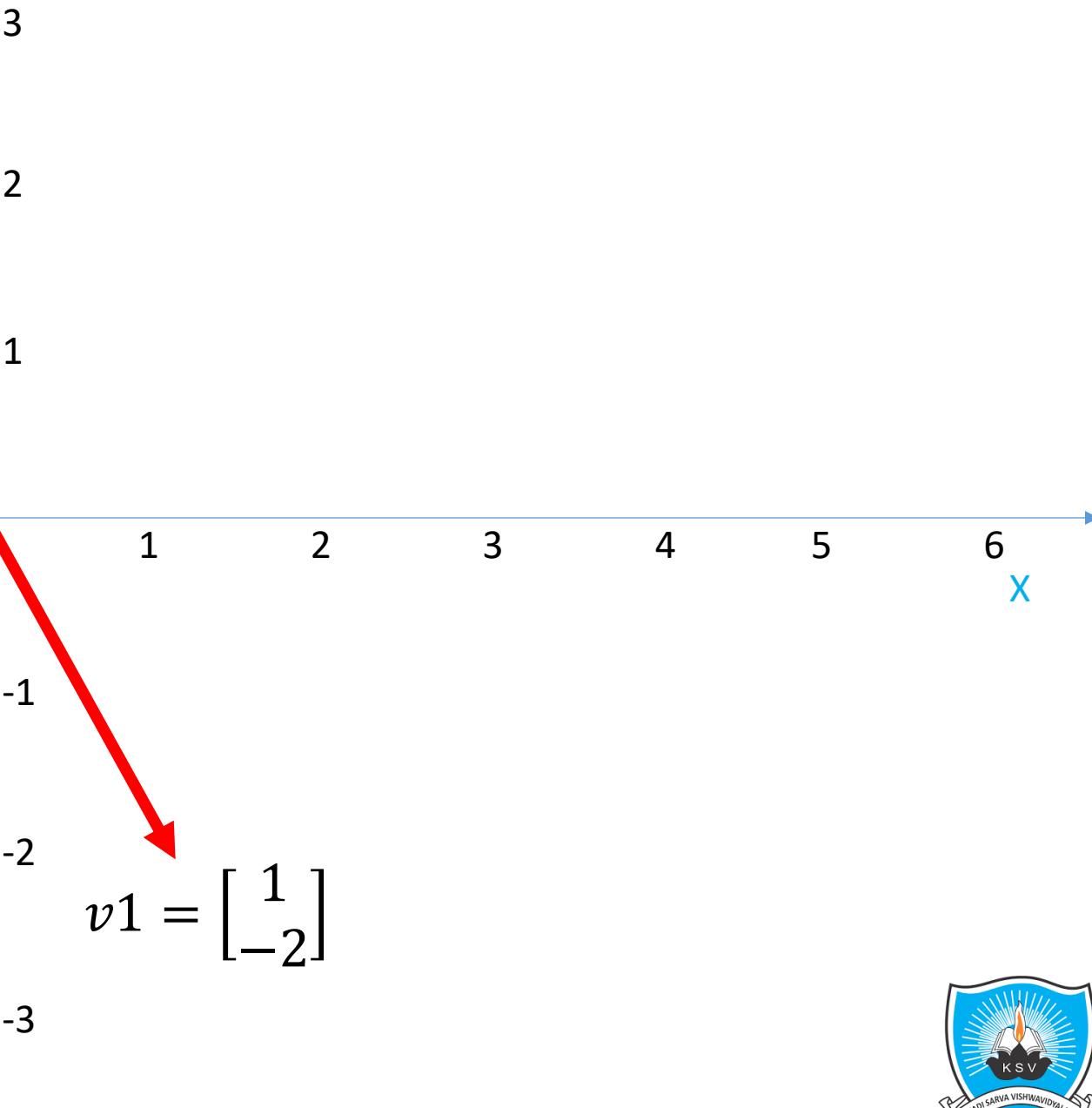


When we multiple a column vector (E.g. v_1) by matrix, we get another vector (E.g. v_2). That means, v_1 is transformed to v_2 .

$$\begin{bmatrix} 2 & 2 \\ 2 & 1 \end{bmatrix} x \begin{bmatrix} 1 \\ -2 \end{bmatrix} = \begin{bmatrix} -3 \\ 0 \end{bmatrix}$$

$$v_2 = \begin{bmatrix} -3 \\ 0 \end{bmatrix}$$

-6 -5 -4 -3 -2 -1



Ex: Transform vector $\begin{bmatrix} 1 \\ -2 \end{bmatrix}$ using matrix $\begin{bmatrix} 2 & 2 \\ 2 & 1 \end{bmatrix}$.

Identity Matrix

$$I = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$$

$$I = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$I \times A = A$$

$$\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} x \\ y \end{bmatrix}$$

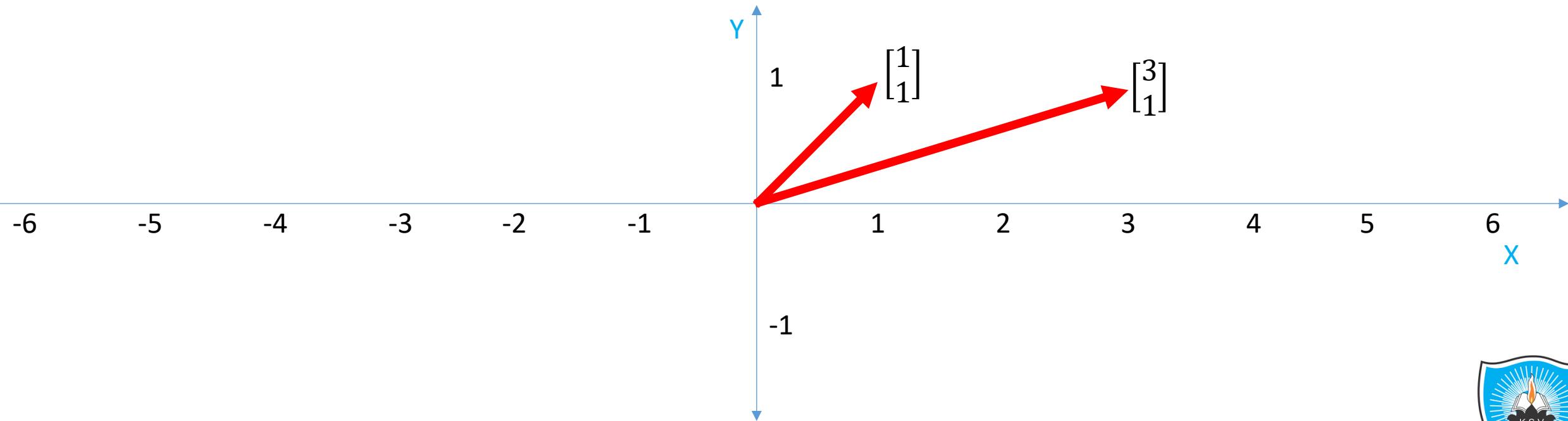
Matrix Inverse

$$A = \begin{bmatrix} 1 & 2 \\ 2 & -1 \end{bmatrix} \quad \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

$$A \times I = \begin{bmatrix} 3 \\ 1 \end{bmatrix}$$

$$A \times A^{-1} = I$$

$$A^{-1} = \begin{bmatrix} \frac{1}{5} & \frac{2}{5} \\ \frac{2}{5} & -\frac{1}{5} \end{bmatrix} \quad \begin{bmatrix} \frac{1}{5} & \frac{2}{5} \\ \frac{2}{5} & -\frac{1}{5} \end{bmatrix} \times \begin{bmatrix} 3 \\ 1 \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$



Complex Conjugate: Negate the imaginary part

$$A = \begin{bmatrix} 2 + 3i & 0 \\ 5 & 3 - i \end{bmatrix}$$

$$A^* = \begin{bmatrix} 2 - 3i & 0 \\ 5 & 3 + i \end{bmatrix}$$

Matrix Transpose:

$$\begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{bmatrix}^T = \begin{bmatrix} 1 & 4 \\ 2 & 5 \\ 3 & 6 \end{bmatrix}$$

$$[A^*]^T = [A^T]^* = A^\dagger$$

Unitary Matrix

$$U U^\dagger = I$$

i.e. U^\dagger is inverse of U

Hermitian Matrix

$$H = H^\dagger$$



Eigen Value & Vector

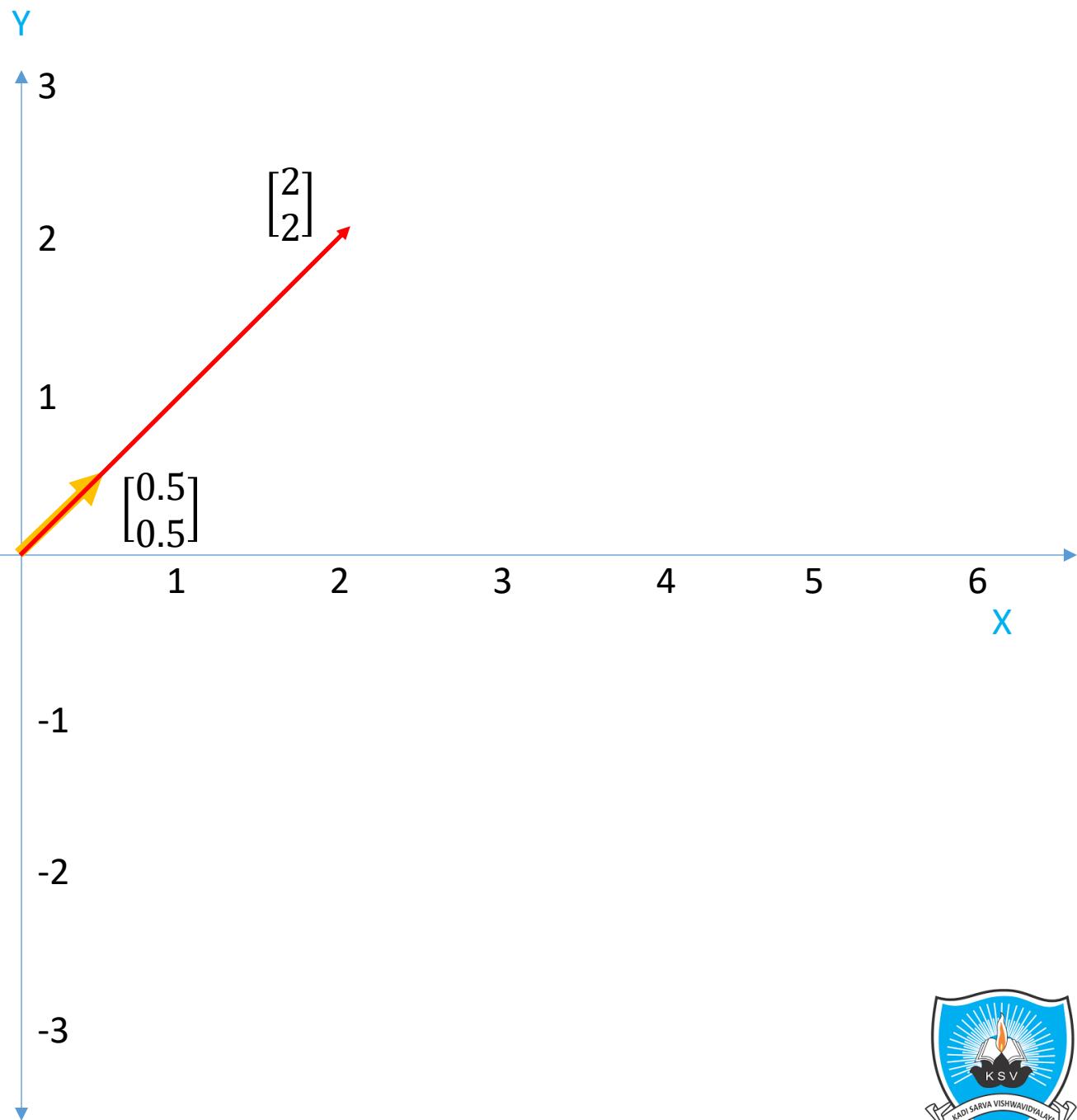
$$\begin{bmatrix} 2 & 2 \\ 2 & 3 \end{bmatrix} \times \begin{bmatrix} 0.5 \\ 0.5 \end{bmatrix} = \begin{bmatrix} 2 \\ 2 \end{bmatrix} = 4 \begin{bmatrix} 0.5 \\ 0.5 \end{bmatrix}$$

$$A\vec{v} = \lambda\vec{v}$$

Eigen Vector : \vec{v} Eigen Value: λ



Ex: Find our Eigen value for $\begin{bmatrix} 2 & 0 \\ 1 & 2 \end{bmatrix} \times \begin{bmatrix} 0 \\ 3 \end{bmatrix}$.



Qubit (Quantum Bit): Dirac (Bra-Ket) Notation

ket notation: $|0\rangle |1\rangle$

$$|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}, |1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \text{ [Vector representation]}$$

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$

$\begin{pmatrix} \alpha \\ \beta \end{pmatrix}$ Element in a complex Hilbert (vector) space.

A Hilbert space is the mathematical framework used to describe the state space of quantum systems. Loosely speaking, it is a complete vector space equipped with an inner product, which allows you to measure angles and lengths of vectors. In quantum computing, these vectors represent quantum states.

bra notation: $\langle 0 | \langle 1 |$

bra (row vector) is a conjugate transpose of a *ket* (column vector).



Qubit (Quantum Bit)

Classical Computer: Bit: Can be either 0 or 1.

Quantum Computer: Qubit: Can be both 0 or 1, at the same time.

Qubit: Any quantum particle that has two states. For example, photon of the light can be polarized horizontally or vertically.

$$|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \quad |1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

Superposition: Quantum particle is in two states simultaneously.

$$|\psi\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} \quad \begin{array}{l} \alpha: \text{How much the qubit is in } |0\rangle \text{ state.} \\ \beta: \text{How much the qubit is in } |1\rangle \text{ state.} \end{array}$$

$$|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

$\alpha = 1$ (100% the qubit is in $|0\rangle$ state)
0% the qubit is in $|1\rangle$ state

$$|1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$\beta = 1$ (100% the qubit is in $|1\rangle$ state)
0% the qubit is in $|0\rangle$ state



Qubit Measurement

When the Qubit is in superposition, it can have both values, 0 and 1. However, when we measure it, it collapses to one of these two states, and gives result of either 0 or 1.

Then, you may ask, what is the role of α and β , then?

α and β are the probability of measuring 0 or 1, respectively.

Probability of measuring $|\psi\rangle$ as 0 is: $|\alpha|^2$

Probability of measuring $|\psi\rangle$ as 1 is: $|\beta|^2$

$$|\psi\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} \frac{\sqrt{3}}{2} \\ \frac{1}{2} \end{pmatrix}$$

Probability of measuring $|\psi\rangle$ as 0 is= $|\alpha|^2 = \frac{3}{4} = 75\%$
Probability of measuring $|\psi\rangle$ as 1 is= $|\beta|^2 = \frac{1}{4} = 25\%$

$$|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

$\alpha = 1$ (100% the qubit is in $|0\rangle$ state)
0% the qubit is in $|1\rangle$ state

$$|1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$\beta = 1$ (100% the qubit is in $|1\rangle$ state)
0% the qubit is in $|0\rangle$ state

$$|\alpha|^2 + |\beta|^2 = 1$$



Dirac Notation

$$|\psi\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix}$$

$$|\psi\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} \alpha \\ 0 \end{pmatrix} + \begin{pmatrix} 0 \\ \beta \end{pmatrix}$$

$$|\psi\rangle = \alpha \begin{pmatrix} 1 \\ 0 \end{pmatrix} + \beta \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

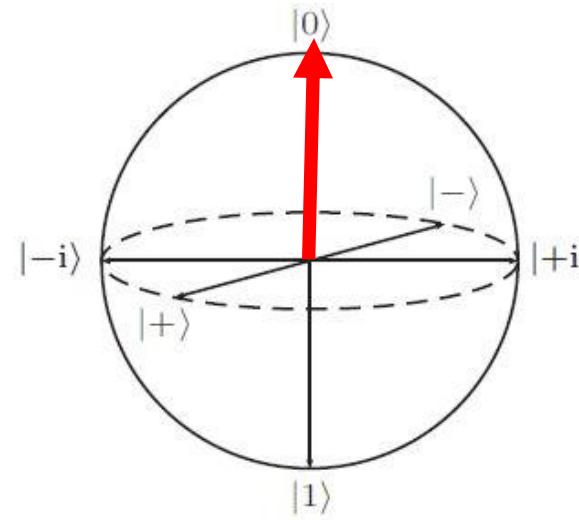
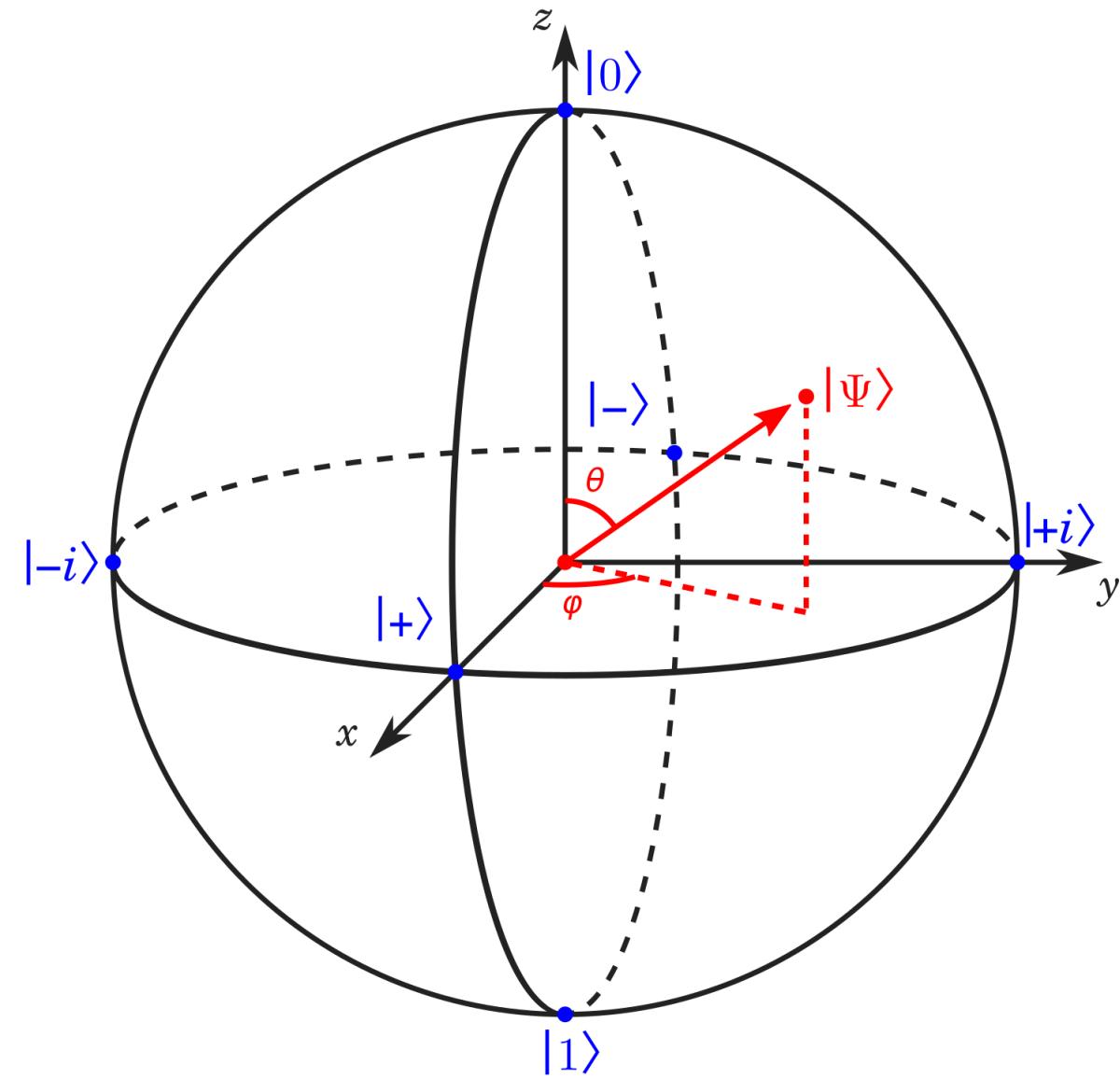
$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$

(Dirac Notation)

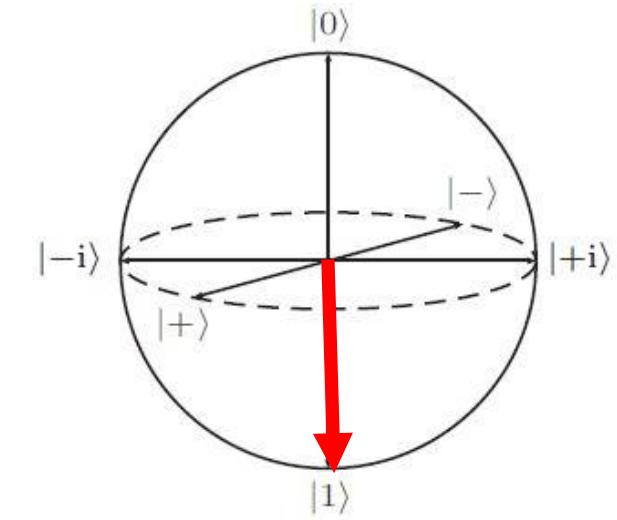
Ex: Convert this into diract notation. $|\psi\rangle = \begin{pmatrix} \frac{1}{2} \\ \frac{2}{2\sqrt{3}} \\ \frac{4}{4} \end{pmatrix}$



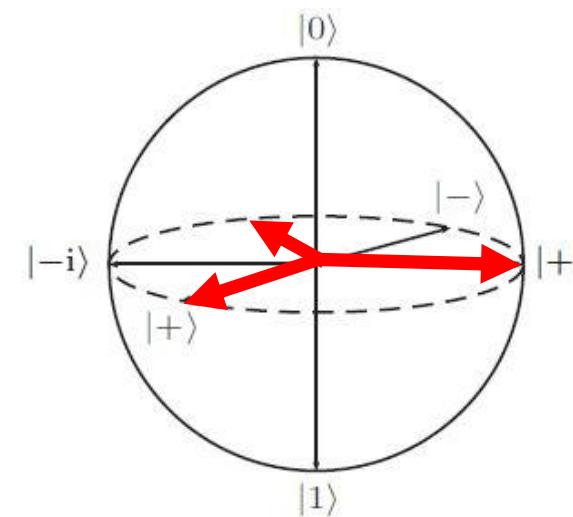
Bloch Sphere



$|0\rangle$

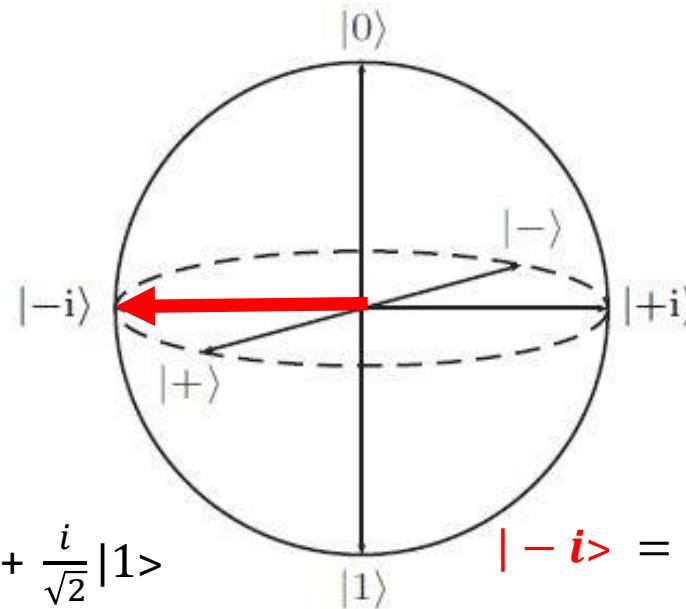
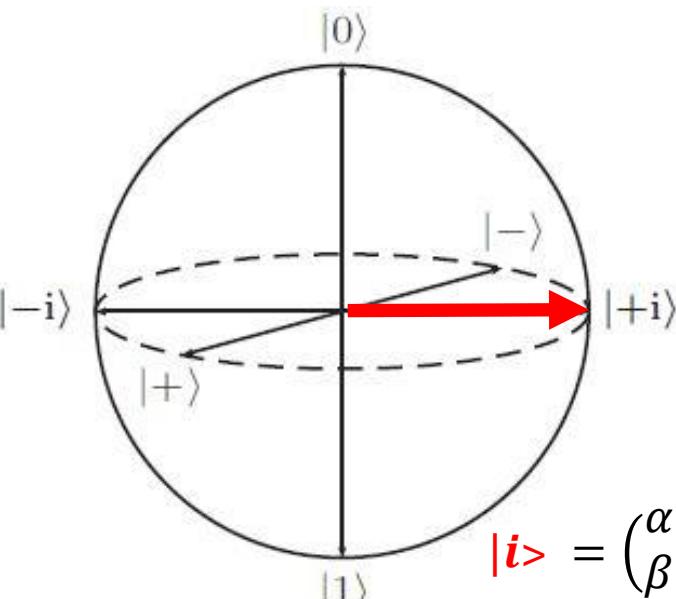
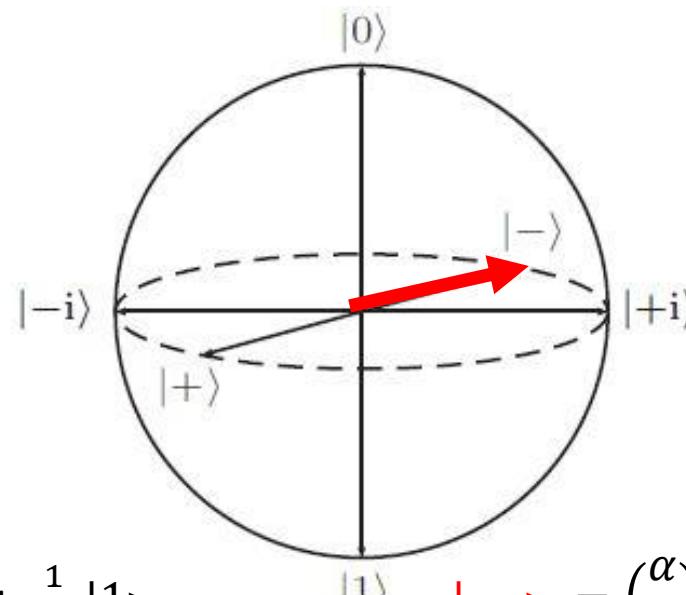
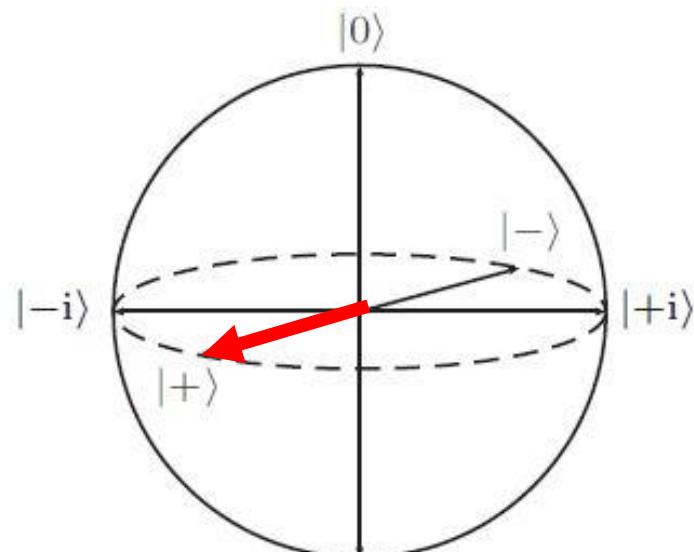


$|1\rangle$



$$|\psi\rangle = \begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{pmatrix}$$

Bloch Sphere



$$|+\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} 1 \\ \frac{1}{\sqrt{2}} \end{bmatrix} = \frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$$

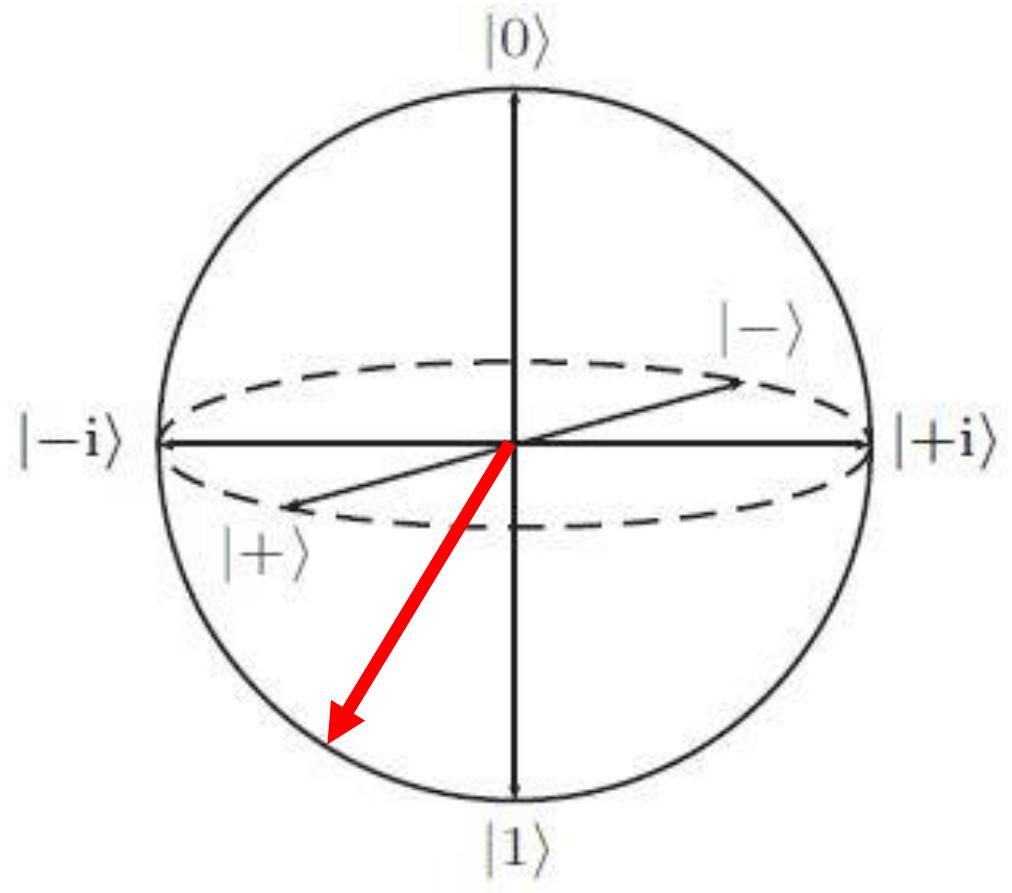
$$|-\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} 1 \\ \frac{1}{\sqrt{2}} \end{bmatrix} = \frac{1}{\sqrt{2}}|0\rangle - \frac{1}{\sqrt{2}}|1\rangle$$

$$|i\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} \end{bmatrix} = \frac{1}{\sqrt{2}}|0\rangle + \frac{i}{\sqrt{2}}|1\rangle$$

$$|-i\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{bmatrix} = \frac{1}{\sqrt{2}}|0\rangle - \frac{i}{\sqrt{2}}|1\rangle$$

Bloch Sphere

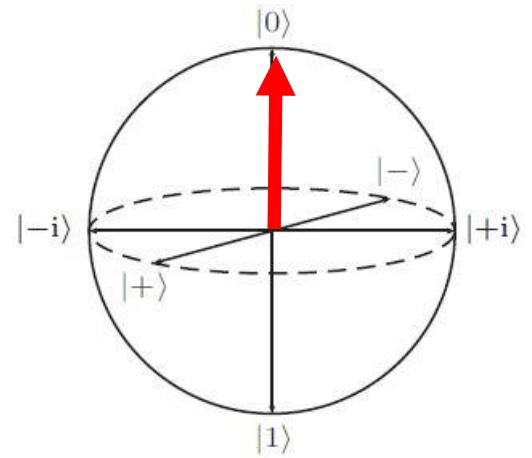
Ex: Locate the state $|\psi\rangle = \begin{pmatrix} \frac{1}{2} \\ \frac{2}{2} \\ \frac{\sqrt{3}}{2} \end{pmatrix}$ on the Bloch sphere



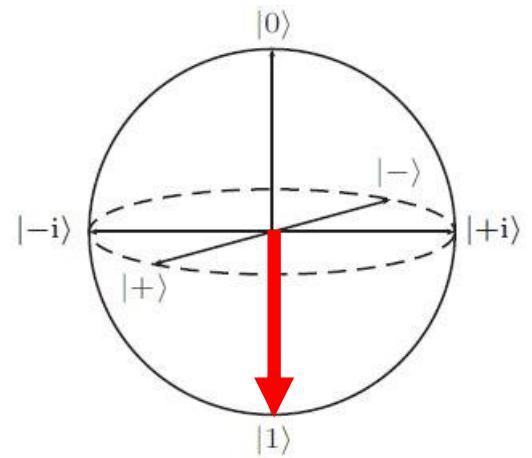
1/4 (25%) Chance of being measured as 0
3/4 (75%) Chance of being measured as 1

Gates

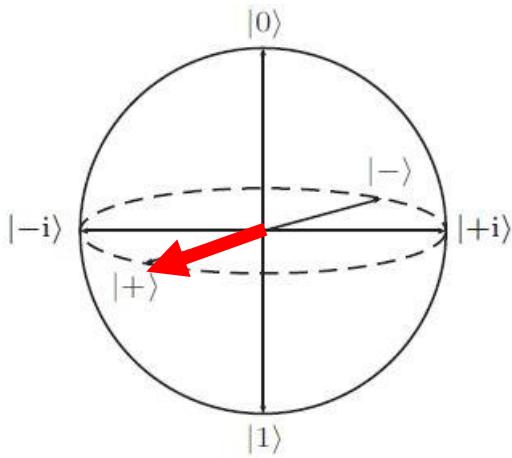




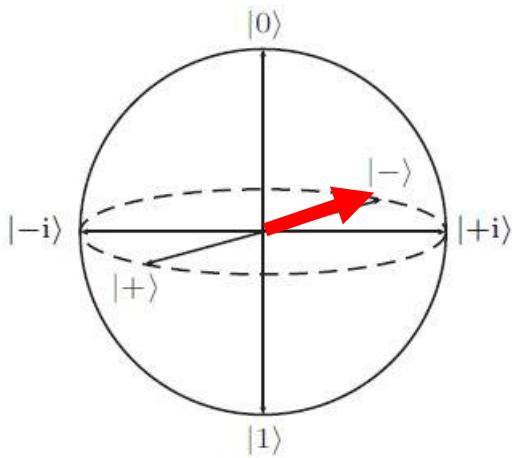
X Gate



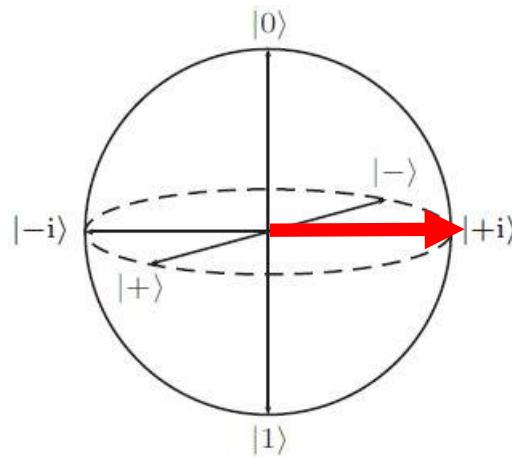
$\text{X } |0\rangle = |1\rangle$



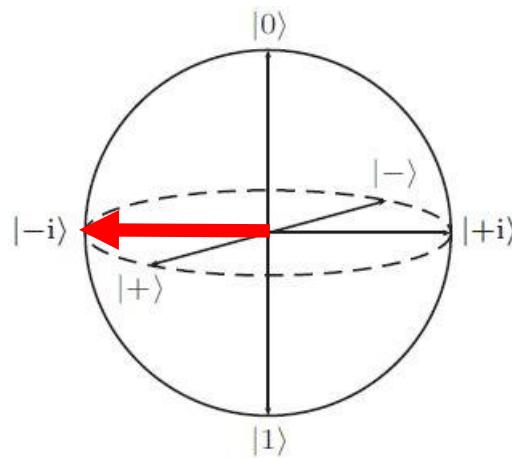
X Gate



$\text{X } |+\rangle = |-\rangle$



X Gate



$\text{X } |i\rangle = |-i\rangle$

X gate flips the qubit 180 (π) radians around the X axis on Bloch Sphere

Y/Z gate flips the qubit 180 (π) radians around the Y/Z axis (respectively) on Bloch Sphere



Gates: X, Y, Z

Y gate flips the qubit 180 (π) radians around the Y axis on Bloch Sphere

Z gate flips the qubit 180 (π) radians around the Z axis on Bloch Sphere

Applying the same gate two times would result into the original state.

This means, X, Y and Z gates are **reversible** or inverse of its own.

$$X = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \quad Y = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} \quad Z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$$



Gates: X, Y, Z

$$X |\psi\rangle = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} \beta \\ \alpha \end{pmatrix}$$

Ex:Prove that applying X gate to $|0\rangle$ results into $|1\rangle$

$$X |0\rangle = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

Ex:Prove that applying X gate to $|1\rangle$ results into $|0\rangle$

$$X |1\rangle = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$



Arbitrary Gate

Let $U = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$ be an arbitrary gate

$$U|0\rangle = \begin{pmatrix} a \\ c \end{pmatrix} \quad U|0\rangle = a|0\rangle + b|1\rangle$$

$$U|1\rangle = \begin{pmatrix} b \\ d \end{pmatrix} \quad U|1\rangle = b|0\rangle + d|1\rangle$$

$$\text{Let } |\psi\rangle = \alpha|0\rangle + \beta|1\rangle \quad U|\psi\rangle = U(\alpha|0\rangle + \beta|1\rangle)$$

$$U|\psi\rangle = \alpha U|0\rangle + \beta U|1\rangle$$



Example:

Let us apply Y gate $\begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix}$ to $|\psi\rangle = \frac{\sqrt{3}}{2}|0\rangle + \frac{1}{2}|1\rangle$

$$Y|\psi\rangle = Y\left(\frac{\sqrt{3}}{2}|0\rangle + \frac{1}{2}|1\rangle\right)$$

$$Y|\psi\rangle = \frac{\sqrt{3}}{2}Y|0\rangle + \frac{1}{2}Y|1\rangle$$

$$Y|\psi\rangle = \frac{\sqrt{3}}{2}\binom{0}{i} + \frac{1}{2}\binom{-i}{0}$$

$$Y|\psi\rangle = \frac{\sqrt{3}}{2}i\binom{0}{1} - \frac{1}{2}i\binom{1}{0}$$

$$Y|\psi\rangle = \frac{\sqrt{3}}{2}i\binom{0}{1} - \frac{1}{2}i\binom{1}{0}$$

$$Y|\psi\rangle = \frac{\sqrt{3}}{2}i|1\rangle - \frac{1}{2}i|0\rangle$$

Ex: Let us apply Z gate $\begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$ to $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$

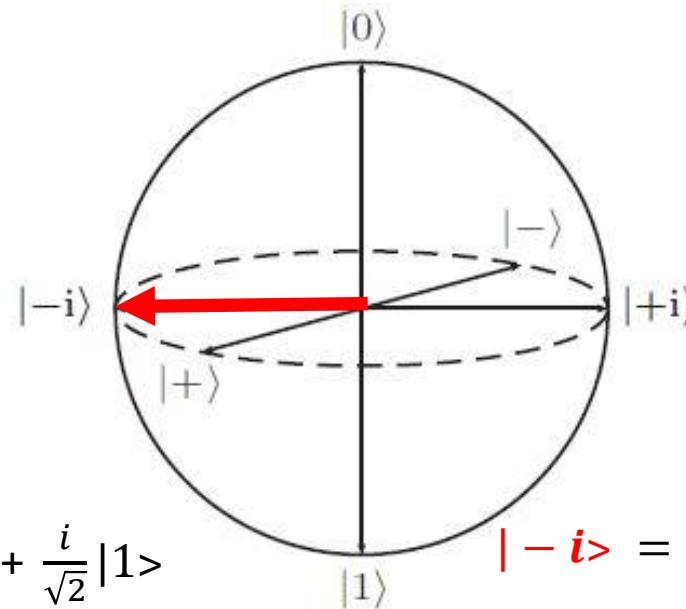
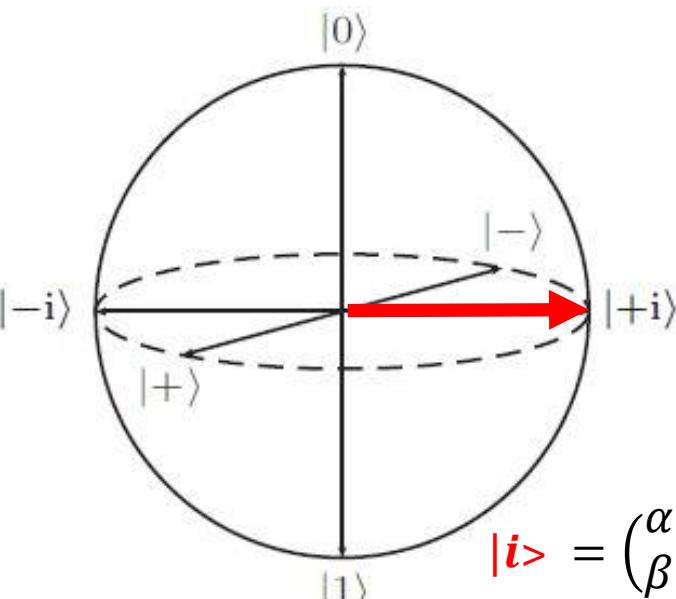
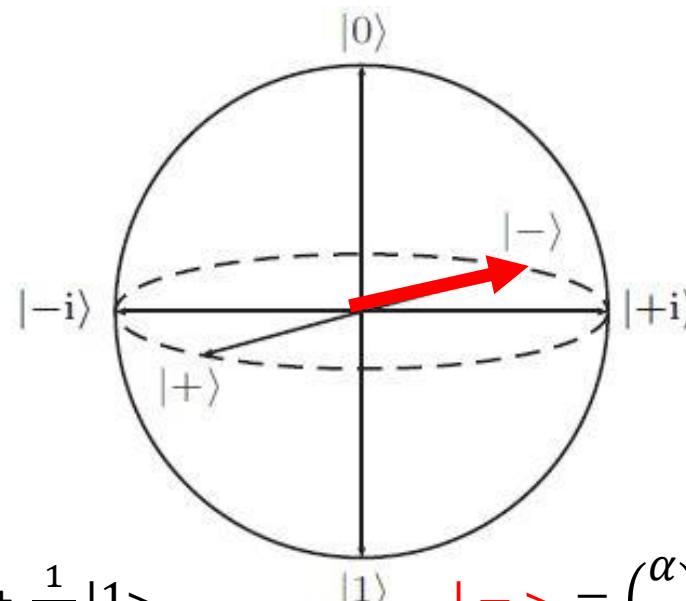
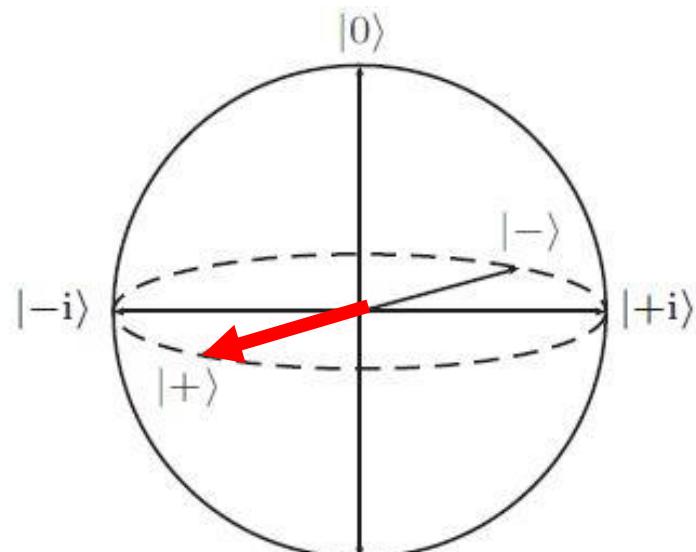
$$\text{ANS} = \alpha|0\rangle - \beta|1\rangle$$



Hadamard (H) Gate



Bloch Sphere



$$|+\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} 1 \\ \frac{1}{\sqrt{2}} \end{bmatrix} = \frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$$

$$|-\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} 1 \\ \frac{1}{\sqrt{2}} \end{bmatrix} = \frac{1}{\sqrt{2}}|0\rangle - \frac{1}{\sqrt{2}}|1\rangle$$

$$|i\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} \end{bmatrix} = \frac{1}{\sqrt{2}}|0\rangle + \frac{i}{\sqrt{2}}|1\rangle$$

$$|-i\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{bmatrix} = \frac{1}{\sqrt{2}}|0\rangle - \frac{i}{\sqrt{2}}|1\rangle$$

Hadamard (H) Gate:

$$H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

$$|+\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} 1 \\ \sqrt{2} \\ 1 \\ \sqrt{2} \end{bmatrix} = \frac{1}{\sqrt{2}} |0\rangle + \frac{1}{\sqrt{2}} |1\rangle$$

$$|i\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} 1 \\ \sqrt{2} \\ 1 \\ \sqrt{2} \end{bmatrix} = \frac{1}{\sqrt{2}} |0\rangle + \frac{i}{\sqrt{2}} |1\rangle$$

$$|-> = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} 1 \\ \sqrt{2} \\ 1 \\ \sqrt{2} \end{bmatrix} = \frac{1}{\sqrt{2}} |0\rangle - \frac{1}{\sqrt{2}} |1\rangle$$

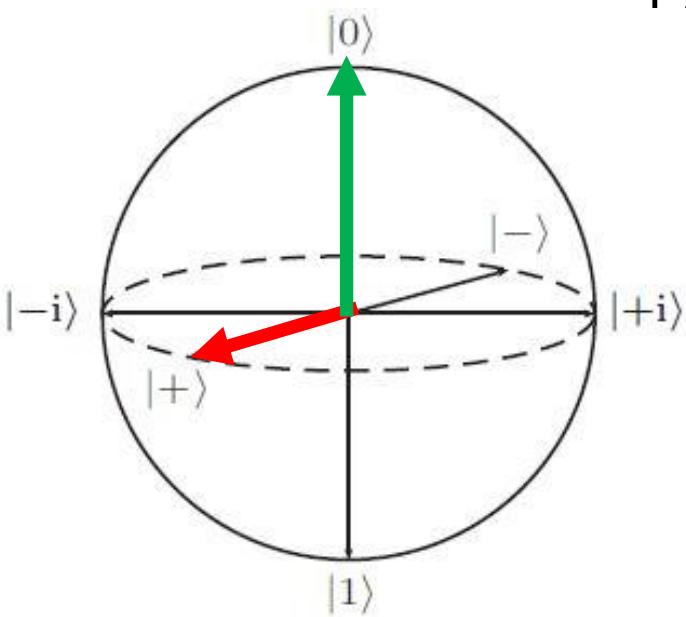
$$|-i\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} 1 \\ \sqrt{2} \\ 1 \\ \sqrt{2} \end{bmatrix} = \frac{1}{\sqrt{2}} |0\rangle - \frac{i}{\sqrt{2}} |1\rangle$$

$$H|0\rangle \rightarrow |+\rangle$$

$$H|+\rangle \rightarrow |0\rangle$$

Hadamard is inverse of its own (Reversible)

$$\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{bmatrix} = \frac{1}{\sqrt{2}} |0\rangle + \frac{1}{\sqrt{2}} |1\rangle = |+\rangle$$



$$H|1\rangle \rightarrow |-\rangle$$

$$H|-\rangle \rightarrow |1\rangle$$

S & T Gate

$$S = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\frac{\pi}{2}} \end{bmatrix} \quad T = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\frac{\pi}{4}} \end{bmatrix}$$



Multiple Qubits

Tensor Product \otimes

Two Qubits in 0 states are represented using $|0\rangle \otimes |0\rangle \Rightarrow |00\rangle$

Two Qubits in superposition

$$(\alpha|0\rangle + \beta|1\rangle) \otimes (\gamma|0\rangle + \delta|1\rangle) = \alpha\gamma|00\rangle + \alpha\delta|01\rangle + \beta\gamma|10\rangle + \beta\delta|11\rangle$$

Probability of measuring $|00\rangle$ is $|\alpha\gamma|^2$

Ex: Combine: $(\frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle) \otimes (\frac{\sqrt{3}}{2}|0\rangle + \frac{1}{2}|1\rangle)$



Multiple Qubits

Ex: Combine: $(\frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle) \otimes (\frac{\sqrt{3}}{2}|0\rangle + \frac{1}{2}|1\rangle)$

$$\frac{1}{\sqrt{2}} \frac{\sqrt{3}}{2} |00\rangle + \frac{1}{\sqrt{2}} \frac{1}{2} |01\rangle + \frac{1}{\sqrt{2}} \frac{\sqrt{3}}{2} |10\rangle + \frac{1}{\sqrt{2}} \frac{1}{2} |11\rangle$$

$$\frac{\sqrt{3}}{2\sqrt{2}} |00\rangle + \frac{1}{2\sqrt{2}} |01\rangle + \frac{\sqrt{3}}{2\sqrt{2}} |10\rangle + \frac{1}{2\sqrt{2}} |11\rangle$$

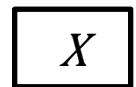
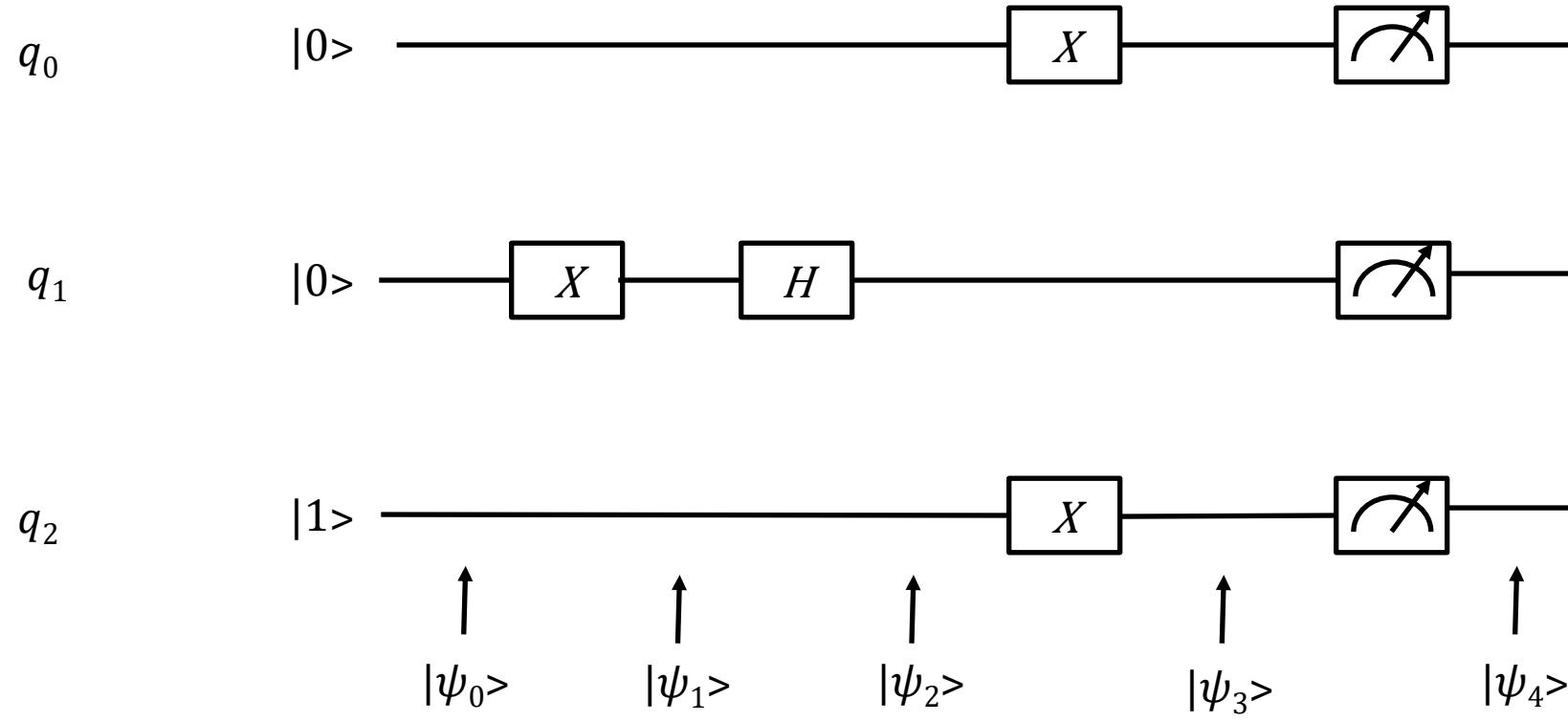
Probability of measuring $|00\rangle$ is $\left|\frac{\sqrt{3}}{2\sqrt{2}}\right|^2 = \frac{3}{8} = 37.5\%$

Probability of measuring $|01\rangle$ is $\left|\frac{1}{2\sqrt{2}}\right|^2 = \frac{1}{8} = 12.5\%$

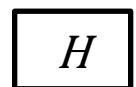
Probability of measuring $|10\rangle$ is $\left|\frac{\sqrt{3}}{2\sqrt{2}}\right|^2 = \frac{3}{8} = 37.5\%$

Probability of measuring $|11\rangle$ is $\left|\frac{1}{2\sqrt{2}}\right|^2 = \frac{1}{8} = 12.5\%$

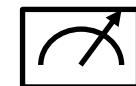




X Gate



Hadamard Gate



Measurement

Quantum Circuits

$$|\psi_0\rangle = |001\rangle$$

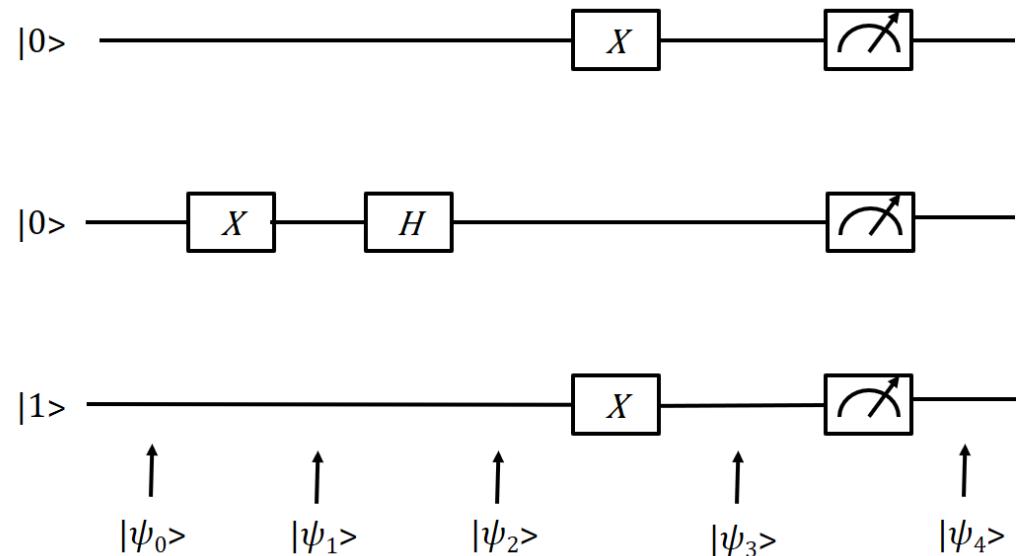
$$|\psi_1\rangle = |0\textcolor{red}{1}1\rangle$$

$$|\psi_2\rangle = |0 - 1\rangle$$

$$|\psi_2\rangle = |0 \otimes \left(\frac{1}{\sqrt{2}}|0\rangle - \frac{1}{\sqrt{2}}|1\rangle \right) \otimes 1\rangle$$

$$|\psi_2\rangle = \frac{1}{\sqrt{2}}(|001\rangle - \frac{1}{\sqrt{2}}|011\rangle)$$

$$|\psi_3\rangle = \frac{1}{\sqrt{2}}(|\textcolor{green}{1}00\rangle - \frac{1}{\sqrt{2}}|\textcolor{green}{1}10\rangle)$$

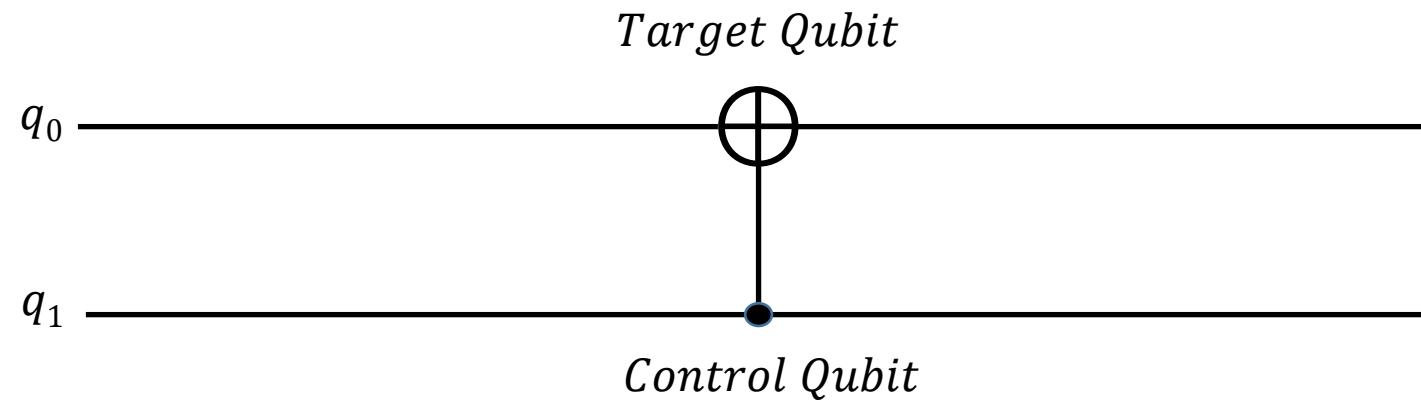


$$\begin{aligned} |\psi_4\rangle &= \\ &|\textcolor{green}{1}00\rangle \text{ (50% Probability)} \\ &|\textcolor{green}{1}10\rangle \text{ (50% Probability)} \end{aligned}$$

Multi-Qubit Gates: CNOT, Toffoli, Controlled Gates



CNOT: Controlled NOT Gate



If the Control Qubit is 1, the Target bit is flipped.

CNOT: Controlled NOT Gate

Ex: Apply CNOT to $\left(\frac{\sqrt{3}}{2} |00\rangle + \frac{1}{2} |01\rangle + \frac{1}{\sqrt{2}} |10\rangle + \frac{1}{4} |11\rangle \right)$,

assuming *first bit as a control* and *second bit as a target bit*

$$CNOT \left(\frac{\sqrt{3}}{2} |00\rangle + \frac{1}{2} |01\rangle + \frac{1}{\sqrt{2}} |10\rangle + \frac{1}{4} |11\rangle \right)$$

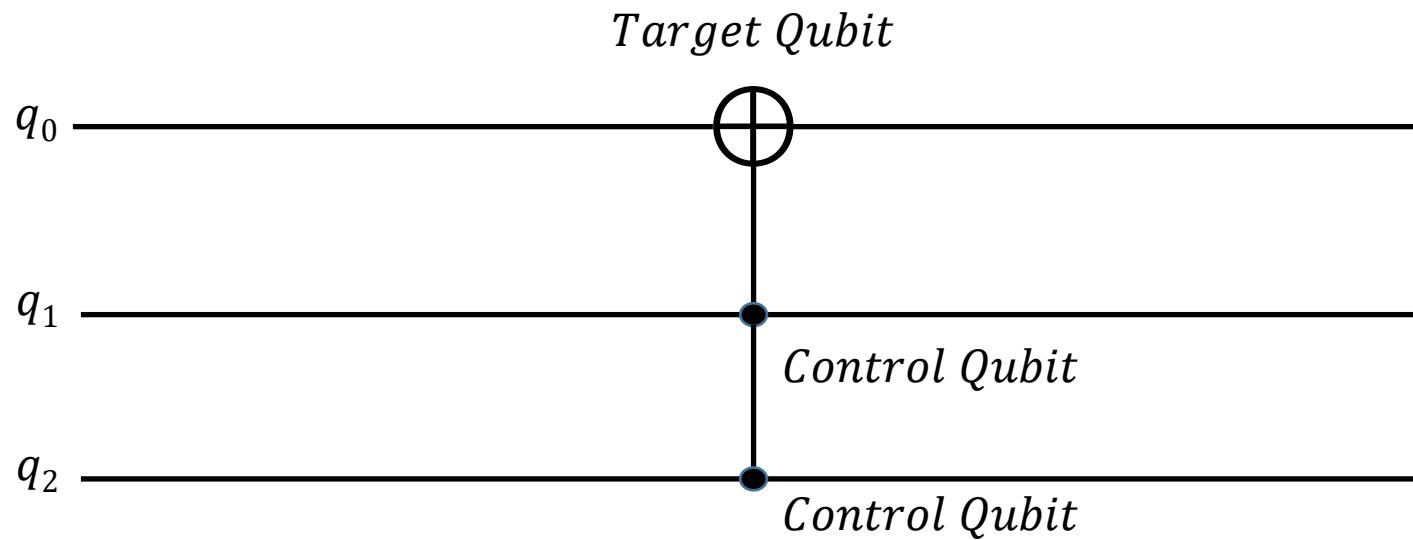
$$\left(\frac{\sqrt{3}}{2} CNOT |00\rangle + \frac{1}{2} CNOT |01\rangle + \frac{1}{\sqrt{2}} CNOT |\textcolor{red}{10}\rangle + \frac{1}{4} CNOT |\textcolor{red}{11}\rangle \right)$$

$$\left(\frac{\sqrt{3}}{2} |00\rangle + \frac{1}{2} |01\rangle + \frac{1}{\sqrt{2}} |\textcolor{red}{11}\rangle + \frac{1}{4} CNOT |\textcolor{red}{10}\rangle \right)$$

Ex: Apply CNOT to $\left(\frac{\sqrt{3}}{2} |001\rangle + \frac{1}{2} |010\rangle \right)$, assuming *third bit as a control* and *second bit as a target bit*



Toffoli Gate: Controlled NOT Gate with 2 control qubits



If all the Control Qubit are 1, the Target bit is flipped.

Toffoli Gate: Controlled NOT Gate with 2 control qubits

Ex: Apply Toffoli to $\left(\frac{1}{\sqrt{2}} \left| 0011 \right\rangle + \frac{1}{\sqrt{2}} \left| 0110 \right\rangle \right)$,

assuming second and third bits as control and fourth bit as a target bit

$$Toffoli \left(\frac{1}{\sqrt{2}} \left| 0011 \right\rangle + \frac{1}{\sqrt{2}} \left| 0110 \right\rangle \right)$$

$$\left(\frac{1}{\sqrt{2}} Toffoli \left| 0011 \right\rangle + \frac{1}{\sqrt{2}} Toffoli \left| 0110 \right\rangle \right)$$

$$\left(\frac{1}{\sqrt{2}} \left| 0011 \right\rangle + \frac{1}{\sqrt{2}} Toffoli \left| 0111 \right\rangle \right)$$



Measuring a qubit

Find out the probability of measuring second qubit as 1 in the state $(\frac{1}{2}|00\rangle + \frac{1}{4}|01\rangle + \frac{e^{i\frac{\pi}{2}}}{\sqrt{2}}|10\rangle + \frac{\sqrt{3}}{4}|11\rangle)$

$(\frac{1}{2}|00\rangle + \frac{1}{4}|01\rangle + \frac{e^{i\frac{\pi}{2}}}{\sqrt{2}}|10\rangle + \frac{\sqrt{3}}{4}|11\rangle)$ There are two super-positions where the second qubit is 1

Let us sum up their probability: $|\frac{1}{4}|^2 + |\frac{\sqrt{3}}{4}|^2$

$$|\frac{1}{16}| + |\frac{3}{16}|$$

$$|\frac{4}{16}|$$

$$|\frac{1}{4}|$$

25%



Measuring a qubit

How do we know what the state collapses to, once we measure a qubit?

Once we measure a qubit, it collapses from superposition and becomes a measurement.

Example: Let us assume that in the given state, we measure the first qubit to be 1.

$$\left(\frac{1}{2}|00\rangle + \frac{1}{4}|01\rangle + \frac{1}{\sqrt{2}}|10\rangle + \frac{\sqrt{3}}{4}|11\rangle\right)$$

As we know, the first qubit collapses to 1, and hence, we may ignore the state where the first qubit is zero. $\frac{1}{\sqrt{2}}|10\rangle + \frac{\sqrt{3}}{4}|11\rangle$

Now to measure the total probability, let us add the remaining probability $|\frac{1}{2}| + |\frac{3}{16}|$

Since, the addition of all the probability should be 1, here it is 11/16, which is not 1. This is due to the fact that we ignored few states.

Here comes a normalization constant A, to bring the total probability to 1.

$$A \left(\frac{1}{\sqrt{2}}|10\rangle + \frac{\sqrt{3}}{4}|11\rangle \right) \quad \left(\frac{A}{\sqrt{2}} \right)^2 + \left(\frac{A\sqrt{3}}{4} \right)^2 = 1 \quad A = 4/11 \quad \frac{4}{11\sqrt{2}} \left(\frac{1}{\sqrt{2}}|10\rangle + \frac{\sqrt{3}}{4}|11\rangle \right) \quad \frac{4}{\sqrt{22}}|10\rangle + \frac{\sqrt{3}}{\sqrt{11}}|11\rangle$$



Measuring a qubit

Example: Let us assume that in the given state, we measure the middle qubit to be 0.

$$|\psi_0\rangle = \left(\frac{1}{2}|000\rangle + \frac{1}{2}|001\rangle + \frac{1}{2}|010\rangle + \frac{1}{2}|101\rangle\right)$$

$$|\psi_0\rangle = \left(\frac{1}{2}|000\rangle + \frac{1}{2}|001\rangle + \frac{1}{2}|010\rangle + \frac{1}{2}|101\rangle\right)$$

$$|\psi_0\rangle = A \left(\frac{1}{2}|000\rangle + \frac{1}{2}|001\rangle + \frac{1}{2}|101\rangle\right)$$

$$\left|\frac{A}{2}\right|^2 + \left|\frac{A}{2}\right|^2 + \left|\frac{A}{2}\right|^2 = 1$$

$$\frac{3A^2}{4} = 1$$

$$A = \frac{2}{\sqrt{3}}$$

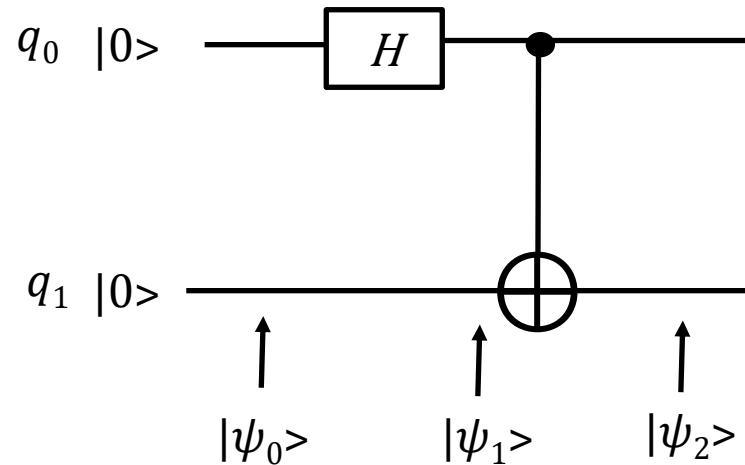
$$|\psi_0\rangle = \frac{2}{\sqrt{3}} \left(\frac{1}{2}|000\rangle + \frac{1}{2}|001\rangle + \frac{1}{2}|101\rangle\right)$$

$$|\psi_0\rangle = \frac{1}{\sqrt{3}}|000\rangle + \frac{1}{\sqrt{3}}|001\rangle + \frac{1}{\sqrt{3}}|101\rangle$$



Entanglement and Bell States





$$|\psi_0\rangle = |00\rangle$$

$$|\psi_1\rangle = H|0\rangle|0\rangle$$

$$|\psi_1\rangle = | + 0 \rangle$$

$$|\psi_1\rangle = \left(\frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle\right) \otimes |0\rangle$$

$$|\psi_1\rangle = \left(\frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|10\rangle\right)$$

Entanglement

$$|\psi_2\rangle = \left(\frac{1}{\sqrt{2}}CNOT|00\rangle + \frac{1}{\sqrt{2}}CNOT|10\rangle\right)$$

$$|\psi_2\rangle = \left(\frac{1}{\sqrt{2}}|00\rangle + \frac{1}{\sqrt{2}}|11\rangle\right)$$

Conclusion: If we measure one of the qubit to 0, then other would collapse into 0.

And If we measure one of the qubit to 1, then other would collapse into 1.

Hence, without at looking both of the qubits, merely, by looking at one of the qubits, we can say the state of the other qubit.

Entangle States

A state is entangled if it can not be factored

A state is entangled if it can not be factored into tensor products of individual qubits

For instance, the state $(\frac{\sqrt{3}}{2\sqrt{5}}|00\rangle + \frac{1}{2\sqrt{5}}|01\rangle + \frac{\sqrt{3}}{\sqrt{5}}|10\rangle + \frac{1}{\sqrt{5}}|11\rangle)$ is not in entangled state as, it can be factored into

$$(\frac{1}{\sqrt{5}}|0\rangle + \frac{2}{\sqrt{5}}|1\rangle) \otimes (\frac{\sqrt{3}}{2}|0\rangle + \frac{1}{2}|1\rangle)$$

But, the state $\frac{1}{\sqrt{2}}(|000\rangle + |011\rangle)$ is in entangled.



Entangled States

Maximally Entangled State

We can certainly determine state of one qubit from other

$$\text{Example: } |\psi_1\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$$

$$|\phi^+\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$$

$$|\phi^-\rangle = \frac{1}{\sqrt{2}}(|00\rangle - |11\rangle)$$

$$|\Psi^+\rangle = \frac{1}{\sqrt{2}}(|01\rangle + |10\rangle)$$

$$|\Psi^-\rangle = \frac{1}{\sqrt{2}}(|01\rangle - |10\rangle)$$

Bell States

Partially Entangled State

By measuring one of the qubits, the amplitude of other qubit is affected

$$\text{Example: } (\frac{\sqrt{3}}{\sqrt{5}}|00\rangle + \frac{1}{\sqrt{5}}|01\rangle + \frac{1}{2\sqrt{5}}|10\rangle + \frac{\sqrt{3}}{2\sqrt{5}}|11\rangle)$$

If we measure the first qubit as 0, the state collapses to:

$$|0\rangle \otimes \boxed{(\frac{\sqrt{3}}{2}|0\rangle + \frac{1}{2}|1\rangle)} \quad \text{Second qubit}$$

If we measure the first qubit as 1, the state collapses to:

$$|1\rangle \otimes \boxed{(\frac{1}{2}|0\rangle + \frac{\sqrt{3}}{2}|1\rangle)} \quad \text{Second qubit}$$



Super Dense Coding



Super Dense Coding

This quantum protocol allows to send two bits of classical information (00, 01, 10, 11) using 1 qubit.

Alice (sender) and Bob (receiver) maximally entangle two qubits.

$$|\psi\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$$

Alice wants to send 00, she does **nothing**: Qubit $|\psi\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$

Alice wants to send 01, she applies **X gate** to her qubit $|\psi\rangle = \frac{1}{\sqrt{2}}(|10\rangle + |01\rangle)$

Alice wants to send 10, she applies **Z gate** to her qubit $|\psi\rangle = \frac{1}{\sqrt{2}}(|00\rangle - |11\rangle)$

Alice wants to send 11, she applies both **X & Z gate** to her qubit $|\psi\rangle = \frac{1}{\sqrt{2}}(|10\rangle - |01\rangle)$

Now, Bob has both the qubit in one of these four states,

$$|\psi\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle), |\psi\rangle = \frac{1}{\sqrt{2}}(|10\rangle + |01\rangle), |\psi\rangle = \frac{1}{\sqrt{2}}(|00\rangle - |11\rangle), |\psi\rangle = \frac{1}{\sqrt{2}}(|10\rangle - |01\rangle)$$



Super Dense Coding

Bob, now applies CNOT as first qubit as control and second qubit as target

Bob, then applies H gate on the left qubit.

$$\begin{array}{lll} |\psi\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle) & \xrightarrow{\text{CNOT}} & |\psi\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |\textcolor{green}{1}\textcolor{red}{0}\rangle) & \xrightarrow{H} & H|+\rangle|0\rangle = |00\rangle \\ |\psi\rangle = \frac{1}{\sqrt{2}}(|10\rangle + |01\rangle) & \xrightarrow{\text{CNOT}} & |\psi\rangle = \frac{1}{\sqrt{2}}(|\textcolor{green}{1}\textcolor{red}{1}\rangle + |01\rangle) & \xrightarrow{H} & H|+\rangle|1\rangle = |01\rangle \\ |\psi\rangle = \frac{1}{\sqrt{2}}(|00\rangle - |11\rangle) & \xrightarrow{\text{CNOT}} & |\psi\rangle = \frac{1}{\sqrt{2}}(|00\rangle - |\textcolor{green}{1}\textcolor{red}{0}\rangle) & \xrightarrow{H} & H|-\rangle|0\rangle = |10\rangle \\ |\psi\rangle = \frac{1}{\sqrt{2}}(|10\rangle - |01\rangle) & \xrightarrow{\text{CNOT}} & |\psi\rangle = \frac{1}{\sqrt{2}}(|\textcolor{green}{1}\textcolor{red}{1}\rangle - |01\rangle) & \xrightarrow{H} & H|-\rangle|1\rangle = |11\rangle \end{array}$$

Bob, now measures the qubits and he'll have $|00\rangle$ or $|01\rangle$ or $|10\rangle$ or $|11\rangle$



Reversible Function



Reversible Function

NOT		AND		OR		XOR	
Input=x	Output=f(x)	Input=x	Output=f(x)	Input=x	Output=f(x)	Input=x	Output=f(x)
0	1	00	0	00	0	00	0
1	0	01	0	01	1	01	1
		10	0	10	1	10	1
		11	1	11	1	11	0

A function f is **reversible**, if given $f(x)$, we can find x

Which of the above gate is reversible?

Check whether the following function is reversible?

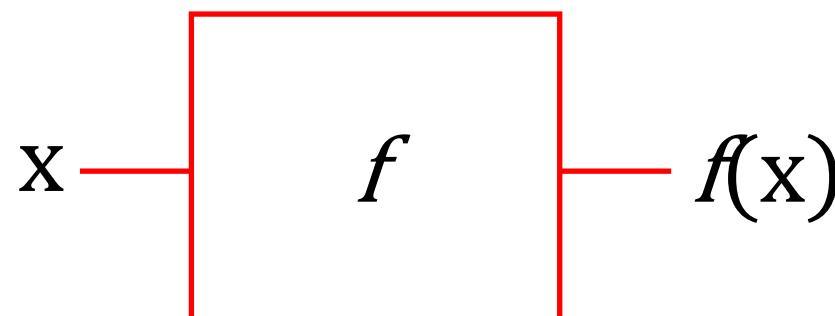
Function F	
Input=x	Output=f(x)
00	01
01	00
10	11
11	10



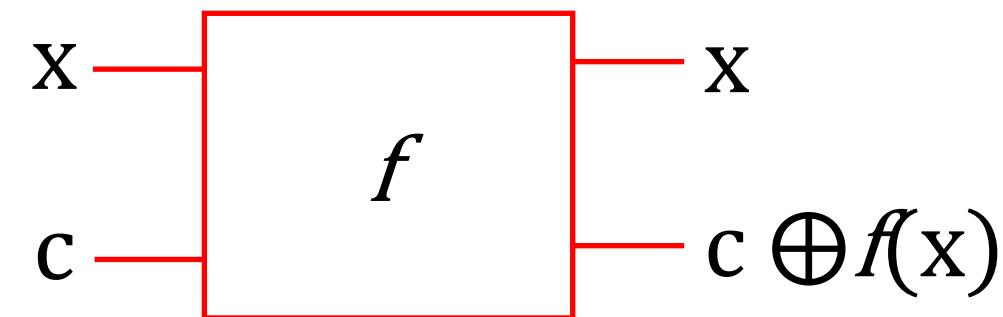
Reversible Function

A function f is **reversible**, if given $f(x)$, we can find x

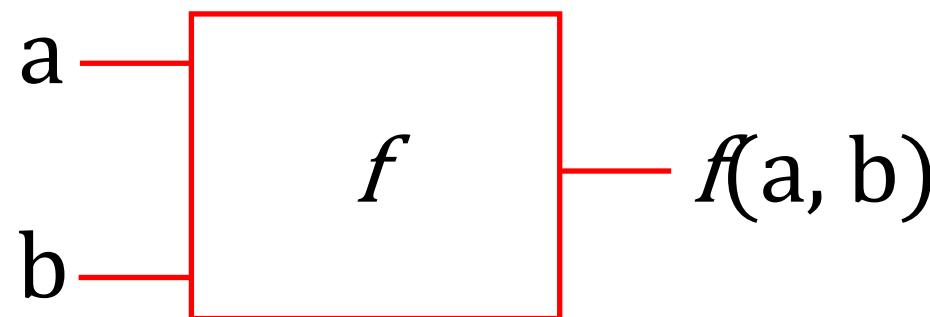
To make a function f is **reversible**, we need add extra information



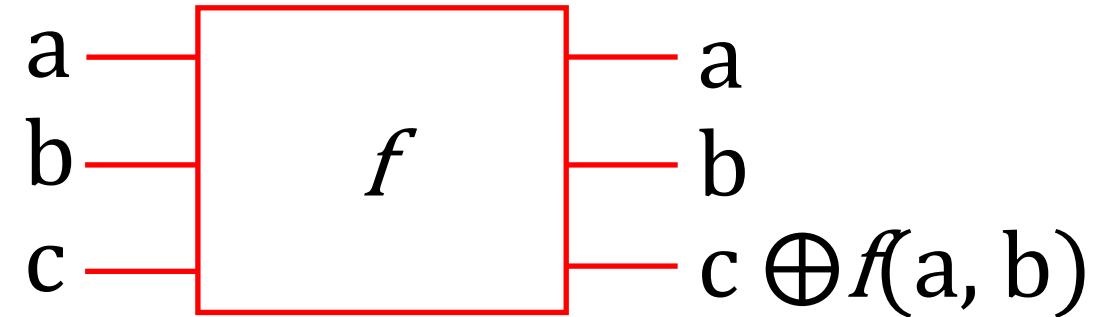
Standard Classical Operations



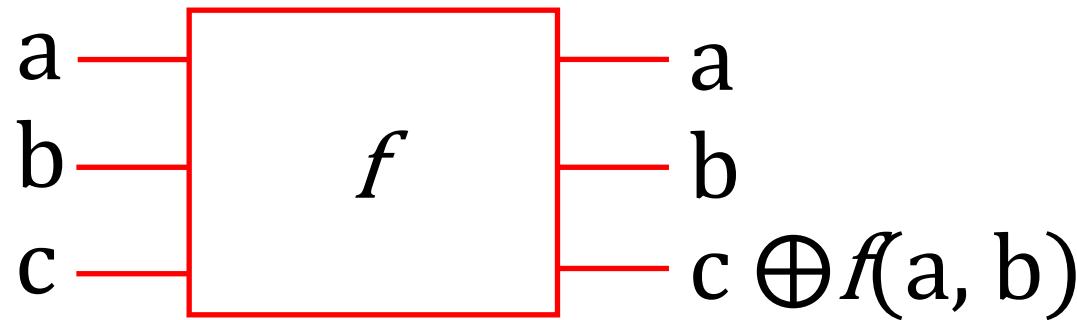
Reversible Classical Operations



Standard OR Gate



Reversible OR Gate



Reversible OR Gate

a	b	c	a	b	$c \oplus f(a, b)$
0	0	0	0	0	0
0	0	1	0	0	1
0	1	0	0	1	1
0	1	1	0	1	0
1	0	0	1	0	1
1	0	1	1	0	0
1	1	0	1	1	1
1	1	1	1	1	0

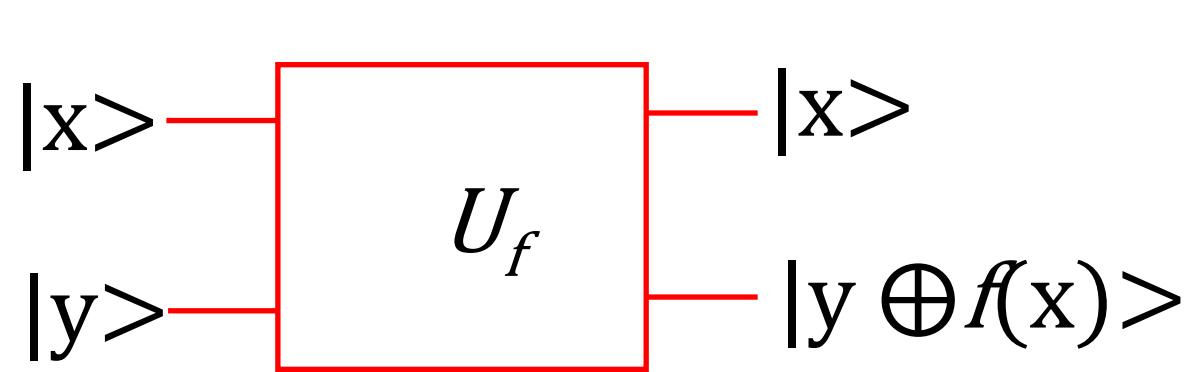
Why all these?

In Quantum Computing, all operations must be **reversible**.

Reversible Function

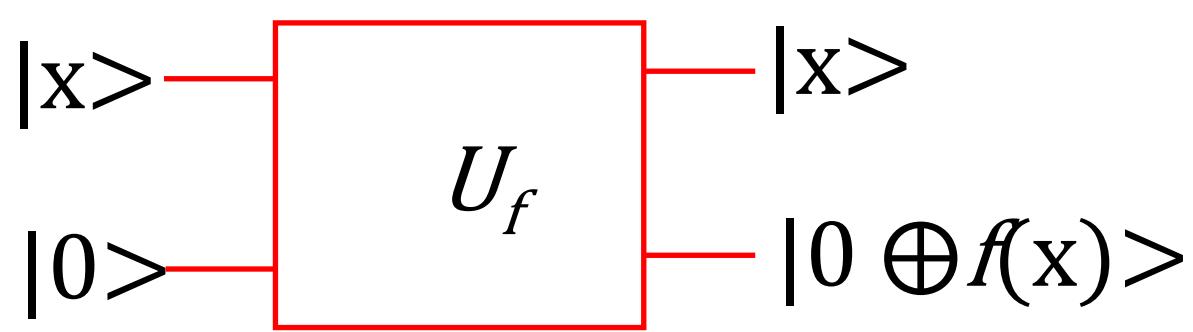
Functions on Quantum Computers / Phase Oracle





Standard Quantum Function / Oracle

$$U_f |x\rangle |y\rangle = |x\rangle |y \oplus f(x)\rangle$$



$$U_f |x\rangle |0\rangle = |x\rangle |0 \oplus f(x)\rangle$$

$$U_f |x\rangle |0\rangle = |x\rangle |f(x)\rangle$$

$$U_f |x\rangle |->$$

$$\text{if } f(x) = 0, \frac{1}{\sqrt{2}}(|x\rangle|0\rangle - |x\rangle|1\rangle) \Rightarrow |x\rangle|->$$

$$U_f |x\rangle \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

$$\text{if } f(x) = 1, \frac{1}{\sqrt{2}}(|x\rangle|1\rangle - |x\rangle|0\rangle)$$

$$U_f \frac{1}{\sqrt{2}}(|x\rangle|0\rangle - |x\rangle|1\rangle)$$

$$\Rightarrow -\frac{1}{\sqrt{2}}(|x\rangle|0\rangle - |x\rangle|1\rangle)$$

$$\frac{1}{\sqrt{2}}(U_f|x\rangle|0\rangle - U_f|x\rangle|1\rangle)$$

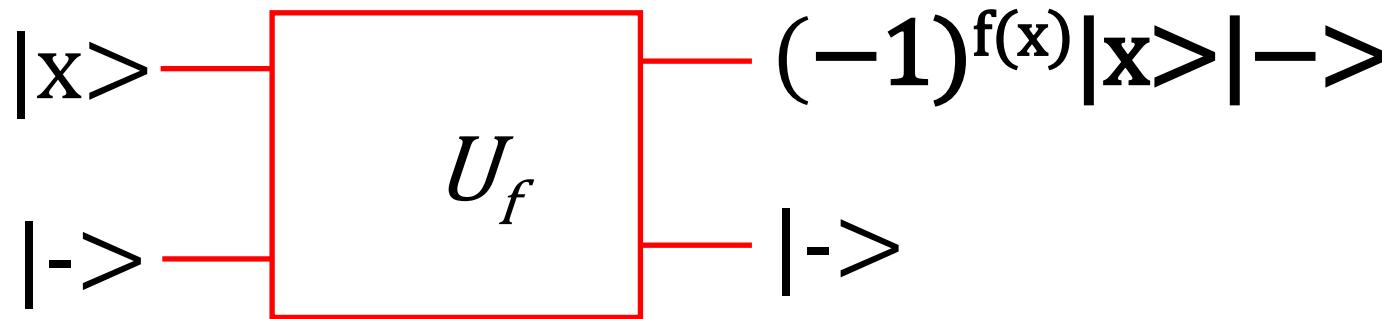
$$\Rightarrow -|x\rangle|->$$

$$\frac{1}{\sqrt{2}}(|x\rangle|f(x)\rangle - |x\rangle|\overline{f(x)}\rangle)$$

$$\text{Combining both, } U_f |x\rangle |-> = (-1)^{f(x)} |x\rangle |->$$

$$U_f |x\rangle |-> = (-1)^{f(x)} |x\rangle |->$$

Phase Oracle



Instead of the function output $f(x)$ being XORed with the $|->$ register a phase of $(-1)^{f(x)}$ was applied to the input to the function ($|x\rangle$)

When we create a function in this way, with the output in this minus state, we call it a *Phase Oracle*

Phase Oracle will be useful in Quantum Algorithms



No-Cloning Theorem

In standard computer, we can simply read a value and copy it to other location/register.

But, in quantum computer, if we have a state, $Let |\psi\rangle = \alpha|0\rangle + \beta|1\rangle$

And if we don't know the value of α and β , then we can not copy the state of to other qubit.



Deutch's Algorithm



Deutch's Algorithm

What problem does the Deutch's Algorithm solve?

Let's say, we have a function f that takes $\{0,1\}$ as input and $\{0,1\}$ as output.

$$f : \{0, 1\} \rightarrow \{0, 1\}$$

Our task is to find out, whether the function f is **balance** or **constant**?

If a function is **constant**, it returns the same output, irrespective of the input: $f(0) = f(1)$

Constant 1	
x	f(x)
0	1
1	1

Constant 0	
x	f(x)
0	0
1	0

Balance function returns 0 for the half of the inputs and 1 for the other half of the inputs.

Balance (NOT)	
X	f(x)
0	1
1	0

Balance (Identity)	
X	f(x)
0	0
1	1



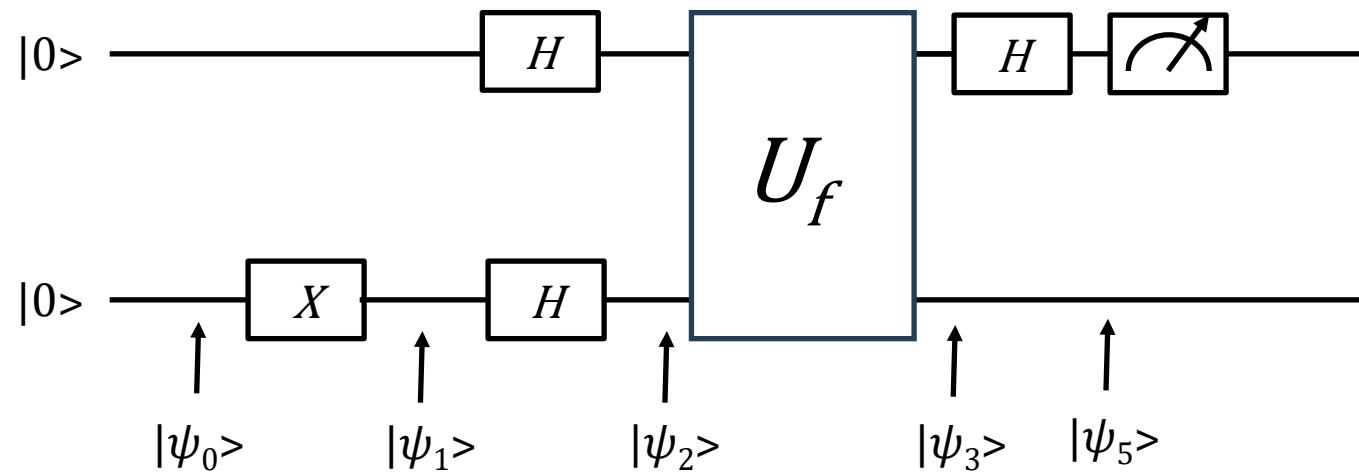
Deutsch's Algorithm

To check whether the function is constant or balance, we need to query the oracle twice. Once with 0 as input and once with 1 as input.

If $f(0) = f(1)$, the function is **constant**

If $f(0) \neq f(1)$, the function is **balance**

With quantum computers, and with Deutsch's algorithm, we only need one query of the function to find out whether the function is **constant** or **balance**!



$$|\psi_0\rangle = |00\rangle$$

$$|\psi_1\rangle = |01\rangle$$

$$|\psi_2\rangle = |+\rangle$$

$$|\psi_2\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) |-\rangle$$

$$|\psi_2\rangle = \frac{1}{\sqrt{2}}(|0\rangle|-\rangle - |1\rangle|-\rangle)$$

$$|\psi_3\rangle = U_f \frac{1}{\sqrt{2}}(|0\rangle|-\rangle - |1\rangle|-\rangle)$$

$$|\psi_3\rangle = \frac{1}{\sqrt{2}}(U_f|0\rangle|-\rangle - U_f|1\rangle|-\rangle)$$

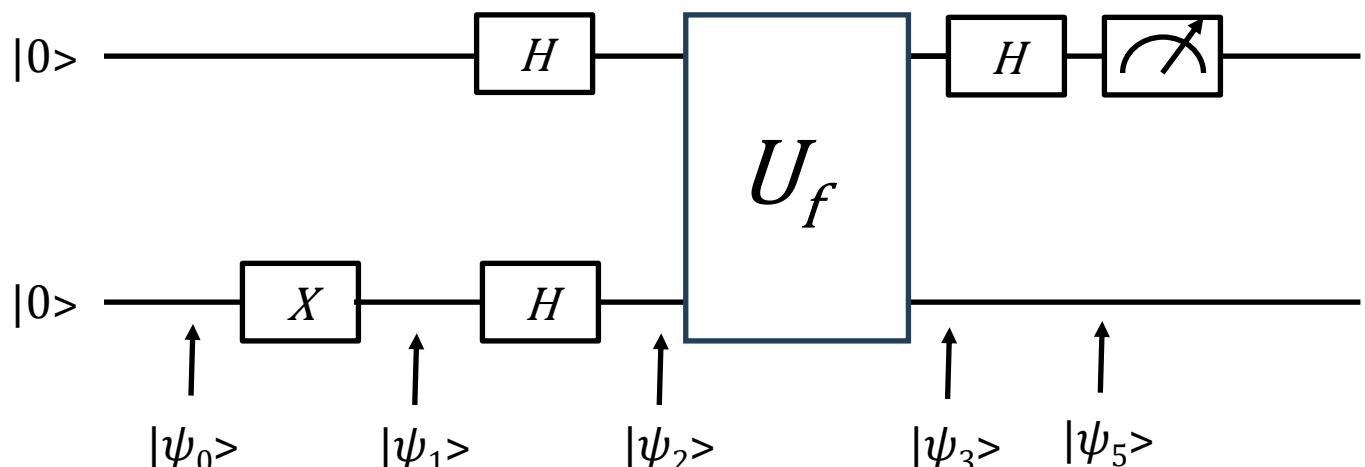
Recall the Phase Oracle, $U_f |x\rangle |-\rangle = (-1)^{f(x)} |x\rangle |-\rangle$

$$|\psi_3\rangle = \frac{1}{\sqrt{2}}((-1)^{f(0)}|0\rangle|-\rangle + (-1)^{f(1)}|1\rangle|-\rangle)$$

We may ignore the $|-\rangle$

$$|\psi_3\rangle = \frac{1}{\sqrt{2}}((-1)^{f(0)}|0\rangle + (-1)^{f(1)}|1\rangle)$$

Deutsch's Algorithm



$$|\psi_3\rangle = \frac{1}{\sqrt{2}}((-1)^{f(0)}|0\rangle + (-1)^{f(1)}|1\rangle)$$

Scenario 1: $f(0) = f(1)$

Scenario 1 (Case 1): $f(0) = f(1) = 0$

$$|\psi_3\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$$

Scenario 1 (Case 2): $f(0) = f(1) = 1$

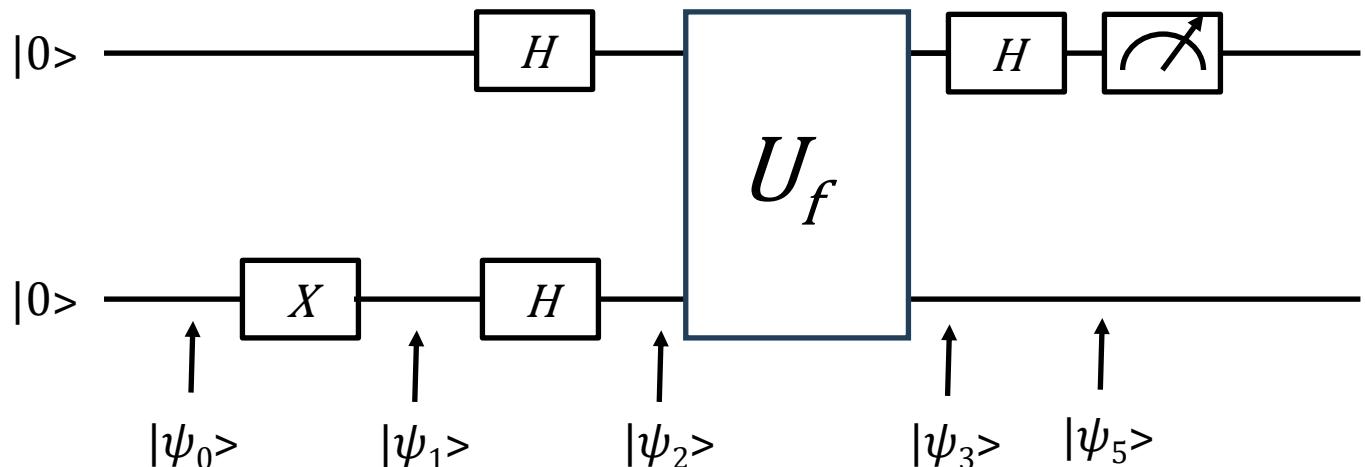
$$|\psi_3\rangle = \frac{1}{\sqrt{2}}(-|0\rangle - |1\rangle)$$

We can factor out the global phase of $|-\rangle$

$$|\psi_3\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$$

That means, if $f(0) = f(1)$, the state becomes $\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$ which is $|+\rangle$

Deutsch's Algorithm



$$|\psi_3\rangle = \frac{1}{\sqrt{2}}((-1)^{f(0)}|0\rangle + (-1)^{f(1)}|1\rangle)$$

Deutsch's Algorithm

Scenario 2: $f(0) \neq f(1)$

Scenario 2 (Case 1): $f(0) = 0, f(1) = 1$

$$|\psi_3\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

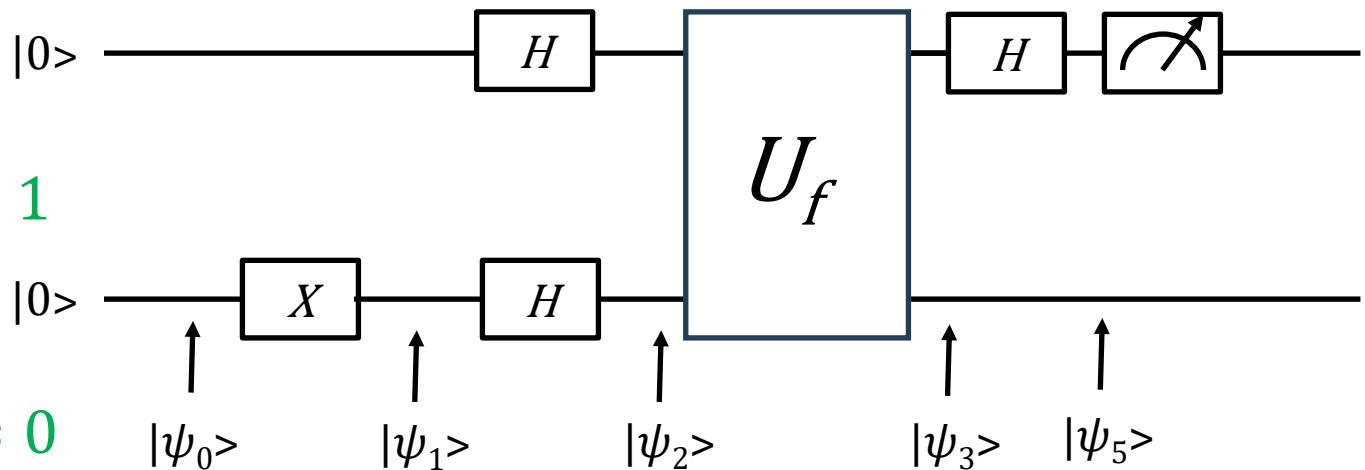
Scenario 2 (Case 2): $f(0) = 1, f(1) = 0$

$$|\psi_3\rangle = \frac{1}{\sqrt{2}}(-|0\rangle + |1\rangle)$$

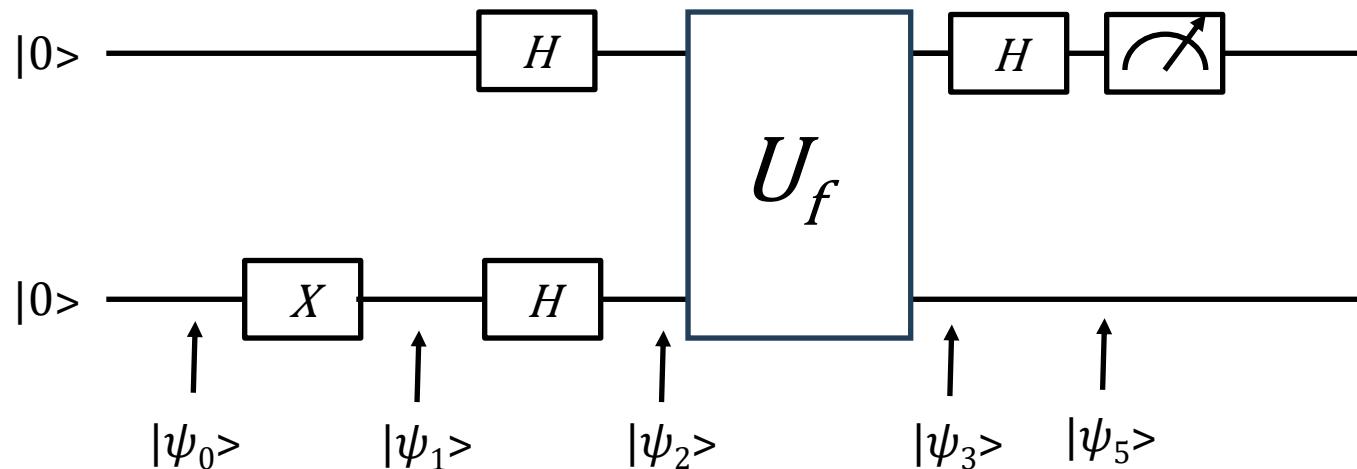
We can factor out the global phase of $|-\rangle$

$$|\psi_3\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

That means, if $f(0) \neq f(1)$, the state becomes $\frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$ which is $|-\rangle$



Deutsch's Algorithm



$$f(0) = f(1)$$

$$|\psi_3\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) = |+\rangle$$

$$|\psi_4\rangle = H|\psi_3\rangle = H|+\rangle = |0\rangle$$

That means, if we measure ($|\psi_4\rangle$) as 0, the function is **constant**

$$f(0) \neq f(1)$$

$$|\psi_3\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) = |->$$

$$|\psi_4\rangle = H|\psi_3\rangle = H|-> = |1\rangle$$

That means, if we measure ($|\psi_4\rangle$) as 0, the function is **balance**.

Deutsch-Jozsa Algorithm



Deutch-Jozsa Algorithm

Deutch-Jozsa algorithm does the same thing that Deutch's Algorithm does, that is to check whether the function is constant or balance

Except, Deutch-Jozsa algorithm works on any number of qubits.

Deutch's Algorithm: $f : \{0, 1\} \rightarrow \{0, 1\}$

Deutch-Jozsa Algorithm: $f : \{0, 1\}^n \rightarrow \{0, 1\}$



Deutsch-Jozsa Algorithm

Constant and Balance Function Revisited:

Constant ZERO	
x	f(x)
000	0
001	0
010	0
011	0
100	0
101	0
110	0
111	0

Constant ZERO	
x	f(x)
000	1
001	1
010	1
011	1
100	1
101	1
110	1
111	1

Balance	
x	f(x)
000	0
001	1
010	1
011	0
100	0
101	1
110	1
111	0

To determine whether f is constant or balanced takes a classical computer at worst $2^{n-1} + 1$ queries of f , where n is number of bits.

With a quantum computer, we just need one query, to determine whether the function is constant or balance.



$$|\psi_0\rangle = |0\rangle^{\otimes n} |-\rangle$$

$$|\psi_1\rangle = H^{\otimes n} |0\rangle^{\otimes n} |-\rangle$$

$$H^{\otimes n} |0\rangle^{\otimes n} = H|0\rangle H|0\rangle H|0\rangle \dots n \text{ times}$$

$$H^{\otimes n} |0\rangle^{\otimes n} = |+\rangle |+\rangle |+\rangle \dots n \text{ times}$$

$$= \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \dots n \text{ times}$$

$$= \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \dots n \text{ times}$$

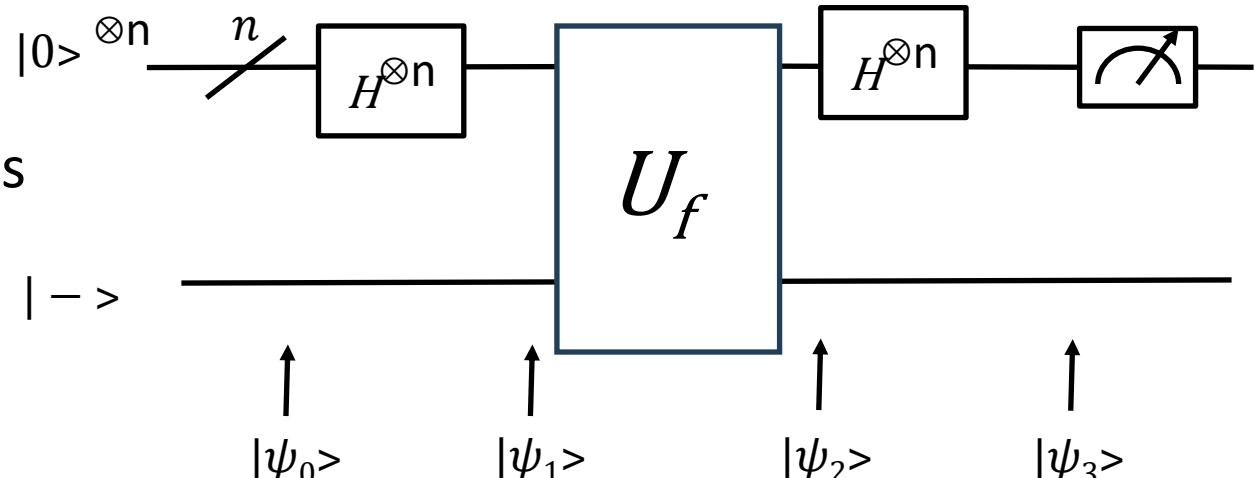
For n=2

$$H^{\otimes 2} |0\rangle^{\otimes 2} = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$$

$$= \frac{1}{\sqrt{2^2}}(|00\rangle + |01\rangle + |10\rangle + |11\rangle)$$

$$H^{\otimes 2} |0\rangle^{\otimes 2} = \frac{1}{\sqrt{2^2}} \sum_{x \in \{0,1\}^2} |x\rangle$$

Deutsch-Jozsa Algorithm



For n=3 $H^{\otimes 3} |0\rangle^{\otimes 3} = \frac{1}{\sqrt{2^3}} \sum_{x \in \{0,1\}^3} |x\rangle$

For any n
$$H^{\otimes n} |0\rangle^{\otimes n} = \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} |x\rangle$$

Deutsch-Jozsa Algorithm

$$|\psi_1\rangle = H^{\otimes n} |0\rangle^{\otimes n} |- \rangle$$

$$|\psi_1\rangle = \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} |x\rangle |- \rangle$$

$$|\psi_2\rangle = U_f \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} |x\rangle |- \rangle$$

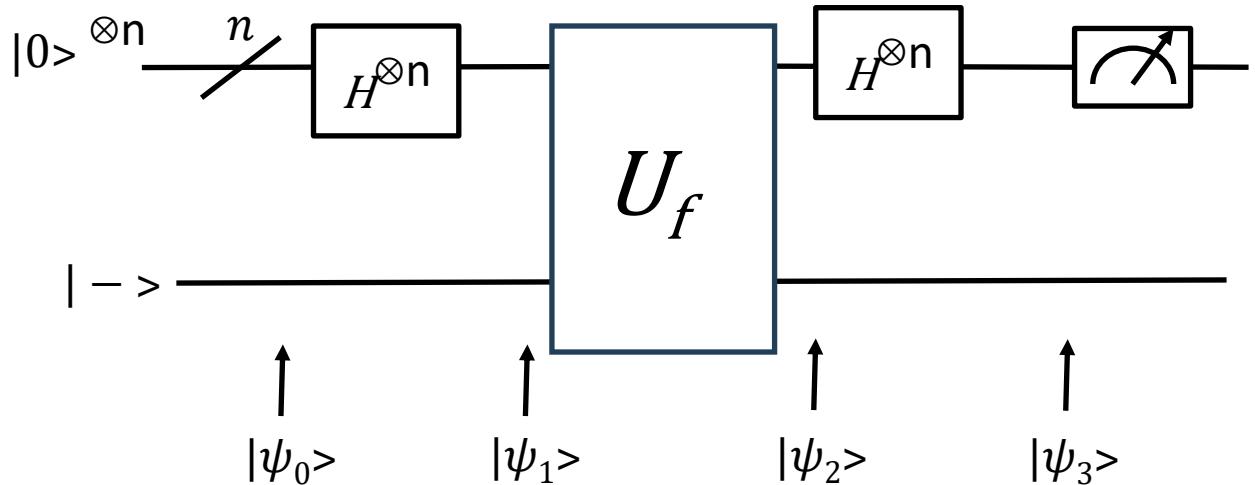
$$|\psi_2\rangle = \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} U_f |x\rangle |- \rangle$$

Recall the Phase Oracle, $U_f |x\rangle |- \rangle = (-1)^{f(x)} |x\rangle |- \rangle$

$$|\psi_2\rangle = \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} |x\rangle |- \rangle$$

We may ignore the $|- \rangle$

$$|\psi_2\rangle = \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} |x\rangle$$



Deutsch-Jozsa Algorithm

$$|\psi_3\rangle = H^{\otimes n} \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} |x\rangle$$

$$|\psi_3\rangle = \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} H^{\otimes n} |x\rangle$$

$$\begin{aligned} H^{\otimes n} |x\rangle &= H|x_0x_1x_2 \dots x_n\rangle \\ &= H|x_0\rangle H|x_1\rangle H|x_2\rangle \dots H|x_{n-1}\rangle \end{aligned}$$

$$H|xi\rangle = \frac{1}{\sqrt{2}}(|0\rangle + (-1)^{xi}|1\rangle), \text{ where } xi$$

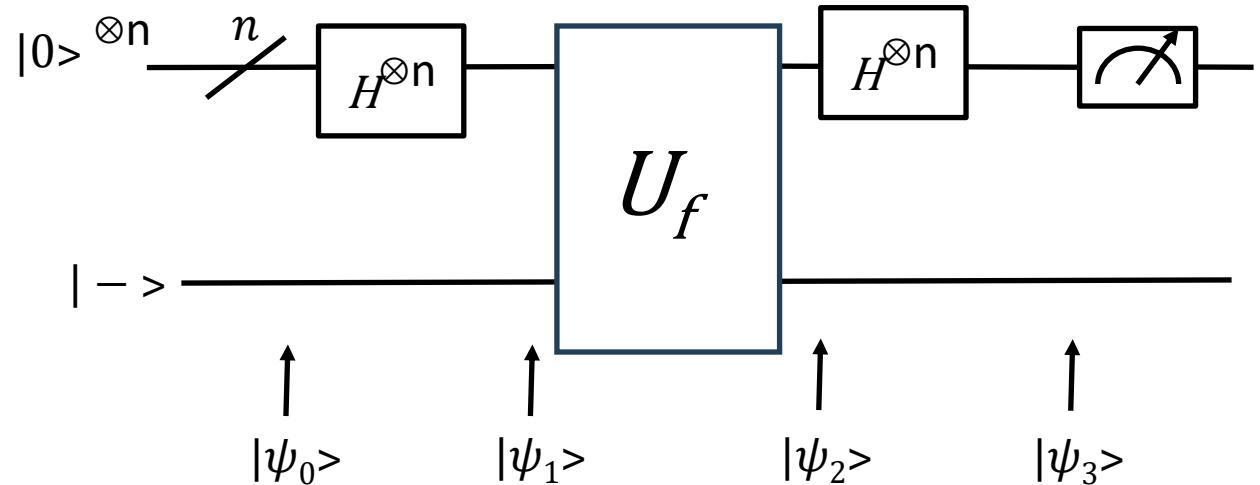
$$\in \{0, 1\}$$

$$H|\mathbf{0}\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) = |+\rangle$$

$$H|\mathbf{1}\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) = |-\rangle$$

$$H^{\otimes n} |x\rangle = \frac{1}{\sqrt{2^n}} \sum_{z \in \{0,1\}^n} (-1)^{x.z} |z\rangle$$

(Derivation skipped)



Deutsch-Jozsa Algorithm

$$|\psi_3\rangle = \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} H^{\otimes n} |x\rangle$$

$$|\psi_3\rangle = \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} \frac{1}{\sqrt{2^n}} \sum_{z \in \{0,1\}^n} (-1)^{x.z} |z\rangle$$

$$|\psi_3\rangle = \sum_{x \in \{0,1\}^n} \frac{1}{2^n} \sum_{z \in \{0,1\}^n} (-1)^{f(x)} (-1)^{x.z} |z\rangle$$

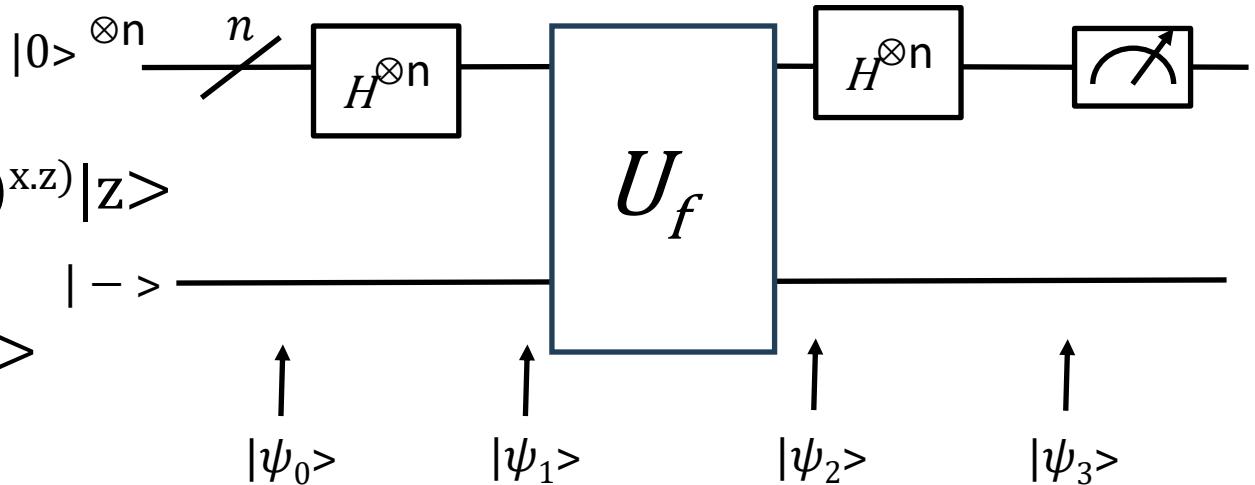
$$|\psi_3\rangle = \sum_{x \in \{0,1\}^n} \frac{1}{2^n} \sum_{z \in \{0,1\}^n} (-1)^{f(x)} (-1)^{x.z} |z\rangle$$

$$|\psi_3\rangle = \sum_{x \in \{0,1\}^n} \frac{1}{2^n} \sum_{z \in \{0,1\}^n} (-1)^{f(x) + x.z} |z\rangle$$

$$|\psi_3\rangle = \sum_{x \in \{0,1\}^n} \frac{1}{2^n} \sum_{z \in \{0,1\}^n} (-1)^{f(x) + x.(0.0.0.0 \cdots 0)} |z\rangle$$

$$|\psi_3\rangle = \sum_{x \in \{0,1\}^n} \frac{1}{2^n} \sum_{z \in \{0,1\}^n} (-1)^{f(x)} |z\rangle$$

If the function f is constant, the value of $f(x)$ would be same, irrespective of input



$|\psi_0\rangle$ $|\psi_1\rangle$ $|\psi_2\rangle$ $|\psi_3\rangle$



Deutsch-Jozsa Algorithm

If the function f is constant, the value of $f(x)$ would be same, irrespective of input

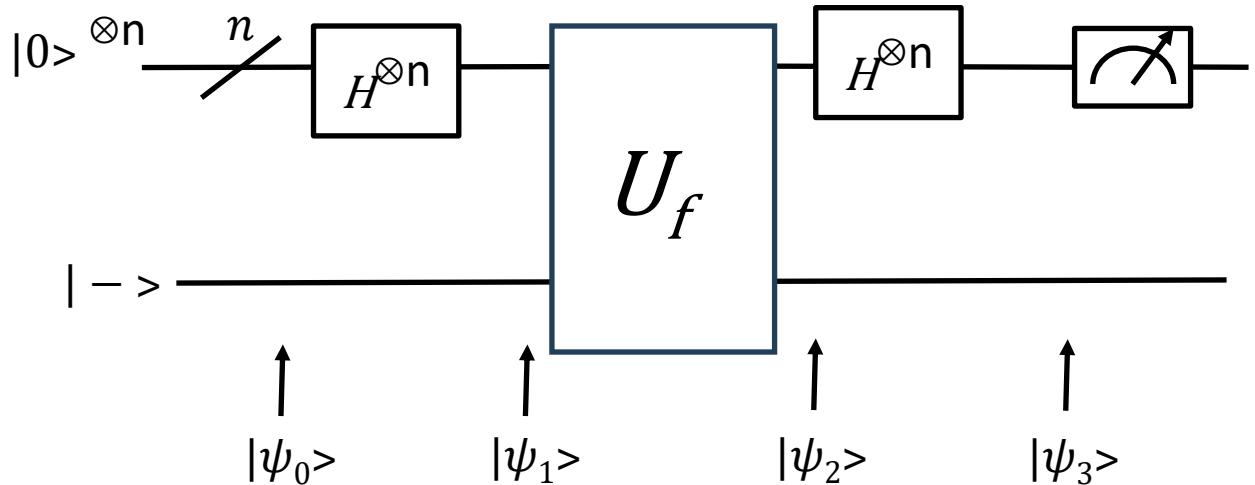
$$\frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} = \frac{1}{2^n} \sum_{x \in \{0,1\}^n} 1, \text{ if } f(x) = 0 \text{ for all } x$$

$$\frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} = \frac{1}{2^n} 2^n, \text{ if } f(x) = 0 \text{ for all } x$$

$$\frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} = 1, \text{ if } f(x) = 0 \text{ for all } x$$

$$\frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} = -1, \text{ if } f(x) = 1 \text{ for all } x$$

If the function f is constant, Amplitude is ± 1



Deutsch-Jozsa Algorithm

If the function f is balance, half the output would result into 0 and other in 1.

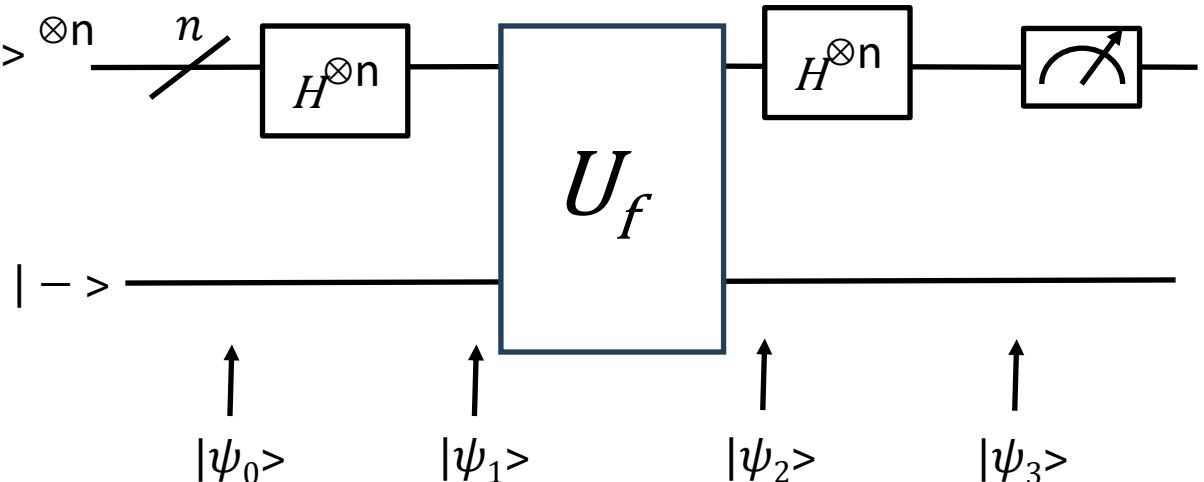
$$\frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} = \frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{f(x)}$$

$$\frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} = \frac{1}{2^n} ((-1)^0 + (-1)^1 + +(-1)^0 + (-1)^1 + \dots (-1)^1)$$

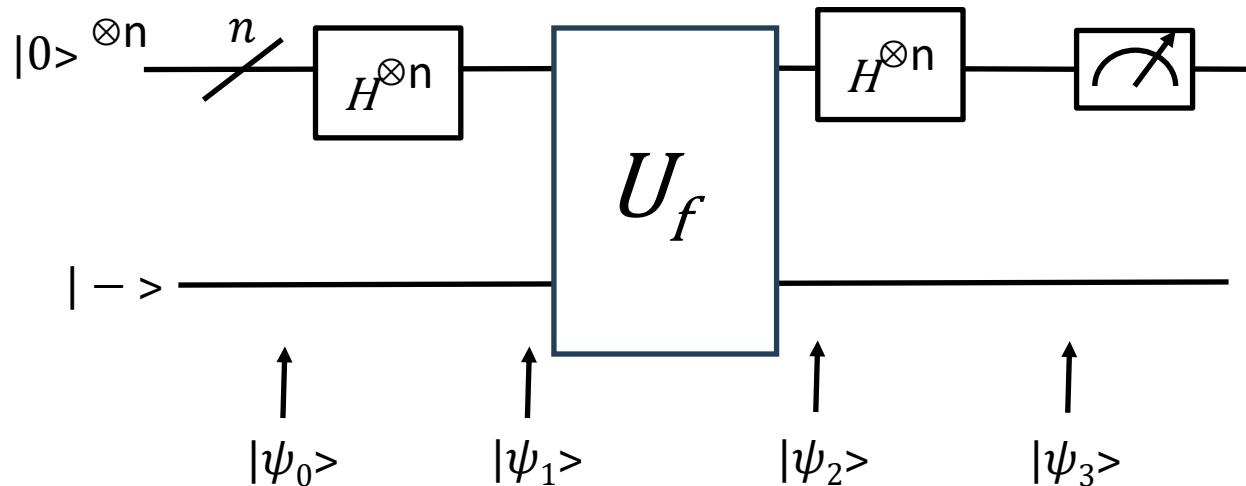
$$\frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} = \frac{1}{2^n} (0)$$

$$\frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} = 0$$

If the function f is balance, Amplitude is 0



Deutsch-Jozsa Algorithm



If the function is constant, Amplitude = ± 1

That is probability of measuring $|0000\dots 0\rangle$ is 1 (100%)

If we measure $|0000\dots 0\rangle$ state then f is constant.

If the function is constant, Amplitude = 0

That is probability of measuring $|0000\dots 0\rangle$ is 0 (0%)

If we measure any other state then f is balanced.

Bernstein-Vazirani Algorithm



Bernstein-Vazirani Algorithm

Imagine we have a function $f: \{0, 1\}^n \rightarrow \{0, 1\}$

$$f(x) = x \cdot s \pmod{2}$$

x is a bit string of length n and s is a secret string of same size. This function returns either 0 or 1. Our task is to find out secret string s .

(mod2) mean we divide the answer of $x \cdot s$ by 2 and take the remainder

For instance, if $x \cdot s = 5$, $x \cdot s \pmod{2} = 1$



Bernstein-Vazirani Algorithm

In Classical Computing,

$$f(0000\dots 0\textcolor{red}{1}) = s_0(0) + s_1(0) + s_2(0) + s_3(0) + \dots + s_{n-2}(0) + s_{n-1}(\textcolor{red}{1}) = s_{n-1}$$

$$f(0000\dots\textcolor{red}{1}0) = s_0(0) + s_1(0) + s_2(0) + s_3(0) + \dots + s_{n-2}(\textcolor{red}{1}) + s_{n-1}(0) = s_{n-2}$$

• • • •

• • •

$$f(000\mathbf{1}\dots00) = s_0(0) + s_1(0) + s_2(0) + s_3(\mathbf{1}) + \dots + s_{n-2}(0) + s_{n-1}(0) = s_3$$

$$f(00\textcolor{red}{1}0\dots00) = s_0(0) + s_1(0) + s_2(\textcolor{red}{1}) + s_3(0) + \dots + s_{n-2}(0) + s_{n-1}(0) = s_2$$

$$f(0\mathbf{1}00\dots00) = s_0(0) + s_1(\mathbf{1}) + s_2(0) + s_3(0) + \dots + s_{n-2}(0) + s_{n-1}(0) = s_1$$

$$f(1000\ldots00) = s_0(1) + s_1(0) + s_2(0) + s_3(0) + \dots + s_{n-2}(0) + s_{n-1}(0) = s_0$$

This means, we need to query the function n times.



$$|\psi_0\rangle = |0\rangle^{\otimes n} |-\rangle$$

$$|\psi_1\rangle = \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} |x\rangle |-\rangle$$

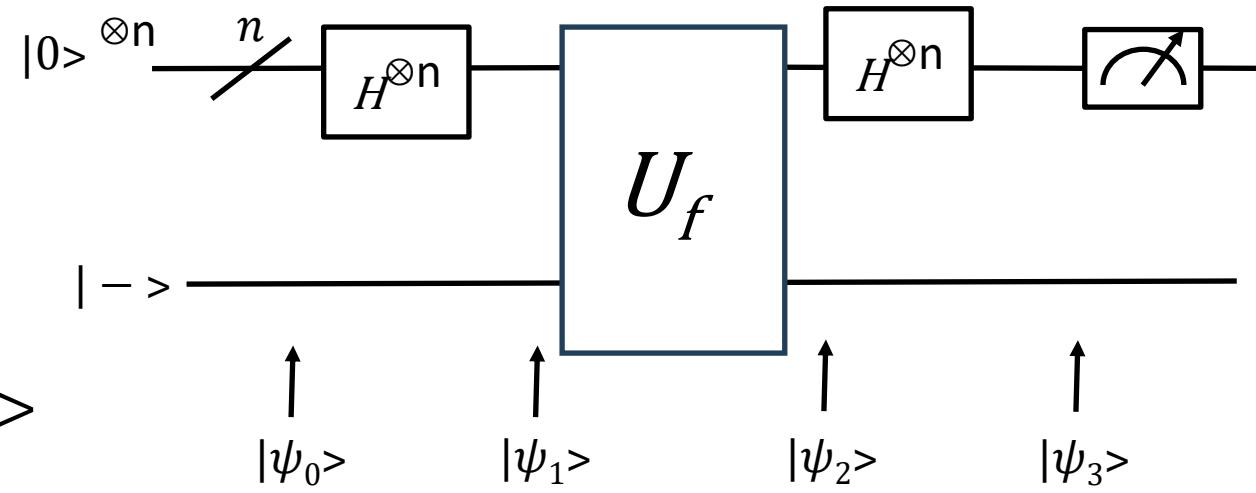
$$|\psi_2\rangle = \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} (-1)^{x \cdot s} |x\rangle$$

$$|\psi_3\rangle = H^{\otimes n} \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} (-1)^{x \cdot s} |x\rangle$$

$$|\psi_3\rangle = \frac{1}{2^n} \sum_{x \in \{0,1\}^n} \sum_{z \in \{0,1\}^n} (-1)^{x \cdot s} (-1)^{x \cdot z} |z\rangle$$

$$|\psi_3\rangle = \frac{1}{2^n} \sum_{x \in \{0,1\}^n} \sum_{z \in \{0,1\}^n} (-1)^{x(s+z)} |z\rangle$$

Bernstein-Vazirani Algorithm



+ is bitwise XOR

Same circuit as Deutch-Jozsa Algorithm

Amplitude of the $|s\rangle$ state:

$$= \frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{(s+s) \cdot x} = \frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{(00\dots 0) \cdot x} = \frac{1}{2^n} \sum_{x \in \{0,1\}^n} 1 = \frac{1}{2^n} 2^n = 1$$

Amplitude of the $|s\rangle$ state = 1

This means, the probability of measuring s after applying the algorithm is 1. That is it. After one query of the function, we can find s by measuring the qubits

